Chapter 1

One Dimension

1.1 One Dimension

This section covers the computing of functions of a single random variable. The following serve as motivating examples,

1. X^2
2. $log(X)$
3. $\frac{1}{X}$
4. $\frac{1}{log(X^2-X)}$
5. $[X-k]^+ - [X-k]^- \equiv X-k$

An interesting feature of the last item, known as Put-Call Parity [6], is that it seems to recover lost information through truncation. If $Z_+ := [X-k]^+$ and $Z_- := [X-k]^-$ then Z_+ and Z_- are maximally correlated. In the language of quantum mechanics the two Z's are *maximally entangled*.

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As a naming convenience it is assumed that X represents a *basic random variable*. A basic random variable is on that is not related to any other random variable within RICO by an explicit function. Conversely the label Z refers to a random variable that is functionally related to some X via an explicit function f such that $Z = f(X)$. When more than one function of X is needed they will be

distinguished by subscripting. Since a random variable is defined by its associated cumulative density function RICO requires random variables to be associated with the following computable expressions

$$P(Z < k) \qquad\qquad P(Z = k) \tag{1.1}$$

for any constant k. In order to incorporate functions of random variables into an algorithmic context is it also necessary to be able to compute

$$P(Z_1 < Z_2) \qquad \text{where} \qquad Z_1 = f_1(X),\ Z_2 = f_2(X)$$

Notice that $P(Z_1 < Z_2) = P(Z_1 - Z_2 < 0)$ and that $Z_1 - Z_2$ is itself a random variable. The computable expressions in (1.1) are sufficient to compute $P(Z_1 < Z_2)$ which may arise in conditional branching statements such as $if(Z_1 < Z_2)\{...\}$. Both branches of an *if* statement may be followed albeit with different associated probabilities. Conditional branching statements will be discussed in a later section.

1.1.1 Plotting $Z = f(X)$

A *continuous* random variable has the property that

$$P(X = k) = 0 \ \forall\ k \in \mathcal{D}(X) \tag{1.2}$$

and it assumed by RICO that a continuous random variable has an associated probability density function, $\rho(X)$. The expression $\mathcal{D}(X)$ refers to the support of X which in the continuous case is the domain of the associated probability density function.

A *discrete* random variable has discrete support such that

$$P(X = k_i) = p_i \ \forall\ k_i \in K = \{k_1, k_2, \ldots, k_n\} \tag{1.3}$$

In general a random variable X may be a mixture of continuous and discretely ditributed probability. In RICO the continuous and discrete aspects of a random variable are represented separately. The continuous aspect of a random variable is more numerically challenging and will be the focus of this section as a special case.

Notice that the discrete aspect of a random variable cannot be ignored even if an original X is purely continuous. Consider the motivating example, $Z = [X - k]^+$ in (5), transforms a continuous X into a *mixed* continuous/discrete Z.

In the next section it is assumed that both X and $Z = f(X)$ are continuous random variables.

Plotting Continuous $Z = f(X)$

To plot the probability density function of $Z = f(X)$ it is assumed that a software module will be employed to render the actual graph for the user and that the module requires an array of pairs of points,

$$\{(z_i, h_i)\}_{i=1}^{n} \tag{1.4}$$

where z_i is a point in the support of Z and $h_i = \rho(z_i)$, the probability density of Z at z_i. The associated probability density function of Z is denoted $\rho(Z)$.

To first approximation the values of each h_i are found by choosing a set of partition endpoints

$$(z_i^p)_{i=1}^{n} \subset \mathcal{D}(Z)$$

so that

$$z_i \in (z_i^p, z_{i+1}^p)$$

and the h_i are approximated as

$$h_i = \frac{p_i}{z_{i+1}^p - z_i^p} \qquad \text{where } p_i = P(z_i^p < Z < z_{i+1}^p) \qquad i = 1, \ldots, n$$

A *partition of* Z is understood to be a partition of the range of $f(X)$.

The relationship between the partition elements is shown in figure 1.1. Assuming $n \geq 2$ the z_i^p partition points are computed in RICO as midpoints of the $(z_i)_{i=1}^{n}$ array

$$z_i^p = \begin{cases} z_1 - \frac{1}{2}(z_1 + z_2) & \text{if } i = 1 \\ \frac{1}{2}(z_{i-1} + z_i) & \text{if } 1 < i \leq n \\ z_n + \frac{1}{2}(z_{n-1} + z_n) & \text{if } i = n + 1 \end{cases}$$

Let the associated partition of Z be

$$Z^p = (-\infty, z_1^p, z_2^p, ..., z_{n+1}^p, \infty)$$

Since Z is a purely continuous random variable the *partition elements* of Z^p are assumed in RICO to be open intervals. The missing endpoints of the partition elements form a set of probability zero. Let the partition elements of Z^p and associated probability values be indexed as

$$\begin{array}{lll} Z_0^p = (-\infty, z_1^p) & p_0 = P(Z_0^p) & \\ Z_i^p = (z_i^p, z_{i+1}^p) & p_i = P(Z_i^p) & i = 1, \ldots, n \\ Z_{n+1}^p = (z_{n+1}^p, \infty) & p_{n+1} = P(Z_{n+1}^p) & \end{array}$$

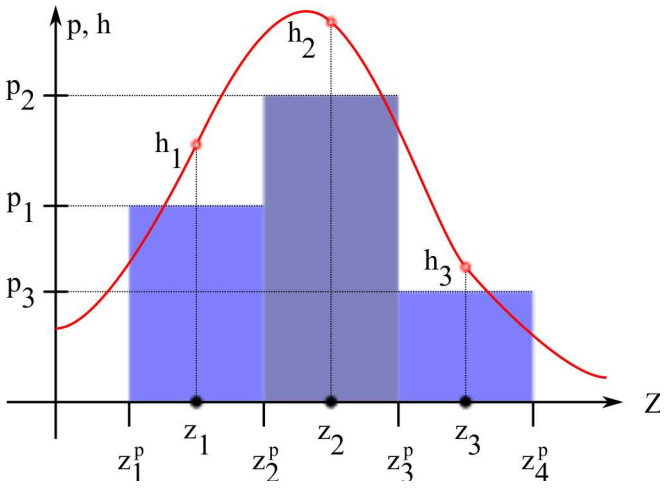

Figure 1.1: Three point Partition of Z

Numerical Considerations

Numerically, real numbers are represented by a finite number of floating point values. Let Ω be the largest representable floating point value. Similarly let ι be the smallest positive floating point value. In RICO, a constant called ω is defined as

$$\omega = min\left\{\Omega, \frac{1}{\iota}\right\} \tag{1.5}$$

and let ϵ be the smallest numerically representable positive probability value. One purpose for defining ϵ is so that so-called *thin tailed* probability distribution such as the Gaussian may be represented using a finite support interval as in

$$\mathcal{D}(X) = (X_{min}, X_{max})$$

where $-X_{min} = X_{max} \approx 10$, depending on ϵ. The X_{min} and X_{max} satisfy the following relations

$$P(X < X_{min}) \leq \epsilon \qquad\qquad P(X_{max} < X) \leq \epsilon$$

In particular all functions of a random variable have domain and range in $(-\omega, \omega)^2$ as represented visually by the shaded region in figure 1.2.

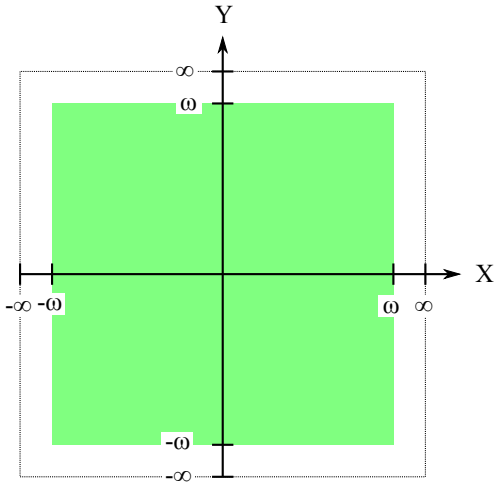

Figure 1.2: Fundamental Numeric Domain

Functions of Random Variables

In RICO there are several types of functions; *primitive*, *composite* and *piecewise monotonic*. Both primitive and composite functions are described elsewhere. An relevant feature of a primitive function such as $x^2, log(x)$, etc. is that it has a *primitive domain*. A primitive domain is a traditional mathematic domain. For example the primitive domain of x^2 is the interval $(-\infty, \infty)$ and for $log(x)$ the primitive domain is $(0, \infty)$. And example of a composite function is $x^2 - x$, a function composed of other composite functions, primitive functions and RICO-supported operations such as addition, subtraction, multiplication, etc.

In RICO, a piecewise monotonic function $f(X)$ of a random variable X is set of *piecewise monotonic function elements*. A piecewise monotonic function element is a domain interval and either another piecewise monotonic function element, a composite function or a primitive function. The set of piecewise monotonic function domain intervals form a non-intersecting set denoted $\mathcal{D}(Z)$ where $Z = f(X)$. The implied partition of the domain space is a set of open intervals containins $\mathcal{D}(Z)$ and denoted Z^p. The Z^p set is described above in the context plotting $f(X)$ and is defined here via

$$\mathcal{D}(Z)_j = z_i^p \in Z^p \qquad \text{for some } i \in 1, \cdots, n \text{ and } j \in 1, \cdots, m$$

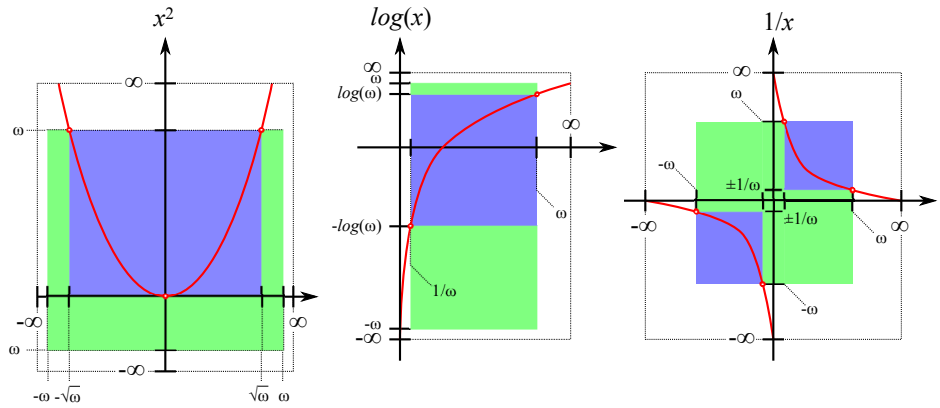

Figure 1.3: Refined Domains of Primitive Functions

For example, if $f(X) = [X - k]^+$ then

$$\begin{aligned} \mathcal{D}(Z)_1 &= (-\infty, k), \quad & \mathcal{D}(Z)_2 &= (k, \infty) \\ f_1(x) &= 0, & f_2(x) &= x \end{aligned}$$

where f_1 and f_2 are primitive functions.

Notice that the recursive nature of the definition of a piecewise monotonic function forms a tree structure denoted $\mathcal{T}(f)$. The leaf nodes of $\mathcal{T}(f)$ are piecewise monotonic elements whose associated functions are either composite of primitive. A critical aspect of the definition of a piecewise monotonic function in RICO is that leaf node functions are monotonic over their associated domain interval. For example the primitive function $f(x) = x^2$ has the associated primitive domain $(-\infty, \infty)$, but the piecewise monotonic function $f(x) = x^2$ is composed of two elements,

$$f(x) = \begin{cases} x^2 \text{ if } x \in (-\sqrt{\omega}, 0) \\ x^2 \text{ if } x \in (0, \sqrt{\omega}) \end{cases} \tag{1.6}$$

where the numerical considerations described in the previous section are taken into account.

Piecewise Monotonic Functions of Random Variables

For the purpose of plotting an other analysis such as computing $P(Z < k)$ for $Z = f(X)$ a composite f must be transformed into a piecewise monotonic f. The

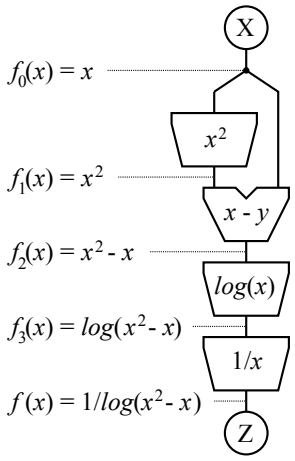

Figure 1.4: Parse Tree of $Z = 1/log(X^2 - X)$

transformation process is referred to as *refinment*. Function refinment serves two purposes. The first is to ensure that the numerical considerations are respected such that the associated function evaluates to a representable floating point value for every element of the domain. The second purpose is to ensure monotonicity within each domain interval.

The refinment for several primitive functions is depicted in figure 1.3 As a guiding example of the refinment process suppose that

$$Z = f(X) = \frac{1}{log(X^2 - X)}$$

where

$$X \sim Normal(0, 1) \text{ and } \mathcal{D}(X) = (-10, 10)$$

The parse tree of f is a composition of primitive functions, composite function and arithmetic operator components is shown in figure 1.4. Transforming the composite f into the piecewise monotonic f is a sequention process that begins at the input X. Several composite function components are identified in figure 1.4 including the indentity function $f_0(x) = x$.

To begin, note that the domain of X is the interval $(-10, 10)$ which is assigned to be the domain of f_0. The next step is to consider the effect of f_0 of the adjacent node x^2 and the y-input of the operation $x - y$. Assume that $f_1(x)$ is refined in isolation of the containing $f(x)$ so that its initial definition is given in equation

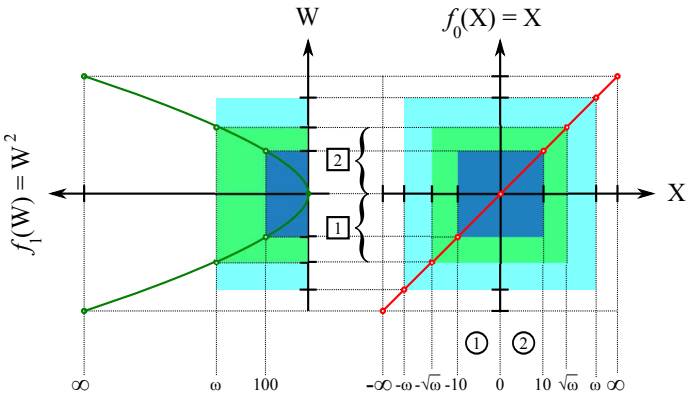

Figure 1.5: Refinement $f_0(X)$ by $f_1(X)$

1.6. The refinment process is to take the range of f_0, $(-10, 10)$, and intersect it with the domain of f_1 into range components $\{(-10, 0), (0, 10)\}$ and find the pre-image under f_0 which is also $\{(-10, 0), (0, 10)\}$. This process is shown in figure 1.5 where $\boxed{1} = (-\sqrt{\omega}, 0)$, $\boxed{2} = (0, \sqrt{\omega})$, $\textcircled{1} = (-10, 0)$, $\textcircled{2} = (0, 10)$.

In isolation the operation $x - y$ accepts any input in $(-\omega, \omega)$ so no further refinement is needed at this stage. Similarly no further refinement of f_2 is imposed by $x - y$. The domain of f_2 is then the set $\{\textcircled{1}, \textcircled{2}\}$ in figure 1.5.

The domain of f_2 in the parse tree 1.4 is refined by the domain of the primitive $log(x)$. Because f_2 involves a 2×1-dimensional transform, namely subtraction, and since the two operands are maximally entangled, the domain is less clearly defined and must be identified numerically. In general any value can be achieved by the transform $x - y$. Referring to the figure 1.6 the domain of $log(x)$ is $\boxed{1} = (1/\omega, \omega)$. The domain of $f_2(X)$ is notated as before,

$$\begin{aligned}\mathcal{D}(f_2) &= \mathcal{D}(f_1) \cap f_2^{-1}(\mathcal{D}(log)) \\ &= (-10, 10) \cap \left((-\sqrt{\omega}, x_1) \cup (x_4, \sqrt{\omega})\right) \\ &= (-10, x_1) \cup (x_2, 10)\end{aligned}$$

where

$$x_1 = min\left\{-f_2^{-1}(1/\omega), -\frac{1}{\omega}\right\} \qquad x_2 = max\left\{+f_2^{-1}(1/\omega), 1 + \frac{1}{\omega}\right\}$$

Numerical considerations dictate that since $f_2^{-1}(1/\omega) < 1/\omega$ the nearest x-value within each domain interval must be found. At this stage the function f_2 is refined

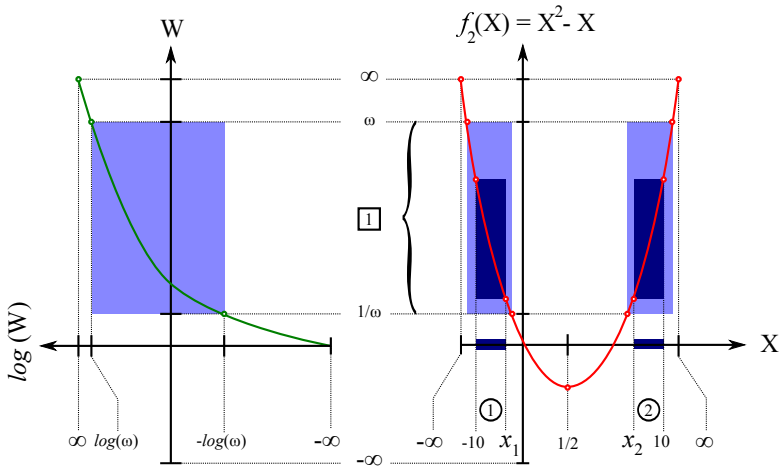

Figure 1.6: Domain of $f_2(X)$

by $log(x)$ and piecewise defined as

$$f_2(x) = \begin{cases} x^2 - x \text{ if } x \in (-10, -\frac{1}{\omega}) \\ x^2 - x \text{ if } x \in (1 + \frac{1}{\omega}, 10) \end{cases}$$

The last step of domain refinement in this example is shown in figure 1.7. Recall from 1.3 that the numerical domain of $1/x$ requires a small exclusion interval $C = (-1/\omega, 1/\omega)$. In isolation the primitive reciprocal function is defined on the piecewise domain $\{(-\omega, -1/\omega), (1/\omega, \omega)\}$. In figure 1.7 the domains are $\boxed{1}$ and $\boxed{2}$ respectively. The pre-images under f_3 are then intersected with the domain of refined f_2 yielding the refined domain of f_3 as

$$\begin{aligned} \mathcal{D}(f_3) &= \mathcal{D}(f_1(X)) \cap f_3^{-1}(\{(-\omega, -1/\omega), (1/\omega, \omega)\}) \\ &= (-10, x_3) \cup (x_4, x_1) \cup (x_2, x_5) \cup (x_6, 10) \end{aligned}$$

In figure 1.7 the refined domain of f_3 is shown as ①, ..., ④. Notice that the intervals (x_3, x_4) and (x_5, x_6) were excluded from the the refined domain of f_2 and that they are very small and thus likely of negligible probability.

Notice that since f does not feed into any other function elements in the example parse tree 1.4 requires no further refinement. The final function of $f(X)$ is shown in figure 1.8 with identified domain { ①, ②, ③, ④ }.

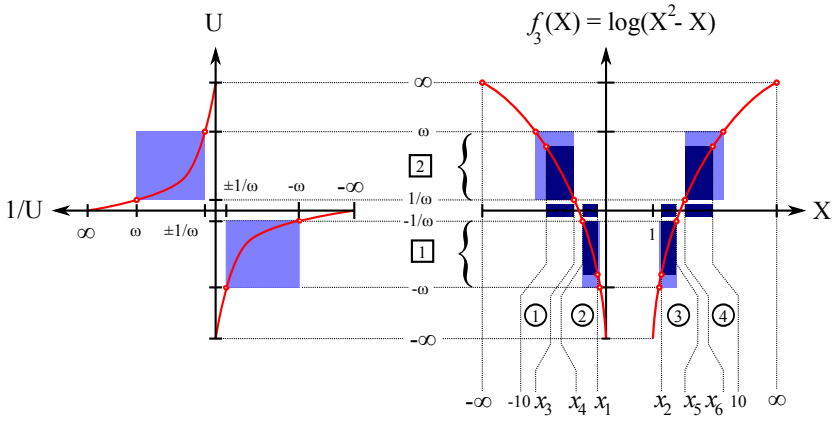

Figure 1.7: Domain of $f_3(X)$

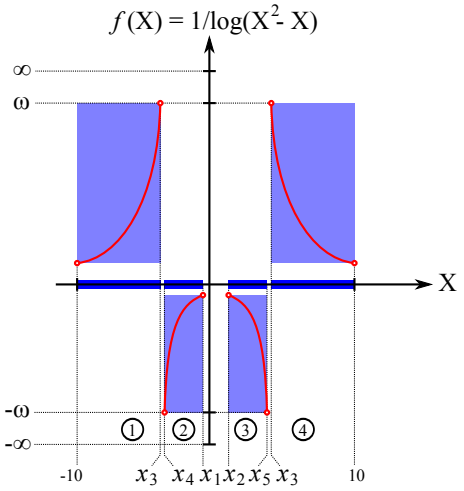

Figure 1.8: $Z = f(X)$ with Domain

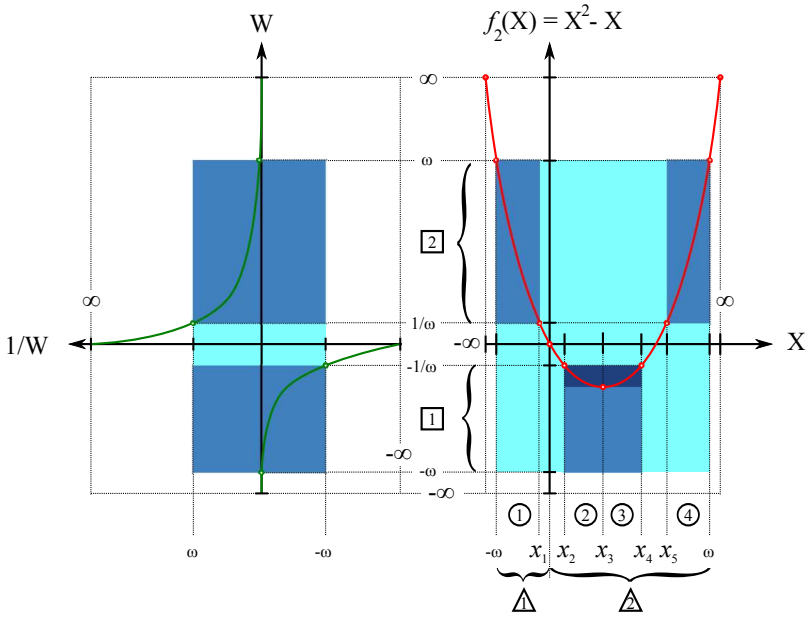

Figure 1.9: Refining $\frac{1}{x^2-x}$

Combining Monotonic Functions of Random Variables

The previous example revealed some important issues in computing functions of random variables. Breaking a function into non-overlapping monotonic fragments begs the question of how to locate interval endpoints. Domain refinement by intersection brought some numerical considerations when values became too small to represent as floating point values. The composition of two bijective functions is a bijective function, but the the same cannot be said for the sum or project of two bijections. The issue of refining a function that is the sum or product of bijections is addressed in this sub-section.

Consider the example,

$$f(x) = \frac{1}{x^2 - x}$$

Since $f(x)$ contains x^2 the initial domain set is $\{(-\omega, 0), (0, \omega)\}$ represented in figure 1.9 by $\triangle 1$ and $\triangle 2$ respectively. Recalling the fundamental domain set of $1/x$ from figure 1.3 as $\{(-\omega, -1/\omega), (1/\omega, \omega)\}$ as $\boxed{1}$ and $\boxed{2}$ respectively.

Continuing with figure 1.9, the refined domain element $\textcircled{1}$ is found by in-

tersecting the pre-image of $\boxed{2}$ under f given domain element $\triangle 1$. Finding the refined domain element (4) seems just a straightforward; the pre-image under f given domain $\triangle 2$ intersected with $\triangle 2$, but only because f happened to be monotonic in (4). Notice that f over domain element $\triangle 2$ is not monotonic. What is required is the intermediate step of refining each original domain element ($\triangle 1$ and $\triangle 2$ in this example) so that f is monotonic over each domain element before further refining f against $1/x$.

Chapter 2

Introduction

2.1 Introduction

The purpose of the work presented herein is to create a means of introducing uncertainty into deterministic models with performance superior to stand-alone Monte Carlo-based sampling techniques.

In the course of developing this material several spreadsheet-based models are analyzed. (See Appendix). A situation believed to be common is that spreadsheet-based models tend to be resistent to structural changes and their sensitivity to inputs is difficult to measure without external tools.

A remedy of the spreadsheet-based modeling issue is to reverse-engineer the spreadsheet to discover the specific model and re-implement this model in a traditional programming environment. This re-implementation process itself only addresses the issue of model brittleness and not the sensitivity issue. The work discussed in this paper recommends implementing a class of models such as those that may be implemented in a spreadsheet environment into a special environment. The environment proposed is called *RICO*, an achronym for *Random input, Correlated output*.

A RICO programming environment allows the programmatic manipulation of random variables as first class computing objects. The practical upshot is that if a given model accepts numeric data, that same model implemented in a RICO environment can substitute any input value with a random variable thereby allowing the study of model input sensisitivity and model response to input uncertainty. When multiple model inputs are replaced with random variables versions of these variables may become correlated within the model even if independent initially.

Numeric model outputs may then be a joint distribution of correlated random variables.

A RICO programming environment is not a specific software product, but a specification for constructing a software programming environment. While developing RICO a collection of software modules are constructed and used to produce many of the numeric and graphical results discussed in this paper. Collectively these software modules consitute a *reference* implementation of RICO. Any code snippets shown all run in at least one of the software modules within the RICO reference implementation.

Not limited to models implementable within a spreadsheet-based environment, RICO defines a number of basic mathematical operations and program flow control statements that are common to a wide variety of models. These operations include addition, subtraction, multiplication, division, exponentiation (both x^y and $exp(x)$) and logarithm. Program flow control statements include conditional statements such as $IF(X \leq Y)$ where either or both X and Y are optionally random variables. Interestingly, if at least one of X and Y is a random variable the conditional IF statement will take *both* not exclusively one or the other of the two implied code paths.

The class of models implementable in a RICO environment as large and several examples are explored. Some examples detailed below energy policy analysis, business cost savings analysis and geometric Black Scholes pricing. Important modeling components include constrained optimization and linear algebra involving random variables.

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   process to a sequence of random variables
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2.1.1 A Basic Model

At its heart a RICO model is a function of some number of inputs that produces some number of output symbolically represented in figure 2.1.

To introduce uncertainty into the model a random variable is identified for one or more inputs. As an example suppose a log-normal looking random variable X such as shown in figure 2.2.

Without loss of generality suppose a basic model with a number of input variables two of which (X and Y) are random and one output value of interest, Z. The Monte Carlo method of analysis requires a random sample for each

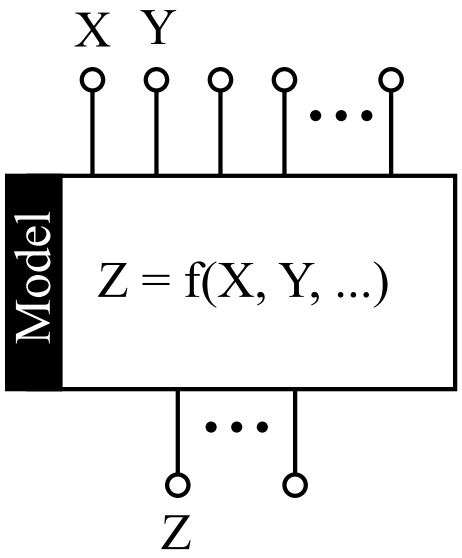

Figure 2.1: A Basic Model

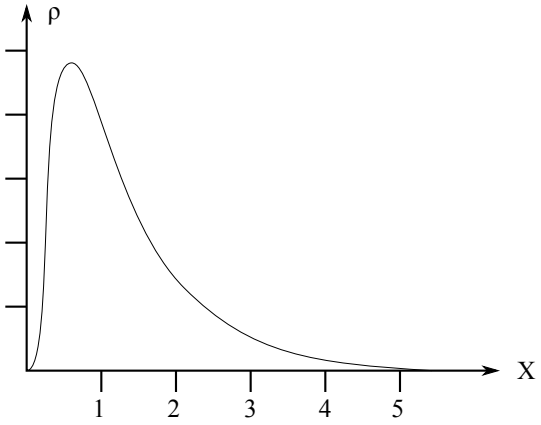

Figure 2.2: An Input Random Variable

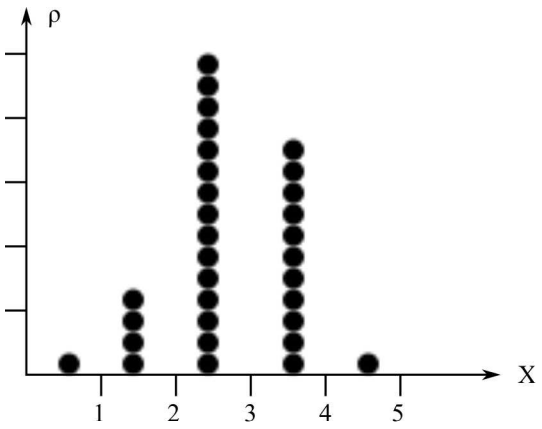

Figure 2.3: Example Output Histogram

random variable is generated. Suppose the random sample for X is denoted $x = (x_1, x_2, ..., x_n)$ and similarly for Y, $y = (y_1, y_2, ..., y_n)$. The model is run n times using the input values x_i and y_i in place of the X and Y inputs respectively for $i \in 1, ..., n$ generating n output values for Z called $z = (z_1, ..., z_n)$.

To understand the output values $z = (z_1, ..., z_n)$ generated by an n-run Monte Carlo method a partition of the space of output values is created by some means and the individual z_i values are counted into *bins*, that is, partition intervals. A possible result with unit-interval bins is shown in figure 2.3

A common exercise in statistics

reference	TO DO

is to postulate a familiy of probability distributions and fit the 'best' one to the output sample z. An example result is shown in figure 2.4

A serious concern is that in the context of an algorithmic model it is challenging to find a family of probability distributions from which to select a 'best' fit for the observed output sample, z. If the output is unimodal a subset of the extensive exponential family may be chosen, but for multi-modal output the choice is less clear.

The alternative offered by RICO is a hybrid approach. The RICO approach is to maintain a symbolic description of input and internal model variables where possible. When a symbolic description is not possible for an internal random variable, a numberic description is used. There is a practical limit on the number

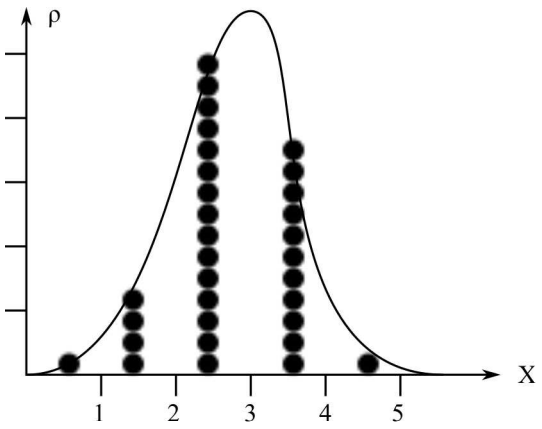

Figure 2.4: Example Output Histogram - Fitted

of dimensions used to describe a numeric random variable and when this limit is surpassed the final techniques applied is traditional Monte Carlo sampling.

A significant advantage that the RICO approach offers is symbolic representation of random variable tails. It is taken as axiomatic that the tail of a random variable is that region that cannot be reliably sampled. For random variables with so-called *thin-tailed* probability distributions, failure to adequately sample from the tails may be inconsequential, but *fat-tailed* distributions such as the Cauchy distribution have a significant amount of probability mass in the tails and this region of the support space cannot be safely ignored. Consider, for example, that the Cauchy distribution, like the St. Petersburg lottery, has no finite expected value.

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Chapter 3

Numeric Computation

3.1 Numeric Representation

Given a sample of a random variable and a partition of the support space as in figure 3.1 a histogram forms a natural summary. In RICO, a method of representing a one dimensional continuous random variable numerically is suggested in figure 3.2. Since random variables may be both discrete and continuously distributed at the same time a pair of parallel arrays is used programmatically,

$$X \sim \begin{cases} X_c = \begin{cases} (x_0 = -\infty, x_1, x_2, ..., x_n, x_{n+1} = \infty) \\ (p_0, p_1, ..., p_n, 0) \end{cases} \\ X_d = \begin{cases} (y_1, y_2, ..., x_m) \\ (q_1, q_2, ..., q_m) \end{cases} \end{cases}$$

where

$$-\infty < x_1 < ... < x_n < \infty$$

The endpoints for the continuously distributed portion of X, X_c, are assumed to be $\pm\infty$ with n partition endpoints between. The p_i values are defined as

$$p_i := \begin{cases} P(x_i < X_c < x_{i+1}) & \text{if } i \in \{1, ..., n-1\} \\ P(X_c < x_1) & \text{if } i = 0 \\ P(x_n < X_c) & \text{if } i = n+1 \end{cases}$$

and the q_j values are defined as

$$q_j := P(X_d = y_j) \text{ for } j \in \{1, ..., m\}$$

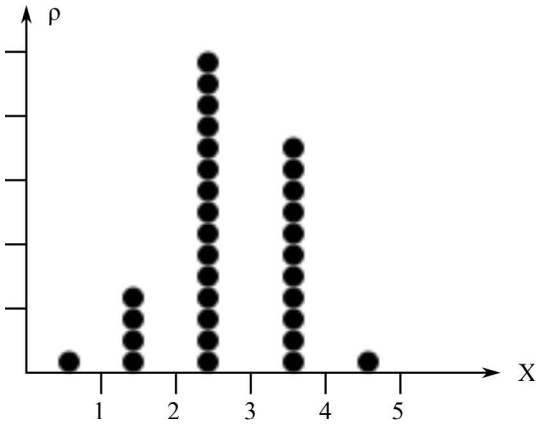


Figure 3.1: Example Sample

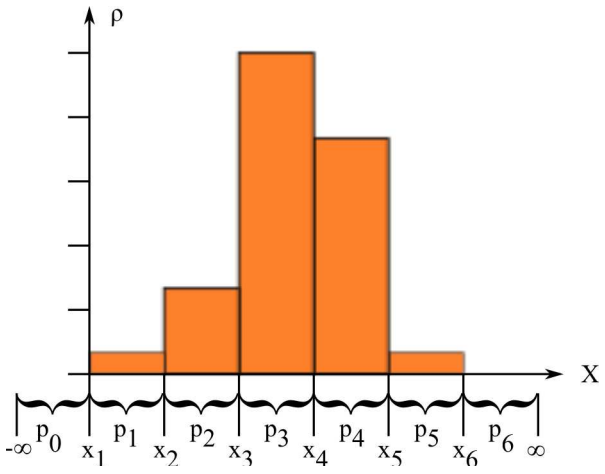


Figure 3.2: 1D Numeric Random Variable

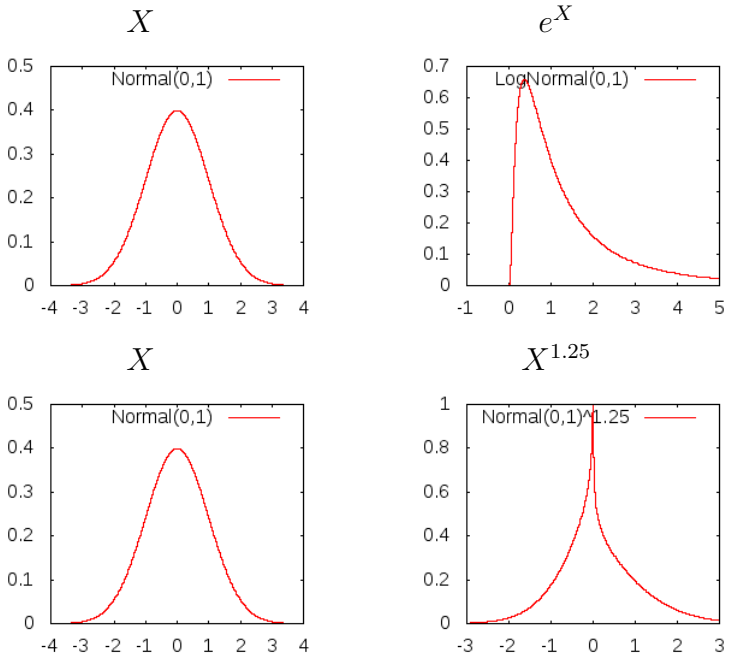


Table 3.1: Standard Normal X, e^X, $X^{1.25}$

For an example of RICO in action, let $X \sim N(0,1)$ numerically represented. New random variables such as e^X and $X^{1.25}$ can be created and graphed, see table 3.1. The code used to generate the graphs is

```
X = NormalNumeric(0,1,1000)
Plot(exp(X))
Plot(X**1.25)
```

Illustrative example of multiplication of two piecewise continuous random variables X and Y. Referring to figure 3.3, the following conditions hold for X and Y,

$$P(1 < X < 2) = 0.2 \qquad P(2 < X < 3) = 0.8$$
$$P(4 < Y < 5) = 0.3 \qquad P(5 < Y < 6) = 0.7$$

and the probability densities are super-imposed on the joint density distribution of (X, Y). For illustrative purposes a partition of the XY random variables is

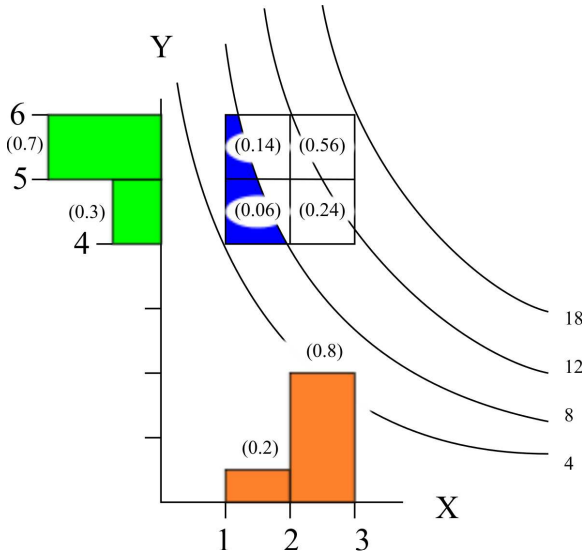


Figure 3.3: Piecewise Continuous X times Y

chosen as $(4, 8, 12, 18)$ to coincide with the corners of the non-zero probability density region of (X, Y). Each partition endpoint $\{4, 8, 12, 18\}$ corresponds to an iso-probability level curve identified in the figure 3.3. Notice that in the case of multiplication the level curves are hyperbolas. To compute the numeric random variable XY given the partition $(4, 8, 12, 18)$, calculate the area between level curves within joint probability rectangles. In particular, to compute $P(4 < XY < 8)$ one must find the fractional area of each of the two shaded rectangles multiplied by the probability contained in each such rectangle. The probability of the two shaded rectangles is $0.14 + 0.06 = 0.2$. To compute the probability of the shaded area using RICO involves forcing the particular partition shown in the figure as follows,

```
X = ContinuousNumeric((1,2,3),(0.2,0.8,0))
Y = ContinuousNumeric((4,5,6),(0.3,0.7,0))
Y.getNumericRandomVariable().force_partition((4,8,12,18))
X*Y
#()
#(4.0(0.11130904824004995)8.0(0.3565452412002493)
   12.0(0.5321457105597007)18.0,)
```

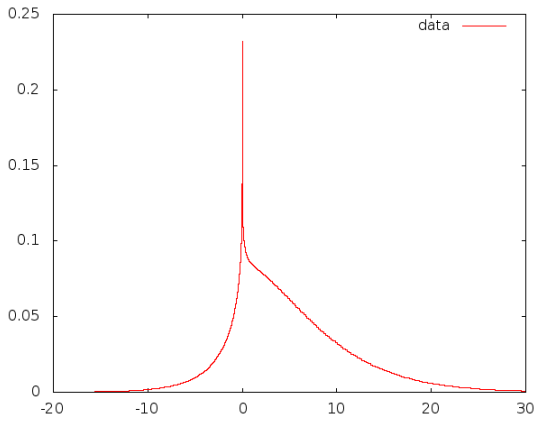


Figure 3.4: N(5,3) x N(1,1)

According to RICO the results of the above numerical example are

$$P(4 < XY < 8) \approx 11\% \quad P(8 < XY < 12) \approx 36\% \quad P(12 < XY < 18) \approx 53\%$$

A larger example of multiplication is XY where $X \sim N(5,3), Y \sim N(1,1)$. The listing follows as the plot is shown in figure 3.4.

```
X = Normal(5,3)
Y = Normal(1,1)
XY = X*Y
Plot().xrange(-20,30).plot(XY).show()
```

Similarly, division of two numeric random variables, this time in one line of code with the result shown in figure 3.5 which happens to be multi-modal.

```
Plot().xrange(-20,30).plot(Normal(7,3)/Normal(1,1)).show()
```

3.2 Correlation

Consider the example in figure 3.2 without preperatory remarks.

The issue in figure 3.2 is that RICO needs to understand that two references to the same underlying variable are 100% correlated. When RICO is asked to

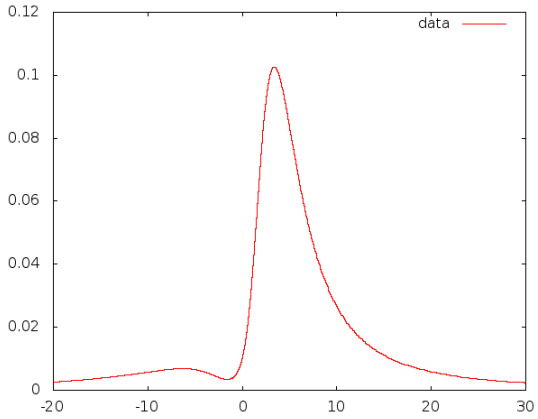


Figure 3.5: N(7,3) / N(1,1)

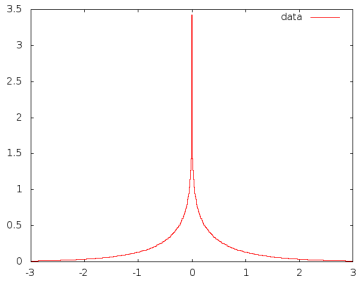
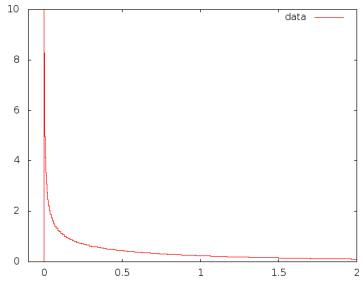
$X * X$	X^2
	

Table 3.2: Standard Normal $X * X$ versus X^2

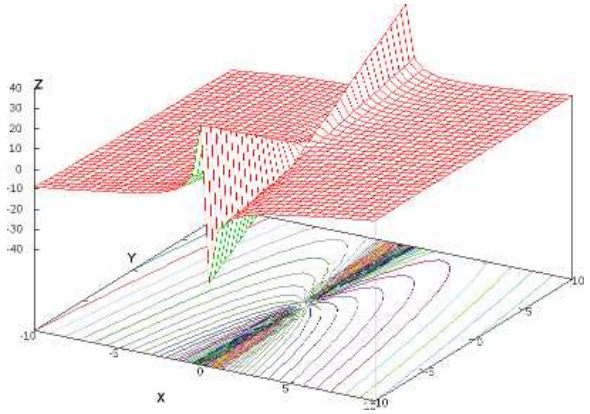


Figure 3.6: X + Y/X

compute X^2 there is no confusion, but $X * X$, without correlation tracking appears as the product of two independent copies of X.

Within the context of an algorithmic model with some random variable inputs intermediate variables are created and mixed with other intermediate variables such that partially correlated expressions are possible,

$$X + XY$$
$$X + Y/X$$

Assuming that $X \perp Y$ the multiplication XY is between independent random variables and is computable is detailed in the previous section. The sum in $X + XY$ is not between independent random variables. Fortunately RICO is equipped with a symbolic processing engine so that the expression $X + XY$ will be factored into $X(1 + Y)$. This expression is computable as a sequence of operations. The expression $1 + Y$ is independent of X so the product $X(1 + Y)$ is between independent random variables.

The expression $X + Y/X$ is not decomposible into a sequence of operations between independent random variables. In such a case the iso-probability level curves are more complicated than the hyperbolas found in the multiplication case. The level curves for a particular partition of $Z = X + Y/X$ are shown in figure 3.6.

In the case of $Z = X + Y/X$, the partitions of X and Y for a partition of the joint (X, Y) space. A subset of the iso-probability level curves associated with the partition of Z are approximated within each (X, Y)-partition rectangle bounded

by (x_0, y_0) and (x_1, y_1), relatively indexed, by approximating the supported probability surface f with a bilinear function of X and Y as

$$f(x, y) = axy + bx + cy + d$$

where the coefficients (a, b, c, d) are

$$\begin{pmatrix} a \\ b \\ c \\ d \end{pmatrix} = \begin{pmatrix} x_0y_0 & x_0 & y_0 & 1 \\ x_1y_0 & x_1 & y_0 & 1 \\ x_0y_1 & x_0 & y_1 & 1 \\ x_1y_1 & x_1 & y_1 & 1 \end{pmatrix}^{-1} \begin{pmatrix} f(x_0, y_0) \\ f(x_1, y_0) \\ f(x_0, y_1) \\ f(x_1, y_1) \end{pmatrix}$$

An iso-probability level curve is then a function $y(x|z)$ for a given partition endpoint z as

$$y(x|z) = \frac{z - d - bx}{ax + c}$$

A necessary step is to compute the fraction of each (X, Y)-partition rectangle intersected each given Z-partition interval. The bilinear approximation admits the following closed form solution,

$$\int y(x|z)\ dx = \frac{log(ax + x)(az - ad + bc) - abx}{a^2} + const$$

As a numerical example suppose X has a partition element bounded by $\{1, 3\}$ and similarly Y has a partition element bounded by $\{-10, 10\}$. If Z has a partition element bounded by $\{0, 1\}$ then figure 3.7 shows that appoximately 8.25% of the probability represented by this partition rectangle is allocated to the $(0, 1)$ partition element of Z. The resulting probability distribution for Z is shown in figure 3.8.

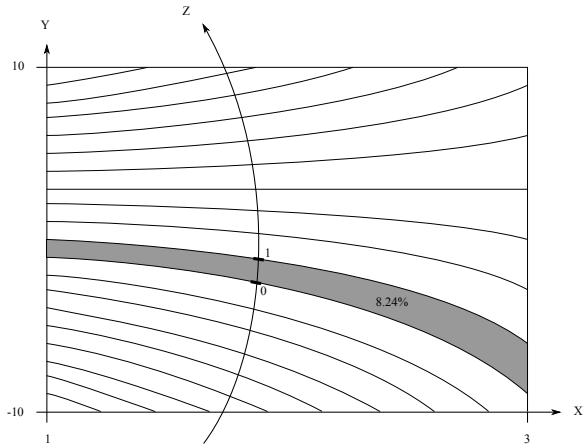


Figure 3.7: A partition element of X + Y/X with level curves

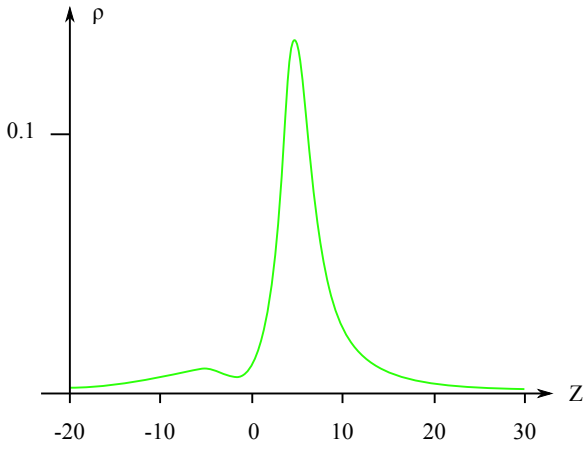


Figure 3.8: $X + Y/X$ where $X \sim N(5, 3), Y \sim N(1, 1)$

Chapter 4

Correlated Operations

Given a random variable A and two real-valued functions f and g such that $X = f(A)$ and $Y = g(A)$ let $Z = h(X, Y)$ where h is some real-value function from $\mathbb{R}^2$. The function $h(x, y) = x + y$ is of particular interest.

Since A may be a mixed random variable the development will by assuming A is discrete, then continuous and finally a general mixed random variable. In all cases X and Y are, by design, 100% correlated through A.

4.0.1 Discrete Operations on Correlated Random Variables

Suppose that,

$$A = ((a_1, a_2, ..., a_n), (p_1, p_2, ..., p_n))$$

where $Pr[A = a_i] = p_i$ for any $i \in 1...n$, that is, A is a discrete random variable. Consequently,

$$X = ((x_1, x_2, ..., x_n), (p_1, p_2, ..., p_n))$$
$$Y = ((y_1, y_2, ..., y_n), (p_1, p_2, ..., p_n))$$

where $x_i = f(a_i)$ and $y_i = g(a_i)$ for each i. Notice that duplicate values of x_i are possible. If it happens that $x_i < x_{i+1}$ for each $i \in 1...n-1$ then X is said to be in *proper form* and similarly for A and Y. To emphasize that X and Y are derived in a pointwise order-preserving manner they may be written in *synchronous* form,

$$X = (x_1, x_2, ..., x_n)$$
$$Y = (y_1, y_2, ..., y_n)$$

where the probability values associated with each x_i and y_i are found in A. The joint probability distribution of X and Y is itself a random variable called XY. Stated in syncrhonous form,

$$XY = ((x_1, y_1), (x_2, y_2), ..., (x_n, y_n))$$

A new random variable $Z = h(X, Y)$ is stated in synchronous form with respect to A as,

$$Z = (h(x_1, y_1), h(x_2, y_2), ..., h(x_n, y_n))$$

To restate Z in proper form requires two steps. The first is to remove duplicates form the range of Z,

$$\mathbf{R}(Z) = \{h(x_i, y_i)\}_{i \in 1..n}$$

The second step is to find the probability associated with each element of the range of Z. Assuming the following proper form of Z as,

$$Z = ((z_1, z_2, ..., z_m), (q_1, q_2, ..., q_m))$$

where $m \leq n$ then,

$$z_j \in \mathbf{R}(Z), \text{ ordered ascending}$$
$$q_j = \sum_{i | z_j = h(x_i, y_i)} p_i$$

For example suppose,

$$A = ((a_1, a_2, a_3), (p_1, p_2, p_3))$$
$$X = (1, 2, 3)$$
$$Y = (1, 3, 2)$$

The joint random variable XY in synchronous form is,

$$XY = ((1, 1), (2, 3), (3, 2))$$

Suppose $Z = h(X, Y)$ where $h(x, y) = x + y$. Then Z in synchronous form with respect to A is,

$$Z = (2, 5, 5)$$

To find the proper form of Z the range is first determined,

$$\mathbf{R}(Z) = \{2, 5\}$$

the proper form of Z is stated as,

$$Z = ((z_1, z_2), (q_1, q_2))$$

where

$$q_1 = p_1$$
$$q_2 = p_2 + p_3$$

since $Z = 5$ is the set $\{(x_2, y_2) = (2, 3), (x_3, y_3) = (3, 2)\}$. Notice that the process of finding Z in proper form is that of integrating iso-value subsets of the joint XY range, the domain of Z.

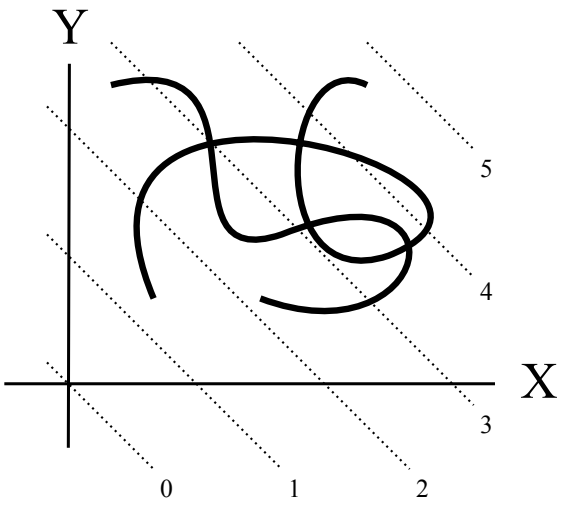


Figure 4.1: Distribution of Continuous XY

4.0.2 Continous Operations on Correlated Random Variables

Suppose A is a real-valued continous random variable. Let P be the probability density function associate with A and write $A \sim P$. The random variables $X = f(A)$ and $Y = g(A)$, in synchronous form, share this association with A, that is, $X\ P$ and $Y\ P$ for any real valued functions f and g. The joint random variable XY is also stated in synchronous form as $XY \sim P$. Finally, $Z = h(XY)$ for any real-valued function h from $\mathbb{R}^2$ is also stated in synchronous form as $Z \sim P$.

To form an interesting example consider that $X = f(A)$ and $Y = g(A)$ form a parametric curve in XY-space and that the probability density P is distributed along this curve. If $Z = h(XY)$ such that $h(x, y) = x + y$ then iso-value contours in XY-space appear as parallel lines with slope -1. In the figure 4.1 the disjoint curves are the single parametric (X, Y) curve and the dotted diagonal lines are the iso-value contours for addition of X and Y. The probability density P of A would appear in the figure perpendicular to the XY plane over the paremetric curve of (X, Y). From the figure it is apparent that $Pr(1 < Z < 5) = 1$. Notice that the iso-value contour labeled 5 in the figure intersects the XY curve at three places so that $Pr(Z = 4)$ is found as the sum of three probability density values in P.

Stating the synchronous form of Z with respect to A is as simple as for X and Y, that is, $Z \sim P$. Finding the proper form of Z may be a more challenging problem. The procedure for computing a numerical approximation to the proper form of Z is detailed in the dissertation [7]. Notice in particular that computing

the proper form of a random variable may be avoided until and observation is required for reasons such as graphing or comparison to unrelated random variables and constant values.

4.0.3 Mixed Discrete / Continuous Operations on Correlated Random Variables

To prepare for the computation of operations on a pair of mixed discrete/continuous random variables dependent on a single common random variable, it is useful to develop the case where one operand is discrete and the other is continuous. This situation can only arise if A is not a purely discrete random variable.

Suppose that A is a continuous random variable such that $A \sim P$ as above. Suppose without loss of generality that X is a discrete random variable and that Y is a continuous random variable. A visual example of a possible joint random variable XY appears in figure 4.2. Included in the figure are the iso-value contours used to compute $Z = h(XY)$ where $h(x, y) = x + y$ as the previous continous example. Notice that this case is not fundamentally different from the previous case where X and Y are both continous. In the figure X has four unique values in its range labeled x_1, x_2, x_3, x_4. It is apparent from the figure that $Pr(1 < Z < 5) = 1$ and that Z is a continuous random variable. The procedure for computing a numerical approximation to the proper form of Z is detailed in the dissertation [7].

4.0.4 Operations on Correlated Mixed Random Variables

If random variable A is *mixed*, that is, containing both continous and discrete probability distributions it is useful to decompose it into discrete and continous components and write,

$$A = A_d \oplus A_c$$

where A_d is a discrete random variable and A_c is a continuous random variable and the $'\oplus'$ operator performs a sum of distribution functions by converting discrete probability to Dirac Delta functions. The components of A are written as,

$$A_d = ((a_1, ..., a_n), (p_1, ..., p_n))$$
$$A_c \sim Q$$

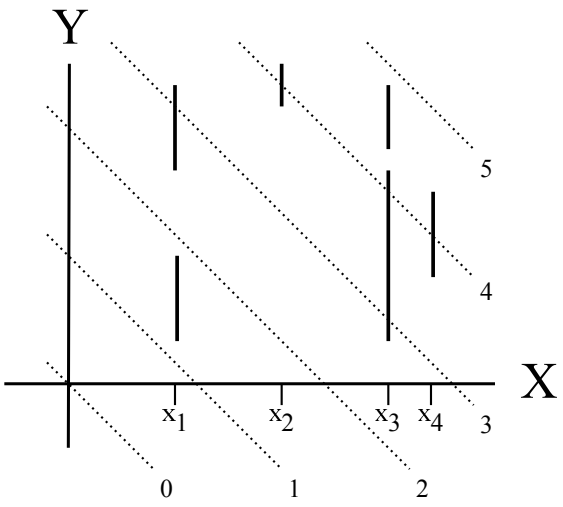


Figure 4.2: Distribution of Discrete/Continous XY

where Q is a conditional probability distribution represented by a continuous probability density function. Notice that if $d = Pr(A_d)$ then $Pr(A_c) = 1 - d$. That is,

$$d = \sum_{i=1..n} p_i$$
$$1 - d = \int_{A_c} dQ$$

where the abuse of integral notation implies that the integral is performed over the range of A_c in the usual sense. A non-trivial mixed random variable then requires that $0 < d < 1$.

Notice in particular that for special case of addition of correlated random variables the operation of addition as in $Z = X + Y$ is that of projecting the XY distribution to the diagonal as shown in figure 4.3.

As an example suppose,

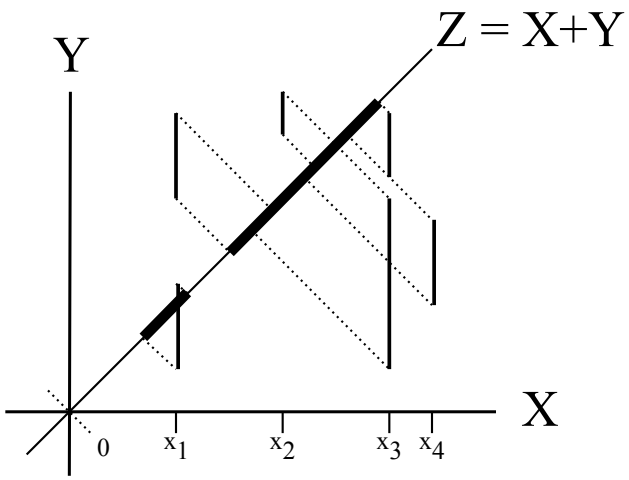


Figure 4.3: Projection of XY-space to $X + Y$-space

$$
\begin{aligned}
A &= \mathbf{U}(-1, 2) \\
f(x) = step(x) &= \begin{cases} 0 & \text{if x} \leq 0, \\ 1 & \text{else} \end{cases} \\
g(y) &= |y| + 1 \\
X &= f(A) \\
Y &= g(A)
\end{aligned}
$$

then in proper form,

$$
\begin{aligned}
X &= ((0, 1), (\frac{1}{3}, \frac{2}{3})) \\
Y &= \mathbf{U}((0, 1, 2), (\frac{2}{3}, \frac{1}{3}))
\end{aligned}
$$

where Y is *multi-uniform* requiring probabilities within partition elements to be specified. Notice that $Pr(X = 0) = \frac{1}{3}$, $Pr(X = 1) = \frac{2}{3}$, $Pr(0 < Y < 1) = \frac{2}{3}$ and $Pr(1 < Y < 2) = \frac{1}{3}$.

Suppose further that $Z = X + Y$. The joint XY figure 4.4 reveals the details. Noticing that the probability is uniformly distributed over the range of XY and

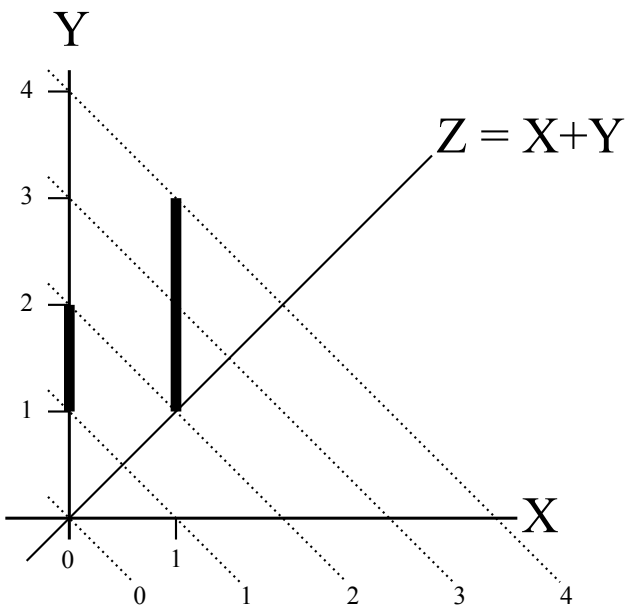


Figure 4.4: Example of Discrete/Continous XY

that the two fragments of that region do not overlap according to the iso-value contours of Z the problem is solved by inspection so that,

$$Z \sim \mathbf{U}(1, 4)$$

To find this result more formally Y is conditioned on the discrete X so that,

$$Y|X = 0 \sim \mathbb{U}(1, 2)$$
$$Y|X = 1 \sim \mathbb{U}(1, 3)$$

then,

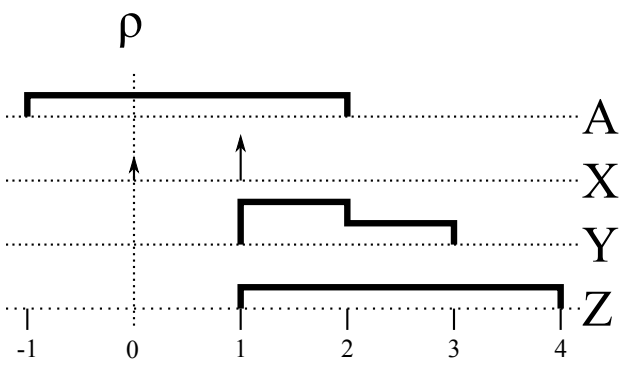


Figure 4.5: Example Discrete/Continous Distributions in Proper Form

$$
\begin{aligned}
Z &= X + Y \\
&= (X + Y | X = 0) \oplus (X + Y | X = 1) \\
&= (0 + Y | X = 0) \oplus (1 + Y | X = 1) \\
&\sim \mathbf{U}(1 + 0, 2 + 0) * Pr(X = 0) + \mathbf{U}(1 + 1, 3 + 1) * Pr(X = 1) \\
&\sim \mathbf{U}(1, 2) * \frac{1}{3} + \mathbf{U}(2, 4) * \frac{2}{3} \\
&\sim \mathbf{U}(1, 4)
\end{aligned}
$$

For visual convenience the proper form distributions of A, X, Y and Z are shown in figure 4.5. Notice that to compute $Z = X + Y$ from the proper forms of X and Y is more challenging than from the synchronous form of XY in figure 4.4 in part because of the otherwise unseen correlation between X and Y through A.

Chapter 5

Constrained Optimization

5.1 Tables and Chairs with Correlated Random Prices

In this final stage of the tables and chairs example we introduce correlated random prices. We follow economic practice by developing a small story around the problem to tie the elements together.

Suppose that a small furniture manufacturer in Portland, Oregon wants to forecast weekly revenue. The manufacturer makes tables and chairs in a small show with a small crew. Using a forecast for demand for tables, chairs and dinette sets the manufacturer derives the likely market prices for tables and chairs. A dinette set is composed of one table and two chairs.

Figure 5.1 shows the independent random variables corresponding to forecast demand for dinette sets (the exponential curve), tables (the tall Gaussian curve) and chairs (the wide Chi-Squared curve). The vertical axis represents probability density and the horizontal axis represents demand for units (in thousands) in the Portland market.

The manufacturer believes that market price and demand for tables are related by the inverse function,

$$P_t = \frac{14 * 80}{D_t + D_d}$$

where D_t is the demand for tables alone and D_d is the demand for dinette sets. Thus $D_t + D_d$ is the total demand for tables. Similarly, the price of chairs is related to the demand for chairs by the inverse function,

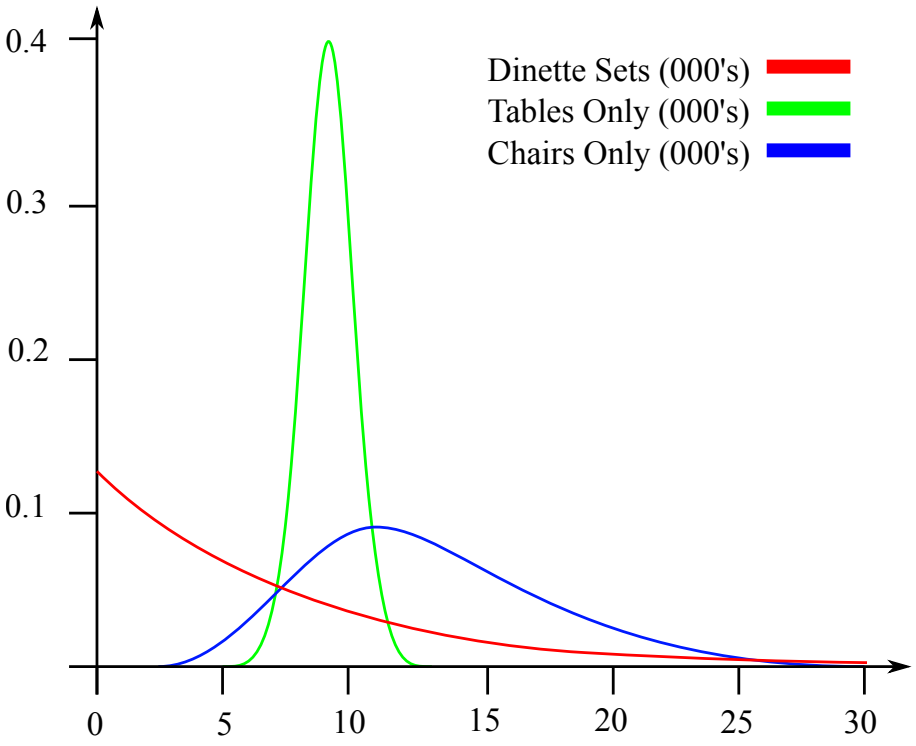


Figure 5.1: Tables, Chairs and Dinette Sets Random Variables

$$P_c = \frac{24 * 45}{D_c + 2D_d}$$

where D_c is the demand for chairs alone and again D_d is the demand for dinette sets. The sale of one dinette set implies the sale of two chairs. The actual functions are immaterial and have been contrived so that the results of this version of the tables and chairs example are comparable to previous versions.

We recognize that P_t and P_c are correlated, but do not need to materialize their joint probability distribution in order to compute revenue results.

The example is data-intensive so we create some prototype software to produce numerical results. Rather than presenting the prototype, written in Python using the Numpy library, we describe the data structures and sequence of operations.

Let our input random variables be,

$$Dt = \{DXt, DPt\}$$
$$Dc = \{DXc, DPc\}$$
$$Dd = \{DXd, DPd\}$$

where,

$$DXt = (DXt_1, \ldots, DXt_{Nt})$$
$$DPt = (DPt_1, \ldots, DPt_{Nt})$$
$$DXc = (DXc_1, \ldots, DXc_{Nc})$$
$$DPc = (DPc_1, \ldots, DPc_{Nc})$$
$$DXd = (DXd_1, \ldots, DXd_{Nd})$$
$$DPd = (DPd_1, \ldots, DPd_{Nd})$$

and we assume that $DPt_{Nt} = DPc_{Nc} = DPd_{Nd} = 0$ as usual for our numeric random variables since probability values are between partition values. We have Nt, Nc and Nd as the number of partition endpoints for each input random variable; tables, chairs and dinette sets respectively.

We form the demand joint probability distribution DP for the input random variables by Cartesian product,

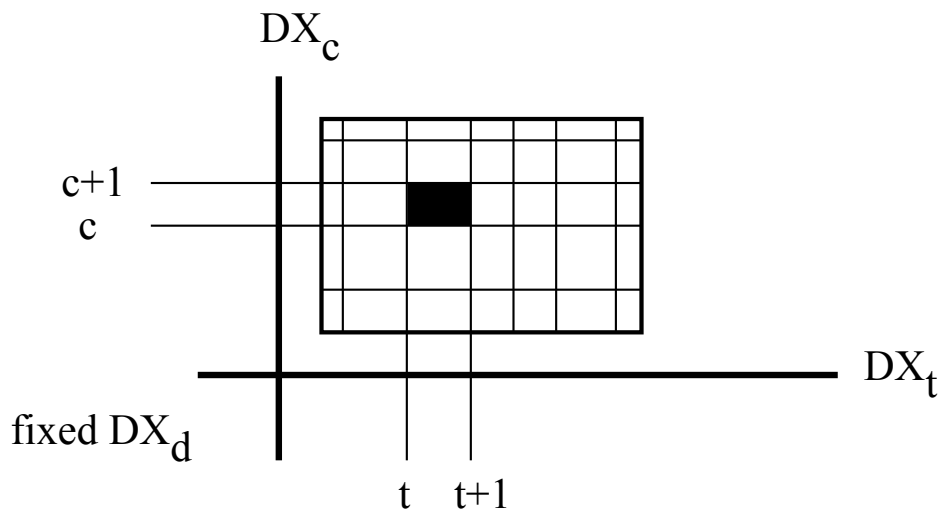


Figure 5.2: One Layer of Demand Probability Array

$$DP = DPt \times DPc \times DPd$$

We separately form parallel 3D arrays for each input demand,

$$\begin{aligned} DT &= DXt \times ones(Nc) \times ones(Nd) \\ DC &= ones(Nt) \times DXc \times ones(Nd) \\ DD &= ones(Nt) \times ones(Nc) \times DXd \end{aligned}$$

where, defining by example,

$$ones(5) = (1, 1, 1, 1, 1)$$

Throughout this example presentation we will tend to use 2D diagrams to represent 3D objects for clarity. In figure 5.2 we represent the demand joint probability DP for a fixed value of DXd.

In an abuse of notation, when no confusion arises, we use t, c and d as both identifiers and indices. We then have the indices for tables, chairs and dinette set random variables,

$$t = 1 \ldots Nt$$
$$c = 1 \ldots Nc$$
$$d = 1 \ldots Nd$$

Now we can refer to a specific point within the demand joint probability array as $DP_{t,c,d}$ or simply DP_{tcd} where the commas are dropped for clarity. The shaded rectangle in figure 5.2 is then the demand-space rectangular block of uniform probability distribution with value DP_{tcd}.

The four 3D arrays DT, DC, DD and DP are all parallel. We will now create other arrays parallel to these. The reason for this parallelism is to ensure that the probability within each block is correctly tracked through each step of the computation process.

The 3D arrays for the prices of tables and chairs are then written,

$$PT = \frac{14 * 80}{DT + DD}$$
$$PC = \frac{24 * 45}{DC + 2DD}$$

where the sums, $DT + DD$ and $DC + 2DD$, are computed element-wise as well as the reciprocal functions. The result is that PT and PC are 3D arrays of size $Nt * Nc * Nd$ and are parallel to the demand and demand probability arrays. Note in particular that the price arrays PT and PC are not random variables. We will describe below how they may be converted to random variable form. This is possible because we have values of prices as vertices of a block of probability distribution and we have the value of the probability uniformly distributed within that block in the 3D array DP.

Now that we have our input prices the manufacturer may apply their optimization and decide what combination of tables and chairs to produce. In this example we have already computed the optimization exhaustively and can therefore partition our demand space into the three output cases A, B and C.

Consider that each point in the 3D demand array represents a particular choice of input values that has associated with it two particular prices, one for tables and the other for chairs. We have already determined the rule for choosing each output case. We know, for example, that if $\frac{2}{3}Pt < Pc$ then output A will be selected and similar rules apply for outputs B and C. This means for each point in our 3D

demand we can assign a Boolean value 1 or 0 where 1 means that point has an associated price for chairs that is larger then two-thirds that of tables. We can thus create a parallel 3D array of Boolean values called a *mask* based on the optimized output selection rules. Let,

$$MA = \frac{2}{3}PT < PC$$
$$MB = \frac{1}{4}PT < PC < \frac{2}{3}PT$$
$$MC = PC < \frac{1}{4}PT$$

Each output is associated with some revenue. Using the price arrays we can form parallel revenue arrays. Let,

$$RA = 45PC$$
$$RB = 24PC + 14PT$$
$$RC = 20PT$$

To convert the revenue arrays, RA, RB and RC into random variables we must first find a partition. We notice that while the revenue arrays are parallel, using the masks we see that any given point in the array space indexed by (t, c, d) is intended to be present in exactly one revenue array. This is because the output A, B and C are mutually exclusive so, for example, the probability of producing \$1000 and \$2000 of revenue using output A can be added to the probability of producing this same range of revenue for output B and for C to arrive at a probability of producing that range of revenue regardless of output choice.

We would like the partition we use for the revenue random variable we are about to produce to span the range of possible revenue values, be fine where there is more revenue information and course where there is less and be so fine overall that numerical artifacts overwhelm the result. We chose for this example to use every $23^r d$ point from each 175-point input demand random variable and rerun the problem on the partition values alone, not the probability values. In the Python code this amounts to a single function call since all the code is in place for the main computation. The result is are smaller versions over the same revenue arrays representing collectively a sample of the possible revenue values this example model produces using the given demand inputs. The steps are as follows,

1. Form one dimensional arrays of valid revenue values for each output.
2. Run the same process as above to generate revenue arrays and output masks. Prepend an s to the name indicating they are small versions due to the reduced partition size.
3. Concatenate the three 1D arrays into a single array called $temp$.
4. Sort the $temporary$ array and remove any duplicates.
5. append the value $-\infty$ to the start of the array and ∞ to the end. Call the result Rx.

In this case the Python code from the prototype sums up the process concisely,

$$temp = concatenate((sRA[sMA], sRB[sMB], sRC[sMC]))$$
$$Rx = concatenate(([-\infty], unique(temp), [+\infty]))$$

where $sRA[sMA]$ returns a one dimensional array from an arbitrary array only for points where the corresponding point in the sMA small output mask array is a 1 and $unique()$ sorts and removes duplicates from an array.

For each (big) output array RA, RB and RC with associated masks MA, MB and MC we create a one dimensional array for the probability distribution that is parallel to the one dimensional partition array Rx.

$$Rap = zeros(Rx)$$
$$Rbp = zeros(Rx)$$
$$Rcp = zeros(Rx)$$

where, defining by example,

$$zeros(5) = (0, 0, 0, 0, 0)$$

The probability arrays, once filled in, will complete the formation of the output revenue random variables,

$$Ra = \{Rx, Rap\}$$
$$Rb = \{Rx, Rbp\}$$
$$Rc = \{Rx, Rcp\}$$

The three output random variables Ra, Rb and Rc are mutually exclusive and since they share a common partition we can add their probability values to find the final output revenue random variable R,

$$R = \{Rx, Rap + Rbp + Rcp\}$$

It remains to describe how to fill in the probability arrays Rap, Rbp and Rcp. We will describe the process for Rap since it is the same for the others.

Given the (big) output revenue 3D array RA, its associated mask MA, the associated 3D probability array DP and the 1D revenue partition Rx we proceed as follows to fill in the zero-valued 1D probability array Rap.

The output revenue 3D array RA together with the associated probability array DP describes a partition of the joint demand (input) space into blocks. Recall that we index the blocks with indices t,c and d so that the (t, c, d) block has uniform probability DP_{tcd} and eight vertices with the following revenues,

$RA_{t,c,d}$	$RA_{t,c,d+1}$	$RA_{t,c+1,d}$	$RA_{t,c+1,d+1}$
$RA_{t+1,c,d}$	$RA_{t+1,c,d+1}$	$RA_{t+1,c+1,d}$	$RA_{t+1,c+1,d+1}$

for some block such that $1 \leq t < Nt$, $1 \leq c < Nc$ and $1 \leq d < Nd$. If all the vertices are *valid*, that is, the associated mask value is 1 for each vertex then figure 5.3 symbolically represents one possible scenario.

The limits of the 3D block projection are the minimum and maximum revenue vertex values. That is,

$$min_{tcd} = Min(RA_{t,c,d}, \ldots, RA_{t+1,c+1,d+1})$$
$$max_{tcd} = Max(RA_{t,c,d}, \ldots, RA_{t+1,c+1,d+1})$$

For the software prototype version of this example we make the assumption that the 3D block probability $DP_{t,c,d}$ is distributed uniformly over the revenue line segment (min, max) so that the density is h_{tcd},

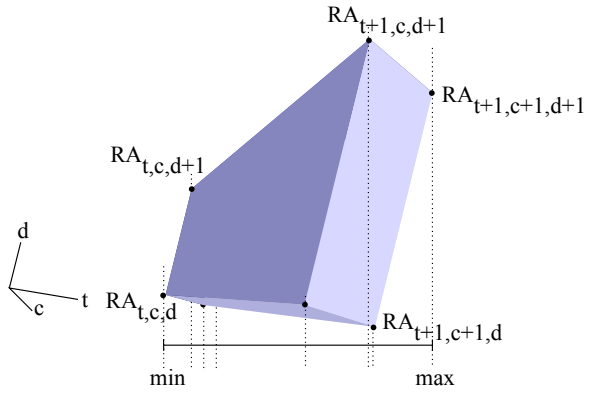


Figure 5.3: Line Projection of 3D Probability Block

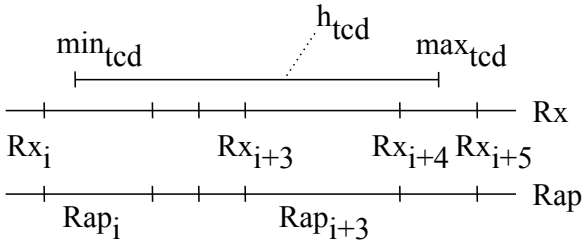


Figure 5.4: Partition Allocation of Probability Line

$$h_{tcd} = \frac{DP_{tcd}}{max_{tcd} - max_{tcd}}$$

where we have assumed that $min_{tcd} < max_{tcd}$. We will address special cases such as when min_{tcd} and max_{tcd} are equal below. Continuing with the general case we now allocate the uniform probability density h_{tcd} to the revenue probability array Rap recalling that Rap is delimited by the partition array Rx. Referring to figure 5.4 we have,

$$\begin{aligned} Rap_i &= Rap_i + (Rx_{i+1} - min_{tcd})h_{tcd} \\ &\dots \\ Rap_{i+3} &= Rap_{i+3} + (Rx_{i+4} - Rx_{i+3})h_{tcd} \\ Rap_{i+4} &= Rap_{i+4} + (max_{tcd} - Rx_{i+4})h_{tcd} \end{aligned}$$

where we have indicated how to compute end cases as well as cases where Rx partitions are spanned by the (min_{tcd}, max_{tcd}) interval.

If the mask MA indicates that some of the vertices of the (t, c, d) block are not valid then we must reduce the amount of block probability DP_{tcd} available for allocation. For example, if 3 of 8 vertices are valid for the (t, c, d) block then the block probability is correspondingly reduced to $\frac{3}{8}DP_{tcd}$ so that the (min_{tcd}, max_{tcd}) interval probability density is,

$$h_{tcd} = \frac{3}{8}\frac{DP_{tcd}}{max_{tcd} - max_{tcd}}$$

If it happened that $min_{tcd} = max_{tcd}$ either because all the valid vertices have the same revenue value or there is only one valid vertex for the (t, c, d) block then the corresponding partition element is located for the Rap array and its value is incremented with the available probability for that block. If it happens that $min_{tcd} = max_{tcd}$ equals an Rx partition endpoint then the available block probability is halved and allocated to the adjacent partition intervals.

We perform the above operations for each output case and combine them into the full revenue random variable and show the result in figure 5.5. The horizontal axis is dollars of revenue and vertical axis is probability density as usual for random variable graphs. We notice that the median value is roughly $2200 because of careful choice of demand inputs and the demand-to-price functional relationship.

Notable features of the optimized revenue in any panel of figure 5.5 is that no matter what happens with the projected demand there is a non-zero minimum revenue (about $300), a strongly likelihood of earning about $2200 and significant possibility of earning considerably more than the median $2200.

We use the machinery developed above to convert the 3D price arrays to random variables. Since there is no optimization involved in computing prices there is no need to generate masks. Since we have some information about the range of prices to expect we choose price partitions directly. In this case each price is partitioned into regularly spaced intervals from 0 to $150. The 3D probability arrays are projected onto the price partitions and the result for each price random variable is shown in figure 5.6. Again we notice that, by design, the price of chairs is has a median price of about $45 and that of tables is about $80 which corresponds with the sharp version of the tables and chairs example.

Notice that the price random variables are marginal probability distributions for a joint probability distribution we have not computed. Since the price for tables and chairs are non-trivially correlated the joint distribution cannot be recovered

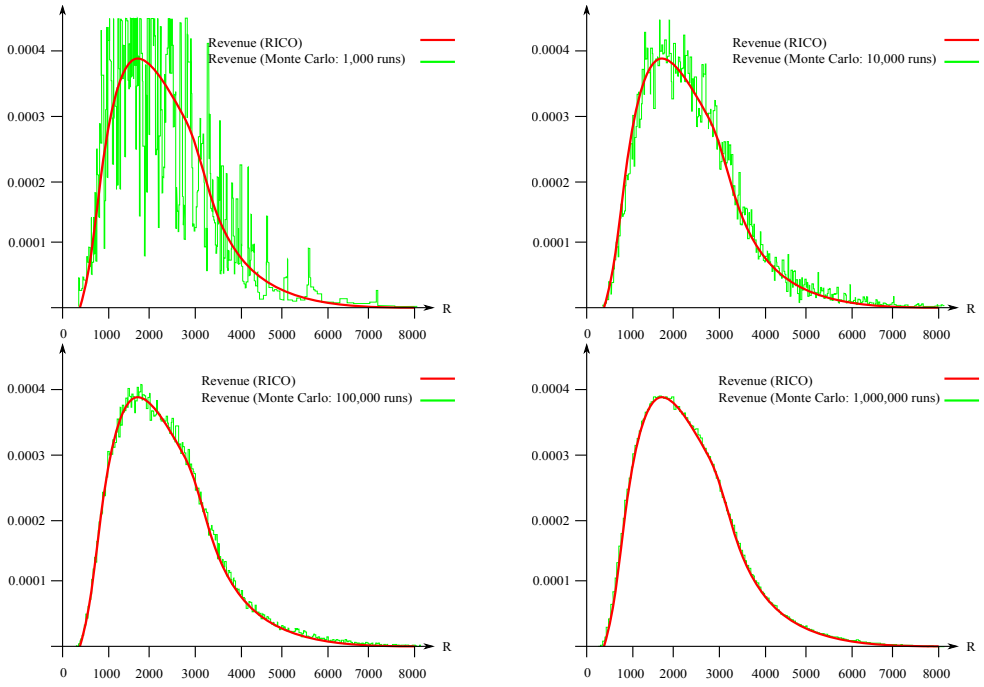


Figure 5.5: Random Revenue from Sales Combinations

from the marginal distributions alone as described in a standard statistics textbook such as [2].

5.1.1 Finding the Joint Price Distribution from the Demand Inputs

The reader will notice that we developed the tables and chairs example with unknown prices in preparation for the introduction of random inputs resulting in correlated prices we described how to proceed with a joint distribution for the two prices, tables and chairs. Then we solved the problem without using, or even finding, the joint price distribution. Instead we used the 3D array created to represent the three demand inputs. For this problem this technique provides directed and satisfactory results.

In this section we revisit the tables and chairs example with the same three demand input, but this time produce the joint price probability distribution. Since so much of this work is devoted to the study of correlated random variables we would be remiss not to include at least one example of same.

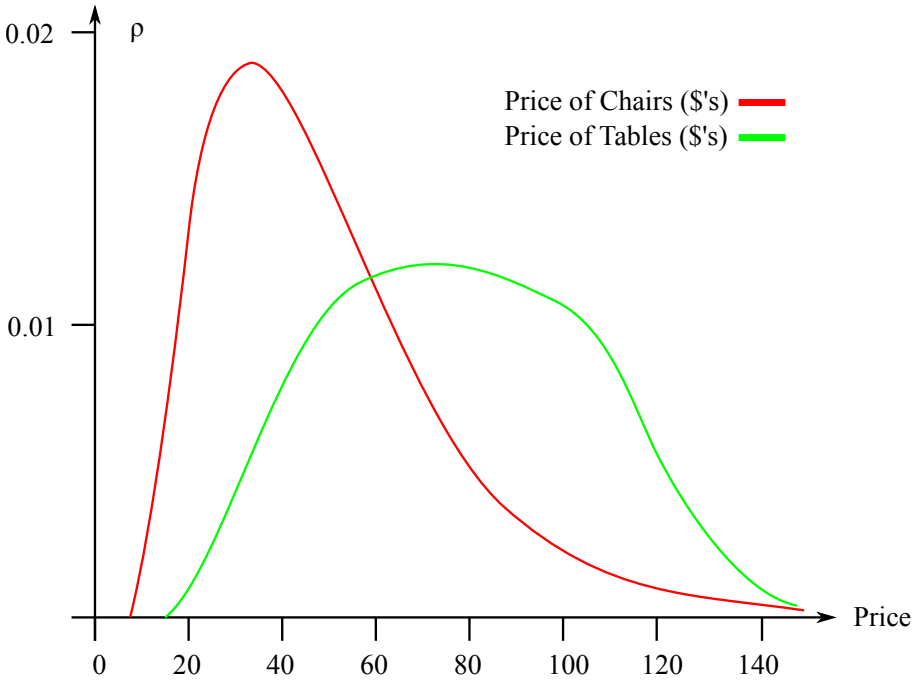


Figure 5.6: Random Variable Table and Chair Prices

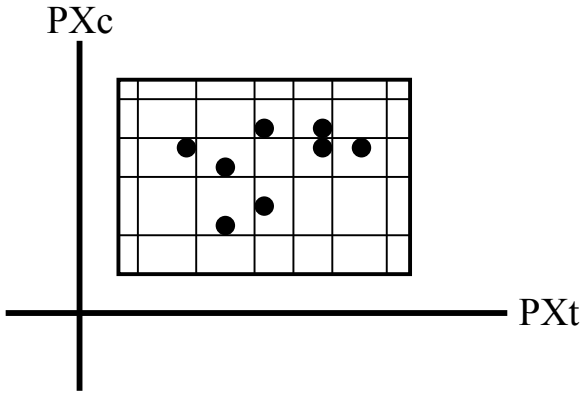


Figure 5.7: Joint Price Partition with Block Vertex Projections

We restart the problem with our three 3D demand arrays, DT, DC and DD. We also have the associated 3D probability array DP. Using the same formulas as before for finding the (correlated) prices of tables and chairs we produce the two 3D price arrays PT and PC respectively.

This brings us to the point where we projected each 3D price array onto a price partition and produced random variable representations of the two prices in figure 5.6. We choose the same price partitions as before, evenly space intervals from \$0 to \$150. This choice allows us to compare the results we are about to obtain with those obtained previously.

Our two price partitions, for PT and PC, describe a 2D partition of the (Pc, Pt)-space. If the number of points in each price partition is Np then we create a 2D array of size Np^2 and initialize it with zero values.

We then realize that our two 3D price arrays PT and PC together describe a 3D lattice of pairs of prices at each vertex surrounding a uniform distribution of probability described by the 3D probability array DP. The 8 vertices of each probability block, each containing the two price values, are projected onto the two dimensional (Pc, Pt)-space. This the 2D analog of our 1D procedure for finding each marginal price random variable by projecting each price block for PT or PC onto the corresponding one dimensional price line. Figure 5.7 shows an example of price block vertex projection.

Since the cluster of vertex projections in figure 5.7 indicate the projection of the 3D price probability block onto the 2D joint price partition we must allocate the block probability accordingly.

Assuming that the cluster of vertex projections represent the limits of the block

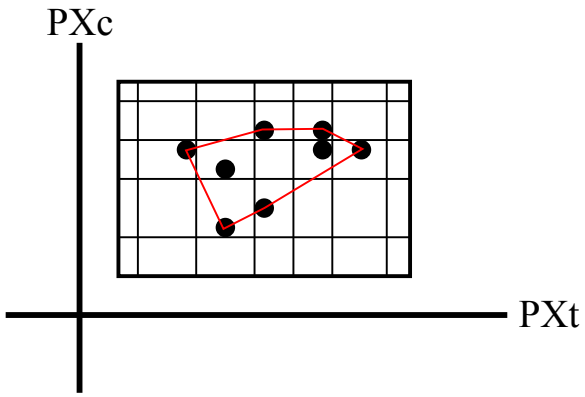


Figure 5.8: Joint Price Partition with Convex Block Projection

projection we can find the convex hull of these points using an algorithm such as the Graham Scan as described in a textbook of computer algorithms such as Corman [5]. We then assume the block probability is distributed uniformly over the interior region of the convex hull and apportion it accordingly to the partition rectangles of the 2D price distribution, called JP. Figure 5.8 shows the convex hull of the projected vertices. The heavy outline of joint price rectangles shows the limits of affected rectangles. Let p be the probability of the projected block and a the area of the convex region, then $h = p/a$ is the probability density. The portion of probability allocated to any given rectangle in the outlined region is h times the area of the rectangle intersecting the convex region.

In our running tables and chairs example we have over 5 million blocks to project so we opt not to engage in a complex computation of multiple rectangle intersections with convex regions associated with each block, at least not for our prototype code. Instead we take a simpler approach as shown in figure 5.9. The heavy outline bounding box represents the Pt and Pc partition limits bounding the block vertex cluster. The shaded inner rectangle represents the rectangular limits of the cluster points. We calculate the probability density of the block probability if distributed uniformly over the inner shaded rectangle and distribute this by area over each intersecting price rectangle.

The results of the calculations of our prototype code for the joint probability distribution of the two correlated prices is shown in figure 5.10. Be aware that the origin is located in the upper-left corner of the graph. The x and y axis are prices of tables and chairs respectively and the vertical axis is probability density.

An top view of the joint probability price distribution is shown figure 5.11. We

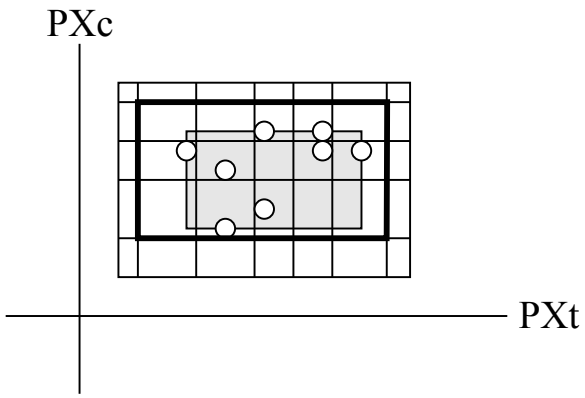


Figure 5.9: Joint Price Partition with Rectangular Block Projection

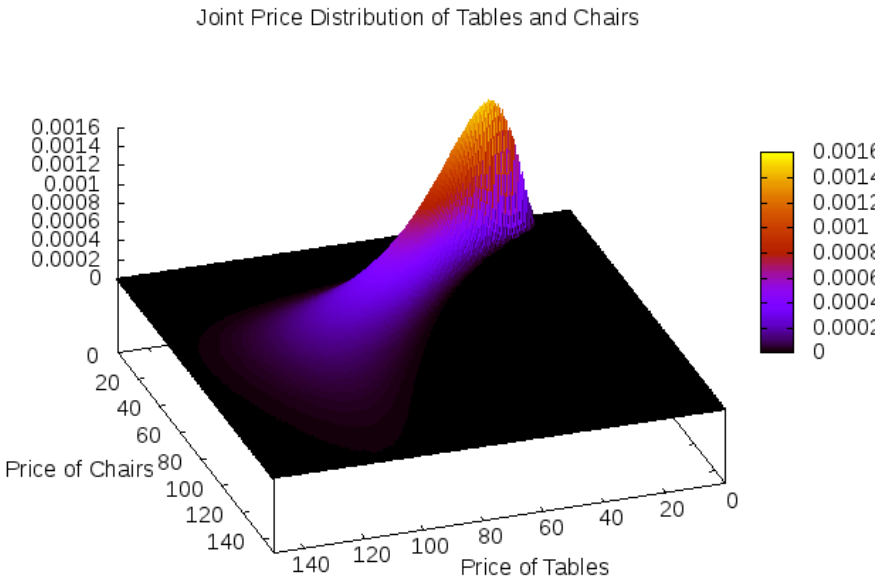


Figure 5.10: Joint Probability Distribution of Table and Chair Prices

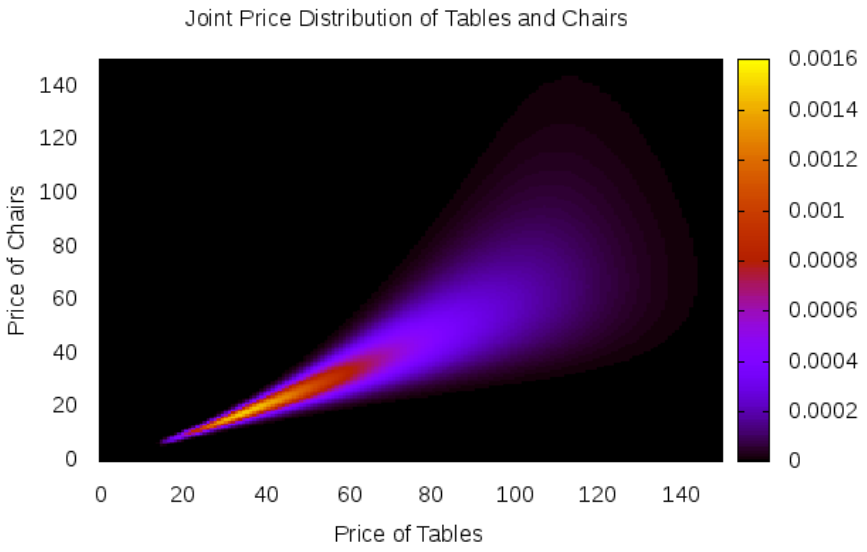


Figure 5.11: Joint Probability Distribution of Table and Chair Prices (Top View)

compare this figure to our original suggestion of the joint probability distribution of prices in figure B.3.

It is interesting to note that the marginal random variable prices we compute with the prototype code are identical to the random variable price we computed previously and show in figure 5.6.

With the joint distribution of prices in hand we are able to address the question of how to compute the probability of each branch in the simplex graph (see figure B.2). In particular we are interested in the probabilities of the branch conditional expressions,

$$
\begin{aligned}
p_t &< p_c \\
\frac{2}{3}p_t &< p_c \\
\frac{2}{3}p_t < p_c &< p_t \\
\frac{1}{4}p_t &< p_c \\
\frac{1}{4}p_t < p_c &< \frac{2}{3}p_t
\end{aligned}
$$

Proceeding as we did from the beginning, but starting with the jointly distributed prices and their partitions P_t and P_c we create 2D arrays PT_2 and PT_2 that are parallel to the 2D joint probability distribution array JP. Taking the last inequality above as an example we divide each expression by p_t to find,

$$\frac{1}{4} < \frac{p_c}{p_t} < \frac{2}{3}$$

In the prototype we form the 2D array expression Qtc as,

$$Qtc = \frac{PC_2}{PT_2}$$

as the element-wise quotient of the two 2D price arrays. We notice that Qtc together with the JP form an improper form two-dimensional random variable, a non-standard usage of the expression. If we choose a 1D partition for Qtc we can project our $\{Qtc, JP\}$ pair onto this partition and find a proper-form random variable, called qtc. As we have mentioned earlier, we address this case by choosing the special partition,

$$Xqtc = (-\infty, 0, \frac{1}{4}, \frac{2}{3}, 1, \infty)$$

The result is,

$$Pqtc = (0, 0.0001012, 0.7206, 0.2388, 0.03024, 0)$$

0011	x_c	x_t	s_W	s_L	b
s_W	5	20	1	0	b_W
s_L	10	15	0	1	b_L
Revenue	p_c	p_t	0	0	

Table 5.1: Tables and Chairs Simplex Tableau for Unknown Prices and Resources

1001	x_c	x_t	s_W	s_L	b
s_W	1	4	$\frac{1}{5}$	0	$\frac{b_W}{5}$
s_L	0	-25	-2	1	$b_L - 2b_W$
Revenue	p_c	p_t	0	0	

Table 5.2: Tableau for Unknown Prices and Resources, State 1001

Combining these into the random variable Q for convenience as,

$$Q = \{Xqtc, Pqtc\}$$

These probability values tell us probability of each simplex directed graph branch and therefore the probability of each result. For example, referring to figure B.2, the probability of taking the first left directed edge under the condition that $p_t < p_c$ is $\mathbb{P}(1 < Q) = 0.03024$. Similarly the probability of reaching result B is $\mathbb{P}(\frac{1}{4} < Q < \frac{2}{3}) = 0.7206$.

5.1.2 Tables and Chairs with Unknowns Prices and Resources

If we allow prices and resources to be described by correlated random variables the impact on the example is to increase the number of branches from each simplex algorithm states and an increase in the number of states. The simplex tableau for unknown prices and resources is shown in table 5.1.

The directed graph for the tables and chairs example with unknown prices and resources is shown in figure 5.12. Notice that there are only $\mathbb{C}(4, 2) = 6$ possible node states in this example. Notice also that there are 5 possible terminal states; manufacture of only tables or only chair limited by either wood resource or labor resource and also the mixed case.

While the conditions present when entering a state are significant, the tableau in each state is denumerable as. We refer to tables 5.2, 5.3, 5.4, 5.5 and 5.6.

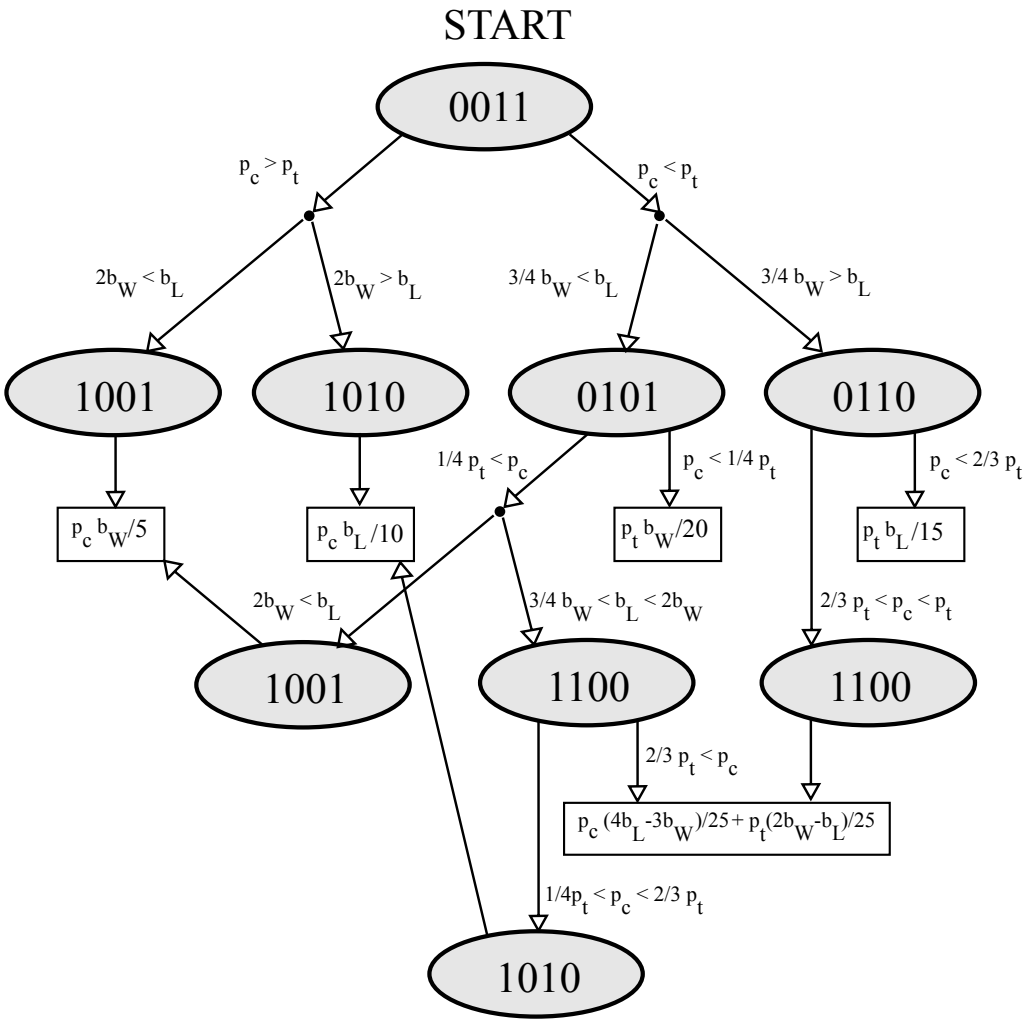


Figure 5.12: Directed Graph for Tables and Chairs with Unknown Prices and Resources

1010	x_c	x_t	s_W	s_L	b
s_W	1	$\frac{3}{2}$	0	$\frac{1}{10}$	$\frac{b_L}{10}$
s_L	0	-25	1	$-\frac{1}{2}$	$b_W - \frac{b_L}{2}$
Revenue	p_c	p_t	0	0	

Table 5.3: Tableau for Unknown Prices and Resources, State 1010

0101	x_c	x_t	s_W	s_L	b
s_W	$\frac{1}{4}$	1	$\frac{1}{20}$	0	$\frac{b_W}{20}$
s_L	$\frac{25}{4}$	0	-$\frac{3}{4}$	1	$b_L - \frac{3}{4}b_W$
Revenue	p_c	p_t	0	0	

Table 5.4: Tableau for Unknown Prices and Resources, State 0101

0110	x_c	x_t	s_W	s_L	b
s_W	$\frac{2}{3}$	1	0	$\frac{2}{30}$	$\frac{b_L}{15}$
s_L	-$\frac{25}{3}$	0	1	-$\frac{4}{3}$	$b_W - \frac{4}{3}b_L$
Revenue	p_c	p_t	0	0	

Table 5.5: Tableau for Unknown Prices and Resources, State 0110

Starting in state 0011 we follow the simplex two-phase decision and first compare the two prices p_c and p_t. We are assuming for clarity, as before, that since we intend to replace p_c and p_t with continuous random variables the probability of equality is zero. In practice we need to check for the possibility of equality. The initial comparisons are,

$$\begin{aligned} & argmax(p_c, p_t) \\ argmin(\frac{b_W}{5}, \frac{b_L}{10}| > 0 \text{ and } & p_c > p_t) \\ argmin(\frac{b_W}{20}, \frac{b_L}{15}| > 0 \text{ and } & p_c < p_t) \end{aligned}$$

where the > 0 condition refers to the requirement that each operand be positive else it is disqualified from the comparison.

In many cases the first or second phase of the simplex algorithm decision is disqualified since it is either non-positive or contradicts a previous assumption.

Since simplex states may be re-entered a computer algorithm can take advan-

1100	x_c	x_t	s_W	s_L	b
s_W	1	0	-$\frac{3}{25}$	$\frac{4}{25}$	$\frac{4b_L - 3b_W}{25}$
s_L	0	1	$\frac{2}{25}$	-$\frac{1}{25}$	$\frac{2b_W - b_L}{25}$
Revenue	p_c	p_t	0	0	

Table 5.6: Tableau for Unknown Prices and Resources, State 1100

tage of this possibility and cache, rather than recompute, certain elements such as the simplex tableau. At each decision point we see again that we are comparing linear combinations of either price or resource variables with zero in the sense that the expression $a < b$ can be rewritten as $0 < b - a$. We have seen in the previous version of the tables and chairs example that each conditional statement results in a filter on the input space so that the simplex algorithm may be viewed as a filtration process. The task of the RICO modeling environment is to determine what portion of the input space passes through each facet of the simplex filtration process to a terminal node and with what probability.

5.2 Beyond the Tables and Chairs Example

In the tables and chairs example we chose correlated random inputs for prices and argued that we could have chosen correlated random inputs for resources instead. We notice in the simplex algorithm the transition decision from one state to the other involves computing the maximum positive revenue impact in the case of prices and then the minimum positive resource impact. These two choices correspond to the two facets of a table pivot as explained above.

It remains to be investigated what happens to this example when some or all of the values of A are unknown. In this example the significance of unknown A values is that the manufacturer is unsure how many resources are consumed by each product.

As explained by Bellman [1], the number of solution states, not to mention the number of internal simplex algorithm states, becomes computationally intractable even for modest problems. We see this for ourselves if the problem has 100 variables and 100 constraints then the number of simplex states is at least $\mathbb{C}\approx(200, 100) = 9.05 \times 10^{58}$. The reference implementation of the AB32 model has many hundreds of variables and several thousand constraint equations.

A way to proceed is for the RICO modeling environment to partially explore the simplex directed graph. The transition from one state to the next using the simplex algorithm involves finding the maximum of a set of linear combinations of prices, in the context of the tables and chairs example, followed by finding the minimum of a set of linear combinations of resource limits assuming the A values are fixed. As we have seem each *choice element* is a linear combination of random variables which are themselves random variables. If we assume there are three choice elements denoted X, Y and Z then the decision,

$$argmax\{X, Y, Z\}$$

results in three probability values,

$$\mathbb{P}(X < \{Y, Z\})$$
$$\mathbb{P}(Y < \{X, Z\})$$
$$\mathbb{P}(Z < \{X, Y\})$$

In this case the simplex algorithm state has three initial branches corresponding to the first decision of the pivot element. These probabilities may be computed explicitly and rather than create a directed graph with all choices listed as we did above with the tables and chairs directed graph, we choose only the most highest probability transition. In this manner we reach a terminal node as in the sharp version of the simplex algorithm.

Since we are able to assign probability values to each transition edge in the directed graph we may apply a choice algorithm to explore other paths based on their likelihood of occurrence. As long as the directed graph remains at least partly unexplored, we suspect this is the case in general, then the random variable results will not have *full probability*, that is, their probability values will sum to less than one. The proximity of the probability sum of a random variable result to unity can be used as a criterion for algorithm termination. That is, if the random variable result is deemed near enough to completion the algorithm can terminate its exploration of the simplex directed graph.

Chapter 6

Black Scholes Construction

6.1 Black Scholes Construction

Consider a financial security such as a stock S with current market price S_0. At some future time T let the price of S be represented by a random variable S_T whose probability distribution is indicated in figure 6.1. Here, no presumption is yet made about the kind of distribution for S_T.

Following Dineen [6] suppose now that S_T satisfies the so-called *no-arbitrage requirement* $S_0 = \mathbb{E}[S_T]$. Suppose an investor, Ivan, holds one unit of stock S as the sole content of his portfolio of assets. Initially his wealth W is then just,

$$W_0 = S_0$$

and at time T his wealth is expressed as,

$$W_T = S_T$$

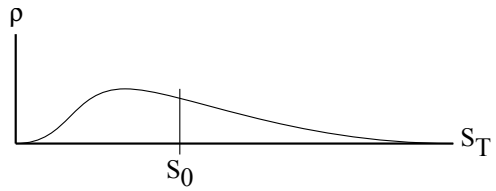


Figure 6.1: Distribution of S_T

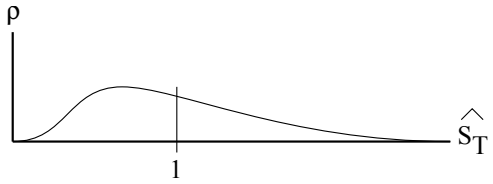


Figure 6.2: Distribution of $\hat{S}_T$

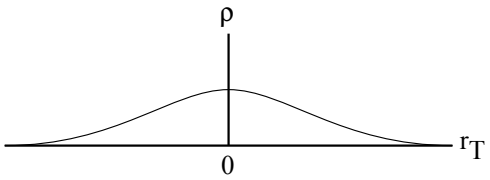


Figure 6.3: Distribution of r_T

Expressing Ivans future wealth in terms of his current wealth and a continuously compounding growth rate r one writes

$$W_T = W_0 e^{r_T}$$

where

$$r_T = log(\frac{S_T}{S_0}).$$

Referring to $\hat{S}_T = S_T/S_0$ as the ***normalized*** stock price rescale the probability distribution S_T to $\hat{S}_T$ so that the mean of $\hat{S}_T$ is one as depicted in figure 6.2.

Since r_T is the log of $\hat{S}_T$, its mean is zero as illustrated in figure 6.3. Notice that if $\hat{S}_T$ is represented by a LogNormal random variable then r_T is represented by a Normal random variable with zero mean.

If instead of holding a unit of stock S, the investor Ivan holds a European-style call option C based on stock S with strike price K exercisable at time T. The value of C at time T is given by

$$C_T = [S_T - K]^+$$

similarly, the related put-option P with the same parameters as S has value P_T at time T with

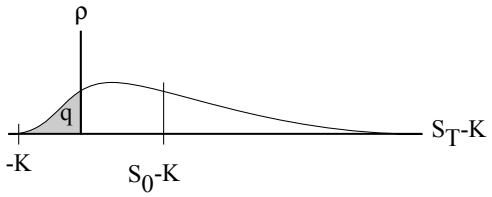


Figure 6.4: Distribution of $S_T - K$

$$P_T = [S_T - K]^-$$

It is helpful to notice that $C_T, P_T \geq 0$. For ease of calculation, the assumption is made that upon maturity all options are settled in cash and that there is no requirement that the *underlying* stock be transferred between parties. A further assumption made here is that all cash values are stated in current (t = 0) dollars. Please note that this runs counter to the not uncommon practice in financial texts, e.g. [6] to include the effect of time-value-of-money in calculations. While the current dollars assumption is equivalent to zero interest rate, the understanding is that all cash may be restated in any chosen year denomination.

To find the probability distribution over the mature value C_T of the call option C at time T with the underlying stock S and strike price K, a sequence of transformations of S_T are necessary. The first transformation is to subtract the strike price K from the distribution S_T. The resulting distribution is illustrated in figure 6.4. The key feature is represented by the shaded region in figure 6.4 is the probability q, such that $q = Pr\{S_T - K \leq 0\}$.

The second transformation, the probability distribution of $[S_T - K]^+$ is illustrated in figure 6.5. Notice that the probability q is now concentrated discretely at zero and represented by a vertical arrow labeled q. Since the remaining portion of the distribution is continuously distributed, the vertical axis still represents probability density. Notice further that the mean μ_{C_T} of this mixed, continuous/discrete probability distribution, is indicated at some positive location in figures. Notice that as $C_T = [S_T - K]^+$,

$$\mu_{C_T} = \mathbb{E}[[S_T - K]^+] = \mathbb{E}[C_T]$$

The significance of μ_{C_T} in figure 6.5 is that this price, based on the no-arbitrage principle, an investor can expect to pay today $(t = 0)$ for a call of this

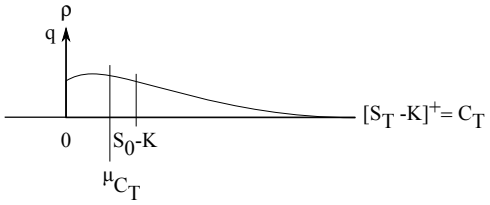


Figure 6.5: Distribution of $[S_T - K]^+$

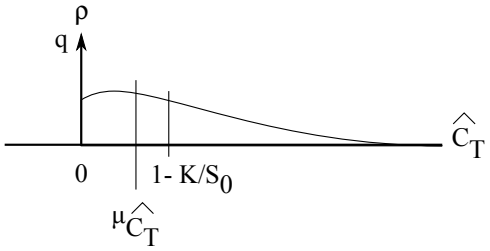


Figure 6.6: Distribution of $\hat{C_T}$

type when S_T represents the commonly accepted probability distribution of the value of stock S at time T. That is,

$$C_0 = \mu_{C_T} = \mathbb{E}[C_T]$$

Normalizing the distribution of the mature call option with respect to the initial value of the underlying S_0 gives the distribution for $\hat{C_T}$, the normalized call option price, illustrated in figure 6.6.

Using the same strike price K and underlying stock S, a put option value at maturity is illustrated in figure 6.7 where value at maturity of the put option is,

$$P_T = \mathbb{E}[[S_T - K]^-]$$

and the no-arbitrage price of the put is,

$$P_0 = \mu_{P_T}$$

Note that the probability concentrated at zero here is $1-q = Pr\{S_T - K < 0\}$. The *put-call parity formula,*

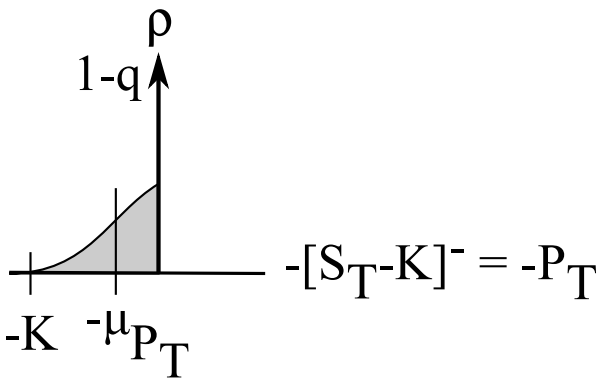


Figure 6.7: Distribution of $-[S_T - K]^-$

$$\mathbb{E}[C_T] - \mathbb{E}[P_T] = S_0 - K$$

is now obtained easily using the notation stated above, noting that

$$\begin{aligned}
\mathbb{E}[C_T] - \mathbb{E}[P_T] &= S_0 - K \\
\mathbb{E}[C_T - P_T] &= S_0 - K \\
\mathbb{E}[[S_T - K]^+ - [S_T - K]^-] &= S_0 - K \\
\mathbb{E}[S_T - K] &= S_0 - K \\
\mathbb{E}[S_T] - K &= S_0 - K \\
S_0 - K &= S_0 - K
\end{aligned}$$

as in Dineen [6], with the exception that Dineen [6] denotes as C_T the expected value of the probability distribution denoted C_T in this paper and similarly for P_T.

There is a geometric interpretation of Put-Call Parity. This can be seen in the perspective figure 6.8, where the probability density axis is perpendicular to the $C_T \times (P_T)$-space. PutCall Parity is realized in that figure as the elementary projection of the joint $C_T \times (PT)$-space to the diagonal $C_T + (P_T)$-space. The analogous algebraic calculation procedure is detailed in the appendix.

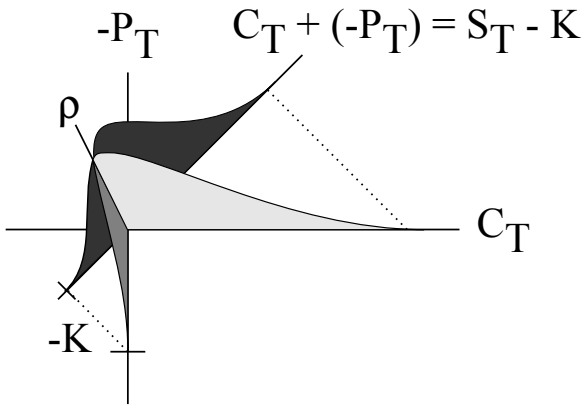


Figure 6.8: Distribution of $C_T - P_T = S_T - K$ in Perspective

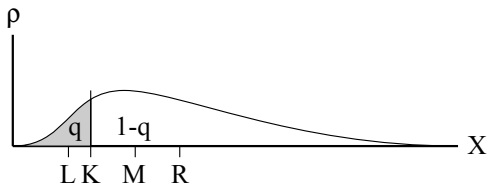


Figure 6.9: Probability Distribution Partitioned at K

6.1.1 Geometric Black-Scholes Pricing

Suppose now that a random variable X has probability density function $\rho(x)$ and is "well behaved" enough to have mean $M = \mathbb{E}[X] < \infty$. Suppose further we have a partition point K as shown in figure 6.9 and let $q = Pr\{X < K\}$ and so $1 - q = Pr\{X \geq K\}$.

Form two new *truncated* random variables from the left and right sections of the K-partitioned random variable X with probability densities,

$$X_L \sim \rho_L(x) = \frac{\rho(x)\mathbf{1}_{X<K}}{q}$$
$$X_R \sim \rho_R(x) = \frac{\rho(x)\mathbf{1}_{X\geq K}}{1-q}$$

as depicted in figure 6.10 where the L and R are expected values of X_L and X_R. Notice that the affine combination of L and R via q, illustrated in figure 6.11, yields the original mean M,

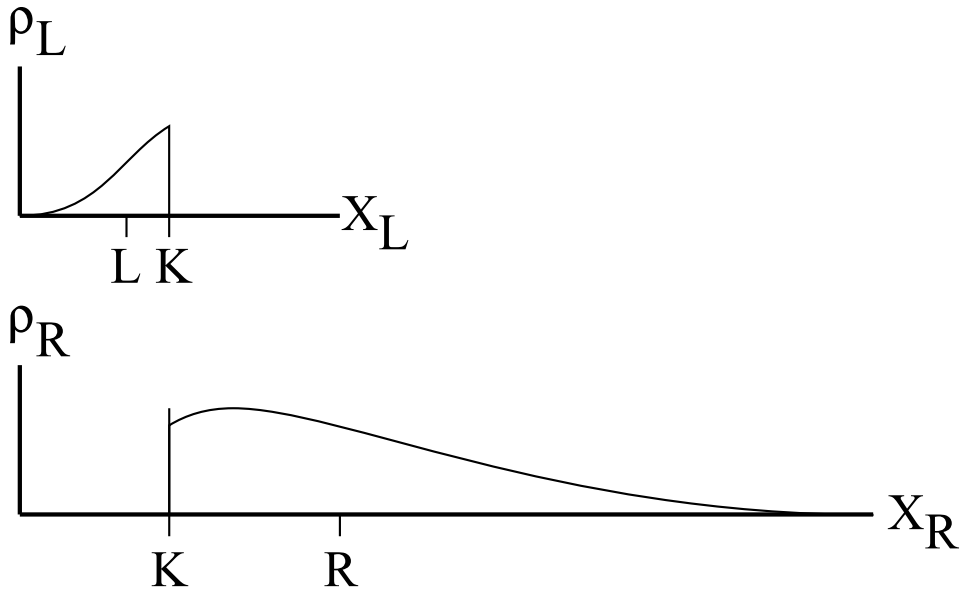


Figure 6.10: Left- and Right-Truncated Random Variables

$$M = q * L + (1 - q) * R$$

Toward Geometric Black-Scholes pricing define new random variables $Y = K + [X - K]^+$ and Z, a discrete random variable such that $Pr\{Z = K\} = 1$ with discrete density $\rho_Z = \delta_{z=K}$. Then the density ρ_Y of Y is

$$\rho_Y(y) = q * \rho_z(y) + (1 - q) * \rho_R(y)$$

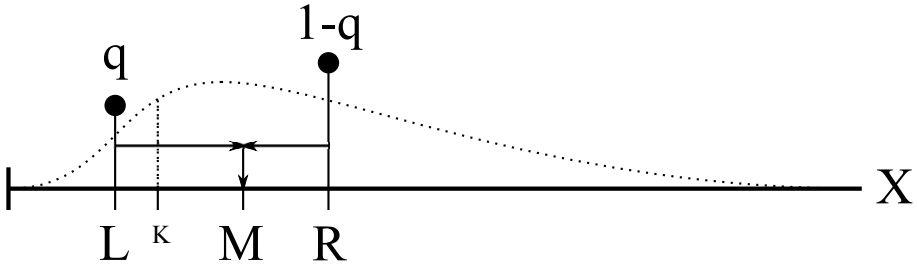


Figure 6.11: Mean Relationship Between Bifurcated Sections

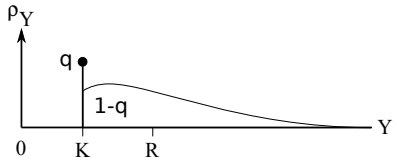


Figure 6.12: Positive Density

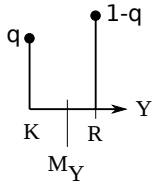


Figure 6.13: Geometric Black-Scholes Pricing

and depicted in figure 6.12. Note in particular that Y is a *mixed* discrete and continuous random variable.

To compute the mean of Y, the X_R portion of Y may be replaced with a discrete density $(1-q)$ at the mean R of X_R, as illustrated in figure 6.13. The mean M_Y of Y is then the affine combination

$$M_Y = \mathbb{E}[Y] = q * K + (1-q) * R$$

Finally Geometric Black-Scholes pricing requires computing the mean of $[X-K]^+$ for some random variable X such that $\mathbb{E}[X] < \infty$ and constant K. Using the notation developed in this section the result is expressed immediately as

$$\mathbb{E}[[X-K]^+] = \mathbb{E}[Y] - K$$

and more conveniently as,

$$\mathbb{E}[[X-K]^+] + K = q * K + (1-q) * R$$

As an illustrative example the traditional Black-Scholes formula is recovered by letting $X = S_T$ be a LogNormally distributed random variable representing the stock price of the underlying asset at some future time T. To use Geometric Black-Scholes to price a European-style call option with maturity time T and

strike price K let $C_T = [X - K]^+$ and find the current price of the call option as $C_0 = \mathbb{E}[C_T]$. Following the steps above symbolically first identify the probability density of S_T,

$$\rho(x) = \frac{1}{x\sqrt{2\pi}\sigma} e^{-\frac{1}{2}(\frac{ln(x)-\mu}{\sigma})^2}$$

Following Dineen [6], the current price S_0 is the mean of S_T,

$$S_0 = \mathbb{E}[S_T] = e^{\mu + \frac{\sigma^2}{2}}.$$

Since S_0 is a known quantity it may be used to solve for the unknown Log-Normal parameter μ,

$$\mu = ln(S_0) - \frac{\sigma^2}{2}$$

and the parameter σ, according to standard interpretations, e.g. Dineen [6], represents the asset volatility and must be given. The option price C_0 is expressed by the formula derived above as

$$\begin{aligned} C_0 + K &= q * K + (1 - q) * R \\ C_0 &= (1 - q) * R - (1 - q) * K \end{aligned}$$

Finding the mean R of the right-truncated LogNormal and its associated probability $(1 - q)$,

$$\begin{aligned} R &= \frac{1}{1-q} \frac{1}{\sqrt{2\pi}\sigma} \int_K^\infty e^{-\frac{1}{2}(\frac{ln(x)-\mu}{\sigma})^2} \, dx \\ &= \frac{1}{1-q} e^{\mu + \sigma^2/2} \Phi\left(\frac{-ln(K) + \mu + \sigma^2}{\sigma}\right) \\ &= \frac{1}{1-q} S_0 \Phi\left(\frac{ln(S_0/K) + \sigma^2/2}{\sigma}\right) \end{aligned}$$

and

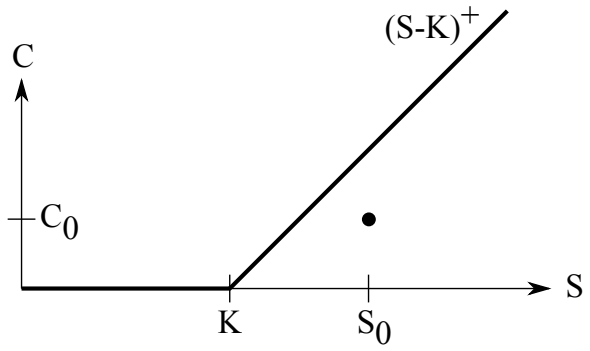


Figure 6.14: Stock-Call Space

$$\begin{aligned} 1 - q &= \Phi\left(-\frac{ln(K) - \mu}{\sigma}\right) \\ &= \Phi\left(\frac{ln(S_0/K) - \sigma^2/2}{\sigma}\right) \end{aligned}$$

where

$$\Phi(z) = Pr(Z <= z) \text{ for } Z \sim Normal(\mu, \sigma^2)$$

the traditional Black-Scholes formula follows immediately,

$$C_0 = S_0\Phi\left(\frac{ln(S_0/K) + \sigma^2/2}{\sigma}\right) - K\Phi\left(\frac{ln(S_0/K) - \sigma^2/2}{\sigma}\right)$$

6.1.2 Portfolio Construction

Given a stock S represented at time $t = 0$ by value S_0 and at $t = T$ by random variable S_T and a European call option C based on S with strike price K at time T the domain of the joint probability distribution of S and C is shown in figure 6.14. The point (S_0, C_0) corresponds to the joint initial price of the stock and call. The broken line is the familiar curve for a call option graph. Perpendicular to the $SC - plane$ is the joint probability density of S and C at time T. The probability density is zero at points away from the broken line.

Since the SC-space represents the price per unit of each financial security there is a dual space where each point represents the number of units held in a hypothetical portfolio. This dual space is referred to here as the portfolio space. Given a portfolio vector π and a point w in SC the dollar value of pi at w is $\pi^T\ w$. Recall that the projection matrix of SC onto a vector v in SC is given by,

$$\Pi = \frac{v\ v^T}{v^T\ v}$$

Notice that given portfolio $\pi = (\pi_s, \pi_c)$ the random variable $Z_T = \pi_s * S_T + \pi_c * C_T$ appears graphically as a hyperplane in SC-space. Notice in particular that if the projection vector in SC-space is π itself then the projection matrix becomes,

$$\Pi = \frac{\pi\ \pi^T}{\pi^T\ \pi}$$

and more to the point the portfolio value V given $x \in SC$-space is,

$$\begin{aligned} V &= \pi^T\ x \\ &= \pi^T \frac{\pi\ \pi^T}{\pi^T\ \pi} x \\ &= \pi^T \Pi x \end{aligned}$$

which means the value of a portfolio, π, may be visualized by representing π as a vector in SC-space, projecting any other point in SC space orthogonally onto π and then computing the inner product of π and the projected point to find the portfolio value.

Suppose, for example, $\pi = (2, 1)$, that is the portfolio contains two shares of stock and one call option with strike price K. The initial price of the stock is S_0 and the initial price of the call is C_0 as usual. This situation is depicted in figure 6.15. The portfolio is represented by Z, the linear combination of S and C. Notice that if $S_T = 0$ then $Z = 0$. If $S_T = K$ then $Z = 2K$ since the call expires out-of-the-money and the value of the portfolio reflects the two shares of stock alone. Geometrically the $Z = 2K$ is found by orthogonally projecting the point $(K, 0)$ to the $\pi = (2, 1)$ vector (the Z-line) and measuring the result with the dual vector, also π to find $V = 2 * K + 1 * 0$. The initial value of the portfolio is shown graphically as Z_0 which is consistent with the computed value

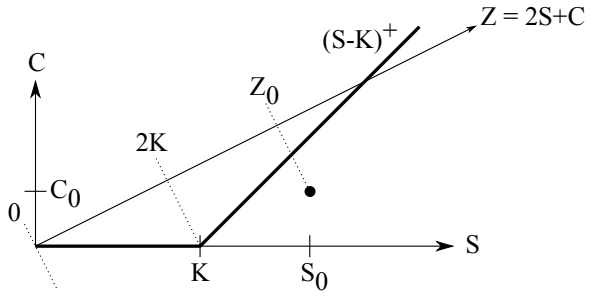


Figure 6.15: Stock-Call Space with (2,1) Portfolio

of $Z_0 = 2 * S_0 + 1 * C_0$. This construction is consistent with the understanding that random variables are represented in a joint space and observed individually through projection.

Orthogonal projection implies the existence of a null space hyperplane. A natural construction is that of a portfolio whose initial value lies in this null space. The initial cost of such a portfolio is zero. Suppose that a call option with strike price K costs C_0 with underlying stock of initial cost S_0 and that $S_0/C_0 = 5$, for example. A zero-cost portfolio is $\pi = (-1, 5)$, that is, sell one share of stock and buy 5 call options. This situation is depicted in figure 6.16. The zero-cost portfolio is represented by $Z = -S + 5C$. Notice that $min(Z) = -K$ and that if the final stock price S_T is $2K$ instead of K the value of the portfolio Z jumps from $-K$ to $3K$, four multiples of K since the portfolio contains 5 calls and one shorted share of stock whereas the terminal stock price rising form 0 to K results in Z falling from 0 to $-K$ since it contains the shorted stock and 5 out-of-the-money call options.

The probability distribution of the example Z is approximated in figure 6.17 where the negative-put-like behavior is shaded and marked with its total probability q assuming S_T is represented by a $LogNormal$ random variable. Notice in particular that the probability distribution of Z is not recognizable as either $LogNormal$ nor truncated $LogNormal$.

6.1.3 Black Scholes Construction: SPY Example

The work on Levy-Stable distributions by Nolan [10] takes the daily prices of a particular stock (ticker: SPY) shown in figure 6.18. The daily returns for SPY are plotted in figure 6.19.

According to Nolan [10] a good fit of the SPY returns is achieved by a mixture

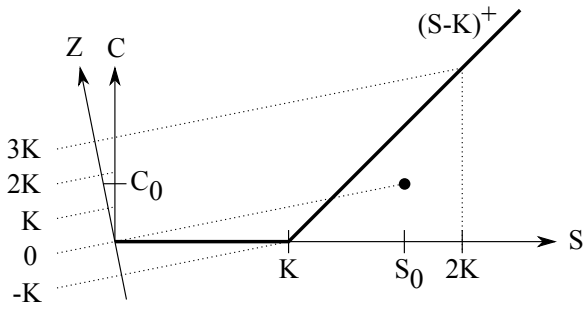


Figure 6.16: Stock-Call Space with Zero-Cost Portfolio

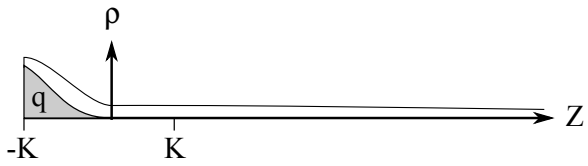


Figure 6.17: Probability Distribution of Zero-Cost Portfolio

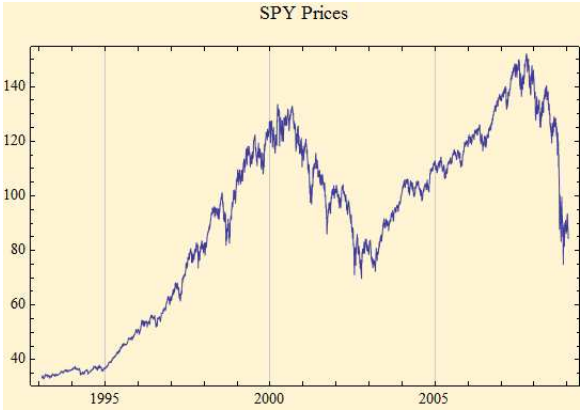


Figure 6.18: Stock Price (ticker symbol: SPY)

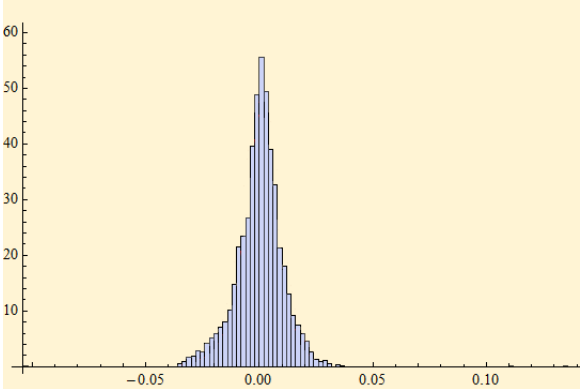


Figure 6.19: Stock Price (ticker symbol: SPY) daily return distribution

of a LogNormal distribution and a Levy-Stable distribution. Using the numeric random variable facilities of RICO to plot $LogNormal(0, \sigma) \times LevyStable(\alpha, \beta, \delta, \gamma)$ for the specific fit parameters cited by Nolan [10],

```
alpha = 1.86034
beta  = -0.0919429
gamma = 0.00600552
sigma = 0.532775
delta = 0.000232571
u     = log(gamma)
LN    = LogNormalNumeric(0, sigma, 100)
LS    = LevyStableNumeric(alpha, beta, delta, gamma, 100)
LNS   = LN * LS
Plot().xrange(-.1, .15).plot(LNS).show()
```

The fit found by Nolan [10] is overlayed on the SPY returns histogram in figure 6.20. Suppose the current price of SPY is \$80/share. Following the Black Scholes construction, the share price of SPY at $T =$ one day in the future is denoted S_T defined as,

$$S_T = 80 \times LNS$$

where LNS is the LogNormal-LevyStable distribution given the fit data found by Nolan [10]. Continuing from the previous code listing, the dsitribution of S_T is computed by RICO as in the following listing with the result is shown in figure 6.21.

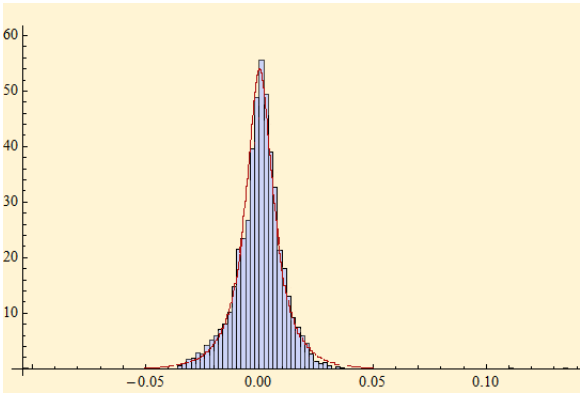


Figure 6.20: Stock Price (ticker symbol: SPY) daily return distribution

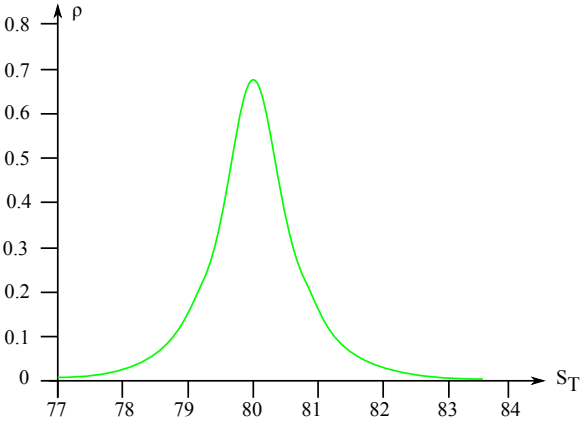


Figure 6.21: SPY one day in future given $80/share price today

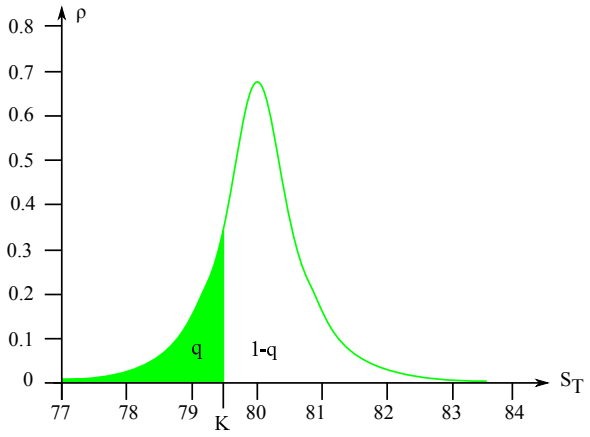


Figure 6.22: SPY one day in future given $80/share price today

```
ST = 80*exp(LNS)
Plot().xrange(77,84).yrange(0,.8).plot(ST).show()
```

Supposing that the strike price for a 1 day option on SPY is $K = \$79.50$. The probability distribution of S_T is then split at K into two pieces as shown in figure 6.22. Let

$$q := P(S_T < K)$$

In this case $q \approx 0.22$ and is represented by the shaded region of figure 6.22. The payoff random variable of a 1-day European-style call option according to the Black-Scholes construction C_T is the following function of S_T

$$C_T = [S_T - K]^+$$

shown in figure 6.23. Notice that C_T is both discrete and continuously distributed. In particular,

$$P(C_T = 0) = q \qquad\qquad P(C_T > 0) = 1 - q$$

The salient point of this example is that the next step of the Black-Scholes construction cannot be completed for ths example. The reason is that the expected value of the continuous portion of C_T is infinite! For the fit parameters used in

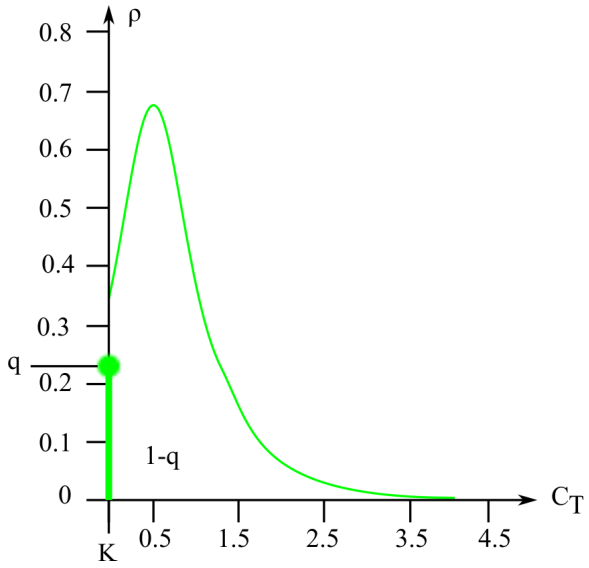


Figure 6.23: Call option payoff on one-day SPY

this example the LNS distribution is *fat tailed* and our example ends without a price for the one-day call option, C_T.

In practice return rates do not necessarily follow a Gaussian distribution. The Black Scholes construction remains computable so long as expected values of truncated distributions are finite. Note in particular that Monte Carlo techniques may not expose this fact.

Appendix A

Exhaustive Bounded Root Finding Method

This section follows the *EveryRoot* method developed by Robert C. Tausworthe [12] under contract with Raytheon, Inc.

Given an objective function f defined on a normalized *investigation* interval, $[0, 1]$, the EveryRoot method finds nearly every root by exhaustive investigation. The EveryRoot method amalgamates several well-known root finding methods. By convention assume that $f_0 = f(0)$, etc.

A.0.4 Specific Assumptions

1. The objective function is presumed continuous on the investigation interval.
2. Given root r it is presumed no other roots exist within interval $[r-xGuard, r+xGuard]$ for a prespecified tolerance value, $xGuard$.
3. A value r is deemed to be a root if $|f(r)| \leq \epsilon f$ for prespecified tolerance value ϵf and $|f(r \pm xGuard)| > \epsilon f$.
4. If successive root estimates r_i and r_{i+1} differ by less than prespecified tolerance value ϵr then r_i is deemed to be a root of f.
5. If an interpolating polynomial L of f differs from the objective function at prespecified points within the interval by less than prespecified tolerance ϵL then L is deemed a *reliable* facsimile of f.

A.0.5 Method Outline

The general approach of the EveryRoot method is to recursively break the investigation interval $[0, 1]$ into subintervals ultimately creating a partition where every subinterval contains exactly one or no roots.

For any subinterval either none, one or both of the endpoints. If an endpoint is a root r then it is centered in a ***guard*** interval, $[r - xGuard, r + xGuard] \cap [0, 1]$ where $r \in \{0, 1\}$ and $[0, 1] - ([r - xGuard, r + xGuard] \cap [0, 1])$ becomes the new investigation subinterval.

Given the endpoints of $[0, 1]$ are not roots, if $f_0 f_1 < 0$ then at least one root exists in interval $(0, 1)$. To find one of the possible roots in $(0, 1)$ a ***bracketing*** method such as Brent's Method [4] or the much more recent Improved Brent's Method [13] is used. The internal root r is found and guarded. The two flanking subintervals $[0, r - xGuard]$ and $[r + xGuard, 1]$ are then investigated.

If $f_0 f_1 > 0$ then the objective function f is fit using an ***optimized*** cubic Lagrange interpolation polynomial L. The optimization comes in the form of two specific function evaluation points within the investigation interval plus the two endpoints. The interpolation method is detailed below. The polynomial L is deemed reliable upon evaluating it at three specific points within the interval where L is most likely to differ from f and finding that the relative error in all three cases is within ϵL tolerance. If L is un-reliable then the investigation interval is bisected into two subintervals.

Suppose L is a reliable facsimile of f and $f_0 f_1 > 0$. If L has an extrema e such that $f_0 L_e < 0$, i.e. on the other side of zero from f_0 and f_1, and subsequently $f_0 f_e < 0$ then the investigation interval is subdivided into $[0, e]$ and $[e, 1]$ so that a bracketing method may be employed. If L_e is *relatively close* to zero then a ***root polishing*** method such as Newton-Raphson or Halley's method maybe employed beginning at some other point in the interval. The method author suggests that if an extrema differs from zero with relative error less than a factor of three (3) of the fit tolerance ϵL then L_e is relatively close to zero. If there is no extrema of L within the interval or none of the extrema are relatively close to zero then the interval is deemed to contain any roots.

A.0.6 The *EveryRoot* Method

```
Write the algorithm, find test cases, then write
                this subsection.
```

TO DO

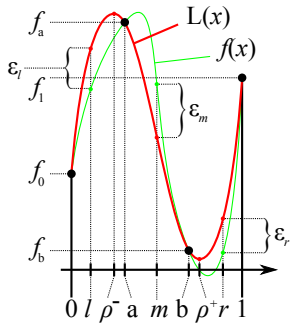


Figure A.1: Example Cubic Lagrange Interpolation

A.0.7 Optimized Cubic Lagrange Interpolation

Given an objective function f over a normalized interval $[0, 1]$ two internal points $0 < a < b < 1$ are chosen for fitting a cubic polynomial L,

$$L(x) = \alpha_0 + \alpha_1 x + \alpha_2 x^2 + \alpha_3 x^3$$

The corresponding four function evaluations are formed into the vector $\mathbf{f}$ where

$$\mathbf{f} = (f_0, f_a, f_b, f_1)^T$$

Evaluating $L(x)$ at points $\{0, a, b, 1\}$ and requiring $L(x)$ to equal $f(x)$ at those points yields the following linear relationship,

$$\begin{pmatrix} f_0 \\ f_a \\ f_b \\ f_1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & a & a^2 & a^3 \\ 1 & b & b^2 & b^3 \\ 1 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \alpha_0 \\ \alpha_1 \\ \alpha_2 \\ \alpha_3 \end{pmatrix}$$

The polynomial coefficients are found by inverting the matrix in the above equation,

$$\begin{pmatrix} \alpha_0 \\ \alpha_1 \\ \alpha_2 \\ \alpha_3 \end{pmatrix} = \frac{1}{D} \begin{pmatrix} D & 0 & 0 & 0 \\ -(D+d) & b & -a & D \\ 2d & -(b+1) & a+1 & -d \\ -d & 1 & -1 & d \end{pmatrix} \begin{pmatrix} f_0 \\ f_a \\ f_b \\ f_1 \end{pmatrix}$$

where $d = b - a,\ D = abd$ and it is assumed for symmetry that $b = 1 - a$.

Figure A.1 depicts an example interpolation with many of the elements discussed in this section present.

As may be found in any elementary calculus text the error of an n^{th}-order polynomial interpolation of $f(x)$ is given by

$$L(x) - f(x) = -\frac{f^{(n+1)}(c)}{(n+1)!}\prod_{i=0}^{n}(x - x_i) \tag{A.1}$$

$$= \frac{f^{(4)}(c)}{4!}\, x(1-x)(x^2 - x + a - a^2) \tag{A.2}$$

where $n = 3$, $x_i \in \{0, a, 1-a, 1\}$ and for some $c \in (0, 1)$. The largest errors occur at the extrema of the fourth-order polynomial $E(x) = x(1-x)(x^2 - x + a - a^2)$ in the error expression. The extrema points are called *ridge-error* points. Solving $E(x)$ for the ridge-error points,

$$\frac{d}{dx}E(x) = 0 \implies x = \left\{\frac{1}{2}\left(1 \pm \sqrt{2a^2 - 2a + 1}\right), \frac{1}{2}\right\}$$

Evaluating $E(x)$ at each ridge-error point yields the unique extrema,

$$\left\{\frac{1}{4}a^2(1-a)^2, -\left(\frac{1-2a}{4}\right)^2\right\}$$

Equating the absolute value of the two ridge-point extrema yields the following candidates for a,

$$2a(1-a) \equiv 1 - 2a \implies a = 1 \pm \frac{1}{\sqrt{2}}$$

Since a is required to lie in the interval $(0, \frac{1}{2})$ so that $0 < a < b < 1$, the ridge-point values are equated in absolute value by the choice

$$a = 1 - \frac{1}{\sqrt{2}} \approx 0.29289322 \tag{A.3}$$

Given this choice of a the ridge-error points are,

$$l = \frac{1}{2}\left(1 - \sqrt{2a^2 - 2a + 1}\right) = \frac{1}{2}\left(1 - \sqrt{2 - \sqrt{2}}\right) \approx 0.117317$$

$$m = \frac{1}{2}$$

$$r = \frac{1}{2}\left(1 + \sqrt{2a^2 - 2a + 1}\right) = \frac{1}{2}\left(1 + \sqrt{2 - \sqrt{2}}\right) \approx 0.882683$$

and the $L(x)$ coefficients are found to be,

$$\begin{pmatrix}\alpha_0\\ \alpha_1\\ \alpha_2\\ \alpha_3\end{pmatrix} = \begin{pmatrix}1 & 0 & 0 & 0\\ -(3+2\sqrt{2}) & 4+3\sqrt{2} & -(2+\sqrt{2}) & 1\\ 4+4\sqrt{2} & -(10+7\sqrt{2}) & 8+5\sqrt{2} & -(2+2\sqrt{2})\\ -(2+2\sqrt{2}) & 6+4\sqrt{2} & -(6+4\sqrt{2}) & 2+2\sqrt{2}\end{pmatrix}\begin{pmatrix}f_0\\ f_a\\ f_b\\ f_1\end{pmatrix} \tag{A.4}$$

Let A_{opt} be the matrix in the above equation. Noticing that

$$L(x) = (1, x, x^2, x^3)\begin{pmatrix}\alpha_0\\ \alpha_1\\ \alpha_2\\ \alpha_3\end{pmatrix} = (1, x, x^2, x^3)\, A_{opt}\, \mathbf{f}$$

it is possible to precompute *ridge-error vectors*, v_l, v_m, v_r corresponding to the ridge-error values l, m, r found above for the sake of computational efficiency. Let

$$\mathbf{v}_k = (1, x_k, x_k^2, x_k^3)\, A_{opt} \qquad \text{for} \qquad k \in \{l, m, r\}$$

The ridge-error vectors are then,

$$\mathbf{v}_l = \frac{1}{4}\left(1+\sqrt{2-\sqrt{2}}, 1+\sqrt{2+\sqrt{2}}, 1-\sqrt{2+\sqrt{2}}, 1-\sqrt{2-\sqrt{2}}\right)^T \tag{A.5}$$

$$\mathbf{v}_m = \frac{1}{4}\left(1-\sqrt{2}, 1+\sqrt{2}, 1+\sqrt{2}, 1-\sqrt{2}\right)^T \tag{A.6}$$

$$\mathbf{v}_r = \frac{1}{4}\left(1-\sqrt{2-\sqrt{2}}, 1-\sqrt{2+\sqrt{2}}, 1+\sqrt{2+\sqrt{2}}, 1+\sqrt{2-\sqrt{2}}\right)^T \tag{A.7}$$

The ridge-error points and vectors are then used to compute the following set of interpolation errors,

$$\epsilon_l = |f(l) - \mathbf{v}_l^T\mathbf{f}\,| \tag{A.8}$$

$$\epsilon_m = |f(m) - \mathbf{v}_m^T\mathbf{f}\,| \tag{A.9}$$

$$\epsilon_r = |f(r) - \mathbf{v}_r^T\mathbf{f}\,| \tag{A.10}$$

The *EveryRoot* method uses the above interpolation errors to determine if $L(x)$ is a reliable facsimile to $f(x)$ over the normalized interval $[0, 1]$. If $L(x)$ is deemed reliable the *EveryRoot* method requires the extrema points of $L(x)$. These are found by solving the quadratic $L'(x) = 0$ and accepting only real roots, that is,

$$L'(x) = 3\alpha_3\, x^2 + 2\alpha_2\, x + \alpha_1 = 0$$

$$\rho = -\frac{\alpha_2}{3\alpha_3} \pm \sqrt{\left(\frac{\alpha_2}{3\alpha_3}\right)^2 - \frac{\alpha_1}{3\alpha_3}}$$

The curvature of $L(x)$ is given by its second derivative,

$$L''(x) = 6\alpha_3\, x + 2\alpha_2$$

When $L''(x)$ is evaluated at the extrema, if any, then the *EveryRoot* method can decide if the extrema is near zero and warrents further investigation. Notice, for example, that in figure A.1 the objective function $f(x)$ crosses zero even though $L(x)$ does not. The example in the figure has grossly exaggerated error values for $\{\epsilon_l, \epsilon_m, \epsilon_r\}$ and $L(x)$ would certainly be rejected by the *EveryRoot* method as un-reliable. An interesting modification to the *EveryRoot* method presents itself in the case of a reliability rejection; subdivide the interval into three parts, not two, namely $\{(0, a), (a, b), (b, 1)\}$. The rationale for this modification is that each of the endpoints has already been evaluated so that $\{f_0, f_a, f_b, f_1\}$ are already in hand.

Notice further from the example depicted in figure A.1 that because $L(x)$ is concave at the upper root, ρ^+, it seems a good place to initialize a root polishing method. In the figure it appears that a root is already found. This is accidental, but makes the point that even a reliable $L(x)$ should not be taken as a duplicate of the objective $f(x)$.

A rule to consider as a modification to the *EveryRoot* method is if

$$|L(\rho)| < max\{\epsilon_l, \epsilon_m, \epsilon_r\}$$

then root polishing should be initiated. A root polisher such as Halley's method, which approximates $f(x)$ with a parabola may best be initiated at the relevant ρ and to bound it by the nearest known neighbors. In the figure the relevant extrema is ρ^+ and the bounding interval is (b, r).

Building the Interpolation Subroutine

The *EveryRoot* algorithm enters the Lagrange interpolation subroutine when $f_0 f_1 > 0$, that is, the end points are on the same side of zero. The evaluations of f_0 and

f_1 are assumed. The Lagrange interpolation subroutine must evaluate $f(x_a) = f_a$ and $f(x_b) = f_b$ to create the α polynomial coefficients. Suppose f_a is such that $f_0 f_a < 0$, that is, on the other side of zero, then the two subintervals described by the partition $[0, a, 1]$, are sent back to the *EveryRoot* method where each is marked as known to contain at least one root.

Continuing with the interpolation subroutine, three more function evaluations at the ridge-error points must be made, namely, $\{f_l, f_m, f_r\}$. The total number of new function evaluations so far is five. If any of the ridge-error evaluations result in a zero crossing the appropriate partition of the investigation interval is sent back to the *EveryRoot* method and each maked as known to contain at least one root. Suppose, for example, only f_l crosses zero such that $f_0 f_l < 0$ then the sub-partitition $[0, l, a]$ is sent back with each containing a root and the remaining $[a, 1]$ is resent to the interpolation subroutine as an investigation interval since $f_a f_1 > 0$.

At this point it is decided whether the interpolation is reliable. If not, there are five new interval endpoints that cannot be recycled so these become natural investigation subintervals, that is, the initial (normalized) interval is partitioned into six subintervals, $[0, l, a, m, b, r, 1]$ and each is re-fed in turn into the interpolation subroutine.

If the interpolation is deemed reliable then extrema are calculated. Either the extrema are real or not and if real inside the interval or not. If there is no extrema within the interval of positive curvature when $f_0 > 0$ and negative if $f_0 < 0$, in other words, close to zero, then the entire interval $[0, 1]$ is marked as containing no roots and the subroutine ends.

If there is an extrema ρ of appropriate curvature then a final function evaluation is made at $f_\rho = f(\rho)$. If $f_0 f_\rho < 0$ then the partition $[0, \rho, 1]$ is sent back each known to contain a root otherwise the two nearest neighbor points are chosen and the roots of the derivative of the objective f are found. For example in figure A.1, ρ^+ is the suggested minima. The nearest neighbors are $\{b, r\}$ and the true minima of f could be left or right of ρ^+.

Suppose the interval $[b, r]$ is suspected of containing the minima. If the endpoints are such that $f'(b) < 0$ and $f'(r) > 0$ then a root bounding method applied to f' over $[b, r]$ is used to find $\hat{\rho}$, the true minima of f. Otherwise the interpolation subroutine is re-entered over interval $[b, r]$ and the flanking intervals $\{[0, b], [r, 1]\}$ are marked as not containing roots. If $\hat{\rho}$ exists and $f_0 f(\hat{r}) > 0$ the interval is declared to no contain a root else the intervals are $\{[b, \hat{\rho}], [\hat{\rho}, r]\}$ each sent to the bounded root method since each contains a root. The possibility that $f(\hat{\rho}) \approx 0$ is considered and if true a single root is returned at $\hat{\rho}$.

Appendix B

Tables and Chairs, Sharply

B.1 An Exemplary 2D Sharp Example

The example presented here is small enough to describe in full detail and incorporate many of the key features of algebraically correlated random variables described in this work. Presented in two stages, first with sharp inputs replicating an example from Gass [8] and the second with algebraically correlated random variable inputs. We will then describe how the example may be generalized.

The basis for the tables and chairs example can be found in Gass [8] wherein a decision must be made by a small furniture manufacturer under resource constraints. The choice is whether to manufacture tables or chair or some combination of both. The goal is to maximize the revenue from the sale of the tables and chairs assuming that all will be sold. The specifics are,

1. There is 400 board-feet of wood available.
2. There is 450 man-hours of labor available.
3. It takes 5 board-feet of wood and 10 man-hours of labor to make a chair.
4. It takes 20 board-feet of wood and 15 man-hours of labor to make a table.
5. Chairs sell for $45 each.
6. Tables sell for $80 each.

Stating the problem in standard form according to Boyd [3] and Greenberg [9],

$$\begin{aligned} \text{maximize} \quad & 45x_c + 80x_t \\ \text{s.t.} \quad & 5x_c + 20x_t \leq 400 \\ & 10x_c + 15x_t \leq 450 \\ & x_c, x_t \geq 0 \end{aligned}$$

where x_c represents the number of chairs to manufacture and sell and x_t represents the number of tables to manufacture and sell.

To form a baseline we will first solve this optimization problem using the simplex method as described in the simplex method section. We will then solve the problem again with the prices kept unknown. At that point we will be ready to introduce random pricing into the problem.

B.1.1 Tables and Chairs with All Inputs Sharp and Known

The simplex method requires all constraints to be stated as equalities so we introduce a slack variable into each inequality. The problem is restated as,

$$\begin{aligned} \text{maximize} \quad & 45x_c + 80x_t \\ \text{s.t.} \quad & 5x_c + 20x_t + s_W = 400 \\ & 10x_c + 15x_t + s_L \leq 450 \\ & x_c, x_t, s_W, s_L \geq 0 \end{aligned}$$

where s_W is the slack variable for the wood resource equation and s_L is the slack variable for the labor resource equation. All variables, x_c, x_t, s_W and s_L are constrained to be non-negative.

Since each slack variable appears exclusively once in the constraint equations and their coefficients are $+1$ they collectively form a basis for the simplex tableau in table B.1

The problem variables are collected into the list $X = (x_c, x_t, s_W, s_L)$ and using the order of this list denote the variables in the current basis with a 1 and the others with a 0 we describe the current simplex state with the binary value,

$$State_0 = 0011$$

	x_c	x_t	s_W	s_L	b
s_W	5	20	1	0	400
s_L	10	15	0	1	450
Revenue	45	80	0	0	

Table B.1: Tables and Chairs Simplex Tableau for State 0011

To pivot the table we find the variable to enter the basis and the basis variable to exit the basis. To find the entering variable we compute the cost impact of each,

$$\begin{aligned} Z_c &= 45 - (5 * 0 + 10 * 0) && = 45 \\ Z_t &= 80 - (20 * 0 + 15 * 0) && = 80 \end{aligned}$$

where Z_c is the cost impact of introducing variable x_c into the basis and Z_t is the cost impact for x_t. Recall that increasing x_c the assumed zero value for a non-basis variable by one unit (one more chair sold, for example) will increase revenue by the price of one chair, $\$45$, and will decrease the slack variables s_W and s_L by 5 and 10 units respectively. Since there is no revenue impact to increasing or decreasing slack variables the revenue impact is zero for each.

The entering variable is selected as,

$$\text{argmax}\{Z_c, Z_t\} \implies x_t$$

We now know that the next simplex state has the form $10??$ because x_c is the entering variable and where the question marks indicate that we do no yet know the exiting variable.

Since x_t is the entering variable we divide $b = (400, 450)$ element-wise by the basis coefficients for x_t namely 20 and 15 and find the minimum non-negative value. In particular one finds,

$$\text{argmin}\{\frac{400}{20}, \frac{450}{15}\} \implies s_W$$

Since $\frac{400}{20} < \frac{450}{15}$ and the former value is associated with basis variable s_W it is chosen as the exit variable. Recall that the reason is because we are allowing the entering variable x_t to increase from zero it forces the basis variable in each equation toward zero. In the first equation a unit increase in x_t is a 20 unit decrease in $s_W = 400$, but only a 15 unit decrease in $s_L = 450$. Since no variable is allowed

	x_c	x_t	s_W	s_L	b
x_t	$\frac{1}{4}$	1	$\frac{1}{20}$	0	20
s_L	$6\frac{1}{4}$	0	$-\frac{3}{4}$	1	150
Revenue	45	80	0	0	

Table B.2: Tables and Chairs Simplex Tableau for State 0101

to be negative the we find which basis variable is driven to zero first by an increase in the entering variable. We now see the new simplex state to be,

$$State_1 = 0101$$

To transform the equations and update the tableau we form the transformation matrix to state 1, B_1 and its inverse as,

$$B_1 = \begin{pmatrix} 20 & 0 \\ 15 & 1 \end{pmatrix} \qquad\qquad B_1^{-1} = \frac{1}{20}\begin{pmatrix} 1 & 0 \\ -15 & 20 \end{pmatrix}$$

recognizing each tableau column as a vector and multiplying on the left by B_1^{-1} we find the new tableau in table B.2.

The two non-basis variables are now x_c and s_W so we find the cost impact for introducing each into the basis,

$$\begin{aligned} Z_c &= 45 - (\frac{1}{4} * 80 + 6\frac{1}{4} * 0) && = 25 \\ Z_w &= 0 - (\frac{1}{20} * 80 - \frac{3}{4} * 0) && = -4 \end{aligned}$$

where Z_w is the cost impact of (re)-introducing s_W into the basis. Since Z_w is negative, s_W it is not eligible to be a basis vector leaving x_c as the only available choice for entering variable.

Dividing b by the vector of coefficients associated with x_c and finding the smallest non-negative value we have,

$$\text{argmin}\{20 \div \frac{1}{4}, 150 \div 6\frac{1}{4}\} = \text{argmin}\{80, 24\} \implies s_L$$

demonstrating the s_L is the exiting variable. The B_2 basis transformation matrix and its inverse become,

	x_c	x_t	s_W	s_L	b
x_c	1	0	-0.12	0.16	24
x_t	0	1	0.08	-0.04	14
Revenue	45	80	0	0	

Table B.3: Tables and Chairs Simplex Tableau for State 1100

$$B_2 = \begin{pmatrix} \frac{1}{4} & 1 \\ 6\frac{1}{4} & 0 \end{pmatrix} \qquad B_2^{-1} = \begin{pmatrix} 0 & 0.16 \\ 1 & -0.04 \end{pmatrix}$$

The simplex state is then,

$$State_2 = 1100$$

and the tableau for state 1100 is shown in table B.3.

We see that the two slack variables are no longer in the basis so they are both zero. This means that at the current state we are using all available resources to manufacture our tables and chairs. We ask if either of the two slack variables should be re-introduced into the basis by calculating the revenue impact for each,

$$\begin{aligned} Z_W &= 0 - (-0.12 * 45 + 0.08 * 80) &&= -1 \\ Z_L &= 0 - (0.16 * 45 - 0.04 * 80) &&= -4 \end{aligned}$$

Since each cost impact is negative we conclude there is no possible way to improve the revenue of the problem and the algorithm terminates with the results,

$$\begin{aligned} x_c &= 24 \\ x_t &= 14 \\ revenue &= 24 * 45 + 14 * 80 = \$2,200 \end{aligned}$$

since $(x_c, x_t) = b$. This means that the optimal choice for the manufacturer is to make 24 chairs and 14 tables which, when sold, will generate a revenue of $\$2,200$.

Figure B.1 shows the resource constraints (diagonal lines), the feasible region (shaded area) and the optimal point, $(24, 14)$. The simplex method starts at the

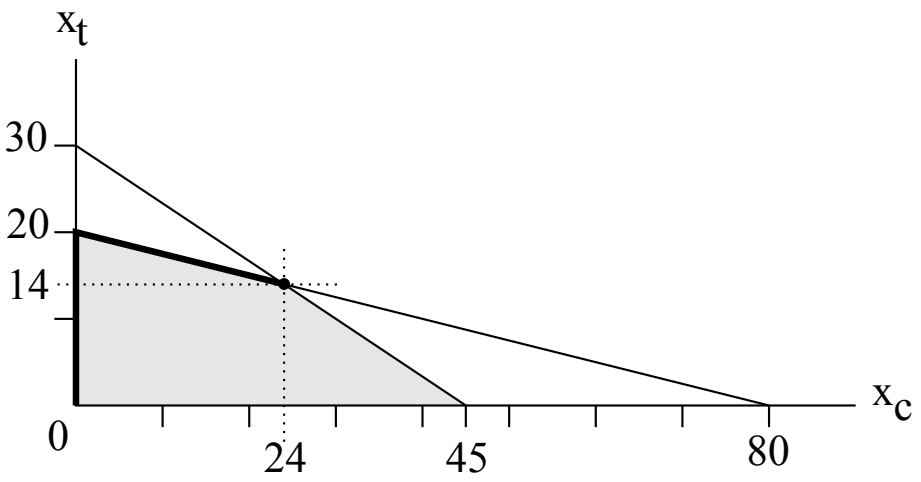


Figure B.1: Tables and Chairs Constraints and Optimal Point

origin in this case and follows the heavy line in figure B.1 from vertex to vertex of the polytope described by the half-space constraints to the optimal vertex. We notice that in this case there is an alternate vertex-path from the origin to the optimal vertex, namely passing through point $(45, 0)$. We will see that the choice of vertex-path is significant in the next version of this example when we leave the prices unknown.

B.1.2 Tables and Chairs with Unknown Prices

As an intermediate step we consider where uncertainty may be injected into the tables and chairs example and decide that leaving the prices sharp, but unknown leads to some revealing results.

A priori there are three places in the tables and chairs example where uncertainty may be injected; the constraint vector b, the price vector p and the constraint matrix A where,

$$A = \begin{pmatrix} 5 & 20 \\ 10 & 15 \end{pmatrix}$$
$$b = \begin{pmatrix} 400 \\ 450 \end{pmatrix}$$
$$p = \begin{pmatrix} 45 & 80 \end{pmatrix}$$

We recognize the values in the A matrix as the amount of resources of each type consumed to manufacture each kind of product, b is the number of resources of each kind available and p is the prices charged for each product.

Suppose that instead of the 5 in matrix A we introduced a random variable. In the context of the tables and chairs example this means that the manufacturer is uncertain about the amount of wood necessary to construct a chair. We assume there is only one kind of chair else we would likely call out the different kinds as different products and give them separate variables. We make similar statements about each value in the A matrix.

Since we are most interested in reflecting the AB32 reference model into the tables and chairs example we elect not to introduce random variables into the A matrix. The reason is that the corresponding A matrix of the AB32 model represents policy features and physical limitations which are assumed for the given policy under consideration.

To choose to introduce random variables into the b or p vectors we must understand how they are used within the simplex algorithm. The simplex method uses a *pivot-table* approach whereby a column and related row within the simplex tableau are chosen and a transition is made to a new state within the algorithm.

To choose the simplex tableau column we compare revenue impacts given a choice of one of the non-basis variables. The values involved in computing the revenue impact of each non-basis variable are prices and products of prices and A matrix coefficients from columns corresponding to the non-basis variables. If we assume that the A matrix values are fixed then we are comparing linear combinations of p vector prices to find the non-basis variable corresponding to the non-negative maximum of revenue impact values.

To choose the simplex tableau row given a column choice we compute the quotient of b and the values in A corresponding to that column. Thus we are comparing linear combinations of b vector values to find the basis variable corresponding to the the non-negative minimum of linear combinations of constraint values.

We thus see that p and b values do not interact directly within the simplex method. We will then choose the p vector for introduction of random variables suggesting, in the tables and chairs example, price uncertainty over the b vector values which would suggest resource uncertainty. We will see that no new insight is gained though choosing b over p or through choosing both for random variable introduction. We will comment below on the choice of A for random variable introduction especially if p or b are chosen as well.

For this intermediate tables and chairs example we have unknown, but sharp

0011	x_c	x_t	s_W	s_L	b
s_W	5	20	1	0	400
s_L	10	15	0	1	450
Revenue	p_c	p_t	0	0	

Table B.4: Tables and Chairs Simplex Tableau for State 0011 with Unknown Prices

prices p_c and p_t for chairs and tables respectively. We make one assumption about these unknown prices; they are positive. The problem may then be stated in standard form as,

$$\begin{aligned} \text{maximize} \quad & p_c * x_c + p_t * x_t \\ \text{s.t.} \quad & 5x_c + 20x_t \leq 450 \\ & 10x_c + 15x_t \leq 450 \\ & x_c, x_t \geq 0 \\ & p_c, p_t > 0 \end{aligned}$$

The initial simplex tableau is shown in table B.4. The only differences from the initial tableau of the first example is the introduction of the state value 0011 into the upper-left corner and the unknown prices p_c and p_t.

Following the steps from the first example we must find the entering non-basis variable by finding,

$$\begin{aligned} &\text{argmax}\{p_c - (5 * 0 + 10 * 0), p_t - (20 * 0 + 15 * 0)\} \\ =&\text{argmax}\{p_c, p_t\} \end{aligned}$$

Since $p_c, p_t > 0$ neither case may be disqualified so we have some possibilities. Either $p_c < p_t$ or $p_c > p_t$ or $p_c = p_t$. Since we intend, in the next example, to introduce continuous random variables in place of p_c and p_t, equality occurs with probability zero so we ignore that case here.

If $p_c > p_t$ then we choose x_c as the entering variable. To find the exiting variable we compute,

$$\begin{aligned} &\text{argmin}\{\frac{400}{5}, \frac{450}{10}\} \\ =&\text{argmin}\{80, 45\} \implies s_L \end{aligned}$$

1010	x_c	x_t	s_W	s_L	b
x_c	1	1.5	0	0.1	45
s_W	0	12.5	1	-0.5	175
Revenue	p_c	p_t	0	0	

Table B.5: Tables and Chairs Simplex Tableau for State 1010 with Unknown Prices

Since we have chosen x_c as the entering variable and s_L as the exiting variable our new state is 1010 and the transition matrix B_{1010} and its inverse B_{1010}^{-1} is,

$$B_{1010} = \begin{pmatrix} 5 & 1 \\ 10 & 0 \end{pmatrix} \qquad B_{1010}^{-1} = \frac{1}{10}\begin{pmatrix} 0 & 1 \\ 10 & -5 \end{pmatrix}$$

The new 1010 tableau is shown in table B.5.

The non-basis variables are x_t and s_L so we compute the revenue of (re)-introducing each of them, respectively, and find the entering variable,

$$\text{argmax}\{p_t - (1.5 * p_c + 12.5 * 0), 0 - (0.1 * p_c - 0.5 * 0)\} = \text{argmax}\{p_t - 1.5 * p_c, -0.1 * p_c\}$$

Since $p_c > 0$ by assumption we have $-0.1 * p_c < 0$ so it must be disqualified as an option. We then ask, under what condition is the first options positive? That is,

$$\begin{aligned} p_t - 1.5 * p_c &> 0 \\ p_t &> 1.5 * p_c \\ \frac{2}{3}p_t &> p_c \end{aligned}$$

Since we have already assumed upon entering this case that $p_c > p_t$ it is not possible for $\frac{2}{3}p_t > p_c$. We therefore terminate the simplex algorithm. Recalling that non-basis variables must be zero we find the following results,

$$\begin{aligned} x_c &= 45 \\ x_t &= 0 \\ revenue &= 45p_c \qquad \text{for } p_t < p_c \end{aligned}$$

0101	x_c	x_t	s_W	s_L	b
x_t	$\frac{1}{4}$	1	$\frac{1}{20}$	0	20
s_L	$6\frac{1}{4}$	0	$-\frac{3}{4}$	1	150
Revenue	p_c	p_t	0	0	

Table B.6: Tables and Chairs Simplex Tableau for State 0101 with Unknown Prices

Returning to our first decision point we now assume $p_c < p_t$ as was the case in the first example. The tableau under the assumption that we transition from state 0011 to state 0101 is shown in table B.6 which we notice is similar to table B.2 except for the unknown prices and the state being recorded in the upper left corner of the tableau.

From state 0101 the non-basis variables are x_c and s_W so we compute the revenue maximizing variables by the usual methods,

$$\begin{aligned} &\text{argmax}\{p_c - (\frac{1}{4}p_t + 6\frac{1}{4} * 0), 0 - (\frac{1}{20}p_t - \frac{3}{4} * 0)\} \\ =&\text{argmax}\{p_c - \frac{1}{4}p_t, -\frac{1}{20}p_t\} \end{aligned}$$

Since $p_t > 0$ the second option of $-\frac{1}{20} * p_t < 0$ and is disqualified. From the first option we conclude that if $p_c < \frac{1}{4} * p_t$ that the simplex algorithm terminates. The results of this termination are,

$$\begin{aligned} x_c &= 0 \\ x_t &= 20 \\ revenue &= 20p_t \qquad\qquad \text{for } p_c < \frac{1}{4}p_t \end{aligned}$$

For the simplex algorithm to not terminate at this part requires that $p_c > \frac{1}{4}p_t$. Recall that we are already operating under the assumption that $p_c < p_t$. Since these two assumptions are compatible (i.e. not impossible) we continue the simplex algorithm. From the calculation for entering variable we conclude that x_c is the entering variable and we must find the exiting variable. Since we have seen this exact situation in the first example we simply recall the result that s_L is the exiting variable and that the resulting tableau is shown in table B.7. Because we

1100	x_c	x_t	s_W	s_L	b
x_c	1	0	-0.12	0.16	24
x_t	0	1	0.08	-0.04	14
Revenue	p_c	p_t	0	0	

Table B.7: Tables and Chairs Simplex Tableau for State 1100 with Unknown Prices

are in the same state as before this new figure is nearly identical to the previous table B.3.

In the first example, once we reached this state (1100) the simplex algorithm terminated. As before we attempt to find the entering variable with the calculation,

$$\begin{aligned} &\text{argmax}\{0-(-0.12p_c+0.08p_t), 0-(0.16p_c-0.04p_t)\} \\ =&\text{argmax}\{0.12p_c-0.08p_t, 0.04p_t-0.16p_c\} \\ =&\text{argmax}\{3p_c-2p_t, p_t-4p_c\} \end{aligned}$$

For the second option to be positive and therefore available for consideration requires that $p_t > 4p_c$ which is to say $p_c < \frac{1}{4}p_t$. However, to reach this state we assumed $p_c > \frac{1}{4}p_t$ so the second option is not available.

If $p_c < \frac{2}{3}p_t$ then the simplex algorithm terminates just as it did in the first example since $45 < \frac{2}{3} * 80 = 53.33...$. The result in this case is,

$$\begin{aligned} x_c &= 14 \\ x_t &= 24 \\ revenue &= 14p_c + 24p_t \qquad\qquad \text{for } \frac{1}{4}p_t < p_c < \frac{2}{3}p_t \end{aligned}$$

If $\frac{2}{3}p_t < p_c$ then the first option for the entering variable is available and the entering variable is found to be s_W.

The exiting variable in this case is then found as,

$$\text{argmin}\{24 \div -0.12, 14 \div 0.08\}$$

Since the first option is negative it is disqualified leaving the second option and therefore the second of the two basis variables x_t as the exiting variable.

Pivoting on the s_W column and the x_t row we have the transition matrix B_{1010} and its inverse B_{1010}^{-1} as,

$$B_{1010} = \begin{pmatrix} 1 & -0.12 \\ 0 & 0.08 \end{pmatrix} \qquad B_{1010}^{-1} = \begin{pmatrix} 1 & 1.5 \\ 0 & 125.5 \end{pmatrix}$$

Applying our transition matrix B_{1010}^{-1} to the 1100 state tableau return us to the 1010 state tableau shown in figure B.5. Since we have determined that state 1010 terminates we have the following results,

$$\begin{aligned} x_c &= 45 \\ x_t &= 0 \\ revenue &= 45p_c \qquad\qquad \text{for } \frac{2}{3}p_t < p_c < p_t \end{aligned}$$

Because we followed a different path between states to arrive at state 1010, the conditions for reaching this state are different. We notice that the conditions for reaching this state directly from the initial 0011 state are $p_t < p_c$ do not intersect the conditions for reaching this state from state 1010 as we have just completed. We combine the two conditions to see that the revenue outcome of $45p_c$ is reached if $\frac{2}{3}p_t < p_c$.

The result of our investigation is that there are three possible cases for any pair of positive prices given that all other values in the tables and chairs example remain the same, that is,

$$\text{Revenue } = \begin{cases} 45p_c & \text{if } \frac{2}{3}p_t < p_c \\ 24p_c + 14p_t & \text{if } \frac{1}{4}p_t < p_c < \frac{2}{3}p_t \\ 20p_t & \text{if } p_c < \frac{1}{4}p_t \end{cases}$$

The results of this example are summarized by the directed graph in figure B.2. The nodes of the graph are the states of the simplex algorithm allied to this example, the edges are marked with the conditions under which the simplex algorithm will follow the edge and the three output cases are labeled A, B and C accompanied by the resulting revenue.

Looking ahead, suppose we were given random variables P_t and P_c representing the price of tables and chairs respectively. Even if P_t and P_c are correlated

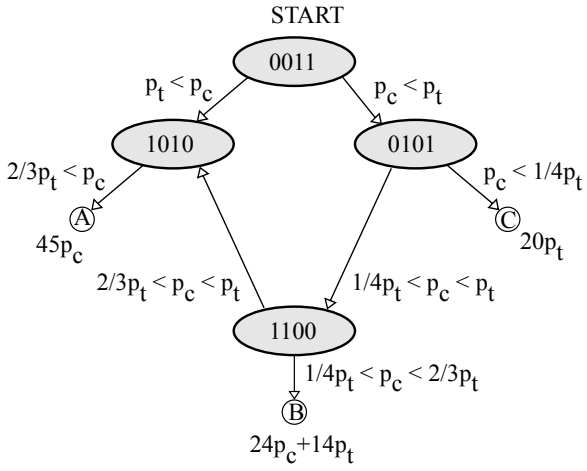


Figure B.2: Tables and Chairs Directed Graph

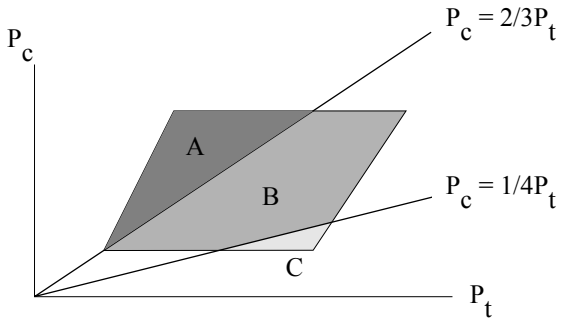


Figure B.3: Tables and Chairs Partitioned Price Probability Space

in some manner they have a joint distribution which we represent as the shaded region in figure B.3. The figure (B.3) has three levels of shading with each region labeled with its outcome (A, B or C) corresponding to the directed graph in figure B.2. We note that if the two price random variables are independent then the shaded region representing the joint probability distribution of P_t and P_c would be rectangular.

Appendix C

Symbolic Expression Transformation

C.1 Simplify

A Rico object may represent a function whose operands are themselves Rico objects. A Rico Object then forms a *parse tree*. The reference to *parsing* stems from a tree of Rico objects being created in response to user interaction or user-provided algorithm. The possibility of algorithmic creation of parse trees drives the focus to expression simplification since algorithms may generate parse trees that involves tens or even tens of millions of terms.

As detailed elsewhere not all forms of algebraically equivalent expressions are equally easy to evaluate numerically. As an example consider $X + XY$ where X and Y are mutually independent random variables. The addition operation between X and XY is not simple since the operands are algebraically correlated. However, the equivalent expression, $X(1 + Y)$ involves the addition and multiplication of mutually independent random variables. Independent combinations of random variables tend to be easier to implement numerically than correlated combinations since there is no correlation component.

Another reason for expending the design cost of creating a function that simplifies Rico parse trees is to reduce the evaluation work load. For example, an expression such as $(X + X^2)/X$ is better represented as the algebraically equivalent $1 + X$. Similarly $log(exp(X))$ is better represented as X. While a user may not type such an expression as $log(exp(X))$, an algorithm may generate it or it may arise as an intermediate step in the simplification process itself.

Rather than attempting to define a *simplification* metric for all admissible expressions, the goal of this section is to demonstrate that the process called *simplification* within the Rico project reliably produces algebraically equivalent output expressions for any admissible input expression and tends to produce simpler equivalent expressions.

The simplification algorithm is detailed below. It employs expression simplification heuristics that will also be detailed. A testing regimine will be detailed and the claim that the testing regimine will guarantee the claim of expression equivanence.

C.1.1 Simplification Overview

When given an arithmetic expression in the form of a Rico object representing an expression parse tree the simplification algorithm will recursively call itself on any operand objects (also called *child* objects) forming a *depth-first* algorighm. This knowledge allows each step in the algorithm to assume that any operand objects are in *simplest form*. The definition of *simplest* in this context is itself recursive. The meaning of *simplest* emerges within the simplification steps.

The simplification algorithm is comprised of several stages, the first is properly called the simplifying step. The purpose of this step is to remove all integer exponents (save 1 and -1) from sums of objects, recover identity operations such as $exp(log(X)) \rightarrow X$ and $X/X \rightarrow 1$, generally convert any expression to the quotient of sums of products to the extend possible.

Under certain circumstances division of polynomials may be attempted in secondary stages of simplification. A conditional last stage of simplification is factoring of common terms.

Another motivating example is,

$$\frac{(x+2)^2+1}{x^2+4x+5} \rightarrow 1$$

The components to the simplification algorithm are described below with the lowest order first.

C.1.2 Rico Objects

Before expression simplification is attempted the admissible expressions must be detailed. For the most part Rico creates it's parse tree as numerical operations

are encountered. If, for example, the user directs Rico to form the square root of a Rico object, the result is a square root object with the starting object as the sole operand. The process is similar for log, exp, subtraction, division, etc. The exceptions are addition, multiplication and exponentiation (or power).

There are three basic kinds of Rico objects to consider during simplification; NaN, numeric, and non-numeric where *NaN* stands for not-a-number. A NaN Rico object arises to indicate a result is undefined such as 0/0, etc. A numeric Rico object is one of the allowed numeric types such as integers, floating point values and fractions (of integers). A non-numeric Rico object represents either a built-in random variable or a function of Rico objects whose result is not numeric and well defined (not NaN).

It is convenient to use compact symbols to represent the three kinds of Rico objects,

ϕ is a NaN Rico object

ν is a numeric Rico object

ρ is a non-numeric Rico object

and capital letters such as X, near end of alphabet represent generic Rico objects.

Simplify Negate

The negation simplifier accepts one operand and obeys the following rules,

$$\begin{aligned} -\{\phi\} &\to \phi \\ -\{\nu\} &\to -\nu \\ -\{\nu \times X\} &\to \{-\nu\} \times X \\ -\{X\} &\to \{-1\} \times X \end{aligned}$$

The $-\{\nu \times X\}$ rule allows number objects to absorb the negation operation. The last rule replaces the negation operation with the product of X and -1, in the conventional number-first order.

Rico Addition

The Rico addition object accepts two or more operands. There are four cases for each operand,

$$\{\phi, 0, \nu, \rho\}$$

where it is understood that a rule that applies to zero is selected before a generic number ν and X applies to any Rico object. Since the rules are applied in order the specific ρ symbol tends not to appear. Rather, the generic X (and subsequent generics) appear most often.

Several features are accommodated in the following rules,

1. Propagate NaN values
2. Discard zeros
3. Perform numeric addition
4. Ensure number objects appear in first position in list of operands
5. Perform coefficient addition
6. Collapse cascaded addition operations

These rules may be encoded as if there are only two operands as follows,

$$\begin{gathered}
\phi + X \to \phi \\
X + \phi \to \phi \\
X + 0 \to X \\
0 + X \to X \\
\nu_1 + \nu_2 \to \{\nu_1 + \nu_2\} \\
X + \nu \to n + X \\
\nu_1 \times X + \nu_2 \times X \to \{\nu_1 + \nu_2\} \times X \\
X + (Y + Z) \to X + Y + Z \\
(X + Y) + Z \to X + Y + Z \\
X + Y \to X + Y
\end{gathered}$$

When there are more than two operands the behaviour is as if an accumulation of two-operand addition operations, i.e. $(X + Y + Z) = ((X + Y) + Z)$.

Rico Multiplication

Within Rico multiplication is handled in a similar manner to addition. Cascades products are flattened, zeros and NaN's dominate expressions and ones are filtered out. The rules are then,

$$
\begin{aligned}
NaN * Y &\rightarrow NaN \\
X * NaN &\rightarrow NaN \\
X * 0 &\rightarrow 0 \\
0 * Y &\rightarrow 0 \\
X * 1 &\rightarrow X \\
1 * Y &\rightarrow Y \\
(X * Y) * Z &\rightarrow X * Y * Z \\
X * (Y * Z) &\rightarrow X * Y * Z \\
(X * Y) * (Z * K) &\rightarrow X * Y * Z * K
\end{aligned}
$$

As in the addition case the NaN dominates the output of any function. Like the NaN value, zero dominates the multiplication output. Note in particular that NaN dominates zero. The next two rules involving one are the equivalent nullification as the zero is for addition. The last three rules detail the expression flattening of cascaded products.

Rico Power

The Power function, X^Y, requires more care so it doesn't provide any automatic flattening of cascaded powers as addition and multiplication do provide. There are still the automatically applied rules for NaN, 0 and 1,

$$
\begin{aligned}
X^{NaN} &\rightarrow NaN \\
NaN^{Y} &\rightarrow NaN \\
NaN^{NaN} &\rightarrow NaN \\
(X1 + \cdots + X3)^{n} &\rightarrow X1^{n} + \cdots + X3^{n} \\
(X1 \times \cdots \times X3)^{Y} &\rightarrow X1^{Y} \times \cdots \times X3^{Y} \\
(X^{Y})^{Z} &\rightarrow X^{Y \times Z} \\
exp(X)^{Y} &\rightarrow exp(X \times Y) \\
n1^{n2} &\rightarrow n3 \\
X^{0} &\rightarrow 1 \\
X^{1} &\rightarrow X
\end{aligned}
$$

The $n1^{n2} \rightarrow n3$ rule is triggered when both operands are numeric values and the rule says to evaluate these two operands into a single numeric value, possibly Nan. The numeric power rule is detailed below.

The rule $(X1 + X2 + ... + X3)^{n} \rightarrow X1^{n} + ... + A3^{n}$ assumes an integer exponent. The distributive law is them applied so that the result is a sum of powers.

The rule $(X1 \times \cdots \times X3)^{Y} \rightarrow X1^{Y} \times \cdots \times X3^{Y}$ distributes powers over each product operand. This rule is in keeping with the overall strategy of exploring legal interactions bewteen operands. A later factoring step for exponents will recover the original form if possible.

The rule $X^{0} \rightarrow 1$ presumes that X is non-numeric. This is no guarantee that X does not represent a random variable with non-zero probability concentrated at zero or either infinity. This rule will need to be revisited when discrete random variables are introduced. Doubtless this rule will be removed at that point.

Rico Numeric Power

The special case results of the power evaluation of two numeric values are detailed in table XXX.

x^y	0	1	$+\infty$	$-\infty$	NaN	$d2$
0	NaN	0	0	NaN	NaN	$\{NaN, 0\} \sim \{-, +\}$
1	1	1	NaN	NaN	NaN	1
$+\infty$	NaN	$+\infty$	NaN	NaN	NaN	$\{0, +\infty\} \sim \{-, +\}$
$-\infty$	NaN	NaN	NaN	NaN	NaN	$\{0, -\infty\} \sim \{-, +\}$
NaN	NaN	NaN	NaN	NaN	NaN	NaN
$d1$	1	x	$sgn(d1)\infty$	0	NaN	$d1^{d2}$

The expression $\{0, -\infty\} \sim \{-, +\}$ in the last column means that 0 is chosen if $d2 < 0$ else $-\infty$ is chosen. There is no possibility for $d2$ to be zero since that case is already addressed.

Equality Test

Several parts of the simplification algorithm rely on the ability to compare two objects for equality. Two Rico objects are *equivalent* if they are of the same Rico type and if they have child objects (operands for functional objects) then, modulo order in the case of addition and multiplication, those child objects are also equivalent.

The following Rico objects (expressions) are then equivalent,

$$X + Y + 3 \equiv 3 + X + Y$$
$$X * 5 \equiv 5X$$

but many algebraically equivalent expressions are not equivalent in the Rico sense. For example,

$$(X + 2)^2 + 1 \neq X^2 + 4X + 5$$

Note especially that Basic random variables are a special type of Rico object. Basic random variables are issued a unique id code upon creation and two Basic random variables are distinguished by id code, not random variable type. For example if $X = Normal(2, 3)$ and $Y = Normal(2, 3)$, then $X \equiv Y$ iff $id(X) = id(Y)$.

Rico Numbers are equivalent if they represent identical values. Floating point Numbers are equivalent if they are within a specific numerical tolerance to compensate for minor least significant bit differences. Fractions in Rico are easily compared by comparing their integer numerator and denominator since they are always represented in lowest form with the numerator carrying the signed value.

C.1.3 Simplify Stage One

Stage One of the simplify algorithm is depth-first. At each node Stage One of the simplify algorithm is called and upon return a transformation appropriate to the current object is called. If the current object is a function and therefore has child objects (operands) then they are assumed to have already been Stage One simplified.

The Stage One object simplifications are detailed by object type. Many simplifications build on each other.

Simplify Subtract

The subtration function in Rico accepts exactly two operands. The Stage One simplification is to transform subtration to addition as follows,

$$A - B \rightarrow A + \{-1\} * B$$

The end result is that a Stage One simplified expression contains no subtraction functions.

Simplify Division

Analagously to the subtraction function, the division function is transformed as,

$$A \div B \rightarrow A * B^{-1}$$

Simplify Square Root

Building on the division function transformation there is replacement of the square root function,

$$\sqrt{A} \rightarrow A^{\frac{1}{2}}$$

Simplify Log

The Stage One simplification huristic is to turn products of sums to sums of products. This extends to the following transformation rules,

$$\begin{aligned} log(exp(X)) &\rightarrow A \\ log(X^Y) &\rightarrow Y * log(X) \\ log(X * Y * Z) &\rightarrow log(X) + log(Y) + log(Z) \end{aligned}$$

In this way, multiplication within a function is exposed outside as addition.

Simplify Exp

By similar reasoning to Stage One log simplification the exp rules are,

$$exp(log(X)) \rightarrow A \qquad exp(X + Y + Z) \rightarrow exp(X) * exp(Y) * exp(Z)$$

In this way, addition within a function is exposed outside as multiplication.

Simplify Addition

A Rico addition object may contain two or more operand objects. The hallmark transformation at this stage is,

$$3X + X \rightarrow 4X$$

Notice that $3X$ is a multiplication object that contains a number object as an operand. It is assumed that the Stage One Multiply simplification has already occurred and that this multiplication object contains zero or one number object operands and furthermore if a number object is present it is the first operand in the list of operands. A further consideration is seen in the following transformation,

$$3XY + X + YX \rightarrow 4XY + X$$

The number object (3) in the above expression is commonly referred to as the coefficient of the expression. This coefficient must be split from the three operand multiply object into a number object and a two operand multiply object. The latter is referred to as the *base* object. Unique base objects is must be identified and the coefficients of any duplicated base objects are summed. Two related special cases are addressed as follows,

$$2XY + X + -2XY \rightarrow X$$
$$2XY + -2XY \rightarrow 0$$

Simplify Multiplication

Multiplication simplification returns NaN if any openand is of type NaN. The distributive law is then applied, as needed, and the result is a sum of Stage One simplified products.

$$X \times NaN \rightarrow NaN$$
$$NaN \times Y \rightarrow NaN$$
$$NaN \times NaN \rightarrow NaN$$
$$(X + Y) \times Z \rightarrow XZ + YZ$$
$$X \times (Y + Z) \rightarrow XY + XZ$$
$$X^Y \times X^Z \rightarrow X^{Y+Z}$$
$$n1 \times n2 \rightarrow n3$$
$$(X \times Y) \times (Y \times Z) \rightarrow XY^2Z$$

The last term applies to an operand that is itself a product of an arbitrary number of operands. For example,

$$(X \times 2 \times Y) \times (Y \times Z) \rightarrow 2XY^2Z$$

C.1.4 Simplify Stage Two

```
This needs to be fleshed out.  First the StageOne
        Simplify needs to be completed.
```

TO DO

Factoring

The Factor tree operation is depth-first. It acts at addition nodes to factor out common nodes from sums of products of nodes. Two example transformations are,

$$AB + AC + D \to A(B + C) + D$$
$$5A + 5^2AB + 5^3ABC \to 5A(1 + 5B(1 + 5C))$$

where the juxtaposition implies multiplication of the separate nodes A, B, C. The first transformation demonstrates that the A node need not be common to all products. The second transformation comes from the concept of Net Present Value.

`Define what happens to fractional powers` TO DO

The expression $AB + AC + BC$ may be factored as $A(B + C) + BC$ or $B(A + C) + AC$ demonstrating that factorization is not uniquely defined. The Factor tree operation is implemented using a *sparse product matrix*.

An example without non-unit exponential is first considered. Given the expression,

$$AB + AC + D$$

A list of uniqe nodes is constructed; $\{A, B, C, D\}$. Using this list to define the columns referred to as a ***header***, a sparse matrix of powers is constructed where each row represents a product of header elements and the rows collectively represent the sum of products. The sparse product matrix for this example is then,

$$\begin{pmatrix} A & B & C & D \\ 1 & 1 & & \\ 1 & & 1 & \\ & & & 1 \end{pmatrix}$$

The A is most common element since it appears in more rows than any other header element it is chosen as the factored element. If another header element, say B had been as common as A then the powers of A and B would be summed and the greater chosen as the factored element. If a tie persist then the first of the tied element in the header list is chosen.

To factor out the A, the sparse matrix rows are partitioned into those including the A element and not. Two new sparse product matrices are created. It

is convenient that the headers of each are simple copies of the original header $\{A, B, C, D\}$. Since the matrices are sparse there is no computational cost beyond a possible header copy operation. The resulting expression may be written as,

$$\begin{pmatrix} A & B & C & D \\ 1 & 1 & & \\ 1 & & 1 & \end{pmatrix} + \begin{pmatrix} A & B & C & D \\ & & & 1 \end{pmatrix}$$

The A expression is now formally factored out and for a reason that will be made clear below zeros are left in the A column of the first matrix,

$$A \times \begin{pmatrix} A & B & C & D \\ 0 & 1 & & \\ 0 & & 1 & \end{pmatrix} + \begin{pmatrix} A & B & C & D \\ & & & 1 \end{pmatrix}$$

The Factor operation recurses until no further factoring is possible as is the case in the example above for both sparse product matrices. The resulting expression is then,

$$A \times (B + C) + D$$

To introduce further considerations a second example is warrented. Consider the expression,

$$5A + 5^2AB + 5^3ABC$$

The corresponding sparse matrix is,

$$\begin{pmatrix} 5 & A & B & C \\ 1 & 1 & & \\ 2 & 1 & 1 & \\ 3 & 1 & 1 & 1 \end{pmatrix}$$

The 5 and the A are equally common and the tie is broken by the sum of powers which is 6 for the 5 node and only 3 for the A node. The 5 is then factored, but

instead of creating two sparse product matrices only one needs be created since the 5 appears in all rows and the expression becomes,

$$5 \times \begin{pmatrix} 5 & A & B & C \\ 0 & 1 & & \\ 1 & 1 & 1 & \\ 2 & 1 & 1 & 1 \end{pmatrix}$$

Since the A is now most common it is factored and only one sparse product matrix results again,

$$5 \times A \times \begin{pmatrix} 5 & A & B & C \\ 0 & 0 & & \\ 1 & 0 & 1 & \\ 2 & 0 & 1 & 1 \end{pmatrix}$$

Now 5 and B are most common, but 5 has more power. Now the reason for the zero's becomes clear because we must leave a one in the place of the first row,

$$5 \times A \times \Big(\begin{pmatrix} 5 & A & B & C \\ 0 & 0 & & \end{pmatrix} + \begin{pmatrix} 5 & A & B & C \\ 1 & 0 & 1 & \\ 2 & 0 & 1 & 1 \end{pmatrix} \Big)$$

The zero row of the first matrix resolves to a one and the 5 of the second matrix is factored,

$$5 \times A \times \Big(1 + 5 \times \begin{pmatrix} 5 & A & B & C \\ 0 & 0 & 1 & \\ 1 & 0 & 1 & 1 \end{pmatrix} \Big)$$

Next the B is factored and the matrices can be replaced by their equivalent expressions with juxtaposition in place of multiplication,

$$5A\Big(1 + 5B(1 + 5C)\Big)$$

Appendix D

Statistics

Probability Distributions

The goal here is to make a fundamental review of statistical distributions for the purposes of modeling information. The context is data sets gathered from various sources that will be processed mathematically into a standadized format for the purposes of modeling some behaviour. The problem is that, in general, the information is uncertain especially if it's an estimate of future values.

Information theory suggests that when in doubt to use probability distributions with the largest entropy. From thermodynamics we understand entropy to be a function of the number of available system states. In this case the reference is to atoms and molecules, but we may understand states to be abstract system states.

The usual definition of entropy involves the logarithm of the number of states. The base of the logarithm is sometimes 2 so that information may be measured in *bits* and sometimes *e* so that information may be measured in *nats*. For the purposes of probability theory entropy of a discrete random variable X is usually defined to be,

$$\sigma_i = \log \frac{1}{p_i} \text{ for each } i \in \text{Support of } X$$

The idea is that probability distributions with higher measures of entropy are more likely to represent real phenomena since these tend to higher entropy as they proceed. A weak reason as stated to be sure, but we've got to start with something.

Bernoulli Distribution

The fundamental probability distribution is the Bernoulli Distribution. Then,

$$X \sim \mathbf{B}(p) \qquad X \in \{0, 1\} \qquad P(X = 1) = p \qquad P(X = 0) = (1 - p)$$

The entropy of $X \sim B(p)$ is then,

$$\mathbf{E}[\sigma] = p \log(\frac{1}{p}) + (1 - p) \log(\frac{1}{1 - p})$$

The notation needs work, but the implication is that we measure the entropy of a probability distribtion as the expected entropy of it's states.

Binomial Distribution

Suppose we have n random variables $X_1, ..., X_n$ so that $X_i \sim \mathbf{B}(p) \, \forall \, i \in 1..n$. If we define,

$$Y = \sum_i X_i \implies Y \sim \mathbf{B}(n, p)$$

then we automatically have the combinatorial implication,

$$\mathbf{P}[X = x] = \binom{n}{x} p^x (1 - p)^{n-x}$$

Simple calculations find that,

$$\mathbf{E}[X] = np$$
$$\mathbf{Var}[X] = np(1 - p)$$

When n is large (say 20 or more) and p not too close to 0 or 1 then the Normal distribution is a reasonable approximation to the binomial,

$$\text{Let } Z_n = \frac{Y - np}{\sqrt{np(1-p)}}$$
$$\lim_{n\to\infty} Z_n = Z \sim \mathbf{N}(0,1)$$

On the other hand there is a different way to find a limit of the binomial distribution.

$$\text{Let } \{Y_n\} \text{ s.t. } np = \lambda \in \mathcal{R}_{++}$$
$$\lim_{n\to\infty} Y_n = Z \sim \mathbf{Poisson}(\lambda)$$

Let's mention without further comment that the Beta distribution is the conjugate prior.

Geometric Distribution

Usually defined as the probability distribution of the number of Bernoulli trials needed to get one success given support $\{1, 2, ..., n\}$.

Since there's only one way this can happen, we have a simple expression,

$$\mathbf{P}[X = x] = (1-p)^{x-1}p \qquad \forall x \in \mathcal{N}$$

And we have the usual statistics,

$$\mathbf{E}[X] = \frac{1}{p}$$
$$\mathbf{Var}[X] = \frac{1-p}{p^2}$$

The *exponential* distribution is the continuous analog of the geometric distribution.

Negative Binomial Distribution

The probability distribution of the number of successes given r failures of Bernoulli trials, each with probability p of success. We say,

$$X \sim \mathbf{NB}(r, p)$$

First we need to understand the "Stars and Bars" trick. If we have s objects to be divided into b non-negative groups then there are,

$$\binom{s+b-1}{b-1}$$

ways to do this. Imagine $s+b-1$ slots into which you put $b-1$ "bars" and the remaining slots are filled with s "stars". This little game answers the question. This may be pictured for $s = 7, b = 4$ as,

$$**||****|* \implies (2, 0, 4, 1)$$

Now the probability mass function and statistics makes sense as,

$$\begin{aligned} \mathbf{P}[X = k] &= \binom{k+r-1}{r-1}(1-p)^r p^k \\ \mathbf{E}[X] &= rp/(1-p) \\ \mathbf{Var}[X] &= rp/(1-p)^2 \end{aligned}$$

Notice we've assumed that failure occurs on the last Bernoulli trial. The definition could be slightly reworked to allow success to occur on the last trial. If $r = 1$ then we have the special case of the Geometric distribution where we have k successes before first failure. It's all point-of-view. Then $\mathbf{Geom}(1-p) = \mathbf{NB}(1, p)$.

If we extend the definition to real numbers we use Gamma functions in our choice function and call the whole thing a *Polya* distribution.

Furthermore, the Poisson distribution is lurking here,

$$\begin{aligned} \text{Let } \{Y_r\} \text{ s.t. } rp/(1-p) = \lambda \in \mathcal{R}_{++} \\ \lim_{r\to\infty} Y_r = Z \sim \mathbf{Poisson}(\lambda) = \lim_{r\to\infty} \mathbf{NB}(r, \frac{\lambda}{\lambda + r}) \end{aligned}$$

Categorical Distribution

The probability distribution of the identification of n objects amoung N individual categories. We sample from space $\Omega = \{1, ..., n\}$. The probability of seeing object x_i is p_i so that $\sum_i p_i = 1$.

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