Seminar P	rogrammie	ersprache	en im Ver	clgeich - Z	Zig
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1. Introduction

The objective of this seminar was getting to know a new language by writing a simple grep like program.

2. Language

- manual memory management
 - allocators have to be manually passed
 - deallocate by defering deinit procedures
 - memory leaks are automatically detected in debug mode
- · minimal standard library
 - no unicode string support, only functions for operating on slices e.g. std.mem
 - ▶ utf-8 string libraries are 3rd party
- exhaustive switch statements
 - useful for enums
- generics are just comptime functions operating on types
- inferred struct literal type . {}
- if enum type is known the type can be omitted, .variant is sufficient
- · errors as values
 - error union explicit or implict
 - ► try: return on error
 - ► catch: handle error
- somewhat immature ecosystem
 - missing regex library
 - ▶ async not available in 0.11 self-hosted compiler

2.1. Introduction

Zig is a general-purpose compiled systems programming language. Andrew Kelley 2015/2016

It is often mentioned as a successor to C.

2.2. Toolchain

Installation was as simple as downloading the release tar archive, extracting the toolchain, and symlinking the binary onto a \$PATH.

The compiler is a single executable called zig, it includes a build system which can be configured in a build.zig file. The configuration is written in zig itself and thus avoids a separate language just for the build system. The toolchain itself is also a C/C++ compiler which allows zig code to directly interoperate with existing C/C++ code from zig.

To start a new project inside an existing directory running zig init-exe will generate the main source file src/main.zig and the build configuration build.zig. The program can then be built and executed by running zig build run.

The zig community also provides a language server called zls, which worked out of the box in neovim. There were some issues with type inference, and completion of member functions of generic types.

2.3. Integers

Zig has concrete integer types with a predetermined size regardless of platform. Common ones are:

- unsigned: u8, u16, u32, u64
- signed: i8, i16, i32, i64.

But it also allows defining arbitrary sized integers like i27, or u3.

Zig arrays have a size known at compile time and are stack allocated, unless specified otherwise:

```
1 const explicit_length = [5]u32{ 0, 1, 2, 3, 4 };
2 const inferred_length = [_]u32{ 5, 6, 7, 8, 9 };
```

2.4. Arrays and slices

Arrays can be sliced using the index operator with an end-exclusive range, returning a slice to the referenced array. By default the slice is a fat pointer which includes a pointer that points to the memory address of the slice inside the array, and the length of the slice.

```
1 const array = [_]u32{ 0, 1, 2, 3, 4 };
2 const slice = array[1..3];
3 std.debug.print("{d}\n", .{slice.*});
```

The output prints:

```
1 { 1, 2 }
```

For better interoperability with C zig also allows sentinel terminated pointers or slices. Most commonly this is used for null-terminated strings:

```
const string: *const [32:0]u8 = "this is a null terminated string";
const fat_pointer_string_slice: []const u8 = string[10..];
const null_terminated_string_slice: [:0]const u8 = string[10..];

const null_terminated_string_ptr: [*:0]const u8 = string[10..];

std.debug.print("size: {}, '{s}'\n", .{@sizeOf([]const u8), fat_pointer_string_slice});

std.debug.print("size: {}, '{s}'\n", .{@sizeOf([:0]const u8), null_terminated_string_slice});

std.debug.print("size: {}, '{s}'\n", .{@sizeOf([*:0]const u8), null_terminated_string_ptr});
```

All three slice/pointer types reference the same data, but in a different way. Variant 1 uses the fat pointer approach described above. Variant 2 also uses the same approach but also upholds the condition that the end of the slice is terminated by a null-byte sentinel. Variant 3 only stores a memory address and relies upon the null-byte sentinel to compute the length of the referenced data when needed. As a result the first and the second slice type have a size of 16 bytes, 8 bytes for the pointer and 8 additional bytes for the length. The sentinel terminated pointer only has a size of 8 bytes, since it doesn't store an additional length field. A limitation of sentinel terminated slices or pointers is that they cannot reference arbitrary parts of an array. Trying to do so fails with the following message:

2.5. Eco system

Compared to C++, Java, or Rust the std library of zig is quite minimal. Neither the zig language itself, nore the std library provide a datatype for strings. String literals are represented a byte slices ([]const

u8), which allows using the whole range of std.mem.* functions to operate on them. There is no unicode support. There are some beginnings of people implementing regex libraries, but none that are ready to be used right now. So I decided on using the rust regex library (rure) through it's C API, since it's a standalone project and not part of some standard library. To include a C library, some modification inside the build.zig configuration file are needed:

```
1 // link all the other stuff needed
2 exe.linkLibC();
3 exe.linkSystemLibrary("util");
4 exe.linkSystemLibrary("dl");
5 exe.linkSystemLibrary("gcc_s");
6 exe.linkSystemLibrary("m");
7 exe.linkSystemLibrary("rt");
8 exe.linkSystemLibrary("util");
9 // link rure itself
10 exe.addIncludePath(LazyPath.relative("rure/regex-capi/include"));
11 exe.addLibraryPath(LazyPath.relative("rure/target/release"));
12 exe.linkSystemLibrary2("rure", .{ .needed = true, .preferred_link_mode = .Static });
```

The dependencies are taken straight from the rure compile script:

```
# N.B. Add `--release` flag to `cargo build` to make the example run faster.
cargo build --manifest-path ../Cargo.toml
gcc -03 -DDEBUG -o iter iter.c -ansi -Wall -I../include -L../../target/debug -
lrure
# If you're using librure.a, then you'll need to link other stuff:
# -lutil -ldl -lpthread -lgcc_s -lc -lm -lrt -lutil -lrure
```

When linking C libraries, zig isn't able to include debug symbols, so crash messages that would normally be informative, only show memory addresses:

```
1 thread 20843 panic: index out of bounds: index 14958, len 14948
2 Unwind error at address `:0x2ebaef` (error.InvalidDebugInfo), trace may be incomplete
```

This is a known bug.

2.6. Compiler bug

With zig 0.11.0 I encountered a bug in the compiler which would affect command line argument parsing. In debug mode arguments were parsed fine, bug in release mode the --ignore-case flag would be parsed as the --hidden flag. All flags are defined as an enum:

```
const UserArgFlag = enum {
    Hidden,
    FollowLinks,
    Color,
    NoHeading,
    IgnoreCase,
    Debug,
    NoUnicode,
}
```

```
10 Help, };
```

The issue was fixed by specifying a concrete type to represent the enum instead of letting the compiler infer the type.

```
@@ -27,12 +27,12 @@ const UserArgKind = union(enum) {
2
        flag: UserArgFlag,
3
    };
4
   -const UserArgValue = enum {
5
   +const UserArgValue = enum(u8) {
        Context,
        AfterContext,
8
9
        BeforeContext,
10
    };
   -const UserArgFlag = enum {
11
   +const UserArgFlag = enum(u8) {
12
13
        Hidden.
        FollowLinks,
14
        Color,
```

3. Development process

Since the scope of the program was pre-determined, the main focus was on performance.

The zig std library provides IterableDir, an iterator for iterating a directory in a depth first manner, but unfortunately that approach doesn't allow filtering of searched directories. To overcome that limitation I mostly copied the std library function for iterating directories and modified it slightly to allow filtering out hidden directories.

3.1. Single threaded optimization

3.1.1. Line by line matching

To keep it simple the first implementation, read the whole file into a single buffer and run a compiled regex pattern match on every line.

3.1.2. Whole text matching for performance

- to avoid allocating iterator for every line
- then avoid using iterator anyway

After some investigation it turned out that initializing the regex search had more overhead than expected, and running the regex pattern match on the whole text instead of every line would improve performance significantly.

3.1.3. Line by line matching with fixed size buffer

- no multiline matches
- easier to implement
- lines need to be iterated anyway for line numbers

Since one of the tests was to search an 8gb large text file, the input would need to be split up into smaller chunks as to avoid running out of memory. This was done using a fixed size buffer which would only load part of the file, searching that buffer upto the last fully included line, then moving

the unsearched parts including possibly relevant context lines to the start of the buffer, and eventually refilling the buffer with remaining data to search.

3.1.4. Whole text matching for performance with fixed size buffer

- the rust regex find functions has extremely high startup overhead on unicode word-character patterns
- avoided to some extend by searching the whole text buffer instead of just line by line slices

To simplify the previous change, the algorithm would once again search the text line by line. So to restore previously achieved performance, the regex pattern matching was once again adjusted to be run on the whole text buffer.

3.2. Parallelization

TODO: mention ripgrep influence

At this point most easy wins in single threaded performance were off the table, so the next easy performance gain would be using multiple threads. The main things that take up time are reading from the file system, and searching the text.

Parallelization was implemented using a thread pool of search workers that would receive file paths using a atomically synchronized message queue. The directory walking remained mostly the same apart from searching files adhoc, they were now sent through the message queue.

The output of multiple threads now had to be synchronized so that lines printed from one file would not be interspersed with other ones. There are two obvious solutions to this problem. One is to use a dynamically growing allocated buffer which stores the entire output of a searched file and then write the entire buffer in a synchronized way when the file is fully searched. This would avoid blocking other threads, but could cause the program to run out of memory if large portions of big files would match a search pattern. The other solution is to just block output of all other threads once a match has been found in a file and then write all lines directly to stdout. This would avoid running out of memory, but could in worst case scenarios cause basically single threaded performance. The final implementation uses a hybrid of the two, each thread has a fixed size output buffer which can be written to without any synchronization. Once the buffer is full access to stdout is locked for other threads until the file is fully searched, but other workers can still access their thread-local output buffers.

Text searching had been parallelized the search workers were now emptying the message queue too quickly, so walking the file system with multiple threads was up next. TODO: Walking the file tree uses a shared stack for depth first searching

4. Conclusion