

RECAP: ✓ ADFGX & ADFGVX cipher.  
 ✓ Permutation Cipher.

• Use of password  
Playfair cipher Use password to prepare the table  
 PASSWORD is the password  
 alphabetically put them: ADOPRSW

Standard

A	B	C	D	E
F	G	H	I/J	K
L	M	N	O	P
Q	R	S	T	U
V	W	X	Y	Z

password

A	D	O	P	R
S	W	B	C	E
F	G	H	I/J	K
L	M	N	Q	T
U	V	X	Y	Z

Encryption:  $\Sigma: P \rightarrow C$

- Pair consecutive characters starting from the 1st character.

MOVE FORWARD :  
 (1) If  $X_1, X_2$  are diagonally opposite on the above table  
 then find  $Y_1, Y_2$ , diagonally opposite for the same rectangle  
 $Y_1$  &  $X_1$  are on the same column  
 &  $Y_2$  &  $X_2$  " " " column.

(2) If  $X_1, X_2$  are on the same row,  
 apply right cyclic 1-shift & get  
 $Y_1, Y_2$

• If  $X_1, X_2$  are on the same  
 column then downward cyclic  
 1-shift is to be applied &  
 get  $Y_1, Y_2$

• If  $X_1 = X_2$ , then insert X between  $X_1$  &  $X_2$  & do  
 pairing as  $X_1X$  &  $X_2X$ .....  
 & then apply the other rules for encryption.

MO → DN  
 VE → WZ  
 FO → AH  
 RW → ED  
 AR → DA  
 D ∴ → VO



## Decryption rules

$y_1 y_2 \dots$

(1')  $= (1)$

(2')  $\odot$  apply left cyclic 1-shift if  $y_1$  &  $y_2$  are on the same row on the table.

(3') apply upward cyclic 1-shift if  $y_1$  &  $y_2$  are on the same column on the table.

## Table

A	D	P	R	S
W	B	C	E	F
G	H	I/J	K	L
M	N	O	Q	T
U	V	X	Y	Z

Decrypt (Playfair cipher)

RD ZP TP DM QA ZP

$$\text{HCF} = \text{GCD}$$

Coprime or relatively prime

$a, b$

$$\text{gcd}(a, b) = 1$$

also denoted by  $(a, b)$

$$d = \text{gcd}(a, b)$$

$$\underline{d|a, d|b}$$

if  $c|a, c|b$ , then  $c|d$ .

Lemma: if  $a = bq + r$  then

$$\text{gcd}(a, b) = \text{gcd}(b, r).$$

(write a proof)

Thm: if  $a, b \in \mathbb{Z}$ , at least one of them is nonzero then  $\text{gcd}(a, b)$  exists &  $\text{gcd}(a, b) = xa + yb$  for some  $x, y \in \mathbb{Z}$ . In fact  $\text{gcd}(a, b)$  is the least positive integer which can be written in the form  $sa + tb$  for  $s, t \in \mathbb{Z}$ .  
 $\rightarrow$  w.l.o.g.,  $b \geq 0$ .



Proof: <sup>consider</sup>  $S = \{sa + tb : sa + tb > 0, s, t \in \mathbb{Z}\} \subset \mathbb{N}$

Well ordering principle  $S$  has a least member

$S \neq \emptyset$  because take  $t=1, s=0$   
 $b \in S$ .

iff  $\gcd(a, -b)$  exists  
 then  $\gcd(a, b)$  exists  
 &  $\gcd(a, b) = \gcd(a, -b)$

let  $d = sa + tb$  be the least member of  $S$ .

Then show that  $d = \gcd(a, b)$ .

Step 1:  $d|a$

Step 2:  $d|b$

Step 3:  $c|a, c|b$   
 $\Rightarrow c|d$ .

if  $d \nmid a$ , then  $a = dq + r \Rightarrow r = a - dq > 0$

$$r = a - (sa + tb)q$$

$$= (1 - sq)a + (-tq)b.$$

since  $r < d$

$r \in S$  a contradiction

Euclidean algorithm

$\alpha, \beta \in \mathbb{Z}, \beta > 0$

$$\alpha = \beta q + r$$

where  $q, r \in \mathbb{Z}$

&  $0 \leq r < \beta$ .

so  $r=0$

so that  $d|a$ .

$\implies d|b$ .

so  $d$  is the  $\gcd(a, b)$

Algorithm. Assume  $b > 0$ . Apply Euclidean algorithm

$$a = bq_1 + r_1$$

if  $r_1 = 0$   $\gcd(a, b) = b$  else.

$$b = r_1q_2 + r_2$$

$$r_1 = r_2q_3 + r_3$$

$\vdots$

$$r_{k-1} = r_kq_{k+1} + r_{k+1}$$

$$r_{k-2} = r_{k-1}q_k + r_k$$



$$r_k = r_{k+1}q_{k+2} + \underbrace{r_{k+2}}_{=0} \quad \& \quad r_{k+2} = 0,$$

For some  $k$ ,  $r_{k+2} = 0 \quad \leadsto \quad \underline{r_1 > r_2 > \dots > r_k}$

Then  $\underline{\gcd(a, b) = r_{k+1}}$

$$\gcd(a, b) = \gcd(b, r)$$

$$\boxed{a = bq + r.}$$

$$r_{k+1} = \gcd(r_k, r_{k+1})$$

$$= \gcd(r_{k-1}, r_k)$$

$$= \gcd(r_{k-2}, r_{k-1})$$

$$\vdots$$

$$= \gcd(a, b)$$

$$\text{Find } \gcd(3026, 445) =$$

$$3026 = 445 \times 6 + 356$$

$$445 = 356 \times 1 + 89$$

$$356 = 89 \times 4 + 0$$

$$\underline{\gcd(3026, 445) = 89}$$

Suppose  $\underline{n > 1}$  fixed.  $m \in \mathbb{N}$ .

Write  $m = \underline{a_k n^k + a_{k-1} n^{k-1} + \dots + a_1 n + a_0} =: p_k(n)$

where  $a_k \neq 0, 0 \leq a_0, a_1, \dots, a_{k-1} \leq n-1$

for some  $k$ .

$$\underline{n^k \leq m \leq \underbrace{n^{k+1} - 1}_{\leq n^{k+1}}}$$

$$p_k(n) \in \mathbb{Z}.$$

$$\deg(p_k(n)) = k.$$

$$\log_n(n^k) \leq \log_n m \Rightarrow k \leq \log_n m \leq k+1$$

$$k = \lfloor \log_n m \rfloor + 1.$$