

Notes on Perfect Squares & k-nary XOR

1. Definition of Perfect Square

A number is a perfect square if it can be written as k^2 for some integer k .

Examples: 0, 1, 4, 9, 16, 25, ...

2. Prime Factorization Property

Every integer can be expressed uniquely as a product of prime factors: $n = p_1^{a_1} \times p_2^{a_2} \times \dots \times p_k^{a_k}$. n is a perfect square if and only if all exponents $a[i]$ are even.

Examples:

- $36 = 2^2 \times 3^2 \rightarrow$ perfect square.
- $18 = 2^1 \times 3^2 \rightarrow$ not a square. (exponent of 2 is odd)

3. Product Check

If we multiply several numbers together, the product is a perfect square iff the combined exponent count of every prime factor is even. In competitive programming, this is often checked using XOR of exponents mod 2 (parity check).

4. Special Cases

- 0 is a perfect square ($0^2 = 0$).
- Negative numbers are not perfect squares in integers.

5. k-nary XOR (Generalized XOR)

Normally XOR is defined in base 2 (binary). But we can generalize XOR into base k . For k -nary XOR, addition is done modulo k without carry. This is useful when tracking exponents modulo k instead of just modulo 2. For example, to check for perfect cubes, we would track exponents modulo 3.

Examples:

- Binary XOR ($k=2$): $1 \oplus 1 = 0$, $1 \oplus 0 = 1$.
- Ternary XOR ($k=3$): $2 \oplus 2 = 1$ (since $(2+2) \bmod 3 = 1$), $2 + 2 + 2 = ((2+2) \bmod 3 + 2) \bmod 3 = 0$ (Freq of 2 is 3 times.)
- General rule: $(a \oplus b) = (a + b) \bmod k$.

6. XOR Hash Trick for Perfect Squares and k-nary Checks

1. Perfect Square Check

- A number is a **perfect square** if every prime factor appears an **even number of times**.
- Example:
 - $36 = 2^2 * 3^2 \rightarrow$ perfect square
 - $18 = 2 * 3^2 \rightarrow$ not a perfect square (2 appears odd times)

2. XOR Hash Trick

- Assign a **random hash value** to each prime number.
- For an array $A[1..n]$, compute **prefix XOR** of the hashes of prime factors modulo 2:
 $\text{prefix}[i] = \text{hash_xor}(A[1..i])$
- To check if the product of a subarray $A[l..r]$ is a perfect square:
 $\text{subarray_xor} = \text{prefix}[r] \wedge \text{prefix}[l-1]$
 $\text{if}(\text{subarray_xor} == 0) \Rightarrow \text{perfect square}$
- XOR cancels even occurrences of primes, leaving 0 if all are even.

3. Handling Zero and Negative Numbers

- **Zero**: Any subarray containing 0 is a perfect square.
- **Negative numbers**: Track the count of negative signs. Product is square if count of negatives is even.

4. k-nary Generalization

- For checking if **each prime occurs a multiple of k times**, define **k-nary XOR**:
 $a \oplus_k b = (a + b) \% k$
- Compute **prefix k-nary sum** for the array.
- Subarray check:
 $\text{subarray_k_xor} = (\text{prefix}[r] - \text{prefix}[l-1] + k) \% k$
 $\text{if}(\text{subarray_k_xor} == 0) \Rightarrow \text{all primes occur multiple of } k \text{ times}$
- Example:
 - Binary XOR ($k=2$) → checks even occurrences.
 - Mod 3 addition ($k=3$) → checks multiple-of-3 occurrences.

5. Summary Table

Check Type	Operation	Condition for Subarray Product
Perfect square	XOR (mod 2)	$\text{subarray_xor} == 0$
Multiple-of-3	Add mod 3	$\text{subarray_mod3} == 0$
General k	Add mod k	$\text{subarray_modk} == 0$

This technique allows fast checking of multiplicities of prime factors in subarrays without computing large products.