SUMMER PRACTICE REPORT CENG 400

STUDENT NAME Sait Elmas

ORGANIZATION NAME METU Computer Engineering Department Parallel Pro-

cessing Laboratory

ADRESS Dept. of Computer Engineering Middle East Technical

University Inonu Bulvari, 06531 Ankara, TURKEY

START DATE 1 August 2022

END DATE 12 September 2022

TOTAL WORKING DAYS 30 days

STUDENT'S SIGNATURE

ORGANIZATION APPROVAL

Contents

1	Introduction		2		
		formation About Project			
	2.1	Analys	sis Phase	3	
		2.1.1	Basic Linear Algebra	3	
		2.1.2	Matrix Factorization	4	
		2.1.3	Matrix Norms (Operator Norms)	5	
		2.1.4	Matrix Norms in Matlab and Julia	6	
	2.2	Design	Phase	6	
	23	Impler	nentation Phase	Q	

1 Introduction

In the summer practice period I worked in the **Parallel Processing Laboratory** of Middle East Technical University, Computer Engineering Department. Active research areas are listed on the web page of laboratory as:

- High Performance Computing
- Development of parallel algorithms and their applications in science and engineering
- Parallel sparse matrix computations
- Parallel application development and run-time environments

Beside these research areas there are courses that are being lectured by the members of laboratory to the undergraduate/graduate level students. Some of the courses are:

- CENG371 Scientific Computing
- CENG478 Introduction to Parallel Computing
- CENG487 Introduction to Quantum Computing
- CENG571 Numerical Analysis
- CENG577 Parallel Computing
- CENG780 Sparse Matrix Computations

After my application for summer practice internship and undergraduate student researcher positions, **Prof. Murat Manguoğlu** has offered the topic of "random sparse matrix factorization".

A matrix (array) that contains "a lot of" zeros is called a **sparse matrix (array)**. Although the ratio of zeros to the non-zeros is not strictly defined for a sparse matrix it will not be wrong to say that for a large matrix if the number of zeros are equal to the number of columns or rows that it is a sparse matrix.

In scientific computations the deterministic matrix algorithms are lack to solve the problems containing sparse matrices. One reason of that is these large matrices containing mostly zeros require many floating point operations that will easily overwhelming the capacity of even very powerful workstations. Another reason of this challenge is large matrices takes so much space to fit in the primary memory so there will be need more that one passes to gain the information from the disc and after the application of the relevant algorithm writing it to the disk back. In this situation the performance is measured by the amount of disk access.

For the first problem mentioned above "random methods" are in use. The idea is to sample the given matrix in such a way that the relatively simple sample contains *nearly* all information of the sparse matrix. That is it spans nearly same linear space with the original sparse matrix space. With this simple representative deterministic matrix methods of software packets can be utilized. For the second problem "parallel processing" algorithms are mainly in focus. This topic has not been studied under this summer practice period.

2 Information About Project

As mentioned in the introduction section "random methods for sparse matrix factorization" has been studied as a solution of the first major problem of sparse matrix computations. I had worked on this method by applying the algorithms that are defined in the article "Finding Structure with Randomness: Probabilistic Algorithms for Constructing Approximate Matrix Decompositions" by N. Halko, P. G. Martinson, J. A. Tropp. For the applications the "Julia Language" utilized.

The progress of this work has been supervised by Prof. Murat Manguoğlu via periodical meetings.

2.1 Analysis Phase

At this stage a review of basic linear algebra, an introductory course about numberical linear algebra and tutorials for julia-language has been studied.

2.1.1 Basic Linear Algebra

In this part some preliminary linear algebra subjects are listed. Although they are very basic it may be desirable to see all the terms that will be used at this report in one place.

This part is heavily depend on the book *Linear Algebra by Stephan H. Friedberg* and some related wikipedia pages.

Definition 2.1 (Vector Space). A **vector space** V over a field F (we mainly work with the fields of real and complex numbers (\mathbb{R}, \mathbb{C}) , so no need to define what a field is!) consist of two sets V (set of vectors), F (the field) and two operations + (addition) and \cdot (scalar multiplication). Addition is defined on the set V that is we want to add the vectors to each other. And scalar multiplication is the relation between vectors and the field \mathbb{R} or \mathbb{C} , all together with the following properties:

For any vector v, v_1, v_2, v_3 in V and for any number a, b in \mathbb{R}

- Commutative vector addition: $v_1 + v_2 = v_2 + v_1$
- Associative vector addition: $v_1 + (v_2 + v_3) = (v_1 + v_2) + v_3$
- There must be a **0 element** satisfying: 0 + v = v + 0 = v
- Every vector v should have an **inverse** e such that v + e = 0
- 1v = v1 = v
- (ab)v = a(bv)
- \bullet $(a+b)v_1 = av_1 + av_2$
- $\bullet \ a(v_1 + v_2) = av_1 + av_2$

 \mathbb{R}^n is an example for a vector space over \mathbb{R} . (1,2,3)+(2,3,4)=(2,3,4)+(1,2,3) for the vectors in \mathbb{R}^3 for instance and $(2*3)\cdot(1,2,3,4)=2\cdot(3,6,9,12)$ in \mathbb{R}^4 . All the properties can be easily verified for this natural vector space.

An other basic example is the set of all complex valued mxn dimensional matrices overt the field of complex numbers \mathbb{C}

Definition 2.2 (Inner Product). Let V be a vector space over a field (\mathbb{R} or \mathbb{C}). An **inner product** on V is a function that takes two vectors and gives a number from the field. It is usually denoted as $\langle x, y \rangle$ or just $x \cdot y$ with the following properties hold:

- (a) $\langle x + z, y \rangle = \langle x, y \rangle + \langle z, y \rangle$
- (b) $\langle cx, y \rangle = c \langle x, y \rangle$
- (c) $\overline{\langle x,y\rangle} = \langle y,x\rangle$. here the bar represents the complex conjugate of the number. If the field is the real numbers than it can be ignored.
- (d) $\langle x, x \rangle > 0$ if $x \neq 0$

An example of inner product is the standard Hermitian product on \mathbb{C}^n as:

$$\langle \boldsymbol{x}, \boldsymbol{y} \rangle = \sum_{j=1}^{n} x_j \bar{y}_j$$
 where $\boldsymbol{x} = (x_1, x_2, ... x_n), \boldsymbol{y} = (y_1, y_2, ..., y_n)$

Definition 2.3 (Inner Product Space). A vector space over a field equipped with an inner product is called an **inner product space**

Definition 2.4 (Norm). Let V be a vector space over real or complex numbers. **Norm** of a vector is a function that assigns a real number to every vector in the vector space. It is denoted as ||x||. The function should satisfy following three properties.

- 1. Triangle inequality: $||x + y|| \le ||x|| + ||y||$
- 2. ||rx|| = |r| ||x||
- 3. $\|\boldsymbol{x}\| \geq 0$ for any \boldsymbol{x} and $\|\boldsymbol{x}\| = 0$ if and only if $\boldsymbol{x} = 0$

In different sources norms are denoted in different symbols and even with different names so it may be confusing while reading technical documents. Here is a list of basic norms with their definitions. For $\mathbf{x} = (x_1, x_2, ..., x_n) \in \mathbb{R}^n$

- l_1 norm (Taxicap norm, Manhattan norm): $\|\boldsymbol{x}\|_1 = \sum_{j=1}^n |x_j|$
- l_2 norm: $||x||_2 = \sqrt{\sum_{j=1}^n |x_j|^2}$. This is the standart Euclidian Distance. With l_2 norm and standart inner product of vectors of \mathbb{R}^n has the canonical relation as $||x||_2 = \sqrt{\langle x, x \rangle}$
- $\mathbf{l_p}$ norm (p norm): $\|\boldsymbol{x}\|_p = \left(\sum_{j=1}^n |x_j|^2\right)^{1/p}$
- l_{∞} norm (Infinity norm, maximum norm) $\|x\|_{\infty} = \max_{i} |x_{i}|$

Definition 2.5. For an inner product space V if the product of two vectors \boldsymbol{x} and bmy, $\langle \boldsymbol{x}, \boldsymbol{y} \rangle = 0$ than we say that \boldsymbol{x} and \boldsymbol{y} are **orthogonal** (perpendicular) to each other. If length of a vector is 1 that is $\|\mathbf{x}\|_2 = 1$ than we call it **unit vector**. If a set of vectors are mutually orthogonal to each other this set is called **orthogonal**. Additionally an orthogonal set with every vector has length 1 is called **orthonormal** set.

All of these terms can be extended to define matrices with its columns are regarded as column vectors. A matrix is called orthogonal if all of its columns are mutually perpendicular and orthonormal if length of each column vector is 1 and they are orthogonal.

Definition 2.6 (Rank of a matrix). The **rank of a matrix** is the number of its columns that are linearly independent.

Definition 2.7 (Orthogonalization). Orthogonalization of a matrix A is finding a matrix say Q, that is orthonormal and that can span the same space with A. Q has rank(A) number of columns.

In Matlab the orth command draws an orthonormal matrix for a given matrix A. As a simple function that computes the orthonormal matrix Q for a given matrix can be scripted as below. The method behind this function is called **Gram-Schmidt** orthogonalization process.

Definition 2.8 (Unitary Matrix). A square matrix A is called **unitary** if $AA^H = A^HA = I$ where A + H denotes the **conjugate transpose** of A.

Definition 2.9 (Normal Matrix). A square matrix A is called **normal** if $AA^H = A^H A$.

Definition 2.10 (Hermitian Matrix). A square matrix A is called **Hermitian** if $A = A^H$.

2.1.2 Matrix Factorization

Writing a matrix as a product of two relatively "simple" matrices may simplify complicated and expensive calculations. Here are some basic matrix factorizations.

Eigendecomposition If a square matrix A is diagonalizable than it can be written as $A = QDQ^{-1}$ where Q contains eigenvectors as its columns and D is a diagonal matrix containing the eigenvectors of A on the diagonal entries. One of the important properties of this decomposition is powers of the matrix A can be calculated as:

$$A^n = (QVQ^{-1})^n = QVQ^{-1}QVQ^{-1}...QVQ^{-1}QVQ^{-1} = QV^nQ^{-1}$$

Evaluating powers of a diagonal matrix is very easy for both human and computer.

QR Decomposition Writing a matrix A as a product of two matrices as A = QR, where Q is orthogonal and R is upper triangular.

Singular Value Decomposition For a $m \times n$ matrix A the SVD is the product of three matrices as $A = U\Sigma V^*$. Here U is $m \times m$ complex unitary matrix, V is $n \times n$ complex unitary matrix and Σ is $m \times n$ rectangular diagonal matrix.

$$A_{m \times n} = Q \cdot \sum_{m \times m} V^* \cdot V^*$$

The diagonal entries $\sigma_i = (\Sigma)_{ii}$ are called singular values of A. They are uniquely determined by the matrix A. They are spaced on the diagonal in descending order..

2.1.3 Matrix Norms (Operator Norms)

Definition 2.11 (Matrix Norm). If we see the set of all $m \times n$ matrices with real of complex entries as a vector space then the definition of the vector norm can be regarded as **matrix norm**. That is A its **norm** or magnitude is denoted as ||A|| and satisfies the following properties.

- 1. $||A|| \ge 0$ for any matrix A.
- 2. ||A|| = 0 if and only if A = 0
- 3. $\|\alpha A\| = |\alpha| \|A\|$ for and $\alpha \in \mathbb{R}$ and $A \in M^{m \times n}$
- 4. $||A + B|| \le ||A|| + ||B||$

Definition 2.12 (Matrix p Norms). These are matrix norms induced by vector p norms. Defined as:

$$\|A\|_p = \frac{\|A\boldsymbol{x}\|_p}{\|\boldsymbol{x}\|_p}$$

Here the vector \boldsymbol{x} lives in \mathbb{K}^n and $A\boldsymbol{x}$ lives in \mathbb{K}^m where \mathbb{K} is real or complex numbers.

• p=1 induces maximum absolute column sum as;

$$||A||_1 = \max_{1 \le j \le n} \sum_{i=1}^m |a_{ij}|$$

• p=2 induces the spectral norm as maximum singular value of the matrix as;

$$||A||_2 = \sigma_{max}(A)$$

• $\mathbf{p}=\infty$ induces the maximum absolute row sum as;

$$||A||_{\infty} = \max_{1 \le i \le m} \sum_{j=1}^{n} |a_{ij}|$$

Definition 2.13 (Frobenius Norm).

$$||A||_F = \left(\sum_{i=1}^m \sum_{j=1}^n |a_{ij}|^2\right)^{\frac{1}{2}}$$

Definition 2.14 ($L_{p,q}$ Norms).

$$||A||_{p,q} = \left(\sum_{i=1}^{m} \left(\sum_{j=1}^{n} |a_{ij}|^{p}\right)^{\frac{q}{p}}\right)^{\frac{1}{q}}$$

In general $L_{p,q}$ norms regards matrices as collection of column vectors and calculates the norm. For example $L_{2,1}$ norm adds all l_2 norms of columns of the given matrix as:

$$||A||_{2,1} = \sum_{i=1}^{m} \left(\sum_{j=1}^{n} |a_{ij}|^2 \right)^{\frac{1}{2}}$$

2.1.4 Matrix Norms in Matlab and Julia

Matlab uses same command for vector and matrix norms: norm(A,p) calculates p norms, norm(A,"fro") calculates Frobenius norm and norm(A, "inf") calculates p_{∞} norm.

```
>> A
2
      A =
      -2
3
             0
                    1
      -3
             -2
                   5
4
                         - 1
            4
                   -2
5
      -3
                         1
            1
                   3
                        -5
6
8
      >> norm(A,1)
      ans = 11
9
10
      >> norm(A, "inf")
11
      ans = 11
12
13
      >> norm(A,2)
14
      ans = 7.7704 %as seen below p2 norm returns the max singular value
15
16
      >> [U, S, V] = svd(A);
17
18
      >> S
      S =
19
      7.7704
                      0
20
                         0
      0 5.9055
                                      0
21
                      4.7711
           0
                                      0
      0
22
                 0
                                 0.9911
                           0
23
24
      >> norm(A, "fro")
25
      ans = 10.9087
26
27
```

In Julia LinearAlgebra pack supplies the norm operator.

```
1
      using LinearAlgebra
      julia > A
2
      4x4 Matrix{Int64}:
3
      -2 0 1 3
4
      -3 -2 5 -1
5
         4 -2 1
      -3
6
          1
              3 -5
8
      julia > opnorm(A) # returns the spectral norm
9
      7.770351825891583
10
11
      julia > norm(A,2) # returns the Frobenius norm
      10.908712114635714
13
14
15
```

2.2 Design Phase

At this section basics of randomized methods and algorithms defined on the reference article are explained.

As mentioned before a low rank approximation of a matrix is basically to draw a subspace of a given linear space that spans approximately the linear space itself. While drawing a subspace different approaches can be utilized and following algorithms are examples of this approaches.

Algorithm 4.1 Randomized Range Finder

Input: An $m \times n$ matrix A, an integer l.

Output: An orthonormal $m \times l$ matrix Q whose range approximates the range of A.

- 1: Draw an $m \times l$ Gaussian random matrix Ω .
- 2: Form the $m \times l$ matrix Y. $Y \leftarrow A\Omega$.

3: Construct an $m \times l$ matrix Q whose columns form an orthonormal basis for the range of Y. $[Q, R] \leftarrow qr(Y)$

At this first algorithm random columns are selected from a given large matrix and QR factorization of the small sample can be calculated easily. Moreover this QR factorization form a basis for more complex calculations.

Algorithm 4.2 Adaptive Randomized Range Finder

Input: An $m \times n$ matrix A, an integer r as oversampling parameter, a floating point number ϵ as approximation tolerance.

Output: An orthonormal $m \times l$ matrix Q whose range approximates the range of A.

```
1: Draw standard Gaussian vectors w^{(1)}, \dots, w^{(n)} of length n.
 2: for i = 1, 2, ..., r do
          y^{(i)} = Aw^{(i)}
 4: end for
 5: j \leftarrow 0
 6: Q^{(0)} \leftarrow []
                                                                                                                        \triangleright m \times n empty matrix
 7: while \max\{\left\|y^{(j+1)}\right\|, \left\|y^{(j+2)}\right\|, ..., \left\|y^{(j+r)}\right\|\} > \epsilon/(10\sqrt{2/\pi}) do
          Overwrite y^{(j)} by (I - Q^{(j-1)}Q^{(j-1)*})y^{(j)}
 9:
          q^{(j)} = y^{(j)} / ||y^{(j)}||
10:
          Q^{(j)} = [Q^{(j-1)}q^{(j)}]
11:
          Draw a standard Gaussian Vector w^{(j+r)} of length n.
12:
          y^{(j+r)} = (I - Q^{(j)}Q^{(j)*})A^{(j+r)}
13:
          for i = (j+1), (j+2), ..., (j+r-1) do
Overwrite y^{(i)} by y^{(i)} - q^{(j)} \langle q^{(j)}, y^{(j)} \rangle
14:
15:
          end for
16:
17: end while
18: Q \leftarrow Q^{(j)}
```

Here a tolerance value ϵ is given to the algorithm and it returns an orthogonal matrix Q whose range approximates the given matrix A with an max error ϵ . The idea is; at each iteration the algorithm adds \mathbf{r} many new vectors from matrix A and maintains the orthonormal matrix Q with every increment. If the desired approximation reached computation halts and returns the current matrix Q.

Algorithm 4.3 Randomized Power Iteration

Input: An $m \times n$ matrix A, integers l and q.

Output: An orthonormal $m \times l$ matrix Q whose range approximates the range of A.

- 1: Draw an $m \times l$ Gaussian random matrix Ω .
- 2: Form the $m \times l$ matrix Y. $Y \leftarrow (AA^*)^q A\Omega$.
- 3: Construct an $m \times l$ matrix Q whose columns form an orthonormal basis for the range of Y. $[Q, R] \leftarrow qr(Y)$

In this version of randomized range finder the matrix A is multiplied with its transpose and q^{th} power of this product is taken. If the singular values of the input matrix A are in some decay algorithms 4.1 and 4.2 work well. However if the singular values are in a flat structure that is do not in a well decay then power iteration reduces the weight of the small singular vectors and increase the weight of dominant singular vectors so that computation is fastened.

Algorithm 4.5 Fast Randomized Range Finder

Input: An $m \times n$ matrix A, integers l

Output: An orthonormal $m \times l$ matrix Q whose range approximates the range of A.

- 1: Draw an $m \times l$ SRFT test matrix Ω .
- 2: Form the $m \times l$ matrix Y. $Y \leftarrow A\Omega$.
- 3: Construct an $m \times l$ matrix Q whose columns form an orthonormal basis for the range of Y. $[Q, R] \leftarrow qr(Y)$

SRFT is subsampled random Fourier transform and can be computed as:

$$\Omega = \sqrt{\frac{n}{l}}DFR$$

where

- **D** is an $n \times n$ diagonal matrix whose entries are independent random variables uniformly distributed on the complex unit circle.
- **F** is the $n \times n$ unitary discrete Fourier Transform (DFT), whose entries take the values $f_{pq} = n^{-1/2}e^{-2\pi(p-1)(q-1)/n}$ for p, q = 1, 2, 3, ..., n
- **R** is an $n \times l$ matrix that samples l coordinates from n uniformly at random; i.e., its l columns are drawn randomly without replacement from the columns of the $n \times n$ identity matrix.

The idea behind this alteration is; the bottleneck of the presented algorithms are the computation of the product $A\Omega$ that requires O(mnl) flops. However with an structured random matrix this can be reduced to O(mnlog(l))

Algorithm 5.1 Singular Value Decomposition

Input: $m \times n$ matrix A, $m \times k$ matrix Q with $||A - QQ^*A|| < \epsilon$ for a desired ϵ .

Output: Matrices U, S, V where U, V are orthonormal and S non-negative diagonal matrices.

- 1: Form the matrix $B = Q^*A$
- 2: Compute an SVD of the small matrix $B = \bar{U}\Sigma V^*$
- 3: Form the orthonormal matrix $U = Q\bar{U}$

At this last algorithm singular value decomposition for the matrix A which is approximated by the matrix Q is being calculated.

2.3 Implementation Phase

At this stage applications of the given algorithms with Julia-language has been utilized. The complete package can be found the following repository with a Jupyter Notebook containing a demonstration. https://github.com/SaElmas/low_rank_approximations

```
using LinearAlgebra
      using Random
2
      using Printf
3
      using Plots
4
5
      using Statistics
6
      using Distributions
      using SparseArrays
      using MatrixMarket
9
      function generate_sampling(n,s,p)
10
      # n : number of columns
11
      # s : size of sampling
12
      # p : probabilities for each column
13
14
      # generate an idendity matrix nxn
15
      # select random s columns from this
16
17
      eye = Matrix(1.0I, n,n);
      total = 0;
18
      if p isa Matrix{Float64}
19
         for i = 1:n
         eye[i,i] = 1 / sqrt(p[i]*s);
21
22
         end
      else
23
        for i = 1:n
24
          total += norm(A[:,i])*norm(B[i,:]);
25
26
27
         for i= 1:n
           pni = norm(A[:,i])*norm(B[i,:]) / total;
```

```
eye[i,i] = 1 / sqrt(pni*s)
29
         end
30
31
       end
       rng = MersenneTwister();
       perm = randperm(rng,s);
       S = eye[:,perm];
35
36
       return S;
       end
37
38
       function basicMatrixMulitplication(A, B, c, P)
39
       # A mxn
40
       # B nxp
41
       # c sampling number c <= n</pre>
42
       # P probabilities contains n float. 0 < pi < 1 and sum(pi) = 1;
43
       _n = size(A);
45
       S = generate_sampling(n,c,P);
       A * S , S' * B;
46
47
48
       # Algorithm 4.1 Randomized Range Finder
49
       # Input A(mxn) matrix, integer 1
50
       # Output Q(mxl) matrix, approximates the range of {\tt A}
51
52
       function rrf(A, 1)
       rng = MersenneTwister()
53
       _n = size(A);
       Om = randn(rng, Float64, (n,1))
       Y = A * Om
56
57
       Q_{-,-} = qr(Y)
       Q = Matrix(Q_)
58
       return 0
59
       end
60
61
       # Algorithm 4.2 Adaptive Randomized Range Finder
62
       # Input A(mxn) matrix, tolerance eps, integer r as oversampling parameter
63
       # Output Q(mxl) orthonormal with tolerance holds with probability 1-min\{m,n\}10^-r
       function arrf(A,eps,r,plot_step)
       # plot_step is used to generat errors and iterations vectors
       # if plot_step is 0 then no vectors are empty
67
       # for a positive value of plot_step the value is used for iteration step.
68
       (m,n) = size(A);
69
       W = zeros(n,r)
70
       Y = zeros(m,r)
71
       Q = zeros(m,1)
72
       j = 0
73
       max_err=0
74
      for i=1:r
75
       w = randn(n,1)
76
       W[:,i] = w
77
78
      Y[:,i] = A*w
       end
79
80
       for i=1:r
81
      ny = norm(Y[:,i])
82
       if ny > eps/(10*sqrt(2/pi))
83
84
       max_err = ny
       end
85
       iteration_step=0;
87
88
       iterations = []
       errors = []
89
90
       while(max_err > eps)
91
       iteration_step +=1;
92
       if plot_step > 0
93
       if iteration_step % plot_step == 0
94
       append!(iterations, iteration_step)
```

```
append!(errors, max_err)
96
       end
97
98
       end
       j += 1
       yj = (1.0I-Q*Q')*Y[:,j]
       qj = yj / norm(yj);
102
       Y[:,j] = yj;
       if j==1
103
       Q[:,j] = qj
104
       else
       Q = cat(Q,qj, dims=2)
106
107
       end
       wjr = randn(n,1)
108
       yjr = (1.0I - Q*Q')*(A*wjr)
109
       Y = cat(Y,yjr,dims=2)
       Y[:,j+r] = yjr
112
       for i = j+1:j+r-1
       yi = Y[:,i]
113
       Y[:,i] = yi - (qj'*yi)*qj
114
       end
115
       max_err = 0
116
       for i = j+1: j+r-1
117
       ny = norm(Y[:,i])
118
       if ny > eps/(10*sqrt(2/pi))
119
       max_err = ny
120
       end
121
       end
123
       end
124
       return Q, iterations, errors
       end
126
       # Algorithm 4.3 Randomized Power Iteration
       # Input A(mxn) matrix, integer 1, power q
       \mbox{\tt\#} Output Q(mxl) matrix, approximates the range of A
129
       function rpi(A, 1, q)
130
131
       rng = MersenneTwister()
       _n = size(A)
       Om = randn(rng, Float64, (n,1))
133
       Y = (A*A')^q*A*Om
134
135
       Q_{-,-} = qr(Y)
       Q = Matrix(Q_{-})
136
       return Q
137
       end
138
139
       # Algorithm 4.4 Randomized Subspace Iteration
140
       # Input A(mxn) matrix, integer 1, power q
141
       # Output Q(mxl) matrix, approximates the range of A
142
       function rsi(A, 1, q)
143
144
       rng = MersenneTwister()
145
       _n = size(A)
146
       Om = randn(rng, Float64, (n,1))
       Y = A * Om
147
       Q = qr(Y)
148
       for i = 1:q
149
       Oshow size(A)
150
151
       Oshow size(Q)
       Y = A * Q
152
       Q = qr(Y)
       end
       return Q
156
       end
       # Algorithm 4.5 Fast Randomized Range Finder
158
       # Input A(mxn) matrix, integer 1
159
       # Output Q(mxl) matrix, approximates the range of A
160
       function frrf(A,1)
161
       m,n = size(A)
162
```

```
D = rucm(n);
163
        F = dftg(n);
164
        R = Matrix(1.0I,n,1)[:,shuffle(1:end)]
165
        Om = sqrt(n/1)*D*F*R
        Y = A * Om
167
        q,r = qr(Y)
168
169
        Q = Matrix(q)
170
       return Q
        end
171
172
        # Algorithm 5.1 Direct SVD
        \mbox{\tt\#} Matrix A(mxn), Q(mxk) matrices with
174
        # with |A-QQ*A| < epsilon
175
        # Matrices U,S,V, U,V are orthonormal, S nonnegative diagonal matrices
176
177
        function direct_svd(A,Q)
179
        B = Q * A
        Uh,S,V = svd(B)
180
        U = Q * Uh
181
        return Q,S,V
182
        end
183
184
        # Uniform Discrete Fourier Transform Generator
185
186
        # Input integer n
        # Output F(nxn) complex matrix
187
        function dftg(n)
188
        F = ones(ComplexF64,n,n)
189
        for i=1:n
190
191
        for j=1:n
        F[i,j] = n^{(-0.5)} \exp(-2 \cdot pi \cdot (i-1) \cdot (j-1)/n)
192
        end
193
        end
194
        return F
195
        end
196
197
198
        # Random Complex matrix whose entries are uniformly distributed
199
        # on unit circle.
        # Random n real numbers in [-pi, pi] that is uniformly distributed chosen
        # those will be considered as angles of complex numbers on unit circle.
201
202
        function rucm(n)
        rng = MersenneTwister(1234)
203
        angles = rand(Uniform(-pi,pi), n,n)
204
        D = ones(ComplexF64, n,n)
205
        for i = 1:n
206
        for j = 1:n
207
        D[i,j] = exp(angles[i,j]im)
208
        end
        end
211
        return D
212
        end
213
214
215
        # Gram-Schmidt orthornormalization
216
        function gso(A)
217
       m,n = size(A)
218
        Q = zeros(m,n)
219
        v1 = A[:,1]
        Q[:,1] = v1 / norm(v1);
221
        for k = 2:n
222
        w = A[:,k]
223
        v = v
224
       for j = 1:k-1
225
        vj = Q[:,j]
226
        v -= (w'*vj)*vj
227
        end
228
        Q[:,k] = v / norm(v)
229
```

```
230 end
231 return Q;
232 end
233
234
```