SWE4001 – System Programming
Module 3: Assembler
Lesson 8 of 9: Assembler Design Options

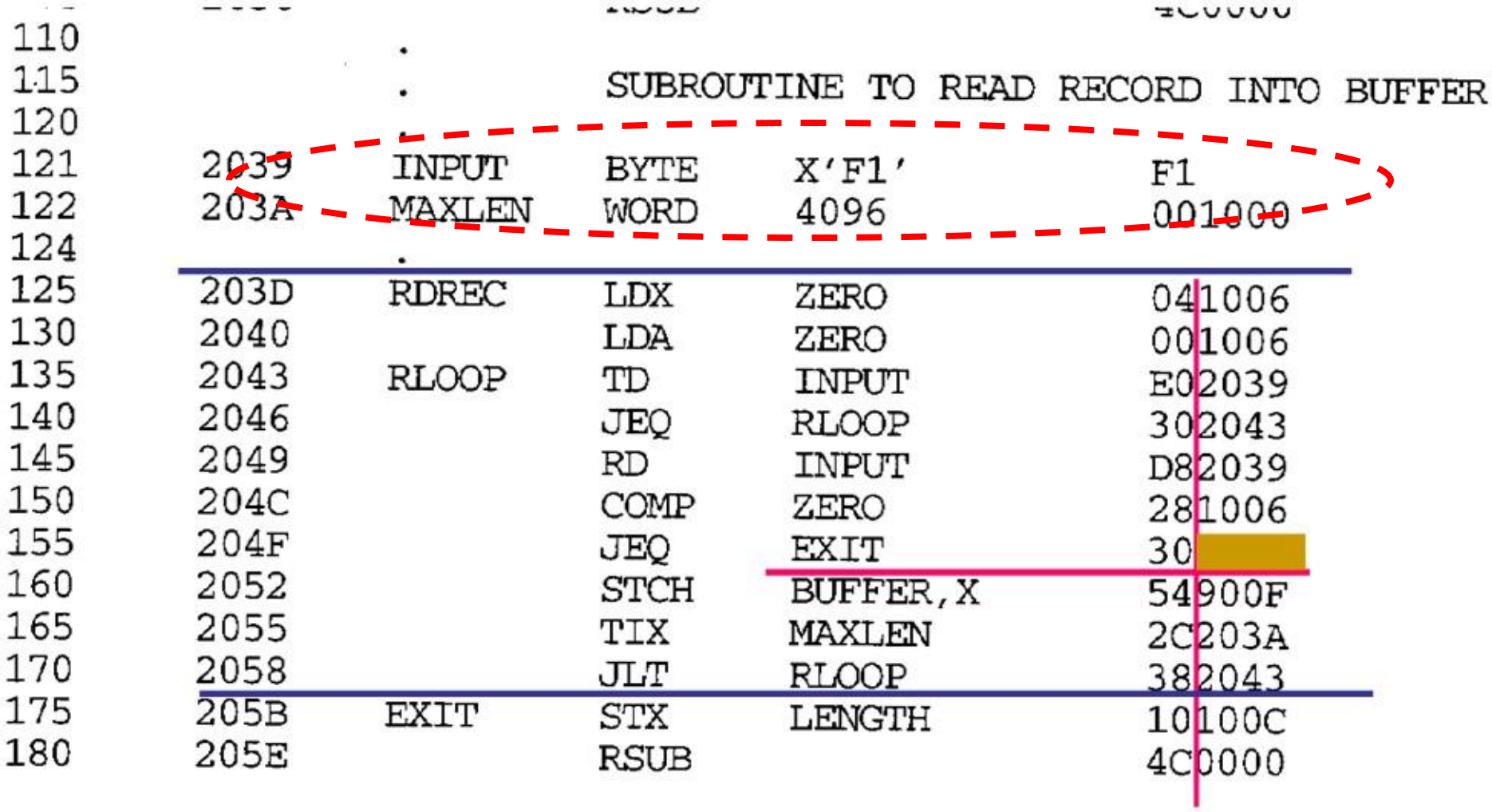
2.4 Assembler Design Options
2.4.1 Two-Pass Assembler

- Most assemblers
  - Processing the source program into two passes.
  - The internal tables and subroutines that are used only during Pass 1.
  - The SYMTAB, LITTAB, and OPTAB are used by both passes.
- The main problems to assemble a program in one pass involves forward references.

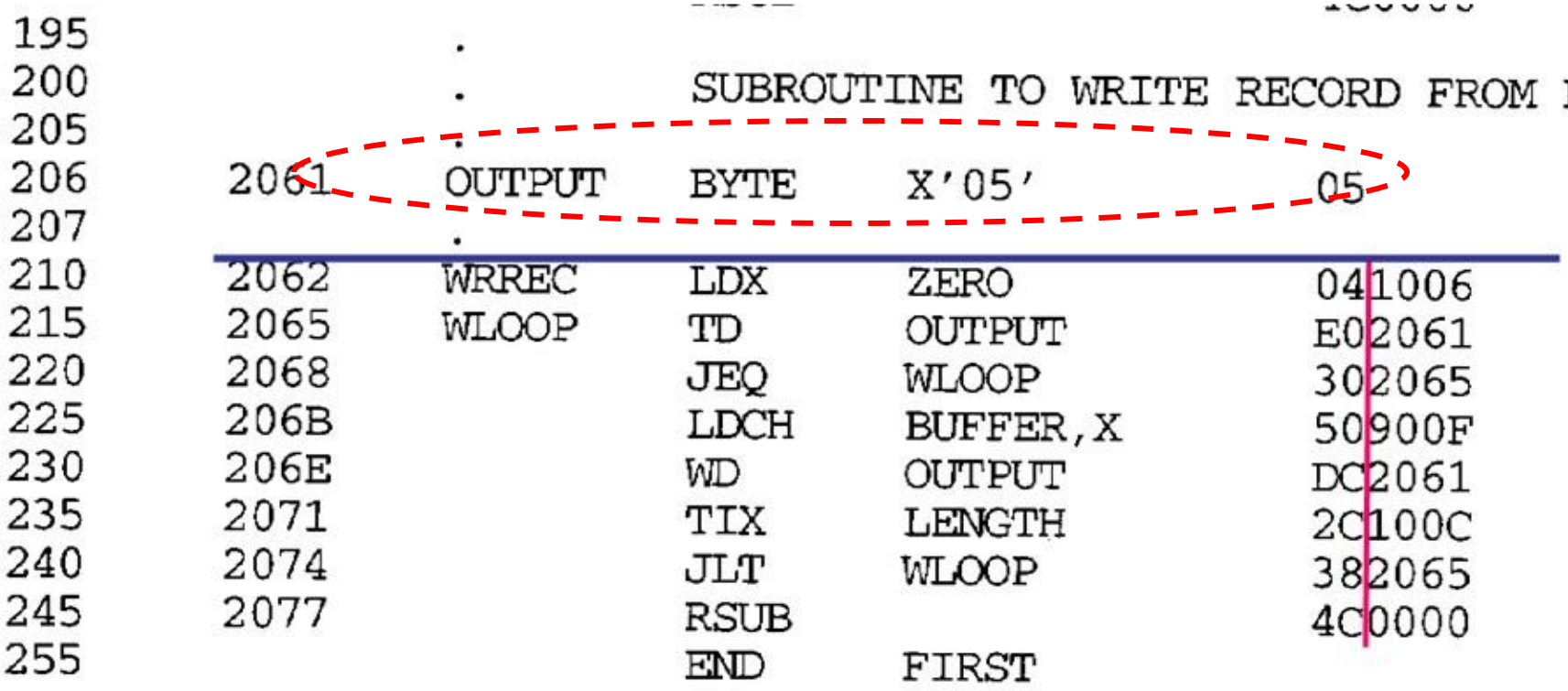
2.4.2 One-Pass Assemblers

- Eliminate forward references
  - Data items are defined before they are referenced.
  - But, forward references to labels on instructions cannot be eliminated as easily.
  - Prohibit forward references to labels.

Line	Loc	Source statement			Object code
0	1000	COPY	START	1000	
1	1000	EOF	BYTE	C'EOF'	454F46
2	1003	THREE	WORD	3	000003
3	1006	ZERO	WORD	0	000000
4	1009	RETADR	RESW	1	
5	100C	LENGTH	RESW	1	
6	100F	BUFFER	RESB	4096	
9		.			
10	200F	FIRST	STL	RETADR	141009
15	2012	CLOOP	JSUB	RDREC	48
20	2015		LDA	LENGTH	00100C
25	2018		COMP	ZERO	281006
30	201B		JEQ	ENDFIL	30
35	201E		JSUB	WRREC	48
40	2021		J	CLOOP	302012
45	2024	ENDFIL	LDA	EOF	001000
50	2027		STA	BUFFER	0C100F
55	202A		LDA	THREE	001003
60	202D		STA	LENGTH	0C100C
65	2030		JSUB	WRREC	48
70	2033		LDL	RETADR	081009
75	2036		RSUB		4C0000



```
110                 .
115                 .       SUBROUTINE TO READ RECORD INTO BUFFER
120                 .
121  2039   INPUT   BYTE    X'F1'           F1
122  203A   MAXLEN  WORD    4096            001000
124                 .
125  203D   RDREC   LDX     ZERO            041006
130  2040           LDA     ZERO            001006
135  2043   RLOOP   TD      INPUT           E02039
140  2046           JEQ     RLOOP           302043
145  2049           RD      INPUT           D82039
150  204C           COMP    ZERO            281006
155  204F           JEQ     EXIT            30
160  2052           STCH    BUFFER,X        54900F
165  2055           TIX     MAXLEN          2C203A
170  2058           JLT     RLOOP           382043
175  205B   EXIT    STX     LENGTH          10100C
180  205E           RSUB                    4C0000
```



```
195
200                .     SUBROUTINE TO WRITE RECORD FROM BUFFER
205                .
206  2061  OUTPUT  BYTE    X'05'          05
207                .
210  2062  WRREC   LDX     ZERO           041006
215  2065  WLOOP   TD      OUTPUT         E02061
220  2068          JEQ     WLOOP          302065
225  206B          LDCH    BUFFER,X       50900F
230  206E          WD      OUTPUT         DC2061
235  2071          TIX     LENGTH         2C100C
240  2074          JLT     WLOOP          382065
245  2077          RSUB                   4C0000
255                END     FIRST
```

Figure 2.18 Sample program for a one-pass assembler.

All variables are defined before they are used.

Two Types

- There are two types of one-pass assembler:
- **Produce object code directly in memory for immediate execution**
  - No loader is needed
  - Load-and-go for program development and testing
  - Good for computing center where most students reassemble their programs each time.
  - Avoids the overhead of writing the object program out and reading it back .
  - For a load-and-go assembler, the actual address must be known at assembly time, we can use an absolute program
- **Produce the usual kind of object program for later execution**

Internal Implementation

- The assembler generate object code instructions as it scans the source program.
- If an instruction operand is a symbol that has not yet been defined
  - The operand address is omitted when the instruction is assembled.
  - The symbol used as an operand is entered into the symbol table.
  - This entry is flagged to indicate that the symbol is undefined yet.

Internal Implementation

- The address of the operand field of the instruction that refers to the undefined symbol
  - Added to a list of forward references associated with the symbol table entry.
- When the definition of the symbol is encountered,
  - The forward reference list for that symbol is scanned,
  - The proper address is inserted into any instruction previously generated.

2.4.2 One-Pass Assemblers

- Load-and-go one-pass assembler
  - The assembler avoids the overhead of writing the object program out and reading it back in.
  - The object program is produced in memory, the handling of forward references becomes less difficult.
  - Figure 2.19(a), shows the SYMTAB after scanning line 40 of the program in Figure 2.18.
  - Since RDREC was not yet defined, the instruction was assembled with no value assigned as the operand address (denote by - - - -).

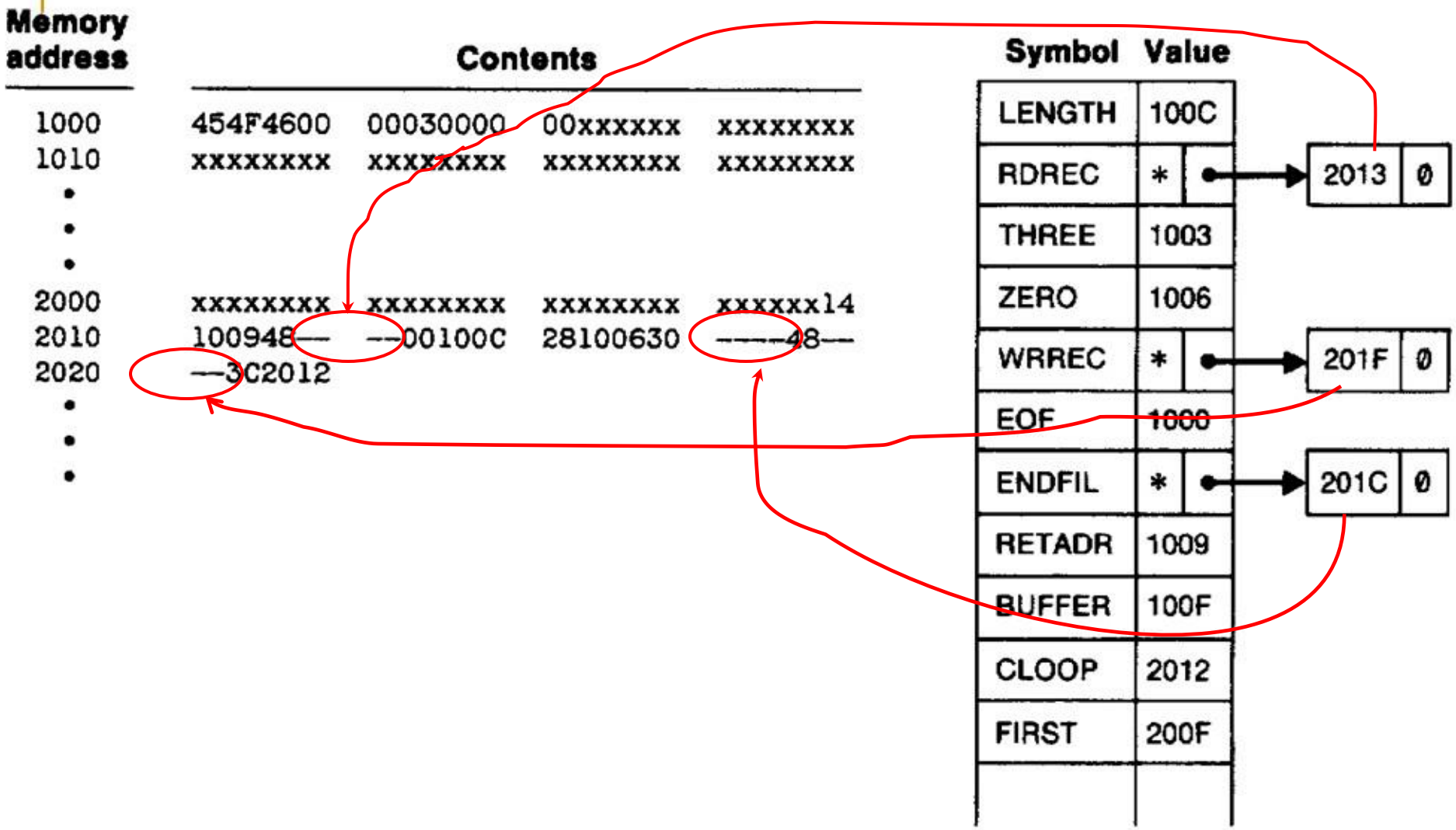


Figure 2.19(a) Object code in memory and symbol table entries for the program in Fig. 2.18 after scanning line 40.

Memory address	Contents			
1000	454F4600	00030000	00xxxxxx	xxxxxxxx
1010	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxxxx
•				
•				
•				
2000	xxxxxxxx	xxxxxxxx	xxxxxxxx	xxxxxx14
2010	10094820	3D00100C	28100630	202448--
2020	--3C2012	0010000C	100F0010	030C100C
2030	48----08	10094C00	00F10010	00041006
2040	001006E0	20393020	43D82039	28100630
2050	----5490	0F		
•				
•				
•				

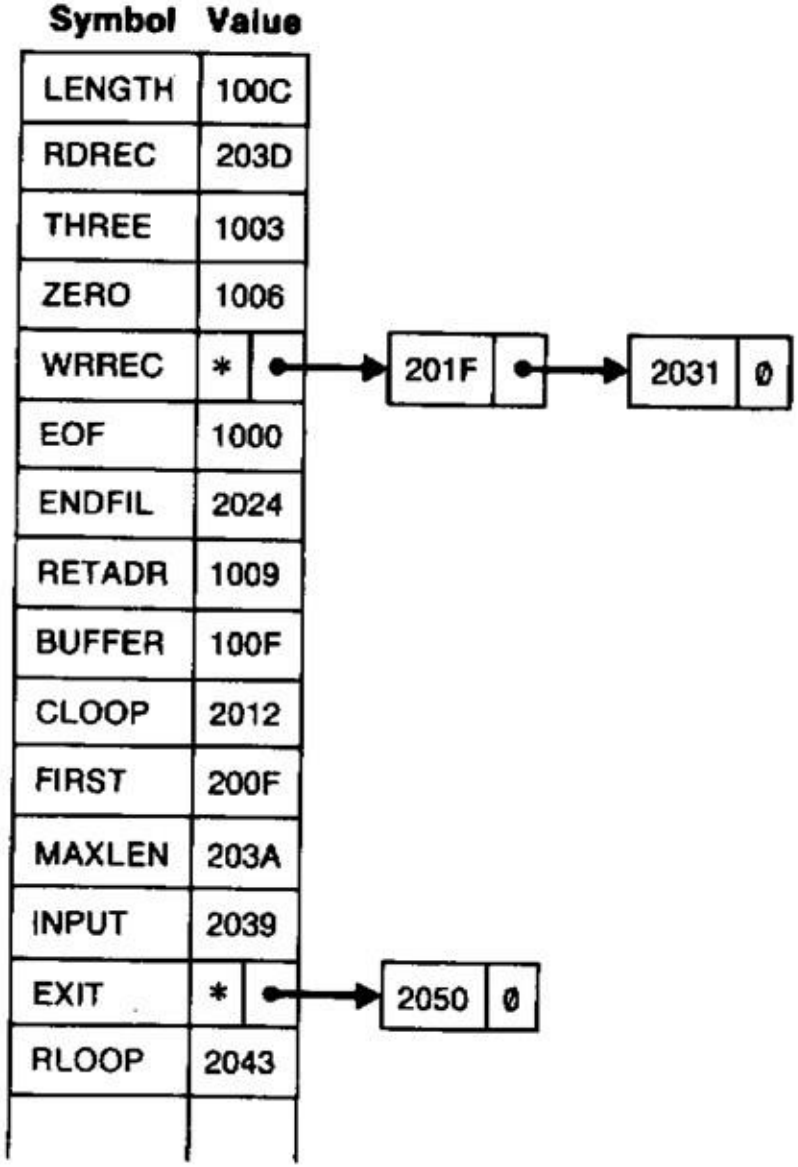


Figure 2.19(b) Object code in memory and symbol table entries for the program in Fig. 2.18 after scanning line 160.

2.4.2 One-Pass Assemblers

- Load-and-go one-pass assembler
  - RDREC was then entered into SYMTAB as an undefined symbol, the address of the operand field of the instruction (2013) was inserted.
  - Figure 2.19(b), when the symbol ENDFIL was defined (line 45), the assembler placed its value in the SYMTAB entry; it then inserted this value into the instruction operand field (201C).
  - At the end of the program, all symbols must be defined without any * in SYMTAB.
  - For a load-and-go assembler, the actual address must be known at assembly time.

2.4.2 One-Pass Assemblers

- Another one-pass assembler by generating OP
  - Generate another Text record with correct operand address.
  - When the program is loaded, this address will be inserted into the instruction by the action of the loader.
  - Figure 2.20, the operand addresses for the instructions on lines 15, 30, and 35 have been generated as 0000.
  - When the definition of ENDFIL is encountered on line 45, the third Text record is generated, the value 2024 is to be loaded at location 201C.
  - The loader completes forward references.

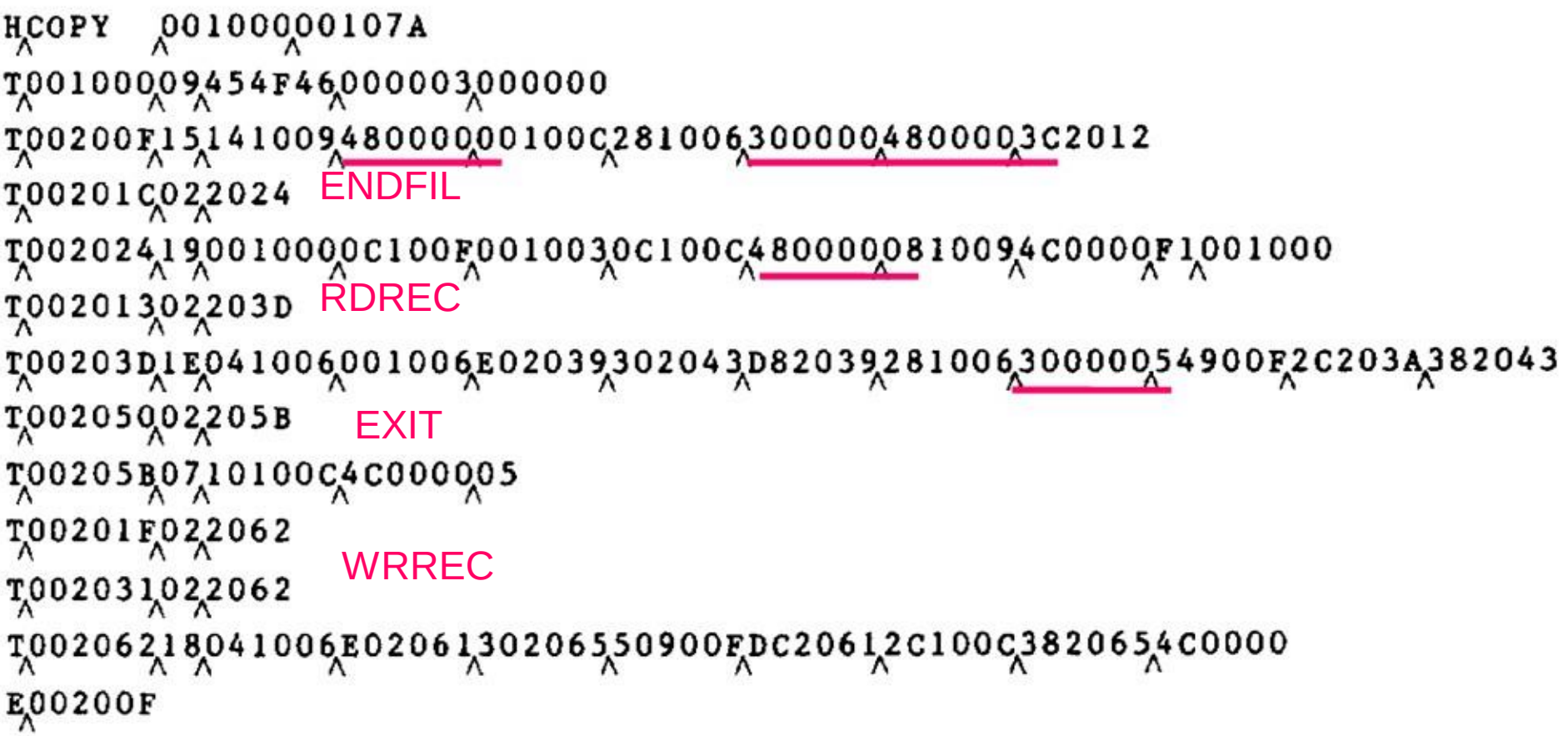


Figure 2.20 Object program from one-pass assembler for program in Fig. 2.18.

2.4.2 One-Pass Assemblers

④ In this section, simple one-pass assemblers handled absolute programs (SIC example).

Multi-Pass Assemblers

- Restriction on EQU and ORG
  - No forward reference, since symbols' value can't be defined during the first pass
- Example:

```
ALPHA   EQU   BETA
BETA EQU   DELTA
DELTA   RESW  1
```

  - Assemblers with 2 passes cannot resolve

Multi-Pass Assembler

- The assembler directives that define symbol requires the R.H.S be defined previously in the source program
- If we use a two-pass assembler, the following symbol definition cannot be allowed.

```
ALPHA   EQU   BETA
BETA    EQU   DELTA
DELTA   RESW  1
```

- This is because ALPHA and BETA cannot be defined in pass 1.
  - we allow multi-pass processing,
  - DELTA is defined in pass 1,
  - BETA is defined in pass 2,
  - ALPHA is defined in pass 3, and the above definitions can
  - be allowed.
- This is the motivation for using a multi-pass assembler.

Multi-Pass Assemblers

- Resolve forward references with as many passes as needed
  - Portions that involve forward references in symbol definition are saved during Pass 1
  - Additional passes through stored definitions
  - Finally a normal Pass 2
- Example implementation:
  - Use link lists to keep track of whose value depend on an undefined symbol

Multi-Pass Assembler Implementation

- Use a symbol table to store symbols that are not totally defined yet.
- For a undefined symbol, in its entry,
  - We store the names and the number of undefined symbols which contribute to the calculation of its value.
  - We also keep a list of symbols whose values depend on the defined value of this symbol.
- When a symbol becomes defined, we use its value to reevaluate the values of all of the symbols that are kept in this list.
- The above step is performed recursively.

Figure 2.21(a): After Pass 1

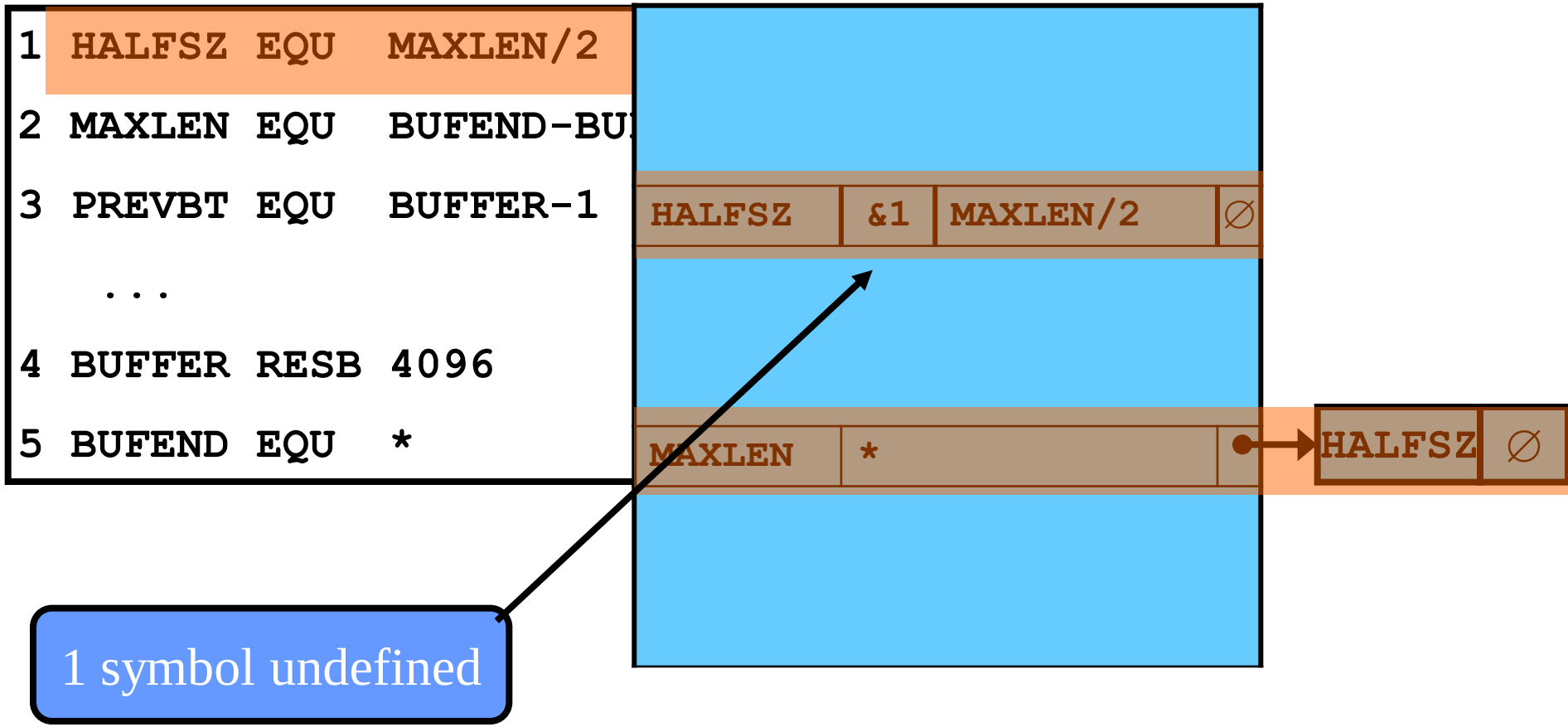


Figure 2.21(c): MAXLEN Defined

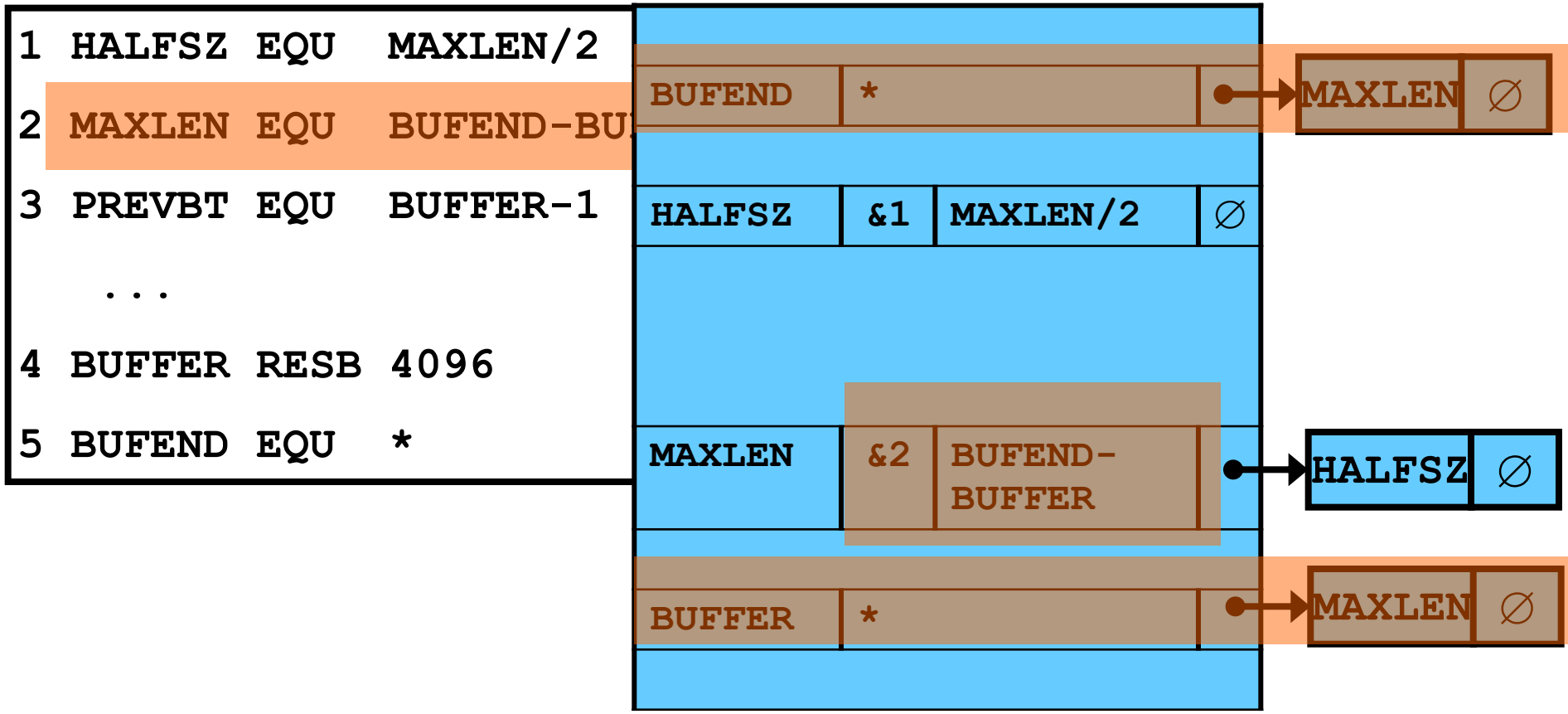


Figure 2.21(d): PREVBT Defined

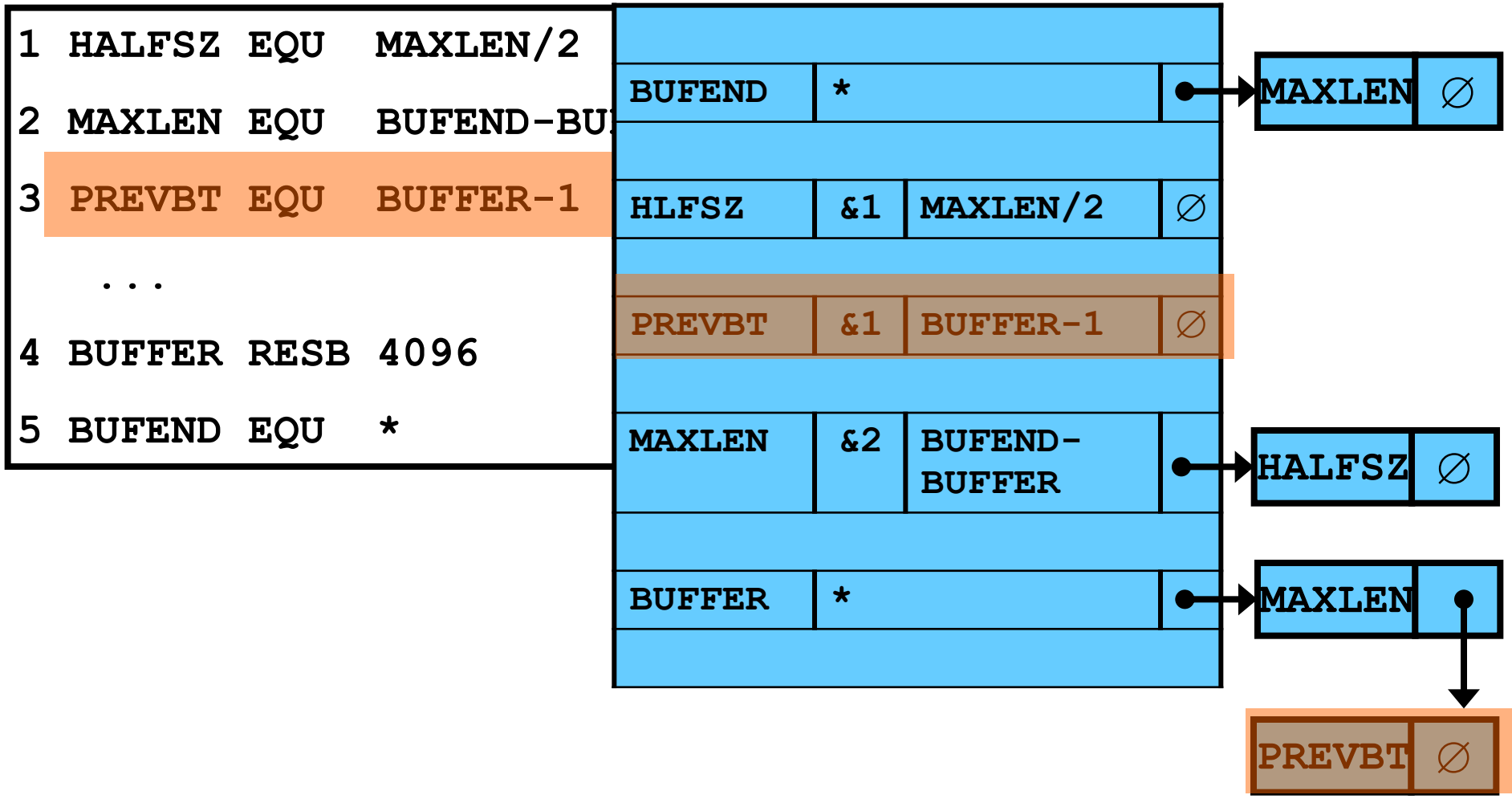


Figure 2.21(e): After Line 4

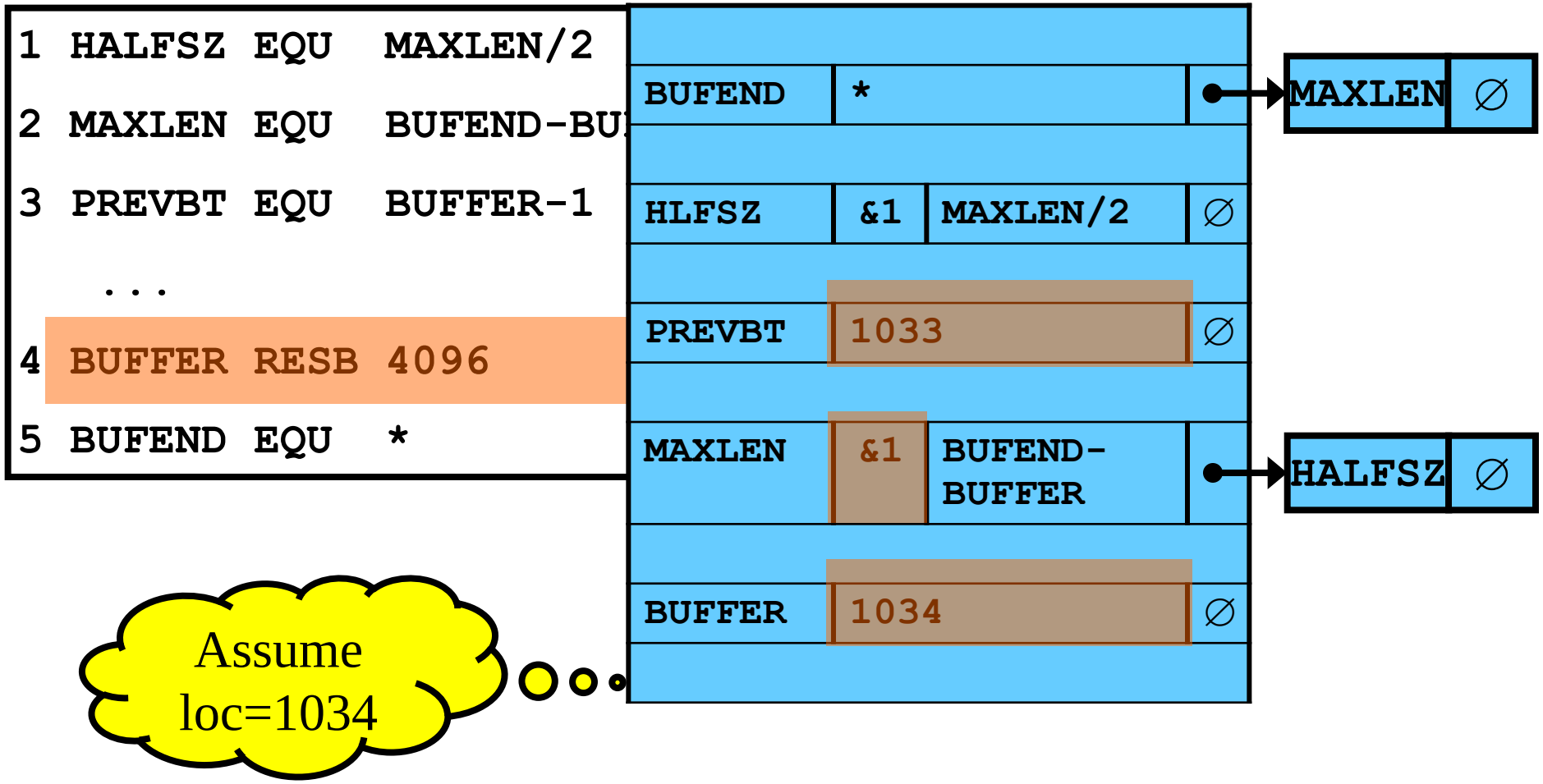


Figure 2.21(f): After Line 5

```
1  HALFSZ  EQU   MAXLEN/2
2  MAXLEN  EQU   BUFEND-BUFFER
3  PREVBT  EQU   BUFFER-1
     ...
4  BUFFER  RESB  4096
5  BUFEND  EQU   *
```