

Link - <https://www.designa.in/17735/uber-oa-ctc-60-l-sde1-sep-5>

-> Given an array of size "N" ; you have to do 3 type of operations sequentially

-> First delete any p elements;

-> Then delete any q-adjacent pair elements.

-> Then delete any r-adjacent triplet elements.

After doing all this :- Sum of the remaining array should be maximum.

Observation :-> If  $Q = 0$   $R = 0$  and  $P \neq 0$  -> solution is to remove p smallest elements in the array the remaining array would give us the maximum sum.

Observation :-> Order of operations does not matter -> you can do any operation anytime. Final result will be the same.

Easier version :-  $Q = 0$   $R = 0$

How to do easier version with dp ->

$dp[i][j]$  = best way to reach till index  $i$  by ignoring  $j$  elements;

$\Rightarrow dp[i][j] = \max(dp[i-1][j-1], b[i] + dp[i-1][j])$

$\Rightarrow$  <https://www.jdoodle.com/ia/1ghl>

Medium - version :-  $R = 0$

$\rightarrow dp[i][j][k]$  = consider first " $i$ " elements but ignore  $j$  elements ;  
and  $k$  pair elements.

$\rightarrow dp[i][j][k] = \max(dp[i-1][j-1][k], b[i] + dp[i-1][j][k], dp[i-2][j][k-1])$

Actual version :-  $P \neq 0$   $Q \neq 0$   $R \neq 0$

$\rightarrow dp[i][j][k][l]$  = consider first " $i$ " elements but ignore  $j$  elements ;  
and  $k$  pair elements and  $l$  triplet elements;[all adjacent type;]

$\rightarrow dp[i][j][k][l] = \max(dp[i-1][j-1][k][l], b[i] + dp[i-1][j][k][l], dp[i-2][j][k-1][l], dp[i-3][j][k][l-1])$

TC -  $O(NPQR)$

Takes  $O(NPQR)$  size;

C++ <https://www.jdoodle.com/ia/1gh7>.

Java. <https://www.jdoodle.com/ia/1gh8>

Py <https://www.jdoodle.com/ia/1ghb>



