

Total: 15 marks

Duration: 25 mins

*open books and notes, no mobile phones, friends disconnected during exam***Be precise, no marks for vague answers**

Roll No:

Q1. Select the correct answer(s) for the following questions. No partial marks. (2×4)

- (i) In a uni-processor system, which of the following statement(s) are true regarding threads of a multi-threaded process?
- (a) More than one thread can be inside system call handler only if they have separate OS stacks.
  - (b) More than one thread can be in RUNNING state if they have separate OS stacks.
  - (c) More than one thread can be inside system call handler if interrupts are disabled during system calls.
  - (d) More than one thread can be inside system call handler if context switching is disabled during system calls.
  - (e) None of the above.

**Ans: a,c**

- (ii) Which of the following statement(s) are true?

- (a) If a process is interrupted during system call handling, the OS will crash.
- (b) If a process is de-scheduled during system call handling, it must resume its execution in OS mode.
- (c) If a process is waiting for an I/O event, it must resume its execution in OS mode.
- (d) A multitasking OS can disable interrupts whenever the CPU is executing in OS mode.
- (e) None of the above.

**Ans: c**

- (iii) Which of the following statement(s) are true?

- (a) Multi-level feedback queue scheduling favors I/O bound processes.
- (b) Starvation problem in multi-level feedback queue scheduling can be addressed by carefully tweaking the value of time quantum.
- (c) Dynamically determining the value of time quantum can address the issue of a malicious user tricking the multi-level feedback queue scheduler.
- (d) Standard deviation of waiting time can be used to measure the fairness of any scheduling algorithm.
- (e) For a given scheduling algorithm, average waiting time monotonically increases with number of processes.

**Ans: a, b, c, d**

- (iv) Which of the following statement(s) are true?

- (a) There can be a common signal handler function for multiple signals.
- (b) If two processes execute `kill(pid, SIGSEGV)` around the same time, one of the calls may fail.
- (c) If two processes execute `kill(pid, SIGSEGV)` around the same time, the target process (with `PID=pid`) may execute the registered signal handler twice.
- (d) If two processes execute `kill(pid, SIGSEGV)` around the same time, the target process (with `PID=pid`) may execute the registered signal handler only once.
- (e) Within a signal handler function, a signal can be generated to another processes using the `kill` system call.

**Ans: a, c, d, e**

Q2. Consider the following C program. (4)

```
1 #include<stdio.h>
2 #include<stdlib.h>
3 #include<signal.h>
4 char *ptr1 = NULL;
5 char *ptr2 = NULL;
6 void sighandler(int signo)
7 {
8     printf("I am in sighandler\n");
9     ptr1 = ptr2;
```

```

10 }
11
12 main()
13 {
14     if(signal(SIGSEGV, &sighandler) < 0)
15         perror("signal");
16
17     ptr2 = (char *)malloc(100);
18     ptr2[0] = 'A';
19     printf("%c\n", *ptr1);
20 }

```

What will be the output of this program when executed in a UNIX/Linux OS? Explain.

**Ans:** The program will print **I am in sighandler** infinitely.

Reason: The program counter (RIP) will point to the instruction that dereferences memory location pointed to by **ptr1**. The memory address **ptr1** is stored in a register (e.g., RAX) before it can be dereferenced. Note that, you have seen this behavior in Assignment-II for page fault handler implementation where if the register state is not saved and restored, the accessed memory location changed when the instruction gets re-executed after exception handling. Therefore, after the page fault and **sighandler** execution, the value of the register involved remain unchanged and causes repeated SEGFAULTs.

Q3. Consider the following C program. (3)

```

1 #include<stdio.h>
2 #include<stdlib.h>
3 #include<unistd.h>
4 #include<fcntl.h>
5 main()
6 {
7     char buf[8];
8     int fd, dfd;
9     fd = open("some.txt", O_RDWR);
10    if(fd < 0){
11        perror("open");
12        exit(-1);
13    }
14    dfd = dup(fd);
15    buf[7] = 0;
16    read(fd, buf, 7);
17    printf("%s\n", buf);
18    read(dfd, buf, 7);
19    printf("%s\n", buf);
20 }
~

```

Assume that **some.txt** is a valid file whose content is: **AAAAAAAABBBBBBBBCCCCCCC** . . . What will be the output of the above program when executed in the UNIX OS? (no explanation needed)

**Ans:**

AAAAAAA

BBBBBBB

The **dup** system call simply create a duplicate file descriptor, the file object pointed to by the two file descriptors remains the same. (See slide-17, IPC.pdf)