Name: Shalini N Reg. No: 805001998

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	Lamportis distributed mutex algorithm					
	P1 -> p3 p0 p1 -	> 08 11 01	ayou in in			
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14 10	PI	PO	p2 []]			
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3 -	add to queue	KO P [3, pg. pi)	Rep (p1 . 3 P3 . P1)			
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6-	peg (8, po. pi) ada p	09 (8. p1 , p8) at	poa (8, p1, p3) add			
7 -	Reg (8 p3, p1) add p	209 (8. ps , ps)	Roa (8, pa, p3) add			
8 -	F	20p(11, po. p1)	Rep (11, p3, p1)			
9						
10-	Rap (11, po, p1)					
11-	Rep(11, p3, p1)					
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4-	Rep (151, p1, & po. p3)	Rel (SI, p1, p2)	Rel (15, pl. p3)			
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18-	Re1 (1, pa.p1) semoro.	901100.	Reils, pa.pl) ronus			
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20	20,10,00					
21-	Pamore pamore	Pop(Q1, p3, p2)				
22 -	Rellan, pi gpa, psu 3		1			
23-	Rq (24, p2, b1)ada	2009 (24. pa Sp1. ps.	829181, p1, p3 Jada			
24-	p.	20010				
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<u>3 -</u>	PI	→ P3 P8 P1 P1 (000) peq (1, p1, p8, p3)	→ P& p P& (000) Req (1. p11. p2)	Req (1, p1, p3)		1 - 1 - 1 - 1 - 1
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15	[cs]				
4-		18. P2. P.3)			
17		Rep (18, p2, p3)			
(8-		<u>ES</u>			
(9					
21 -	Rel (20. PI & P2, P33) p 09 (20	2, pa, {p1, p3}) pag (22, p1, p3)			
24-		p. pl. p2) pag (22, p2, p3)			
23 -	Rep (2.	2. p2. p1) R99 (21, p3			
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21-	Rep (21, p3, p1)				
28-	pop(28.p1.p2)(000) p2p(8	28, p1, p2)			
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