

EL9343 Homework 01

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Question 01

Prove the *Transpose Symmetry* property, i.e., $f(n) = O(g(n))$ if and only if $g(n) = \Omega(f(n))$.

Solution

To prove that $f(n) = O(g(n))$ if and only if $g(n) = \Omega(f(n))$

Let us consider the Left hand Side –

If $f(n) = O(g(n))$, then for that condition to satisfy, based on the definition of big-Oh notation,

There exists a constant c , and $\Rightarrow f(n) \leq c \cdot g(n)$

And for c to be positive ($c > 0$) $\Rightarrow g(n) \geq \frac{1}{c} \cdot f(n)$

Upon inspecting the equation $[g(n) \geq \frac{1}{c} \cdot f(n)]$, we can notice that it closely aligns with the definition of Omega notation (Ω),

Which implies that $g(n) = \Omega(f(n))$

Therefore, upon evaluation, **it is proved that $f(n) = O(g(n))$ if and only if $g(n) = \Omega(f(n))$**

Question 02

Problem 3-1 in CLRS textbook.

Asymptotic behavior of polynomials

where $a_d > 0$, be a degree- d polynomial in n , and let k be a constant. Use the definitions of the asymptotic notations to prove the following properties.

2.a

If $k \geq d$, then $p(n) = O(n^k)$.

Solution

Let $c_1 = \frac{a_d}{2}$, $c_2 = \frac{3 \cdot a_d}{2}$, $0 \leq c_1 \cdot n^d \leq p(n) \leq c_2 \cdot n^d$ for all $n \geq n_0$ (1)

For a given function $g(n)$, denote $O(g(n))$ as the set of functions given by –

$O(g(n)) = \{f(n) : \text{there exists positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq c \cdot g(n) \text{ for all } n \geq n_0\}$

From the given statement, $f(n) = p(n)$ and $g(n) = n^k$, where k is a constant.

Given $k \geq d \Rightarrow n^d \leq n^k \Rightarrow c_2 \cdot n^d \leq c_2 \cdot n^k \dots (2)$

From equations (1) and (2), $0 \leq p(n) \leq c_2 \cdot n^d \leq c_2 \cdot n^k$ for all $n \geq n_0$

Hence, $p(n) = O(n^k)$

2.b

If $k \leq d$, then $p(n) = \Omega(n^k)$.

Solution

For a given function $g(n)$, denote $\Omega(g(n))$ as the set of functions given by –

$\Omega(g(n)) = \{ f(n) : \text{there exists positive constants } c \text{ and } n_0 \text{ such that } 0 \leq c \cdot g(n) \leq f(n) \text{ for all } n \geq n_0 \}$

From the given statement, $f(n) = p(n)$ and $g(n) = n^k$, where k is a constant.

Given $k \leq d \Rightarrow n^k \leq n^d \Rightarrow c_1 \cdot n^k \leq c_1 \cdot n^d \dots (3)$

From equations (1) and (3), $0 \leq c_1 \cdot n^k \leq c_1 \cdot n^d \leq p(n)$ for all $n \geq n_0$

Hence, $p(n) = \Omega(n^k)$

2.c

If $k = d$, then $p(n) = \Theta(n^k)$.

Solution

For a given function $g(n)$, denote $\Theta(g(n))$ as the set of functions given by –

$\Theta(g(n)) = \{ f(n) : \text{there exists positive constants } c_1, c_2 \text{ and } n_0 \text{ such that}$

$0 \leq c_1 \cdot g(n) \leq f(n) \leq c_2 \cdot g(n) \text{ for all } n \geq n_0 \}$

From the given statement, $f(n) = p(n)$ and $g(n) = n^k$, where k is a constant.

From equation (1),

$0 \leq c_1 \cdot n^d \leq p(n) \leq c_2 \cdot n^d$ for all $n \geq n_0$

Hence, $p(n) = \Theta(n^k)$

2.d

If $k > d$, then $p(n) = o(n^k)$.

Solution

For a given function $g(n)$, denote $o(g(n))$ as the set of functions given by –

$o(g(n)) = \{ f(n) : \text{for any positive constants } c > 0, \text{ there exist a constant } n_0 > 0 \text{ such that}$

$0 \leq f(n) < c \cdot g(n) \text{ for all } n \geq n_0 \}$

From the given statement, $f(n) = p(n)$ and $g(n) = n^k$, where k is a constant.

Given $k > d \Rightarrow n^d < n^k \Rightarrow c_2 \cdot n^d < c_2 \cdot n^k \dots (4)$

From equations (1) and (4), $0 \leq p(n) \leq c_2 \cdot n^d < c_2 \cdot n^k$ for all $n \geq n_0$

Hence, $p(n) = o(n^k)$

2.e

If $k < d$, then $p(n) = \omega(n^k)$.

Solution

For a given function $g(n)$, denote $\omega(g(n))$ as the set of functions given by —
 $\omega(g(n)) = \{f(n) : \text{for any positive constants } c > 0, \text{ there exist a constant } n_0 > 0 \text{ such that } 0 \leq c \cdot g(n) < f(n) \text{ for all } n \geq n_0\}$
 From the given statement, $f(n) = p(n)$ and $g(n) = n^k$, where k is a constant.
 Given $k < d \Rightarrow n^k < n^d \Rightarrow c_1 \cdot n^k < c_1 \cdot n^d \dots (5)$
 From equations (1) and (5), $0 \leq c_1 \cdot n^k < c_1 \cdot n^d \leq p(n)$ for all $n \geq n_0$
 Hence, $p(n) = \omega(n^k)$

Question 03

Problem 3-2 in CLRS textbook.

Relative asymptotic growths

Indicate, for each pair of expressions (A,B) in the table below, whether A is O , o , Ω , ω , or Θ of B .
 Assume that $k \geq 1$, $\epsilon > 0$, and $c > 1$ are constants. Your answer should be in the form of the table with “yes” or “no” written in each box.

Solution

A	B	O	o	Ω	ω	Θ
$lg^k n$	n^ϵ	Yes	Yes	No	No	No
n^k	c^n	Yes	Yes	No	No	No
\sqrt{n}	$n^{sin n}$	No	No	No	No	No
2^n	$2^{n/2}$	No	No	Yes	Yes	Yes
$n^{lg c}$	$c^{lg n}$	Yes	No	Yes	No	Yes
$lg(n!)$	$lg(n^n)$	Yes	No	Yes	No	Yes

Question 04

For each algorithm write down all the possible formulas that could be associated with it.

Solution ¶

For the sake of understanding, I will tabulate them

Formula	A1	A2	A3	A4	A5
a	No	Yes	Yes	Yes	Yes
b	No	No	Yes	Yes	No

Formula	A1	A2	A3	A4	A5
c	Yes	No	Yes	Yes	Yes
d	Yes	No	No	Yes	Yes
e	No	No	Yes	No	No
f	No	No	Yes	Yes	No
g	No	No	Yes	Yes	No

Question 05

For the following algorithm: Show what is printed by the following algorithm when called with $\text{MAXIMUM}(A, 1, 5)$ where $A = [10, 8, 6, 4, 2]$? Where the function PRINT simply prints its arguments in some appropriate manner.

Solution

The final output will be — 8 10 6 10 2 4 4 10