

SACHIN GOYAL  
150020069

Q4.  $G$  is a PRG. ~~There is no Discriminator  $D$  such that~~

$$\Rightarrow \Pr[D(x) = 1] - \Pr[D(G(x)) = 1] \leq \text{negligible}$$

for all Discriminators  $D$ .

$$G_1(s) = G(s, s_2, \dots, s_{|s|-1}) \parallel s_{|s|}$$

~~\* Assume  $G_1(s)$  is not a~~

New  $G_1(s \parallel b) = G(s) \parallel b$  — (1)  $b$  is a random bit.

~~\* Since  $G_1$  is a~~

\* Assume  $G_1(s)$  is not pseudo random generator

$\Rightarrow$  There exists  $D_1$  for  $G_1$  such that

$$\Pr[D_1(x) = 1] - \Pr[D_1(G_1(s)) = 1] > (n)^{-k} \quad (2)$$

(assuming probabilistic polynomial time (PPT) adversaries).

Main  $\rightarrow$  Let us construct a subroutine ~~for~~ for  $D_1$ .

~~$D_2(x) = D_1(x \parallel b)$  where  $b$  is a random bit~~

Subroutine  $D(x) = D_1(x \parallel b)$  where  $b$  is a random bit

$$\therefore \Pr[D(x) = 1] - \Pr[D(G(s)) = 1]$$

$$= \Pr[D_1(x || v) = 1] - \Pr[D_1(G(s) || v)]$$

$$= \Pr[D_1(x') = 1] - \Pr[D_1(G_1(s || v))]$$

$$[\text{since } G_1(s || v) = G(s) || v]$$

$$= \Pr[D_1(x') = 1] - \Pr[D_1(G_1(s'))]$$

(From our assumption of  $G_1$  not being PRG (eq<sup>n</sup> 2))

$$\Rightarrow \Pr[D_1(x') = 1] - \Pr[D_1(G_1(s'))] > (n+1)^{-k}$$

$$\Rightarrow \Pr[D(x) = 1] - \Pr[D(G(s)) = 1] > (n+1)^{-k}$$

not negligible.

$\Rightarrow$  CONTRADICTION

$\Rightarrow$  WRONG ASSUMPTION

$\therefore G_1$  is a pseudo random generator



Q1

~~me so~~

Encryption is perfectly secret if

$$\Pr[C=c | M=m_i] = \Pr[C=c | M=m_j]$$

$\forall i, j \ i \neq j$

$$E_{m_K}(m) = (K+m) \bmod 5$$

① take  ~~$m=0$~~   $m=0$   
 $C=0$  for  $K=0$  &  $K=5$

$$\therefore \Pr[C=0 | M=0] = 2/5$$

② take  $m=2$

$$\therefore C=0 \text{ for } K=3$$

$$\therefore \Pr[C=0 | M=2] = 1/5$$

$$\therefore \Pr[C=0 | M=0] \neq \Pr[C=0 | M=2]$$

NOT A PERFECTLY SECURE

Q2

Again  $\Pr[C=c | M=m_i] = \Pr[C=c | M=m_j]$

Let  $l=2$ ,  $C=11$

$\forall i, j \ i \neq j$

①  $\Pr[C=11 | M=00] = \frac{1}{4}$

$$K=11$$

$$\Pr[c = 11 | m = 01] = 0$$

because  $K = 10$   
not valid

$\therefore$  Not perfectly secure

[Q3]

~~to~~  $\text{negl}_1$  is negligible

$$\Rightarrow \text{negl}_1(n) < \frac{1}{K(n)} \quad \forall n > N$$

$$\text{let } \text{negl}_1(n) < \frac{1}{p(n)q(n)} \quad \forall n > N$$

where  $K(n)$  is some polynomial

where  $q(n)$  is some positive polynomial

$$\rightarrow \text{negl}_2 = p(n) \text{negl}_1(n)$$

$$\rightarrow \text{negl}_2(n) < \frac{p(n)}{p(n)q(n)} = \frac{1}{q(n)} \quad \forall n > N$$

$\Rightarrow \text{negl}_2(n)$  is also negligible