

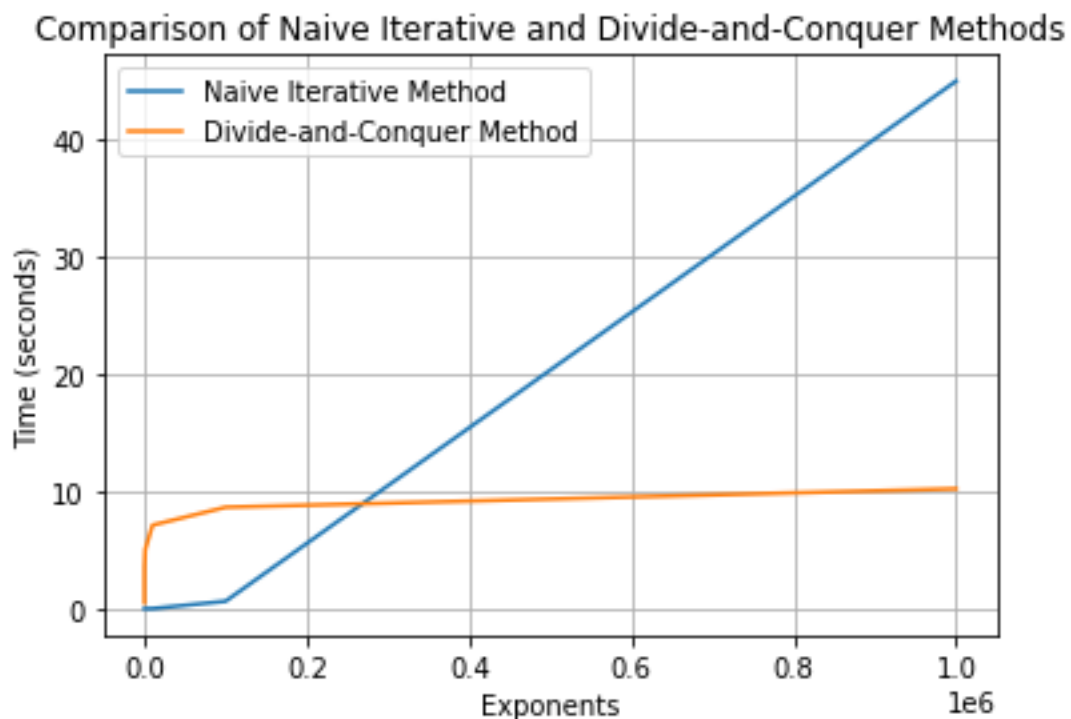
# Analysis\_Assignment1\_Report

## Question1:

B) We have 2 approaches here the Iterative (naive) one and the Divide and conquer one:

Iterative	Divide and conquer
$\Theta(n)$ as we see in the graph and in the code that the function has a linear growth.	$\Theta(\log n)$
	<i>will be proved using master's theorem and we can see it also using the graph.</i>
	<u>Recurrence relation</u> $\begin{cases} T(n) = \Theta(1) & \text{If } (n=1) \\ T(n) = T(n/2) + \Theta(1) + \Theta(1) & \text{Otherwise} \end{cases}$

C)



D)

Proof (Master's Theorem):

$$a=1; b=2; f(n)=\Theta(1)$$

$$n^{**(\log-b(a))} = n^{**(\log-2(1))} = 1$$

$$\text{since } f(n) = \Theta[n^{**(\log-b(a))}] = \Theta[n^{**(\log-2(1))}] = \Theta(1)$$

$$\text{then } T(N) = \Theta[(n^{**\log-b(a)}) * \log(n)] = \Theta(\log(n))$$

Question2

B)

Recurrence relation

$$\{ T(n) = \Theta(1) \quad \text{If } (n=1)$$

$$T(n) = 2T(n/2) + n \quad \text{Otherwise}$$

}

Proof (Master's Theorem):

$$a=2; b=2; f(n)=n$$

$$n^{**(\log-b(a))} = n^{**(\log-2(2))} = n$$

$$\text{since } f(n) = \Theta[n^{**(\log-b(a))}] = \Theta[n^{**(\log-2(2))}] = \Theta(n)$$

$$\text{then } T(N) = \Theta[(n^{**\log-b(a)}) * \log(n)] = \Theta(n * \log(n))$$

C)

