CS8792 Cryptography and Network Security Basic Concepts in Number Theory

Unit-II

Prime Numbers

Euler's Theorem

Primality Testing

Chinese Remainde

Discrete

# CS8792 Cryptography and Network Security

Basic Concepts in Number Theory

## Session Objectives

CS8792 Cryptography and Network Security Basic Concepts in Number Theory

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Prime Numbers

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Primality Testing

Chinese Remainde

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- To learn about prime numbers
- To check a number is prime or not
- To learn Chinese remainder theorem

## Agenda

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- 1 Prime Numbers
- 2 Euler's Theorem
- 3 Primality Testing
- 4 Chinese Remainder Theorem
- 5 Discrete Logarithms

### Presentation Outline

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#### Prime Numbers and Factorization

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- prime numbers only have divisors of 1 and self
- to factor a number n is to write it as a product of other numbers: n=a x b x c
- note that factoring a number is relatively hard compared to multiplying the factors together to generate the number
- the prime factorisation of a number n is when its written as a product of primes eg.  $91=7\times13$ ;  $3600=2^4\times3^2\times5^2$

$$a=\prod_{p\in P}P^{a_p}$$

■ Relatively Prime Numbers: two numbers a, b are relatively prime if they have no common divisors except 1

#### Fermat's Theorem

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#### Fermat's Theorem:

$$a^{p-1} mod p = 1$$

- where p is prime and gcd(a,p)=1 also known as Fermat's Little Theorem
- useful in public key and primality testing

## Euler Totient Function $\phi(n)$

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- when doing arithmetic modulo n
- **complete set of residues** is: 0..n-1
- reduced set of residues is those numbers (residues) which are relatively prime to n
- eg for n=10, complete set of residues is 0,1,2,3,4,5,6,7,8,9 reduced set of residues is 1,3,7,9
- number of elements in reduced set of residues is called the Euler Totient Function  $\phi(n)$

## Euler Totient Function $\phi(n)$

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Discrete Logarithms ■ in general need prime factorization, but

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### Euler's Theorem

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Discrete Logarithm ullet a generalisation of Fermat's Theorem  $a^{\phi(n)} ullet$  mod old N=1

- where gcd(a,N)=1
- eg. a=3; n=10;  $\phi$  (10)=4;
- hence  $3^4 = 81 = 1 \mod 10$
- $a=2; n=11; \phi(11)=10;$
- hence  $2^{10} = 1024 = 1 \mod 11$

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## **Primality Testing**

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- often need to find large prime numbers
- traditionally sieve using trial division
- ie. divide by all numbers (primes) in turn less than the square root of the number
- only works for small numbers
- alternatively can use statistical primality tests based on properties of primes
- for which all primes numbers satisfy property
- but some composite numbers, called pseudo-primes, also satisfy the property

## Miller Rabin Algorithm

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Discrete Logarithms a test based on Fermat's Theorem

algorithm is:

TEST (n) is:

1 Find integers **k**, **m**, **k**> **0**, **m** odd, so that  $(n-1)=2^k$ .m

2 Select a random integer a, 1 < a < n-1

**3** if  $a^m \mod n = 1$  then return ("maybe prime");

**4** for j = 0 to k - 1 do

**5** if  $(a^{2^{j}m} \mod n = n-1)$  then return(" maybe prime")

6 return ("composite")

## Miller Rabin Algorithm

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Discrete Logarithms Is 561 prime?

- **1** Find 561  $1 = 2^k$ . m
- **2** Choose a, 1 < a < n 1
- 3 Compute  $b_0 = a^m \mod n$
- 4 if  $b_0 = +1 \implies$  n is a composite number else if  $b_0 = -1 \implies$  n may be a prime number
- **5** Compute  $b_i = b_{i-1}^2$  , check for composite or prime
- 6 Repeat step number 5

## Miller Rabin Algorithm

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■ 
$$561-1=2^k$$
.m  
 $\frac{560}{2^2}=280$ ;  $\frac{560}{2^3}=140$ ;  $\frac{560}{2^3}=70$ ;  $\frac{560}{2^4}=35$ ;  $\frac{560}{2^5}=17.5$ 

$$\bullet$$
 560=  $2^4.35$ ; k=4; m=35

$$b_0 = 2^{35} \mod 561 = 263$$

3 Is 
$$b_0 = \pm 1 \mod 561$$

$$b_1 = b_0^2 = 263^2 \mod 561 = 67$$

$$b_3 = 67^2 \mod 561 = 1$$

561 is a composite number

Solve: Is 53 a prime number?

#### Probabilistic Considerations

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- if Miller-Rabin returns "composite" the number is definitely not prime
- otherwise is a prime or a pseudo-prime
- chance it detects a pseudo-prime is < 1/4
- hence if repeat test with different random a then chance n is prime after t tests is:
- Pr(n prime after t tests) =  $1 4^{-t}$ eg. for t=10 this probability is > 0.99999

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Discrete Logarithms Chinese Remainder Theorem: If  $m_1, m_2, ..., m_k$  are pairwise relatively prime positive integers, and if  $a_1, a_2, ..., a_k$  are any integers, then the simultaneous congruences

$$x \equiv a_1 \pmod{m_1}$$
,  $x \equiv a_2 \pmod{m_2}$ , ...,  $x \equiv a_k \pmod{m_k}$ 

have a solution, and the solution is unique modulo m, where

$$m = m_1 m_2 \cdots m_k$$
.

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Discrete Logarithm

#### To compute $X \pmod{M}$

- first compute all  $a_i = A \mod m_i$  separately
- determine constants  $c_i$ , where  $M_i = M/mi$
- then combine results to get answer using:

$$X \equiv (\sum_{i=1}^k a_i c_i) \mod M$$

$$c_i = M_i \times (M_i^{-1} \mod m_i)$$
 for  $1 \le i \le k$ 

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Discrete Logarithms What's x such that:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x\equiv 2 \; (mod \; 7)?$$

$$X \equiv (\sum_{i=1}^{k} a_i c_i) \mod M$$
;  $c_i = M_i \times (M_i^{-1} \mod m_i)$ 

$$X = a_1.M_1.M_1^{-1} + a_2.M_2.M_2^{-1} + a_3.M_3.M_3^{-1} \mod M$$

$$M_1.M_1^{-1} \equiv 1 \mod m_1$$

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Discrete Logarithm Using the Chinese Remainder theorem:

$$a_1 = 2$$
;  $a_2 = 3$ ;  $a_3 = 2$ ;  $m_1 = 3$ ;  $m_2 = 5$ ;  $m_3 = 7$ ;

$$M = m_1 \times m_2 \times m_3 = 3 \times 5 \times 7 = 105$$

$$M_1 = M/m_1 = 105/3 = 35$$

• 2 is an inverse of 
$$M_1 = 35 \pmod{3}$$
  
(since  $35 \times 2 \equiv 1 \pmod{3}$ 

• 
$$M_1.M_1^{-1} \equiv 1 \mod m_1 \implies 35. M_1^{-1} \equiv 1 \mod 3$$

$$\blacksquare$$
 gcd(35,3);gcd(3,2);gcd(2,1); gcd(1,0)= 1

■ 
$$35 = 11 \times 3 + 2 \implies 2 = 35 - 11 \times 3$$

$$3 = 1 \times 2 + 1 \implies 1 = 3 - 1 \times 2$$

$$1 = 3 - 1 \times 2$$
= 3 - (35 - 11 \times 3) = -1 \times 35 + 12 \times 3

■ 
$$1 = -1 \times 35 + 12 \times 3$$
;  $-1 \times 35 \equiv 1 \mod 3$   
⇒  $2 \times 35 \equiv 1 \mod 3$ ; **2** is inverse of **35 mod 3**

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Discrete Logarithms Using the Chinese Remainder theorem:

- $M_2 = M/m_2 = 105/5 = 21$ 
  - 1 is an inverse of  $M_2 = 21 \pmod{5}$  (since  $21 \times 1 \equiv 1 \pmod{5}$ )
- $M_3 = M/m_3 = 105/7 = 15$ 
  - 1 is an inverse of  $M_3 = 15 \pmod{7}$  (since  $15 \times 1 \equiv 1 \pmod{7}$
- $\blacksquare$  So , X  $\equiv$  2 x 2 x 35 + 3 x 1 x 21 + 2 x 1 x 15 = 233  $\equiv$  23 (mod 105)
- So answer:  $X \equiv 23 \pmod{105}$

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#### Primitive Root

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- Primitive root: if p is prime, then successive powers of a 'generate' the group mod p
- these are useful but relatively hard to find

#### Powers of Mod 19

 $Z_{19}$ 

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a	$a^2$	$a^3$	$a^4$	$a^5$	$a^6$	$a^7$	$a^8$	$a^9$	$a^{10}$	$a^{11}$	$a^{12}$	$a^{13}$	$a^{14}$	$a^{15}$	$a^{16}$	$a^{17}$	$a^{18}$
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
2	4	8	16	13	7	14	9	18	17	15	11	3	6	12	5	10	1
3	9	8	5	15	7	2	6	18	16	10	11	14	4	12	17	13	- 1
4	16	7	9	17	11	6	5	1	4	16	7	9	17	11	6	5	1
5	6	11	17	9	7	16	4	1	5	6	11	17	9	7	16	4	1
6	17	7	4	5	11	9	16	1	6	17	7	4	5	11	9	16	1
7	11	1	7	11	1	7	11	1	7	11	1	7	11	1	7	11	1
8	7	18	11	12	1	8	7	18	11	12	1	8	7	18	11	12	1
9	5	7	6	16	11	4	17	1	9	5	7	6	16	11	4	17	1
10	5	12	6	3	11	15	17	18	9	14	7	13	16	8	4	2	- 1
11	7	1	11	7	1	11	7	1	11	7	1	11	7	1	11	7	1
12	11	18	7	8	1	12	11	18	7	8	1	12	11	18	7	8	1
13	17	12	4	14	11	10	16	18	6	2	7	15	5	8	9	3	1
14	6	8	17	10	7	3	4	18	5	13	11	2	9	12	16	15	- 1
15	16	12	9	2	11	13	5	18	4	3	7	10	17	8	6	14	1
16	9	11	5	4	7	17	6	1	16	9	11	5	4	7	17	6	1
17	4	11	16	6	7	5	9	1	17	4	11	16	6	7	5	9	1
18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1

For the prime number 19 the **primitive roots** are **2**, **3**, **10**, **13**,

## Discrete Logarithms

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- The inverse problem to exponentiation is to find the discrete logarithm of a number modulo p
- That is to find i such that  $b = a^i \pmod{p}$
- This is written as  $i = dlog_a b \pmod{p}$
- If a is a primitive root then it always exists, otherwise it may not, eg.
- The discrete logarithm does not always exist, for instance there is no solution to  $2^x \equiv 3 \pmod{7}$ .
- There is no simple condition to determine if the discrete logarithm exists.
- Whilst exponentiation is relatively easy, finding discrete logarithms is generally a hard problem



## Discrete Logarithms

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Discrete Logarithms For example, consider  $Z_{23}$ 

To compute  $3^4$  in this group, we first compute  $3^4$ =81, and then we divide 81 by 23, obtaining a remainder of 12. Thus  $3^4$ =12 in the group  $Z_{23^*}$ 

Discrete logarithm is just the inverse operation. For example, take the equation  $3^k\equiv 12 \pmod{23}$  for k. As shown above k=4 is a solution, but it is not the only solution. Since  $3^{22}\equiv 1 \pmod{23}$ , it also follows that if n is an integer then  $3^{4+22n}\equiv 12\times 1^n\equiv 12 \pmod{23}$ . Hence the equation has infinitely many solutions of the form 4+22n.

## Summary

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- concept of groups, rings, fields
- modular arithmetic with integers
- Euclid's algorithm for GCD & Inverse
- finite fields GF(p)
- **polynomial** arithmetic in general and in  $GF(2^n)$
- Fermat's and Euler's Theorems
- Primality Testing
- Chinese Remainder Theorem
- Discrete Logarithms