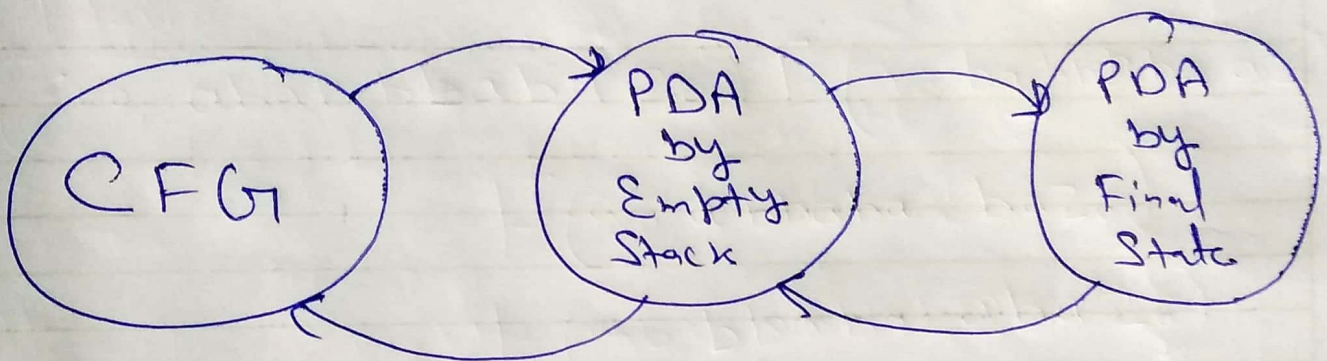


CFG to PDA :**PDA**

There are three classes of Languages accepted by PDA.

- 1) CFL defined by CFG
- 2) Language accepted by Empty Stack
- 3) " " " Final State.



$$G = (V, \Sigma, P, S) \Rightarrow P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$$

$$Q = \{q_i\}$$

$$\Sigma = \{ \text{Input symbol (Terminals)} \}$$

$$\Gamma = (V \cup \Sigma)$$

δ is defined by

$$q_0 \rightarrow q$$

$$\begin{array}{l} q_0 \rightarrow q \\ Z_0 \rightarrow S \\ F \rightarrow \phi \end{array}$$

i) For each non Terminal

$$\delta(q, v, A) = \{ (q', \alpha) \mid A \rightarrow v\alpha \}$$

$$Z_0 \rightarrow S$$

$$F \rightarrow \phi$$

$$\delta(q, v, A) = \{ (q', \alpha) \mid A \rightarrow v\alpha \}$$

is in Production

2) For each Terminal **PDA**

(16)

$$\begin{aligned} \delta(q, a, a) &= \{(q, \wedge)\} \mid a \in \Sigma \\ \delta(q, b, b) &= \{(q, \wedge)\} \mid b \in \Sigma \end{aligned}$$

Ex 1:- Construct PDA equivalent to following CFN.

$$S \rightarrow 0BB$$

$$B \rightarrow 0S \mid 1S \mid 0$$

$$\underline{\text{Sol}^n} = G = (\{S, B\}, \{0, 1\}, P, S)$$

$$\Rightarrow (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$$

$$\Rightarrow \{q, \{0, 1\}, \{S, B, 0, 1\}, \delta, q, S, \emptyset\}$$

1) For each Non-terminal:

δ :-

$$\text{i) } \delta(q, \wedge, S) = \{(q, 0BB)\}$$

$$\text{ii) } \delta(q, \wedge, B) = \{(q, 0S), (q, 1S), (q, 0)\}$$

2) For each Terminal:

$$\text{iii) } \delta(q, 0, 0) = \{(q, \wedge)\}$$

$$\text{iv) } \delta(q, 1, 1) = \{(q, \wedge)\}$$

In PDA There are 4 Transition.

Ans

PDA

(17)

Ex2: Construct PDA equivalent following grammar.

$$S \rightarrow 0S1 \mid A$$

$$A \rightarrow 1A0 \mid S \mid \epsilon$$

Soln: PDA = $(\{q\}, \{0, 1\}, \{S, A, 0, 1\}, \delta, q, S, \phi)$

* There is no Final state because we construct acceptance of Empty Stack.

δ !:- For ^{Each} Non-Terminal.

$$\delta(q, \wedge, S) = \{(q, 0S1), (q, A)\}$$

$$\delta(q, \wedge, A) = \{(q, 1A0), (q, S), (q, \epsilon)\}$$

δ !:- For ^{Each} Terminal:

$$\delta(q, 0, 0) = \{(q, \wedge)\}$$

$$\delta(q, 1, 1) = \{(q, \wedge)\}$$

Ans

PDA

(12)

Ex 3: Construct PDA equivalent to following CFN.

$$S \rightarrow aABB \mid aAA$$

$$A \rightarrow aBB \mid a$$

$$B \rightarrow bBB \mid A$$

Soln: PDA = $(\{q\}, \{a, b\}, \{S, A, a, b\}, \delta, q, S, \phi)$

δ : For each Non Terminal:

$$\delta(q, \Lambda, S) = \{(q, aABB), (q, aAA)\}$$

$$\delta(q, \Lambda, A) = \{(q, aBB), (q, a)\}$$

$$\delta(q, \Lambda, B) = \{(q, bBB), (q, A)\}$$

δ : For each Terminal:

$$\delta(q, a, a) = \{(q, \Lambda)\}$$

$$\delta(q, b, b) = \{(q, \Lambda)\}.$$

Ans