LU-15: USING FIRST ORDER LOGIC

LU Objectives To explain assertions and queries in FOL LU Outcomes CO:3 Represent knowledge base using assertions

Assertions and Queries

- Sentences are added to a knowledge base using TELL, exactly as in propositional logic. Such sentences are called assertions.
- For Ex: we can assert that

John is a king, Richard is a person, and all kings are persons:

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TELL(KB, King(John)).
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TELL(KB, Person(Richard)).

TELL(KB, \forall x King(x) \Rightarrow Person(x)).

Assertions and Queries

- We can ask questions of the knowledge base using ASK.
- For Ex,
 ASK(KB, King(John))
 returns true.
- Questions asked with ASK are called queries or goals.

The kinship domain

- It is the domain of family relationships
- The objects in our domain are people
- There are two unary predicates, Male and Female.
- Kinship relations are parenthood, brotherhood, marriage, and so on. They are represented by binary predicates: Parent, Sibling, Brother, Sister, Child, Daughter, Son, Spouse, Wife, Husband, Grandparent, Grandchild, Cousin, Aunt, and Uncle
- We use functions for Mother and Father, because every person has exactly one of each of these

The kinship domain

- For example, one's mother is one's female parent:

 ∀ m, c Mother (c)=m ⇔ Female(m) ∧ Parent(m, c).
- One's husband is one's male spouse:
 ∀ w, h Husband(h,w) ⇔ Male(h) ∧ Spouse(h,w).
- Male and female are disjoint categories: $\forall x \, \text{Male}(x) \Leftrightarrow \neg \text{Female}(x)$.
- Parent and child are inverse relations:
 ∀ p, c Parent(p, c) ⇔ Child (c, p) .
- A grandparent is a parent of one's parent:
 ∀ g, c Grandparent (g, c) ⇔ ∃p Parent(g, p) ∧ Parent(p, c)
- A sibling is another child of one's parents:
 ∀ x, y Sibling(x, y) ⇔ x = y ∧ ∃p Parent(p, x) ∧ Parent(p, y)

Axioms and Theorems

- Each of the sentences in previous slide can be viewed as an axiom of the kinship domain.
- Axioms are commonly associated with purely mathematical domains. But they are needed in all domains.
- They provide the basic factual information from which useful conclusions can be derived.
 Our kinship axioms are also definitions

Axioms and Theorems

- The axioms define the Mother function and the Husband, Male, Parent, Grandparent, and Sibling predicates in terms of other predicates.
- Our definitions "bottom out" at a basic set of predicates (Child, Spouse, and Female) in terms of which the others are ultimately defined.
- There is not necessarily a unique set of primitive predicates;
- we could equally well have used Parent, Spouse, and Male. In some domains, there is no clearly identifiable basic set.

Axioms and Theorems

- Some logical sentences are theorems—that is,
 they are entailed by the axioms.
- For example, consider the assertion that siblinghood is symmetric:

 \forall x, y Sibling(x, y) \Leftrightarrow Sibling(y, x)

- In fact, it is a theorem that follows logically from the axiom that defines siblinghood.
- If we ASK the knowledge base this sentence, it should return true

Why need theorems?

- From a purely logical point of view, a knowledge base need contain only axioms and no theorems, because the theorems do not increase the set of conclusions that follow from the knowledge base.
- But from a practical point of view, theorems are essential to reduce the computational cost of deriving new sentences.
- Without theorems, a reasoning system has to start from first principles every time.

- Numbers are perhaps the most vivid example of how a large theory can be built up from
- a tiny kernel of axioms.
- Consider the theory of natural numbers or non-negative integers.
- It uses a predicate NatNum that will be true of natural numbers;
- we need one constant symbol, 0; and we need one function symbol, S (successor).

Natural numbers are defined recursively:

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NatNum(0).
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 \forall n NatNum(n) \Rightarrow NatNum(S(n)).

We also need axioms to constrain the successor function:

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\forall n 0 \neq S(n).
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$$\forall$$
 m, n m \neq n \Rightarrow S(m) \neq S(n).

 We can define addition in terms of the successor function:

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\forall m \ NatNum(m) \Rightarrow + (0,m) = m.
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$$\forall$$
 m, n NatNum(m) \land NatNum(n) \Rightarrow + (S(m), n) = S(+(m, n))

- To make sentences about numbers easier to read, we allow the use of infix notation.
- We can also write S(n) as n+ 1, so the second axiom becomes
 - \forall m, n NatNum(m) \land NatNum(n) \Rightarrow (m + 1) + n = (m + n) + 1.
- This axiom reduces addition to repeated application of the successor function.
- The use of infix notation is an example of syntactic sugar.

- Similarly we can define multiplication as repeated addition, exponentiation as repeated multiplication, integer division and remainders, prime numbers, and so on.
- Thus, the whole of number theory can be built up from one constant, one function, one predicate and four axioms

Sets

The set domain representations using FOL or Predicate Logic:

- The empty set is a constant written as {}.
- Unary Predicate : Set
- Binary Predicate: x ∈ s ((x is a member of set s), s₁ ⊆ s₂ (set s1 is a subset of set s2)
- Binary function: $s_1 \cap s_2$ (the intersection of two sets), $s_1 \cup s_2$ (the union of two sets), $\{x \mid s_2\}$) (Set resulting from adjoining element x to s)

Sets - Axioms

1. The only sets are the empty set and those made by adjoining something to a set:

$$\forall$$
 s Set(s) \Leftrightarrow (s={}) \lor (\exists x, s2 Set(s2) \land s={x|s2}).

2. The empty set has no elements adjoined into it. In other words, there is no way to decompose {} into a smaller set and an element:

$$\neg \exists x, s \{x \mid s\} = \{\}.$$

3. Adjoining an element already in the set has no effect:

$$\forall x, s x \in s \Leftrightarrow s = \{x \mid s\}$$
.

Sets - Axioms

4. The only members of a set are the elements that were adjoined into it.

$$\forall x, s x \in s \Leftrightarrow \exists y, s2 (s=\{y \mid s2\} \land (x=y \lor x \in s2))$$
.

5. A set is a subset of another set if and only if all of the first set's members are members of the second set:

$$\forall$$
 s1, s2 s1 \subseteq s2 \Leftrightarrow (\forall x x \in s1 \Rightarrow x \in s2).

6. Two sets are equal if and only if each is a subset of the other:

$$\forall$$
 s1, s2 (s1 =s2) \Leftrightarrow (s1 \subseteq s2 \land s2 \subseteq s1)

Sets - Axioms

7. An object is in the intersection of two sets if and only if it is a member of both sets:

$$\forall$$
 x, s1, s2 x \in (s1 \cap s2) \Leftrightarrow (x \in s1 \land x \in s2).

8. An object is in the union of two sets if and only if it is a member of either set:

$$\forall$$
 x, s1, s2 x \in (s1 \cup s2) \Leftrightarrow (x \in s1 \vee x \in s2).

Lists

- Lists are similar to sets.
- The differences are LIST that lists are ordered and the same element can appear more than once in a list.
- Nil is the constant list with no elements;
- Cons, Append, First, and Rest are functions; and Find is the predicate that does for lists what Member does for sets.
- List? is a predicate that is true only of lists.
- The empty list is [].
- The term Cons(x, y), where y is a nonempty list, is written [x|y]. The term Cons(x, Nil) (i.e., the list containing the element x) is written as [x].
- A list of several elements, such as [A,B,C], corresponds to the nested term Cons(A, Cons(B, Cons(C, Nil))).

- The first order axioms in this section are much more concise than prepositional logic, capturing in a natural world
- WKT the wumpus agent receives a percept vector with five elements. The corresponding first-order sentence stored in the knowledge base must include both the percept and the time at which it occurred; otherwise, the agent will get confused about when it saw what.

- Typical percept sentence would be
 Percept ([Stench, Breeze, Glitter, None, None], 5).
- Here, Percept is a binary predicate, and Stench and so on are constants placed in a list.
- The actions in the wumpus world can be represented by logical terms:

Turn(Right), Turn(Left), Forward, Shoot, Grab, Climb.

 To determine which is best, the agent program executes the query

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ASKVARS(∃ a BestAction(a, 5)), which returns a binding list such as {a/Grab}.
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- The raw percept data implies certain facts about the current state.
- For Ex:

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\forall t, s, g, m, c Percept ([s, Breeze, g,m, c], t) \Rightarrow Breeze(t), \forall t, s, b, m, c Percept ([s, b, Glitter,m, c], t) \Rightarrow Glitter (t),
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 These rules exhibit a trivial form of the reasoning process called perception.

- Simple "reflex" behavior can also be implemented by quantified implication sentences.
- For example,
 ∀ t Glitter (t) ⇒ BestAction(Grab, t) .
- Objects in the environment are squares, pits, and the wumpus
- Adjacency of any two squares can be defined as
 ∀ x, y, a, b Adjacent ([x, y], [a, b]) ⇔
 (x = a ∧ (y = b 1 ∨ y = b + 1)) ∨ (y = b ∧ (x = a 1 ∨ x = a + 1)).

- A unary predicate Pit is true for squares containing pit
- Since there is exactly one wumpus, a constant Wumpus is just as good as a unary predicates.
- The agent's location changes over time, so we write At(Agent, s, t) to mean that the agent is at square s at time t.
- We can then say that objects can only be at one location at a time:

 \forall x, s1, s2, t At(x, s1, t) \land At(x, s2, t) \Rightarrow s1 = s2.

Diagnostic Rules - The wumpus world

- These rules generate hidden causes from observed effects. They help to reduce hidden facts in the world.
- Diagnostic Rule for finding pit
 - 1. If square is breezy the agent can deduce some adjacent square must contain pit
 - \forall s Breezy(s) \Rightarrow \exists r Adjacent (r, s) \land Pit(r).
 - 2. If square is not breezy the agent can deduce no adjacent square contains a pit
 - \forall s \neg Breezy(s) $\Rightarrow \neg \exists r \text{ Adjacent } (r, s) \land \text{Pit}(r)$.

Diagnostic Rules - The wumpus world

Combining two rules we obtain the biconditional sentence

 \forall s Breezy(s) \Leftrightarrow \exists r Adjacent (r, s) \land Pit(r).

Causal Rules - The wumpus world

- Causal rules reflect the assumed direction of causality in the world. That is these rules are used to infer effect from cause.
- Some hidden property of the world causes certain percepts to be generated.
- Ex: 1) a pit causes all adjacent squares to be breezy \forall r Pit(r) \Rightarrow [\forall s Adjacent (r, s) \Rightarrow Breezy(s)]
- 2) If all squares adjacent to a given square are pitless, the square will not be breezy
- $\forall s[\forall r \ Adjacent (r, s) \Rightarrow \neg Pit(r)] \Rightarrow \neg Breezy(s)]$

Quantification- The wumpus world

 In first-order logic we can quantify over time. For example, the axiom for the arrow becomes

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\forall t HaveArrow(t + 1) \Leftrightarrow (HaveArrow(t) 
 \land \negAction(Shoot, t)).
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