

HW 2 theory

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Problem 1: Input: n - an integer higher than 2
Recursive definition of algorithm:

$$j = 3, \dots, n$$

$$f_1 = 1$$

$$f_2 = 1$$

$$f_j = f_{j-1} + f_{j-2}$$

output: single integer: f_n

bash-like pseudo-code

FibonacciSeq(n)

f1= 1

f2= 1

i=0

while $i \leq n$

do

 if $f1 > f2$ then

$f1 = f1 + f2$

```

        f1
    else
        f2= f2 + f1
    f2
fi
i=1+i
done
if f2 < f1 then
    return f1
else
    return f2
fi

```

Problem 2:

input: w is an n by n adjacency matrix, with element w_{jk} containing the weight of edge from vertex k to j , or INF if there is no edge, we always have $w_{kk} = 0$

Recursive definition of the algorithm::

$$j = 1, \dots, n$$

$$d_j = \infty$$

$$d_1 = 0$$

$$d_j = \min_{(1 \leq k \leq j)} d_k + w(kj)$$

output: integer d_j for all j in $1, \dots, n$

bash-like pseudo-code

```

BellmanFord(w)

for j in n
do
     $d[0,j] = \text{INF}$ 
done

for k in n
do
     $d[k,k]=0$ 
done

for j in n
do
    for k in n
    do
        if  $[w[k,j] \neq \text{INF}]$  then
             $d[k,j] = d[k-1,j]$ 
            if  $[d[k,j] > d[k-1] + w[kj]]$  then
                 $d[j] = d[k-1] + w[k,j]$ 
            fi
        fi
    fi
fi

```

```

do
    if [w[k,j] != INF] then
        d[k,j] = d[k-1,j]
        if [d[k,j] > d[k]+w[k,j]] then
            if 1 ≤ k ≤ j then
                d[j]=d[k-1]+w[k,j]
            fi
        fi
    fi
done
return d[j]

```

done

problem 3 Input: L - a postive integer n - a postive integer V - an array of n postive integers w - an array of n postive integers, each $i=1$

recursive defintuon of the algorithm:

$m(jk) = 0$ for $j = 0, \dots, L$ and $k = 0$

$m(jk) = 0$ for $j = 0$ and $k = 1, \dots, n$

$m(jk) = m_{j, k-1}$ if $j - w_k < 0$ for $j = 1, \dots, K$ and $k = 1, \dots, n$

$m(jk) = \max(m_{j - w_k, k-1} + V_k, m_{j, k-1})$ if $j - w_k \geq 0$

for $j = 1, \dots, L$ and $k = 1, \dots, n$

a single integer: $m_{L,n}$

```

bash-like pseudo-code

unknownAlgo(rhime(L,n,V,w))

for j in range(0,L)
do
    m[j,0]=0
done

for k in range(0,n)
do
    m[0,k]=0
done

for j in k
do
    for k in n
    do
        if [m[j,k]==m[j,k-1]] then
            if [j - w[k] < 0] then
                m[j,] = j- w[k]
            fi
        fi
    fi
fi

```

```

done

done

for j in L
do
    for k in n
    do
        if  $[m[j,k] < m[[j-wk,k-1]+V[k]],m[j,k-1]]]$  then
             $m[j,k] = m[[k-wk,k-1+V[k]],m[j,k-1]]]$ 
            if  $[j-wk \geq 0]$  then
                 $m[,wk]=j-wk$ 
            fi
        fi
    done
done

done

return  $m[j,k]$ 

```