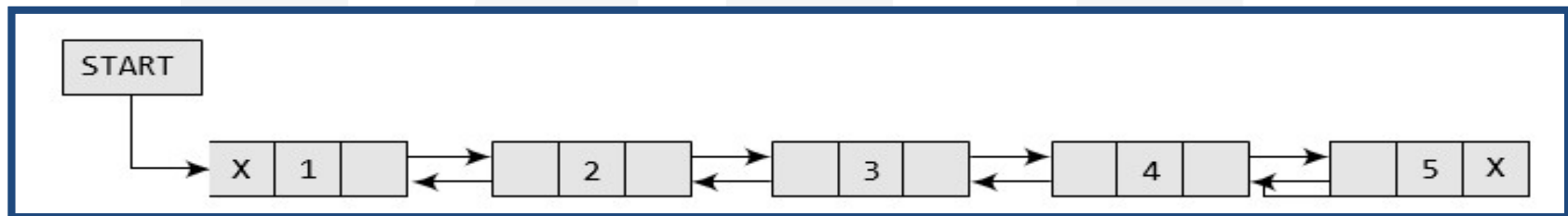


Doubly link list

A doubly linked list or a two-way linked list is a more complex type of linked list which contains a pointer to the next as well as the previous node in the sequence.



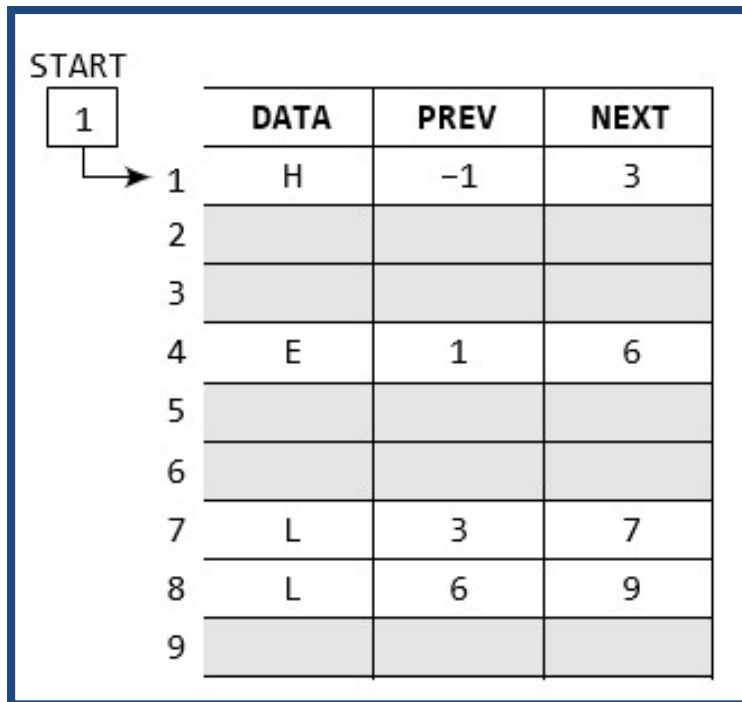
Doubly link list

In C, the structure of a doubly linked list can be given as,

```
struct node
{
    int data;
    struct node *prev;
    struct node *next;
};
```

The PREV field of the first node and the NEXT field of the last node will contain NULL. The PREV field is used to store the address of the preceding node, which enables us to traverse the list in the backward direction.

Memory representation of doubly link list



- START is used to store the address of the first node
- We reach a position where the NEXT entry contains -1 or NULL. This denotes the end of the linked list.
- When we traverse the DATA and NEXT in this manner, we will finally see that the linked list in the above example stores characters that when put together form the word HELLO.

Insertion in doubly link list

Case 1: The new node is inserted at the beginning.

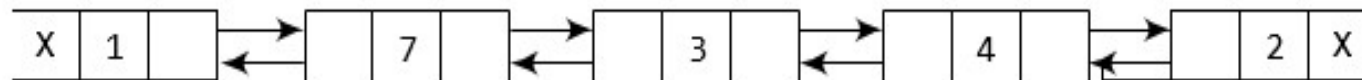
Case 2: The new node is inserted at the end.

Case 3: The new node is inserted after a given node.



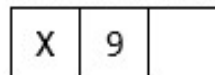
Case 1: The new node is inserted at the beginning.

Consider the doubly linked list shown in Fig. 6.39. Suppose we want to add a new node with data 9 as the first node of the list. Then the following changes will be done in the linked list.

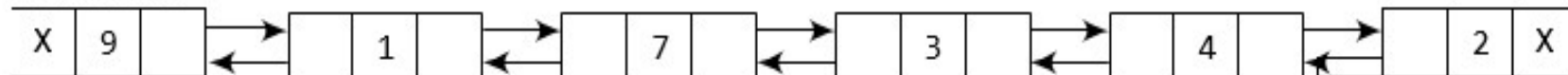


START

Allocate memory for the new node and initialize its DATA part to 9 and PREV field to NULL.



Add the new node before the START node. Now the new node becomes the first node of the list.



START

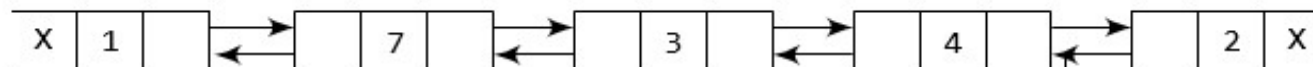
Case 1: The new node is inserted at the beginning.

```
Step 1: IF AVAIL = NULL
        Write OVERFLOW
        Go to Step 9
    [END OF IF]
Step 2: SET NEW_NODE = AVAIL
Step 3: SET AVAIL = AVAIL -> NEXT
Step 4: SET NEW_NODE -> DATA = VAL
Step 5: SET NEW_NODE -> PREV = NULL
Step 6: SET NEW_NODE -> NEXT = START
Step 7: SET START -> PREV = NEW_NODE
Step 8: SET START = NEW_NODE
Step 9: EXIT
```

- In Step 1, we first check whether memory is available for the new node. If the free memory has exhausted, then an OVERFLOW message is printed.
- Otherwise, if free memory cell is available, then we allocate space for the new node. Set its DATA part with the given VAL and the NEXT part is initialized with the address of the first node of the list, which is stored in START.
- Now, since the new node is added as the first node of the list, it will now be known as the START node, that is, the START pointer variable will now hold the address of NEW_NODE.

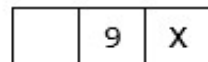


Case 2: The new node is inserted at the end.

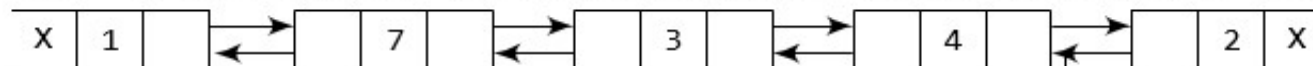


START

Allocate memory for the new node and initialize its DATA part to 9 and its NEXT field to NULL.



Take pointer variable PTR and make it point to the first node of the list.



START, PTR

Move PTR so that it points to the last node of the list. Add the new node after the node pointed by PTR.



START

PTR

Suppose we want to add new node with data 9 as the last node of the list. Then the following changes will be done in the linked list.

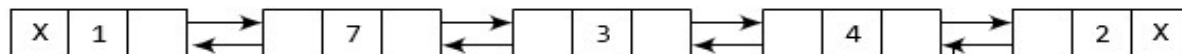
Case 2: The new node is inserted at the end.

```
Step 1: IF AVAIL = NULL
        Write OVERFLOW
        Go to Step 11
    [END OF IF]
Step 2: SET NEW_NODE = AVAIL
Step 3: SET AVAIL = AVAIL -> NEXT
Step 4: SET NEW_NODE -> DATA = VAL
Step 5: SET NEW_NODE -> NEXT = NULL
Step 6: SET PTR = START
Step 7: Repeat Step 8 while PTR -> NEXT != NULL
Step 8:     SET PTR = PTR -> NEXT
    [END OF LOOP]
Step 9: SET PTR -> NEXT = NEW_NODE
Step 10: SET NEW_NODE -> PREV = PTR
Step 11: EXIT
```

- In Step 6, we take a pointer variable PTR and initialize it with START.
- In the while loop, we traverse through the linked list to reach the last node.
- Once we reach the last node, in Step 9, we change the NEXT pointer of the last node to store the address of the new node.
- Remember that the NEXT field of the new node contains NULL which signifies the end of the linked list. The PREV field of the NEW_NODE will be set so that it points to the node pointed by PTR (now the second last node of the list).

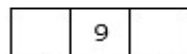


Case 3: The new node is inserted after a given node.

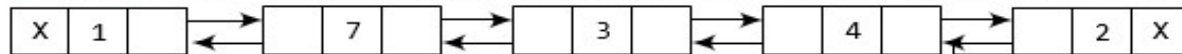


START

Allocate memory for the new node and initialize its DATA part to 9.

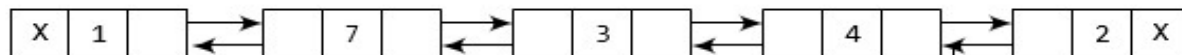


Take a pointer variable PTR and make it point to the first node of the list.



START, PTR

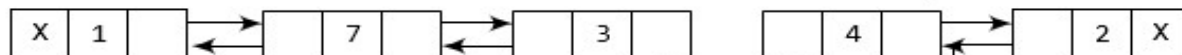
Move PTR further until the data part of PTR = value after which the node has to be inserted.



START

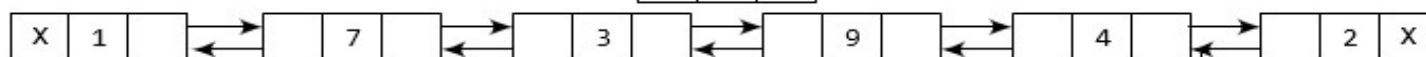
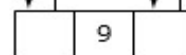
PTR

Insert the new node between PTR and the node succeeding it.



START

PTR



START

Case 3: The new node is inserted after a given node.

```
Step 1: IF AVAIL = NULL
        Write OVERFLOW
        Go to Step 12
    [END OF IF]
Step 2: SET NEW_NODE = AVAIL Step
3: SET AVAIL = AVAIL -> NEXT Step
4: SET NEW_NODE -> DATA = VAL Step
5: SET PTR = START
Step 6: Repeat Step 7 while PTR -> DATA != NUM
Step 7:     SET PTR = PTR -> NEXT
    [END OF LOOP]
Step 8: SET NEW_NODE -> NEXT = PTR -> NEXT
Step 9: SET NEW_NODE -> PREV = PTR
Step 10: SET PTR -> NEXT = NEW_NODE
Step 11: SET PTR -> NEXT -> PREV = NEW_NODE
Step 12: EXIT
```

- In Step 5, we take a pointer PTR and initialize it with START. That is, PTR now points to the first node of the linked list.
- In the while loop, we traverse through the linked list to reach the node that has its value equal to NUM. We need to reach this node because the new node will be inserted after this node.
- Once we reach this node, we change the NEXT and PREV fields in such a way that the new node is inserted after the desired node.

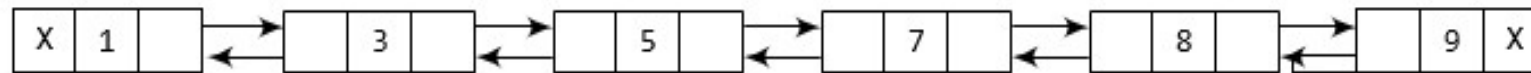
Deletion in doubly link list

- Case 1: The first node is deleted.
- Case 2: The last node is deleted.
- Case 3: The node after a given node is deleted.



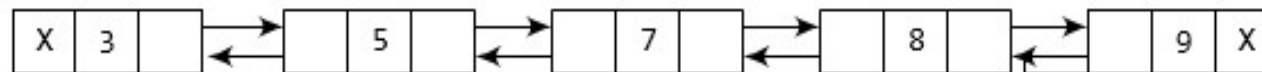
Case 1: The first node is deleted in doubly link list

Consider the doubly linked list shown in Fig. 6.47. When we want to delete a node from the beginning of the list, then the following changes will be done in the linked list.



START

Free the memory occupied by the first node of the list and make the second node of the list as the START node.



START

Case 1: The first node is deleted in doubly link list

```
Step 1: IF START = NULL
        Write UNDERFLOW
        Go to Step 6
    [END OF IF]
Step 2: SET PTR = START
Step 3: SET START = START -> NEXT
Step 4: SET START -> PREV = NULL
Step 5: FREE PTR
Step 6: EXIT
```

- In Step 1 of the algorithm, we check if the linked list exists or not. If $START = NULL$, then it signifies that there are no nodes in the list and the control is transferred to the last statement of the algorithm.
- However, if there are nodes in the linked list, then we use a temporary pointer variable PTR that is set to point to the first node of the list.
- For this, we initialize PTR with START that stores the address of the first node of the list.
- In Step 3, START is made to point to the next node in sequence and finally the memory occupied by PTR (initially the first node of the list) is freed and returned to the free pool

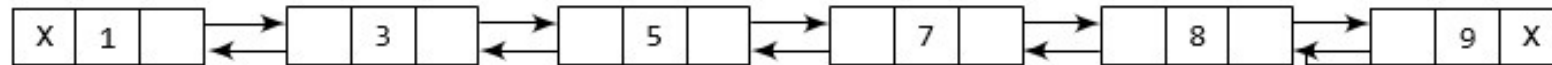


Case 2: The last node is deleted in doubly link list



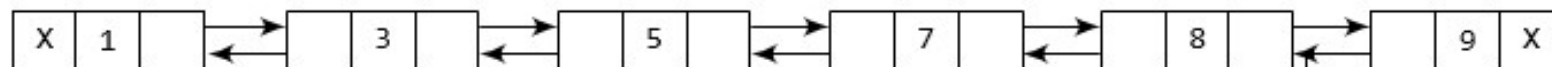
START

Take pointer variable PTR that points to the first node of the list.



START, PTR

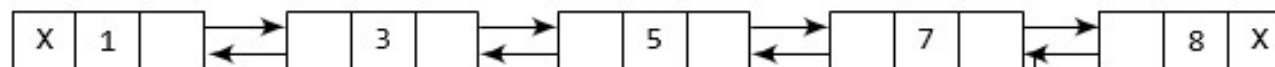
Move PTR so that it now points to the last node of the list.



START

PTR

Free the space occupied by the node pointed by PTR and store NULL in NEXT field of its preceding node.



START

Case 2: The last node is deleted in doubly link list

```
Step 1: IF START = NULL
        Write UNDERFLOW
        Go to Step 7
    [END OF IF]
Step 2: SET PTR = START
Step 3: Repeat Step 4 while PTR -> NEXT != NULL
Step 4:     SET PTR = PTR -> NEXT
    [END OF LOOP]
Step 5: SET PTR -> PREV -> NEXT = NULL
Step 6: FREE PTR
Step 7: EXIT
```

- In Step 2, we take a pointer variable PTR and initialize it with START. That is, PTR now points to the first node of the linked list.
- The while loop traverses through the list to reach the last node. Once we reach the last node, we can also access the second last node by taking its address from the PREV field of the last node.
- To delete the last node, we simply have to set the next fi of second last node to NULL, so that it now becomes the (new) last node of the linked list. The memory of the previous last node is freed and returned to the free pool.

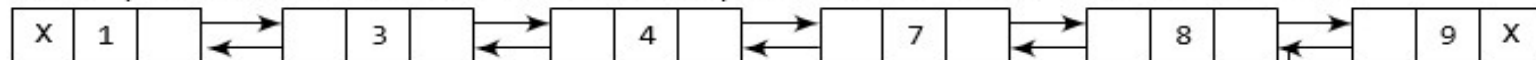


Case 3: The node after a given node is deleted.



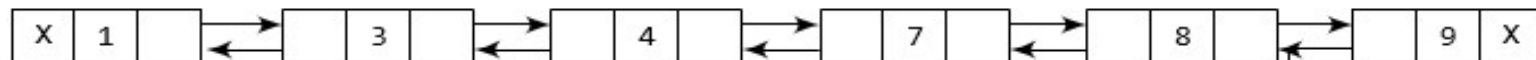
START

Take pointer variable PTR and make it point to the first node of the list.



START, PTR

Move PTR further so that its data part is equal to the value after which the node has to be inserted.



START

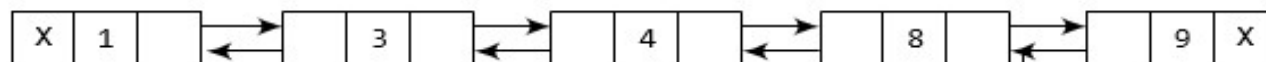
PTR

Delete the node succeeding PTR.



START

PTR



START

Case 3: The node after a given node is deleted.

```
Step 1: IF START = NULL
        Write UNDERFLOW
        Go to Step 9
    [END OF IF]
Step 2: SET PTR = START
Step 3: Repeat Step 4 while PTR -> DATA != NUM
Step 4:     SET PTR = PTR -> NEXT
    [END OF LOOP]
Step 5: SET TEMP = PTR -> NEXT
Step 6: SET PTR -> NEXT = TEMP -> NEXT
Step 7: SET TEMP -> NEXT -> PREV = PTR
Step 8: FREE TEMP
Step 9: EXIT
```

- In Step 2, we take a pointer variable PTR and initialize it with START. That is, PTR now points to the first node of the doubly linked list. The while loop traverses through the linked list to reach the given node.
- Once we reach the node containing VAL, the node succeeding it can be easily accessed by using the address stored in its NEXT field.
- The NEXT field of the given node is set to contain the contents in the NEXT field of the succeeding node. Finally, the memory of the node succeeding the given node is freed and