

Lecture 13: Finite Automata – What is it and what are its types?

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Automata is a Plural of Automaton (i.e. Automatic)

self-controlled machine

- ✓ It's an other method for defining languages
- ✓ It's a Graphical Method
- ✓ Finite Automata (FA) is also called Finite Automata Machine (FAM)

Token changes Position on the input of certain number by Dice

State

Letters in alphabet

Tokens

Transition

Direction

Initial State

Final State

Normal State



Lecture 13: Finite Automata – What is it and what are its types?

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Limited

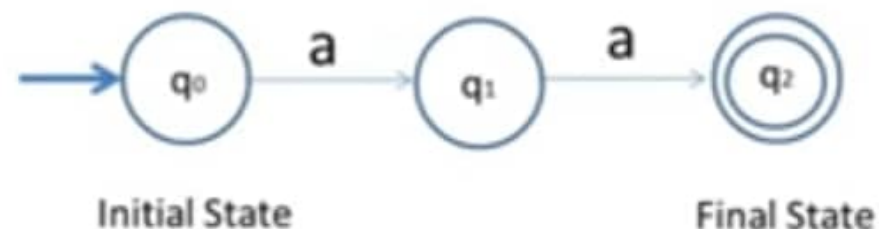
محدود

(Jis chez ki koi had ho)

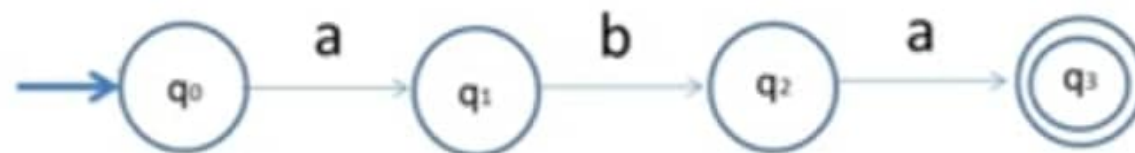
Q: Number of States
Q₀: Initial State
F: Final State
Σ: Letters in Alphabet
δ: Transitions / Movements



L1 = {aa}



L2 = {aba}



Lecture 13: Finite Automata – What is it and what are its types?

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What are types?

FA: $\{Q, Q_0, F, \Sigma, \delta\}$

- Q: Number of States
- Q_0 : Initial State
- F: Final State
- Σ : Letters in Alphabet
- δ : Transitions / Movements

FA / FAM

Deterministic Finite Automata (DFA)

Only one output



DFA

Non-Deterministic Finite Automata (NFA)

Can be any output
/ Output is not determined



NFA

- ✓ Finite Automata is also named as Finite Automata Machine (FAM)

What is it named as Finite Automata?

Automata: Plural of Automaton (i.e. automatic, self controlled machine)

Finite: Finite number of states / letters/ transitions

Lecture 14: Deterministic Finite Automata - What is DFA and how to draw DFA?

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Next state is known
(Has only one next state)

Plural of Automaton (i.e. automatic, self controlled machine)

Finite number of states, letters and transitions

Q: Number of States
Q₀: Initial State
F: Final State
Σ: Letters in Alphabet
δ: Transitions / Movements



Symbols:

DFA



NFA



L1 = {aa}



Initial State

Final State



Initial State

Final State



Initial State

Final State

Lecture 14: Deterministic Finite Automata - What is DFA and how to draw DFA?

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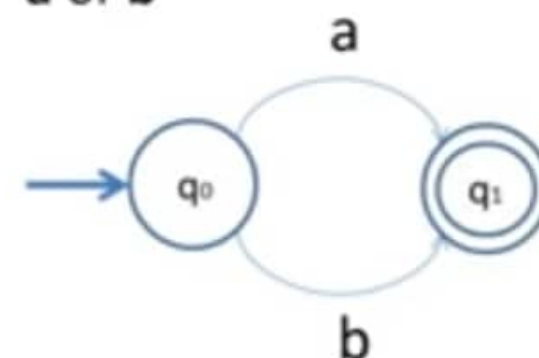


loop

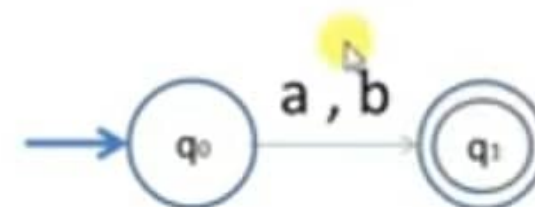


Contains **a or b**

$a+b$



Starts with a and contains any number of b's in end

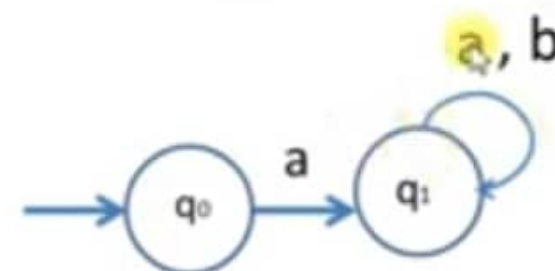


Contains any number of **a's** or **b's**



Stars with **a**

$$R = a(a+b)^*$$



Definition:

- ✓ It's a type of Finite Automata that used to define languages in which next state is already known
- ✓ It consists of following 5 Parts;

FA: $\{Q, Q_0, F, \Sigma, \delta\}$

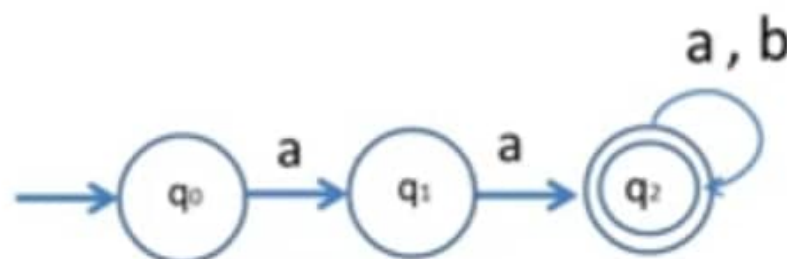
Q: Number of States

Q_0 : Initial State

F: Final State

Σ : Letters in Alphabet

δ : Transitions / Movements



How to draw Deterministic finite automata

DFA for a language that starts with **aa**

$$R = aa(a+b)^*$$

- Circle/Ovals for Initial State, Normal State and Final State
- Arrow for Transition with letters above
- Loop and OR
- Only one output; single letter can not go to multiple state from one state

What is it named as Deterministic Finite Automata?

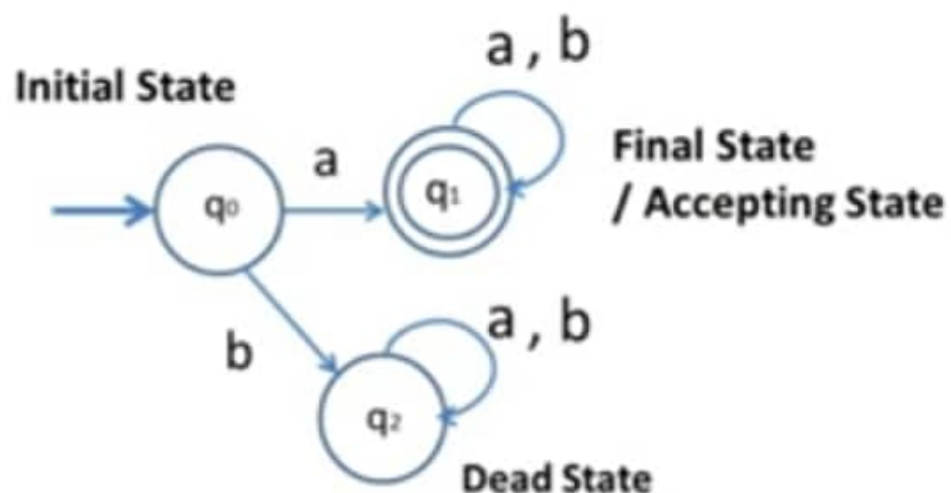
Automata: Plural of Automaton (i.e. automatic, self controlled machine)

Finite: Finite number of states / letters/ transitions

Deterministic: Output is already known (has only one next state)

Starts with a

$$R = a(a+b)^*$$



q_0 : Initial State

q_1 : final State

q_2 : Dead State (Dead End State)

Initial State

It is the state that the machine naturally starts in before it reads any input. It is called as **Entry Point**.



Final State

It is the state where the machine halts when it has no input left. It is also called **Accepting State**

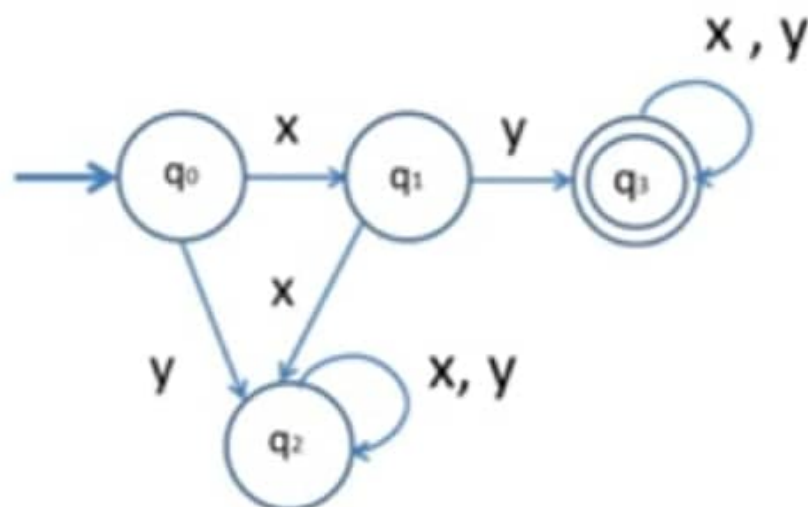


Dead State / Dead End State

It is also called as **Rejecting State** and **Trap State**. Once the machine enters a dead state, there is no way for it to reach an accepting state

Construct an FA which recognizes the set of all strings defined over $S = \{x, y\}$ starting with the prefix 'xy'.

$$R = xy(x+y)^*$$



xyxx

xyyyy

xyxxyy

yx

xx

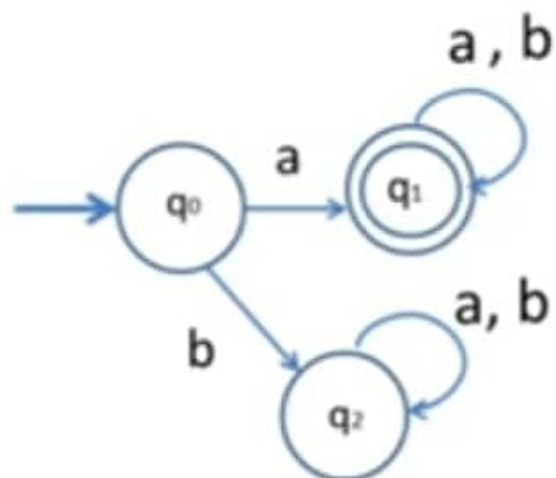
q₀: Initial State

q₃: Final State (Accepting State)

q₂: Dead State (Dead End State/Trap State/Rejecting State)

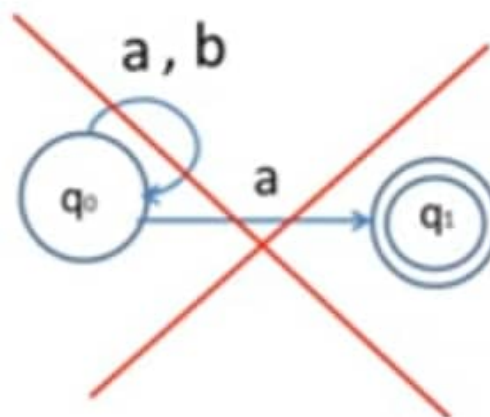
Starts with a

$$R = a(a+b)^*$$



Ends with a

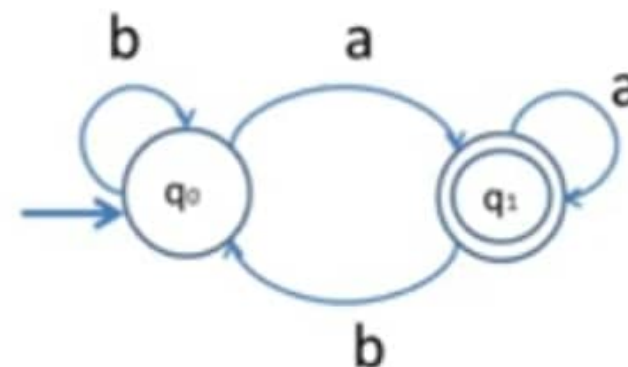
$$R = (a+b)^*a$$



a
aa
aaa
ba
baa
bbaa
aba
abba
babbaa

DFA:

One letter can not go to many state from one state.
(i.e., Only one output)



EVEN-EVEN

Even number of a's and Even number of b's

R = $[aa + bb + (ab+ba)(aa+bb)^*(ab+ba)]^*$

E =

RE =

$\Sigma = \{a, b\}$

aabb

bbaa

abab

baba

aa

bb

abaaab

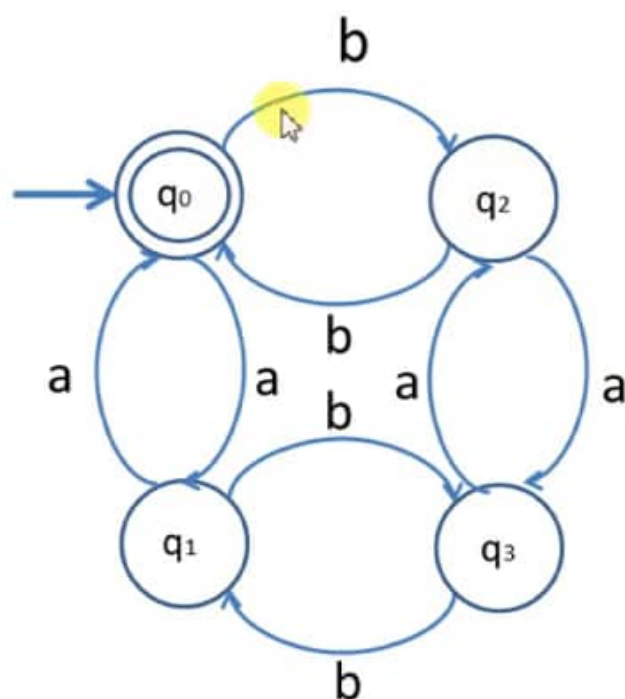
ab**bb**ab

Even Numbers

0, 2, 4, 6, 8, ..

Odd Numbers

1, 3, 5, 7, 9, ...

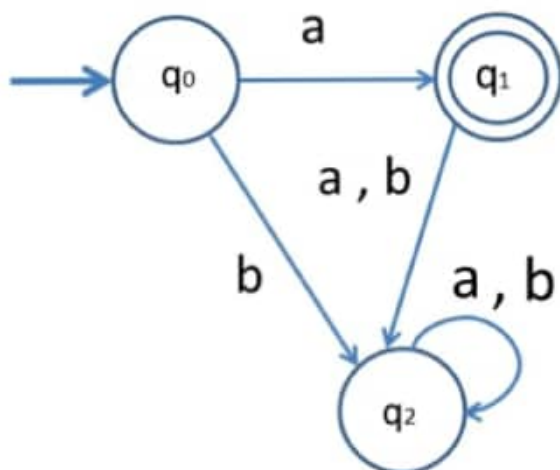


RE : Regular Expression

DFA : Deterministic Finite Automata

Q: Draw fa that accept exactly a

$R = a$



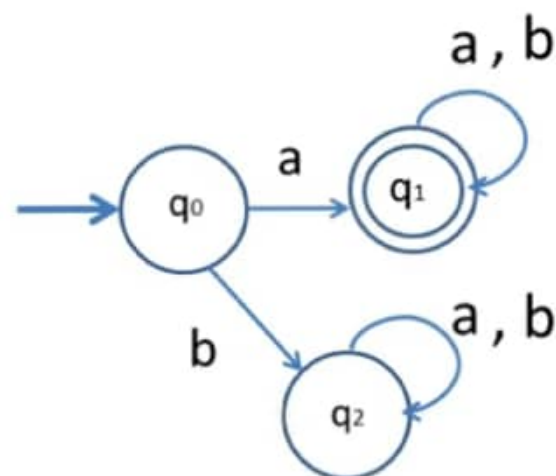
q_0 : Initial State

q_1 : final State

q_2 : Dead State (Dead End State)

Q: Draw fa that accepts all words starting with a

$R = a(a+b)^*$



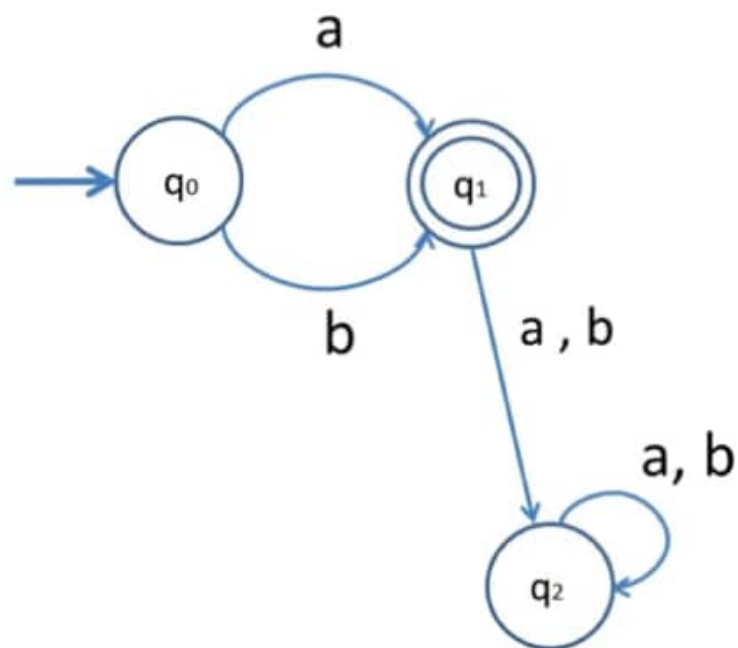
q_0 : Initial State

q_1 : final State

q_2 : Dead State (Dead End State)

Q: Draw fa that accepts exactly a or b

$R = a+b$

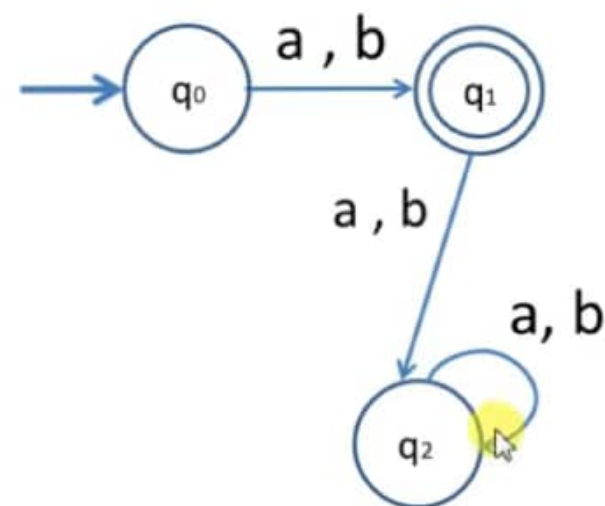


q_0 : Initial State

q_1 : final State (accepting state)

q_2 : Dead State (Dead End State / trap / rejecting)

Comma , = OR



q_0 : Initial State

Q: Draw fa for the language that have b as a second letter over $\Sigma = \{a, b\}$

$$R = (a+b)b(a+b)^*$$

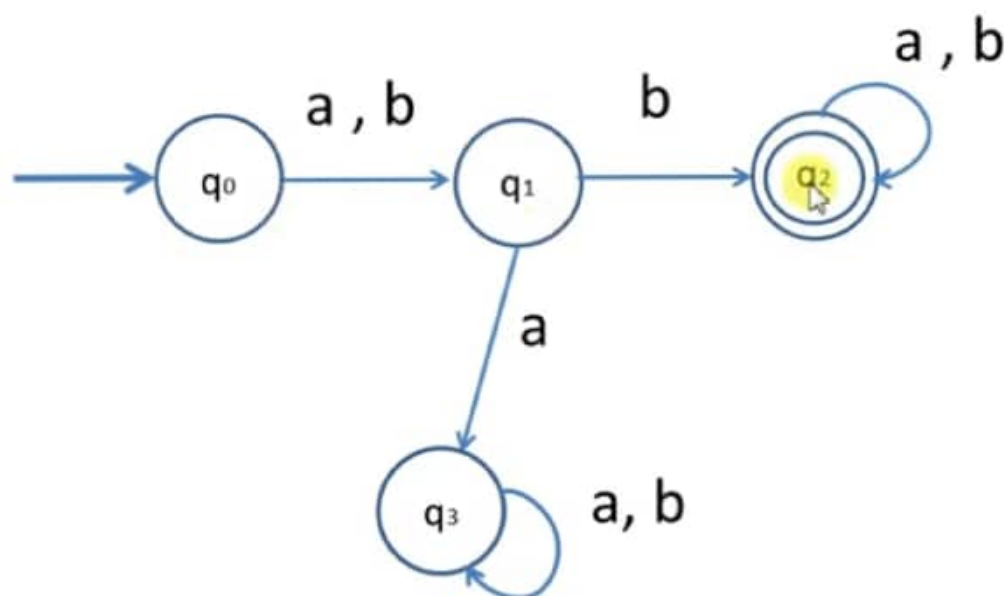
ab

bb

abaa

bbaa

abaababa



q₀: Initial State

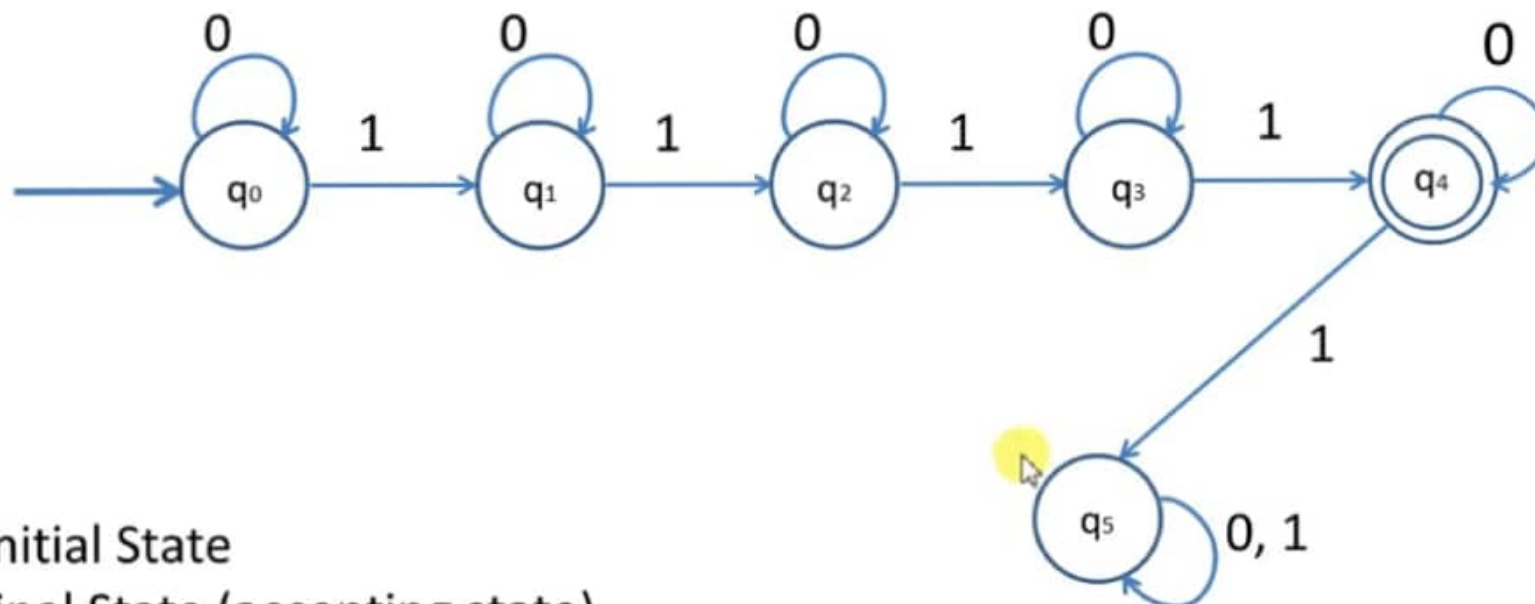
q₂: final State (accepting state)

Q: Draw fa for the language that have exactly 4 ones in every string over $\Sigma = \{0, 1\}$

$$R = 0^*1 0^*1 0^*1 0^*1 0^*$$

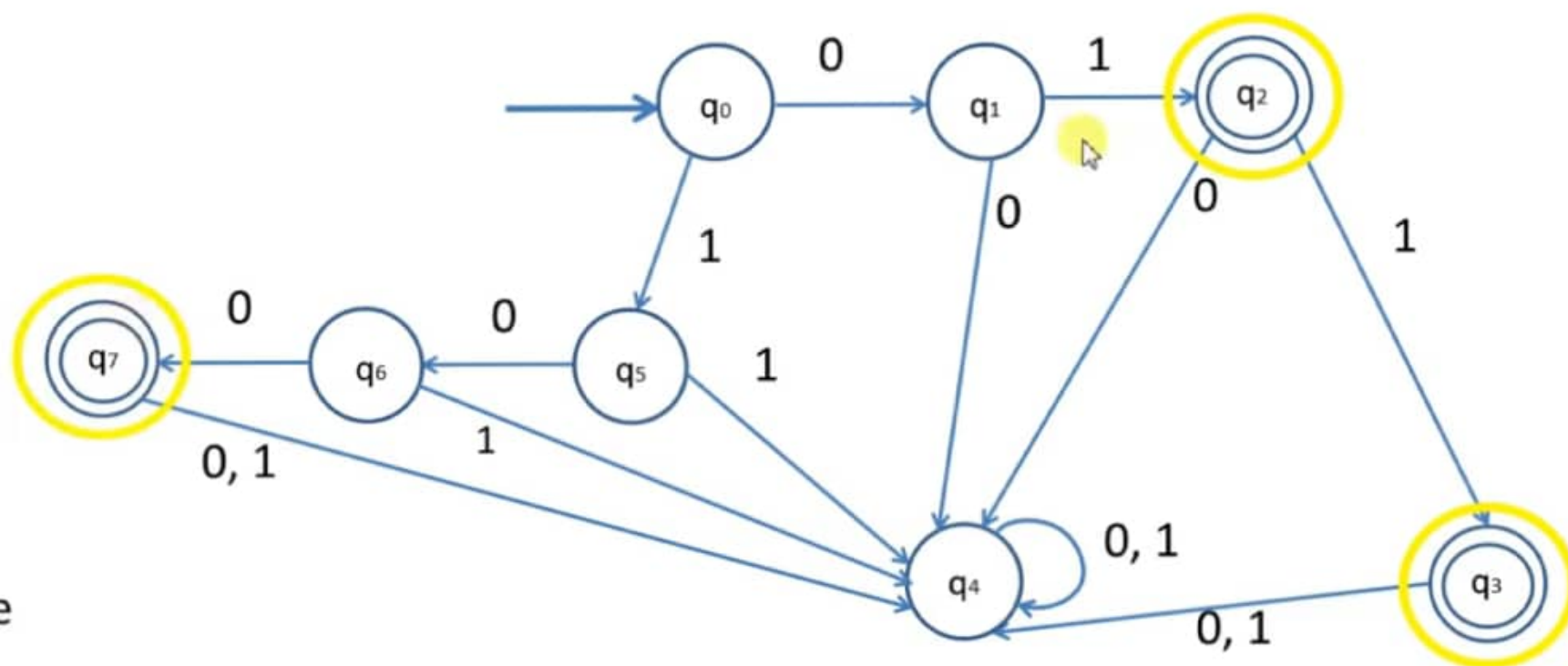
1111
01111
011110
010101010

q_0 : Initial State
 q_4 : final State (accepting state)



Q: Draw fa for the language that only accept $L = \{01, 011, 100\}$ over $\Sigma = \{0, 1\}$

$R = 01 + 011 + 100$

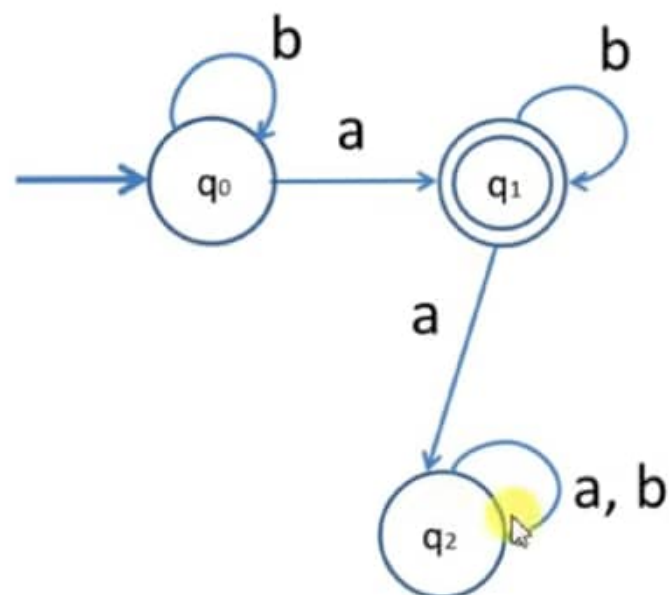


q_0 : Initial State
 q_2, q_3, q_7 : final State

Draw fa for the language $L = \{ w.n_a = 1 \mid w \in (a,b)^* \}$

$R = b^* a b^*$

a
ab
ba
bab
bbbabbb



q_0 : Initial State

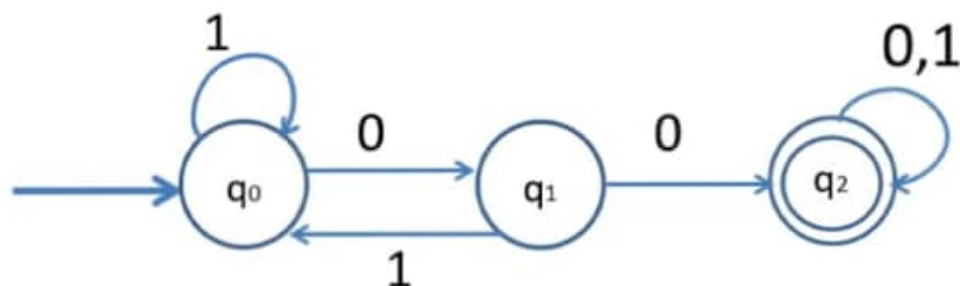
q_1 : final State

q_2 : Dead State (Dead End State)

Q: $L = \{ w \in \{0, 1\}^* \mid w \text{ contains } 00 \text{ as a substring} \}$

$R = (0+1)^* 00 (0+1)^*$

00
100
1001
101001
010001
11100111
10001



q_0 : Initial State

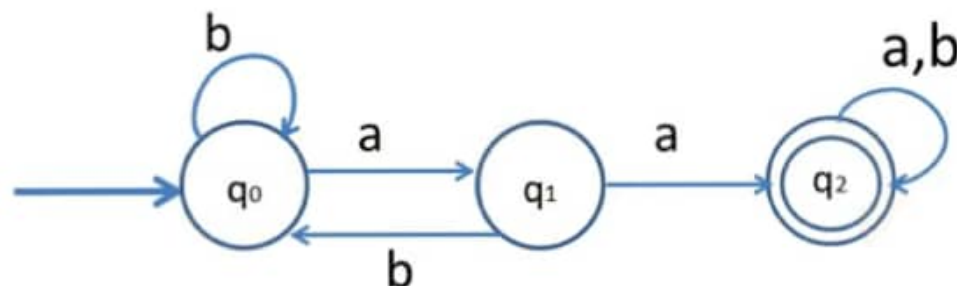
q_2 : final State (Accepting State)

Q: Draw fa that accept all strings with double aa in somewhere over $\Sigma = \{a, b\}$

$$R = (a+b)^* aa (a+b)^*$$

$$R = (0+1)^* 00 (0+1)^*$$

aa
baa
baab
babaab
abaaab
bbbaabbb
baaab

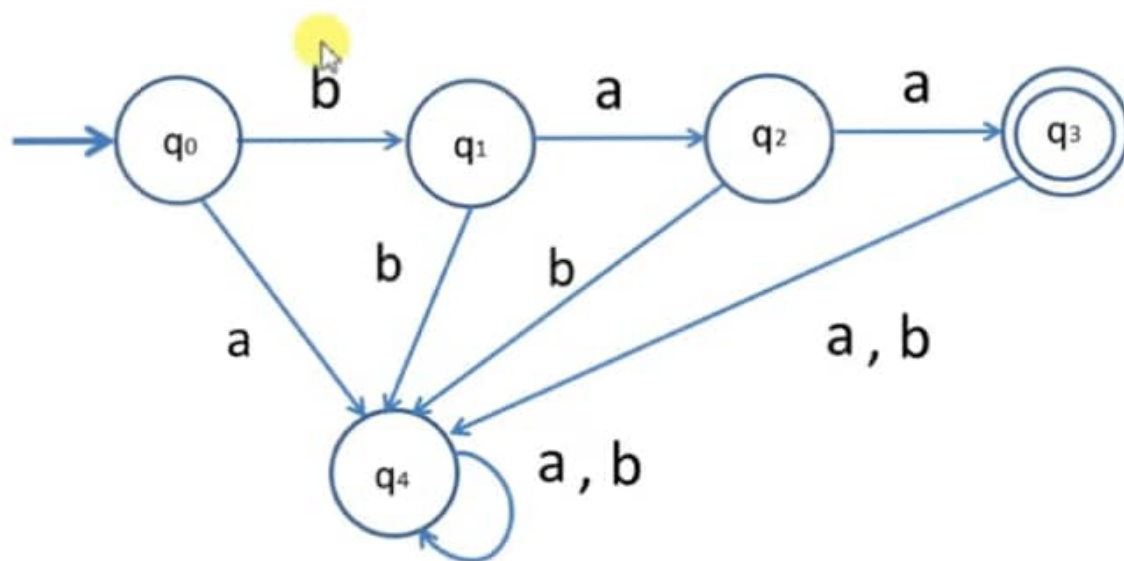


q_0 : Initial State

q_2 : final State (Accepting State)

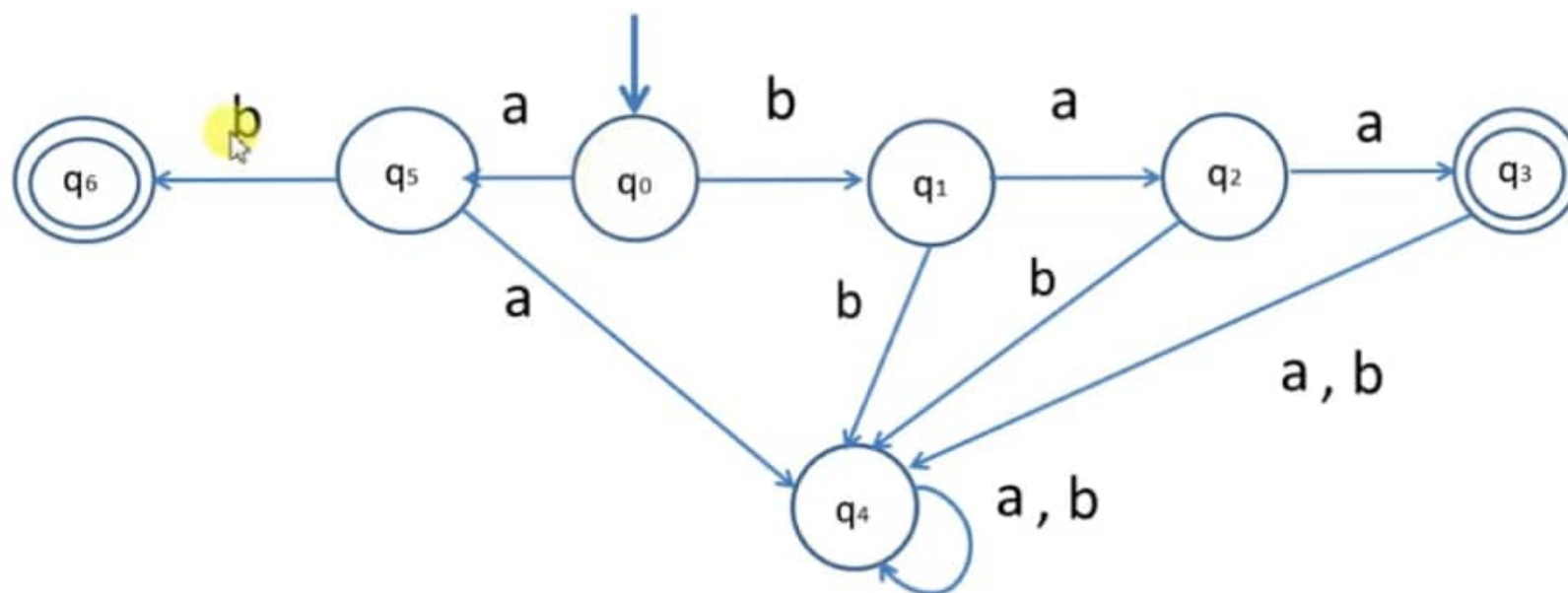
Q: Draw fa that contains exactly baa

R = baa



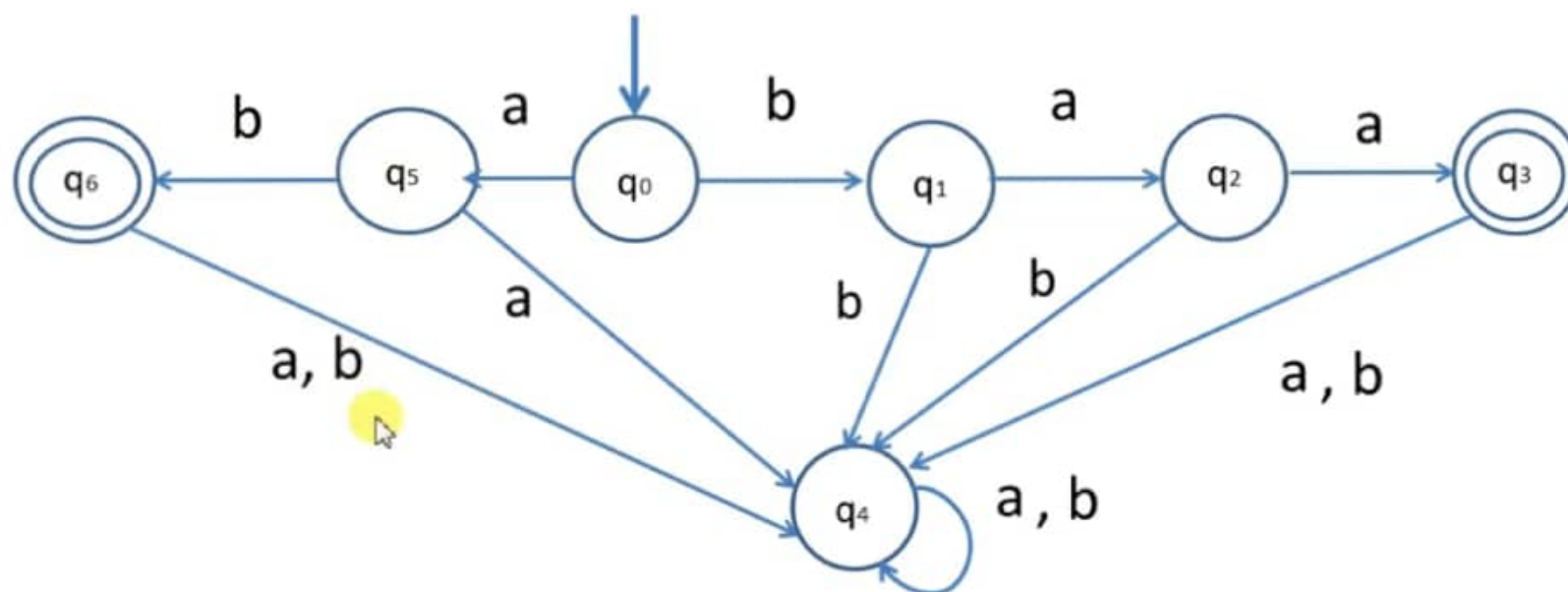
Q: Draw fa that accept exactly baa and ab

R = baa + ab



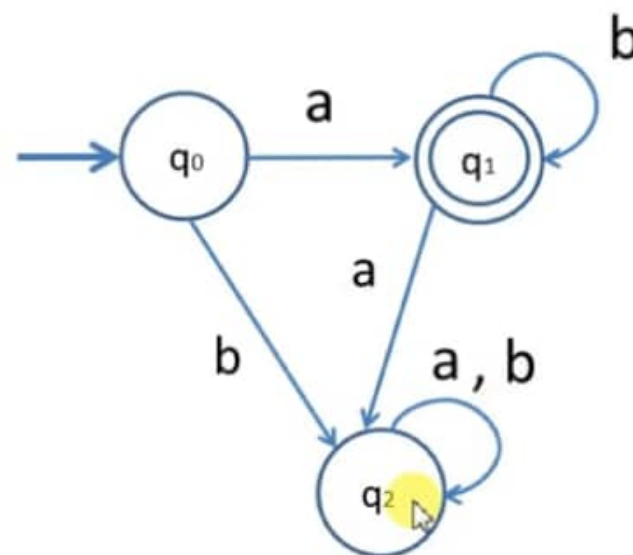
Q: Draw fa that accept exactly baa and ab

R = baa + ab



Q: Draw fa starting with a and contains any number of b's in end

$R = ab^*$



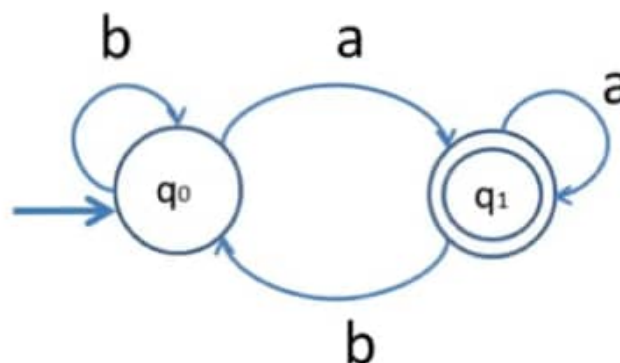
q_0 : Initial State

q_1 : final State (Accepting State)

q_2 : Dead State (Dead End State/Trap/Rejecting)

Q: Draw fa accept all words ending with a

$$R = (a+b)^*a$$



a
aa
aaa
ba
baa
bbaa
aba
abba
babbaa

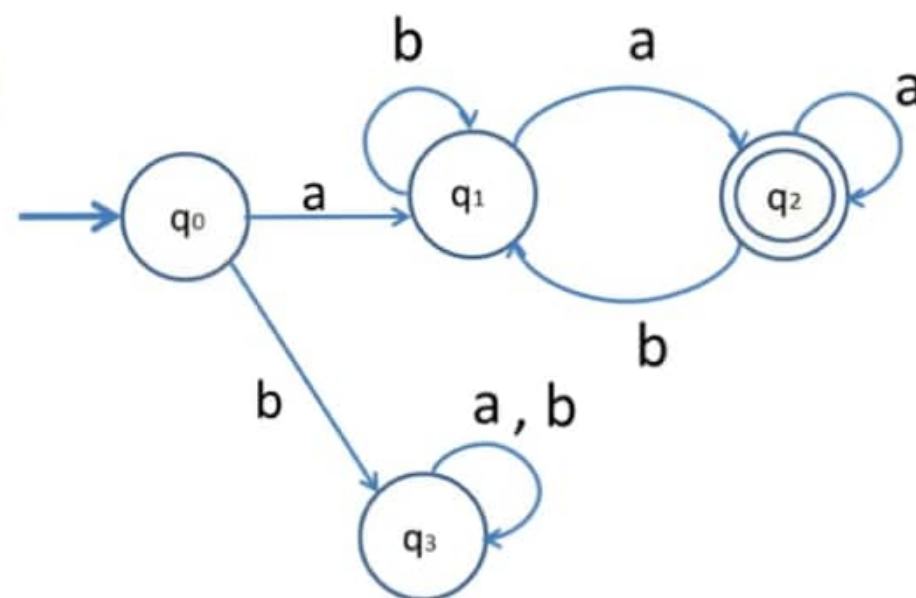
q₀: Initial State

q₁: final State (Accepting State)

Q: Draw fa that accept all words starting and ending with a over $\Sigma = \{a, b\}$

$$R = a(a+b)^*a$$

q_0 : Initial State
 q_2 : final State
 q_3 : Dead End State

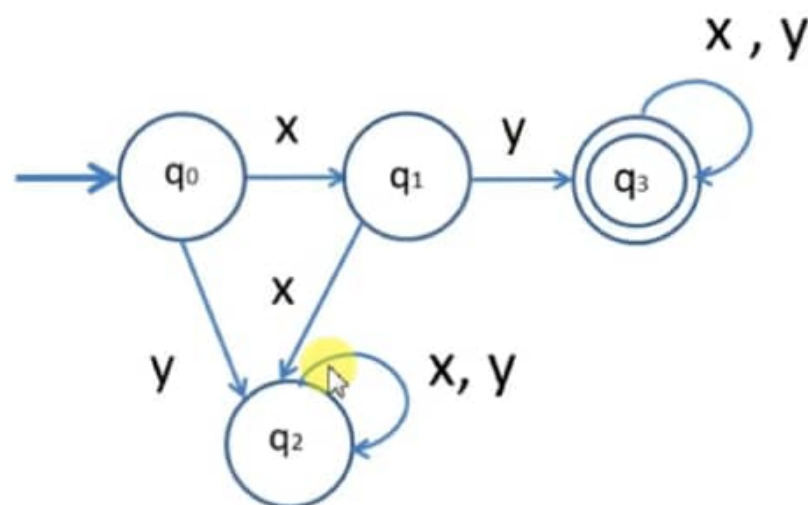


Lecture 17: DFA Important Examples and Solutions

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Q: Construct an FA which recognizes the set of all strings defined over $S = \{x, y\}$ starting with the prefix 'xy'.

$$R = xy(x+y)^*$$



xyxx
xyyyy
xyxxyy

yx
xx

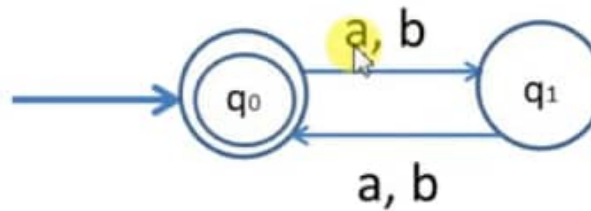
q₀: Initial State

q₃: Final State (Accepting State)

q₂: Dead State (Dead End State/Trap State/Rejecting State)

Q: Draw fa that accept all strings with even length over $\Sigma = \{ a , b \}$

$$R = (aa + ab + ba + bb)^*$$

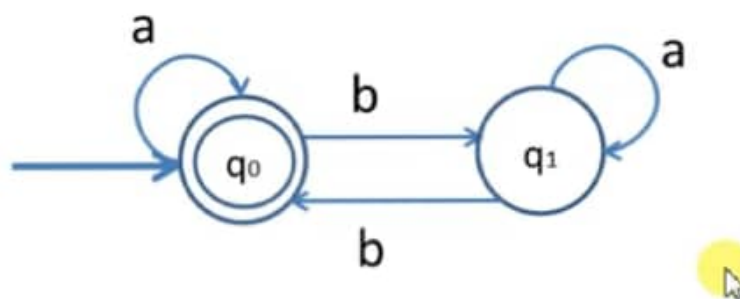


q_0 : Initial State

Q: Draw fa that accept all strings with even no of b's over $\Sigma = \{a, b\}$

$$R = a^* + (a^* ba^* ba^*)^*$$

bb
aaaa
baba
bbaa

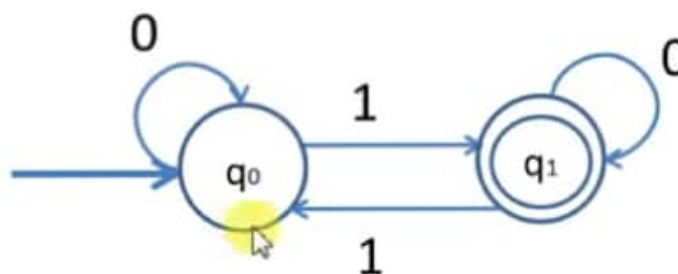


q_0 : Initial State

q_1 : Final State (Accepting State)

Q: Draw fa that accept all strings with odd no of 1's over $\Sigma = \{0, 1\}$

$$R = 0^*10^* (10^*10^*)^*$$



q_0 : Initial State

q_1 : Final State (Accepting State)