

Stacks -- Procedures -- Macros

UCS1502

MICROPROCESSORS AND INTERFACING

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AP/CSE



Learning Objectives

1. To understand the stack of 8086
2. To understand the difference between procedures and macros

Overview

- ▶ Stack
- ▶ Procedure
- ▶ Types of Procedures
- ▶ Advantage and disadvantages
- ▶ Macro

Stack -Introduction

- ▶ A section of memory is set aside for storing return addresses.
- ▶ It is also used to save the contents of registers for the calling program while a procedure executes.
- ▶ Used to hold data or addresses that will be acted upon by a procedure.
- ▶ Stack segment register holds the segment base address of the stack segment. Its size is 64K
- ▶ Stack Pointer register is used to hold the offset of the last word written on the stack.



```
; 8086 PROGRAM fragment showing the initialization  
; of stack segment register and stack pointer register
```

```
STACK_SEG SEGMENT STACK  
    DW 40 DUP(0)  
STACK_TOP LABEL WORD  
STACK_SEG ENDS
```

```
CODE SEGMENT  
    ASSUME CS:CODE, SS:STACK_SEG  
    MOV AX, STACK_SEG ; Initialize stack  
    MOV SS, AX ; segment register  
    LEA SP, STACK_TOP ; Initialize stack pointer  
    : ; Continue with program  
    :  
CODE ENDS  
END
```

Fig. 5.8 Required program additions when using a stack.

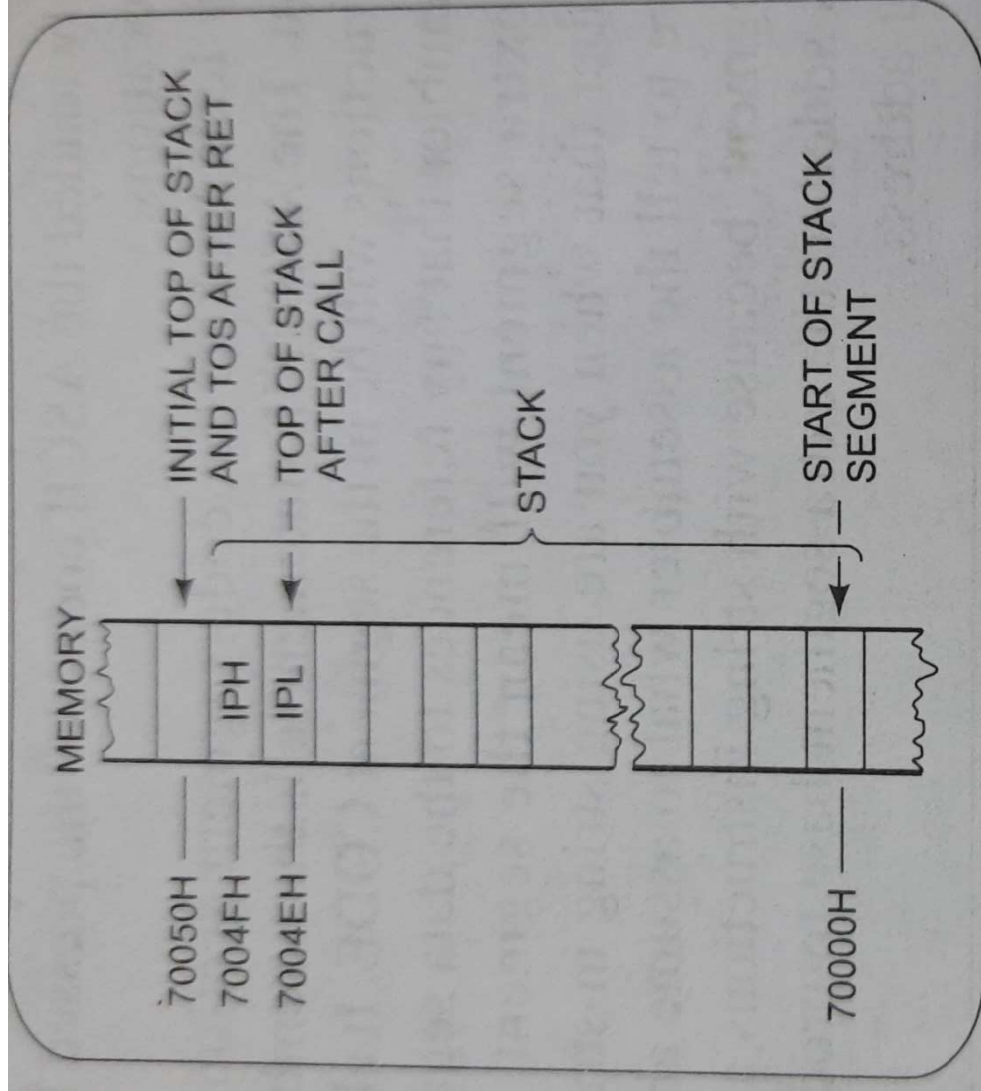
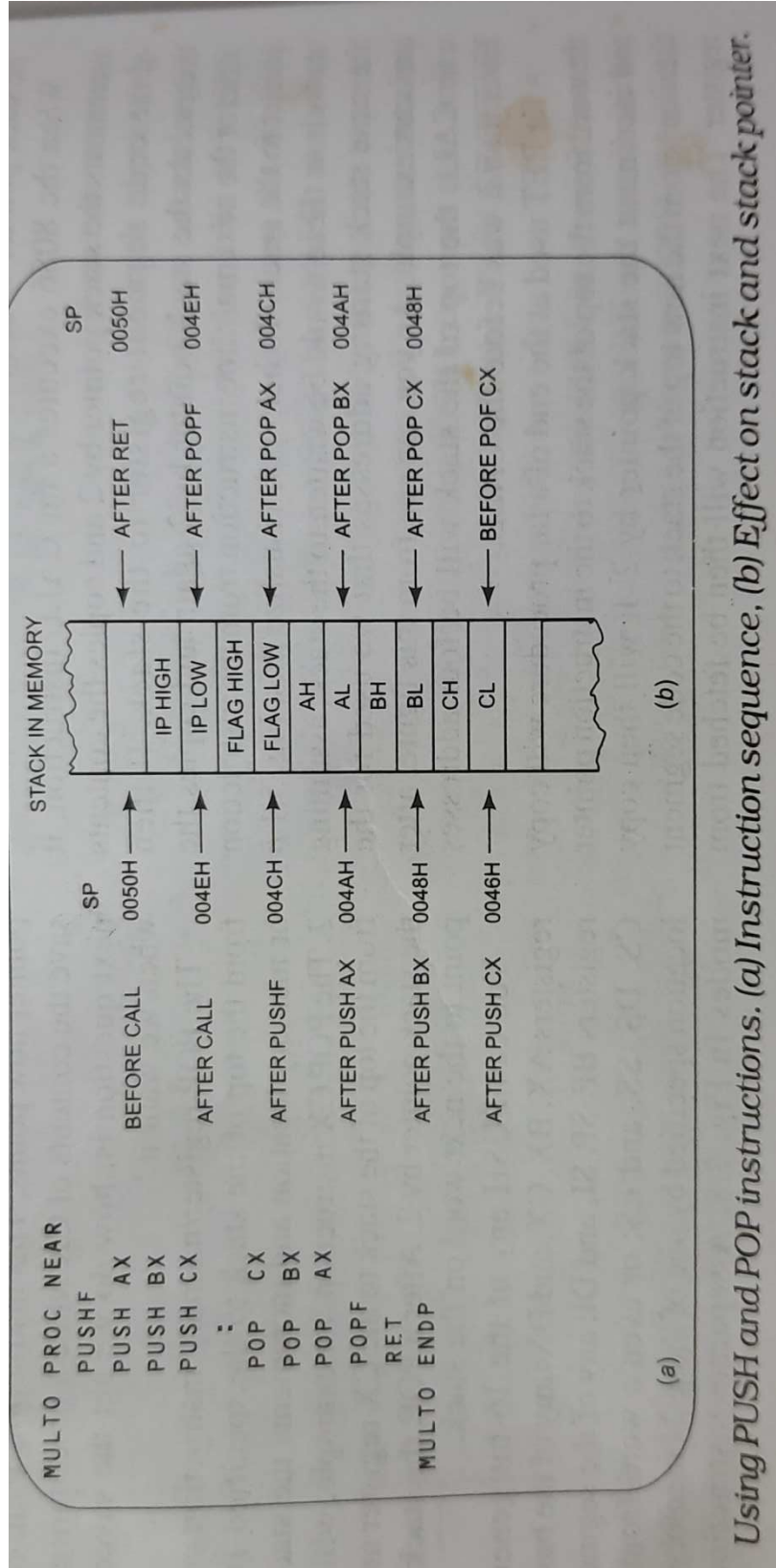


Fig. 5.7 Stack diagram showing how the return address is pushed onto the stack by CALL.



Using PUSH and POP instructions. (a) Instruction sequence, (b) Effect on stack and stack pointer.

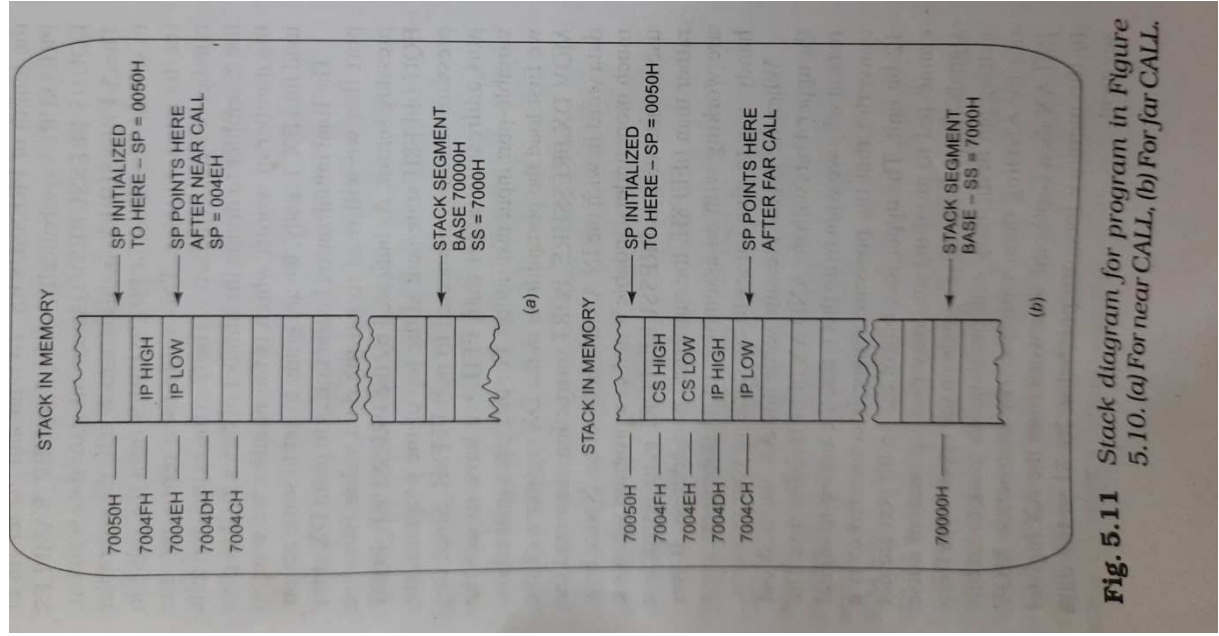


Fig. 5.11 Stack diagram for program in Figure 5.10. (a) For near CALL, (b) For far CALL.

Procedure

- ▶ To avoid writing the sequence of instructions in the program each time, we write a sequence of a separate subprogram called procedure.
- ▶ Direct Within-Segment Near call
- ▶ Indirect Within-Segment Near call
- ▶ Inter-segment Far call



```

CODE    SEGMENT
        ASSUME CS:CODE, DS:DATA, SS:STACK_SEG
        :
        :
        CALL MULTIPLY_32
        :
CODE    ENDS

PROCEDURES SEGMENT
        MULTIPLY_32 PROC FAR
            ASSUME CS:PROCEDURES
            :
            :
        MULTIPLY_32 ENDP
PROCEDURES    ENDS

```

Fig. 5.24 Program additions needed for a far procedure.

CALL = CALL

Within segment or group, IP relative

| Opcode | DispLow | DispHigh | Operation |
|--------|---------|----------|--|
| E8 | 19 | | $IP \leftarrow IP + \text{Disp16} - (\text{SP}) \leftarrow \text{return link}$ |

Within segment or group, Indirect

| Opcode | mod 010 r/m | Operation |
|--------|-------------|--|
| FF | 16 | $IP \leftarrow \text{Reg16} - (\text{SP}) \leftarrow \text{return link}$ |
| FF | 21 + EA | $IP \leftarrow \text{Mem16} - (\text{SP}) \leftarrow \text{return link}$ |

Inter-segment or group, Direct

| Opcode | offset-low | offset-high | seg-low | seg-high |
|--------|------------|-------------|---------|----------|
| 9A | 28 | | | |
| | | | | |

Operation: $CS \leftarrow \text{segbase}$
 $IP \leftarrow \text{offset}$

Inter-segment or group, Indirect

| Opcode | mod 100 r/m | mem-low | mem-high |
|--------|-------------|---------|----------|
| FF | 37 + EA | | |
| | | | |

Operation: $CS \leftarrow \text{Segbase}$
 $IP \leftarrow \text{Offset}$

(a)

RET = Return from Subroutine

| Opcode | Clocks | Operation |
|--------|--------|----------------------|
| C3 | 8 | intra-segment return |
| CB | 18 | intra-segment return |

Return and add constant to SP

| Opcode | DataL | DataH | Operation |
|--------|-------|-------|---------------------------|
| C2 | 12 | | intra-segment ret and add |
| CA | 17 | | intra-segment ret and add |

(b)

Fig. 5.6 8086 CALL and RET instruction formats.
(a) CALL. (b) RET. (Intel Corporation)

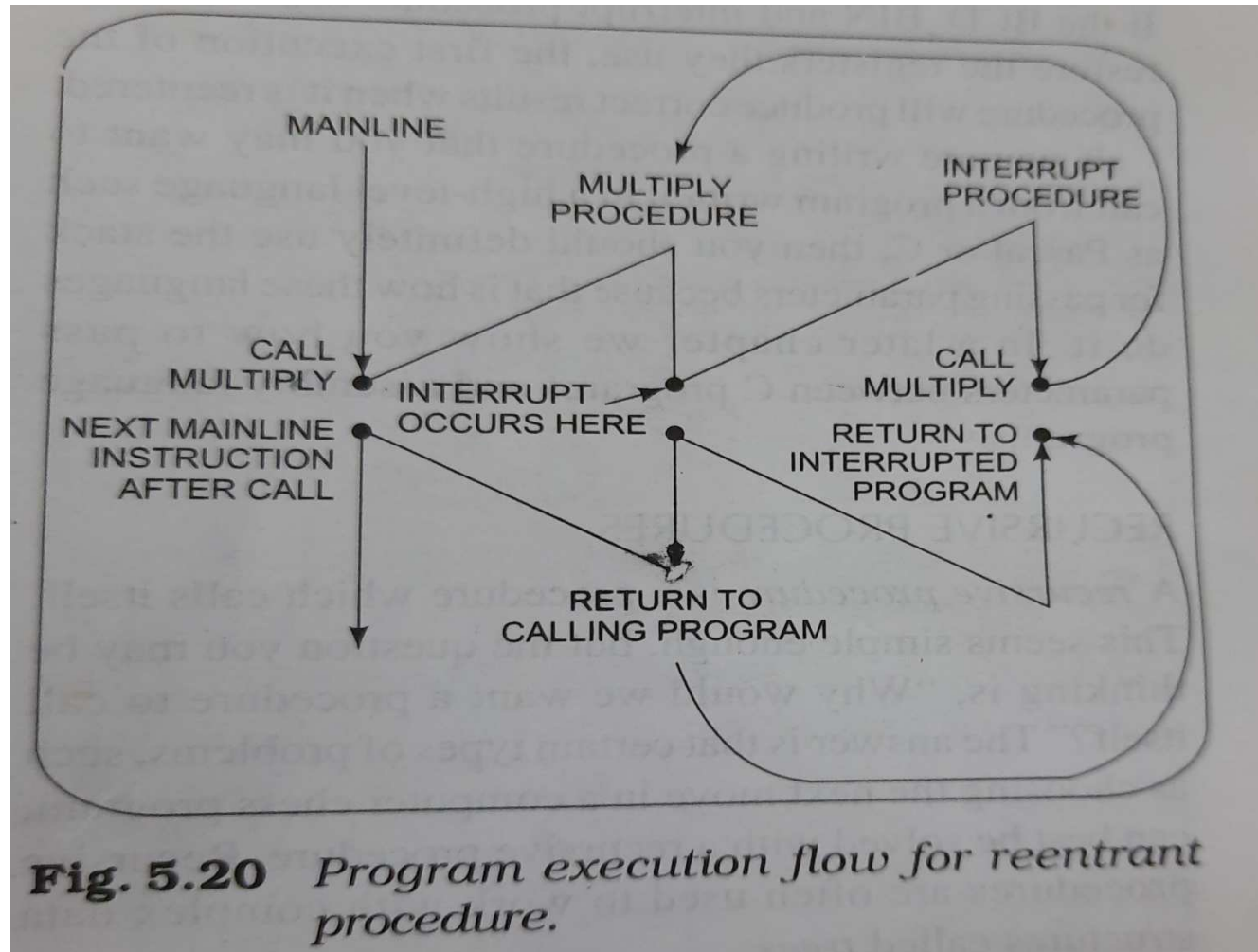
Passing parameters to and from procedures

- ▶ Four ways
 - ▶ In Registers
 - ▶ In dedicated memory locations accessed by name
 - ▶ With pointers passed in registers
 - ▶ With the stack



Fig. 5.17 Example program passing parameters on the stack.

Reentrant Procedure



Recursive Procedure

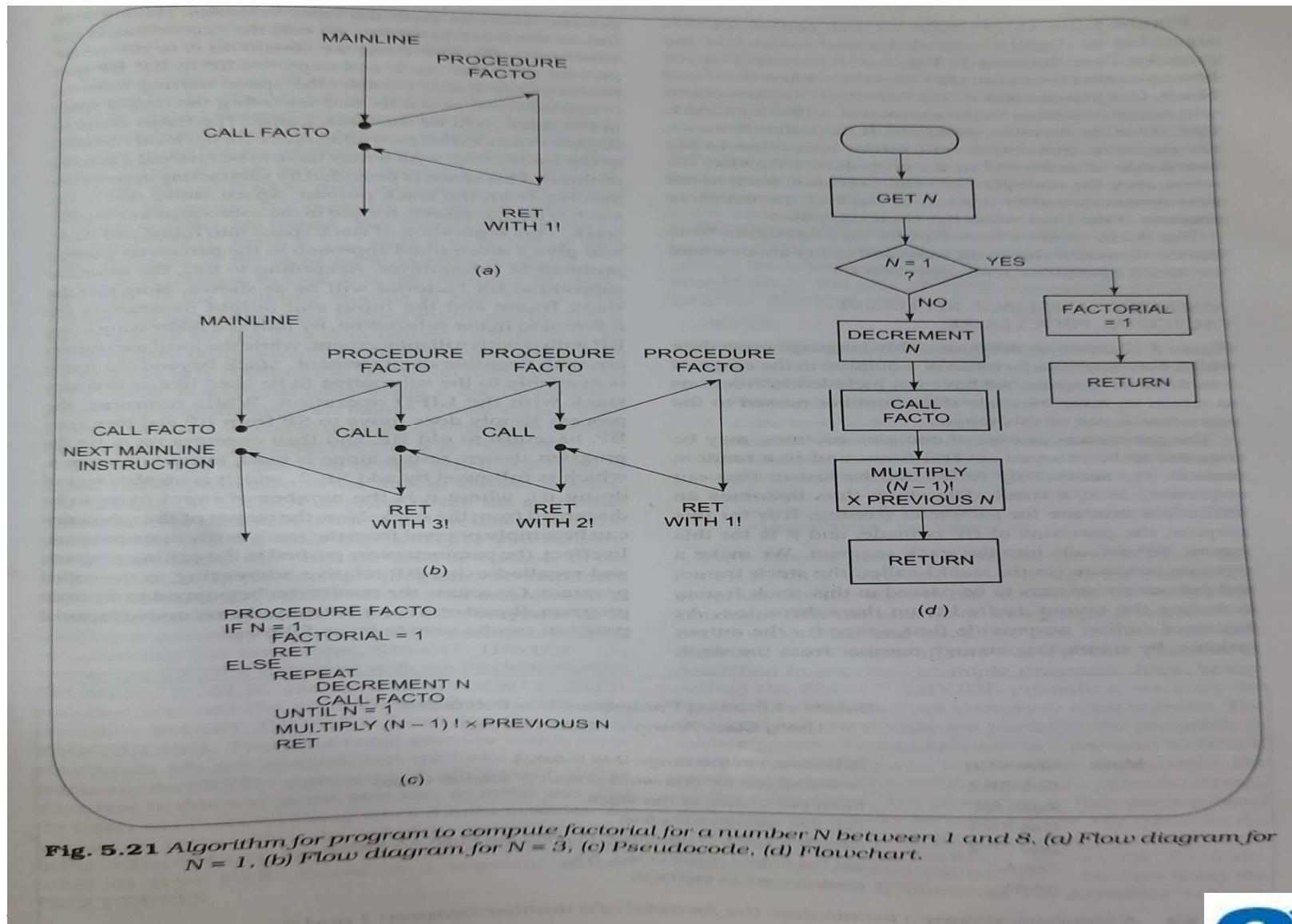


Fig. 5.21 Algorithm for program to compute factorial for a number N between 1 and 8, (a) Flow diagram for $N = 1$, (b) Flow diagram for $N = 3$, (c) Pseudocode, (d) Flowchart.

Advantage and disadvantage of procedure

- ▶ **Advantage:**
 - ▶ Machine codes for the group of instructions in procedure should be put in memory only once
- ▶ **Disadvantage:**
 - ▶ Need for a stack
 - ▶ Overhead time to call and return



Macro

- ▶ A macro is a group of instructions we bracket and give a name at the start of our program
- ▶ Each time we call it, the assembler will insert it in the program.



Advantage and disadvantage of macro

- ▶ Advantage:
 - ▶ No need for a stack
 - ▶ No Overhead time to call and return
- ▶ Disadvantage:
 - ▶ Machine codes for the group of instructions in procedure should be put in memory multiple times.



```
PUSH-ALL MACRO
PUSHF
PUSH AX
PUSH BX
PUSH CX
PUSH DX
PUSH BP
PUSH SI
PUSH DI
PUSH DS
PUSH ES
PUSH SS
```

```
ENDM
```

```
BREATH_RATE PROC FAR
ASSUME CS:PROCEDURES, DS:PATIENT_PARAMETERS
PUSH_ALL ; Macro call
MOV AX, PATIENT-PARAMETERS : Initialize data
MOVE DS, AX ; segment reg
```

```

MOVE_ASCII MACRO NUMBER, SOURCE, DESTINATION
    MOV CX, NUMBER      ; Number of characters to be moved in CX
    LEA SI, SOURCE       ; Point SI at ASCII source
    LEA DI, DESTINATION ; Point DI at ASCII destination
    CLD                 ; Autoincrement pointers after move
    REP MOVSB           ; Copy ASCII string to new location
ENDM

```

```

MOV CX, 03DH      ; Number of characters to be moved in CX
LEA SI, BLOCK_START ; Point SI at ASCII destination
LEA DI, BLOCK_DEST ; Point DI at ASCII destination
CLD               ; Autoincrement pointers after move
REP MOVSB        ; Copy ASCII string to new location

```

Check your understanding?

- ▶ What is the status of SP when PUSH is executed?
- ▶ What is the difference between procedure and macro?
- ▶ What increases the overhead time to call and return?



Summary

- ▶ Stack
- ▶ Procedure
- ▶ Types of Procedures
- ▶ Advantage and disadvantages
- ▶ Macro

Reference

- ▶ Douglas V Hall, “Microprocessors and Interfacing, Programming and Hardware”, TMH, 2012.



Thank you

