#### CSCI-GA.1170-003/004 Fundamental Algorithms

February 6, 2019

#### Problem Set 2

Lecturer: Yevqeniy Dodis Due: Thursday, February 14

#### Problem 2-1 (Sort recurrences)

8 Points

Sort the following recurrences in increasing order of growth of the corresponding functions. Justify (very) briefly.<sup>1</sup>

- (a) T(n) = 16T(n/4) + n;
- (b) T(n) = 2T(n/4) + n;
- (c)  $T(n) = 3T(n/5) + \log n$ ;
- (d)  $T(n) = 9T(n/3) + n^2$ ;
- (e) T(n) = T(n/3) + 10;
- (f)  $T(n) = 9T(n/3) + n^3$ ;
- (g)  $T(n) = 8T(n/2) + n^3$ .

## Problem 2-2 (Methods for Solving Recurrences)

12 points

Consider the recurrence  $T(n) = T(\lceil n/4 \rceil) + T(\lceil n/3 \rceil) + n$  with T(1) = 1.

- (a) (4 Points) Using a recursion tree, determine a tight asymptotic upper bound on T(n).
- (b) (4 Points) Prove your upper bound using induction.
- (c) (4 Points) Using a suitable variable change, solve the recurrence  $U(n)=3U(\lceil n^{1/3}\rceil)+7$  with U(2)=1.

## Problem 2-3 (The Same Outcome in Different Ways) 15 Points

Consider the following recurrence  $T(n) = 4T(n/2) + n^2 \log n$ , T(1) = 1.

- (a) (2 points) Can the master's theorem, as stated in the book, be applied to solve this recurrence? If yes, apply it. If not, formally explain the reason why.
- (b) (4 points) Solve the above recurrence using the recursion tree method.
- (c) (4 points) Formally verify that your answer from part (b) is correct using induction.
- (d) (5 points) Solve the above recurrence exactly using domain range substitution.

<sup>&</sup>lt;sup>1</sup>For this entire homework assignment, you may ignore the fact that the argument to T may not be an integer.

# Problem 2-4 (Faster Mergesort)

#### 10 points

- (a) (6 points) Suppose you have some procedure FASTMERGE that given two sorted lists of length m each, merges them into one sorted list using  $m^c$  steps for some constant c > 0. Write a recursive algorithm using FASTMERGE to sort a list of length n and also calculate the run-time of this algorithm as a function of c. For what values of c does the algorithm perform better than  $O(n \log n)$ .
- (b) (4 points) Let A[1...n] be an array such that the first  $n \sqrt{n}$  elements are already in sorted order. Write an algorithm that will sort A in substantially better than  $O(n \log n)$  steps.