## FIT3155 - Assignment 1

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## Comments to the marker: Q1

This implementation behaves similarly to if the pattern was appended to both the front and end of the reference text and if the z arrays were computed accordingly

## pat\$ref\$pat

The z prefix and z suffix values can be computed for **each index in the reference text** as per the above setup using a standard z algorithm.

Then each m length region of the reference text is considered (where m is the length of the pattern) and if  $m - (z_i + z_suffix_i) = 2$  and if the characters swapped were a match, it would detect a transposition in the text.

The complete matches are found by just checking cases where z\_i = m

However, the implementation makes use of a number of optimizations to implement the above logic efficiently

- 1. Uses a basic z algorithm on the pattern to compute the z prefix and z suffix for the **pattern** only. This takes away the need to slice the array later, to separate the z values for the reference only.
- 2. Then uses this information to compute the z array for the prefix and the suffix of the **reference** with respect to the pattern in a single loop. Saving time and space
- The reference text is looped over and for every m letters, it is checked if
  Z\_i = m (corresponding to a full match)
  m (z\_i + z\_suffix\_i) = 2 (Only two characters mismatch) and if the mismatched characters match if swapped, this would correspond to a transposition.

The implementation uses O(n+m) space and time, where n is the length of the reference text and m is the length of the pattern.