

ElGamal crypto system:

$G = \langle g \rangle$ Alice & Bob picks a & ~~b~~
& g^a & ~~g^b~~ are made public.

Bob wants to send message M .
He picks k & computes g^k & $(g^a)^k = g^{ak}$ & Mg^{ak}
 (g^k, Mg^{ak}) to Alice.

Alice computes $(g^k)^a = g^{ak}$
& computes $(Mg^{ak})g^{-ak} = M$.

Example: $G = \mathbb{Z}_{29}^*$ units of \mathbb{Z}_{29}
 $= \langle 2 \rangle$

$\mathbb{F}_q \setminus \{0\} = \mathbb{F}_q^*$
is a cyclic group.

2 is public

$M=12$

$a=3$ & public: $2^3 = 8$

~~$b=5$ & public: $2^5 = 32 = 3$~~

$k=5$
 $2^k = 32 = 3$
 $(2^k, M2^{ak}) = (3, 12 \times 2^{3 \times 5}) = (3, 12 \times 8^5)$
 $= (3, 5)$ sent to Alice

Alice computes $(2^k)^a = (2^5)^3 = 3^3 = 27$

$27 \times 14 = 1 \pmod{29}$
 \parallel
 $(-2) \times 14 = -28$

$(M2^{ak})2^{-ak} = 5 \times 27^{-1} =$
 $= 5 \times 14$

$$= 70$$

$$= 12$$

For a gp G if $a, b \in G$

Find the smallest non-negative integer x
s.t. $a^x = b$

This eqn may not have a soln.

If there is a soln, then the smallest non-negative integer x is denoted by $\log_a b$ & it is called the discrete logarithm of b wrt a .

E.g.: $\mathbb{Z}_5 = \{0, 1, 2, 3, 4\}$

$$\mathbb{Z}_5^\times = \{1, 2, 3, 4\} = \{2, 2^2=4, 2^3=3, 2^4=1\} = \langle 2 \rangle$$

$b \in \mathbb{Z}_5^\times$ $2^x = b$ has solns.

& $2^x = 4$ then $x = 2, 6, \dots$

$$\log_2 4 = 2 \text{ in } \mathbb{Z}_5^\times$$

$$3, 3^2=4, 3^3=2, 1$$

$$3^y = 4$$

$$\log_3 4 = 2$$

$$\log_3 1 = 0$$

$$\log_4 2 \text{ in } \mathbb{Z}_5^\times ?$$

$$2 \notin \langle 4 \rangle = \{4, 1\}$$

$$4^x = 2 \text{ has no soln.}$$

$$\log_4 2 \text{ does not exist}$$

$$G = S_3 = \{ (1), (1\ 2), (1\ 3), (2\ 3), (1\ 2\ 3), (1\ 3\ 2) \}$$

$$(1\ 2)^2 = 1 = (1\ 3)^2 = (2\ 3)^2$$

$$(1 \ 2 \ 3)^2 = (1 \ 3 \ 2)$$

$$(1 \ 2 \ 3)^3 = 1$$

$$(1 \ 3 \ 2)^2 = (1 \ 2 \ 3)$$

$$(1 \ 3 \ 2)^3 = 1$$

$$\log_{(1 \ 2 \ 3)} (1 \ 3 \ 2) = 2$$

$$\log_{(1 \ 2)} (1 \ 3 \ 2) \text{ does not exist}$$

Observe in a cyclic group G with respect to g a generator of the gp, discrete log exists for every element. (Since any $b \in \langle g \rangle$, $\exists x$

$$\text{s.t. } \underline{\underline{g^x = b}}$$

To find x , we can find $g, g^2, g^3, \dots, g^{n-1}$ & compare b with each of them. ($n = \text{order of } g$)

Can we find $\log_g b$ with lesser operations

Shank's Baby-step Giant-step algorithm

$$G = \langle g \rangle \quad \text{order } g \text{ is } n.$$

Have to find $\log_g a$ (i.e. finding x s.t. $g^x = a$)

$$m = \lceil \sqrt{n} \rceil \quad \text{least integer larger than or equal to } \sqrt{n}.$$

$$x = mq + r \quad \text{where } 0 \leq r \leq m-1$$

$$0 \leq q \leq m-1$$

$$a = g^x = g^{mq+r}$$

$$\boxed{g^{mq} = a g^{-r}}$$

$$\begin{aligned} &\text{as } m(m-1) + m - 1 \\ &= m^2 - m + m - 1 \\ &= m^2 - 1 \approx n \end{aligned}$$

Baby-step Compute

$$B = \{ (ag^{-r}, r) \mid 0 \leq r \leq m-1 \} \subset G \times \{0, 1, \dots, m-1\}$$

Suppose $\boxed{ag^{-s} = 1}$ for $r=s$

Then $a = g^s$ so that $\log_g a = s$.

Suppose $ag^{-r} \neq 1$ for $r \in \{0, 1, \dots, m-1\}$
 Then go to the giant step.

Compute $g^m, g^{2m} = (g^m)^2, g^{3m}, \dots$

Find q (smallest) s.t. $g^{mq} = ag^{-s}$ for some $s \in \{0, 1, \dots, m-1\}$

Then $a = g^{mq+s}$
 so that $\log_g a = \underline{mq+s}$.

$2\sqrt{n} + 2\sqrt{n} \times \sqrt{n} = \underline{2\sqrt{n} + n}?$

$n=61 \quad G = \mathbb{Z}_{61}^\times = \{1, \dots, 60\} \pmod{61}$

$= \langle 2 \rangle = \{2, 4, 8, 16, 32, 3, 6, 12, 24, 48, 35\}$

$g=2, g^7=31$

$m = \lceil \sqrt{61} \rceil = 8$

$a_1 = 6$

$a_2 = 7$

$\log_2 7$

$\log_2 6$

$ag^{-r} \mapsto \begin{pmatrix} 6 & 0 \\ 0 & 1 \end{pmatrix} (6 \times 31, 1), (6 \times 31^2, 2)$

$$\begin{array}{c} \parallel \\ 3 \times (2 \times 31) \\ (3, 1) \end{array}$$

$$(3 \times 31, 2) = (82, 2),$$

$$\begin{array}{c} \parallel \\ (1+2) 31 \\ 31+1 = 32 \end{array}$$

$$(32 \times 31, 3) = (16, 3), \quad (16 \times 31, 4) = (8, 4), \quad (8 \times 31, 5) = (4, 5),$$

$$\parallel \\ 16 \times 2 \times 31$$

$$(4 \times 31, 6) = (2, 6), \quad (2 \times 31, 7) = (1, 7)$$

$$a g^{-7} = 1 \text{ so } a = g^7 \quad \therefore \quad \log_g a = 7$$

$$\boxed{\log_2 6 = 7}$$

$$\text{Next, } \log_2 7 = ?$$

$$(7, 0), \quad (7 \times 31, 1) = ((1+4) \times 31, 1) = (31+2, 1) = (33, 1)$$

$$(33 \times 31, 2) = (1 + \underbrace{2 \times 16} \times 31, 2) = (31+16, 2) = (57, 2),$$