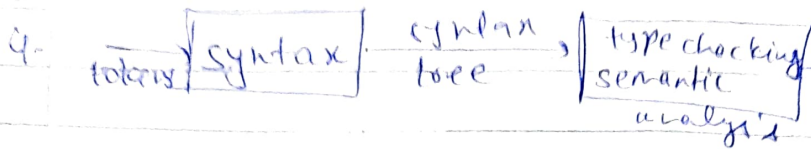


use :- find about domain name info. by using whois

The website to get detailed info about domain name info including its owner, date of registration, expiry, name of registrar, contact info etc.

- Finding IP add.



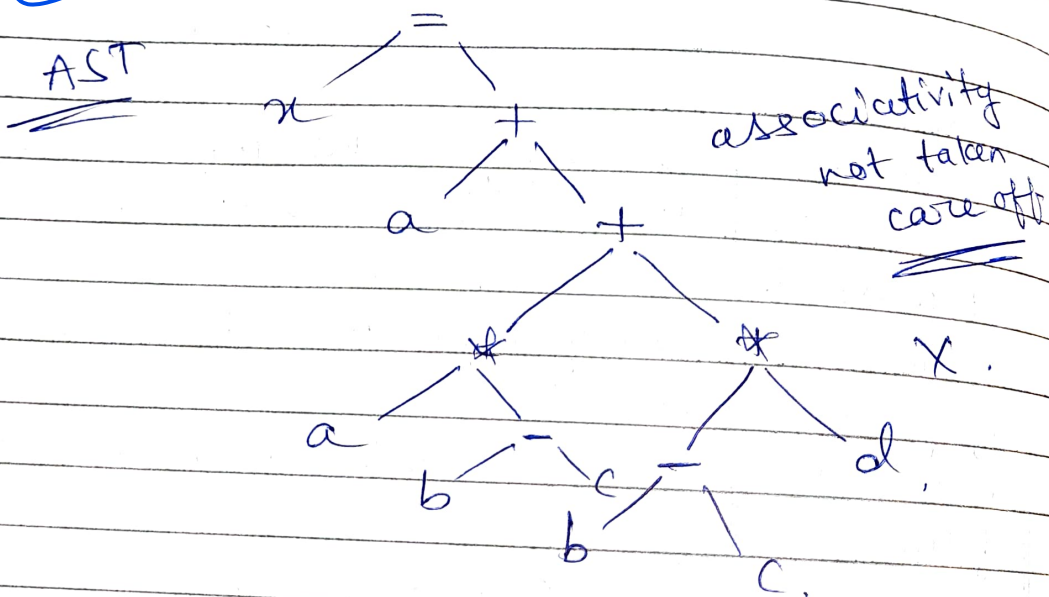
Intermediate Code Representation

- Intermediate Code Representation -

1) Syntax tree.

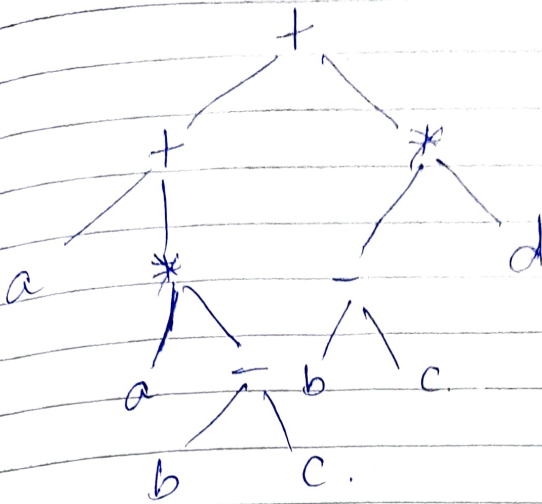
2) DAGs (Directed Acyclic Graphs)

ex :- $x = a + a * (b - c) + (b - c) * d$



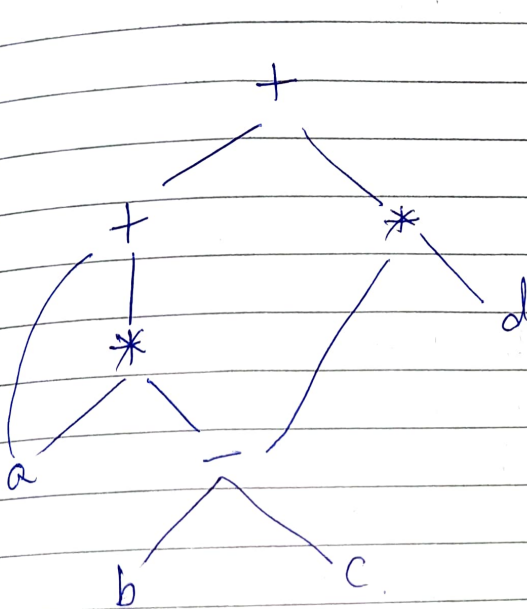
DAG

* '+' left associative



AST

DAG



common subexpressions removed.

3AC (3 address code) { at most 3 addresses }

- { 3 quadruples
- { 3 triples
- { indirect triples

ex: $a = y + x + z$

$t_1 = y + x$

$t_2 = t_1 + z$

$a = t_2$

3AC for given expression

if ($x < y$) then add1 else add2 } more than 3 addresses

C me IR nhi hota.
Java has ICA thing!!

Page No.	
Date	

→ Quadruples -

4 columns / 4 cells

	operators.	arg1	arg2	result
0	-	c		t1
1	*	b	t1	t2
2	-	c		t3
3	*	b	t3	t4
4	+	t2	t4	t5
5	=	t5		a
6				

Why * before +

ex: $a = b * -c + b * -c$

≡ ' - ' then

hence Priority
than Binary

→ triples

	op	arg1	arg2
0	-	c	
1	*	b	(0)
2	-	c	
3	*	b	(2)
4	+	(1)	(3)
5	=	a	(4)

→ Indirect triples

35	(0)
36	(1)
37	(2)
38 101	(3)
39 108	(4)
40 109	(5)
!	!

SSA! - ICG with quadruples & triples, take any expression.

low level
virtual machine

ass-SLR parser
generators

SSA (Single Static Assignment) - [LLVM uses this]

$$p_1 = a + b$$

$$q_1 = p_1 - c$$

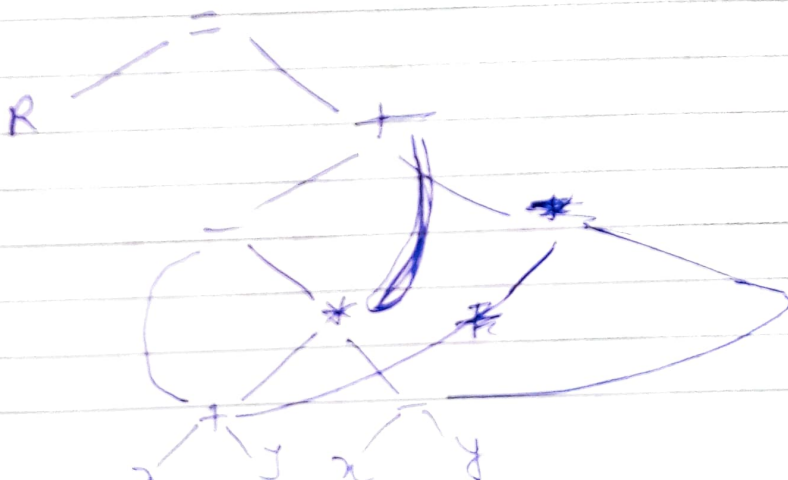
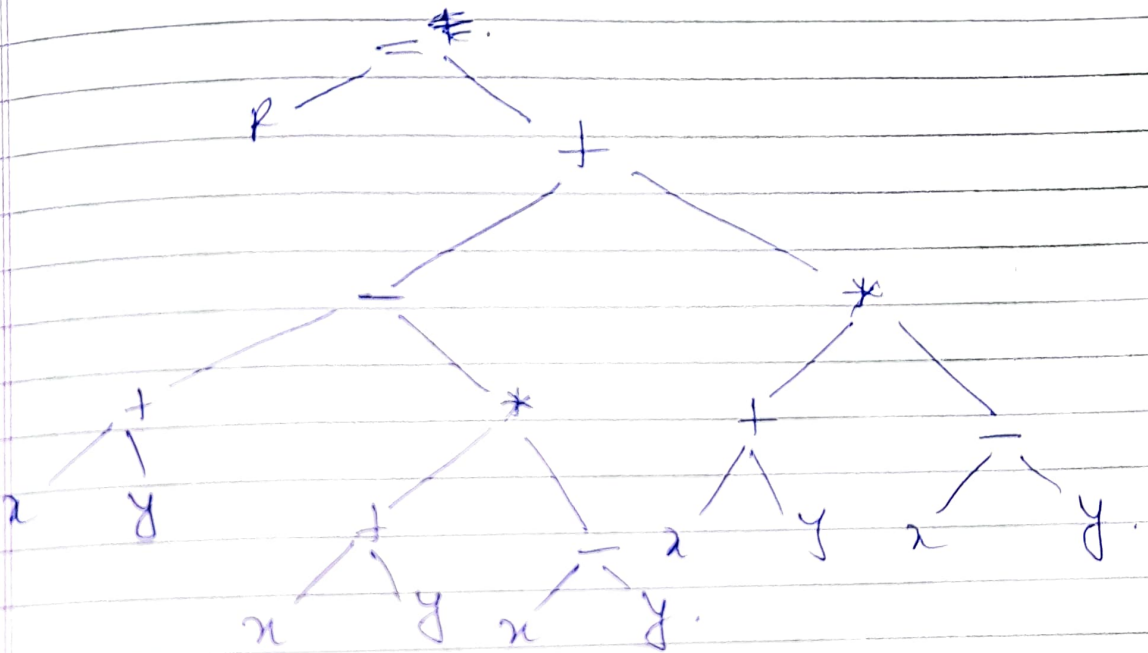
$$p_2 = q_1 * d$$

$$p_3 = e - p_2$$

$$q_2 = p_3 + q_1$$

useful for Dataflow Analysis
at the time of optimization,
or CG.

$$R = \text{exit} \quad ((x+y) - ((x+y) * (x-y))) + ((x+y) * (x-y))$$



$$((x+y) - ((x+y) * (x-y))) + ((x+y) * (x-y))$$

	operator	arg 1	arg 2	result
0	+	x	y	t1
1	-	x	y	t2
2	*	t1	t2	t3
3	+	x	y	t4
4	-	t4	t3	t5
5	+	x	y	t6
6	-	x	y	t7
7	*	t6	t7	t8
8	+	t5	t8	t9
9	=	t9		AR

	operator	arg 1	arg 2
0	+	x	y
1	-	x	y
2	*	(0)	(1)
3	+	x	y
4	-	(3)	(2)
5	+	x	y
6	-	x	y
7	*	(5)	(6)
8	+	(5) (4)	(7)
9	=	(8)	R

Finding IP address

- You can use ping command at your prompt.
- This command is

5) Hybrid cloud migration -
techniques -

1) lift & shift.

2) Cloud bursting.

Process:

Pros: cost optimization, flexibility, managing.

Cons: Network latency.

CC

3AC

Arrays

$$1D \quad x = a[i]$$

$$\text{base} + i \times w$$

$$t_1 = i \times w$$

$$t_2 = a[t_1]$$

{base + t₁}

3AC

To find the address of any element in a 2-Dimensional array there are the following two ways:

Row Major Order

Column Major Order

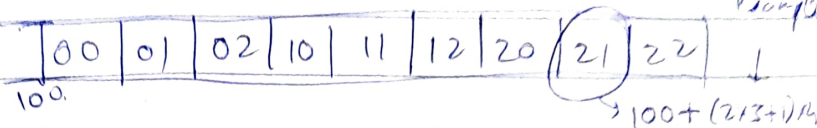
2D Arrays

00	01	02
10	11	12
20	21	22

3x3

$x = a[i][j]$
 $a[\text{row}][\text{col}]$

row major



$$L = a[i][j] = \text{base} + (i \times C + j) * W$$

\downarrow \downarrow \downarrow
 col. width of each element.

$t_1 = i \times C$
 $t_2 = t_1 + j$
 $t_3 = t_2 \times W$
 $t_4 = a[t_3]$
 $x = t_4$

~~NOT~~ $a + t_3$

col major

00	10	20	01	11	21	02	12	22
----	----	----	----	----	----	----	----	----

$$\text{base} + (j \times C + i) * W$$

\downarrow
 row

3D Arrays

$a[i][j][k]$

$a[r][c][b]$

$$\text{Loc } a[i][j][k] = \text{base} + (i \times C \times b + j \times b + k) * W$$

\downarrow \downarrow \downarrow \downarrow
 col. last dim. width

$t_1 = i \times C$
 $t_2 = t_1 \times b$
 $t_3 = j \times b$
 $t_4 = t_2 + t_3$

$t_5 = t_4 + k$
 $t_6 = t_5 * W$
 $t_7 = a[t_6]$
 ~~t_8~~ $x = t_7$

Q

$$t_0 = i \times 1024$$

$$t_1 = j \times 32$$

$$t_2 = k \times 4$$

$$t_3 = t_1 + t_0$$

$$t_4 = t_3 + t_2$$

$$t_5 = X[t_4]$$

what is c & b?

$$\begin{aligned} j \times 32 & \times 1024 \\ j \times 32 & \times 1024 + k \times 4 \\ t_5 & = a[j \times 1024 + k \times 4] \end{aligned}$$

Ans

$$t_5 = X[i \times 1024 + j \times 32 + k \times 4]$$

$$= X[4(i \times 256 + j \times 8 + k)]$$

$$b = 8,$$

$$c =$$

$$(i \times 8 \times 32$$

$$+ j \times 8 + k) \times 4$$

$$C = 32, b = 8$$

• Boolean expression-

operator - &, !, ~

precedence - not > and > or

ex:- a and b or not c

$$t_1 = \text{not } c$$

$$t_2 = a \text{ and } b$$

$$t_3 = t_2 \text{ or } t_1$$

} 3AC

• Conditional statements -

→ if, while, for etc.

→ Jump or goto statements having addresses

if (a < b) then ~~x~~ ¹ else 0

2 ways of doing it

way ① - 100: if ($a < b$) goto 103
 101: $t_1 = 0$ → else part 1st
 102: goto 104 → After
 103: $t_1 = 1$ else, exit
 104:

way ② - L1: if ($a < b$) goto L_t
 $t_1 = 0$
 goto exit
 L_t: $t_1 = 1$.
 exit: ~~exit~~

ex: 2. if ($a < b$) then $X = Y + 1$ else $X = Y - 1$

3PL - 100: if ($a < b$) goto 104
 101: ~~$X = Y - 1$~~ $t_1 = Y - 1$
 102: ~~goto 105~~ $X = t_1$
 103: ~~$X = Y + 1$~~ goto 106
 104: ~~[exit or blank]~~ $t_1 = Y + 1$
 105: $X = t_1$
 106:

Eg value
 $X = a + b$
 ↓ addⁿ
 ↓ SD
 $t_1 = a + b$
 $X = t_1$

ex 3 - while ($c < d$) do
 if ($a < b$) then $X = Y + 1$ else $X = Y - 1$.

100: if ($c < d$) goto 102
 101: goto ~~102~~ → 109
 102: if ($a < b$) goto 106
 103: $t_1 = Y - 1$
 104: $X = t_1$
 105: goto ~~108~~ 100
 106: $t_1 = Y + 1$
 107: $X = t_1$

108: goto 100
 109: exit