

ECU - Execution control unit

RCU - Register control unit

EU - Execution unit

EBP - Base pointer

ESP - stack pointer

SI - source index

DI - destination index

CS - code segment

DS - data segment

SS - stack segment

ES - external segment

~~for EAX = ESI~~

H - High L - low

EAX - Accumulator register

EBX - Base register

ECX - Count register

EDX - Data register

ALU - Arithmetical logic unit

CU - Control unit

IP - Instruction pointer

FSB - frontside Bus

BIU - Bus interface unit

IOB - internal data Bus

1) editor

2) preprocessor

3) compiler

4) Assembler

5) linker

6) loader

i intermediate file

obj object file

asm assembly file / .S

.exe

load on RAM

T - Trap
 O - Diskin
 I - Interrup
 S - Output overlay
 R - Reg
 P - Auxiliary states
 C - carry

REG

AH	AL
BH	BL
CH	CL
DH	DL

BP	SP
CS	DS
ES	FS
SS	GS

ECU

CS	4
Text	
Data	
Stack	

EV

OS	
Text	
Data	
Stack	

RAM

Hello.exe	
Text	
Data	
SymbolsTable	

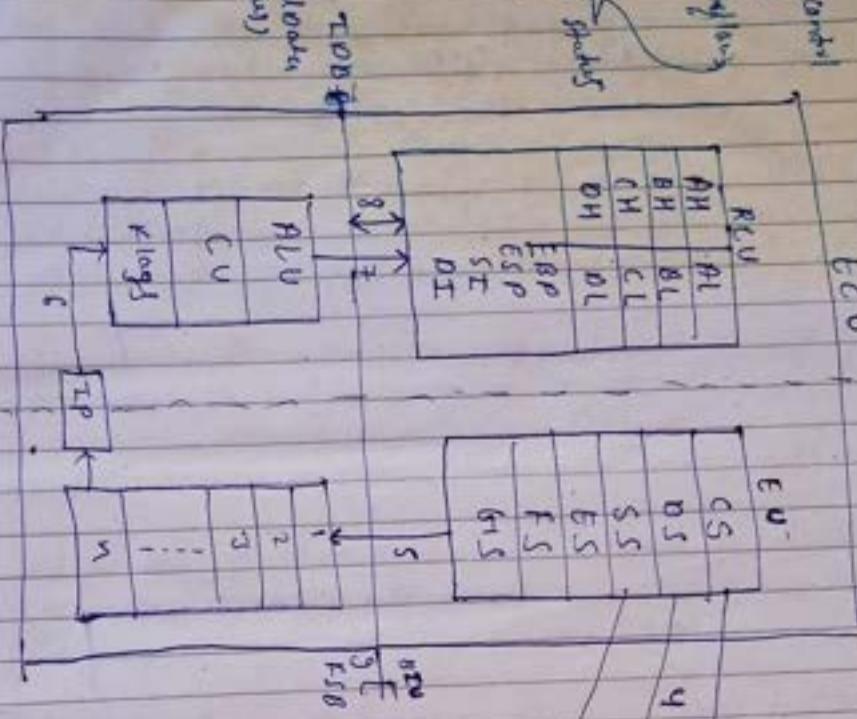
loader

Hello.obj	
Text	
Data	
SymbolTable	

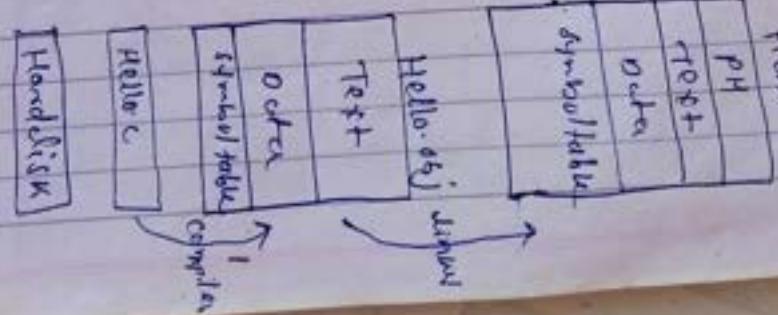
virtual

1	
Compiler	

(Internal data bus)



mother board



Computer Architecture

tool chain:- it is a set of programming tools which are used in chain to convert our program into executable file.

editor:- editor is a program which is used to write something on it and which is used to edit that written things (eg: wordpad, notepad)

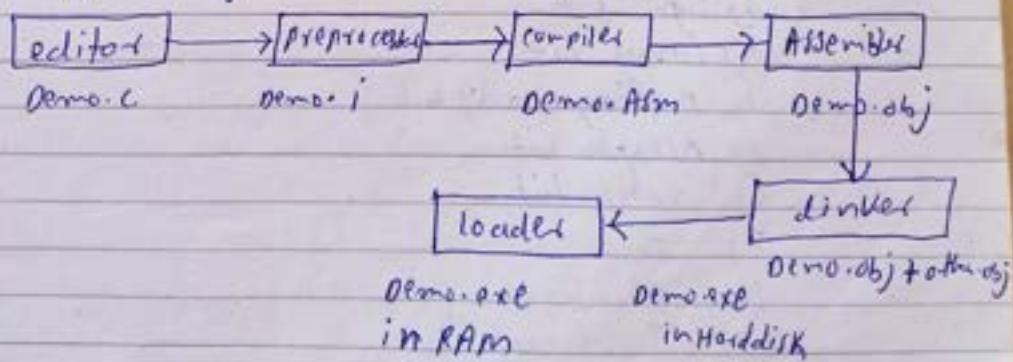
Compiler - compiler is a program which is used to convert high level programming language (C) into machine dependent programming language.
eg: GND C compiler, Turbo C, Borland, visual studio

Assembler:- it is a program which is used to convert machine dependent language into machine understandable language (binary) opasm, MASM, TASM

Linker:- it is a program which is used to link our object file with other object file

loader - it is a part of operating system which is used to load executable file from harddisk into RAM

Diagram of X86 tool chain



* Types of CPU registers -

1) General purpose registers -

- i) EAX (Accumulator register)
- ii) EBX (Base register)
- iii) ECX (Count register)
- iv) EDX (Outer register)

2) Index registers

- i) ESI (Source index)
- ii) EDI (Destination index)

3) Pointer registers

- i) ESP (Stack pointer)
- ii) EBP (Base pointer)
- iii) IP (Instruction pointer)

4) Segment registers

- i) CS (Code segment)
- ii) DS (Data segment)
- iii) SS (Stack segment)
- iv) ES (Extra segment)
- v) FS, GS (Additional forms)

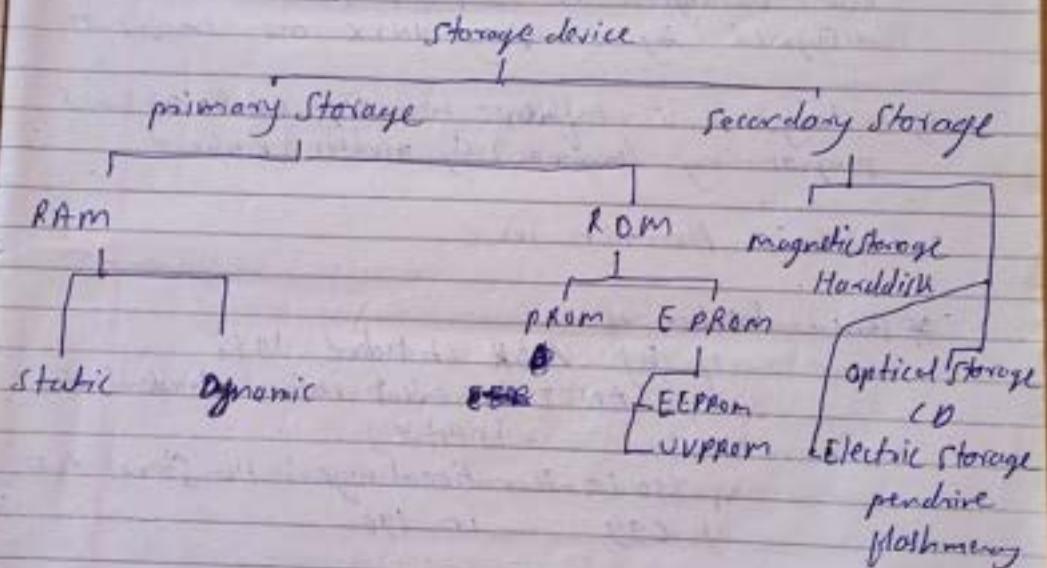
5) Flag registers -

- 1) Parity bit
- 2) Sign bit
- 3) Carry bit
- 4) Auxiliary carry bit
- 5) overflow bit
- 6) zero bit
- 7) Trap
- 8) Direction
- 9) Interrupt

* Components of Computer :-

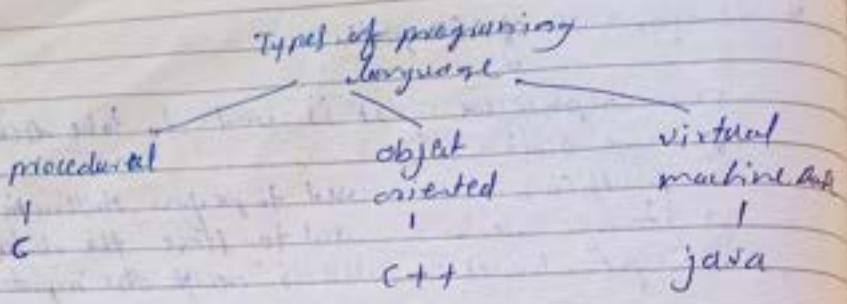
- 1) microprocessor :- it is used to take decisions and action
- 2) multi-co-processor - used to perform arithmetic computation
- 3) Storage device :- used to store the data
- 4) input devices :- used to accept the input
eg. keyboard, mouse
- 5) output devices :- used to display the output
printers, monitor

* storage devices in computer:-



* programming language -

it is a language which is used to communicate with machine (computer)



C programming language:

1) it is developed and designed by Dennis Ritchie from Unix in AT&T Bell it is designed by developing UNIX OS in 1969-1973.

2) C language is influence from BCPL (Basic combined programming language) by Martin Richard.

3) unix is written in c.

* Standardization of C

1) By K&R standard 1978

2) ANSI (American National Standardization Institute)

3) ISO (International organization Standardization)

4) C99 in 1999

* Characteristics of C language

1) high level programming language

2) procedural programming language

3) general purpose ——————

4) block structured ——————

5) case sensitive

6) compiled

7) platform dependent

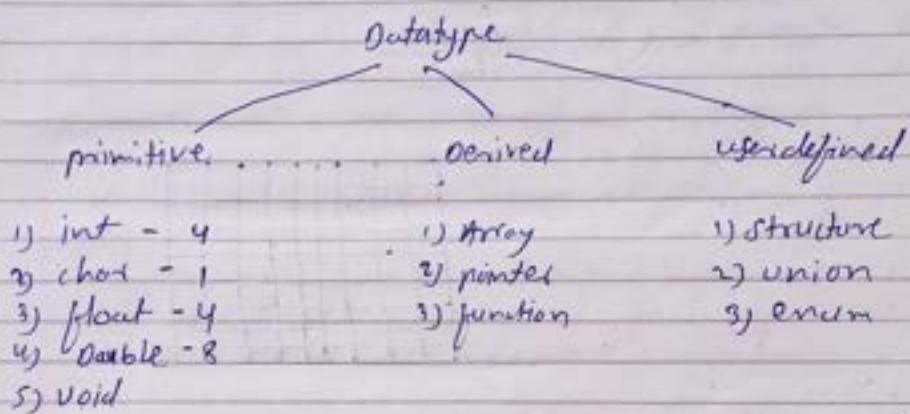
virtual
machine
1
java

was created
it is
in 1969-1972

is combined
language.

- 1) free-flow programming language
- 2) it uses top-down approach
- 3) it is standardize language

* Datatype in "C"



* Datatype qualifiers

it is used to provide new quality for our
Datatype

- 1) const (non-modifiable value)
- 2) volatile (Compiler optimization is removed)

* Datatype modifiers:-

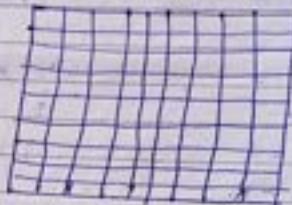
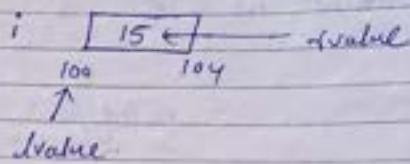
it is used to modify the range & size
of datatype

- 1) signed
- 2) unsigned
- 3) short
- 4) long

* Lvalue & Rvalue of variable

- 1) Lvalue :- it is the location of that variable
- 2) Rvalue :- it is the reflect of that location

int i = 15



* Declaration of variable:

it is a location at which the size of variable is decided by the compiler but memory is not allocated yet

* Definition of variable

it is a location at which actual memory gets allocated.

- 18 / 15

* Storage class:-

- 1) This concept is language independent concept means it is not specific to any programming language
- 2) The concept of storage class can be used in C & C++ also without any change in the concept
- 3) The concept of storage class is applicable for the variables which are used in our program or the functions which are defined in a program
- 4) There are 4 main storage classes in C & C++
 - i) auto
 - ii) extern
 - iii) static
 - iv) register
- 5) If we want to understand the concept of storage class then we have to consider 5 different characteristics such as
 - i) memory region in which our variable is allocated
 - ii) there are multiple memory option such as
 - i) stack section
 - ii) data section (BSS or nonBSS)
 - iii) static segment
 - iv) CPU registers

According to the type of storage class memory can be allocated in any of the above memory locations.

2) Default value :-

It indicates by default initialization is performed by the memory manager if the variable not initialized

Default value of variable can be

- i) it may contains some garbage value
- ii) it may contains 0 or 0.0 to '0'

3) scope:-

Scope of a variable is depend on the region in which our variable is accessible.

Scope of a variable can be :-

- 1) file scope
- 2) program scope

4) lifetime:-

Lifetime of variable is the gap in between memory allocation and deallocation of that variable

Lifetime of a variable can be :-

- 1) local lifetime
- 2) global lifetime

5) linkage:-

Linkage of variable indicates the regions in which our variable can be linked

There are 3 types of linkage such as

- 1) no linkage
- 2) internal linkage
- 3) external linkage

17 auto storage class:-

- i) memory - stack
- ii) default value - garbage
- iii) scope - Block scope or function scope
- iv) lifetime - local lifetime
- v) linkage - no linkage

- a) Any variable which is declared or defined inside any block or function then the storage class of that variable is considered as auto.
- b) for local variables auto is the default storage class.
- c) as it is default storage class use of auto keyword is optional.
- d) function argument, local variable as a default auto storage class
- e) to avoid the drawback of garbage valued auto variable should be initialised explicitly

e.g: void fun (int i)

```
    {  
        int j=11;  
        int k;  
    }
```

in above function i, j, k considered as auto variables.
This variable also can be defined as an
auto int j=11;
auto int k;

in above syntax auto keyword is not use because it is optional.

memory for this auto variable is allocated inside the stack frame of a function in which it is defined.

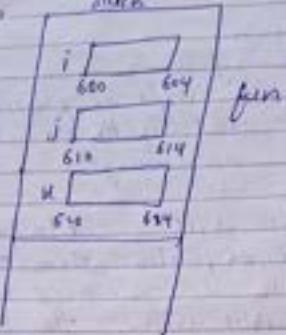
With example

Diagrammatic representation:

EBP → 700

Stack

ESP → 500



In above diagram stack frame of function func is displayed which contains the memory for all local variables.

According to the standard processor architecture cur. Stack grows downward.

The starting address of stack frame is stored in EBP and end address is stored in ESP

2) register storage class:-

i) memory :- CPU registers

ii) default value → garbage

iii) scope :- block scope or function scope

iv) lifetime - local lifetime

v) linkage - no linkage

a) The concept of No register storage class is almost same onto storage class.

b) The only difference betn auto and register storage class is the memory allocation region.

Ques 2

- e) if any of the variable is frequently required in our program then that variable should be declared as register variable.
- d) when we try to access any variable the value of that variable gets copied from RAM to CPU registers.
- e) if the memory for variable is already allocated inside the registers then time which is required to copy a value from RAM to CPU registers gets minimized.
- f) Register storage class is considered as a request because memory gets allocated inside CPU registers.
- g) if there is no CPU register available then the request for register storage class gets denied and then that variable gets default storage class as auto.
- h) whether the CPU register gets allocated or not is decided at runtime.
- i) there is no such direct technique by using we can predict the register storage class is applied or not.
- j) if we create the variable as a register variable then we can not fetch the address of that variable
eg: void fun () {

```
register int i=0;
for(i=0; i<1000; i++)
    printf("%d", i);
```

y

function func
removes for

candidate

need in exp

in scope

is almost

Her Storage

in above example i variable which is loop counter
is considered as a register variable.

As this loop counter required 1000 times in
our program due to which instead of accessing
from RAM we can store it into CPU register.

```
eg: register int i=11; // allowed  
register int arr[10]; // not allowed  
register float f=3.14f; // not allowed  
register double d=3.14; // not allowed  
register char ch='a'; // allowed
```

According to above syntax we can conclude that
register storage class can be applied only for
integral datatypes.

means it is only applicable for short and int but
it is not applicable for float and double.

we can not create derived datatype as array
which contains register variables

```
eg: register int i=11; //  
int main ()  
{  
    // code  
}
```

the above syntax is not a compiletime error
but it is a logical problem.

if we declare the variable as global variable then
memory for that variable gets allocated through
out the program and it is illogical to acquire

points

i) in
extern
variables

global
variables
used
in
functions

all that
by far

int but
all

day

of

then
through
will

the CPU registers throughout the program.

- # extern storage class:-
 - i) memory - data
 - ii) default value - 0 or \000
 - iii) scope - throughout program scope
 - iv) lifetime - global (throughout program)
 - v) linkage - external linkage

If we declare or define any variable as a global variable then default storage class gets allocated for that variable in system.

The concept extern storage class is generally used if our project contained multiple program files.

Eg: // first.c

```
int i=11; // variable definition
void fun() // function definition
{
    int n;
}
```

// second.c

```
int main()
{
```

```
    extern int i; // variable declaration
    printf("%d", i); // 11
    extern void fun(); // function declaration
    fun(); // call
}
```

In above project there are 2 .c files as first.c and second.c.

if we want to write the project in multiple files
then main function should be in one
file only and in other file we can write
anything.

to compile this multifile program we have to
use command of
gcc first.c second.c &t

if we want access the contents from some other file
then we have to declare that variable as extern or provide
as a extern in our file

if it is not declared by using extern keyword
then compiler generates error

by using the extern keyword our compiler search
externally for that variable.

as like about application we can write multiple
files in one project & we can access the content
one file into other file by using extern storage
class.

in this program variable u is local variable to
the function fun due to which we can not
access that variable by using extern keyword.

4. static Storage class:-

- i) memory - data section of static segment
- ii) default value - 0,00 \$ \0
- iii) scope - file scope
- iv) lifetime - global lifetime (throughout program)
- v) linkage - internal linkage

C Static variable

19/7/15

- 1) ~~Temporary~~ In case of auto storage class memory is allocated for the variable inside the stack frame of that function.
- 2) when the function terminates memory gets automatically free due to which we can not pollute values of that variable.
- 3) if we want persistent local values then it should be allocated outside the stack frame.
- 4) for this purpose static storage class is used which is used to preserve the values of local variables across the function calls.

5) void fun ()

```
int i=10;
i++;
if (i==15)
{
    fun();
}
```

```
int main ()
{
```

```
    fun();
    return 0;
}
```

6) in above program inside fun function there is one local variable named as i which is initialized to 10.

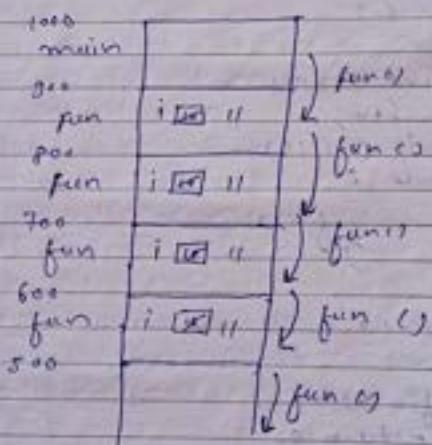
7) this function gets called recursively from that function again and again

i) for every function call separate stack frame gets created which contains value of $i = 10$

ii) due to which the condition which is written inside that function will not get executed successfully because memory for local variable i gets allocated for every function call.

iii) due to this drawback our program abnormally terminates when memory for stack gets full if it generates an runtime exception or stack overflow

iv) we can avoid this problem by declaring that i variable as static.



(2) to avoid this problem our modified code should be

void fun()

Static int $i = 10;$

$i++;$

stack frame self created

which is written
I get created
for local variable
fun call

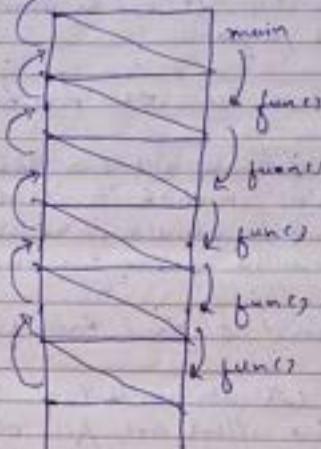
grows downwards
stack gets first and
at stack overflow
overwriting that

if (i < 5)

 y = func();

y =

0.5



- * There are two types of static variables as
 - ↳ local static
 - ↳ global static

e.g.: static int global = 10;
 int x = 10;

void func()

{

 Static int local = 10;

 y

 int main ()

{

 func();
 return 0;

y

in above program local is considered static local variable
global is considered as static global variable

The similarity b/w these two types of variable is it maintains its value throughout the execution of program.

and difference b/w these two types of variable is

i) Local static variable is accessible only inside the block in which it is return.

ii) Global static variable is a variable which is accessible inside file in which it is written.

* Difference global & global static variable

i) Lifetime of both the variables is throughout the program but there are few differences b/w of

a) global variable can be accessible outside the file by using extern keyword.

b) global static variable can not be accessible outside the file because it has an internal linkage.

* Symbol table:-

i) This table maintains metadata about the data of our program.

ii) This variable table maintains all the identifiers which are used in our programs. This table gets created by the compiler as well as operating system.

is controlled
global variable

scope of variable
throughout the

symbol table

variable only visible
return

variable which
is not been written

variable

throughout the
different block of

outside the file

to be accessible
can introduce

about file

identifier
this
as well as

- g) all the information which are required for identifier
is maintained in symbol table

Name of identifier	Address (virtual)	Value	from	To	Annotation

- a) this table contains the entries such as name of
the identifier used in the program
- b) the address which is allocated by the compiler
which is virtual address
- c) the value of the variable if it is initialised
- d) the starting address of variable is accessible
- e) end address till which we can access our variable
- f) If variable has a reference variable which is used
in C++ then its name stored in another column.

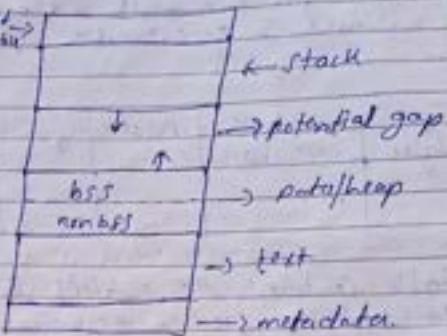
* layout of program

Below diagram explains the layout of an
executable file.

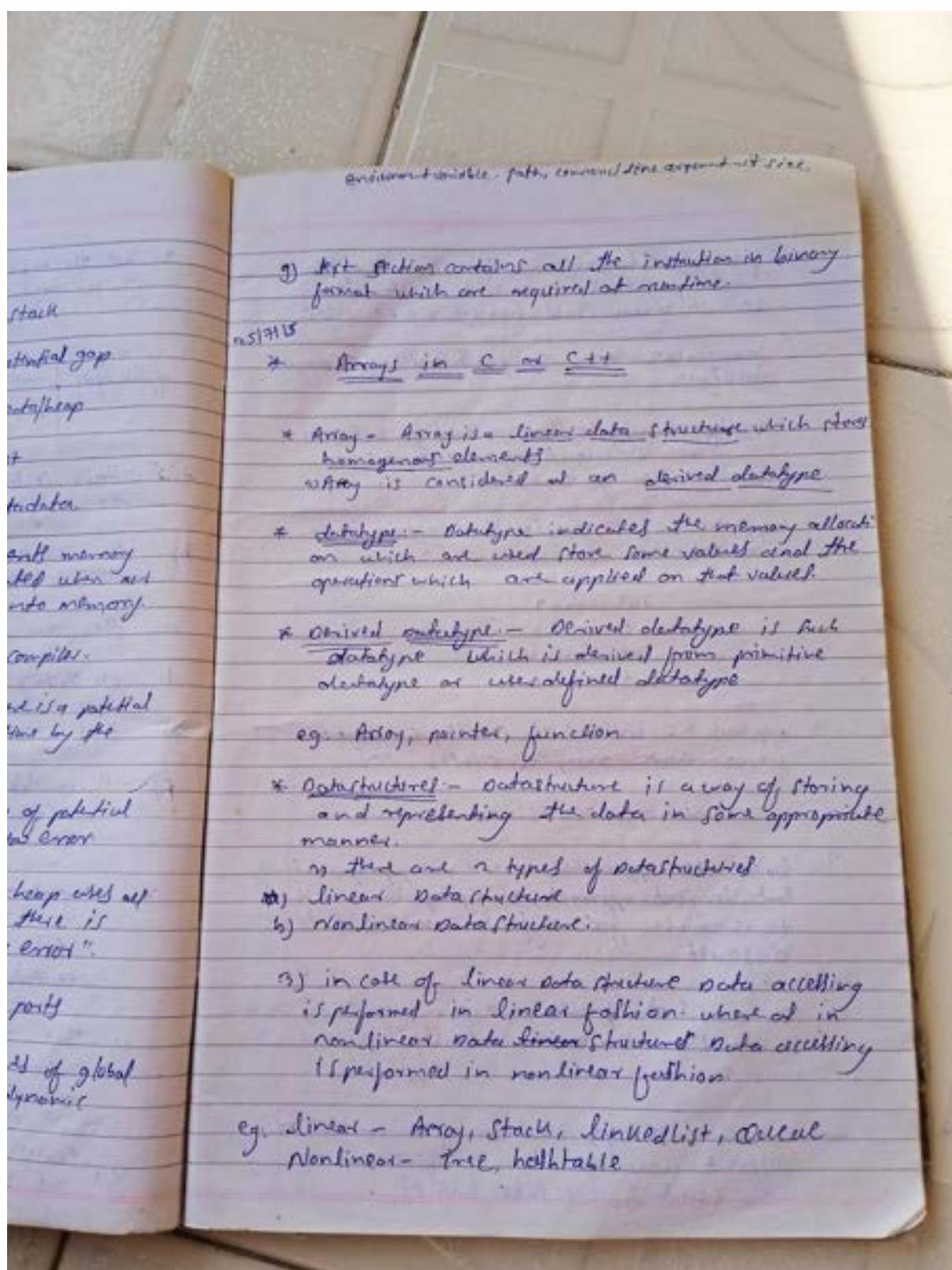
- 1) The topmost part of our layout containing the
memory for environment variables & commandline
argument
- 2) Environment Variable are such variable whose values
are set by running the program.

Democ. ELF

commandline arguments
& Environment Variables



- 1) after the commandline arguments memory for stack gets allocated when an executable file gets loaded into memory.
- 2) size of stack gets fixed by the compiler.
- 3) below the memory of stack there is a potential gap which is used at a runtime by the stack or by data sections.
- 4) if our stack uses all the memory of potential gap then there is stack overflow error.
- 5) and if our data sections or heap uses all the memory of potential gap there is "there is no enough memory error".
- 6) data section is divided into 2 parts bss and rodata which is used to store values of global variable and the memory has dynamic memory allocation techniques.



- if we want to store multiple elements of same datatype then we have to create array because array stores homogeneous elements.

- Homogeneous means every element is of same datatype

e.g. if i want to store marks of 5 students in mathematics then there are 2 programmatical approaches of

approach 1: we can create 5 different integer variables such as

```
int marks1;  
int marks2;  
int marks3;  
int marks4;  
int marks5;
```

approach 2: instead of creating 5 different variables we can create array of 5 integers as

```
int marks[5];
```

- Both the approaches are working programmatically but in first approach as a programmer we have to remember five different names which is difficult.

- To avoid this problem we can create an array in which we have to remember only one name.

- In above 2 approaches if we create N different variables then there are N different symbol table entries

elements of some
variable
and elements

it is of some

of 5 shall be
are 7

first integer

different variables
of

programmatically,
we have
with 15

to an array
you

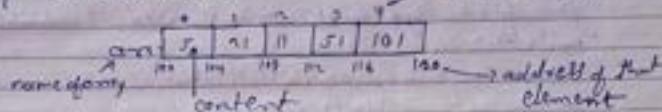
to N

But if we create array of 10 elements then
there is only one symbol table entry

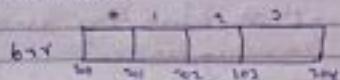
- there are 2 types of array of
 - 1) single dimensional array
 - 2) multi dimensional array

* diagrammatic representation of array

1) int arr[5]; *index of array*



2) char arr[4];



- 3) if we want to access the element of array then
we have to use 2 things of
 - 1) name of the array
 - 2) index of that element

- 4) According to above diagram if we want to access
element 51 then Syntax should be

arr[3];

Similarly if i want to initialize character 'A' in
3rd index of array arr then Syntax should be

arr[2] = 'A';

- 5) there are multiple ways in which we can initialize
the contents of array

eg. int arr[5] = { 10, 20, 30, 40, 50 };

b) int arr[5];
arr[0] = 10;
arr[1] = 20;
arr[2] = 30;
arr[3] = 40;
arr[4] = 50;

The first way initialization is called as member initialization list.

The next way of initialization is called as member by member initialization technique.

6) if we know all the elements of array then we can use member initialization list but if we want to initialize only some of the elements then we have to use member by member initialization technique.

7) eg. int arr[5];
arr[1] = 20;
arr[2] = 40;

In above syntax we want to initialize only 2 members due to which member by member initialization technique is used.

8) If array is uninitialized then the content in that array depend on storage class of that array.

Eg. int arr[4]; // extern storage class
int main ()
{ }

int arr[4]; // auto storage class
return 0;

3)

ii) At both the arrays are uninitialized contents of array arr is zero and content arr[0] is garbage ^{like} auto storage class.

- (a) if first element of array is initialize by using member initialization list then all the remaining elements of that array gets its default value of 0, 10, or 0 etc.

int arr[5] = {5, 3};

0	1	2	3	4
100	104	108	112	116

This above concept is not applicable if we are using member by member initialization technique

- ii) in case of member initialization technique we can not initialize random element of

int arr[5] = h, 4, 3, 2; // error

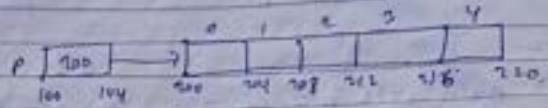
- 13) There are 2 ways in which we can allocate the memory for array

- static memory allocation.
- dynamic memory allocation

Ans int arr[5]; // static allocation

0	1	2	3	4
100	104	108	112	116

// dynamic memory allocation
int *p = (int *) malloc (5 * sizeof (int));



- 11) In above case memory for array is allocated on heap and our pointer p holds the base address of that allocated memory.
- 14) In case of dynamic memory allocation all the values are by default set to zero.

- 15) In case of static memory allocation size of array should be compile time integral constant otherwise compiler generates an error.

e.g. int arr[5]; // allowed
a) int no=5;
int arr[nos]; // not allowed
it is not allowed because no is variable.

b) const int no=5;
int arr [no]; // error
it is allowed on some of the platforms and
some of the compilers

c) #define MAX 5
int arr [MAX]; // allowed
it is allowed because after preprocessing our
syntax should be
int arr [5];

- d) int arr[]; // not allowed
- e) int arr[3+2-1]; // allowed

if (int);

4
16 320

is allotted on
the Base address

using all the
zero.

for size of
legal constant
array.

i variable.

and

by all

while specifying the size of array we can use
any expression that will get evaluated to
compiletime integral constant.

16) Specifying the size of array is optional if we are
using member initialization list

int arr[] = {1, 2, 3, 4, 5}; // allowed

in this syntax compiler calculates no of elements
which are initialize and depend on which it
decides the size of array.

17) Name of array is considered as Base address of
that array

As Base address is the only entity which is available
to the compiler to access every element of
array due to which we can not modify base
address of array.

int arr[4] = {10, 20, 30, 40};

10	20	30	40
100	104	108	112

printf("%d", arr); // 100
arr++; // error

18) after understanding concept of array we have to
read the Syntax of array properly as

int arr[5];

arr is one dimensional array
which contains 5 elements
each element is of type integer.

`int brr[3][4];`

brr is 2 dimensional array
which contains 3 one dimensional arrays
each one dimensional array contains 4 elements
each element is of type integer.

`int arr[3][2][4];`

arr is 3 dimensional array
which contain 3 two dimensional array
Each 2 dimensional array contains 2 one
dimensional array
Each one dimensional array contains 4 elements
and each element is of type integer.

"pointers in C & C++" 16.7.15

1) pointer is a language independent concept means it exactly same in C & C++.

e.g.
is already

2) pointer is considered as a system programming primitive in C & C++.

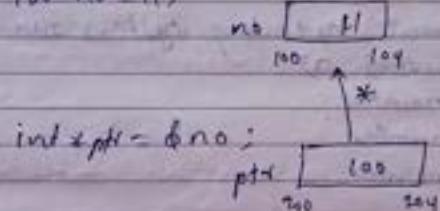
3) pointer is a derived datatype which can be derived from any primitive datatype, any user defined datatype or any derived datatype.

4) pointer - pointer is a variable which stores address

& the capability of that variable is to fetch the content whose address is stored in that variable.

already

e.g. int no=11;



⇒ operators used in pointers:-

'&' is a unary operator which is called as address operator which is used to refine address of any type of variable.

'*' operator is used to refine the virtual address of the variable.

'*' : it is unary operator which is called as content of operator. This operator is used to fetch the contents whose address is stored in pointer variable.

Some important point about pointers

- 1) As pointer is a variable which store address and additiblty is always integers due to which size of any pointer is always 4 byte
- 2) Means size of pointer is 4 byte irrespective of type of pointer
- 3) pointer is also consider as a normal variable just the extra capability is added to our pointer which is used to fetch the contents where address is stored in pointer

Ans
Ques

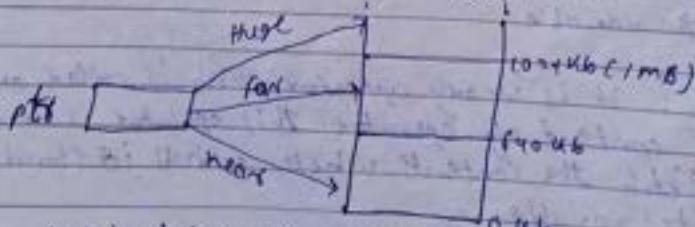
Type of pointers:-

- 1) In real mode operating system there are 3 types of pointers as
 - 1) near pointer
 - 2) far pointer
 - 3) huge pointer
- 2) Real mode OS are 8 bit operating system which can access only one 1mb of memory from RAM
- 3) In Real mode OS only single process run at a time.

0-640kb - near pointer

640kb-1024kb - far pointer

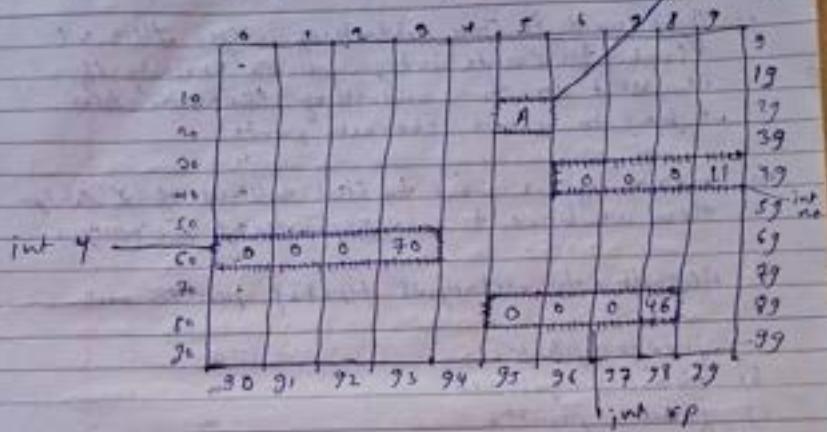
1024kb-above - huge pointer



but in today's OS which work in protected mode there is no such type of pointer that OS uses the concept of flat addressing scheme where there is no partition of memory

due to which today's compiler and C.R. don't support
concept type of pointers.

* Enclosed representation of pointers in RAM



above diagram is 16 bytes portion of RAM every small block indicate 1 byte memory when we declare any of the variable in our program memory gets allocated in this manner

According to this diagram our pointer gets sequential 4 bytes of memory like a normal integer variable

This pointer stores only address but by looking at the contents of the pointer we can not predict it is pointer or normal variable.

U type
concept of
memory

Some important point about pointers

- 1) As pointer is a variable which store address and additiblty is always integers due to which size of any pointer is always 4 byte
- 2) Means size of pointer is 4 byte irrespective of type of pointer
- 3) pointer is also consider as a normal variable just the extra capability is added to our pointer which is used to fetch the contents where address is stored in pointer

Ans
Ques

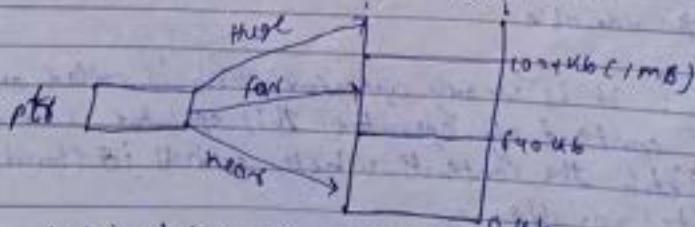
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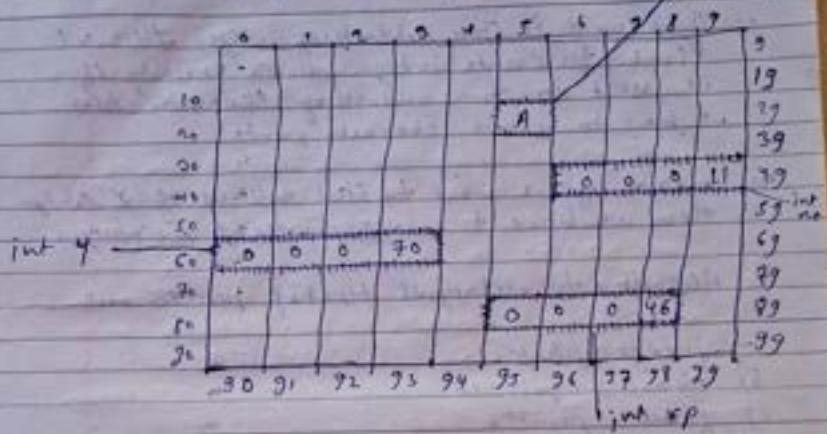
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due to which today's compiler and C.R. don't support
concept type of pointers.

* Enclosed representation of pointers in RAM



ans 3

char ch = 'A';

int y = 70;

int no = 11;

int *p = &no;

which can
be written as

above diagram is 16 bytes portion of RAM every small
block indicate 1 byte memory when we declare
any of the variable in our program memory gets
allocated in this manner

According to this diagram our pointer gets sequential
4 bytes of memory like a normal integer variable

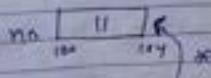
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U p to
concept of
memory

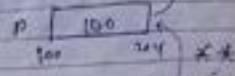
Types of pointer depend on datatype.

- 1) if we want to store address of integer then we have to create integer pointers. Similarly if we want to store address of character then we have to create character pointers.
- 2) intone if we want to store address of X datatype then we have to create pointers of X datatype.
- 3) According to dataype types of pointers are:
 - a) integer pointers
 - b) character pointers
 - c) float pointers
 - d) double pointers
 - e) structure pointers
 - f) union pointers
- 4) all this type of pointers require 4 bytes of memory irrespective of type of pointer.
- 5) if we store address of pointer then it is considered as pointer to pointers.
- 6) To understand this concept considered below example

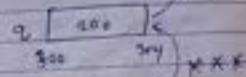
int no = 10;



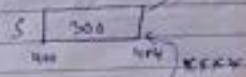
int *p = &no;



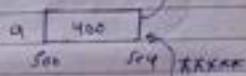
int **pp = &p;



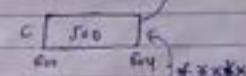
int ***ps = &p;



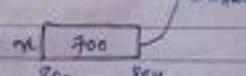
int ****pd = &c;



int *****pc = &d;



int *****n = &d;



p; // 100

no; // 11

*p; // 11

&p; // 100

**p; // 500

***d; // 300

****a; // 11

d s; // 400

a(***) ; // 100

a(xp) ; // 100

(***c) ; // 100

**z ; // 11

a(**z) ; // 100

*****d ; // 1000

*****e ; // 100

a(**xxc) ; // 100

a(**xxd) ; // 400

* Arithmetic operations on pointers:-

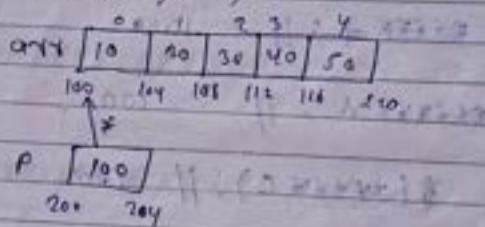
1) As pointer is a variable like some other variable we can perform all arithmetic operations on pointers.

2) After understanding the concept of pointer we have to understand arithmetic operations and expected result on our code.

a) Addition:-

i) addition of pointer with integral constant.

```
int arr[5] = {10, 20, 30, 40, 50};  
int *p = &arr;
```



```
printf("%d", p); // 100  
printf("%d", *p); // 10
```

```
p = p + 2;
```

```
p = 100 + 2;
```

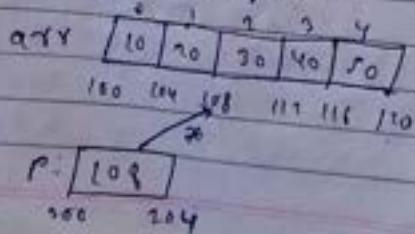
```
p = 100 + 2 * sizeof(int));
```

```
p = 100 + 2 * 4;
```

```
p = 100 + 8;
```

```
p = 108
```

modified diagram



printf("%d", p); // 108

printf("%d", *p); // 108

Some other
operations

multiple
assignment
is not done

constant
of

The above syntax uses the concept of addition of
pointer with integral Constant.

When we add any integral constant with pointer
then we have to multiply that constant with
Size of pointed type.

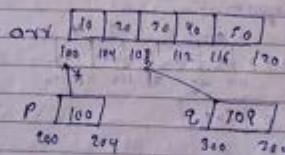
In our example pointed type is integer and
pointed datatype which is array is also integer.

ii) Addition of two pointers:-

Addition of two pointers is not allowed it
generates compilation error because after addition
it may generate some invalid addresses.

e.g.: int arr[5] = {10, 20, 30, 40, 50};
int *p = &(arr[0]);

int *q = &(arr[2]);



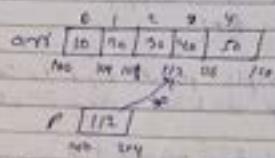
p + q; // Not allowed
compilation error.

2) Subtraction:-

is subtracting integral constant from pointer

e.g. int arr[5] = {10, 20, 30, 40, 50};

int *p = &arr[0];

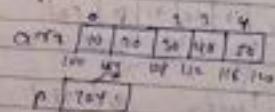


printf("%d", p); // 10
printf("%d", *p); // 10

$$P = 102 - 72 \\ P = 102 - 2 * sizeof(int);$$

$$P = 102 - 2 * 4 \\ P = 102 - 8 \\ P = 104$$

modified diagram:-



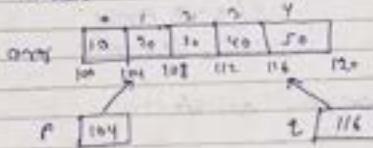
printf("%d", p); // 104
printf("%d", *p); // 104

int from pointer
26, 24, 20, 32

1);

ii) Subtract one pointer from another pointer.

Eg: int arr[5] = {10, 20, 30, 40, 50};
int *p = &arr[1];
int *q = &arr[4];



$$\begin{aligned} q - p &= \\ &= 116 - 104 = \frac{(116 - 104)}{(\text{size of int})} \\ &= \frac{12}{4} \\ &= 3 \end{aligned}$$

when we subtract two pointers then we have to divide it by size of pointer type.

if we want to perform subtraction of pointers then
1) both pointers should be same point
2) both pointers must point to same memory
which is related in nature

iii) multiplication :-

multiplication of two pointers and multiplication of integral constant not allowed

p * q ; // Compilation error
p * 2 ; // Compilation error

d) division:-

division of 2 pointers and dividing
pointer with integral Constant is not
allowed

$P/2; // \text{Error}$
 $P/2; // \text{Error}$

e) increment operator:-

we can apply increment operator ++ on
pointer

f) decrement operator:-

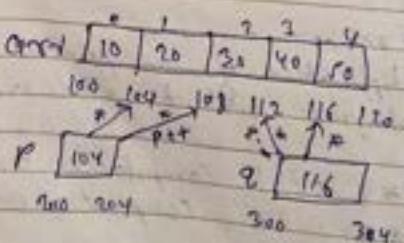
we can also apply decrement operator
-- on pointer

$++$ & $--$ operators internally not hand
operators they used addition and subtraction
technique indirectly

int arr[5] = {10, 20, 30, 40, 50};

int *P = &(arr[1]);

int *Q = arr[4];



$p++ ; //$

$p = p + 1;$
 $p = 104 + 1;$
 $p = 104 + 1 * \text{sizeof}(int);$
 $p = 104 + 4$
 $p = 104 + 4$
 $p = 108;$

$q-- ; //$

$q = q - 1;$
 $q = 116 - 1;$
 $q = 116 - 1 * \text{sizeof}(int);$
 $q = 116 - 1 * 4;$
 $q = 116 - 4;$
 $q = 112;$

1/8/15
*

void pointer - it is considered as generic pointers because that pointer can hold address of any datatype.

↳ meant void pointer can store address of integer, address of float, address of double, address of char.

↳ if pointed datatype is not known then we can use the concept of generic pointer i.e (void*).

e.g.: int no = 10;
void *p;
 $p = &no;$
 $\text{printf}("The value is %d", (\text{int} *)p);$

↳ according to above syntax at the time of dereferencing the content we have to use specific typecasting.

5) we can assign any type of pointer to void pointer directly.

eg: int no = 10;
int *iptr = &no;

char ch = 'A';
char *cptr = &ch;

void *p;

p = iptr; // allowed
p = cptr; // allowed

6) we can not apply any pointer arithmetic operations on void pointer

7) eg: void *p, *q;

p+3; // error
q-p; // error

* Memory allocation :-

There are 2 ways in which we can allocate memory

- 1) Static memory allocation
- 2) Dynamic memory allocation

Static memory allocation is also called compile time memory allocation or early memory allocation

Dynamic memory allocation is also called of runtime memory allocation or late memory allocation

17/15

to

In both the types of memory allocation memory is allocated in main memory i.e. RAM

There are some differences in static memory allocation and dynamic memory allocation such as

Static

1) memory is allocated at runtime, how much memory should be allocated is decided at compile time.

2) memory allocated inside stack or data

3) the default value of allocated memory is dependent storage class

4) there can be memory shortage & memory wastage

5) if there is no enough memory then our program is unable to execute.

6) there is no need to use memory special function for memory allocation or deallocation

7) there is no need to free the memory

dynamic

memory is allocated at runtime but how much memory should be allocated is decided at runtime

memory is allocated inside heap

default value is always zero

there can not be memory shortage & memory wastage

if there is no enough memory then our program may abort at runtime

we have to use some special function such as malloc(), calloc(), realloc(), new(), free(), delete

it is programmer's responsibility to free the allocated memory

8) we can not increase or decrease the allocated memory

we can increase or decrease the allocated memory by using realloc.

9) once memory get allocated statically it remains till in main memory till the lifetime of that storage class ends.

After the life of dynamic memory gets completed we can free that memory.

10) it requires less time to access the data.

it requires more time to access the data.

SP11S

memory allocation and deallocation function

i) malloc :- memory allocation

a) this is a basic function which is use to allocate memory dynamically.

this function accept one parameter which is no of bytes that we want to allocate and it returns the base address of allocated memory.

prototype

```
void *malloc (size_t no);
```

a) size_t is a typedef for unsigned int

```
typedef unsigned int size_t;
```

a) a malloc is a library function which internally calls brk() - system call and that brk is system call internally calls growing algorithm.

or decrement
memory by
2)

if our malloc function fails to allocate memory
it returns NULL otherwise it returns a valid
base address of allocated memory.

e.g. dynamic
completed
cc func

malloc → bres → memory

grossregion - algorithm
bres - system call

at time
state

malloc

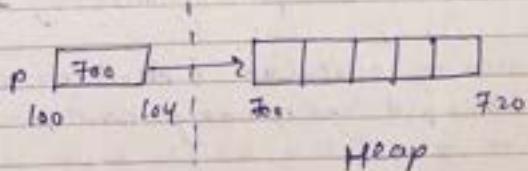
e.g. if we want to allocate memory for five integers
array then syntax should be

int *p = NULL;
p = (int*) malloc (5 * sizeof(int));

↑ no of elements
base address of allocated memory ↓ typecasting function name ↓ size of each element

is no
it
memory

diagram:-



and int

in internal
system

in above diagram memory for pointer p is allocated
in either stack or data section depend on its
storage class and the actual memory to stored
the data is allocated inside Heap

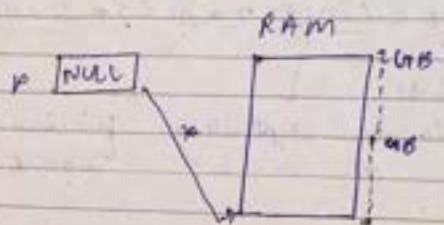
the below syntax indicate the static memory allocation
of same memory.

int arr[5];

If the above dynamic memory allocation fail it returns NULL and otherwise it returns valid Block allocated from Heap.

* Null pointer:-

- 1) null pointer is such pointer which stores zeroth address of our memory.
- 2) generally at this address there is source code of operating system.
- 3) if pointer is uninitialized then it may generate runtime accident to avoid that we have to initialize with it NULL.



* internal concept of memory management

- 1) memory allocation & deallocation is performed and managed by memory manager of o.s.
- 2) memory management is just a unit of operating system which internally maintains record of allocated memory and unallocated memory.
- 3) memory manager maintaining 2 data structures
 - 1) free list of memory.
 - 2) used list of memory.

6.8.2
with
n Heap

pointer
memory

u. code

garbage
box

rent

mixed and

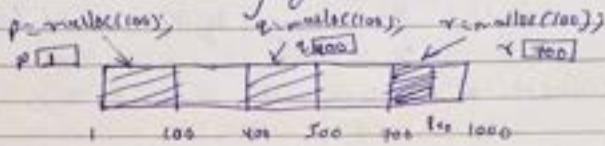
operating
ed of
memory

isolate

- used list of memory contains all the information about currently allocated memory while the free list of memory contains the information about such a memory which is unallocated

- 4) Both two Datastructure contains 2 columns
 - 1) Base address
 - 2) no of bytes

which are currently free or in used



U L M		F L M	
Base address	no of bytes	Base address	Size
1	100	100	300
400	100	500	200
700	100	800	200

- 5) when we call malloc(), calloc(), realloc() function entry gets removed from free list of memory and it's get added in used list of memory
- 6) when we call free() function entry gets removed from ULM and it gets added in free list of memory

** Calloc : Calculated allocation

- 1) The internal mechanism of malloc & calloc is pretty same
- 2) generally calloc function is used to allocate memory for array
- 3) This function is called calculated alloc because it performs internal calculation to decide the no of bytes should be allocated
- 4) void *calloc (size_t no, size_t size);
- 5) first parameter of calloc indicate no of elements and second parameter indicates size of each element
- 6) if we want to allocate memory for array of float which contains 10 elements then Syntax should be

```
float *p = NULL;
p = (float *)calloc (10, sizeof (float));
```
- 7) we can use the same logic by using malloc function as

```
p = (float *)malloc (10 * sizeof (float));
```
- 8) in old operating systems the memory which is allocated by malloc contains garbage and memory which is allocated by calloc contain zero in it

2) But in today's operating system there is no such difference.

3) realloc(): - Reallocation -

1) This function is generally used to increase the size of already allocated memory or decrease it if size

```
void *realloc(void *p, size_t size);
```

2) first parameter of realloc is base address of already allocated memory
2nd parameter is the new size that we want to allocate
that size can be greater than the previous or less than previous size.

e.g. int *p = NULL;
 $p = (\text{int} *)\text{malloc}(5 * \text{sizeof}(\text{int}));$
int *q = NULL;
 $q = (\text{int} *)\text{realloc}(p, 7 * \text{sizeof}(\text{int}));$

3) in above syntax we are using q as a new pointer which is used to store address of expanded memory

4) at this place we can use p pointer albeit if that expanded memory is not available contiguously then it returns NULL.

5) due to which we lost the base address of previously allocated memory. To avoid this problem we can take new pointer of q;

6) There are 2 scenarios in which we can use realloc()

- i) To increase the existing memory
- ii) To decrease the existing memory

3) if we want to increase already allocated memory then memory manager checks whether that much amount of memory is available immediately after that memory and if it is available then it returns the same base address.

4) But if it is not available then it allocates the memory sequentially in another memory slot and returns the new base address and frees previously allocated memory implicitly.

5) But all the scenarios are not applicable if we want decrease the existing memory.

6) Some extra behaviours in realloc:-

It generally realloc() function is used to reallocate existing memory but there are 2 scenarios in which realloc() function works in different manners:

a) if first parameter of realloc is NULL then it behaves as a malloc.

e.g.: int *p=NULL;
p=(int*)realloc(NULL, 40);

b) in above syntax new 40 bytes memory get allocated by and its base address gets copied inside p.

b) int *p=NULL;
p=(int*)malloc(40);

`P = (int *)malloc(4, 0);`

by using above code the syntax the allocated memory by malloc gets free.

** `free()`:

- 1) it is a good programming practice to free the dynamically allocated memory after its used gets completed.
- 2) if we hold that memory throughout the execution of program then it becomes memory leakage.
- 3) To free that physically allocated memory we have to use `free()` function.

`void free(void *p);`

1) return value of `free` is void but it accept one parameter which is the base address of existing memory which can be deallocated by `malloc`, `calloc`, `realloc`.

2) when we call `free()` function memory manager of that operating system used list of memory table to search for its base address.

3) if that base address is found then that entry gets removed from list of memory table and it gets added in free list of memory table.

4) if we pass base address to `free` function that address should not be modified address bcoz memory manager unable to search in used list of memory table.

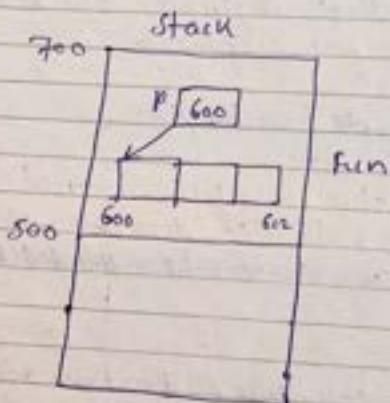
* `allocac` :- This function is same as `malloc`, but it is used to allocate memory inside the stack frame of function in which that `allocac` function called.

In case of `allocac` there is no need to free the memory explicitly because that memory is freed implicitly when our function terminates.

Void fun()

int *p = NULL;

p = (int *)allocac(3 * sizeof(int));



* function pointers:-

1) Function pointer is such a pointer which stores the address first instruction of our function.

2) When we create a function pointer we can store address of any function whose prototype gets matched with prototype of that function pointer.

at malloc()
empty initial
list that

ed to free
e memory is
minuted.

nt));

store
nation

fatigue

int add (int, int);
↓ ↓ ↓ ↓ ↓ ↓
int (*P) (int, int);

3) according to above syntax p is a pointer which
may store address of a such function which accepts
two parameters first is integer and 2nd is integer
and that function returns integer.

4) At this point only the pointer gets created but
it contains nothing.

5) As we want to have address of function then there
is no need to use address of operator because
name of the function is internally consider as
its base address.

P = add;
P(10, 20);

6) when we use this p pointed by putting 10 & 20
the function gets called implicitly where address
is stored in pointer p.

eg: int add (int no, int no); //10000
 {
 printf ("add");
 }
 int Sub (int no, int no); //20000
 {
 printf ("Sub");
 }
 int mult (int no, int no); //30000
 {
 printf ("mult");
 }

```
int main()
{
    int (*p)(int, int);
    p = add;
    p(10, 20); // 1000 add
    p = sub;
    p(30, 20); // 2000 sub
    p = mult;
    p(5, 2); // 3000 mult
    return 0;
}
```

7) the concept of function pointers used in libraries such as dynamic link library where we have to receive address of function due to which acceptor should be function pointer.

8) like normal function member memory for function pointer is always 4 bytes but it stores the address from text section of a process

* program which is used to demonstrate the pointers which may point to any section of a process (text, data, stack, heap).

```
int global1 = 10;
int global2;

void fun(int n)
{
    printf("fun");
}
```

```

int main()
{
    int local = 1;
    int *a = &global1;
    int *b = &global2;
    int *c = &local;
    int *x = (int*)malloc(9);
}

void (*y)(int**);
y = fun;

y
return 0;

```

Hello.exe



8/8/15 Dynamic memory allocation

* one dimensional array :-

```
#include <stdio.h>
#include <stdlib.h>
```

```
int main()
{
```

```
    int element = 0, i = 0;
    int *p = NULL;
```

```
    printf("Enter no of elements");
    scanf("%d", &element);
    // allocate memory
    p = (int *)malloc(element * sizeof(int));
    // if memory is not allocated
    if (p == NULL)
    {
        printf("Unable to allocate memory");
        return -1;
    }
```

```
// accept elements from user
    printf("Enter element");
    for (i = 0; i < element; i++)
    {
        scanf("%d", p + i);
    }
```

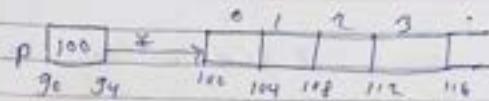
// print the element

```
for (i = 0; i < element; i++)
{
    printf("%d", p[i]);
    // p[i] = *(p + i)
```

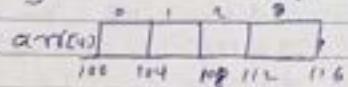
// free allocated memory

```
free(p);
```

}



Static memory allocation diagram :-



In above program we have to allocate memory dynamically for one dimensional array.

for that purpose we have to accept no of elements from the user & then allocate the memory by using malloc.

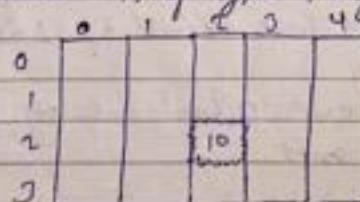
after use of that memory gets completed we have to free the allocated memory by using free function

* 2 Dimensional array:-

Consider the static memory allocation syntax as

arr[4][5]

according to above Syntax arr is a 2 dimensional array, which contains 4 one dimensional array each one dimensional array contains 5 elements each element is of type integer.

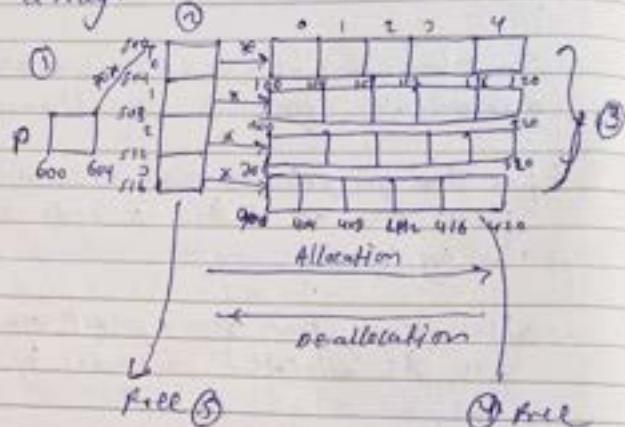


and $[2][2] = 10;$

According to above concept we can conclude that

2-D array is collection of one dimensional array.

As we know dynamic memory allocation for one dimensional we can use that same concept repeatedly to allocate memory for 2-D array.



According to above diagram for memory allocation we have to follow 3 steps as

Step 1: allocate memory for pointer p which is used to hold base address of array of pointers.

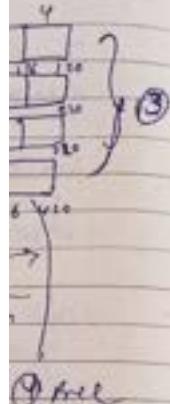
Step 2:- allocate memory for array of integer pointers where each element of that array holds the base address of 1-D array.

Step 3:- allocate actual memory to store our data elements.

To free that allocated memory which is allotted by using above 3 steps we have to follow reverse approach.

dimention

constant
same
memory



```
#include <stdio.h>
#include <stdlib.h>
```

```
int main ()
```

{

```
    int row=0, col=0, i=0, j=0;
```

Step 1: int *xp = NULL;

```
    printf ("Enter rows & cols ");
    scanf ("%d %d", &row, &col);
```

Step 2:

```
    int p = (int**) malloc (row * sizeof (int *));
```

Step 3:

```
    for (i=0; i<row; i++)
```

{

```
        p[i] = (int*) malloc (col * sizeof (int));
```

alloc

allocation

which
way

legal
way

not

allowable
weise

```
    printf ("Enter the elements ");
```

~~scanf~~

```
    for (i=0; i<row; i++)
```

{

```
        for (j=0; j<col; j++)
```

{

```
            scanf ("%d", &p[i][j]);
```

}

```
    printf ("Your entered elements are ");
```

```
    for (i=0; i<row; i++)
```

{

```
        for (j=0; j<col; j++)
```

{

```
            printf ("%d", p[i][j]);
```

}

// below code is used to free deallocated memory
allocated memory

Step 4:

```
for(i=0; i<n; i++)
    free(p[i]);
```

Step 5:

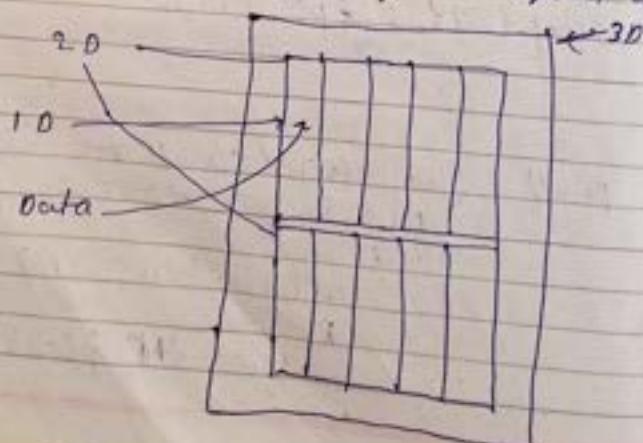
```
free(p);
```

9

* 3 - Dimensional array:-
consider the below syntax of

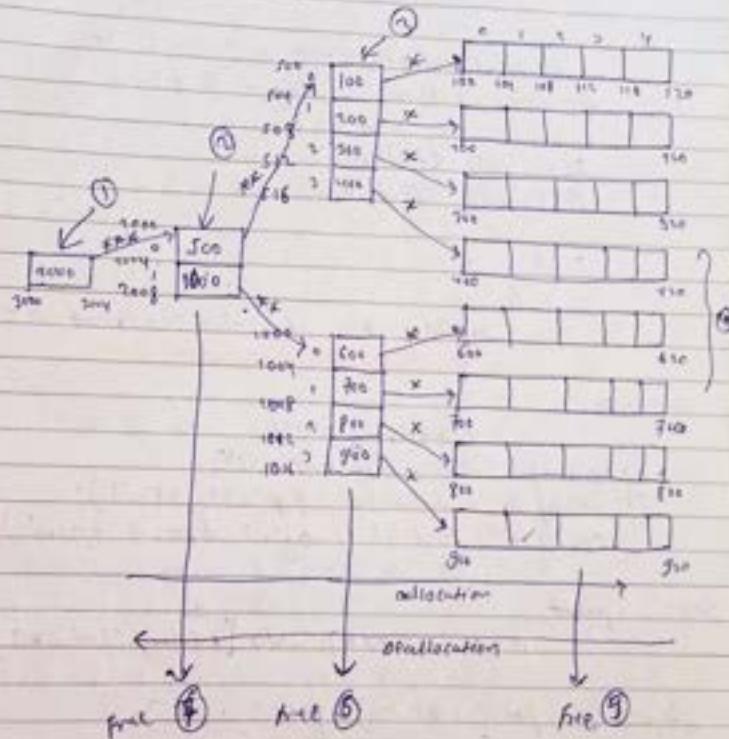
`int arr[2][4][5];`

This array is considered as 3 dimensional array which contains 2 two dimensional array. Each 2 dimensional array containing 4 one dimensional array, each one dimensional array containing 5 elements. Each element is of type integer. Below diagram is layout of above syntax of.



ed memory

if we want to allocate flat memory dynamically
then our memory representation should be



and
allocated
being
one
H

4 of
-30

```
int main()
{
    int first, second, third = 0;
    step1: int ***p = NULL;
    int arr[first][second][third]
```

memory allocation step - 1, 2, 3, 4

memory deallocation step - 5, 6, 7

```
#include <stdio.h>
#include <stdlib.h>
```

```
int main()
{
    int first=0, second=0, third=0;
    i=0, j=0, u=0;
    step1: int ***p = NULL;
    printf("Enter dimensions");
    // scanf("1d*fd*f", &p[i][j][u]);
    scanf("fd*f*fd", &first, &second, &third);
```

step2: print

```
P=(int**)malloc(sizeof(int**));
step3: for(i=0; i<first; i++)
    {
        p[i]=(int*)malloc(second*sizeof(int));
        for(j=0; j<second; j++)
            p[i][j]=(int*)malloc(third*sizeof(int));
```

step4: for(i=0; i<first; i++)
 {
 for(j=0; j<second; j++)
 p[i][j]=(int*)malloc(third*sizeof(int));

y

```
printf("Enter element %d");
for (i=0; i<first; i++)
{
    for (j=0; j<second; j++)
    {
        for (k=0; k<third; k++)
        {
            if ("Third", &p[i][j][k])
                y
        }
    }
}
```

// same logic used print the elements

steps:

```
for (i=0; i<first; i++)
{
    for (j=0; j<second; j++)
        free p[i][j];
}
```

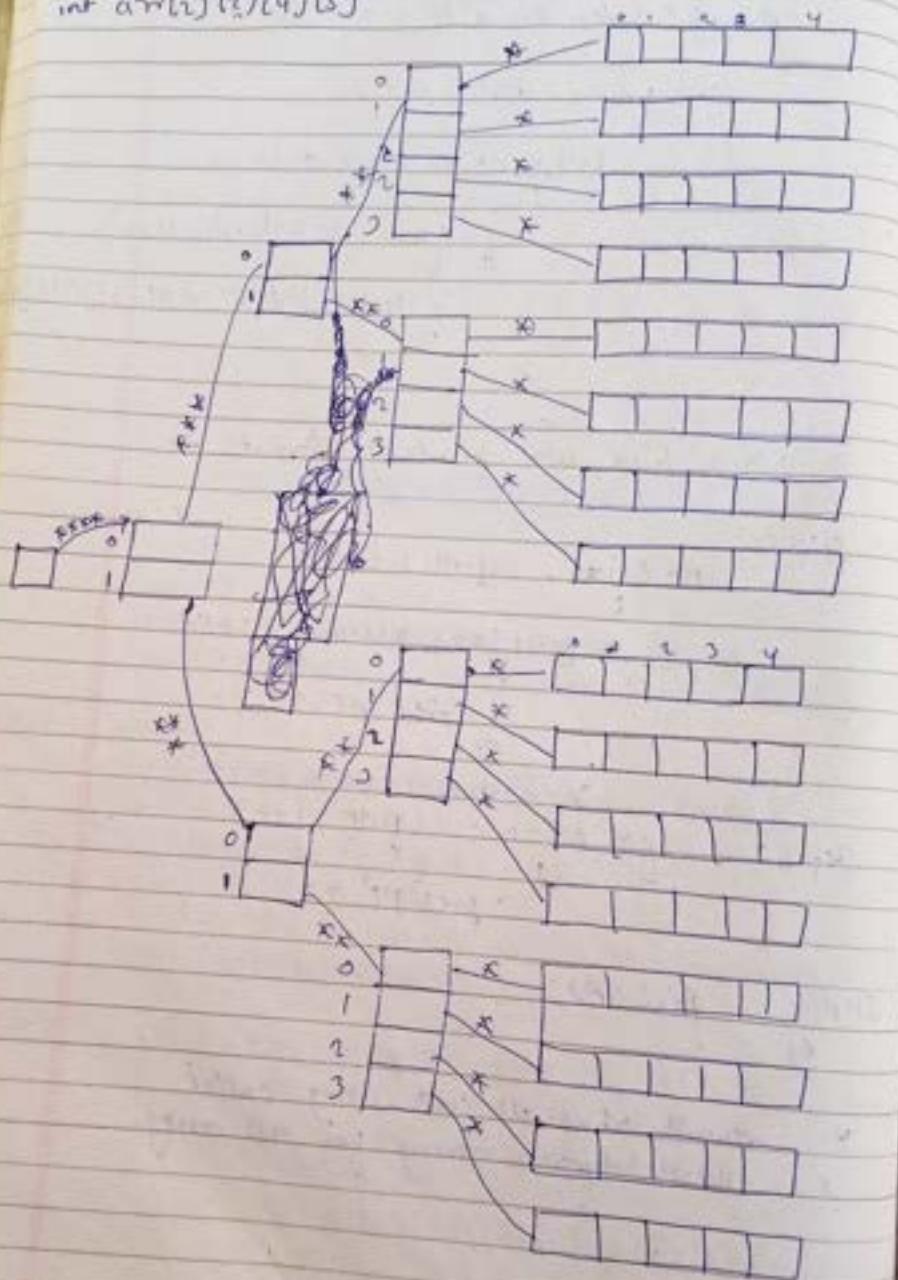
Step 6: for (i=0; i<first; i++)
{
 free p[i];
}

Step 7 free (p);

* rewrite above all code using calloc
* allocate dynamic memory for 40 array

rd + sizeof(int));

int arr[2][2][4][5]



```
#include <stdio.h>
#include <stdlib.h>

int main()
{
    int first=0, second=0, third=0, fourth=0,
        i=0, j=0, k=0, l=0;

    int *** p = NULL;
    printf("Enter dimensions");
    scanf("%d%d%d%d", &first, &second, &third, &fourth);

    p = (int *** )malloc(first * sizeof(int **));
    for(i=0; i<first; i++)
    {
        p[i] = (int **)malloc(second * sizeof(int *));
        for(j=0; j<second; j++)
        {
            p[i][j] = (int *)malloc(third * sizeof(int));
            for(k=0; k<third; k++)
                p[i][j][k] = (int *)malloc(fourth * sizeof(int));
        }
    }
}
```

y y

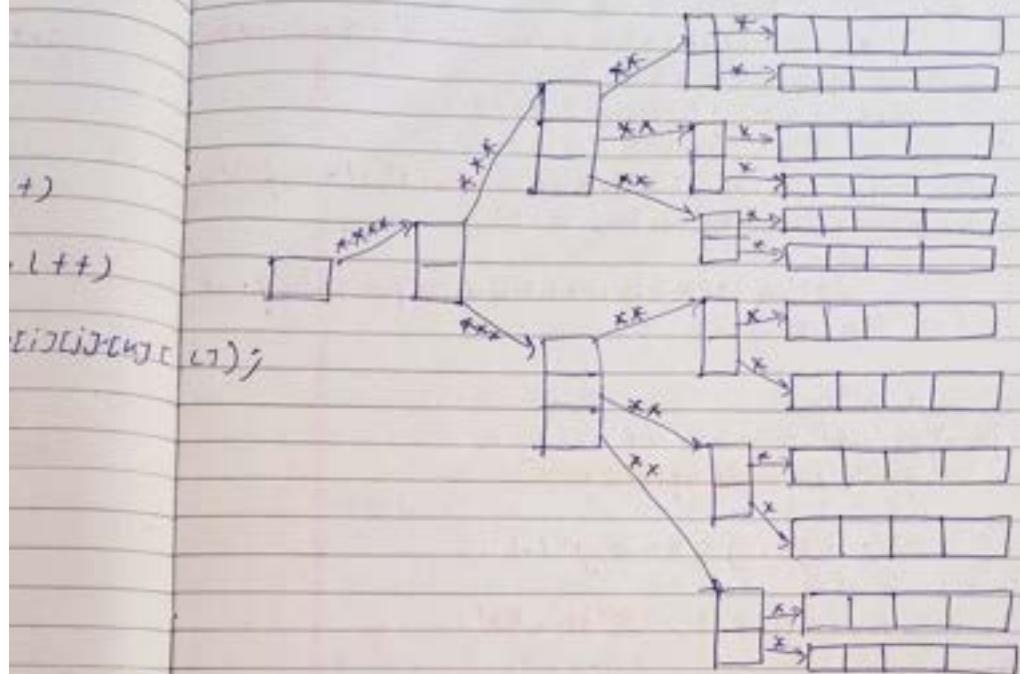
9

```
printf("Enter elements\n");
for(i=0; i<first; i++)
{
    for(j=0; j<second; j++)
    {
        for(k=0; k<third; k++)
        {
            for(l=0; l<fourth; l++)
            {
                printf("%d %d %d %d\n", i, j, k, l);
            }
        }
    }
}
```

y y y

Y/N

int arr[2][3][4]



```
#include <cselib.h>
#include <stlib.h>
```

```
int main()
{
    int first, second, third, four;
    int i, j, k, l;

    int ***p = NULL;
    printf("Enter dimension");
    scanf("%d %d %d %d", &first, &second, &third, &four);
    p = (int ***)malloc(first * sizeof(int **));
}
```

for(i=0, i<first, it++)

$f * l[i] = (\text{int} \ldots) \text{malloc}(\text{second} * \text{sizeof}(\text{int}))$

```
for(j=0; j < second; j++)
```

$p[i][j] = (\text{int} *)\text{malloc}(\text{third} * \text{size}(\text{int}))$

for (k=0; j < Third; k++)

$p[i][j][k] = (\text{int} *) \text{malloc}(\text{fourth_Size}[\text{is}] * \text{Size}[\text{is}])$

3

```
printf("enter elements");
```

```
for ( i=0; i<first; i++ )
```

```
for(j=0; j < 8 and; j++)
```

```
for (ix=0; ix<third; ix++)
```

$$\text{for } L = 0 : 1 \leq \alpha - \beta \leq 1/2$$

`Scoref("f-d", &P[i;j;j][i;j]);`

3 3

Effort memory

for i=0; i<first; i++)

free(p);
for(j=0; j<n; j++)

~~h free(pE[i][j]);~~

for (u=0; u<third; u++)

```
for (t=0; !l[found]; t++);
```

```
    free (p[i][j][u]);  
for (t=0; t<fourth; t++)
```

~ ~ Free (ptr[j][j][w][t])

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* Array and pointers *

1) One Dimensional array:-

Name of array is internally considered as Base address of that array.

when we access any element of array then that normal syntax gets converted into pointer syntax.

Consider below one dimensional array as

```
int arr[5] = {50, 40, 30, 20, 10};
```

when we create this kind of array symbol table contains its name address and scope.

Consider the below syntax as

```
arr[3];
```

can be converted as.

```
* (arr + 3);  
* (100 + 3);  
* (100 + 3 * sizeof(int));  
* (100 + 12);  
* (112);  
20
```

Above syntax indicated that when we access any element of the array it gets converted into the pointer syntax.

All the below syntax gives the same result as

$A[B]$	$Arr[3]$
$*(A + B)$	$* (Arr + 3)$
$* (B + A)$	$* (3 + Arr)$
$B[A]$	$3[Arr]$

\Rightarrow

9.8.15

* Array and pointers *

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* (100 + 3 * sizeof(int));  
* (100 + 12);  
* (112);  
20
```

above Syntax indicated that when we access any element of the array it gets converted into the pointers Syntax.

all the below Syntax gives the same result as

A[B];		Arr[3];	}
* (A + B)		* (Arr + 3);	
* (B + A);		* (3 + Arr);	
B[A];		3[Arr];	

20

* why array index starts from zero.

we can access every element of array by using its name and appropriate index.

Consider the below syntax of

```
int arr[4] = { 10, 20, 30, 40 };
```

arr	0	1	2	3
	10	20	30	40
	100	104	108	112

If we start array index from one instead of zero then our syntax to access the first element should be arr[1]:

This syntax gets converted into pointed format as

```
*(&arr + 1);  
*(arr + 1 * 4);  
*(100 + 4);  
*(104);
```

according to above syntax due to starting index one we can directly access the element which is at address 104.

The element at address 104 is not the first element of our first element stored it is 100. due to the starting index 1 we can not access the element at address 100

hence array index start from zero.

Two dimensional array

Consider the below syntax of a two-dimensional array of integers:

for this array which is multidimensional array we can draw its diagram in two ways as:-

1) Diagrammatic representation.

	0	1
Rows	0	10 20
	1	30 40
	2	50 60

← col →

2) memory layout / memory representation:-

*	*	*	*
10 20 30	40	50	60

There are 2 ways in which the contents of array get stored in memory

1) Row major storage

2) Column major storage

In case of row major storage contents of multi-dimensional array get stored in Row by Row format.

whereas, in column major representation contents get stored column by column format.

The storage type technique is depend on the programming language used.

C, C++ uses row-major storage mode and
Python uses column-major storage
(some of the version)

Accessing mechanism of the contents array
is irrespective of the storage mechanism.

Row-major:-

10	20	30	40	50	60
----	----	----	----	----	----

Column-major:-

10	30	50	20	40	60
----	----	----	----	----	----

Similarly as one dimensional array we can
also convert the normal syntax of accessing
the elements into the pointed syntax.

Consider if i want to access the element at

$\text{arr}[2][1]$;

The above normal syntax can be converted into
the pointed syntax of

$\text{arr}[1]\text{x} // (\text{arr}[1])[1] \rightarrow \text{x}[1]; \text{x}[1];$
 $\ast(\&(\text{arr}+2)+1)$

e

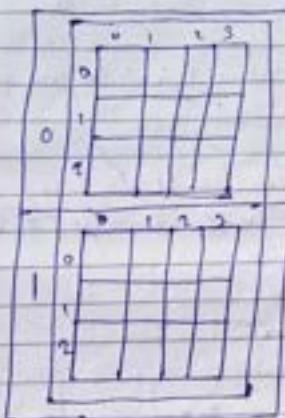
In general terms our arr.array syntax
can be converted into pointed syntax of

$\text{x}[A][B],$
 $\ast(\ast(A+\&)\&B)$
 $\ast(\ast(A+\&)+B)$
 $\ast(B+(\ast(A+\&))) ;$

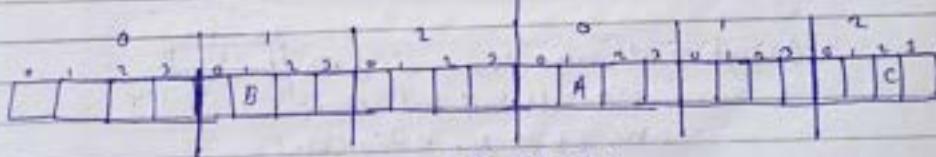
* 3 dimensional array:-

int arr[2][2][4];

Consider below 3 dimensional array. its memory representation should be



its memory representation should be



A // Arr[1][0][1];

B // Arr[0][1][1];

C // Arr[1][2][2];

H.W

problem on pointer 1

problem on C

problem on array

problem on Control instructions

problem on storage

* Constant and pointers:-

Constant is considered as datatype qualifier
in C & C++

If const keyword is used while declaring the
variable then we can not change the
contents of that variable.

Consider the below syntax as

int no = 10;
no = 11; // allowed
no++; // allowed
no--; // allowed

Const int no = 10;
no = 11; // not allowed
no++; // not allowed
no--; // not allowed

According to the concept of constant the
variable which is initialized at the time of
declaration can not be changed throughout
execution of program.

Like normal constant variable we can also
create constant array.

Constant array should be by using initialization list otherwise we can not initialize it.

Ex: `const int arr[4] = {10, 20, 30, 40};`

arr is one dimensional array which contains 4 elements and each element is of type constant integer.

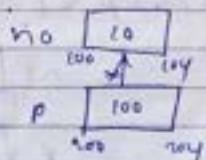


`arr[1]++;` // not allowed
`arr[2] = 15;` // not allowed

As the concept of constant is applied on any datatype due to which it is also applicable for pointers.

There are some scenarios to understand the concept of constant with pointers.

Scenario 1:- `int no = 10;`
`int *p = &no;`

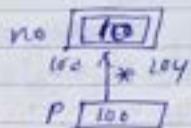


p is a pointer which holds address of integer
`no = 11;` // allowed
`no ++;` // allowed
`p = &no;` // allowed
`p++;` // allowed
`(*p) = 11;` // allowed

In this scenario pointed type is nonconstant
addressed pointer type is nonconstant

Scenario 2:-

```
const int no = 10;  
const int *p = &no;
```



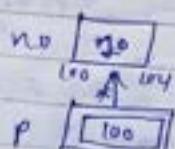
p is pointer which holds address of integer constant

no++; //not allowed
no +=; //not allowed
p = &x; //allowed
p++; //allowed
(pp) = 11; //not allowed

In this scenario pointed type is constant but
pointer type is nonconstant

Scenario 3:-

```
int no = 10;  
int *const p = &no;
```



p is ~~nonconstant~~ pointer which holds address of integers.

nd

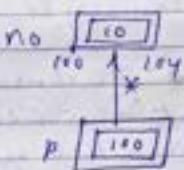
$no = 11; //$ allowed
 $no++;$ // allowed
 $p = &n;$ // not allowed
 $p++;$ // not allowed
 $(*p) = 11;$ // allowed

in this scenario pointed type is non-constant but
pointer type is constant

Scenario 4:-

```
const int no = 10;  
const int* const p = &no; // both some  
int const* const p = &no;
```

Ans



$no = 11;$ // not allowed
 $no++;$ // not allowed
 $p = &n;$ // not allowed
 $p++;$ // not allowed
 $(*p) = 11;$ // not allowed

in this scenario pointed type is constant as well as
pointer type is constant

* complicated statement reading in C

1) int no=10;

no is a variable of type int which is initialized to value 10.

2) int arr[4];

arr is a one dimensional array, which contains 4 elements, each element is of type integer.

3) struct Demo obj[5];

obj is a one dimensional array, which contains 5 elements, each element is of type struct demo.

4) int arr[3][5];

arr is a 2 dimensional array, which contains 3 one dimensional arrays, each one dimensional array contains 5 elements, each element is of type integer.

5) int arr[3][2][4];

arr is 3 dimensional array, which contains 3 two dimensional arrays, each 2 dimensional array contains 2 one dimensional arrays, each one dimensional array contains 4 elements, each element is type integer.

6) `int *p;`

initials

`p` is a pointer which can hold address of integer

7) `char *p;`

`p` is a pointer which can hold address of character

8) `void *p;`

`p` is a pointer which can hold address of any datatype

9) `char **p;`

`p` is a pointer which holds address of character pointer

* eg:

`char ch = 'A';`

`char *q = &ch;`

`char **p = &q;`

`ch [A]`

`100 ↗ p = 101`

`q [100]`

`104 ↗ pc = 108`

`p [104]`

`100 ↗ np`

10) `const char **p;`

`p` is a pointer which holds address of such a character pointer,

that pointer holds address of constant character

`const char ch = 'A';`

`const char *q = &ch;`

`const char **p = &q;`

`A`

`100 ↗ p = 101`

`a [100]`

`200 ↗ pc = 204`

`p [200]`

`300 ↗ np`

`ch++;` // X `np = &ch;` // X

`ch = 'B';` // X `np++;` // X

`*q = 'B';` // X `*np = ch2;` // X

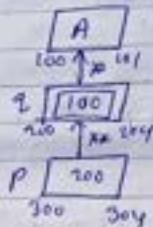
`(*q)++;` // X `(*np)++;` // X

15/8/15

1) char * const * p;

p is a pointer which holds
address of constant pointer and
that pointer holds address of character.

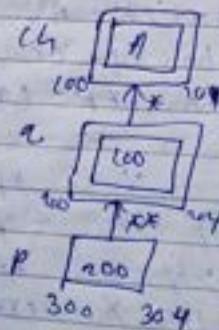
```
char ch = 'A';
char *const q = &ch;
char *const *p = &q;
```



2) const char * const * p;

p is a pointer which holds
address of constant pointer,
and that pointer holds address of character constant.

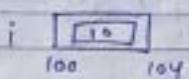
```
const char ch = 'A';
const char *const q = &ch;
const char *const *p = &q;
```



18/8/15

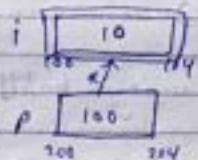
* const int i = 10;

i is a variable of type constant integer, which is currently initialize to 10.



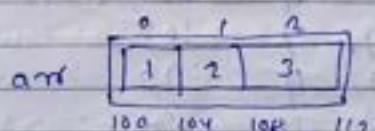
* Const int *p = &i;

p is a pointer which holds address of integer constant, and currently it holds address of variable i.



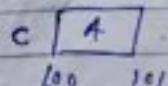
* const int arr[3] = { 1, 2, 3 };

arr is one dimensional array, which contains 3 elements, each element is of type constant integer, and that array is initialize to values 1, 2, 3.



* char c = 'A';

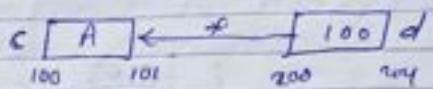
c is a variable of type character which is initialize to value A.



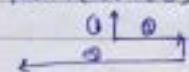
deep copy by value variable

* char *d = &c;

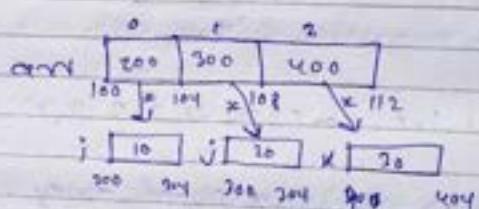
d is a pointer which holds address of character
and currently it holds address of variable c



* int *(arr[3]) = {&i, &j, &k}; i=10, j=20, k=30

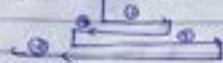


① arr is a one dimensional array, @ which contains 3 elements, @ each element is type of int *
currently that array holds address of variable i, j, k.



* arr;	//	100
*arr;	//	100
arr[i];	//	200
*arr[0];	//	108
*(arr[1]);	//	20
*(arr[3]);	//	30

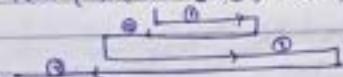
* int (*arr)[3];



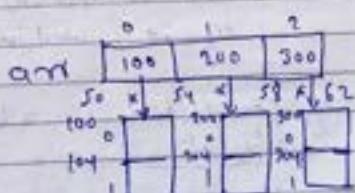
- ① arr is,
- ② pointed which holds address of
- ③ one dimensional array which contains 3 elements
- ④ each element is of type integer.



* int (*arr[2])[2];



- ① arr is one dimensional array which contain 3 elements
- ② each element is of type pointer
- ③ which holds address of one dimensional array which holds address of
- ④ one dimensional array which contain 2 elements
- ⑤ each element is of type integer.



* Rules of Statement reading by Greedy parsing algorithm

The rule which we are using to read the statement are internally used by compiler for tokenization activity.

Step

- 1) while reading start from identifier name.
- 2) traverse the statement on right hand side till we get closing circular bracket
- 3) again traverse that statement on left hand side till we get opening circular bracket
- 4) repeat steps 2 & step 3 till our statement ends

* void fun();

fun is a function which accept nothing and return nothing.

* void fun(c);

fun is a function which nothing and while returns nothing (if it is used in c)

fun is a function which accept any no of argument and it returns nothing
if it is used in C++

this type of convention is allowed in C++ only for function declaration.

* int fun(int, char);

fun is a function which accept two parameters

first parameter is value of integer
and second parameter is value of character
and that function return integer.

According to above program main is considered as a called function & fun is called function

function calling techniques

There are 3 types of function calling conventions available in C & C++ dependent on the contents of parameters such as

- ↳ call by value
 - ↳ call by address
 - ↳ call by reference

- if call by value : in this type function calling taking we have to pass the value to called function.

→ call by address : in this type of function calling technique we have to pass addresses of global or local variable . acceptor in calleee function should be pointers because we are sending the addresses

3) call by reference (only in C++):
↳ this concept is similar as

call by address concept in this concept we have to use concept of reference in C++ which is consider as another name to the variable.

* function returning mechanism:-

All three function calling mechanisms there are 3 ways in which we can return something from the function.

1) return of a value:-

In this returning mechanism we have to return some value which should be stored in some another variable of caller function.

2) return of a address:-

In this technique we can return address of any global or local variable. In this case accepted in called function should be pointer.

3) return as a reference (only in C++):-

In this technique we can return name of the variable which is accepted in called function in some another variable as reference variable.

* int fun (int, char)

In this function we are using concept of call by value and return by value technique.

* int * fun (int);

fun is a function which accept one parameter i.e. value of integer and that function returns address of a integer.

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```

    caller           called
    int main()
    {
        int *p = NULL;
        int q = 111;
        p = func(q);
        pf("d", *p);
        return 6;
    }

    int func(int x)
    {
        int y;
        y = x + 1;
        return y;
    }

```

above code uses concept of call by value and return by address technique

In above programs no variable is taken as global variable because we want to return address of that variable.

if we take that variable as a local variable while returning that variable scope of that variable gets terminated.

* int func(int *, char *);
fun is a function which accepts two parameters
first is address of integers, and is address of characters
and that function returns integer

all by

1-6

```
callcc  
int main ()  
{  
    int no = 11, x = 0;  
    char ch = 'A';  
    x = fun (&no, &ch);  
    return x;  
}  
  
int fun (int *p, char **q)  
{  
    *q = 'B';  
    *p = 22;  
    return 10;  
}
```

in above program we are using call by address technique and return by value technique

* int fun(int a[]);

fun is a function which accept one parameter
ie base address of array or
address of array of integers.

int fun (int *a);

this syntax is also same as above syntax

int main ()

{

int arr[3] = {10, 20, 3};

fun(arr);

return 0;

}

* int **fun();

fun is a function which accept nothing and it
return address of integer pointer

eg: int main()

{ int no=1;

int *p = &no;

int **func()

{

 return &p;

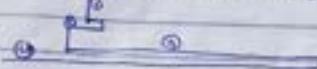
}

 return no;

}

address
e

* int (*p)(int, char);

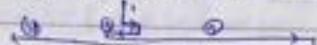


parameters

- (1) p is
- (2) pointer which holds address of
- (3) such a function which accept two parameters
- i.e. first is value of integer, 2nd is value of character
- (4) and that function returns integer.

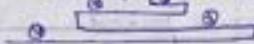
ex

* int (*p)(int **);



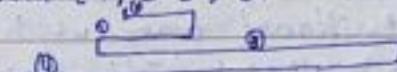
if

* int (*daytab)[13]; // int *a[13];



- (1) daytab is
- (2) a pointer which holds address of
- (3) one dimensional array which contains 13 elements
- (4) and each element is of type integer.

* void (*p[3])(int, int);



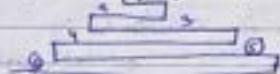
- (1) p is a one dimensional array which contain 3 elements and each element is of type pointer
- pointer which holds address of such function each function accept two parameters and function return nothing.

* above program is used to create array of function pointers

```
void fun (int x, int y);  
void func (int x, int y);  
void sum (int x, int y);  
void (*F[3])(int, int) = {fun, func, sum};  
F[1](10, 20); // call func
```

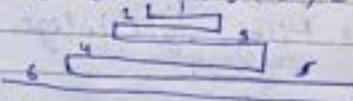
```
for (int i=0; i<=2; i++)  
    F[i](10, 20); // call fun, func, sum
```

* `char (*(*x[3]))();`



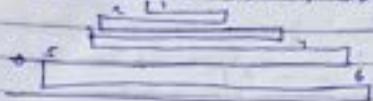
- (1) X is a function which accept nothing
- (2) which returns pointer to
- (3) one dimensional array which contains 3 elements
- (4) each element is of type pointer
- (5) which holds address of function which accept nothing
- (6) and returns characters.

* `char (*(*x[3]))()[5]`



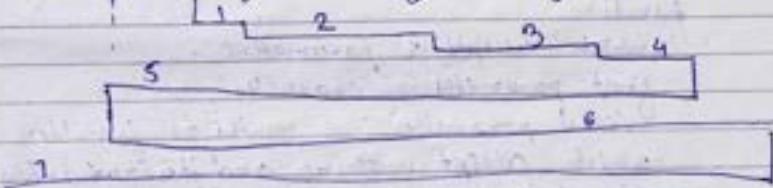
- (1) X is one dimensional array which contain 3 elements (2) each element is of type pointer which holds address
- (3) such a function which accept nothing
- (4) and which returns pointer to
- (5) one dimensional array which contain 5 elements (6) and each element is of type character.

* int * (* arr[5])(());



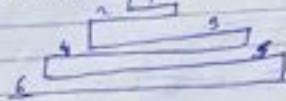
- Ques 3
Ans
- 1) arr is one dimensional array which contains 5 elements
 - 2) and each element is of type pointer
 - 3) which holds address of such a function which accept nothing
 - 4) that function returns pointer to
 - 5) such function which accept one parameter
 - 6) i.e. int
 - 7) that function returns address of integer

* void (*bsd_signal(int sig,void (*fun)(int)))(int);



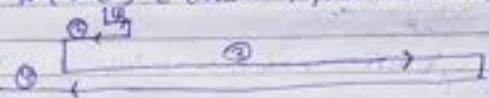
- Ques 4
Ans
- (1) ~~bosd~~ bsd_signal is a function which accept a parameter
 - (2) its first parameter is integer,
 - (3) and 2nd parameter is pointer to function which accept one parameter as a integer and it returns nothing
 - (4) and bsd_signal function returns pointer to
 - (5) such a function which accept one parameter i.e. integer
 - (6) and that function returns nothing

* float (* b[5])(void);



b is a function which accept nothing which returns pointed to an array which contains 5 elements and each element is of type pointer which holds address of such fun which accept nothing and return float

* void * (*c) (char, int());



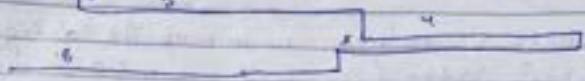
c is a pointer which holds address of such function which accepts 2 parameters first parameter is character second parameter is a pointer to function which accept nothing and it returns integer and whole function returns address of any datatype.

* void * (*a[5])(char * const, char * const);



a is one dimensional array which contains 5 elements each element is of type pointer, which holds address of such a function, which accept two parameters first parameter is constant pointer which holds address of character, and parameter is also constant pointer which holds address of character and whole fun returning address of any datatype

* `Void * (xp) (int, char **(x)(char *, char **));`



P is a pointer
which holds address of such a function
which accepts two parameters
first parameter is integer
and point parameter is pointer to such a function
which internally accepts two parameters
where first parameter is address of character,
and parameter is address of character pointer of
that function returns address of character pointer
and whole function return address of any datatype

16/8/15

* user-defined Datatypes in C.

- 1) user-defined datatype is one of the datatype in C which contains contents are depend on user of the programming language i.e. programmer.
- 2) depend on the programmatic requirement we have to design our user-defined datatype.
- 3) There are 3 user-defined datatypes in C such as:
 - i) structure.
 - ii) union
 - iii) enumeration.
- 4) Actual contents of structure, union & enumeration are directly depend on the programmer.
- 5) Before using any of the datatype first we have to think about its content and then we have to declare that appropriate datatype.
- 6) if we create single variable of any datatype then it is manageable but if we create multiple such variable then it is not manageable to avoid this problem we can use derived datatype i.e. array.
- 7) The major restriction on array is the homogenous elements of that array.
- 8) means in array we can store elements of same datatype.
- 9) If we want to store related element together under one name then we have to use any of the user-defined datatype.

say if we also want to store marks of 10 student
then we create array of integers

If we want to store division of 10 student then
we have to create array of characters

If we want to store birth year of student then we
have to create array of integers

All above practical approaches are correct but we
can not map this 3 different datastructure together

means we can retrieve division of student, marks of
student, birth year of student simultaneously
but if we want to perform that activity then we have
to create structure which contains name of student,
marks of student, it's division etc.

The above technical example indicates need of userdefined
datatype.

Some logical concept of userdefined datatype is used
in C++ which is named as class.

* important points about structure and union

structure - structure is userdefined datatype which
contains all related things under single name

Structure can contain homogenous or heterogeneous
elements means structure can contains elements of
different datatypes also.

- 1) Structure can contain any of the primitive datatype.

eg: struct demo

{

int i;

char c;

float f;

double d;

}

in above syntax struct is keyword. demo is name of
that structure. i, c, f, d are members of that
structure and ; after end closing } curly bracket
indicate end of structure declaration.

2) Structure can contain any derived datatype instead
of array, pointer etc.

eg: struct demo

{

int arr[5];

char *p;

}

3) Structure can contain any user-defined datatype of
its member means it contain structure &
union inside it

struct outer

{

int i;

struct inner

{

int j;

}

}

~~eg. str~~
in above example inner named structure is considered as an inner structure.

Before using any of the structure first we have to declare that structure.

Eg: struct demo

```
    {  
        int i;  
        float f;  
    };
```

when we declare the structure memory is not allocated for that structure members

at the time of structure declaration new entry gets inserted inside symbol table by the compiler.

that symbol ^{entry} contains name of the structure and no of bytes which are required for that structure.

of
&

* Actual memory for the structure will be get allocated when create object or variable of that structure.

```
struct demo obj1, obj2;
```

obj1, obj2 are objects of struct demo.

Q) we can create object of variable of a structure at the time of declaration itself.

Eg: struct demo

```
    {  
        int i;  
        float f;  
        obj;  
    };
```

6) we can not initialize any member of structure at the time of declaration because at that point memory is not allocated

eg: struct Demo

```
int i = 10; // error  
{  
    j;
```

7) if we try to initialize that member it generates compilation error

8) we can initialize members of structure after creating variable of that structure.

9) we can initialize the member by using member by member technique. Such as

eg: struct Demo

```
{  
    int i;  
    int j;  
}  
struct Demo obj;  
obj.i = 10;  
obj.j = 20;
```

10) we can also initialize content of structure by using initialization list.

eg: struct Demo

```
{  
    int i;  
    char ch;  
}  
struct Demo obj = {10, 'a'};  
struct Demo obj = { .ch = 'a' } correct; // dependent
```

Structure
that

- 11) we can use initialization list at the time of declaration also

struct demo

```
int i;  
char ch;  
obj = 10, 'a';
```

creates

if we use initialization list for initialization purpose
we have to initialize all the members in sequence
in which they are declared.

12)

- if we want to initialize specific elements in initialization list the same compiled below this syntax of:

eg:

struct demo

```
int i;  
char ch;  
y;  
struct demo obj = { .ch = 'a' };
```

- 13) Every name of member of a structure should be different to avoid ambiguity while accessing the members.

if names are same then it generates compile time error

eg: struct Demo

```
{  
    int i; // compile time error  
    char i;  
};
```

dent

14) in same program we can use some identifiers
name as well as local or global variable
as well as structure member

eg:

```
int i = 11;  
struct Demo  
{  
    int i = 10;  
};  
obj = 4103;
```

what code is allow according to every compiler
because the variable name of i inside the
structure is not directly accessible by its name
but it is accessible by its structures variable
name.

15) like normal statement reading ^{practise}, we can also
read structure declaration at

eg: struct Demo
{
 int i;
 char ch;
};
struct Demo obj;

obj is a variable of type Demo, which is of user
defined datatype i.e structure, which contains
two elements as first is integer and 2nd
is character.

16) As array is derived datatype we can create array
of user defined datatype i.e structure or
Union.

variables

e.g: struct Demo

```
{  
    int i;  
    int j;  
};  
struct Demo arr[3];
```

arr is one dimensional array which contains 3 elements, each element is of type Demo, where Demo is a structure which internally contains two elements where 1st element is integer and 2nd is also integer.

- 17) we can initialize above array in multiple ways such as

- i) struct Demo arr[3] = {{10,20},{30,40},{50,60}};
- ii) arr[0].i = 10, arr[0].j = 20;
- iii) arr[1].i = 30, arr[1].j = 40;
- iv) arr[2].i = 50, arr[2].j = 60;

- 18) if members of structure are not initialized then it contains it's default values according to the storage class of structure variable not structure declaration.

e.g: void fun()

```
{  
    struct Demo  
    {  
        int i;  
        int j;  
    };  
    struct Demo obj;  
    printf("%d", obj.i); // garbage value  
}
```

struct Demo

{
 int i;

 obj;}

struct Demo

int main()

{

 ff("Id", obj.i); //o

 return 0;

}

19) if we have structure variable to access its members
then we have to use . dot (direct accessing
operator)

e.g: struct Demo

{

 int i,j;

};

struct Demo obj;

obj.i=10;

obj.j=20;

20) if we have structure pointed which
holds address of structure object
then to access its member
we have to use → operator
for (indirect access operator)

eg. Struct Demo

```
{  
    int i;  
    int j;  
}
```

Struct Demo obj = {10, 20};

Struct Demo *ptr;
ptr = &obj;

printf("%d %d", ptr->i, ptr->j);

abcs
etting

- 1) we can use assignment operator to assign contents of ~~from~~ one structure variable into other structure variable.

But this copy of one structure variable to other structure variable is shallow copy.

eg: Struct Demo

```
{  
    int i;  
    j;
```

Struct Demo obj1, obj2;

obj1.i = 11;

obj2 = obj1;

11 ←

obj1 = obj1

printf("%d", obj2.i); 11 11

- 2) we can not use any comparison operator, any relational operator or any equality operator to compare two variables of same structure.
means if we use "==" operator to compare two variables it generates compilation error.

```

eg: struct Demo
{
    int i,j;
};

struct Demo obj1={10,20};
struct Demo obj2={10,20};

if (obj1==obj2) // compilation error
{
    printf("equal");
}
else
{
    printf("not equal");
}

```

23) as equality operator is not applicable for structure variable then we can apply equality operator on every member of that structure and then we can compare that structure successfully.

```

eg: if ((obj1.i==obj2.i)&&(obj1.j==obj2.j))
{
    printf("equal");
}
else
{
    printf("not equal");
}

```

24) allocate memory for structure of variable in two ways:

- static allocation
- dynamic allocation

eg. Struct Demo

{

 int i, j;

y;

struct Demo obj; // static memory allocation

struct Demo *p = (struct Demo *) malloc

(sizeof(struct Demo));

// dynamic memory allocation

Ques) we can create structure inside structure in two ways

1) struct Demo

{

 int i;

 y;

 struct outer

{

 int i;

 int j;

 struct Demo var;

 } obj;

2) struct outer

{

 int i, j;

 struct Demo

{

 int i;

 y var;

} obj;

in both the above scenarios we are using concept
of nested structure.

But in first scenario our inner structure i.e
Demo is declared independently outside the
outer structure due to which we can also create

object of ~~not~~ Demo structure independently

but in 2nd scenario Demo structure is declared inside the outer structure due to which we can not create object of Demo structure independently.

- 26) if we are using the concept of nested structure in which the inner structure is declared inside outer structure then it is necessary to create variable of inner structure immediately after the memory for inner structure is not allocated

eg: struct outer
 { int i, j;
 struct inner
 {
 int x, y;
 };
 yobj;

in this scenario size of obj is 8 bytes only because variable of inner structure is not created inside the outer structure.

- 27) we can use same identifier name in case of nested structure

eg: struct outer
 { int i, j;
 struct inner
 { int i;
 y var;
 };
 y obj;

in this eg. member of outer structure i can be accessed
obj.i whereas member of inner structure i can be accessed
obj.var.i.

78) we can not create variable of same structure inside that structure.

Eg: struct demo

```
int i;  
struct demo obj; // error
```

y;

about syntax generated compilation error because when compiler try to create variable obj inside structure Demo it generates error.

Error is generated because at that point compiler is unable to find out sizeof structure demo because it not yet completely declared.

79) we can not variable of same structure inside structure but we can create pointer ^{create} of some structure inside that structure.

because Compiler easily predict size of pointer as it is always 4 bytes.

Eg: struct demo

h

```
int i,j;  
struct demo *p;
```

y;

// allowed

This type of structure is called as self-referential structure.

18 8 15

* unnamed structure :-

- 1) unnamed structure is a structure which is considered as normal structure but without its name if we program contains unnamed structure then we have to create variable of that structure at the time declaration itself.

ex: struct

```
{  
    int i, j;  
}
```

- 2) in this example obj is considered as unnamed structure.

- 3) after this structure declaration we can not create any variable that structure because structure identifier is missing.

- 4) in case of unnamed structure it's symbol table entry is not created by the compiler.

* Similarly as unnamed structure we can create named union also.

Union is considered as an undefined datatype which contains any primitive datatype, any user defined datatype, any derived datatype.

Syntactical way of writing union & structure both are same but internal memory allocation of structure & union is different.

ex: union demo

```
h  
{  
    int i;  
    char ch;  
}
```

18.8.15

function which is
without its name
should have
structure at the

in this Union there are 2 members of integer
and character. according to its size integer
requires largest memory according to its size but
to which only 4 bytes get allocated for our
variable obj.

5) when we create variable of union memory for
only few largest datatype get allocated due to which
we can only initialise element one element of
at a time

6) if we initialise multiple elements then that's
union contains the value which is initialised
at the end

e.g.:
obj.i = 97;
obj.ch = 'b';

according to above syntax character 'b' is the last
value in our Union due to which the previous value
i.e. 97 gets override with the ASCII value of 'b'

7) if we set only single value in union and if we
try to retrieve by using some different member
name of union then it gives same value
which is initialised previously

e.g. obj.i = 97;

printf("%c", obj.ch); // a

above printf generate output of a because 97
is the ASCII value of character a.

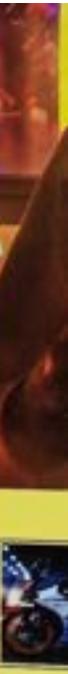
1st announcement
we can not
become struct

1-symbol table
file

can create

real datatype
int, any user
datatype.

there both are
declaration of



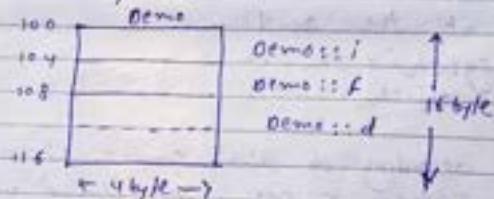
After designing any structure we can draw
it's diagram in two ways as
i) Diagrammatic representation
ii) memory layout

e.g. struct Demo

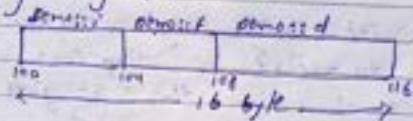
```
int i;  
float f;  
char ch;
```

3 obj: 1 byte

i) Diagrammatic representation



ii) memory layout



* Memory allocation strategy for structure or union

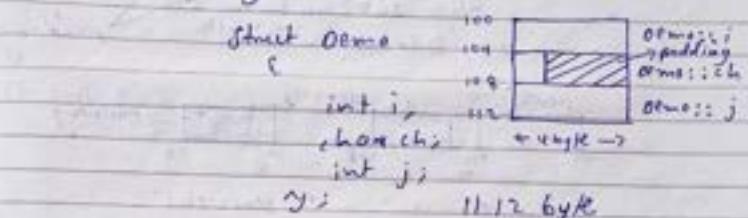
i) memory allocation strategy is purely depend on type of compiler and memory management strategies used by operating system.

ii) generally on 32 bit compiler and 32 bit o.s every member memory of structure and union gets allocated if word size format.

we can draw
as

- 3) WORD is concept which used in Assembly lang. which indicates 4 byte size.
 - 4) means every element of structure gets it's memory in form of an word size.
 - 5) due to this concept accessing mechanism of Structure or Union becomes faster.
- * Consider the below structure we explain the concept of padding

↑
16 byte
↓



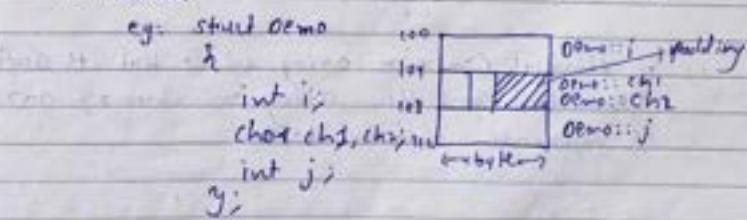
padding:- padding is the memory allocation which are allocated due to memory management strategies but that memory locations are not accessible to the program.

we use union

depend on
agent

2 bit o.s.
and unions

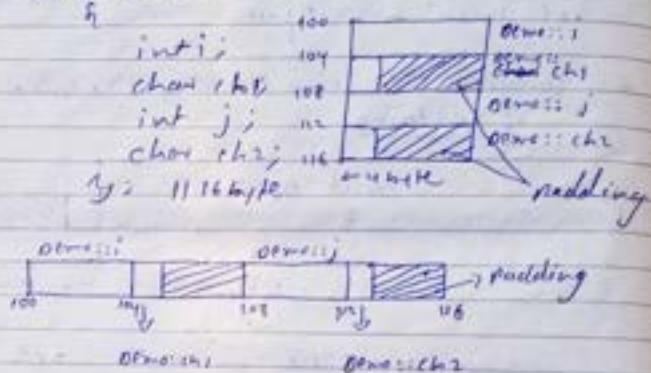
Generally if our structure contains character as a determinable then the concept of padding gets introduced.



In this structure 2nd member is character due to which now 4 byte memory get allocated allocated that character for that 4 bytes 1st one byte is useful for that character

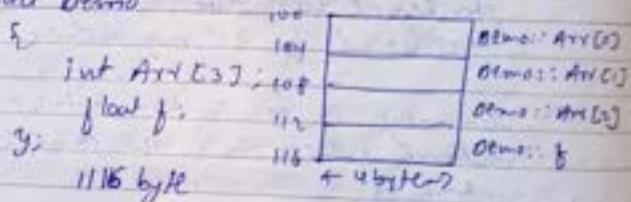
offer that char there is another character named as chr & this new char require one byte of memory we have 2 byte of memory free so to which from that three bytes 1 byte get allocated for that char & remaining 2 byte padding.

e.g. struct Demo



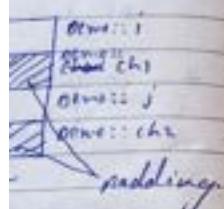
* Structure contain array:-

struct Demo



- 1) if Structure Contain array & it has its data member consider array member of array is separate data members.

Not character
does require one
memory free due
to 1 byte self
using 2 byte



→ padding
16

12

11

10

9

8

7

6

5

4

3

2

1

0

not its data
w of array

eg: struct Demo
{

int i;

int j;

char obj[3];

obj[0]	obj[1]	obj[2]	obj[i]
0	1	2	obj[i]
3	4	5	obj[i]

Sizeof(obj) = 12 byte

Sizeof(obj[0]) = 1 byte

Sizeof(obj[i]) = 1 byte

Sizeof(obj[0..i]) = 4 byte

eg: struct Demo
{

int arr[2];

int j;

y; obj; // 2 byte

104		obj::arr[0]
105		obj::arr[1]
106		obj::arr[2]
107		obj::arr[3]
108		obj::arr[4]
109		obj::arr[5]
110		obj::j
111		

- a) In this structure a dimensional array, it uses the concept now mentioned representation we have to apply that representation technique while drawing the layout of that structure.

eg: struct Demo
{

char ch1, ch2, ch3; int j;

y; // 2 byte

char ch1	char ch2	char ch3	padding
101	102	103	104 105 106

eg: struct Demo

{ char ch1, ch2, ch3; }

y; // 2 byte.

char ch1	char ch2	char ch3
101	102	103

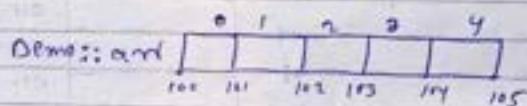
Size of structure may vary if we change the
compiled on the basis of it and written by standard memory
management strategy

- 3) If structure contain all characters of its
data members then the rule of word
size is not applicable for that structure.
- 4) According to this concept above structure
requires 3 bytes of memory without any padding
in eg:-

Eg:- Struct Demo

```
char arr[5];  
} // Structure
```

- 5) To this a structure like above rule is applicable
because it contains array of characters.



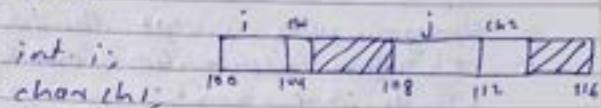
- 6) According to above all the examples which are
concern with size of structure there is a
concept of padding.
- 7) As every structure containing the padding means
our every structure variable waste our memory
- 8) Due to this concept of padding it may generate
some problems on embedded platform where
memory is less.
- 9) To avoid this drawback of padding we have to
use custom ^{padding} ~~padding~~ technique.

~~eg - #pr~~

- 10 To apply this concept of packing we have to use preprocessor directive named as `#pragma pack`

eg. struct demo

```
{  
    int i;  
    char ch1;  
    int j;  
    char ch2;  
};
```

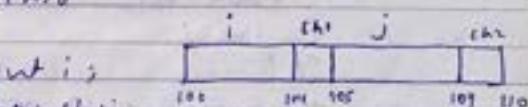


- 11 This structure requires 16 bytes memory from that 16 bytes 6 bytes memory gets wasted for padding means actual 10 bytes memory is useful for storing a date to avoid this wastage we can rewrite our structure as

`#pragma pack(1)`

struct demo

```
{  
    int i;  
    char ch1;  
    int j;  
    char ch2;  
};
```



- 12 According to this diagram concept size of our structure is 10 bytes

- 13 when we write `#pragma pack` in our structure request of 1 byte packing is applicable to only structure which is declared after that prototype.

14 if we want to stop packing request we have to write syntax at the end of declaration as -
#pragma pack()

15 #pragma pack concept should be use before structure declarations otherwise there is no use of that concept

16 if we use #pragma pack at the time of creating a variable of that structure then its concept is not applicable because how much memory shall be allocated for that structure is already decided at the time declaration itself.

e.g. struct demo

{

```
int i;  
char ch;  
char ch2;  
int j;
```

}

#pragma pack(1) ← not applicable

```
struct Demo obj;
```

```
sizeof(obj); // 12 byte
```

17 we can provide any parameter to #pragma pack of 1 or 2 or 4 or 8 or 16.

18 if we provide some other values often compiler ignore that parameter and use it's default packing concept of 4 bytes.

we have
declaration

19 The concept of packing is also applicable for union also.

* Binary Numbering System and its Concept

before
it no

of creating
concept
many should
only

17 Binary numbering system is such a numbering system which is used internally to store data in memory.

2) To understand the concept of binary numbering system we have to consider its base and range. Base of binary numbering system is 2 and range is 0 to 1.

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0	2^9	2^{10}
256	128	64	32	16	8	4	2	-1	512 1024

Consider integer variable no, which declared as
int no = 22;

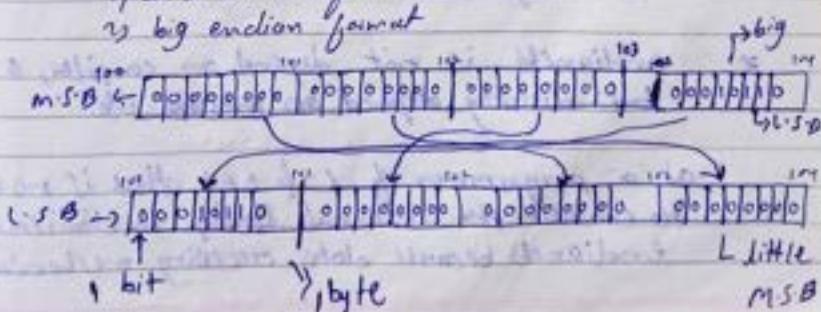
no 22
102 104

ble

To store that no in memory it require 4 bytes according to storage mechanism of int datatype there are 2 ways in which we can store our data in memory of.

pack of

LLC ignore
packing



* L.S.B - least significant Bit

it is a bit having least magnitude.

* least significant Byte:-

If if byte contains least significant bit then it is called as least significant byte

* m.s.B - most significant Bit

it is a bit containing highest magnitude

* most significant Byte:-

If byte contains most significant bit then it is called as most significant byte

* little endianess:-

It is an endianess in which L.S.B is stored at lower address and m.s.b is stored at higher address.

e.g. intel processor family, some of the processor family uses little endianess.

* Big endianess :-

It is an endianess in which m.s.b is stored at lower address and L.S.B is stored at higher address.

e.g. arm processor

and some of the processor arm family uses Big endianess.

* endianess is not depend on compiler, o.s but directly depend on processor use

as a programmer of C/C++ there is no need to consider the internal storage mechanism of endianess because data accessing mechanism

of any of the endianess is exactly same

If we interact with Assembly language then Endianess important means endianess is just way of storing the data internally.

* Bitfield in Structure or Union

1) According to previous topic of Program pack we can save the memory which got wasted due to the padding.

↳ Consider the below layout of a structure which is used to store date, month & year.

struct calendar

{

 int date;

 int month;

 int year;

};

3) As this structure contains 3 integer variable it requires 12 bytes of memory i.e. 96 bits.

4) As this structure is used to store date the range of every data element should be fixed.

5) Our date variable can store values of 1 to 31
month variable stores values of 1 to 12
year variable stores values of 2000 to 2040
as a range is fixed instead of taking 4 byte
i.e. 32 bit for every variable we can only
take required no of bits.

7) to store date are 6 bit enough
to store month 5 bit are required we can
store year values 13 bit enough

8) To apply this concept our modified syntax should be

eg. `#pragma pack(1)`

struct calendar

```
int date:6;  
int month:5;  
int year:13;
```

when we calculate size of structure we get 3 byte
means we can save 9 byte.

* Some important points about bitfield

1) we can not set negative or zero value as a bitfield

2) we can not set bitfield to float or double

3) we can not set bitfield to an array

4) we can not set bitfield value which is greater
than size of that datatype

Let's see some examples

```
int i:0; -① // error  
int j:-3; -② // error  
float f:6; // error - ③  
double d:7; // error - ④  
int arr[7]:4; // error - ⑤
```

eg: int n:35; // error -①

all the above syntax generates error

- * If structure contains bitfield then we can not apply & (address of operator) on the members of that structure.

eg: struct demo

```
{  
    int i; j;  
};  
struct demo obj;  
printf("Enter value");  
scanf("%d", &obj.i); // error
```

int *p = &obj.i; // error
int *p = &obj; // allowed

12/8/15

All the above points are applicable as it is for union also means we can use only concept such as anonymous, bitfield for union also only the major difference between

union & structure is according to memory allocation

- when we create structure which contains some data elements
- when we try to fetch address of that data it shows in sequential manner

eg: struct demo

```
{  
    int i;  
    int j;  
    int k;  
};  
struct demo obj;
```

demo

100		demo::i
104		demo::j
108		demo::k
112		

← 4 bytes →

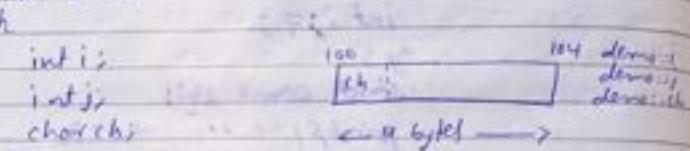
22. 3. 15

Consider base address of that object is 100 when output of following printf statements

```
printf("Id", &obj[i]) ; 11100  
printf("Id", &obj[j]) ; 11104  
printf("Id", &obj[k]) ; 11108
```

If we consider union and its data members then address of every element is same because there is no separate memory allocation for every element

Eg: Union demo



3.

Union demo obj;

Consider base address of obj is 100 then output of below printf is

```
printf("Id", &obj.i) ; 11100  
printf("Id", &obj.j) ; 11100  
printf("Id", &obj.ch) ; 11100
```

* User-defined Datatypes:-

Enumeration:- 1) Similar to Structure Union, enum is also considered as user-defined datatype 2) Enum is also called as an Enumeration Constant because the value get stored in enum is always constant 3) As like Structure Union we can also create enum as user-defined datatype

Eg: struct demo

```
struct demo {  
    int i, j, k;  
};
```

enum hello

```
enum hello {  
    i, j, k;  
};
```

- 4) According to above syntax demo is considered as a structure which contains 3 integers co-ordinates at i, j, k
- 5) Similarly hello is name of enum where i, j, k are considered as enumeration constants or integral constants

5

when

calculator
create

enum i
enum j
enum k

of below

is allo
red as
ed inside
rule can

a
i,j,k;
contined

- 6) if structure contains multiple elements then separate memory is allocated for every elements of structure but in case of enum there is no separate memory allocation is required. 7) Generally memory for enums is allocated inside text section of process. 8) According to above we cannot apply size of operator on variable of enum.
eg. enum demo {a,b,c};

- 9) In this enum every variable a,b,c are not initialized with any value due to which 1st variable of enum gets its default value start from zero & every other variable gets values incremented by 1.

a; 110 b; 111 c; 112

10) enum demo {a=5, b, c};

In this enum 1st variable is initialized with value 5 due to which concept of default initialization is not used.

11) As 1st variable is initialized with value 5 due to which next variable gets its next value i.e. 6 & so on.

a; 115 b; 116 c; 117

12) enum demo {a=7, b,c,d,a};

Above enum generates error because a variable declared twice in enum which is not allowed because it generally ambiguity means same named variable is not allowed in enum.

13) enum demo {a=11, b, c=11};

In this enum both variables a & c gets its value as 11 which is allocated means same value can be used in enum.

14) enum demo {a=3+2, b, c};

In this case 1st variable is initialized with only expression which gets evaluated is 5 due to which remaining variable gets its value 6 & 7.

a; 115 b; 116 c; 117

* To initialize enum variable we can use any expression that should evaluate to an integral constant.

enum demo(a=-6, b, c);

To enum variable we can store -ve value also.

a; 11-7 b; 11-6 c; 11-5

we can not fetch value of any variable from
text section

7) enum demo { };

This enum is considered as unnamed enum and similar
like structure union.

8) enum demo { a,b,c } obj;

In this enum obj is variable of type
enum demo.

According to the above syntax our object
variable can be initialized with a values if enum
variable only otherwise it generates error.

obj = 1; // error

obj = b; // allowed

9) enum demo { a,b,c };

a = 15; // error

b++ ; // error

Above two statements generates compilation error
because we cannot change value of enumeration
variable.

10) enum demo { a,b,c };

int *p = &a; // error

As we can not fetch address of enum variable
it generates compilation error.

11) enum demo { a = 'a', b = 'c' };

above syntax is allowed because 1st variable a
gets initialized to ASCII value of a i.e 97
which is integral constant

a; // 97

b; // 98

c; // 99

similar

type

object
of enum
s

variable
iteration

variable

variable a
o g f d

12) enum demo {a,b,c};
struct hello
{
 int i;
 enum demo obj;
};

Any structure/union contains object of enumeration constant. This member of structure should be initialized with value of enum only.

13) struct demo obj;
obj.j = 11;
obj.obj = 'a';
obj.obj = 15; //error

*typedef of the last #/p>

4*) string in c:-

String is not considered as a direct datatype

- 1) In C++ string is considered as an existing hierarchy class which contains some predefined functions.
- 2) In java String is considered as datatype. In C & C++ string is consider as character array which ends with '\0'.
- 3) '\0' is considered as a delimiter for the string.
- 4) There are multiple ways in which we can create string in C & C++ as

e.g.
1) char arr[] = { 'H', 'E', 'L', 'L', 'O', '\0' };
2) char buf [] = "Hello";
3) char *p = "Hello";

In this 3 ways we can create string in C or C++

5) if we use 1st way for creation of a string, we have to insert '\0' at the end explicitly. And this explicit insertion of '\0' is not need in case of Syntax 2 & Syntax 3 because it gets implicitly added by compiler.

6) If string is created by using member member initialization list then '\0' is expected at the end of string.

7) But if string is enclosed in " " then there is no need to insert '\0' at the end.

arr [H E L L O \0]
 | | | | | | | |
 0 1 2 3 4 5 6 7

buf [H E L L O \0]
 | | | | | | | |
 20 21 22 23 24 25 26

P [30+] → [H E L L O \0]
 | | | | | | | |
 30 31 32 33 34 35 36

According to all the above diagrammatic representation every string requires 6 bytes of memory '\0' is initiated at the end of every string.

* strlen:- it is a string library function which accept base address of string & it counts no of characters from that base address till we get first '\0'.

* sizeof:- it is a unary operator which gives the no of bytes which should be allocated for every datatype. sizeof operator returns the no of bytes by taking help of symbol table, datatype table.

g, we
atty. But
and in
gets

members
of at

there is

elation

'o' is

with
no
will be

the
every
by
de.

int i = 10;
sizeof(i); // 4 info fetched from symbol table
sizeof(int); // 4 info fetched from datatype table
Consider the above syntax of string & predict the
o/p

sizeof(arn); // 6
strlen(arn); // 5
sizeof(buf); // 6
strlen(buf); // 5
sizeof(p); // 4 due to pointer
strlen(p); // 5
sizeof(*p); //
strlen(ptr); // 3

As we have syntax we can create string from
that 3 syntax. we can reinitialize the string
by using 3rd syntax only of
p = "world";

arr[] = { 'a', 'b', 'c', 'd' }; // error
arr = "world"; // error

In above all syntax are we can only change
contents of string in 1st syntax or
arr[1] = 'x';

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* C Preprocessor

- 1) Preprocessor is a first component of ~~xxc~~ tool chain which gets involved after editor in a tool chain.
- 2) Preprocessor is a program which accepts in high level programming language such as C and C++ and convert that input into expanded source code.
- 3) instead of scanning whole program preprocessor scans only such lines which starts with # symbol.
- 4) According to this '#' symbol gives direction to the preprocessor due to which '#' is called as preprocessor directive symbol.
- 5) This preprocessor program performs multiple task such as,
 - 1) Tokenization and whitespace removal
 - 2) file inclusion
 - 3) Macro expansion
 - 4) Conditional compilation
- 6) After performing all this task our expanded source code is still in high level programming language then that source code is passed towards the compiler for remaining processing.

A.1) Tokenization and whitespace removal:-

- 1) Tokenization is a first process which is performed by preprocessor.
- 2) the process of tokenization our preprocessor except statement of comments from our source code.

- 1) That statement gets converted into multiple ~~with spaces~~ and that word is called of token. and above process is called of tokenization.
- 2) After performing the process of tokenization over them gets forwarded towardsly the compiler.
- 3) in this phase of preprocessor, preprocessor removes extra whitespaces which may contain normal space in it in our program, that all things gets removed by the preprocessor.
- 4) if a programmer writes a comment in a program then all that comment gets removed by the postprocessor.
- 5) after passing through this phase of a preprocessor our program is consider as a normal file according to an operating system i.e "unformatted uniform stream of bytes".

Hello.c

```
#include <stdio.h>
int main() // main function
{
    printf("Hello");
}
```

Hello.i

```
#include <stdio.h> int main () { printf ("Hello"); }
```

↓ ↓
 add contents from tokenization
 in .i file file inclusion

* 2) file inclusion :-

- " This is a preliminary task of preprocessor which is used to include the contents of header file in our source file.
- " Header file is a just a file which is normal **.txt** file which contains:
 - i) prototypes of a function
 - ii) declarations of user-defined datatype or structures, enums, unions, classes
 - iii) typedefs which are used in our program
 - iv) some user-defined macros (**#define**)
 - v) header file inclusion for other headerfile.
- " when preprocessor gets **#include** keyword in our program then the preprocessor search that specified header file and copy the content of that header file in our source file.

eg: Hello.c

```
#include <stdio.h>
int main()
{
    printf("Hello world");
}
```



Hello.i

```
int printf(const char *format, ...);  
int scanf(const char *format, ...); } #3  
Contents of  
stdio.h
#define NULL 0
int main();
```



```
printf("Hello world");
```

- macro which includes file
- normal
- datatype
of
any program
`#define`)
header file.
- inout
put
readoff
file.
- * of
h
- (i) According to above code snippet after preprocessing the contents of stdio.h header file gets added in our source code.
 - (ii) There are two syntactical ways in which we can include our header file such as:
 - (i) `#include <filename>`
 - (ii) `#include "filename"`
 - (iii) if filename is written within angular bracket the preprocessor search that file in compile specified or system specified path.
 - (iv) if file name is written in double code then compiler search that file in current directory of a process.
 - (v) if that file is not found in a current directory then preprocessor search that file in compile specified path.
- * current directory :- current directory of process is a directory in which our parent process resides.
- * parent directory :- parent directory of a process is a directory in which our file resides.
- ### MACRO Expansion:-
- i) There are 2 types MACRO such as:
 - (i) predefined macro - provided by compiler
 - (ii) userdefined macro - provided by programmes

3) Every type of macro is divided into 2 types

a)

i) object like macro -

e.g. #define MAX 10

ii) function like macro -

e.g #define MULT(X,Y) X * Y

4) we have to use macro for 2 reasons given as

i) A using macro program become much more readable

ii) if we want to change value of any

macro then there is no need to scan whole program but we just change value of one position it gets automatically replaced in whole program by preprocessor itself.

5) Macro is divided in two parts such as

macro template & macro definition

e.g. #define MAX 10

↓ →
macro template macro definition

6) To perform this macro expansion talk preprocessor creates one data structure named as macro expansion table which contains 2 columns in it.

7) first column contains macro template & 2nd column containing macro definition.

Macro template	Macro definition
MAX	10
NULL	0

2. Function

- i) while performing the task of macro expansion preprocessor replace every macro template of our program with its definition by taking help of macro expansion table.

e.g. Hello.c

```
#define INTEREST 9  
void calculate (int no)  
{  
    int finalAmount = 0;  
    finalAmount = (no * INTEREST)/100;
```

↓
Hello.i

```
void calculate (int no)  
{  
    int finalAmount = 0;  
    finalAmount = (no * 9) / 100;  
}
```

- ii) Below eg explains the concept of function like macros

Hello.c

```
#define mult(x,y) x*y  
int main()  
{
```

int no=0;

no = MULT(10,20);

↓

Hello.i

```
int main()
```

{

int no=0;

no = 10*20;

↓

eg. There is no such concept of scope of macro.

We can define same macro multiple time which is called as redefinition of macro.

eg:-

void fun1()

{

#define MAX 10

printf("Id", MAX); 10

gun();

printf("Id", MAX); 10

3

void gun()

{

#define MAX 20

printf("Id", MAX); 20

2

int main()

{

fun();

return;

2

(i) whenever redefine the macro in a program just is no separate entry gets added in macro expansion table but it only replace old definition with new definition.

(ii) when preprocessor gets #define keyword newly gets added in macro expansion table and also it gets #undef keyword that entry gets removed from macro expansion table.

MACRO
in
MACRO

after getting ~~getting~~ ~~#undef~~ we can not use that MACRO because it gets removed from MACRO expansion table

neid fun()

#define MAX 10
pf ("1d", MAX); //10

undef MAX
pf ("1d", MAX); //error

we can also write macro which contains multiple lines in it.

If we want to include next line in our macro then at ~~at~~ end of previous line there should be \ at the end.

eg:

#define MAX 10
printf("Hello");

in this example printf statement is also included in our MACRO.

here

* Some predefined macros in C and C++

definitions

i) stringized macro:-

This macro is used to convert our parameter into string.

eg: #define Demo(st) #st

int main

{

 pf (Demo(Helloworld));

 ↓
 printf("Helloworld");

any
less

ii) Concatenation macro

This macro is used to concat input arguments

e.g. #define DEMO(X,Y) X##Y

int var=10;

printf("%d", DEMO(var));

printf("%d" var);

iii) Macro which give filename of a program

e.g. printf("%s", --FILE--);
// return (filename)

iv) Macro which gives the line no

14 e.g.

15 printf("%d", --LINE--); 1115

16
↓
(line no)

v) Macro which gives the system time

printf("%s", --TIME--);

// System time
(11:20 AM)

vi) Macro which gives current DATE

printf("%s", --DATE--); // System date

23 Aug 2015

vii) Macro which is used to check our compiler is Standard compiler or not

~~if~~ if(--STDC == 1)

printf("std compiler");

↓

* Miscellaneous directive of preprocessor.

1) #pragma pack(1)

This macro is used to instruct compiler about memory allocation techniques.

2) #pragma warning -nul

This macro is used to display the warning or hide the warning
nul = return value required warning

3) #pragma startup fun

this pragma startup macro is used to call the function before calling main
in this syntax fun function gets called before main

```
void fun() {  
    // code  
}
```

4) #pragma exit fun

this exit pragma directive is used to call the function after terminating main function

84 Conditional Compilation:

1) Concept of conditional is used to instruct the compiler that which part of code should be compiled and which part of code should not be compiled to use this concept there are multiple macros available such as

```
#if  
#endif  
#else  
#elif
```

eg. int main()

{
int no = 5;

#if (no == 5)

 printf("Hello");

Hello

 printf("world");

world

in the above program printf("world") statement is not
the compile by compiler because condition is false.

eg. #define MAX 10

#ifdef MAX

 printf("Hello");

#ifndef #endif

~~#ifdef~~ and ~~#ifndef~~

* #ifdef and #ifndef are used to check whether
the specified macro is defined or not in program
or not.

at above syntax of MAX macro is define

due to which the below printf statement is compiled
successfully.

* operator in "C"

1) Expression :- Expression is a statement which gets internally evaluated by the processor.

Expression containing more than one operator & more than one operand.

2) Operator :-

operator is considered as a symbol which decides the action taken on the identifiers.

3) operands :-

operands are such entity on which we have to perform the operation.

*) Types of Expression:-

There are two types of expression according no of operators in it.

1) Simple Expression:-

If expression contains single operator in it then it is called as simple expression.

Eg. $n = 15;$
Operand \leftarrow ~operator

2) Compound Expression

If expression contains more than one operator in it then it is called as compound expression.

Eg: $A + B = C;$

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* precedence of an operator:-

If expression contains more than one operator in it then precedence of an operator decides the sequence in which subexpression gets evaluated.

* Associativity of an operator:-

If expression contains multiple operators in it having same precedence then associativity of an operator decides in which sequence we have to evaluate all subexpressions i.e. L → R or R → L.

* Types of operators:-

operators are classified according to two types

1) Types of operators according to no of operands

according to the no of operands operators are divided into 3 types as:

1) unary operator - requires single operand

2) binary operator - requires two operands

3) ternary operator - requires three operands

2) Types of operators it depends on its behaviour or operations

1) Arithmetic operators:-

Eg: +, -, *, /, %

2) Relational operators:-

Eg: <, ≤, ≥, ≠, ==

3) Equality operators.

Eg: ==, !=

operator is
decided
by

multiple
then
which sequence
is i.e.

two types
of operators

1) one

2) operand
operator

view of

4) logical operators:-

&&, ||, !

5) bitwise operators:-

&, ^, ~

6) miscellaneous operators

, -r, sizeof, (), [], .

7) Shorthand assignment operators:-

+ =, -=, *=, /=, %=, <<=, >>=

8) Shorthand operators

++ , --

9) Shift operators:-

<<, >>

precedence & Associativity table

precedence	operator	type	Associativity
1	<code>>>> + - ~ * &</code>	miscellaneous	L→R
2	<code>! ~</code>	unary	R→L
3	<code>* / %</code>	multiplicative	L→R
4	<code>+ -</code>	Additive	L→R
5	<code><< >></code>	shifter operator	L→R
6	<code><= >= == !=</code>	relational	L→R
7	<code>= += -= *= /=</code>	equality operators	L→R
8	<code>&</code>	bitwise and	L→R
9	<code>^</code>	bitwise EX-OR	L→R
10	<code> </code>	bitwise OR	L→R
11	<code>&&</code>	logical and	L→R
12	<code> </code>	logical or	L→R
13	<code>? :</code>	Conditional	R→L
14	<code>= += -= *= /= <= >= &= ^= =</code>	Assignment	Right → Left
15	<code>,</code>	comma	L→R

associativity

17 miscellaneous operators :-

$\rightarrow R$

$\rightarrow L$

$\rightarrow R$

miscellaneous operators are sub-operators which are provided by language designers
a) () - this operator generally used to increase precedence of operator
eg: Ans = (a+b) * c

$\rightarrow R$

$\rightarrow R$

$\rightarrow R$

in above expression according to the normal precedence rule multiplication operator evaluate first but due to the () bracket first + operator evaluates first

$\rightarrow R$

$\rightarrow R$

if any expression contains any sub expression in a () bracket then that sub expression gets evaluated first without considering the precedence rule.

$\rightarrow R$

operator
Expression - Ans = $x + \underline{(y-z)}$ ← Sub Expression
 |
 |
 operator
 specifies

$\rightarrow R$

$\rightarrow R$

in above compound expression ; if subexpression i.e. $(y-z)$ gets evaluated first before + operator.

$\rightarrow L$

b) [] :- Subscript operator

$\rightarrow L$

this operator generally used in 'C' programming language if we want to declare an array or if we want to access any element of array.

eg: int arr[7];
 printf("%d", arr[2]);

c) ' . ' operator (direct access operator)

it is used to access the members of structure
and union if we have variable of that structure

d) ' -> ' operator (indirect access operator)

it is used to access members of structure
and union if we have pointer which point to
structure or union

Eg: struct Demo

{

int i, j;

y;

struct Demo obj;

struct Demo *p = &obj;

obj.i = 11;

p->i = 21;

(*p).i = 21;

i) we can use . & -> operator in different ways

such as obj.i = 21;

p->i = 21;

(*obj)->i = 11;

(*p).i = 21;

(e) ++ & -- (Shorthand increment or decrement)

Eg: int no = 10;

no++;

↓
no = no + 1

no--;

↓
no = no - 1

increment and decrement operator is divided into 2 types of

future
value

done
'to'

↳ Postincrement of ~~post-decrement~~: post-decrement
↳ Preincrement of pre-decrement

This two types are applicable ~~only~~ if increment or decrement operators used in expression

If it used independently then both this types of pre & post belong in a same manner

Eg. the below eg. indicates difference b/w preincrement and postincrement operator of

1) int no = 10;
no++; // 11

2) int no = 10;
++no; // 11

3) int no = 10; ans = 0;
ans = no++;
printf("%d%d", ans, no); // 10, 11

4) int no = 10; ans = 0;
ans = ++no;
printf("%d%d", ans, no); // 11 11

According to above four example in first & 2nd example there is no difference preincrement and postincrement but if same concept is used in Expression then preincrement & postincrement are considered as different concept.

* all the concept are applicable for decrement operator also. *

2) unary operators

all this operators are unary operator because this operator requires single operand.

a) +, - : These operators are used to assign sign to our variable.

e.g. int no = 10; ans =

ans = +no;

ans = -no;

b) ! ~

both operators called as not operators but ! is considered as logical not operator whereas ~ is considered as bitwise not operator.

logical not operator not going to change any of the data whereas bitwise operator change state of our data.

e.g. if (!no < 10)

{

 // code

}

int no = 10;

ans = ~no;

as this is a bitwise operator first we have to consider bit representation of value 10.

0 0 0 0 1 0 1 0
~ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓
1 1 1 1 0 1 0 1

(C) * & & (dereference operator and addressof operator)

* operator is called as addressof operator which is used to return the virtual address of any identifier used in our program.

& operator gives the virtual address which is available in symbol table.

virtual address is real address which is provided by the compiler at the time of compilation but it is not an existing address.

* i.e. dereference operator which is used to return the contents whose address is stored in any variable.

Consider below syntax & symbol table to understand the concept of & and * operator.

int x = 15; $x \boxed{15}$
 400 404

int no = 10; $no \boxed{10}$
int xp = &x; $no \uparrow 100$
 p $\boxed{100}$
 no 204

symbol table:-

Name	address	value	size	from	to	annotation
x	400	15	4	6	4	-
no	100	10	4	6	4	-
p	200	100	4	6	21	-

when we use name of identifier only then prompted search its name in name column of symbol table and return the contents of value column.

1) `printf("1d", 7); // 15`
if we use `d` operator before identifier name
then processor search that name in name
column & return the contents of address
column.

e.g. `printf("1d", 8no); // 100`

2) if we use `*` dereference operator then first
we have to search the name in name column,
fetch the value of corresponding that name

That value should be present in address column of
symbol table and returns its corresponding
contents from value column.

e.g. `printf("1d", *p); // 10`

3) `sizeof` is unary operator which returns its result
by considering the contents of symbol table or
datatype table.

e.g. `sizeof(no); // 4` Consider symbol table
`sizeof(double); // 8` Consider datatype
table

int no = 10;
`printf("1d", sizeof(++no)); // 4`
`printf("1d", no); // 10`

The expression written inside `sizeof` operator is evaluated
but its result not reflected in original
operator.

3) multiplicative operators:-

(a) * - this is binary operator which gives multiplication of two operands.

e.g. int ans = 0, a = 5, b = 2;
ans = a * b; // 10

(b) / - division operator:-

ans = 10 / 3 // 3

(c) % - it gives the remainder which gets generated after division

e.g. no = 10 / 3 ; // 1
no = -10 / 3 ; // -1
no = -10 / -3 ; // -1
no = 10 / -3 ; // 1
no = 10.0 / 3 ; // error

mod (%) operator is only applicable for integral constant it is not applicable for floating point variables

4) Additive operators:- (+, -)

no = 10 + 3 ; // 13
no = 10 - 3 ; // 7

30 8 12

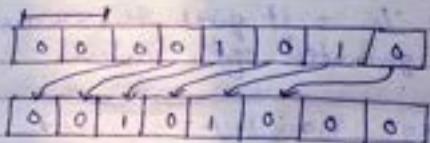
5) Shift operators:- This type of operators are used to shift bits of our data in left direction or right direction.

Generally these shift operators are used to manipulate our data & modify that data.

i) \ll Left shift operator:-

int no = 10;

no = no \ll 2; // no = 10



If we shift out no in left direction by n no of bits then n empty bits created on right side.

Always empty bits should be initialize with zero only.

We can directly calculate result of shift operator by using formula of $x \ll y$

$$x \ll y = x \times 2^y$$

In our eg. x is 10 & y is 2 due to which

$$\begin{aligned} 10 \ll 2 &= 10 \times 2^2 \\ &= 10 \times 4 \\ &= 40 \end{aligned}$$

We can not apply any shift operator on floating point no. because floating point no are not stored in sequential format.

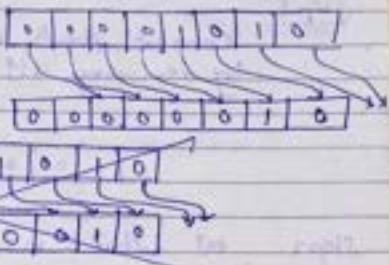
12

Type of conversion

2) \gg right shift

int no = 10;

no = no \gg 2;



if we shift our no to right side magnitude of
our no gets decreased

if we want to calculate result by using formula
then consider

$$[x \gg 2 = x/y]$$

in our example x is 10 and y=2 due to this

$$10 \gg 2 = 10/2$$

$$= 10/4$$

$$\therefore 2$$

* negative no representation of RAM

1) if we store the no then that no gets stored
by directly converting into its binary format

2) if our data is -ve then processor performs some
special task while storing that no into RAM

e.g.: Consider no as -10

int no = -10; // how stored in memory

Step 1 convert that no into its binary representation by removing its sign.

0 + 1 0 0 0 1 0 1 0

Step 2 as that no is +ve number first we have to calculate its one's complement by toggling every bit of a number.

1 1 1 1 0 1 0 1

Step 3 apply the concept of 2's complement on above data.

in case of 2's complement we have to add one in our actual data.

1 1 1 1 0 1 1 0 1

+ 1

1 1 1 1 0 1 1 0
Signbit ←

the output which is generated by the above step gets stored in memory.

after step 3 the first bit of no. one which is part of our actual data but that bit is used for representing whether our no is +ve or -ve.

representation

if first bit is one then our no is consider
-ve otherwise consider +ve.

the below steps are performed when we read the no.
whose first bit is one i.e. sign bit is on.

we have
toggling

Step 1 Consider that no as it is which gets stored in
a memory.

1	1	1	1	0	1	1	0
---	---	---	---	---	---	---	---

Step 2 apply one's complement logic on above no.

0	0	0	0	1	0	0	1
---	---	---	---	---	---	---	---

in above

Step 3 apply concept of 2's complement of above no

0	0	0	0	1	0	0	1
---	---	---	---	---	---	---	---

of one

+ 1

0	0	0	0	1	0	1	0
---	---	---	---	---	---	---	---

as this no becomes 10 when we try to print
that no on screen it display -10
because 1'st bit of no is one when we receive
that no from RAM.

* difference b/w signed & unsigned variables

Step

Sign bit

the

Signed & Unsigned are considered of datatype
modifiers due to which range of our datatype
gets modified.

Consider the example of character which requires
8 bit.

if our character is signed then we can write
only 7 bits from that 8 bits because one bit
is reserved for signed representation

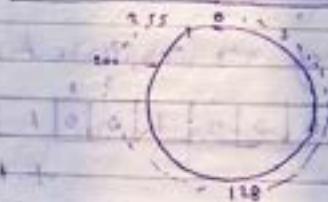
char datatype
/ \
signed unsigned

consider the example of character as.

unsigned character

7	6	5	4	3	2	1	0
1	0	0	1	0	0	0	0

actual data



-128

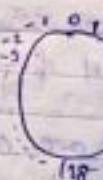
0	1	0	0	1	0	0	0
---	---	---	---	---	---	---	---

2) Signed character

signedbit ->	1	0	1	0	0	0	0
--------------	---	---	---	---	---	---	---

actual data

signedbit is on



signedbit is off

-128

6. Relational :- This operator are binary operator which are to check relationship b/w two operator result of this two operator is either true or false

```
if (10 > 7)
{
    printf("Yes");
}
else
{
    printf("No");
}
```

In above example relational operator is used in if condition as and condition is true our control comes inside if block.

7. Equality operator:-

This freq has == & != are used to check whether contents of both the operators are equal or not.

Similar as the concept of relational operator it returns either true or false depend on equality.

8.9.10 bitwise (&, |, ^) AND, OR, EX-OR :-

all this operators are use to manipulate bits of our data.

op1	op2	&		^	0	1
1	1	1	1	0	0	0
0	1	0	1	1	1	1
1	0	0	1	1	1	0
0	0	0	0	0	0	1

It is not truth table

int no1 = 15, no2 = 21, Ans = 0;

convert no1 into binary

2	15	2	21
2	7	2	10
2	3	2	5
2	1	2	2
2	0	2	1
		2	0
		1	

1) Ans = no1 & no2

no1 = 00001111

00001111

no2 = 00010101

00010101

Ans = 00001001

00001001

printf("%d", Ans); // 5

2) Ans = no1 | no2

no1 =

00001111

no2 =

00010101

Ans =

00011111

3) Ans = no1 ^ no2

00001111

00010101

00010101

printf("%d", Ans); // 26

Note that two last

11.12) Logical operators

example to demonstrate concept of precedence & associativity

$$Ans = A + B * C \& X / Y - (Z + T) / R$$

$$\text{out}1 = Z + T$$

$$Ans = A + (B * C) \& X / Y - \text{out}1 / R$$

$$\text{out}2 = B * C$$

$$Ans = A + \text{out}2 \& X / Y - (\text{out}1 / R)$$

$$\text{out}3 = \text{out}1 / R$$

$$Ans = (A + \text{out}2) \& X / Y - \text{out}3$$

$$\text{out}4 = A + \text{out}2$$

$$Ans = \text{out}4 \& X / Y - \text{out}3$$

$$\text{out}5 = Y - \text{out}3$$

$$Ans = (\text{out}4 \& X / Y) / \text{out}5$$

$$\text{out}6 = \text{out}4 \& X$$

$$Ans = \text{out}6 / \text{out}5$$

$$\text{out}7 = \text{out}6 / \text{out}5 \quad Ans = \text{out}7$$

logical operators apply on the condition this operators not applicable as any of the operator

there are 2 logical operators in C such as

&&, || both this operators are binary operators
about which it requires two operands.

both the operands can be an expression the result of and operator is depend on truth or falsity of operand 1 & operand 2

to apply the concept of logical operators we have to consider truth table of

op1\op2	00	11
T	T	T
F	F	T
F	T	F
F	F	F

in case of logical operators preprocessor internally uses short circuit evaluation rule as

i) if operation to be performed is logical
or then

a) if first operand is true then we have to check result of 2nd operand

b) if result of first operand is false then there is no need to check the result of 2nd operand because overall result becomes false

ii) if operation to be performed is ||

a) if result of first operand is true then there is no need to check the result of 2nd operand because overall result is true

b) if result of first operand is false then we have to check result of 2nd operand.

e.g. int no1=10, no2=7, no3=6;

(i) if ((no1==10) && (no2==6))

{ pf("yes"); }

|| yes

(ii) if ((no1==10) || (no2==6))

{ pf("yes"); }

|| not evaluated

value = 117

	no1	no2	value
1	T	F	1
2	F	T	1
3	T	T	1
4	F	F	0

is conditional operator :-

This is only one ternary operator in C
means it requires 3 operands.

This operator is formed of an conditional operator
because everything which can be written by using
if else construct can be written by using
ternary operator.

Consider the below example of if else condition
it corresponds to conditional operator.

```
if (no == 5) {  
    A = 7;  
} else {  
    A = 6;  
}
```

Similarly as above example we can also convert
nested if else into one ternary operator such as.

```
if (no == 7) {  
    if (x == 5) {  
        A = 5;  
    } else {  
        A = 6;  
    }  
}
```

↳ If no is 7 then A is 5
↳ If no is 7 then A is 6

```
else {  
    if (x == 7) {  
        B = 7;  
    } else {  
        B = 6;  
    }  
}
```

$(x = -7) : ((y = -5) : (A = 5) : (A = 6)) : ((x = -7) : ((B = 6) : (B = 7)))$

14) Assignment & short hand operator

If we want to assign value to one operand
and then to another operator then assignment
operator should be use.

e.g. int no1=10
no2=no1
assignment

Here are some standard assignment operator like
etc

$C = C \times D$ $A = A / B$ $A = A + B$ $A = A - B$ $A = A \ll 2$
 $C *= D$ $A /= B$ $A += B$ $A -= B$ $A \ll= 2$

$B = B / 2$

$B /= 2$

15) ',' Comma operator:-

Comma is a such operator which is not always
considered as a operator but some place it is
considered as a separator.

In below example , is worked as a separator.
e.g. int a[] = {10, 20, 30, 40, 50};

17
7);

fun(10, 20)

void add(int a, int b);

The below example is use of operator

int a = 10, 20, 30, 40; // 10

[int a = 10]

int b = (10, 20, 30, 40); // 40

[int b = 40]

24/10/22
Continue

TypeDef:-

- 1) TypeDef is not a storage class.
- 2) TypeDef is concept which is used in C & C++ which is used to assign some different name to an existing datatype.
- 3) By using typeDef we can not create new datatype.
- 4) When we create typeDef its entry gets added into datatype table which is used to manage all information concern with the datatype.
- 5) `typedef int Integer;`
According to above syntax Integer is considered as new name for datatype int.
- 6) After this typeDef statement can use Integer as a datatype.
`Integer no = 10;`
Above statement is considered as
`int no = 10;`
- 7) We can apply concept of typeDef for any primitive datatype, any user-defined datatype, any user-defined datatype.

Eg: struct demo

```
i  
int i;  
int j;  
int k;  
y;
```

in C & C++
parent name

new datatype

A added
manage
datatype.

is consider

Integer

any

- 8) If we want to create variable of structure demo then we can create variable of struct demo obj;
- 9) According to this syntax we have to write struct keyword & name of struct everytime to avoid this we can create it's typedef as
`typedef struct demo newtype;`
- 10) According to this Syntax newtype is new name for existing datatype named as struct demo after this syntax we can create variable of that structure of newtype obj;
- 11) we can create typedef for structure & union at the time of declaration itself
`eg: typedef struct demo
{ int i;
int j;
} new-type;`
- 12) newtype is not object for struct demo but it is typedef for struct demo As it is a typedef we can create variable of structure as new-type obj;
- 13) If we are creating typedef for structure & union at the time of declaration then we can create that structure as unnamed structure also
`eg: typedef struct
{ int i, j;
} new-type;`
- 14) According to above syntax new-type is typedef for that unnamed structure

18) Using `typedef` we can create multiple names for some datatype such as

```
typedef struct demo  
{  
    int i;  
} newtype1, newtype2;
```

In this Syntax, `new-type1` & `new-type2` are considered as new-type names for structure `demo`.

- a. we can compare the concept of `#define` with `typedef`
eg `typedef int *iptr1;`
`#define iptr2 int *`

`iptr` is considered as datatype like `int *` where as in case of `#define` `int *` is an replacement text for keyword `iptr`.

```
iptr1 p,q;  
iptr2 x,y;
```

According to above Syntax `p,q` are considered as `2 variables of type int*`

where as `x` is a variable which is considered as `int*` & `y` variable is only considered as `int`.

Accordingly to above difference, we can conclude that `typedef` is applicable for every variable that we are creating but `#define` is not applicable for all.

- * There are some major differences between `#define` & `typedef` just consider it yourself.

and for

- 1) all typedef are handled by compiler whereas #define handled by preprocessor
↳ entry for typedef gets added in datatype table whereas entry for #define macro gets added in macro expansion table
- 2) Replacement of macro is just blind replacement whereas replacement of typedef is logical replacement

```
typedef int (*FP)(int,int)
FP(X+Y)
int add (int i, int j)
{
    return (i+j);
}
int sub (int i, int j)
{
    return (i-j);
}
```

considered

considered
as int

```
X = add;
Y = sub;
X(10,20); // call to function add
Y(20,10); // call to function sub
```

In above syntax FP is not considered as a variable of funⁿ pointer but it is considered as new datatype which is similar as function pointer which stores address of such a function which accepts 2 parameters both are integers & it return integer.

#define

module.h

Basic description of kernel module.

```
1 // Hello2.c  
2 #include <linux/module.h>  
3  
4 int init_module(void)  
5 {  
6     printk(KERN_INFO, "module is loaded");  
7     return 0;  
8 }  
9  
10 void cleanup_module(void)  
11 {  
12     printk(KERN_INFO, "module is removed");  
13 }  
14  
15  
16  
17  
18  
19  
20
```

1) this header file is required for the flags which are used in our module.

init-module and cleanup-module's function prototype is written in that header file.

2) this header file contains some predefined symbols of kernel which are used at run time.

4) this function gets automatically invoke by insmod utility at the time of loading the kernel module after insmod command memory gets allocated inside kernel address space and the specific .ko file gets loaded into that memory.

after invoking this function name of c_kernel_module gets added in the list of modules which is maintained by kernel at location /proc/modules.

6) it is not possible to perform user interaction from the kernel module program but the states of kernel module can be maintained in kernel log file.

printk function is used to write the data in that particular file.

7) this file gets automatically created when boot up at location /var/log/kernlog
printk is variable no of argument function in which the first argument should be jiffies that is superiority of the msg and 2nd argument is the message that we want to write.

8) if we return 0 value for init function then it indicates the success. if there is any failure at the time of loading kernel module return value is -1.

this return value is not used for programmes but the corresponding error msg gets displayed in the kernel log file.

9) cleanup-module is a keyword which is used for such a function which gets implicitly called at a time of removing module from memory
when we rmmod command is invoked this function gets invoked.

generally this function is used to deallocate the resources for the kernel module.

used in

when

of

mod

le

located

No file

~~Chapter 11~~
Section 11.2
initfunc in above program init module and cleanup module is the name of function which gets called by insmod and rmmod command

we can also provide any userdefined name for this function but we have to register that user defined name with module_init and module_exit macros to use that macro successfully in our program we have included <linux/init.h> header file

we can also used __init and __exit flag at a time of defining this function.

this flags are used to optimize the memory for our kernel module due to this flag memory gets immediately deallocated when the execution of that function gets completed.

26. module 3 : inside kernel module we can create local variables which can be used in our kernel module.

datatype and name of variable is same normal c programming technique.

static int no __initdata = 2;

in this syntax due to the static we can not access variable the outside the file

int is the datatype of that variable

no is the name of that variable

__initdata is the macro which indicates that the memory for that variable should be allocated at the time of loading the kernel module

int *arr
 int arr[2][3][4];

```

int ***p = (int ***)malloc(2 * sizeof(int));
for (i=0; i<2; i++)
{
  int **p = (int **)malloc(3 * sizeof(int));
  for (j=1; j<3; j++)
    {
      int *p = (int *)malloc(4 * sizeof(int));
    }
}
  
```

1770
 590 30 → 150
 70 → 9 D COUNT

~~$\frac{30}{10} \times 150$~~ 1 E O U COUNT

Local
 all.
 mal

decay

ABCDEF ← 2 3 4 last row
 B → C E A D

A B C D E F G H I J K L M

Z Y X W V U T S R Q P O N

That
 ended

He 110	Demolition	System
100	200	400
*	*	*
600	604	608
700	704	708
800	804	808
900	904	908
1000	1004	1008
1100	1104	1108
1200	1204	1208
1300	1304	1308
1400	1404	1408
1500	1504	1508
1600	1604	1608
1700	1704	1708
1800	1804	1808
1900	1904	1908
2000	2004	2008
2100	2104	2108
2200	2204	2208
2300	2304	2308
2400	2404	2408
2500	2504	2508
2600	2604	2608
2700	2704	2708
2800	2804	2808
2900	2904	2908
3000	3004	3008
3100	3104	3108
3200	3204	3208
3300	3304	3308
3400	3404	3408
3500	3504	3508
3600	3604	3608
3700	3704	3708
3800	3804	3808
3900	3904	3908
4000	4004	4008
4100	4104	4108
4200	4204	4208
4300	4304	4308
4400	4404	4408
4500	4504	4508
4600	4604	4608
4700	4704	4708
4800	4804	4808
4900	4904	4908
5000	5004	5008
5100	5104	5108
5200	5204	5208
5300	5304	5308
5400	5404	5408
5500	5504	5508
5600	5604	5608
5700	5704	5708
5800	5804	5808
5900	5904	5908
6000	6004	6008
6100	6104	6108
6200	6204	6208
6300	6304	6308
6400	6404	6408
6500	6504	6508
6600	6604	6608
6700	6704	6708
6800	6804	6808
6900	6904	6908
7000	7004	7008
7100	7104	7108
7200	7204	7208
7300	7304	7308
7400	7404	7408
7500	7504	7508
7600	7604	7608
7700	7704	7708
7800	7804	7808
7900	7904	7908
8000	8004	8008
8100	8104	8108
8200	8204	8208
8300	8304	8308
8400	8404	8408
8500	8504	8508
8600	8604	8608
8700	8704	8708
8800	8804	8808
8900	8904	8908
9000	9004	9008
9100	9104	9108
9200	9204	9208
9300	9304	9308
9400	9404	9408
9500	9504	9508
9600	9604	9608
9700	9704	9708
9800	9804	9808
9900	9904	9908
10000	10004	10008

T	T	T	T	T	T	F
T	F	E	T	E	T	T
F	T	F	F	F	T	T
F	F	F	F	F	E	E

odd	$1+1 = 0$	1
	$1+0 = 1$	0
	$0+1 = 1$	0
	$0+0 = 0$	0

$$\text{int } P_{-10} \begin{array}{|c|} \hline 10 \\ \hline \end{array} 2^0 \cdot \begin{array}{l} 1010 - 10 \\ 10100 - 20 \\ \hline 01011 \end{array} \quad 3^3$$

int *ph = &p; 100
 int **ptph = &ph; 000
 **ptph = 200; ph

int const xpt; $\tau = \infty$
const int xpt = 4;

`const int const x[4];`

A hand-drawn illustration on lined paper. It features a large circle in the center. Inside the circle, there is a sketch of a face with a cross-like mark over its eyes. To the left of the circle is a small square shape, and to the right is another small square shape.

7e006
0x7e00

const int wptrs = 0 0 0 0 0 0

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00000000000000000000000000000000

wptrs = 101

10100000000000000000000000000000

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10100000000000000000000000000000

100101

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011010

+ 1

011011

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10000000000000000000000000000000

101000

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01010100000000000000000000000000

+ 11111111111111111111111111111111

01100100000000000000000000000000

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