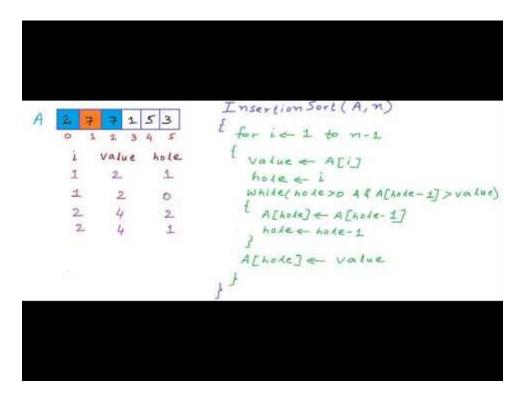
## **Insertion sort**

Insertion sort is a simple sorting algorithm that works the way we sort playing cards in our hands.

## Algorithm:

```
// Sort an arr[] of size n
insertionSort(arr, n)
Loop from i = 1 to n-1
    a) Pick element arr[i] and insert it into sorted sequence arr[0...i-1]
```



- Time Complexity: O(n<sup>2</sup>)
- Auxiliary Space: O(1)
- Boundary Cases: Insertion sort takes maximum time to sort if elements are sorted in reverse order. And it takes minimum time (Order of n) when elements are already sorted.
- Sorting In Place: Yes
- Stable: YesOnline: Yes

Algorithmic Paradigm: Incremental Approach

An **online algorithm** is one that can process its input piece-by-piece in a serial fashion, i.e., in the order that the input is fed to the algorithm, without having the entire input available from the start.

Uses: Insertion sort is used when number of elements is small. It can also be useful when input array is almost sorted, only few elements are misplaced in complete big array.

## **Next**

- What is Binary Insertion Sort?
- · How to implement Insertion Sort for Linked List?

## **Concepts Problems**

Given an array arr[], a pair arr[i] and arr[j] forms an inversion if arr[i] < arr[j] and i > j. For example, the array {1, 3, 2, 5} has one inversion (3, 2) and array {5, 4, 3} has inversions (5, 4), (5, 3) and (4, 3).

If we take a closer look at the insertion sort code, we can notice that every iteration of while loop reduces one inversion. The while loop executes only if i > j and arr[i] < arr[j]. Therefore total number of while loop iterations (For all values of i) is same as number of inversions. Therefore overall time complexity of the insertion sort is O(n + f(n)) where f(n) is inversion count. If the inversion count is O(n), then the time complexity of insertion sort is O(n). In worst case, there can be  $n^*(n-1)/2$  inversions. The worst case occurs when the array is sorted in reverse order. So the worst case time complexity of insertion sort is  $O(n^2)$ .

Merge sort can also be used to count inversion

*Insertion sort* is faster for small input size (n) because *Quick Sort* has extra overhead from the recursive function calls. Insertion sort is also more stable than Quick sort and requires less memory.