

Sparse Matrix

→ Very few non-zero element in Matrix

row	column	no. of non-zero element	← coordinate list
Total rows	Total column	total num. of non-zero element	

Compressed sparse row (30% of memory is reduced)

$A[x, y, \dots]$ → element

$IA[0, 1, 5, \dots]$ → Row & element

$JA[5, 8, 2, \dots]$ → column & element

→ Addition of sparse matrix

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 6 & 0 \\ 0 & 0 & 0 \end{bmatrix} \quad B = \begin{bmatrix} 0 & 5 & 0 \\ 0 & 4 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

	0	1	2		0	1	2
Total rows →	3	1	2	→ row	3	1	2
Total column →	3	1	2	→ column	3	2	2
	2	1	6	→ element	2	5	4

Total non-zero element

	0	1	2	3
C	3	1	1	2
	3	1	2	2
	3	1	5	10


```
* struct Element
{
    int i;
    int j;
    int val e;
}
```

```
* struct sparse
{
    int m;
    int n;
    int num;
    struct element *e;
```

```
int main()
{
    struct sparse and s;
    create(&s);
    // create(struct sparse *s)
    // #take dimension in s → m & n
    // #take number of input in s → num
    s → e = new Element[s → num]
    for (i = 0; i < s → num; i++)
    {
        scanf("%d %d %d", &s → e[i].i,
            &s → e[i].j,
            &s → e[i].e);
    }
```

* Adding.

Make two sparse matrix s1, s2

in function add

```
add (struct sparse *s1, struct sparse *s2)
{
```

```
    if (s1 → m != s2 → m || s1 → n != s2 → n) // dimension is different
        return;
```

```
    sum = new Sparse // struct sparse in heap
```

```
    sum → m = s1 → m; // giving structure
```

```
    sum → n = s1 → n // dimensions
```

```
    sum → e = new Element[s1 → num + s2 → num]
```


// Adding algorithm

$i=0, j=0, k=0;$
↳ for s1 ↳ for s2 ↳ for sum

while ($i < s1 \rightarrow \text{num}$ & $j < s2 \rightarrow \text{num}$)

{

if ($s1 \rightarrow e[i] \cdot i < s2 \rightarrow e[j] \cdot i$)

// if row of s1 is less than s2 copy s1 to sum

sum $\rightarrow e[k++] = s1 \rightarrow e[i++];$ // copy complete structure

// if row of s2 is less than s1 copy struct s2 to sum

elseif ($s1 \rightarrow e[i] \cdot i > s2 \rightarrow e[j] \cdot i$)

sum $\rightarrow e[k++] = s2 \rightarrow e[j++];$

else

// if column of s1 is less than s2 copy struct s1 to sum

if ($s1 \rightarrow e[i] \cdot j < s2 \rightarrow e[j] \cdot j$)

sum $\rightarrow e[k++] = s1 \rightarrow e[i++];$

elseif ($s1 \rightarrow e[i] \cdot j > s2 \rightarrow e[j] \cdot j$) // if column of s2 is less than s1 copy s2 to sum

sum $\rightarrow e[k++] = s2 \rightarrow e[j++];$

else

if both row & column are equal first copy entire structure of s1

sum $\rightarrow e[k++] = s1 \rightarrow e[i++];$

sum $\rightarrow e[k++] \cdot x = s2 \rightarrow e[j++] \cdot x;$

Then add value of structure of s2 to sum structure value

structure value

friend function is global (no need of scope resolution operator)

sparse operator $+ (sparse \& s)$,

istream \rightarrow for input

ostream → for output

reference	function name
↓	↓

istoeam should
return reference
var is

format \rightarrow istream & operator \gg (istream &is, sparse &s)

ostream & operator << (const team &os, sparse &s)

cin >> class name

~~code~~ 'code' << class name

Sparse SparseOperator + (Sparse & s)

↑ return type

~~X X X~~

Polynomial Representation

$$p(x) = 3x \overset{\text{EXP}}{\textcircled{5}} + 2x^2 + 5x^2 + 2x + 7$$

↑
coeff

2EXP

coeff	3	2	5	2	7
exp	5	4	2	1	0

struct poly

3

```
int n;
```

Struct Term * t;

3;

Struct Form 5

```
int coeff;
```

```
int Exp;
```

3.

Algo →

i) ~~poly~~ input >> number of elements.
(in Polydata type variable)
in variable n of poly)

ii) Assign heap memory to t as array
 Eg. $p \cdot t = \text{new Term}[p \cdot n]$

evaluation of $3x^2 + 5x$

$$= 3 \times x \times x + 5 \times x$$

iii) for $i=0$ till $i < p.n$

input $(p.t[i].coeff)$

input $(p.t[i].Exp)$

• Evaluation of $3x^2 + 5x$

• Addition of polynomial (see in repo for code)

