Insertion Sort

Introduction

Insertion Sort is a simple sorting algorithm that builds the final sorted list one element at a time. It is efficient for small datasets or partially sorted lists and is often used in combination with more complex algorithms.

Logic

- 1. Start with the second element (index 1) and compare it with the previous elements.
- 2. If the current element is smaller (or larger, depending on the desired order) than the previous element, shift the previous element one position to the right.
- 3. Repeat step 2 until you find the correct position for the current element.
- 4. Insert the current element into its correct position in the sorted part of the list

C++ implementation of Insertion sort: Insertion Sort. Python implementation of Insertion sort: Insertion Sort. Java implementation of Insertion sort: Insertion Sort. Javascript implementation of Insertion sort: Insertion Sort.

Complexity

- Time Complexity:
 - Best Case: O(n) (when the list is already sorted)
 - Average Case: O(n^2)
 - Worst Case: O(n^2)
- Space Complexity:
 - O(1) (constant extra space)

Advantages

- Simple and easy to implement.
- Efficient for small datasets or nearly sorted lists.

Considerations

- Inefficient for larger datasets due to its quadratic time complexity.
- Not suitable for datasets with large inversions (elements that are far from their sorted position).

Disadvantages/Limitations

- The number of comparisons and shifts can be relatively high, especially for larger datasets.
- There are more efficient sorting algorithms available for larger and unsorted datasets.

Edge Cases

- Insertion Sort's performance is relatively better when the initial list is partially sorted.
- It's an adaptive algorithm; its performance improves if the data is almost sorted.

External Links

To learn more about Insertion Sort, explore these resources: - Insertion Sort - Wikipedia - Sorting Algorithms: Insertion Sort - Geeksfor Geeks - Insertion Sort Visualization - Visualgo