

String-Search Algorithms Performance Comparison

The algorithms included are:

- Naive Substring Search
- Knuth-Morris-Pratt (KMP)
- Aho-Corasick
- Boyer-Moore

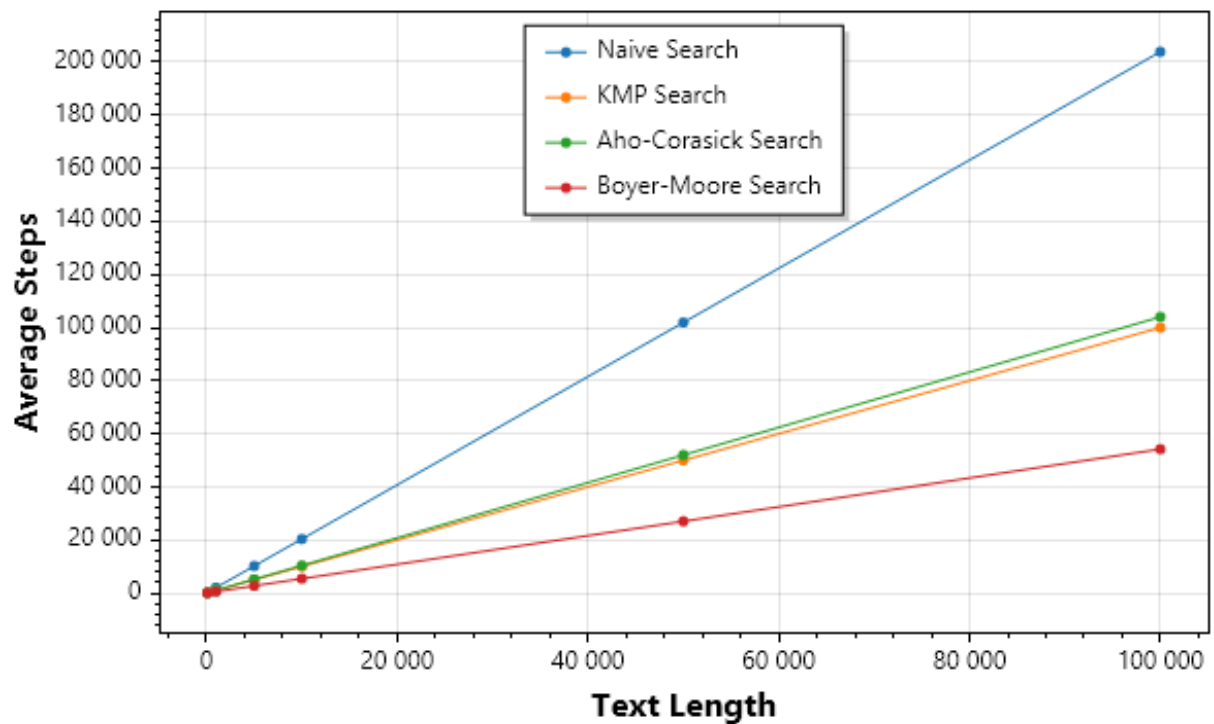
Testing Suites

Text Lengths: 100, 1000, 5000, 10000, 50000, 100000 characters

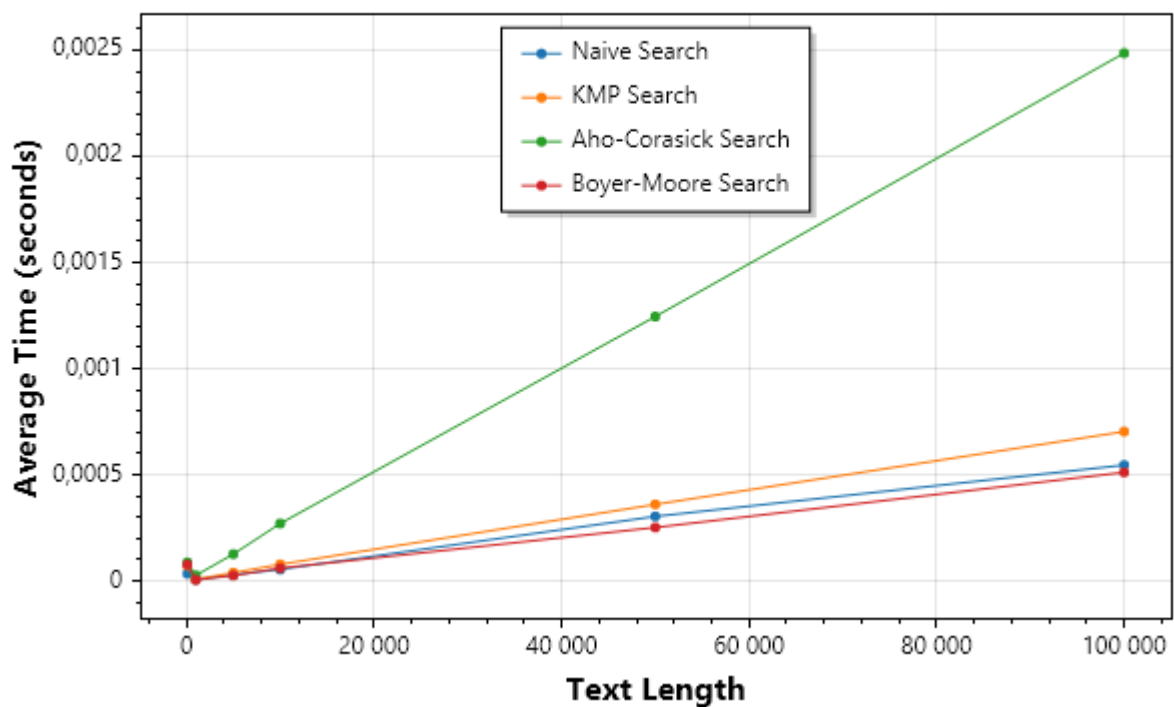
Pattern Length: 2 characters

Number of Trials: 5 per configuration

Performance Comparison of Substring Search Algorithms - Steps



Performance Comparison of Substring Search Algorithms - Time

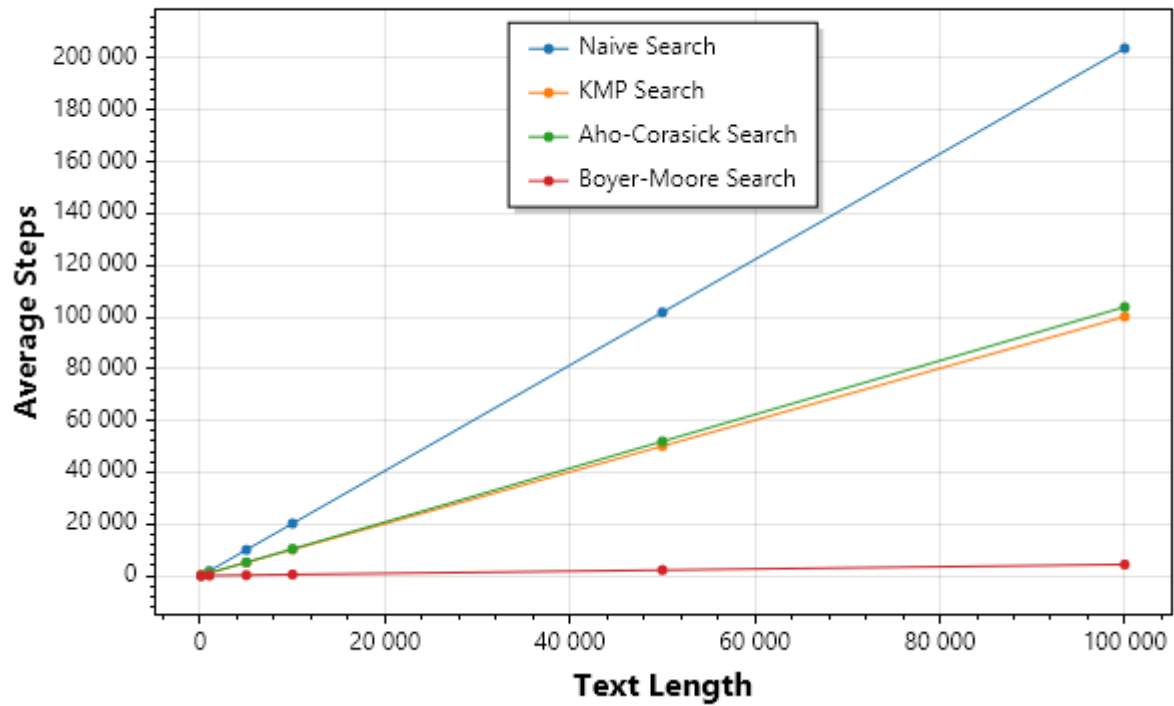


Text Lengths: 100, 1000, 5000, 10000, 50000, 100000 characters

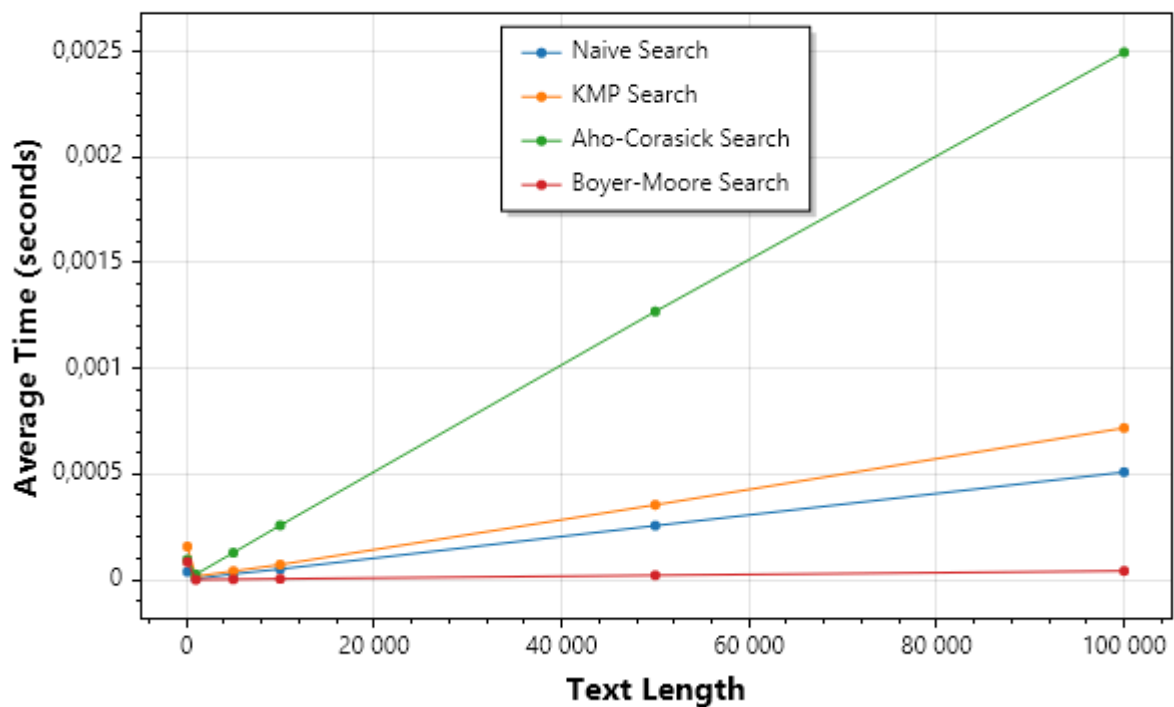
Pattern Length: 100 characters

Number of Trials: 5 per configuration

Performance Comparison of Substring Search Algorithms - Steps



Performance Comparison of Substring Search Algorithms - Time



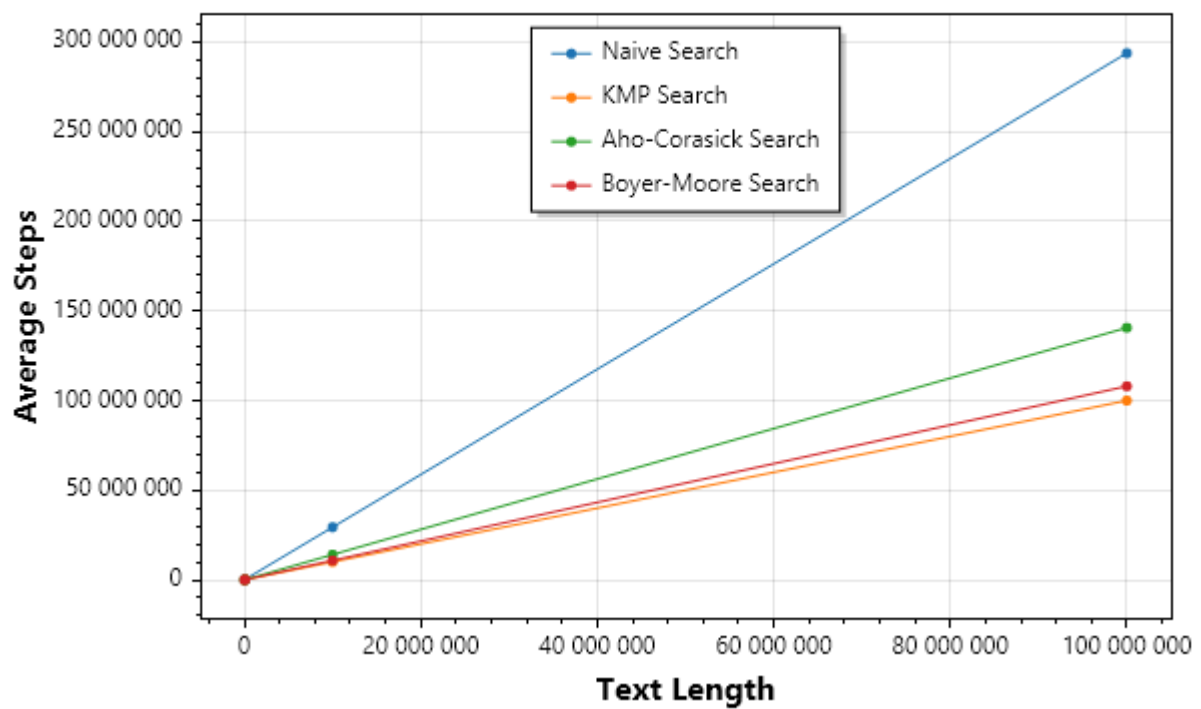
Text Lengths: 10,000, 10,000,000, 100,000,000, 1,000,000,000 characters

Pattern Length: 5 characters

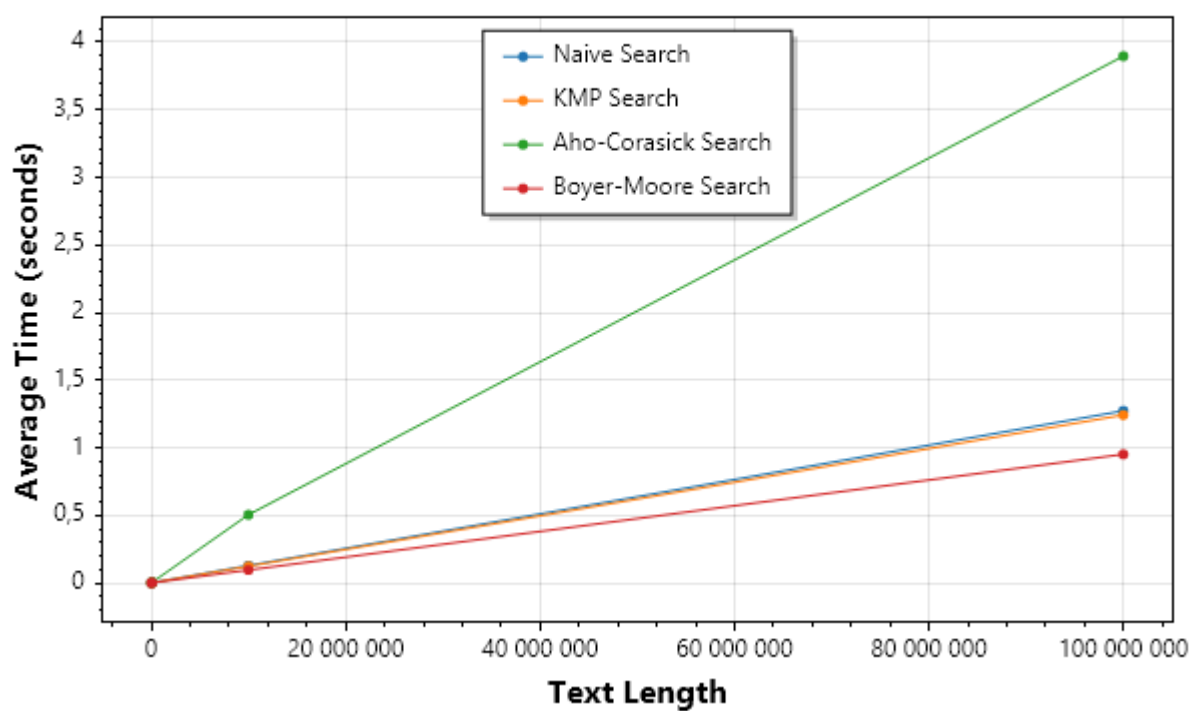
Number of Trials: 1 per configuration

Alphabet: 'ab'

Performance Comparison of Substring Search Algorithms - Steps



Performance Comparison of Substring Search Algorithms - Time



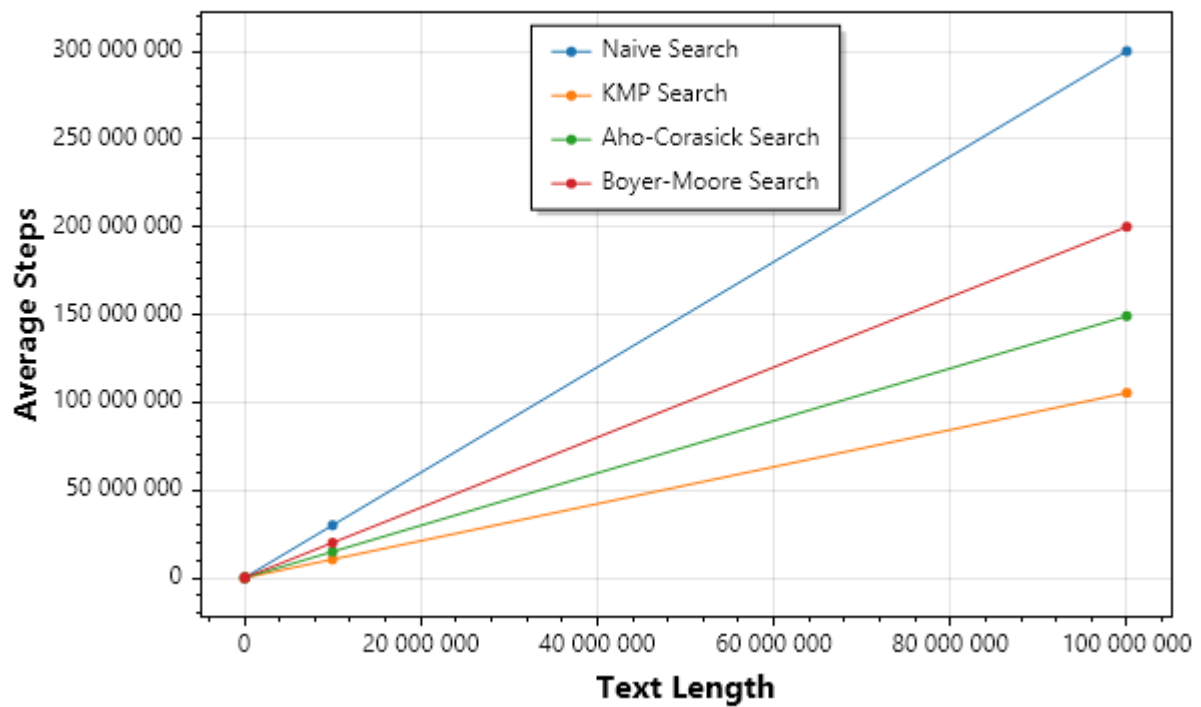
Text Lengths: 10,000, 10,000,000, 100,000,000, 1,000,000,000 characters

Pattern Length: 5000 characters

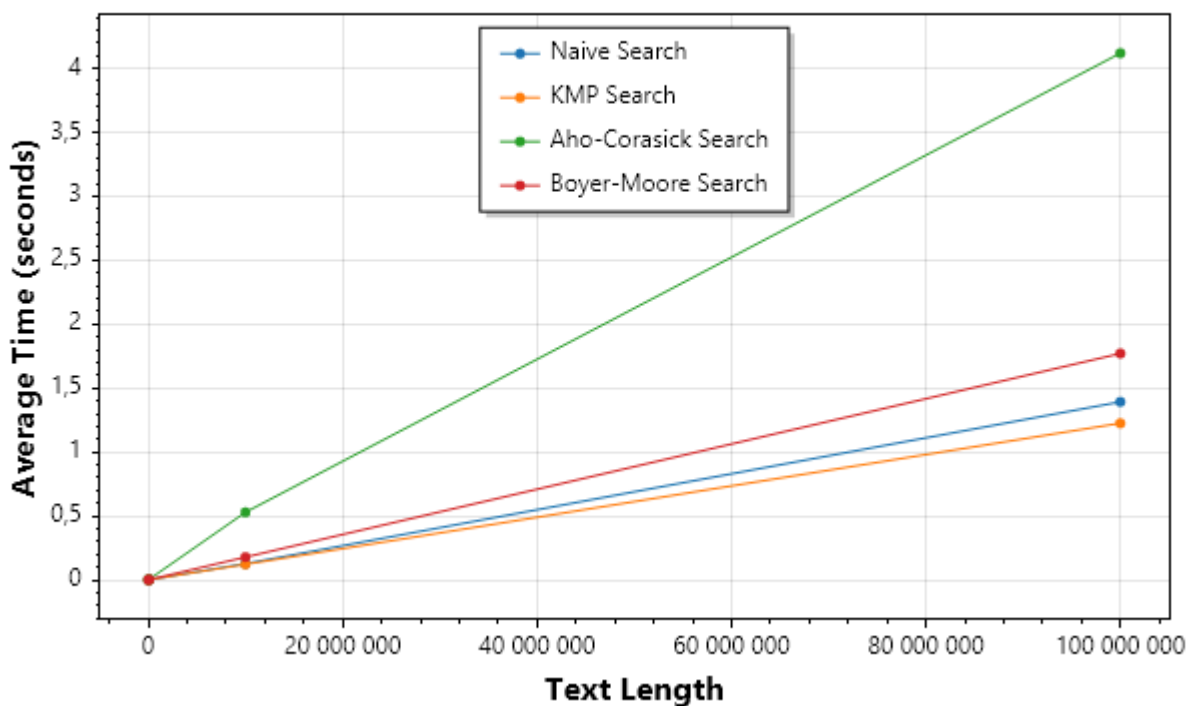
Number of Trials: 1 per configuration

Alphabet: 'ab'

Performance Comparison of Substring Search Algorithms - Steps



Performance Comparison of Substring Search Algorithms - Time



Complexity Analysis

1. Naive Substring Search:

- **Steps:** The Naive Search has a worst-case complexity of $O(n \times m)$, where n is the text length and m is the pattern length, as it checks each position in the text for a potential match.
- **Time Complexity:** The naive approach tends to perform worse with longer texts and larger patterns, which matches with increased execution time in tests with longer text lengths.

2. Knuth-Morris-Pratt (KMP):

- **Steps:** KMP optimizes by avoiding redundant checks, achieving a complexity of $O(n+m)$. The preprocessing of the pattern to build the partial match table takes $O(m)$, and the search through the text takes $O(n)$.
- **Time Complexity:** KMP's performance should remain consistent regardless of pattern length, as indicated by the tests showing improved performance in larger configurations compared to the naive approach.

3. Aho-Corasick:

- **Steps:** Designed for searching multiple patterns, Aho-Corasick has a complexity of $O(n+k)$, where k is the total length of all patterns. Building the automaton is typically more costly, but the search itself is linear with respect to the text length.
- **Time Complexity:** Aho-Corasick's performance benefits from large patterns and texts, with results showing efficient performance in multi-pattern matching scenarios in large text sizes.

(at least it is supposed to be so? because i took realization from ur python file and remade it in C#, maybe this strange time-consuming operations are the consequences of using C# structures?)

4. Boyer-Moore:

- **Steps:** With a complexity averaging $O(n/m)$ but potentially $O(n \times m)$ in the worst case, Boyer-Moore is highly efficient when the pattern is long and contains fewer common letters with the text. The bad-character heuristics improve its performance significantly.
- **Time Complexity:** The Boyer-Moore algorithm shows the fastest times for certain large text configurations, especially when the pattern is short.