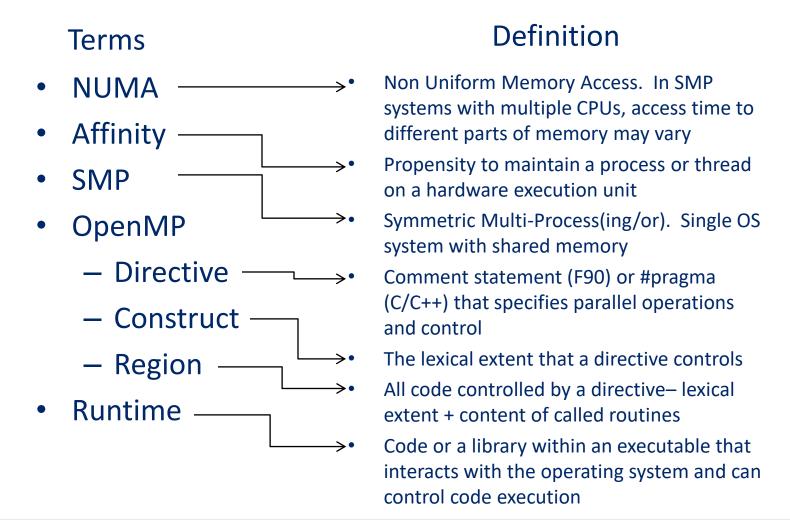
COMP 364 / 464 High Performance Computing

OpenMP:

Multi-threaded and task parallelism

Computing Terminology



What is OpenMP (Open Multi-Processing)

- De-facto standard for high-performance parallel programming on Symmetric Multi-Processor (SMP) Systems
- It is an API (Application Program Interface) for designing and executing parallel Fortran, C, and C++ programs
 - Based on threads, but
 - Higher level than POSIX threads (*Pthreads*)
 http://www.llnl.gov/computing/tutorials/pthreads/#Abstract
- Implemented by:
 - Compiler Directives (or #pragma's)
 - Runtime Library (interface to OS and Program Environment)
 - Environment Variables
- Compiler option required to interpret/activate directives
- http://www.openmp.org/ has tutorials and description
- Directed by OpenMP ARB (Architecture Review Board)

OpenMP -- Overview

- Standard developed in the late 1990s
- The "language" is easily grasp. You can start simple and expand.
- Lightweight from system perspective
- Very portable GNU GCC and vendor (Intel) compilers.
- OMP compilers generate the parallel code but developers must define the parallelism and communication mechanisms.
 - Time spent finding parallelism can be the most difficult part. The parallelism may be hidden.
 - Writing parallel OpenMP code examples is relatively easy.
 - Developing parallel algorithms and/or parallelizing existing serial code is much harder.
 - Expert level requires awareness of scoping and synchronization.

OpenMP Executable Runs on an SMP*

- Shared Memory Systems:
 - One Operating System
 - Instantiation of ONE process
 - Threads are forked (created) from within your program
 - Multiple threads on multiple cores

^{*} SMP = Symmetric Multi-Processor: The execution of the operating system has equal access to any of the "processors"

OpenMP History

Primary OpenMP participants

AMD, Cray, Fujitsu, HP, IBM, Intel, NEC, PGI, Oracle, MS, TI, CAPS, NVIDIA ANL, LLNL, cOMPunity, EPCC, LANL, NASA, ORNL, RWTH, TACC

OpenMP Fortran API, Version 1.0 Published October 1997

OpenMP C API, Version 1.0
 Published October 1998

OpenMP 2.0 API for Fortran
 Published 2000

OpenMP 2.0 API for C/C++ Published 2002

OpenMP 2.5 API for C/C++ & F90 Published 2005

OpenMP 3.0 Tasks
 Published May 2008

OpenMP 3.1 Published July 2011

OpenMP 4.0 Affinity, Offload, SIMD Published July 2013

OpenMP 4.5
 Published Nov 2015

Advantages/Disadvantages of OpenMP

Pros

- Shared Memory Parallelism is easier to learn than distributed MP
- Coarse-grained or fine-grained parallelism
- Parallelization can be incremental
- Widely available, portable
- Converting serial code to OpenMP parallel can be easier than converting to multi-process, distributed parallel
- SMP hardware is ubiquitous now: Supercomputers and your desktop/laptop

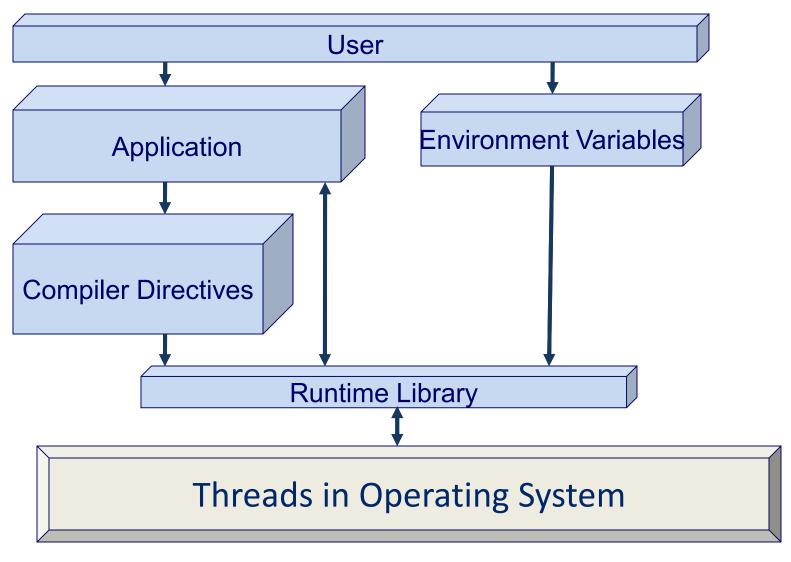
Cons

- Scalability limited by memory architecture:
 - Non-uniform memory access (NUMA) can be devastating.
 - Thread synchronization dependent upon fast access to global variables in main memory.
- Available on SMP systems "only"
- Beware: "Upgrading" large serial code may be time-consuming ... but it always will be!
- Not directly part of the language.

OpenMP Parallel Directive

- Supports parallelism by Directives (Fortran) and Pragmas (C/C++)
- Unlike others that require base language changes and constructs (Pthreads)
- Unlike MPI which supports parallelism through communications library between processes.

OpenMP Architecture



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OpenMP Syntax

OpenMP Directives: sentinel, construct and clauses

```
C #pragma omp construct [clause [[,]clause]...]
F90 !$omp construct [clause [[,]clause]...]
```

Example

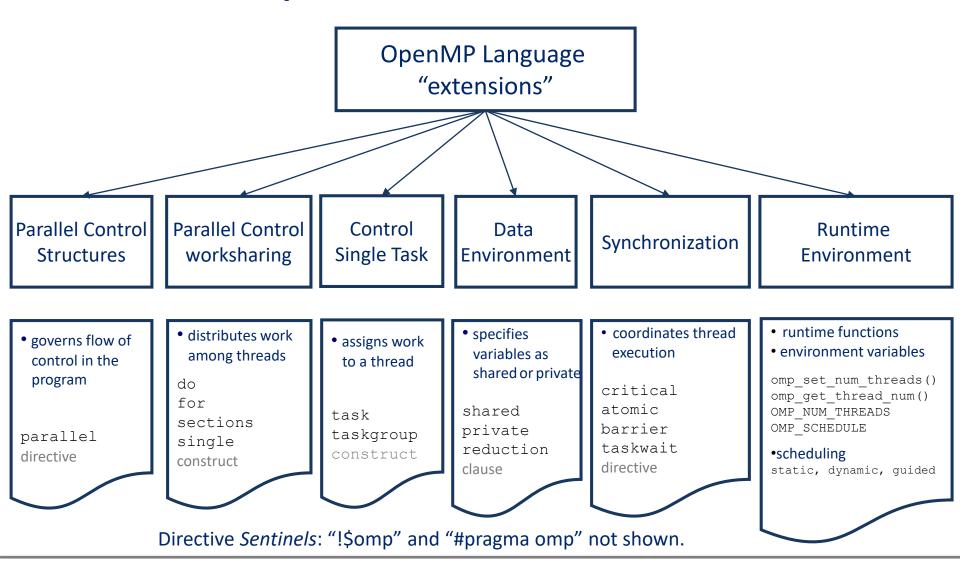
```
C #pragma omp parallel num_threads(4)
F90 !$omp parallel num threads(4)
```

Function prototypes and types are in the file:

```
C #include <omp.h>
F90 use omp lib
```

Most OpenMP constructs apply to a "structured block", that
is, a block of one or more statements with one point of entry
at the top and one point of exit at the bottom

OpenMP Constructs



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OpenMP Directives

• OpenMP directives begin with special comments/pragmas that a OpenMP-aware compiler can interpret. Directive sentinels are:

```
F90 !$OMP
C/C++ # pragma omp
```

Syntax: sentinel parallel clauses

uses defaults when clauses not present

```
!$OMP parallel
...
!$OMP end parallel
```

```
# pragma omp parallel {...}
```

- Fortran parallel regions are enclosed by enclosing directives.
- C/C++ parallel regions are enclosed by curly brackets.

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Parallel Region & Thread Number

```
1 #include <omp.h>
2 ...
3 int ierr = 0;
4 #pragma omp parallel
5 {
6    int numThreads = omp_get_num_threads();
7    int threadId = omp_get_thread_num();
8    ierr = do_work( threadId, numThreads );
9 }
```

Every thread can inquire the total number of threads (numThreads in line 6) and get a unique value for the thread number [0,K-1)

Parallel Region & Thread Number

```
#ifdef OPENMP
1
   #include <omp.h>
   #endif
5
   int ierr = 0;
6
   #pragma omp parallel
7
      int numThreads = 1;
8
   #ifdef OPENMP
10
      numThreads = omp get num threads();
11 #endif
12
      ierr = work(numThreads);
13 }
```

OpenMP code can be disabled and ignored by the compiler. All sentinels and directives ignored when OMP disabled by compiler. Pre-processor flag useful to remove OMP API function calls and header. _OPENMP automatically defined when OMP enabled.

Parallel Region & Worksharing

Use OpenMP directives to specify Parallel Region, worksharing constructs, and mutual exclusion

parallel
end parallel

Use parallel ... end parallel for F90 Use parallel {...} for C

parallel for
parallel sections

Code block

for

sections

single

master

critical

atomic

Each Thread Executes

Worksharing: splice loop

Worksharing: splice calls

Only one thread (first there)

Only master thread (0)

Execute one thread at a time

Update one thread at a time

A single worksharing construct (e.g. **for**) may be combined on a parallel directive line.

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Parallel Region

```
#pragma omp parallel
{
    #pragma omp for
    for ( i = 0; i < n; i++) {
        work(i);
    } // <- implicit barrier
} // <- implicit barrier</pre>
```

- In above example the **for** loop iterations are split among the threads via the **for** worksharing construct.
- Which iteration(s) are executed by which threads? Controlled by schedule construct. Default is to split the iterations into K=#threads chunks, each n/K long.

OpenMP Combined Directives

- Combined directives
 - parallel for and parallel sections
 - Same as parallel region + one for worksharing construct

```
#pragma omp parallel for
for (i= 0; i< 100; i++) {
    a[i] = b[i];
} // <- implicit barrier</pre>
```

- Fixed (known) trip count and increment required
- no break
- limited C++ throw
- continue ok (since it doesn't impact other iterations)
- Basic rule of thumb: if it can be done with SIMD, it can be threaded with OpenMP-for

Parallel Region

worksharing (WS) constructs: for, sections, and single

- WS Threads execution their "share" of statements in a PARALLEL region.
- for worksharing may require run-time work distribution and scheduling

```
1 #pragma omp parallel for
2 for (i= 0; i< n; i++) {
3    a[i] = b[i] + c[i];
4 } // <- implicit barrier</pre>
```

```
Line 1: Team of threads formed (parallel region).

Line 2-4: Loop iterations are split among threads.

Implied barrier at "}".
```

Each loop iteration must be independent of other iterations!

Replicated vs. Workshare Constructs

- Replicated: Work (code) blocks are executed by all threads.
- Worksharing: Work is divided among threads ... no overlap.

```
code1();
 #pragma omp parallel
                            #pragma omp parallel for
                                                                   #pragma omp for
   {code}
                            for (i=0; i<4*n; i++)
                                                                   for (i=0; i<4*n; i++)
                                 {code}
                                                                       code();
                                                                   code3();
                                                                } // end parallel
                            i=[0,n) i=[n,2n)
                                           i=[2n,3n)
                                                     i=[3n,4n)
                                                                code1 code1
                                                                                code1
                                                                                         cdde1
code code code
                             code code
                                                       code
                                             code
                                                               i=[0,n) i=[n,2n) i=[2n,3n) i=[3n,4n)
                                                                code code
                                                                                code
                                                                                        code
                                                                code3 code3
                                                                                code3
                                                                                         code3
   Replicated
                                   worksharing
                                                                           Combined
```

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#pragma omp parallel

OpenMP Worksharing Scheduling

Clause Syntax: parallel for schedule(schedule-type[,chunk-size])

Schedule Type

Schedule (static, chunk)

- Threads receive chunks of iterations in thread order, round-robin.
 (Divided "equally" if no chunk size ... N / p)
- Good if every iteration contains same amount of work ... <u>uniform</u> workload.
- May help keep parts of an array in a particular processor's cachegood between **parallel for**'s.

Schedule (dynamic, chunk)

- Thread receives chunks as it (the thread) becomes available for more work ... a queue.
- Default chunk size may be 1
- Good for load-balancing non-uniform workloads.

OpenMP Worksharing Scheduling

Schedule (guided, chunk)

- Thread receives chunks as the thread becomes available for work.
- Chunk size decreases exponentially, until it reaches the chunk size specified (default is 1).
- Balances load and reduces number of requests for more work.
- (I have never found this useful.)

Schedule (runtime)

- Schedule is determined at run-time by the OMP_SCHEDULE environment value.
- Useful for experimentation

OpenMP Worksharing Scheduling

For example, loop with 100 iterations and 4 threads

• schedule (static)

Thread	0	1	2	3	
Iteration	Iteration 1-25		51-75	76-100	

• schedule (dynamic, 15) (one possible outcome)

Thread	0	1	3	2	1	3	2
Iteration	1-15	16-30	31-45	46-60	61-75	76-90	90-100

 Dynamic scheduling has overhead since assigning the "next" chunk requires thread communication (access to shared data object)

OpenMP worksharing -- Sections

sections

- Blocks of code are split among threads task parallel style
- A thread might execute more than one block or no blocks
- Implied barrier

```
#pragma omp sections
{
    #pragma omp section
    { TASK1(); }
    #pragma omp section
    { TASK2(); }
    #pragma omp section
    { TASK3(); }
}
// <- implicit barrier</pre>
```

OpenMP worksharing - Single

- **single** (or master)
 - Block of code is executed only by first thread (single) or thread-0 (master)
 - Implied barrier at '}' after single ... no barrier after master.

```
1 int global_count = 0; // Defined globally.
2 ...
3 #pragma omp parallel shared( global_count )
4 {
5    fool();
6    #pragma single
7    {
8       global_count++;
9       printf("%d\n", global_count);
10    }
11    foo2();
12 }
```

OpenMP Data scoping

#pragma omp directive-name [clause [[,]clause]...]

- private (variable list)
 - Each thread has its own copy of the specified variable
 - Variables are undefined after worksharing region
- shared (variable list)
 - Threads share a single copy of the specified variable
- default (type)
 - A default of private, shared, or none can be specified
 - Note that loop counter(s) of worksharing constructs are always private by default; everything else is shared by default

OpenMP Data Scoping

- firstprivate (variable list)

firstprivate

lastprivate

- Like **private**, but copies are initialized with the master thread's value
- lastprivate(variable list)
 - Like private, but final value is copied out to master thread's copy
 - for: last iteration; sections: last section
- reduction (op:variable)
 - Each thread has its own copy of the specified variable
 - Can appear only in reduction operation
 - All copies are "reduced" back into the original master thread's variable.
 - Example reduction (+:sum) adds all values together.

OpenMP Data Scoping

- for and parallel for constructs
 - index variable is automatically private
- Automatic storage variables
 - private, if declared in a scope inside the construct (e.g. ordinary local variables declared inside functions)

OpenMP worksharing - Single

- shared Variable is shared (seen) by all processors.
- **private** Each thread has a private instance of the variable.
- Defaults: All loop indices are private, all other variables are shared.

```
#pragma parallel for shared(a) private(t1,t2)
for (i=0; i<1000; i++){
   t1 = f[i];
   t2 = g[i];
   a[i] = sqrt( t1*t1 + t2*t2 );
}</pre>
```

OpenMP Data Scoping

```
int sum = 0;
#pragma omp parallel for reduction(+:sum)
for (i=0; i<N; i++) {
   sum = sum + a[i];
// Each thread's copy of sum is added
// to original sum at end of loop
printf("sum= %f\n",sum);
#pragma omp parallel for lastprivate(temp)
for (i=0; i<N; i++) {
   temp = f[i];
printf("f[N-1] == f\n", temp);
// temp is equal to f[N-1] at end of loop
```

- nowait clause
 - Threads encounter a barrier synchronization at end of worksharing constructs.
 - Specifies that threads can proceed to the next task without waiting.

```
#pragma omp ... nowait
```

```
#pragma omp parallel for nowait
for (i=0; i<N; i++) {
   b[i] = foo(a[i]);
}
// Threads can start foo2() before others are finished with foo()

#pragma omp parallel for
for (i=0; i<N; i++) {
   d[i] = foo2(c[i]);
}
// All threads block (or sync) here.</pre>
```

- barrier directive explicitly adds synchronization point
 - All threads must reach the barrier and wait until the last gets there.

#pragma omp barrier

```
#pragma omp parallel
   const int thread id = omp get thread num(),
             num threads = omp get num threads();
   const int N per thread = N / num threads;
   int istart = thread id * N per thread,
         iend = istart + N per thread;
   for (i = istart; i < iend; i++)</pre>
     b[i] = foo( a[i] ); // Fixed cost per iteration
   // Don't let any threads past this point yet to avoid race
   // scenario on b[].
   #pragma omp barrier
   #pragma omp for schedule(dynamic) nowait
   for (i=0; i<N; i++) {
     b[i] += foo2(c[i]);
// All threads block (or sync) here implicitly.
```

- critical section permits access to only one thread at a time
 - All threads execute the critical section but they access the section with mutual exclusion ... only one can be inside the code at a given time.
 - This avoids race conditions but is very expensive.
 - #pragma omp critical

```
InitializeOueue( allWorkItems );
#pragma omp parallel
  bool notEmpty = true; // This is a private variable for each thread.
   while ( notEmpty )
      workItemType workItem;
      // Pop a task off the queue one thread at a time. This avoids
      // race condition on queue state.
      #pragma omp critical
         notEmpty = PopQueue( &workItem );
      // Do some work (if necessary)
      if ( notEmpty )
         foo(workItem);
// All threads block (or sync) here.
```