



Bonus Class - Bit Manipulation

Special class

Bitwise Operators:-

int a = 0
 ~a → 1
 -1 → two's complement



a	b	a/b
0	0	0
0	1	0
1	0	0
1	1	1

a	b	OR a/p
0	0	0
0	1	1
1	0	1
1	1	1

a	b	a ^ b
0	0	0
0	1	1
1	0	1
1	1	0

10 → 2

Decimal \rightarrow Binary $\underline{\underline{\hspace{1cm}}}$
Binary \rightarrow Decimal $\underline{\underline{\hspace{1cm}}}$ already

sign

1's comp

2's comp

int a = 0

~a = (-1)

a = 000 - - - 0000

~a = 1 1 1 1 - - - 1 1 1 1

-ve

1's → 000 - - - 00
2's →

0000 - - - 01

→ (-1)

→ Bitwise op → $&, |, ^, \sim, <<, >>$

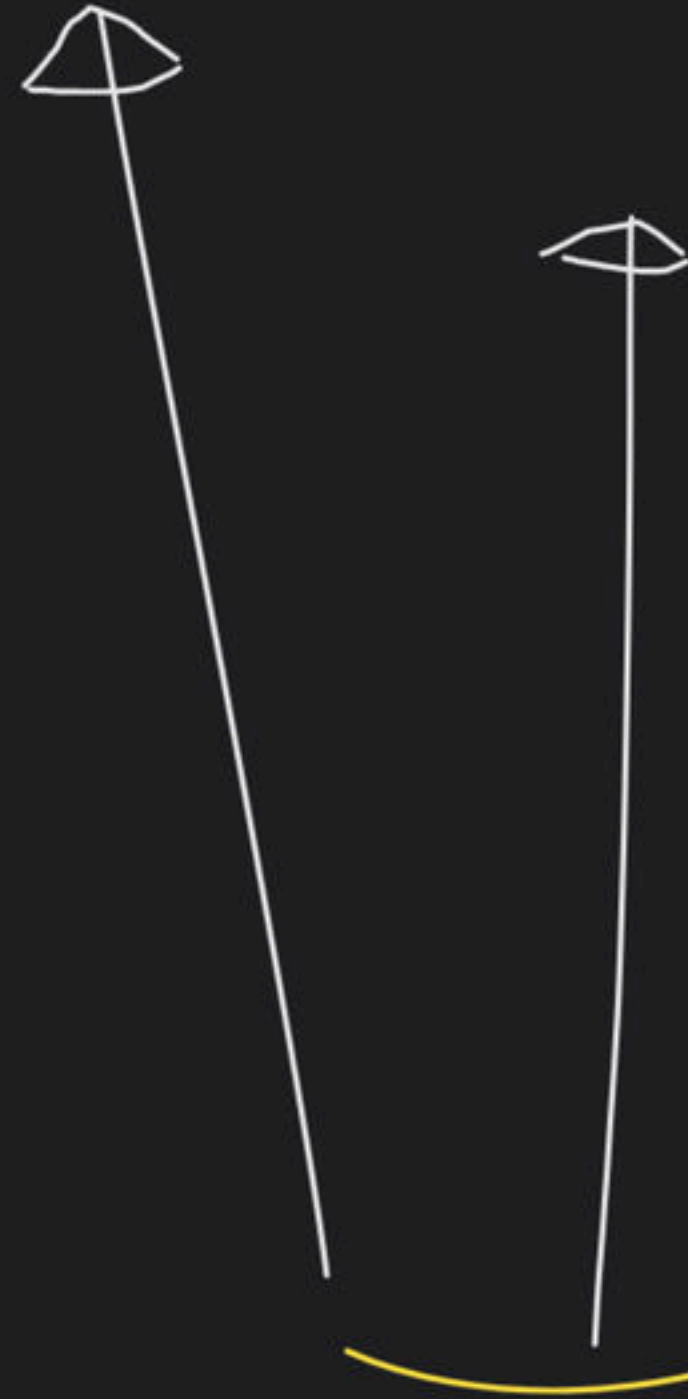
→ sign bit

→ 1's comp

→ 2's comp

→ D → B

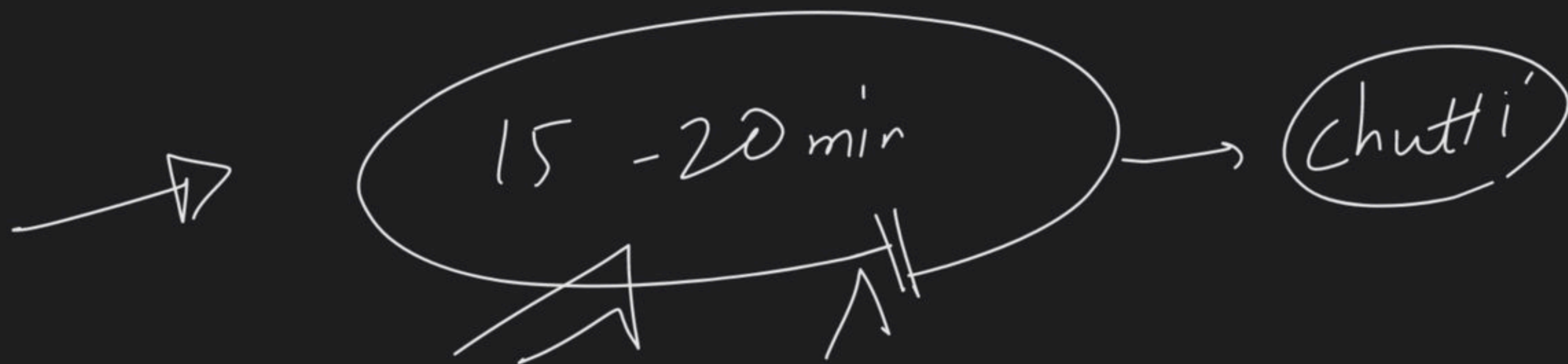
→ B → D

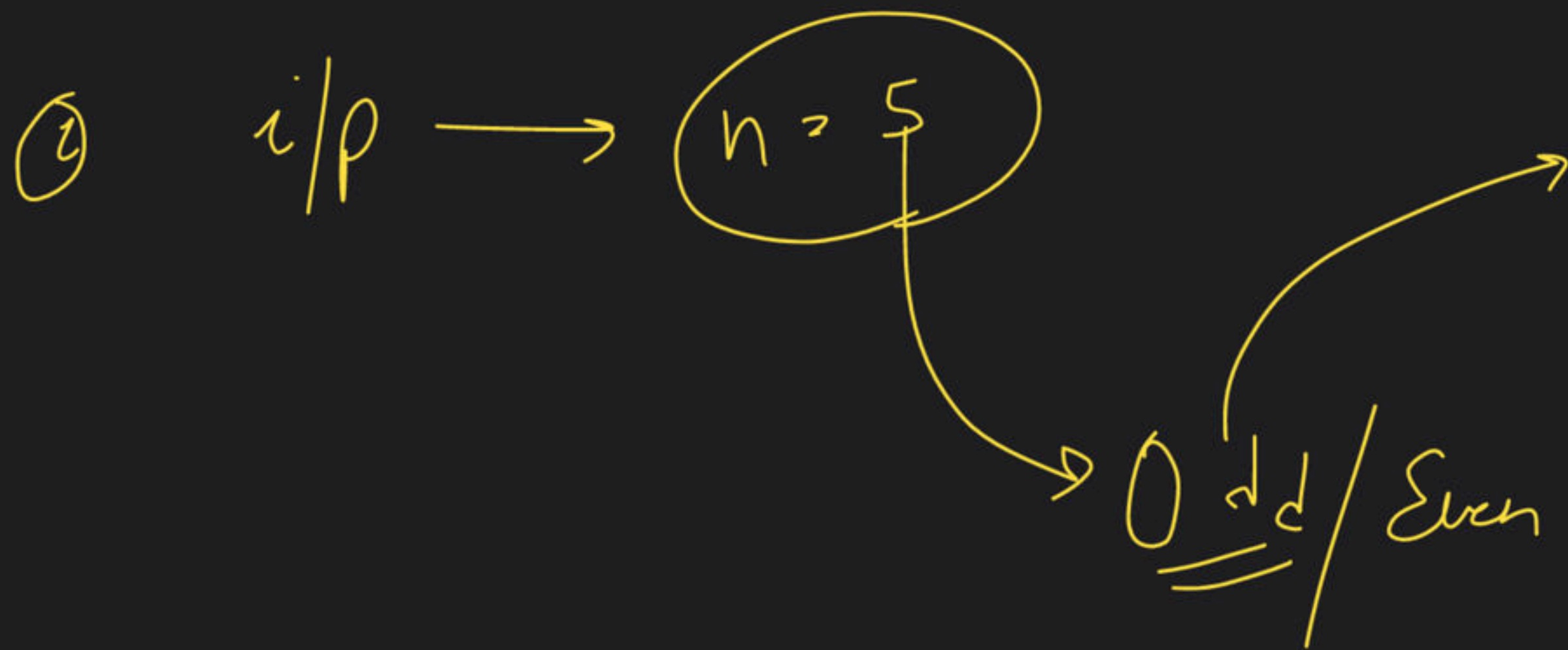


Rewrite

→ This is not the

Last Class



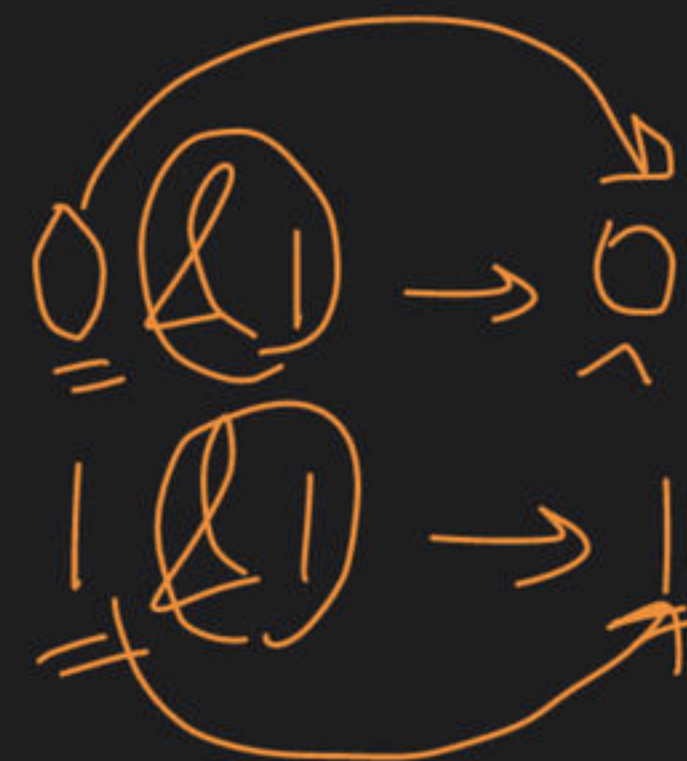


Binary

3 \rightarrow 11^{LSB}

5 \rightarrow 101

7 \rightarrow 111



n

21

0

\rightarrow even

0/0

$\div 2 = 0 \rightarrow \text{Even}$
 $\div 2 = 1 \rightarrow \text{odd}$
 \uparrow

11 $\rightarrow 0$

→ Mask

int mask = 1 << i
int ans = n & mask

get i^{th} bit

→ n = 10

i = 3
3rd

00000 --- 000010

& 00000 --- 00001 → mask

00000 --- 00001

LBB

1 & 1
0

mask → 1

0000 --- 1

000001000

<< 3

& mask → 0 → bit 0

non zero → bit 1

change i^{th} bit

1

Set i^{th} bit

0

clear i^{th} bit

Set i^{th} bit

2nd bit 1 < n < K_{rd}

n = 1010 , ~~n~~ $i = 2$

1st 2nd 3rd 4th

0000 ——— 0000 | 0 1 0

mask 0000 00 . 0000 0100 OR

0000 000 ——— 0000 1110

0
1
—
1

↪ create mask¹
mask = 1 << 2
1 << i

0000 1 << 2
0000 100

ans = n | mask

clear i^{th} bit

$\hookrightarrow n=10, \quad i=1$

0000 — — 0000 1010

11th
1 0th

mask
2

1111 — — 1111 101

0000 0000 1000

(8)

data lagana

21

set

or 1

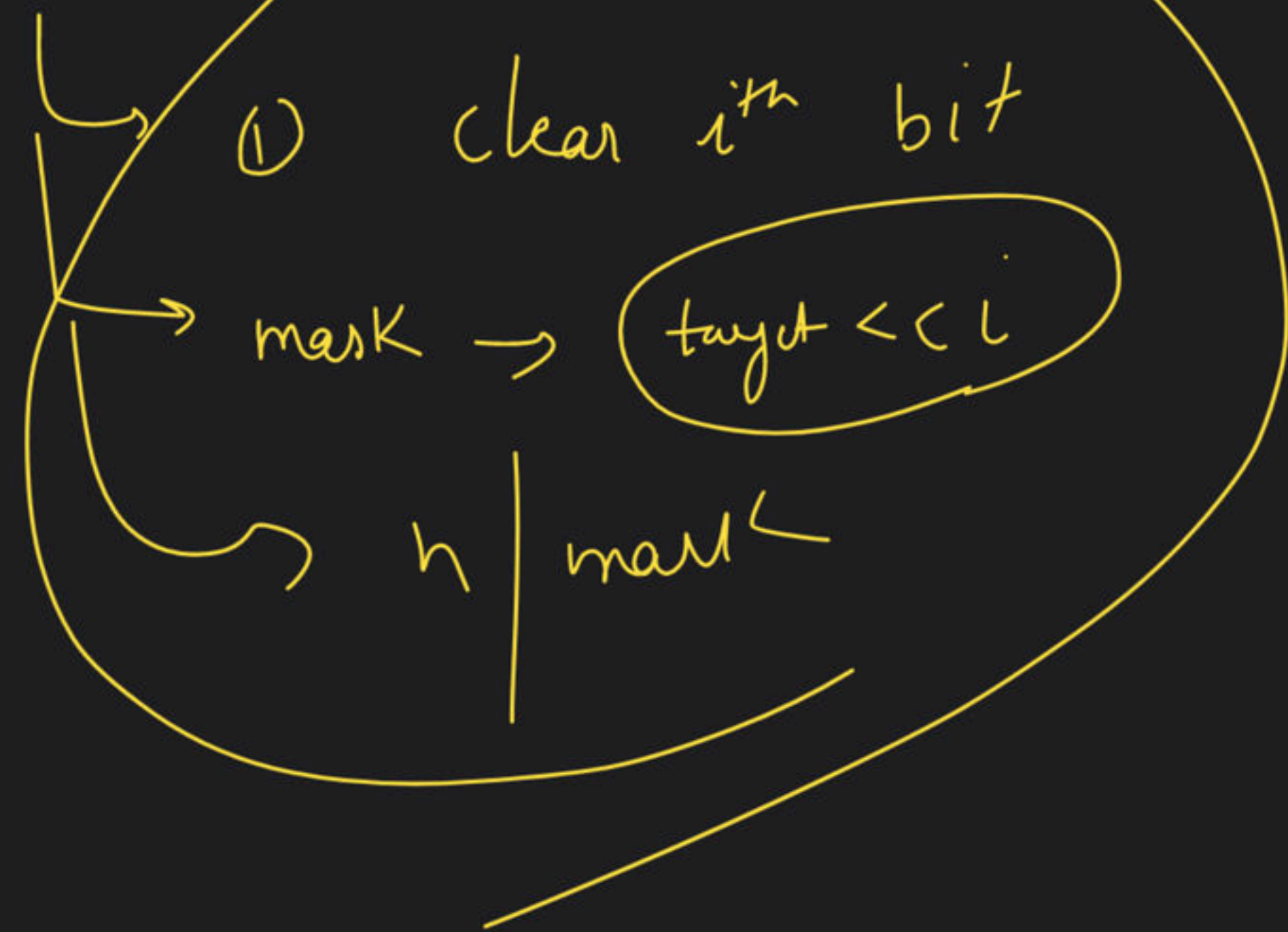
(0000 — — 0010)

$i < 1$

$(1 < i)$

11
—
0

Update i^{th} bit



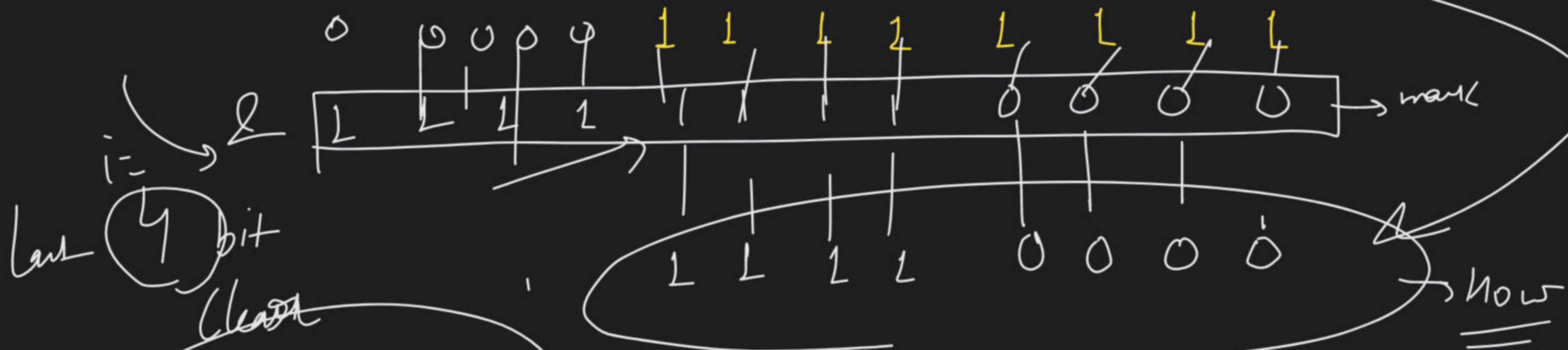
$$\textcircled{4} \sim \cancel{F} \wedge \cancel{F} \sim \cancel{F} \wedge \cancel{F}$$

$$= \textcircled{4}$$



~~$n = 13$~~

$11111110000 \leftarrow (-1 < i)$



- ① create mark = $(-1 < i)$
- ② $n = n \& \text{mark}$

--- 1 1 1 1 1 1 0 0 0 0

1 1 1 1 1 1 1 1 $\leftarrow 7$

-1 0

00000 - . 000 111

0000 - - 000 100

(4)

→ Check power of 2

$$\left(\frac{16}{2} \rightarrow \frac{8}{2} = 4 \right. \\ \left. \frac{2}{2} \right)$$

i/p → $n = 16$

2, 4, 8, 16, 32, 64

$2^4 \rightarrow \sqrt[4]{n}$

128, 256

2572, 1024

$$\begin{array}{lcl}
 8 & \rightarrow & 1000 \\
 7 & \rightarrow & \underline{0111} \& \\
 & & \underline{0000}
 \end{array}$$

$$\begin{array}{lcl}
 2^1 \rightarrow 2 & \rightarrow & \underline{10} \\
 2^2 \rightarrow 4 & \rightarrow & \underline{100} \\
 2^3 \rightarrow 8 & \rightarrow & \underline{1000} \\
 2^4 \rightarrow 16 & \rightarrow & \underline{10000}
 \end{array}$$

$$\begin{array}{lcl}
 16 & \rightarrow & 10000 \\
 15 & \rightarrow & \underline{01111} \& \\
 & & \underline{0000}
 \end{array}$$

$$\begin{array}{lcl}
 n & \rightarrow & 4 \rightarrow 100 \\
 n-1 & \rightarrow & 3 \rightarrow \underline{011} \& \\
 & & 000
 \end{array}$$

Count set bit

$$n \& (n-1) == 0$$

it is a power of 2

→ Count Set bits

8 → 0000 - - - 001000

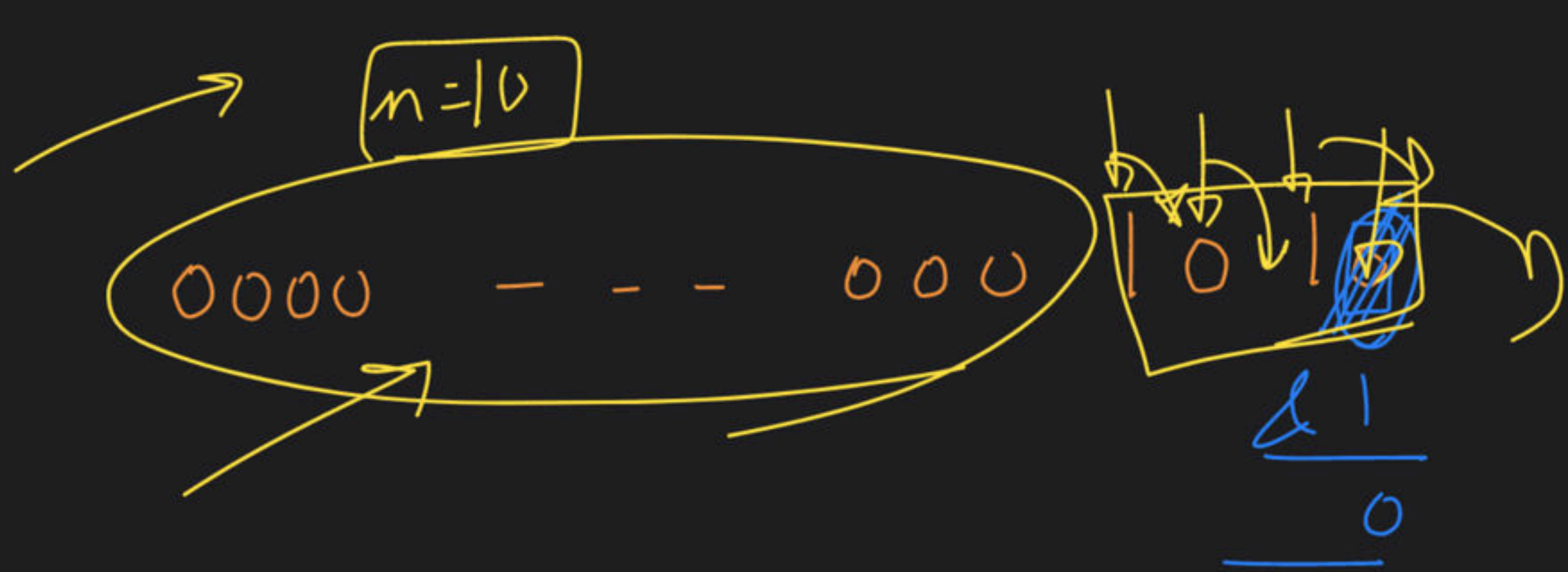
no. of set bit = 1

7 → 0000 - - - 00111

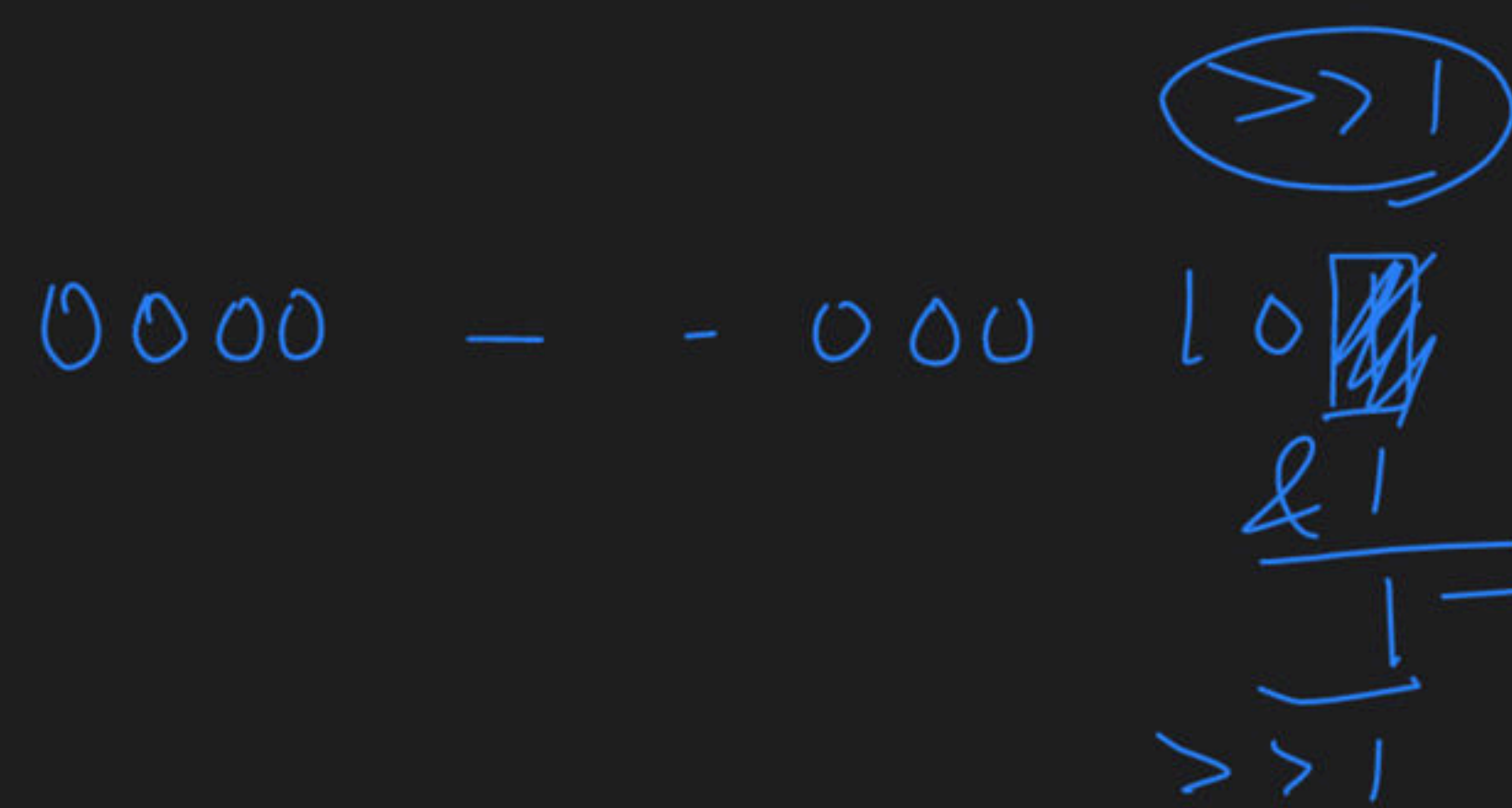
no. of set bit = 3

10 → 00000 - - - 1010

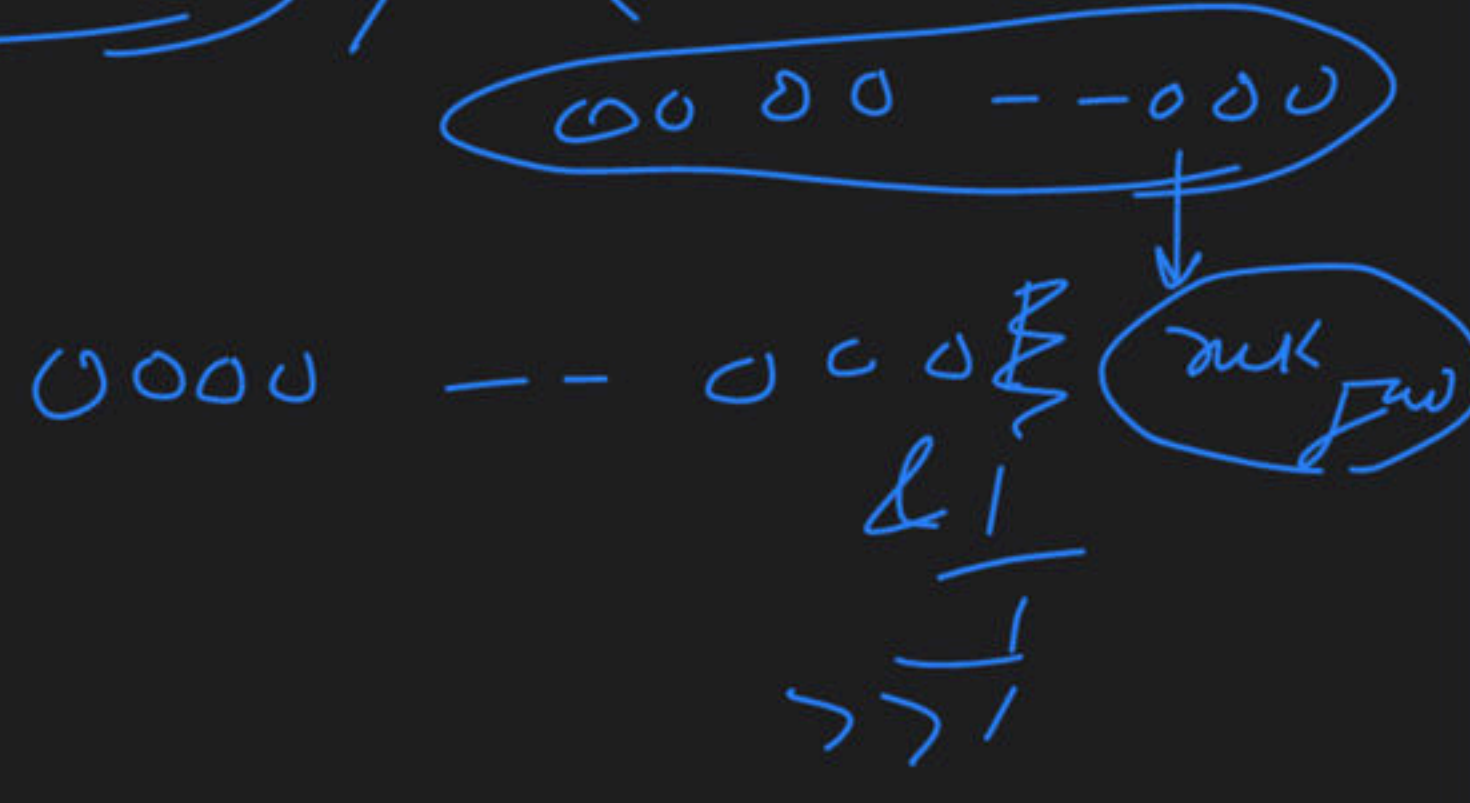
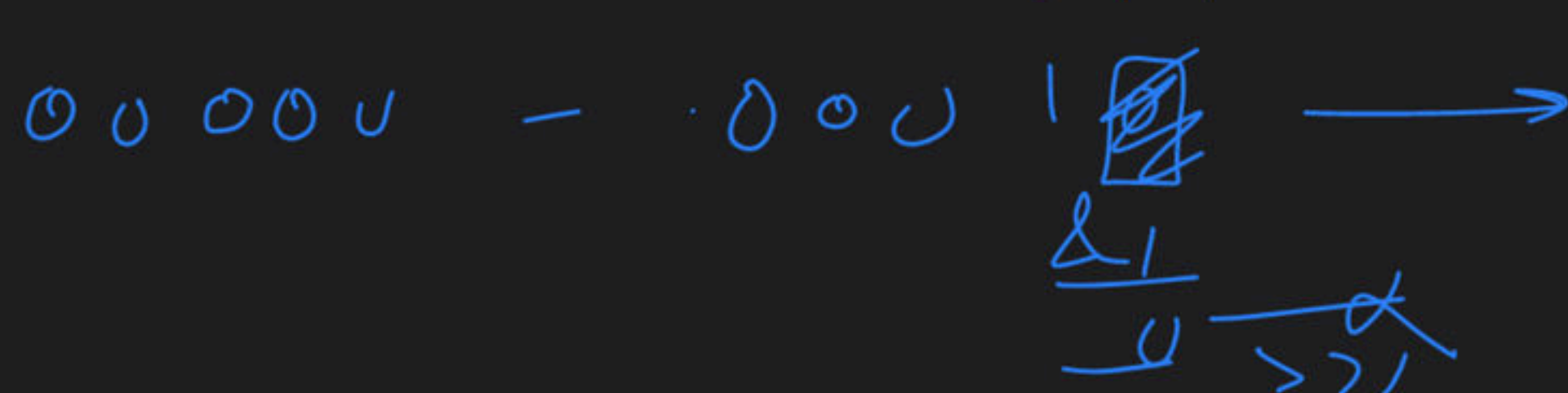
no. of set bit = 2



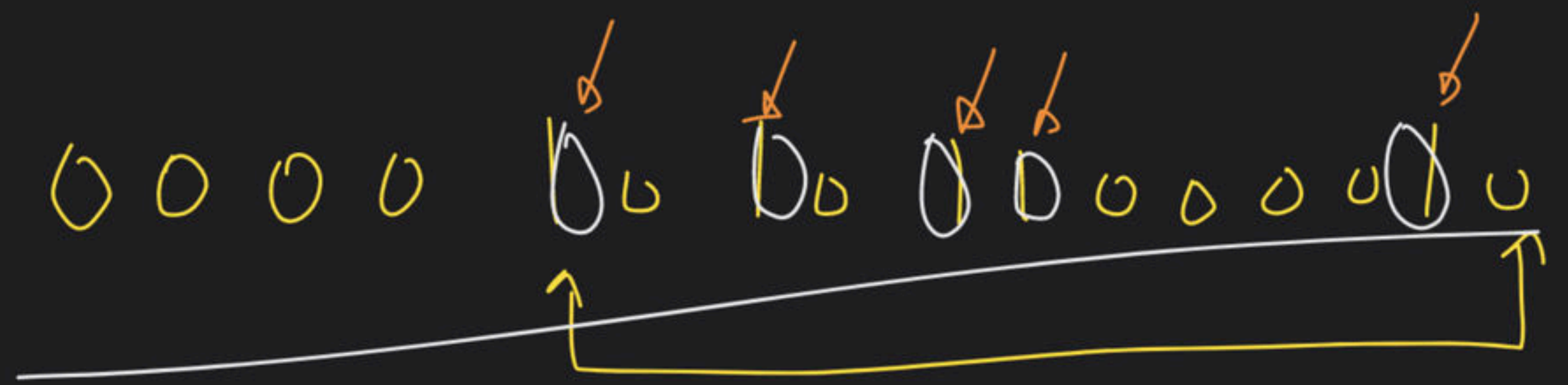
0



count inc / 2



fast



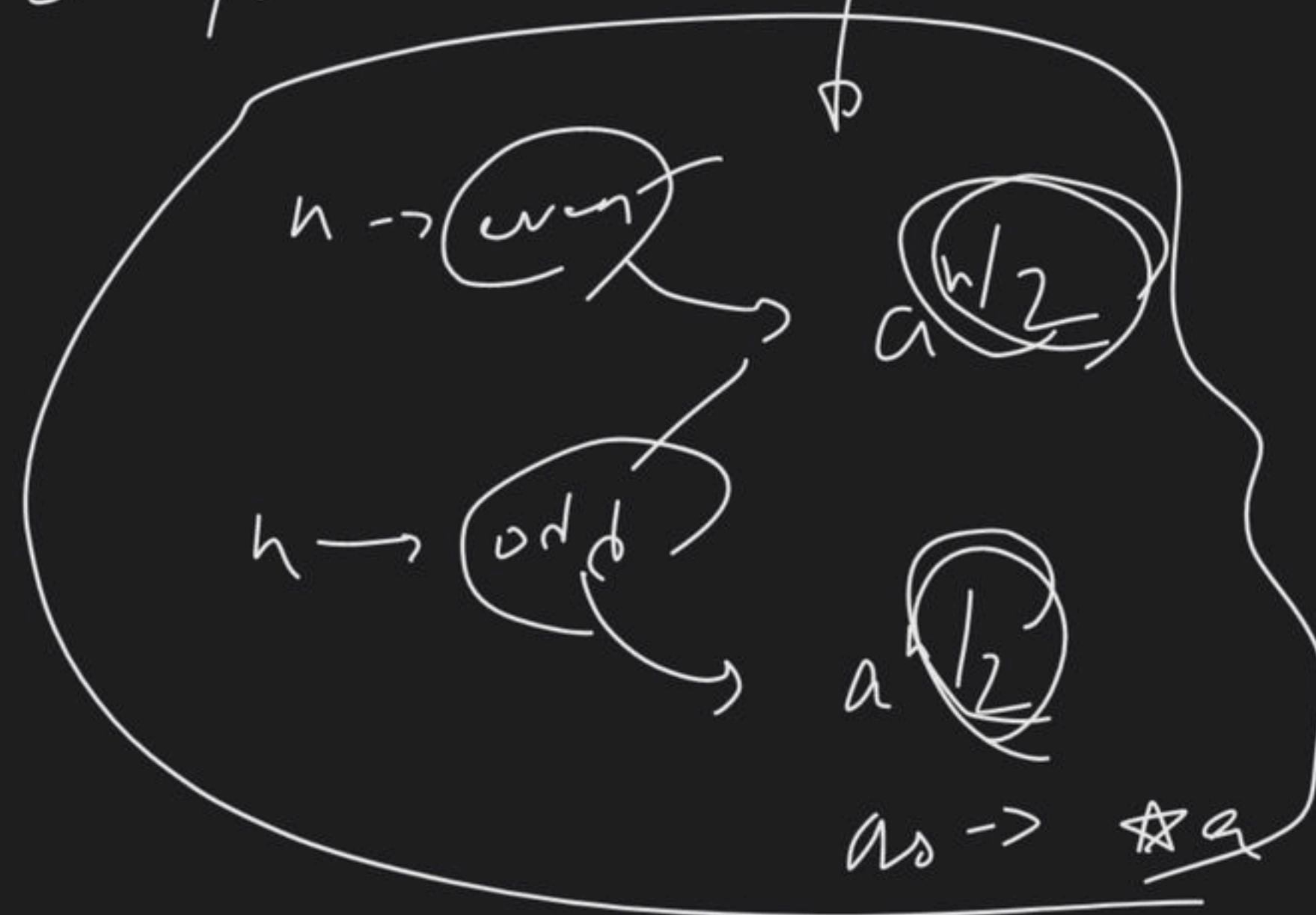


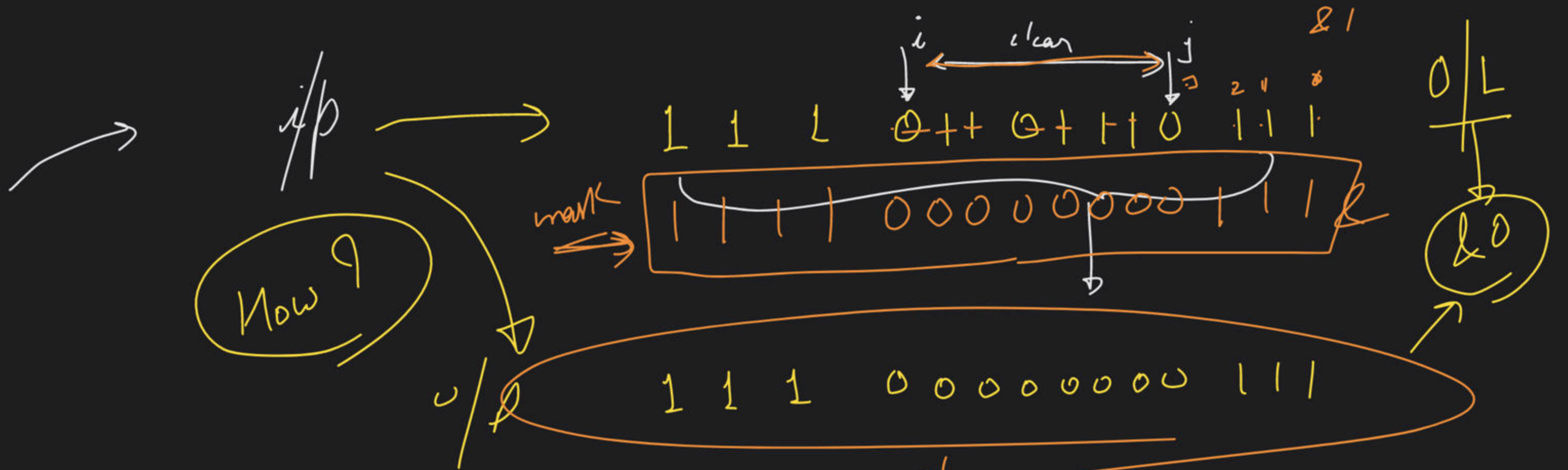
fast Exponentiation

n/w
 $\&1$
 $>>1$

a^n

$n \rightarrow \text{even/odd}$



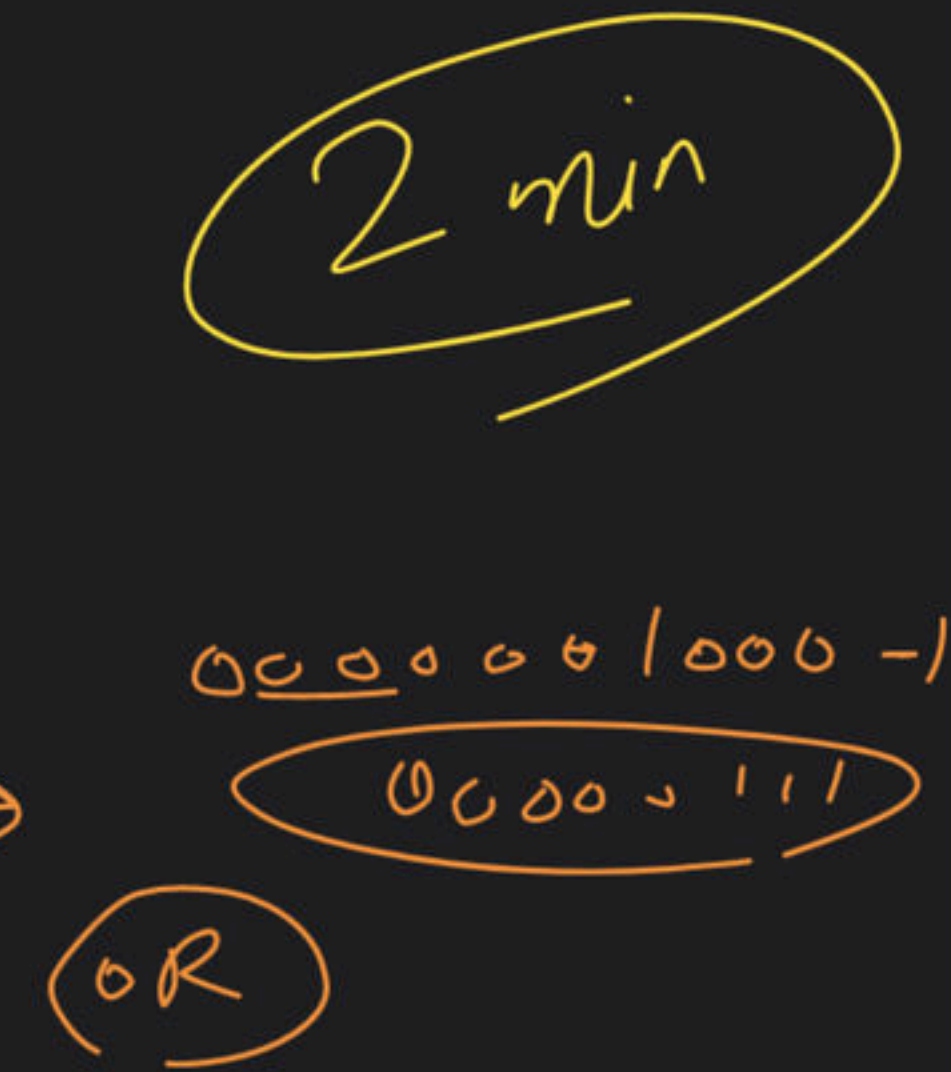


Equations:

$$a = (-1 << (i+1))$$

$$b = (L << j) - 1$$

a/b



15

$f(15, 2, 1)$

0000 — 0000

i	j
1	1
2	1
3	1

0000 — 000

1	0	0	1
---	---	---	---

↪ g



$a \rightarrow$

$a = (-1 \ll (i+1))$

$(1 \ll j) - 1$
 $b \nearrow$

00000000

$\begin{array}{c} | \\ \ll j \\ \ll 3 \end{array}$

00000000 1000

-1

\rightarrow 00000000 111

$2 \rightarrow 10 \rightarrow \overset{n-1}{1}$

$4 \rightarrow 100 \rightarrow 11$

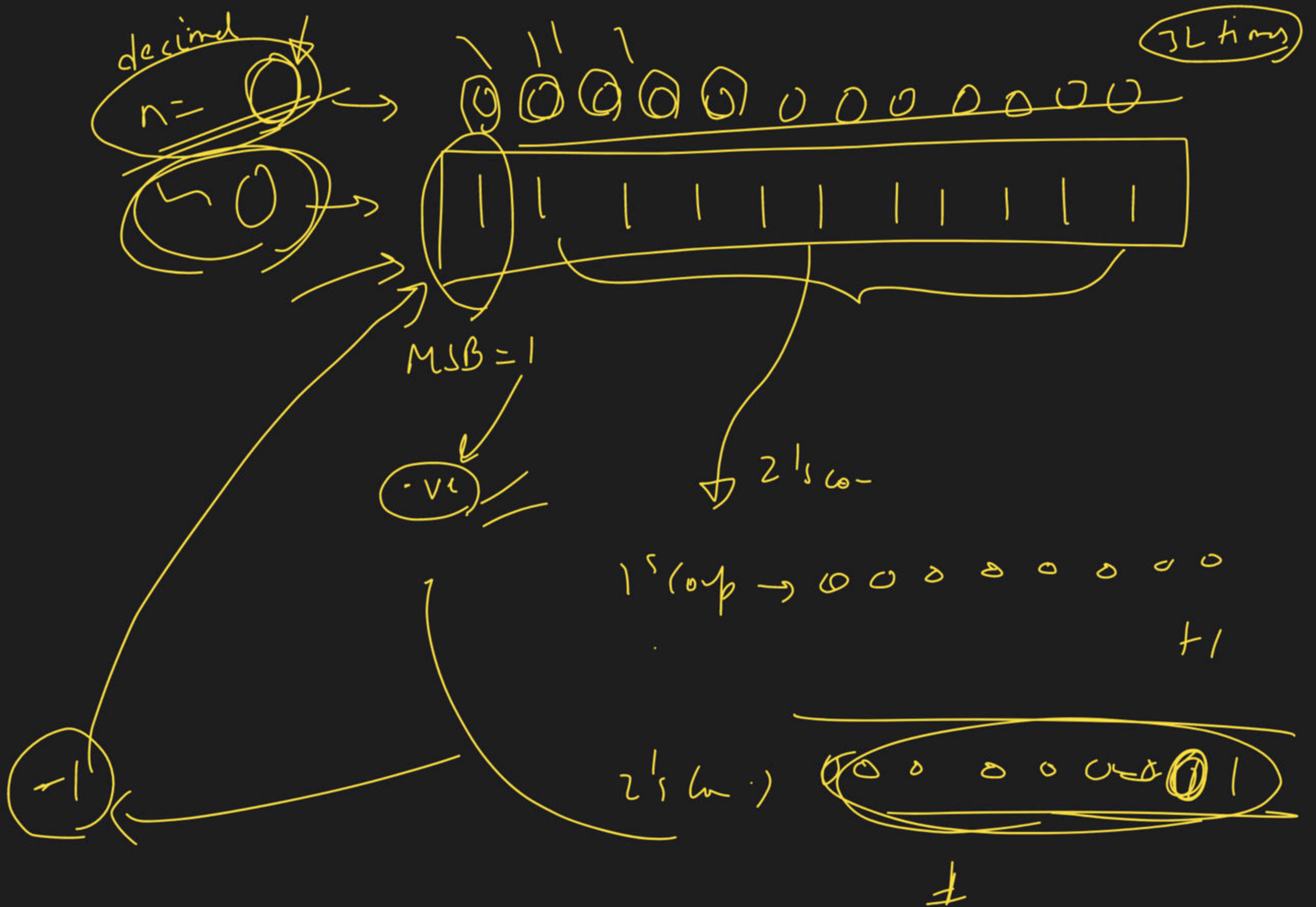
$8 \rightarrow 1000 \rightarrow 111$
 \wedge

$16 \rightarrow 10000 \rightarrow 1111$
 \nearrow

32

11111
 $3'$

$a/b \rightarrow \text{mark}$

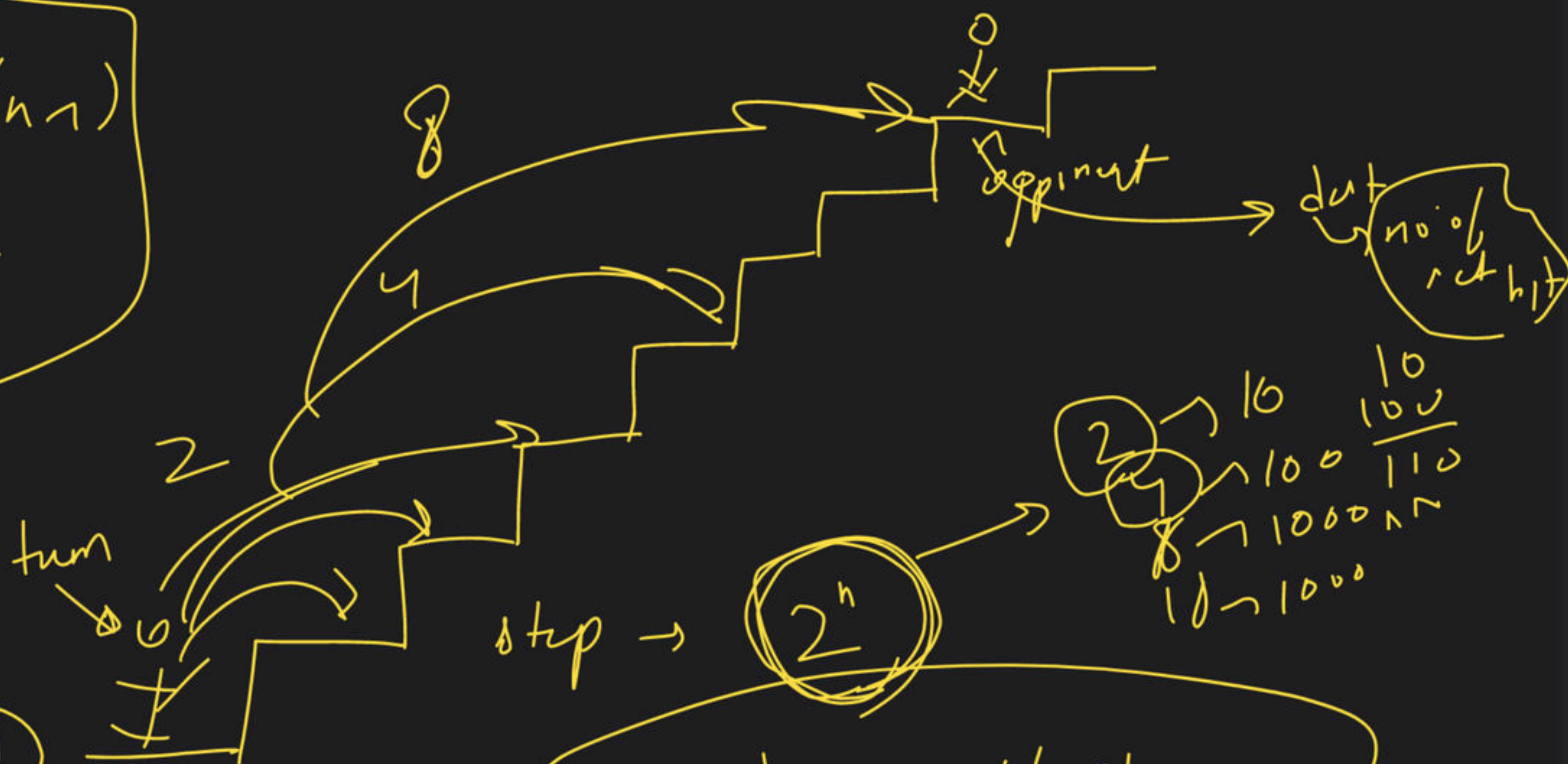




```
while (n != 0)
{
    n = n << 1;
    cnt++;
}
```

$2^n \rightarrow 1$ set bit
 $1 \rightarrow 1$ set bit

$n = n \< (n-1)$



step $\rightarrow 2^n$

min no. of steps to reach opponent

2 $\rightarrow 10$
4 $\rightarrow 100$
8 $\rightarrow 1000$
10 $\rightarrow 10000$

total set bit

last set bit

$2 + 4 + 8$
 $10 + 100 + 1000 \Rightarrow 1110$

→ Subsequence Using Bitmasking

$n \rightarrow \text{length of string}$

str = "abc"

Power set
or
Subsequence

" "
a
b
c
ab
bc
ac
abc

2^n

already
solved

→ Recursion

inc/exc pattern

1

a	b	c
✓	* ₀	x

→ a

a	b	c
x	✓	x

→ b

a	b	c
x	x	✓

→ c

a	b	c
✓	✓	x

→ ab

a	b	c
✓	x	✓

→ ac

a	b	c
x	✓	✓

→ bac

a	b	c
✓	✓	✓

→ abc

a	b	c
✓	✓	✓

→ " "

"abc"

length $\rightarrow n$

Power set \rightarrow

2^n

$$n=3$$

$$2^n = 2^3 = 8$$

$$0 \rightarrow 2^n - 1$$

$$0 \rightarrow 8 - 1$$

$$0 \rightarrow 7$$

Binary

32

0 \rightarrow 000

1 \rightarrow 001

2 \rightarrow 010

3 \rightarrow 011

4 \rightarrow 100

5 \rightarrow 101

6 \rightarrow 110

7 \rightarrow 111

abc \rightarrow ""

abc \rightarrow c

abc \rightarrow b

abc \rightarrow bc

abc \rightarrow a

abc \rightarrow ac

abc \rightarrow ab

abc \rightarrow abc

for (0 \rightarrow $2^n - 1$)

num

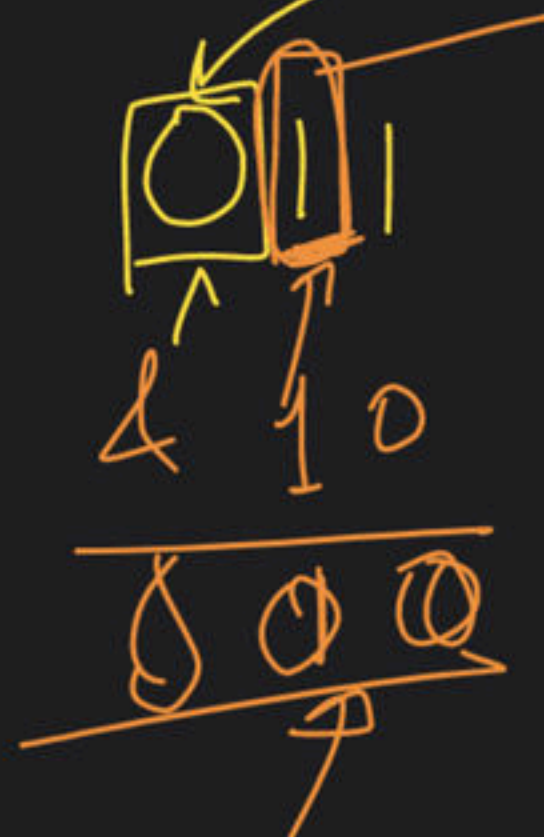
for (0 \rightarrow n)

{ //index

}

}

3 →
num



str → a b c

i = 0

i = 1

temp

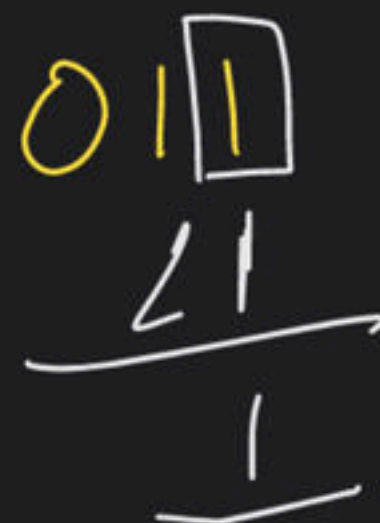
0 1 1

a b c

a b c

i = 2

3 →
num



max
1 < 0

i = 0

a b c

a

0 1 1
1 0 0
0 0 1

i = 1

1 < 1

0 1 0

1 < 2

1 0 0

