

Queue

- Work on FIFO (First in First Out)
- Insert from rear end and delete from front end

Queue ADT

Data:

1. Space for storing element

2) Front → For deletion

3) Rear → For Insertion

Operation

enqueue()

dequeue()

isEmpty()

isFull()

First()

Last()

- Queue can be implemented using

→ Array

→ Linked List

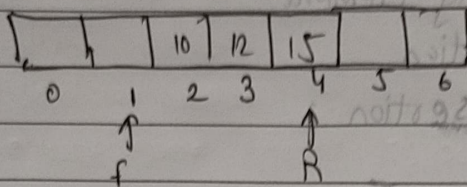
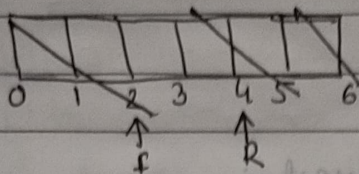
- Insert Queue using one pointer

→ Insertion will take order of 1 time

→ Deletion will take n time as the rest of the element needs to be shifted as deletion of any element will make vacant space

• Queue Using 2 pointer .

→ Initially $Front = Rear = -1$



Empty.

$if (front == Rear)$

→ Here insert using Rear pointer

→ Delete using front pointer (i.e. delete element of address at which it stands)

• Struct Queue

int size;

int front;

int Rear;

int *q;

int i;

* Struct Queue

```
{
    int size;
    int front;
    int Rear;
    int *Q;
}
```

```
};
```

int main()

```
{
    Struct Queue q;
    printf("Enter size");
    scanf("%d", &q.size);
    q.Q = (int*) malloc (q.size * sizeof(int));
    q.front = q.Rear = -1;
}
```

Void enqueue (Queue *q, int x)

```
{
    if (q.Rear == q.size - 1)
        printf("Queue is full");
    else
```

```
    q.Rear++;
    q.Q[q.Rear] = x;
}
```

Void dequeue (Queue *q) //dequeue

```
{
    int x = -1;
    if (q.front == q.Rear)
        printf("Queue is Empty");
    else
```

```
    {
        q.front++;
        x = q.Q[q.front];
    }
```

```
    return x;
}
```


* Drawback

→ Deleted spaces cannot be reused.

Solutions

1) Resetting pointer.

if (Front == Rear)

front = Rear = -1;

2) Circular Queue.

• whenever front is pointing should be empty.

→ Using mod value can be inserted.

Rear = (Rear + 1) % size

		→ Next index of Rear
0	$(0 + 1) \% 7$	1
1	$2 \% 7$	2
2		
3		
4		
5		
6	$7 \% 7$	0

* Queue Using Linked list

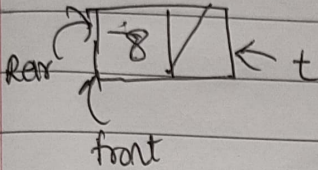
→ For Empty
if (front == NULL)

→ Full

Node *t = New Node

if (t == NULL)

→ for one node



* Double ended Queue.

↳ Insertion as well as deletion from both front & rear is possible

Insert

Delete

✓

✓

✓

✓

* Priority Queue

→ Element	A	B	C	D	E	F
Priority	1	1	2	1	2	3

Priority Queue

Q₁ [A | B | D]

→ Next two will be pop out completely.

Q₂ [C | E]

Q₃ [F]

* When ~~priority~~ is element itself is priority.

element 8, 8, 3, 5, 1, 3, 2

Methods to solve such issues of Queue

1) Insert in same order.
Delete Max priority by searching it.

2) Insert in increasing order of Priority.
Delete last element of array

* Method 1 for priority queue

Search $\rightarrow O(n)$

Shift $\rightarrow O(n)$

$O(2n) = O(n)$

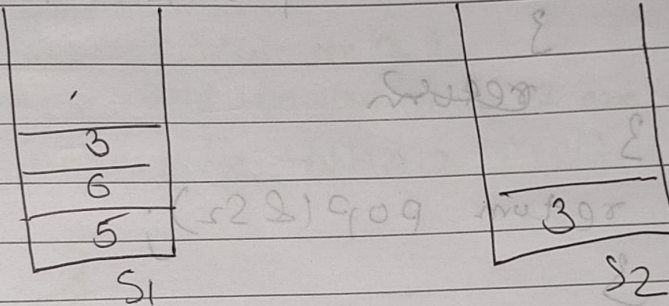
* Method 2

Insert $\rightarrow O(n)$

Delete $\rightarrow O(n)$

* Implementation of queue using 2 stack

• To delete 5 shift 3, 6 from S_1 to S_2



Logic

\rightarrow If S_2 is empty transfer element from S_1 to S_2 and then delete S_2 element

\rightarrow so it will act as queue instead of stack

\rightarrow If element are there in S_2 stack delete normally

```
enqueue(int x)
{
    push(&S1, x);
}
```

```
int dequeue()
{
```

```
    int x = -1
```

```
    if (isEmpty(S2))
```

```
    {
        if (isEmpty(S1))
```

```
        {
            printf("Queue Empty");
```

```
            return x;
        }
```

```
    }
```

```
    else
```

```
    {
```

```
        while (!isEmpty(S1))
```

```
            push(&S2, pop(&S1));
```

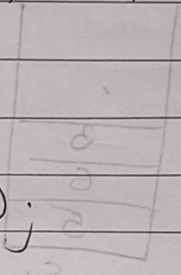
```
    }
```

```
    return
```

```
    }
```

```
    return pop(&S2);
```

```
}
```



Stack

IT element are there in stack