

Divide and Conquer: Solving Recurrences in DAA

Master Theorem, Substitution Method, Recursion Tree

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Introduction

- **What is a Recurrence Relation?**

A recurrence relation expresses the running time of a recursive algorithm.

- **Why Solve Recurrences?**

To analyze the time complexity of recursive algorithms like Merge Sort, Quick Sort, and Binary Search.

- **Agenda**

- Understand three key techniques for solving recurrences:
 - **Master Theorem**
 - **Substitution Method**
 - **Recursion Tree**

1. Substitution Method

- The Substitution Method solves a recurrence by guessing the solution and using mathematical induction to prove it.
- $T(n) = \begin{cases} T(n/2) + c, & \text{if } n > 1 \\ 1, & \text{if } n = 1 \end{cases}$
 - Pros: Precise control
 - Cons: Guessing the bound can be tricky

2. Master Theorem

- The Master Theorem provides a shortcut to solve recurrences of the form $T(n) = aT(n/b) + f(n)$ for divide-and-conquer algorithms.

- $a \geq 1, b > 1$

- Solution: $T(n) = n^{\log_b a} \cdot U(n)$ Where, $U(n)$ depends on $h(n) = \frac{f(n)}{n^{\log_b a}}$

- Pros: Fast and Direction
- Cons: Doesn't work for all recurrences

if	$h(n)$	
	$n^r, r > 0$	$O(n^r)$
	$n^r, r < 0$	$O(1)$
	$(\log n)^i, i \geq 0$	$(\log_2 n)^{i+1}$
		$i+1$

$$T(n) = T(n/2) + c$$

- $T(n) = T(n/2) + c; a = 1, b = 2, f(n) = C$
- $T(n) = n^{\log_b a} U(n)$
- $T(n) = n^{\log_2 1} U(n) = n^0 U(n) = U(n)$
- $U(n)$ depends on $h(n)$ $h(n) = \frac{f(n)}{n^{\log_b a}} = C = (\log n)^0 \cdot C$

if	$h(n)$	$U(n)$
	$n^r, r > 0$	$O(n^r)$
	$n^r, r < 0$	$O(1)$
	$(\log n)^i, i \geq 0$	$(\log_2 n)^{i+1}$
		$\frac{\quad}{i+1}$

answer: $O(\log_2 n)$

$$T(n) = 8T(n/2) + n^2$$

- Class Activity
- Answer: $O(n^3)$

3. Recursion Tree

- The Recursion Tree Method breaks down a recurrence into levels to visually sum the total work across recursive calls.
- Good for understanding how costs accumulate
- Helps form guesses for substitution
 - Pros: intuitive, helps with complex recursions
 - Cons: Can get messy for large recurrences

Thank You