

我们为什么要考斐波那契数列？

一道上机笔试题的解析，兼谈技术面试与编程基本功

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大纲

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题目

笔试规则

- 时间不限，语言不限，使用的工具不限；
- 可以 Google 可以百度，但不可向他人求助；
- 每个问题的代码运行时间应该不超过 1 分钟；
- 从 6 个大问题 14 个小问题中任选 N 个小问题解决；
- 如果能成功解决了 2 个或 2 个以上小问题，则进入下一轮技术终面，否则面试流程结束。

题目

在数学上，著名的斐波那契数列以递归的方法来定义：

$$F_n = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ F_{n-1} + F_{n-2} & n \geq 2 \end{cases}$$

1. 试算出 F_{100} 的准确整数值。
2. 请问斐波那契数列中，第一个有 1000 位数字的是第几项？
3. 试算出第 10 亿项除以 1000000007 的余数，即 $F_{10^9} \bmod (10^9 + 7)$ 。
4. 试算出第 10^{1000} 项除以 1000000007 的余数，即 $F_{10^{1000}} \bmod (10^9 + 7)$ 。

这道题目以后，收到很多质疑：

- 题目太简单了，一个循环就出来了。
- 题目太简单了，一个递归就出来了。
- 题目太简单了，一个矩阵快速幂就全秒了。
- 题目太难了，搬转工程师不应该考算法。
- 题目太难了，不应该考数学。
- 题目太难了，不应该考递归。

经典小学奥数题：

某人上台阶，可以一步上一级，也可以上两级。请问上 100 级的台阶，一共有多少种走法？

如果考算法

经典递归算法题：

用一些 1×2 的砖，铺满 $2 \times n$ 的矩形区域，一共有多少种铺法？

失败的案例

直接递归, 卒

典型:

```
1 def fib(n):  
2     if n == 0 or n == 1:  
3         return n  
4     else:  
5         return fib(n-1) + fib(n-2)
```

直接递归, 卒

简洁型, 一行:

```
def fib(n): return fib(n-1) + fib(n-2) if (n > 2) else n
```

直接递归，卒

- 直接接递归，无论机器多猛，一般在 $n = 45$ 左右就是极限了；
- $n = 100$ 递归到太阳熄灭都不会有结果；
- 一般对算法的复杂度没有概念，仅仅会用简单的递归函数；
- 我们是重度使用 Scala 函数式编程的团队，最怕的就是这种写法；
- 把第一个小问题设计为 $n = 100$, 就是为了暴露这个问题。

不能正确处理大数运算，卒

Java:

```
1  public class Fib {
2
3      public static long fib(int n) {
4          long a = 0L, b = 1L, c = 1L;
5          for (int i = 1; i <= n; i++) {
6              c = b;
7              b = a + b;
8              a = c;
9          }
10         return c;
11     }
12
13     public static void main(String args[]) {
14         System.out.println(fib(100)); // 3736710778780434371
15     }
16 }
```

不能正确处理大数运算，卒

JavaScript:

```
1 function fib(n) {  
2     for (i = 0, x = 0, y = 1; i < n; i++) {  
3         var z = y  
4         y = x + y  
5         x = z  
6     }  
7     return x  
8 }  
9  
10 console.log(fib(100)) // 354224848179262000000
```

不能正确处理大数运算，卒

- 有人抱怨这里故意埋大整数溢出的坑考他们。
- 能够正确地处理比较明显的数值溢出，是每一位程序员应该具备的基本素质。
- 就算考虑实际的项目，很多时候仍然需要注意数值的溢出。
 - 曾经把某些业务数据的统计值类型错误地设计成 32 位整数，导致溢出。
 - 广告业务，一张月报表的广告请求数，很容易超过 21 亿的。
 - 直接把所有统计值的类型设计成 64 位整数行不行？

不能正确处理大数运算，卒

- 有不少人选意识到了大整数溢出的问题，会直接用数组之类的实现大数加法。
- 严格来说，也可以归结为一种“错误”。
- 不熟悉自己使用的编程语言，重复制造无意义的轮子。

不能正确处理大数运算，卒

- 第一个小问题设计成 $n = 100$ 而不成 $n = 50$,
- 制造大于 64 位无符号整数的结果。
- 用 Python 做题的人可能在毫无意识的情况下避开了这个坑。

直接用通项公式计算，卒

$$F_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right]$$

直接用通项公式计算，卒

- 一般是算不出正确答案，为什么？
- 犯这个错的人，一般数学比较好，但对工程意义上的“计算”没有概念的；
- 我们不考查数学能力，即使知道了这个公式也没什么作用；

两个数加来加去加错，卒

```
def fib(n):  
    if n == 0 or n == 1:  
        return n  
    else:  
        (a, b) = (1, 1)  
        for i in range(n):  
            b = a + b  
            a = b  
        return b
```

```
fib(10) # 1024 WTF!!!!
```

两个数加来加去加错，卒

- 这类错误难以一一列举;
- 基本属于久没写代码，或者不会写代码;
- 一想好简单，一写就错;
- 一错就说题目难。

其他错误

- 一上来就用动态规划 (DP), 但是实现错误的
 - 实际上这个问题不属于典型的动态规划问题,
 - 最多只能算是动态规划的极端版,
 - 一般刷过算法题,
 - 但对算法理解过浅,
 - 不能处理实际工程项目里更加普遍更加复杂的算法问题。
- 一上来就祭出矩阵快速幂, 因姿势不对而失败。这种一般是
 - 理论知识较好,
 - 但工程能力稍欠,
 - 实现已有算法能力不够强

参考解法

直接循环：简单粗暴，快速有效

需要解决的问题：

- 计算逻辑；
 - 要求用笔死算，90% 以上没问题；
 - 要求写代码算，90% 以上凉凉！
- 大数运算。
 - Scala 有 BigInt,
 - Java 有 BigInteger,
 - JavaScript 有 BigInt,
 - Python 默认的 Int 就是,
 - Haskell 有 Integer

直接循环：简单粗暴，快速有效

Scala:

```
1  object FibLoop extends App {  
2      def fib(n: Int) = {  
3          var (a, b, c) = (BigInt(0), BigInt(1), BigInt(1))  
4          var i = 1  
5          while (i <= n) {  
6              c = b  
7              b = a+b  
8              a = c  
9              i += 1  
10         }  
11         a  
12     }  
13     println(fib(100)) // 354224848179261915075  
14 }
```

直接循环：简单粗暴，快速有效

Java:

```
1  import java.math.BigInteger;
2
3  public class FibJava {
4
5      public static BigInteger fib(int n) {
6          BigInteger a = BigInteger.valueOf(0L);
7          BigInteger b = BigInteger.valueOf(1L);
8          BigInteger c = BigInteger.valueOf(1L);
9          for (int i = 1; i <= n; i++) {
10             c = b;
11             b = a.add(b);
12             a = c;
13         }
14         return a;
15     }
16
17     public static void main(String args[]) {
18         System.out.println(fib(100)); // 354224848179261915075
19     }
20 }
```

直接循环：简单粗暴，快速有效

Python:

```
1  def fib(n):
2      a, b = 0, 1
3      i = 1
4      while (i <= n):
5          a, b = b, a + b
6          i = i + 1
7      return a
8
9  print(fib(100)) # 354224848179261915075
```

直接循环：简单粗暴，快速有效

JavaScript:

```
1 function fib(n) {  
2     for (i = 0, x = 0n, y = 1n; i < n; i++) {  
3         y = x + y  
4         x = y - x  
5     }  
6     return x  
7 }  
8  
9 console.log(fib(100)) // 354224848179261915075n
```

直接循环：简单粗暴，快速有效

Haskell:

```
1  -- 本人不会!
```

递归：函数式的解

Scala:

```
1 object FibTailRec extends App {  
2   def fib(n: Int) = {  
3     @scala.annotation.tailrec  
4     def f(x: BigInt, y: BigInt, i: Int): BigInt =  
5       if (i == 0) x else f(y, x + y, i - 1)  
6     f(0, 1, n)  
7   }  
8   (1 to 100).map(x => (x, fib(x))).foreach(println)  
9 }
```

递归：函数式的解

Java:

```
1  import java.math.BigInteger;
2
3  public class FibJava {
4
5      private static BigInteger _fib(BigInteger x, BigInteger y, int i) {
6          return (i == 0) ? x : _fib(y, x.add(y), i - 1);
7      }
8
9      public static BigInteger fib(int n) {
10         return _fib(BigInteger.valueOf(0L), BigInteger.valueOf(1L), n);
11     }
12
13     public static void main(String args[]) {
14         System.out.println(fib(100)); // 354224848179261915075
15     }
16 }
```

递归：函数式的解

Python:

```
1 def fib(n):  
2     def f(x, y, i): return x if not i else f(y, x + y, i-1)  
3     return f(0, 1, n)  
4  
5 print(fib(100)) # 354224848179261915075
```


递归：函数式的解

JavaScript

```
1  'use strict';
2
3  function fib(n) {
4      let f = (x, y, i) => (i === 0) ? x : f(y, x+y, i-1)
5      return f(0n, 1n, n)
6  }
7
8  console.log(fib(100)) // 354224848179261915075n
```

递归：函数式的解

Haskell:

```
1 fib n = fib' 0 1 n where
2     fib' x y 0 = x
3     fib' x y i = fib' y (x + y) (i-1)
4
5 main = do print $ fib 100 -- 354224848179261915075
```

其他解法

Haskell 著名的斐波那契数列生成式:

```
1  fibs = 1 : 1 : (zipWith (+) fibs (tail fibs))
```

```
1  ~ $ ghci
2  GHCi, version 8.6.5: http://www.haskell.org/ghc/  :? for help
3  λ> fibs = 1 : 1 : (zipWith (+) fibs (tail fibs))
4  λ> take 100 $ fibs
5      [1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597,
6        2584, 4181, 6765, 10946, 17711, 28657, 46368, 75025, 121393, 196418,
7        317811, 514229, 832040, 1346269, 2178309, 3524578, 5702887, 9227465,
8        14930352, 24157817, 39088169, 63245986, 102334155, 165580141, 267914296,
9        433494437, 701408733, 1134903170, 1836311903, 2971215073, 4807526976,
10       7778742049, 12586269025, 20365011074, 32951280099, 53316291173,
11       86267571272, 139583862445, 225851433717, 365435296162, 591286729879,
12       956722026041, 1548008755920, 2504730781961, 4052739537881, 6557470319842,
13       10610209857723, 17167680177565, 27777890035288, 44945570212853,
14       72723460248141, 117669030460994, 190392490709135, 308061521170129,
15       498454011879264, 806515533049393, 1304969544928657, 2111485077978050,
16       3416454622906707, 5527939700884757, 8944394323791464, 14472334024676221,
17       23416728348467685, 37889062373143906, 61305790721611591,
18       99194853094755497, 160500643816367088, 259695496911122585,
19       420196140727489673, 679891637638612258, 1100087778366101931,
20       1779979416004714189, 2880067194370816120, 4660046610375530309,
21       7540113804746346429, 12200160415121876738, 19740274219868223167,
22       31940434634990099905, 51680708854858323072, 83621143489848422977,
23       135301852344706746049, 218922995834555169026, 354224848179261915075]
24  λ> head $ drop 99 fibs
25      354224848179261915075
```

其他解法

```
1  λ> print $ take 140 $ zip [1..] fibs
2  [(1,1),(2,1),(3,2),(4,3),(5,5),(6,8),(7,13),(8,21),(9,34),(10,55),(11,89),(12,144),
3  (13,233),(14,377),(15,610),(16,987),(17,1597),(18,2584),(19,4181),(20,6765),(21,10946),
4  (22,17711),(23,28657),(24,46368),(25,75025),(26,121393),(27,196418),(28,317811),(29,514229),
5  (30,832040),(31,1346269),(32,2178309),(33,3524578),(34,5702887),(35,9227465),(36,14930352),
6  (37,24157817),(38,39088169),(39,63245986),(40,102334155),(41,165580141),(42,267914296),(43,433494437),
7  (44,701408733),(45,1134903170),(46,1836311903),(47,2971215073),(48,4807526976),(49,7778742049),
8  (50,12586269025),(51,20365011074),(52,32951280099),(53,53316291173),(54,86267571272),(55,139583862445),
9  (56,225851433717),(57,365435296162),(58,591286729879),(59,956722026041),(60,1548008755920),
10 (61,2504730781961),(62,4052739537881),(63,6557470319842),(64,10610209857723),(65,17167680177565),
11 (66,27777890035288),(67,44945570212853),(68,72723460248141),(69,117669030460994),(70,190392490709135),
12 (71,308061521170129),(72,498454011879264),(73,806515533049393),(74,1304969544928657),(75,2111485077978050),
13 (76,3416454622906707),(77,5527939700884757),(78,8944394323791464),(79,14472334024676221),(80,23416728348467685),
14 (81,37889062373143906),(82,61305790721611591),(83,99194853094755497),(84,160500643816367088),
15 (85,259695496911122585),(86,420196140727489673),(87,679891637638612258),(88,1100087778366101931),
16 (89,1779979416004714189),(90,2880067194370816120),(91,4660046610375530309),(92,7540113804746346429),
17 (93,12200160415121876738),(94,19740274219868223167),(95,31940434634990099905),(96,51680708854858323072),
18 (97,83621143489848422977),(98,135301852344706746049),(99,218922995834555169026),(100,354224848179261915075),
19 (101,573147844013817084101),(102,927372692193078999176),(103,1500520536206896083277),(104,2427893228399975082453),
20 (105,3928413764606871165730),(106,6356306993006846248183),(107,10284720757613717413913),(108,16641027750620563662096),
21 (109,26925748508234281076009),(110,43566776258854844738105),(111,70492524767089125814114),(112,114059301025943970552219),
22 (113,184551825793033096366333),(114,298611126818977066918552),(115,483162952612010163284885),(116,781774079430987230203437),
23 (117,1264937032042997393488322),(118,2046711111473984623691759),(119,3311648143516982017180081),
24 (120,535835925499096640871840),(121,8670007398507948658051921),(122,14028366653498915298923761),
25 (123,22698374052006863956975682),(124,36726740705505779255899443),(125,59425114757512643212875125),
26 (126,96151855463018422468774568),(127,155576970220531065681649693),(128,251728825683549488150424261),
27 (129,407305795904080553832073954),(130,659034621587630041982498215),(131,1066340417491710595814572169),
28 (132,1725375039079340637797070384),(133,2791715456571051233611642553),(134,4517090495650391871408712937),
29 (135,7308805952221443105020355490),(136,11825896447871834976429068427),(137,19134702400093278081449423917),
30 (138,30960598847965113057878492344),(139,50095301248058391139327916261),(140,81055900096023504197206408605)]
```

其他解法

Scala, 忠诚的 Haskell 追随者:

```
1  val fibs: LazyList[BigInt] = BigInt(1) #:: BigInt(1) #:: fibs.zip(fibs.tail).map(x => x._1 + x._2)

1  ~ amm
2  Loading...
3  Welcome to the Ammonite Repl 2.0.4 (Scala 2.13.1 Java 1.8.0_242)
4  @ val fibs: LazyList[BigInt] = BigInt(1) #:: BigInt(1) #:: fibs.zip(fibs.tail).map(x => x._1 + x._2)
5  fibs: LazyList[BigInt] = LazyList(...)
6  @ fibs(99)
7  res1: BigInt = 354224848179261915075
8
9  @ fibs.take(10)
10 res2: LazyList[BigInt] = LazyList(1, 1, 2, 3, 5, 8, 13, 21, 34, 55)
```

其他解法

$$\begin{array}{lcl} fibs & = & 1, \quad 1, \quad 2, \quad 3, \quad 5, \quad 8... \\ (tail \ fibs) & = & 1, \quad 2, \quad 3, \quad 5, \quad 8, \quad 13... \\ fibs' & = & 1 + 1, \quad 1 + 2, \quad 2 + 3, \quad 3 + 5, \quad 5 + 8... \end{array}$$

第 2 问, 易如反掌

```
1 fibs
2 .zipWithIndex
3 .dropWhile(_._1.toString.length < 1000)
4 .head
5 ._2 + 1
```

```
1 ~ amm
2 Loading...
3 Welcome to the Ammonite Repl 2.0.4 (Scala 2.13.1 Java 1.8.0_242)
4 @ val fibs: LazyList[BigInt] = BigInt(1) #:: BigInt(1) #:: fibs.zip(fibs.tail).map(x => x._1 + x._2)
5 fibs: LazyList[BigInt] = LazyList(...)
6 @ fibs.zipWithIndex.dropWhile(_._1.toString.length < 1000).head._2 + 1
7 res1: Int = 4782
```

第 3 问，快速矩阵幂？

快速矩阵幂？杀鸡用牛刀。

暴力尾递归，你值得拥有：

```
1 def fib(n: Int) = {  
2   @scala.annotation.tailrec  
3   def f(x: Int, y: Int, i: Int): Int =  
4     if (i == 0) x else f(y, (x + y) % 1000000007, i - 1)  
5   f(0, 1, n)  
6 }  
  
1 Scala >> val t = System.currentTimeMillis ; fib(1000000000) ; println(s"time escape: ${System.currentTimeMillis-t}ms")  
2 time escape: 2519ms  
3 t: Long = 1589459096621L  
4 res10_1: Int = 21
```

不到 3 秒！

第3问, C++

C++ 暴力循环:

```
1  #include <iostream>
2
3  int fib(int n) {
4      int a = 0, b = 1, c = 0;
5      for (int i = 0 ; i < n ; i++) {
6          c = b;
7          b = (a + b) % 1000000007;
8          a = c;
9      }
10     return a;
11 }
12
13 int main() {
14     std::cout << fib(1000000000) << "\n";
15     return 0;
16 }
```

```
1  ~ $ g++ -O2 fibonacci.cc -o fibonacci-cpp
2  ~ $ time ./fibonacci-cpp
3  21
4  ./fibonacci-cpp  2.47s user 0.00s system 99% cpu 2.473 total
```

2.47 秒!

第 3 问, Rust

Rust 尾递归:

```
1 fn main() {  
2     fn fib_m(n: i32) -> i32 {  
3         fn _fib(x: i32, y: i32, i: i32, n: i32) -> i32 {  
4             if i == n { x } else { _fib(y, (x+y) % 1000000007, i+1, n) }  
5         }  
6         _fib(0, 1, 0, n)  
7     }  
8     println!("fibM(10^9) = {}", fib_m(1_000_000_000));  
9 }
```

```
1 ~ $ time ./fibonacci_rs  
2 fibM(10^9) = 21  
3 ./fibonacci_rs 2.53s user 0.00s system 99% cpu 2.500 total
```

2.50 秒!

第 3 问, JavaScript

JavaScript 循环:

```
1 function fibM(n) {  
2     for (i = 0, x = 0, y = 1; i < n; i++) {  
3         var z = y  
4         y = (x + y) % 1000000007  
5         x = z  
6     }  
7     return x  
8 }  
9  
10 console.log(fibM(1000000000))
```

```
1 ~ $ time node fibonacci-loop.js  
2 21  
3 node fibonacci-loop.js 4.08s user 0.01s system 99% cpu 4.088 total
```

4.1 秒!

第 3 问, JavaScript

JavaScript 尾递归:

```
1  'use strict';
2
3  function fibM(n) {
4      let f = (x, y, i) => (i === 0) ? x : f(y, (x+y) % 1000000007, i-1)
5      return f(0, 1, n)
6  }
7
8  console.log(fibM(1000000000)) // RangeError: Maximum call stack size exceeded
```

栈溢出!

第 3 问, Python

Python 循环:

```
1  def fib(n):
2      a, b = 0, 1
3      i = 1
4      while (i <= n):
5          a, b = b, (a + b) % 1000000007
6          i = i + 1
7      return a
8
9  print(fib(1000000000))
```

第 3 问, Python

Python 循环测试

```
1 ~ $ time python2.7 fibonacci-loop.py
2 21
3 python2.7 fibonacci-loop.py 57.04s user 0.01s system 99% cpu 57.103 total
4 ~ $ time python fibonacci-loop.py
5 21
6 python3.7 fibonacci-loop.py 108.94s user 0.02s system 99% cpu 1:49.00 total
7 ~ $ time python3.8 fibonacci-loop.py
8 21
9 python3.8 fibonacci-loop.py 112.31s user 0.01s system 99% cpu 1:52.33 total
10 ~ $ time python3.9 fibonacci-loop.py
11 21
12 python3.9 fibonacci-loop.py 120.74s user 0.01s system 99% cpu 2:00.81 total
13 ~ $ time pypy fibonacci-loop.py
14 21
15 pypy fibonacci-loop.py 3.11s user 0.01s system 99% cpu 3.134 total
16 ~ $ time pypy3 fibonacci-loop.py
17 21
18 pypy3 fibonacci-loop.py 3.13s user 0.01s system 99% cpu 3.150 total
```

发生了什么?

第 3 问, Python

Python 尾递归:

```
1 def fib(n):  
2     def f(x, y, i): return x if not i else f(y, (x + y) % 1000000007, i-1)  
3     return f(0, 1, n)  
4  
5 print(fib(1000000000)) # RecursionError: maximum recursion depth exceeded
```

栈溢出!

快速矩阵幂

$$\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \times \begin{bmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{bmatrix} = \begin{bmatrix} a_{11}b_{11} + a_{12}b_{21} & a_{11}b_{12} + a_{12}b_{22} \\ a_{21}b_{11} + a_{22}b_{21} & a_{21}b_{12} + a_{22}b_{22} \end{bmatrix}$$

$$\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \times \begin{bmatrix} b_{11} \\ b_{21} \end{bmatrix} = \begin{bmatrix} a_{11}b_{11} + a_{12}b_{21} \\ a_{21}b_{11} + a_{22}b_{21} \end{bmatrix}$$

快速矩阵幂

$$\begin{aligned}\begin{bmatrix} F_{n+1} \\ F_n \end{bmatrix} &= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \times \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} \\ &= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^2 \times \begin{bmatrix} F_{n-1} \\ F_{n-2} \end{bmatrix} \\ &= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^3 \times \begin{bmatrix} F_{n-2} \\ F_{n-3} \end{bmatrix} \\ &= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n \times \begin{bmatrix} F_2 \\ F_1 \end{bmatrix}\end{aligned}$$

$$\begin{bmatrix} F_{n-1} & F_n \\ F_n & F_{n+1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n = A^n$$

快速矩阵幂

```
1  val Q = 1000000007L
2
3  final case class Matrix(a11: Long, a12: Long, a21: Long, a22: Long) {
4    /**
5     * 取模矩阵乘法
6     * @param x 被乘数
7     * @param m 模
8     * @return
9     */
10   def *(x: Matrix, m: Long = Q): Matrix =
11     Matrix(a11 = (a11 * x.a11 + a12 * x.a21) % m,
12            a12 = (a11 * x.a12 + a12 * x.a22) % m,
13            a21 = (a21 * x.a11 + a22 * x.a21) % m,
14            a22 = (a21 * x.a12 + a22 * x.a22) % m)
15   /**
16    * 矩阵幂 (取模), m = Q
17    * @param p 幂
18    * @return
19    */
20   def ^(p: Int): Matrix = {
21     require(p > 0, "parameter `p` must be greater than 0")
22     Iterator.continually(this).take(p).reduceLeft(_ * _)
23   }
24 }
```

快速矩阵幂

```
1  val Q = 1000000007L
2
3  final case class Matrix(a11: Long, a12: Long, a21: Long, a22: Long) {
4    /**
5     * 取模矩阵乘法
6     * @param x 被乘数
7     * @param m 模
8     * @return
9     */
10   def *(x: Matrix, m: Long = Q): Matrix =
11     Matrix(a11 = (a11 * x.a11 + a12 * x.a21) % m,
12            a12 = (a11 * x.a12 + a12 * x.a22) % m,
13            a21 = (a21 * x.a11 + a22 * x.a21) % m,
14            a22 = (a21 * x.a12 + a22 * x.a22) % m)
15   /**
16    * 矩阵幂 (取模), m = Q
17    * @param p 幂
18    * @return
19    */
20   def ^(p: Int): Matrix = {
21     require(p > 0, "parameter `p` must be greater than 0")
22     Iterator.continually(this).take(p).reduceLeft(_ * _)
23   }
24 }
```

快速矩阵幂

```
1  val Q = 1000000007L
2
3  final case class Matrix(a11: Long, a12: Long, a21: Long, a22: Long) {
4    def fastPow(n: Int): Matrix = {
5      @scala.annotation.tailrec
6      def f(z: Matrix, x: Matrix, ys: List[Char]): Matrix =
7        ys match {
8          case '0' :: xs => f(z * z, x, xs)
9          case '1' :: xs => f(z * z, x * z, xs)
10         case _      => x
11       }
12      val ZERO = Matrix(1, 0, 0, 1)
13      f(A, ZERO, n.toBinaryString.reverse.toList)
14    }
15
16    def ^(p: Int): Matrix = fastPow(p)
17  }
18
19  def fib(n: Int) = (A^(n)).a12
```

快速矩阵幂

```
1  Scala >> def fib(n: Int) = A.fastPow(n).a12
2  defined function fib
3
4  Scala >> val t = System.currentTimeMillis ; fib(1000000000) ; println(s"time escape: ${System.currentTimeMillis-t}ms")
5  time escape: 0ms
6  t: Long = 1589518862649L
7  res11_1: Long = 21L
8
9  Scala >>
```

不到 1ms !

快速矩阵幂

```
1  object CatsMonoidFib extends App {
2
3      import cats._
4
5      val Q = 1000000007
6
7      final case class Matrix(a11: Long, a12: Long, a21: Long, a22: Long)
8
9      class MatrixMultiplicationMonoid(m: Int) extends Monoid[Matrix] {
10
11          override def empty: Matrix = Matrix(1, 0, 0, 1)
12
13          override def combine(x: Matrix, y: Matrix): Matrix =
14              Matrix(a11 = (x.a11 * y.a11 + x.a12 * y.a21) % m,
15                    a12 = (x.a11 * y.a12 + x.a12 * y.a22) % m,
16                    a21 = (x.a21 * y.a11 + x.a22 * y.a21) % m,
17                    a22 = (x.a21 * y.a12 + x.a22 * y.a22) % m)
18      }
19
20      implicit val monoid: MatrixMultiplicationMonoid = new MatrixMultiplicationMonoid(Q)
21
22      val A: Matrix = Matrix(0, 1, 1, 1)
23
24      def fib(n: Int): Long = monoid.combineN(A, n).a12
25
26      val logger: Logger = Logger(getClass)
27
28      logger.info("fib(10^9) = {}", fib(1000000000))
29  }
```

超级快速矩阵幂

$$A^{p+q} = A^p \cdot A^q$$

$$A^{p \cdot q} = (A^p)^q$$

$$\begin{aligned} A^{p^q} &= A^{p^{q-1} \cdot p} \\ &= \left(A^{p^{q-1}} \right)^p \\ &= \left(\left(A^{p^{q-2}} \right)^p \right)^p \\ &= \left(\left(\left(A^{p^{q-3}} \right)^p \right)^p \right)^p \\ &= \underbrace{\left(\left(\left(A^p \right)^p \dots \right)^p \right)^p}_q \end{aligned}$$

超级快速矩阵幂

```
1  val Q = 1000000007L
2
3  final case class Matrix(a11: Long, a12: Long, a21: Long, a22: Long) {
4    /**
5     * 超级矩阵幂 (取模),  $M^{(p^q)}$ ,  $m = Q$ 
6     * @param p 幂
7     * @param q 超幂
8     * @return
9     */
10   def ^^ (p: Int, q: Int): Matrix = {
11     require(p > 0, "parameter `p` must be greater than 0")
12     require(q > 0, "parameter `q` must be greater than 0")
13     (1 until q).foldLeft(this ^ p)((z, _) => z ^ p)
14   }
15 }
16
17 def fibPQ(p: Int, q: Int) = (A ^^ (p, q)).a12
```

超级快速矩阵幂

```
1  def fibPQ(p: Int, q: Int) = (A^(p, q)).a12
2
3  val t = System.currentTimeMillis
4
5  (1 to 1000).foreach { _ =>
6    val _ = fibPQ(10, 1000)
7  }
8
9  println(s"""F(10^1000) = fibPQ(10, 1000) = ${fibPQ(10, 1000)}""")
10 println(s"""F(10^9)    = fibPQ(10, 9)    = ${fibPQ(10, 9)}""")
11 println(s"time escape: ${System.currentTimeMillis - t}ms")
```

```
1  F(10^1000) = fibPQ(10, 1000) = 552179166
2  F(10^9)    = fibPQ(10, 9)    = 21
3  time escape: 527ms
```

思考题

快速幂能做到多快

问题 1: Euler Project 122(<https://projecteuler.net/problem=122>)

The most naive way of computing n^{15} requires fourteen multiplications:

$$n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n \cdot n = n^{15}$$

But using a “binary” method you can compute it in 6 multiplications:

$$n \cdot n = n^2$$

$$n^2 \cdot n^2 = n^4$$

$$n^4 \cdot n^4 = n^8$$

$$n^8 \cdot n^4 = n^{12}$$

$$n^{12} \cdot n^2 = n^{14}$$

$$n^{14} \cdot n = n^{15}$$

快速幂能做到多快

问题 1: Euler Project 122(<https://projecteuler.net/problem=122>)

However it is yet possible to compute it in only **5** multiplications:

$$n \cdot n = n^2$$

$$n^2 \cdot n = n^3$$

$$n^3 \cdot n^3 = n^6$$

$$n^6 \cdot n^6 = n^{12}$$

$$n^{12} \cdot n^3 = n^{15}$$

We shall define $m(k)$ to be the minimum number of multiplications to compute n^k ; for example $m(15) = 5$.

For $1 \leq k \leq 200$, find $\sum m(k)$.

不重复尾数的个数

问题 2:

我们用 G_n 来表示那菲波契数列每一项截取最后 n 位数字组成的新数列。比如,

$$G_1 = 0, 1, 1, 2, 3, 5, 8, 3, 1, 4, 5, 9, 4, 3, 7, 0, 7, 7, 4, 1, 5, 6, 1, 7, 8 \dots$$

$$G_2 = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 44, 33, 77, 10, 87, 97, 84 \dots$$

$$G_3 = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987 \dots$$

已知 G_1 中一共有 10 个不重复的元素, G_2 中一共有 100 个不重复的元素, 而 G_3 中一共只有 750 个不重复的元素。

记 G_n 中所有不重复元素的个数为 $g(n)$. 试计算 $g(4)$, $g(5)$, $g(6)$ 和 $g(7)$.

完。

谢谢大家！