Problem 2: GodBolt [24 points]

Scribe: [Scribe's name here]

Go to <u>godbolt.org</u>. This is a website that will compile source code from several different programming languages into one of many target assembly languages, and color code the correspondence between the two programs.

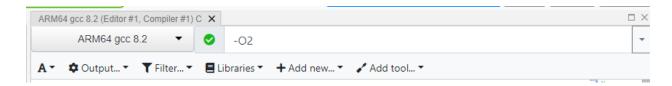
In the left hand window, select "C" is the source language and paste in the following code:

```
#include <stdlib.h>
void scale_array(unsigned long long *arr, unsigned long size) {
   int i=0;
   while(i<size) {
      unsigned long long temp = arr[i];
      temp *= 100;
      arr[i] = temp;
      i++;
   }
}</pre>
```

Your window should look like this (with perhaps rows highlighted as different colors):

```
COMPILER
                        Add... ▼
                                                                         Become a Perfor
                                More ▼
     EXPLORER
C source #1 X
                                                                                   \square \times
                                                              C
A ▼ B Save/Load + Add new... ▼ Vim
      #include <stdlib.h>
  2 ∨ void scale_array(unsigned long long *arr, unsigned long size) {
          int i=0;
  3
           while(i<size) {
  4
  5
               unsigned long long temp = arr[i];
               temp *= 100;
  6
  7
               arr[i] = temp;
  8
               i++;
  9
 10
```

In the right hand window, select "carm64g820" or "ARM64 gcc 8.2" for the architecture and "-O2" for the compiler options:



After a moment, the right window should be populated with ARM assembly code. You can see which C statements each instruction corresponds to by matching the colors between the two windows.

Because this compiler is using the full Arm assembly language (not just the LEG subset described in lecture), there are a few instruction variants you may be unfamiliar with (note that we do **NOT** expect you to be familiar with these instructions outside the scope of this problem). The first is:

```
add x1, x2, x3, lsl 4
```

This instruction combines an add instruction and an 1s1 instruction, which could be translated to:

```
x1 = x2 + (x3 << 4)
```

The second new instruction looks like:

```
str x0, [x1], 4
```

This instruction simply combines a store instruction with an add instruction. It is equivalent to:

```
sturw x0, [x1,#0]
add x1, x1, #4
```

1. In the space, below paste the compiled assembly and explain how each instruction is implementing the specified C code. In particular, describe how the assembly multiplies an element by 100 without the use of an explicit multiply instruction. [12 points]

```
cbz
              x1, .L1
                                 // branch to end if size is zero
              x1, x0, x1, lsl 3 // x1 points to end of array (&arr[size])
       add
.L3:
                                 // start of loop
       ldr
              x3, [x0]
                       // x3 = arr[i]
              x2, x3, x3, 1s1 1 // x2 = arr[i]*3
       add
       add
              x2, x3, x2, 1s1 3 // x2 = x2*8+arr[i] (i.e. <math>arr[i]*25)
       lsl
              x2, x2, 2
                                // x2 = x2*4 (i.e. arr[i]*100)
                               // arr[i] = arr[i]*100; i++
       str
              x2, [x0], 8
              x0, x1
                                 // if at end of array...
       cmp
              .L3
                                 // ...don't execute loop again
       bne
.L1:
       ret
                                 // return
```

Change the source code above so that rather than multiplying each element of the array by 100, it instead multiplies each element by 127.

2. Briefly (30 words or fewer) describe the change in the resulting assembly and how it accomplishes the multiplication. [12 points]

Rather than doing several shifts and adds, multiplying by 127 can be done by shifting to the left 7 (i.e. multiplying by $2^7=128$) and then subtracting the original value once (i.e. x*127 == x*128 - x).

Problem 3: Putting it All Together [24 points]

Scribe: [Scribe's name here]

Convert the following C code to ARM. Assume the system requires memory alignment, and that the assembler will not fix opcodes. Use labels when possible. The return value should be in register X2. Assume:

- Starting address of the **data** array is stored in register **X0**, and consists of **512 elements**.
- Assume all operations will be within the 32 bit unsigned integer limits.

Code explanation (you don't need to read this to answer the question):

The ABC370 factory is responsible for producing packs of 20 pencils. However, this factory is not always perfect, which means some packs will contain less than the desired amount of 20 pencils. In the case where there are less than 20 pencils per pack, we simply multiply by 4. Return the number of pencil packs not changed in register X2.

These questions will help you fill out the below [10 points]:

- How much padding (in bytes) is needed between the pack_id array and num_penc?
 4 bytes
- 2. How much padding is needed at the end of the pencil_pack struct?
 0
- How large is each pencil_pack struct instance?
 32 bytes
- 4. Where can the address of "data" be found at the start of the check_pencil_arr, assuming we're following the ARM ABI? X0 (from question description)
- 5. Given the skeleton code below, determine what register is holding:
 - a. i X3
 - b. good_count X2 (from question description)
 - c. data[i].num penc X5

```
C
                                                               ARM
struct pencil_pack{
    char pack_id[20]; //20B
    int64 t num penc; //8B
};
int32 t
check_pencil_arr(
                                             func:
           struct pencil_pack data[]) {
                                                   MOVI
                                                           X2,
                                                                  #0x00
 int32_t good_count = 0;
                                                   MOVI
                                                           Х3,
                                                                  #0x00
                                             loop:
 for(int32_t i = 0; i < 512; ++i){
                                                   CMPI
                                                           Х3,
                                                                  #512
     if(data[i].num_penc < 20){</pre>
                                                   B.EQ
                                                           end
                                                   LSL
                                                           X4,
                                                                   Х3,
                                                                         #5
                                                           Х4,
                                                                   Х4,
                                                                          X0
         data[i].num_penc *= 4;
                                                   ADD
                                                           X5,
                                                                         #24]
                                                   LDUR
                                                                  [X4,
                                                           X5
                                                   CMPI
                                                                  #20
     } else {
        ++good_count;
                                                   B.LT
                                                           smaller
     }
                                                   ADDI
                                                           Х2,
                                                                   X2,
                                                                         #1
  }
                                                           forend
                                             smaller:
                                                                        #2
                                                   LSL
                                                          X5,
                                                                  X5,
 return good_count;
                                                   STUR
                                                                 [X4,
                                                                        #24]
                                                          X5
                                             forend:
                                                   ADDI
                                                           X3
                                                                  X3
                                                                       #1
                                                       loop
                                             end:
```

Note:

- Indexing into an array can be thought of as pointer arithmetic arr[i] = *(arr + i)
- Since each struct takes 32 bytes, this means the **ith** element of pencil_pack is at a memory address of **pencil_pack** (base address) + 32 (struct size in Bytes) * i (iter)
- From lab 2, we know that a << b = a * 2^b (i.e. a * 32 = a * 2⁵ = a << 5)