

Observations: $b^n - 1 = (b-1)(b^{n-1} + b^{n-2} + \dots + 1)$

$$b^{mn} - 1 = (b^m - 1)(b^{m(n-1)} + \dots + 1)$$

$$x^n - 1 = (x-1)(x^{n-1} + \dots + 1)$$

$$x^{mn} - 1 = (x^m - 1)(x^{m(n-1)} + x^{m(n-2)} + \dots + 1)$$



no roots of $x^{m(n-1)} + x^{m(n-2)} + \dots + x^m + 1 = g(x)$
is an m th root of unity.

Q. Prove if $d = \gcd(m, n)$ & $a > 1$ an integer, then
 $\gcd(a^m - 1, a^n - 1) = a^d - 1$.

Soln: set $m = dm_1$, $n = dn_2$

$$a^m - 1 = (a^d - 1) f_{m_1}(a) \quad \& \quad a^n - 1 = (a^d - 1) f_{n_1}(a)$$

$$f_{m_i}(a) = a^{d(m_i-1)} + a^{d(m_i-2)} + \dots + 1 \quad i=1,2$$

$$\gcd(f_{m_1}(a), f_{n_1}(a)) = 1 ?$$

$$\gcd(f_{m_1}(x), f_{n_1}(x)) = 1$$

Observe

Roots of $f_{m_1}(x)$ are $m = m_1 d$ roots of unity which
not d roots of unity

|| $f_{n_1}(x)$ are $n = n_1 d$

$$s f_{m_1}(a) + t f_{n_1}(a) = 1$$

for some $s, t \in \mathbb{Z}$
(e.g. when 1
is the gcd
of $f_{m_1}(a)$ & $f_{n_1}(a)$).

$$\begin{aligned} & \rightarrow s(x) f_{m_1}(x) + t(x) f_{n_1}(x) = 1 \quad \text{for } m, n, s, t \in \mathbb{Z}[x] \\ & \text{put } x=a \\ & s(a) f_{m_1}(a) + t(a) f_{n_1}(a) = 1 \quad (\text{initially they are in } \mathbb{Q}[n]) \\ & \text{so that } f_{m_1}(a) \text{ \& } f_{n_1}(a) \text{ are co-prime.} \end{aligned}$$

Silver Pohlig - Hellman exponentiation cipher.

let p be prime & $e \in \mathbb{N}$ s.t. $0 < e < p$ &
 $\gcd(e, p-1) = 1$ \rightarrow called exponent.

[E.g., $p=29$, $e=3, 5, \dots$]

$$\begin{aligned} \mathcal{E} : \mathbb{Z}_p &\rightarrow \mathbb{Z}_p \\ m &\mapsto m^e \end{aligned}$$

$$\begin{aligned} \mathcal{D} : \mathbb{Z}_p &\rightarrow \mathbb{Z}_p \\ m &\mapsto m^d \end{aligned}$$

$$\text{so that } \mathcal{D}(\mathcal{E}(m)) = m \quad \underbrace{d=?}$$

$$|U(\mathbb{Z}_p)| = p-1$$

$$\begin{aligned} \therefore (m^e)^d &= m \\ \therefore m^{de} &= m \end{aligned}$$

$$\therefore \boxed{de \equiv 1 \pmod{p-1}}$$

$$\text{since } \gcd(e, p-1) = 1 \quad \exists \text{ s \& t}$$

$$\text{s.t. } \boxed{se + t(p-1) = 1}$$

$$\text{so } \underline{se \equiv 1 \pmod{p-1}}$$

$$\begin{aligned} m^{se} &= m^{1-t(p-1)} \\ &= m \cdot (m^{p-1})^{-t} \\ &= m \cdot 1^{-t} \\ &= \underline{m} \end{aligned}$$

Example Plain text HELLO

$$A = \{A=0, \dots, Z=25, !=26, . = 27, ? = 28\}$$

$$\underline{p=29 \text{ \& char } e=3.}$$

$$\underline{d=19} \quad (\text{so that } \underline{ed \equiv 1 \pmod{28}})$$

The private key is known to the receiver only.

The public key is in public domain.

After encryption the message is sent & any one can see it.

The receiver can decrypt the message using the private key.

So these keys are generated by the receiver & the public key is announced.