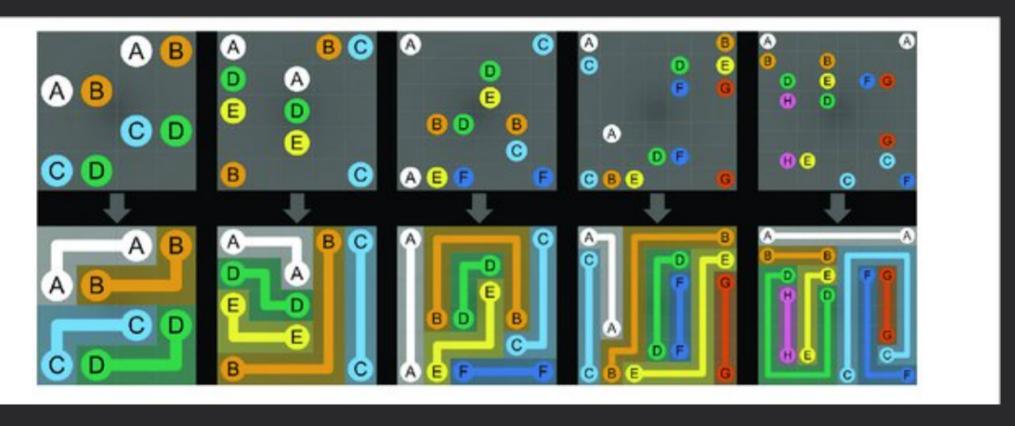
Flow Free - a circuit design problem



In this programming assignment you'll be expected to build a solver for the Flow Free game. Its origins can be traced back to 1897, a puzzle published by the professional puzzler Sam Loyd, in a column he wrote for the *Brooklyn Daily Eagle (more details in the Matematica Journal).* This puzzle was popularised in Japan under the name of Numberlink.

Game Rules



The game presents a **grid** with <u>colored dots</u> occupying some of the squares. The objective is to <u>connect dots of the same color</u> by drawing *pipes* between them <u>such that the entire grid is occupied</u> by pipes. However, <u>pipes may not intersect</u>.

The difficulty is determined by the size of the grid, ranging from 5x5 to 15x15 squares.

The Algorithm

Each possible configuration of the Flow Free grid is called a *state*. Each position in the grid can be free or contain a color. The Flow Free Graph $G=\langle V,E\rangle$ is implicitly defined, i.e. the graph is not completely loaded in memory at once, instead, it starts only with the initial state, and the remainder of the graph is discovered while searching for a solution. The vertex set V is defined as all the possible configurations (states), and the edges E connecting two vertexes are defined by the legal movements (**right**, **left**, **up**, **down**) for **each color**. The code provided makes sure that only legal moves are generated:

<u>A move is legal only if it extends a color pipe</u>. A color cannot be placed if it's not adjacent to the pipe of the same color. Furthermore, to decrease the number of edges in the Graph, the pipe can only be started from one of the two initial colors, making one of them the start and the other the goal of the pipe.

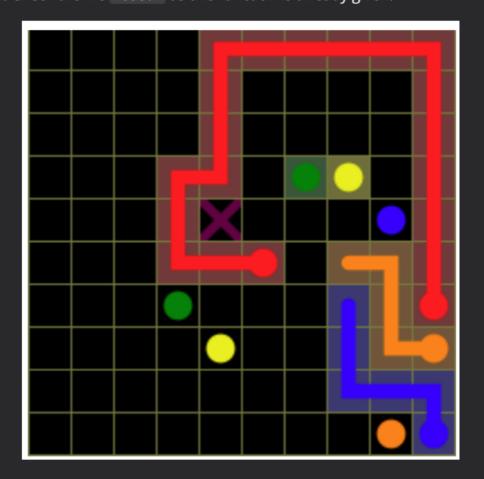
All edges have a weight of 1.

Your task is to find the path traversing the Flow Free Graph from the initial state (vertex) leading to a state (vertex) where all the grid cells contain a color, and therefore finding a solution to the puzzle. A path is a sequence of movements painting a new cell in the grid. You are going to use Dijkstra to find the shortest path first, along with some game-specific optimizations to speed up your algorithm.

The initial node (vertex) corresponds to the initial configuration (Algorithm 1, line 1). The algorithm selects a node to expand (generate the possible pipe extensions, line 5) along with an ordering of which is the *next color* the algorithm should try to connect (line 6). Therefore, when a node is expanded, it has at most 4 children (right, left, up, down). Once all the children have been added to the priority queue, the algorithm will pick again the next node to extend (line 5), as well as the *next color to consider*, until a solution is found. Different color orderings can be tried with the command line options (type ./flow -h to see all options). Each ordering has a different impact on the number of nodes needed by Dijkstra to find a solution. By default, the search will expand a maximum of 1GB of nodes (line 12). This is checked by a variable named max_nodes.

```
GRAPHSEARCH(Graph, start)
     node \leftarrow start
 1
 ^{2}
     pq \leftarrow priority\ Queue\ Containing\ start\ node\ Only
 ^{3}
     while frontier \neq empty
 4
     do
 5
        node \leftarrow pq.pop()
        nextColor \leftarrow select \ next \ color \ in \ ordering
 6
 7
        for each move direction
 8
        do
 9
            if move direction for nextColor is valid
10
               then
11
                     newNode \leftarrow applyMoveDirection(node)
12
                     if number of nodes in pq \geq max\_nodes
13
                        then
                              result = SEARCH\_FULL
14
15
                              break
16
17
                     if newNode is a dead end
18
                        then
19
                              deleteNewNode
20
                              continue
21
22
                     if newNode is the solution
23
                        then
24
                              solution = NewNode
25
                              result = SEARCH\_SUCCESS
26
                              break
27
28
                     pq.push(newNode)
29
30
31
     free priority queue
     return result
32
```

When you applyMoveDirection you have to create a new node that points to the parent, updates the grid resulting from applying the direction chosen, updates the priority of the node, and updates any other auxiliary data in the node. The priority of the node is given by the length of the path from the root node, i.e. how many grid cells have been painted. Check the file node.h as this function is already given.



A *dead-end* is a configuration for which you know a solution cannot exist. The figure above shows an example of a dead-end: the cell marked with X cannot be filled with any color. The problem is unsolvable with the current configuration, and we do not need to continue searching any of its successors. You can <u>recognize dead-end</u> cells by looking for a **free cell** in the grid that is **surrounded by three completed path segments (colored) or walls**. The current position of an in-progress pipe and goal position have to be treated as free space. Detecting unsolvable states early can prevent lots of unnecessary search. The search will eventually find a solution. If we recognize dead-ends and do not generate their successors (pruning), the search for a solution will likely take less time and memory, as you'll avoid exploring all possible successors of an unsolvable state.

You might have multiple paths leading to a solution. **Your algorithm should consider the possible actions in the following order**: left (0), right (1), up (2) or down (3).

<u>Make sure you manage the memory well.</u> When you finish running the algorithm, you have to free all the nodes from the memory, otherwise you will have memory leaks. You will notice that the algorithm can run out of memory fairly fast after expanding millions of nodes.

Check the files utils.h and queue.h where you'll find many of the functions in the algorithm already implemented. Other useful functions are located directly in the file search.c, which is the only file you need to edit to write your algorithm inside the function game_dijkstra_search. For dead-end detection, you need to look at extensions.c, function game_check_deadends. Look for the comment FILL IN THE CODE. All the files are in the folder src/.

Deliverable 1 - Dijkstra **Solver** source code

You are expected to hand in the source code for your solver, written in C. Obviously, your source code is expected to compile and execute flawlessly using the following makefile command:

make

generating an executable called

flow

Remember to compile using the optimization flag <code>gcc -03</code> for doing your experiments, it will run twice as quickly as compiling with the debugging flag <code>gcc -g</code> (see <code>Makefile</code>). The provided Makefile compiles with the optimization flag by default, and with the debugging flag if you type <code>make debug=1</code>. For the submission, please <code>do not remove the -g option from your Makefile</code>, as our scripts need this flag for testing. Your program must not be compiled under any flags that prevent it from working under <code>gdb</code> or <code>valgrind</code>.

Your implementation should be able to solve the *regular* puzzles provided. To solve the *extreme* puzzles, you'll need further enhancements that go beyond the time for this assignment, but feel free to challenge yourself if you finish early and explore how you would solve the extreme puzzles.

The Code Base

You are given a base code. You can compile the code and execute the solver by typing ./flow <puzzleName> .
You are going to have to program your solver in the file search.c. Look at the file and implement the missing part in the function called game_dijkstra_search. Once you implement the search algorithm, go to the file called extensions.c and implement the function called game_check_deadends

You are given the structure of a node in <code>node.*</code> files, and also a priority queue <code>queues.*</code> implementation. Look into the <code>engine.*</code> and <code>utils.*</code> files to know about the functions you can call to perform the search.

In your final submission, you are free to change any file, but make sure the command line options remain the same.

Input

In order to execute your solver use the following command:

```
./flow [options] <puzzleName1> ... <puzzleNameN>
```

for example:

```
./flow puzzles/regular_5x5_01.txt
```

Will run the solver for the regular 5 by 5 puzzle, and report if the search was *successful*, the *number of nodes* generated and the *time* taken. if you use flag -q (quiet) it will report the solutions more concisely. This option can be useful if you want to run several puzzles at once and study their performance.

If you append the option <code>-A</code> it will *animate* the solution found. If you append the option <code>-d</code> it will use the *dead-end detection* mechanism that you implemented. Feel free to explore the impact of the other options, specifically the *ordering* in which the colors are explored. By default, the color that has fewer free neighbors (most constrained), is the one that is going to be considered first.

All the options can be found if you use option -h:

```
$./flow -h
usage: flow_solver [ OPTIONS ] BOARD1.txt
BOARD2.txt [ ... ] ]
Display options:
 -q, --quiet
                          Reduce output
  -d, --diagnostics
                          Print diagnostics when search unsuccessful
 -A, --animation
                          Animate solution
 -F, --fast
                          Speed up animation 4x
 -C, --color
                          Force use of ANSI color
 -S, --svg
                          Output final state to SVG
Node evaluation options:
 -d, --deadends
                          dead-end checking
Color ordering options:
 -r, --randomize
                          Shuffle order of colors before solving
  -c, --constrained
                          Disable order by most constrained
Search options:
  -n, --max-nodes N
                          Restrict storage to N nodes
                          Restrict storage to N MB (default 1024)
  -m, --max-storage N
Help:
 -h, --help
                          See this help text
```

Output

Your solver will print the following information if option -q is used:

- 1. Puzzle Name
- 2. SearchFlag (see utils.c, line 65-68 to undestand the flags)
- 3. Total Search Time, in seconds
- 4. Number of generated nodes
- 5. A final Summary

For example, the output of your solver ./flow -q ../puzzles/regular_* could be:

```
../puzzles/regular_5x5_01.txt s 0.000 18
../puzzles/regular_6x6_01.txt s 0.000 283
../puzzles/regular_7x7_01.txt s 0.002 3,317
../puzzles/regular_8x8_01.txt s 0.284 409,726
../puzzles/regular_9x9_01.txt s 0.417 587,332
5 total s 0.704 1,000,676
```

These numbers depend on your implementation of the search, the ordering you use, and whether you prune dead-ends. If we use dead-end pruning $./flow -q -d ../puzzles/regular_*$ we get the following results

```
../puzzles/regular_5x5_01.txt s 0.000 17
../puzzles/regular_6x6_01.txt s 0.000 254
../puzzles/regular_7x7_01.txt s 0.001 2,198
../puzzles/regular_8x8_01.txt s 0.137 182,136
../puzzles/regular_9x9_01.txt s 0.210 279,287
5 total s 0.349 463,892
```

Remember that in order to get full marks, your solver has to solve at least the regular puzzles.