# ARTIFICIAL ÎNTELLÎGENCE.

### LAB4

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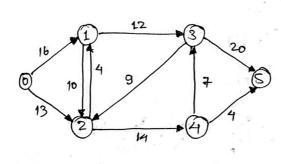
ATM: To study, understand and implement applications of BFS (Breadth First Search) le DFS (Depth Fint Search).

· Breadth First Search:

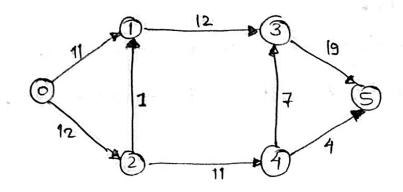
PROBLEM STATEMENT: Given a graph which represents a flow network where every edge has a capacity. Also given two vertices source's" le sink 't' in the graph, find the maximum possible flow from s to t with following constraints:

- · Flow on an edge doesn't exceed the given capacity of the edge.
- · Incoming flow is equal to outgoing flow for every vester except s and t.

PROBLEM SOLVINGS



The maximum possible flow in the above graph is 23.



#### ALGORITHM :

The following is a simple idea of Food-Fulkerson Algorithms

- 1) start with initial flow as 0.
- 2) while there is a augmenting path from sounce to sink.

  add this path flow to flow.
  - 3) return flow.

Residual Graph of a flow retwork is a graph which indicates additional possible flow. Every eage of residual graph hos a value called residual capacity which is equal to the original capacity of the edge minus current flow.

- · Residual capacity is v if there is no edge between two vertices.
- . We can initeate the residual graph as original graph as there is no initeal flow and initially residual capacity is equal to original capacity.
- · BFS builds the parent arr [] using which we traverse through the found parth and find possible flow through this path by finding minimum restual capacity away the path.
- · we need to update residual capacities, we subtract path flow from all edges along the path and we add path flow along the reverse edges.

```
# Maximum flow in a network:
from collections import defaultdict
class Graph:
    def init (self, graph):
        self.graph = graph
        self. ROW = len(graph)
    def BFS(self, s, t, parent):
        visited = [False] * (self.ROW)
        queue = []
        queue.append(s)
        visited[s] = True
        while queue:
            u = queue.pop(0)
            for ind, val in enumerate(self.graph[u]):
                if visited[ind] == False and val > 0:
                    queue.append(ind)
                    visited[ind] = True
                    parent[ind] = u
                    if ind == t:
                        return True
        return False
    def FordFulkerson(self, source, sink):
        parent = [-1]*(self.ROW)
        \max flow = 0
```

```
while self.BFS(source, sink, parent) :
            path flow = float("Inf")
            s = sink
            while(s != source):
                path flow = min (path flow, self.graph[parent[s]][s])
                s = parent[s]
            max_flow += path_flow
            v = sink
            while(v != source):
                u = parent[v]
                self.graph[u][v] -= path flow
                self.graph[v][u] += path flow
                v = parent[v]
        return max_flow
graph = [[0, 16, 13, 0, 0, 0],
        [0, 0, 10, 12, 0, 0],
        [0, 4, 0, 0, 14, 0],
        [0, 0, 9, 0, 0, 20],
        [0, 0, 0, 7, 0, 4],
        [0, 0, 0, 0, 0, 0]]
g = Graph(graph)
source = 0; sink = 5
print ("The maximum possible flow is %d " % g.FordFulkerson(source,
sink))
```



Depth filst Search:

PROBLEM FORMULATION:

Lexicographic sosting of a given set of keys. Given a set of strings, redurn them in alphabetical order.

## PROBLEM SOLVING:

Lexi cographic sorting of a set of teys can be accomplished with a simple trie-based algorithm by:

- · Insert all keys that a type.
- · point all beys in the thie by performing pre-order traversal on trie to get output in alphabetically increasing order.

#### ALGORITHM:

To print the strings in algorithmic alphabetical order we have to first insert them in the trie and then perform traversal to print in alphabetical order.

The node of the contain an index[] array which stores the index position of all the strings of arr[].

ending at that wade.

Except for the trie's left node all the other nodes have the stre O for the index[] array.

```
#lexicographic sorting
class Trie:
   def init (self):
        self.key = None
        self.character = [None] * 26
def insert(head, s):
    # start from the root node
   curr = head
   for c in s:
        key = ord(c) - ord('a')
        # create a new node if the path doesn't exist
        if curr.character[key] is None:
            curr.character[key] = Trie()
        # go to the next node
        curr = curr.character[key]
    # store key in the leaf node
   curr.key = s
# Function to perform preorder traversal on a given Trie
def preorder(curr):
    # return if Trie is empty
    if curr is None:
        return
   for i in range (26):
        if curr.character[i]:
            # if the current node is a leaf, print the key
            if curr.character[i].key:
                print(curr.character[i].key)
            preorder(curr.character[i])
```

```
if name == ' main ':
    # given set of keys
   words = [
        'lexicographic', 'sorting', 'of', 'a', 'set', 'of', 'keys',
'can', 'be',
        'accomplished', 'with', 'a', 'simple', 'trie', 'based',
'algorithm',
        'we', 'insert', 'all', 'keys', 'in', 'a', 'trie', 'output',
'all',
        'keys', 'in', 'the', 'trie', 'by', 'means', 'of', 'preorder',
        'traversal', 'which', 'results', 'in', 'output', 'that',
'is', 'in',
        'lexicographically', 'increasing', 'order', 'preorder',
'traversal',
        'is', 'a', 'kind', 'of', 'depth', 'first', 'traversal'
   ]
   head = Trie()
    for word in words:
        insert(head, word)
   preorder(head)
■ dfs.py × ⊕
```

