

# **Artificial Intelligence**

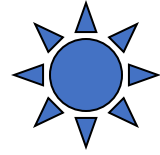
## **DFS and BFS**

**Zahoor Tanoli (PhD)**

**CUI Attock Campus**

# You will learn

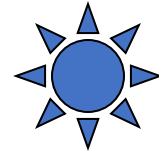
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- What is search and why it is needed
- What is uninformed searches
- Understanding and implementing DFS and BFS
- Time and space complexity
- Informed search introduction

# Simple Search Algorithm

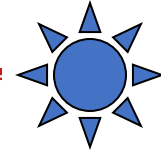
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The process of deciding  
what actions and states to consider  
E.g., driving Attock → Lahore  
in-between states and actions defined  
States: Some places in Attock & Lahore  
Actions: Turn left, Turn right, go straight,  
accelerate & brake, etc.

# Search Algorithm

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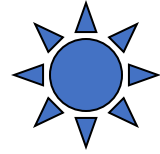


Because there are many ways to achieve the same goal

- Those ways are together expressed as a tree
- Multiple options of unknown value at a point,
  - the agent can examine different possible sequences of actions, and choose the best
- This process of looking for the best sequence is called ***search***
- The best sequence is then a list of actions, called ***solution***

# Search Algorithm

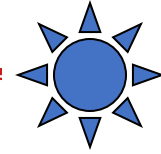
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- Defined as
  - taking a problem
  - and returns a solution
- Once a solution is found
  - the agent follows the solution
  - and carries out the list of actions – execution phase
- Design of an agent
  - “Formulate, search, execute”

# Simple Search Algorithm

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A search node is a path from some state X to the start state, e.g., (XBAS)

The state of a search node is the most recent state of the path, e.g. X.

Let Q be a list of search nodes, e.g. ((XBAS) (CBAS)...) )

Let B be the start state

1. **Initialize Q with search node (S) as only entry; set Visited = (S)**
2. **If Q is empty, fail. Else, pick some search node N from Q**
3. **If state (N) is a goal, return N (we've reached the goal)**
4. **(otherwise) Remove N from Q**
5. **Find all the descendants of state(N) not in Visited and create all the one-step extensions of N to each descendant.**
6. **Add the extended paths to Q; add children of state(N) to Visited**
7. **Go to step 2.**

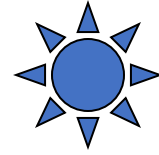
## Critical Decisions:

**Step 2: Pick N from Q**

**Step 6: adding extension of N to Q**

# Implementing the Search Strategies

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## Depth-First

Pick first element of Q

Add path extensions to front of Q

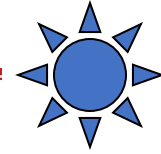
## Breadth-first

Pick first element of Q

Add path extensions to end of Q

# Testing for the goal

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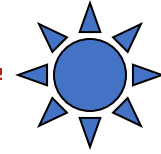


- This algorithm stops (in step 3) when  $\text{state}(N)=G$  or, in general, when  $\text{state}(N)$  satisfies the goal test.
- We could have performed this test in step 6 as each extended path is added to  $Q$ . This would catch termination earlier and be perfectly correct.
- However, performing the test in step 6 will be incorrect for the optimal searches. I have chosen to leave the test in step 3 to maintain uniformity with these future searches.



# Definition

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## Visited *(not expanded)*

a state  $M$  is first visited when a path to  $M$  first gets added to  $Q$ . In general, a state is said to have been visited if it has ever shown up in a search node in  $Q$ . The intuition is that we have briefly “visited” them to place them on  $Q$ , but we have not yet examined them carefully.

## Expanded

a state  $M$  is expanded when it is the state of the search node that is pulled off of  $Q$ . At that point, the descendants of  $M$  are visited and the path that led to  $M$  is extended to the eligible descendants. In principle, a state may be expanded multiple times. **We sometimes refer to the search node that led to  $M$  (instead of  $M$  itself) as being expanded.** However, once a node is expanded we are done with it. We will not need to expand it again. In fact, we discard it from  $Q$ .

# Visited States

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Keeping track of visited states generally improves time efficiency when searching graphs, without affecting correctness.

*Space may still be a problem to keep track of all visited states*

If all we want to do is find a path from the start to the goal, there is no advantage to adding a search node whose state is already the state of another search node

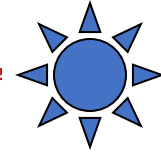
This would be the ideal way of avoiding loops in search graphs/trees

Any state reachable from the node the second time would have been reachable from that node the first time

Note that, when using Visited, each state will only ever have at most one path to it (search node) in Q

# Implementation Issues

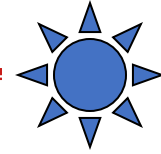
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- Although we will be speaking about Visited list, this is never the preferred implementation – needs continuous checking if any node is present in the list
- If the graph states are known ahead of time, as an explicit set then space is allocated in the state itself to keep a mark; which makes both adding to Visited and checking if a state is Visited a constant time operation – like if we have a data structure we can just add a flag bit
- Alternatively, as is more common in AI, if the states are generated on the fly, then a hash table may be used for efficient detection of previously visited states – it still require some space and constant access time
- Note that, in any case, the incremental space cost of the Visited list will be proportional to the number of states – which can be very high in some problems

# Terminology

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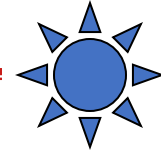
**Heuristic** – The word generally refers to a “rule of thumb” something that may be helpful in some cases but not always. Generally held to be in contrast to “guaranteed” or “optimal”

**Heuristic Function** – In search terms, a function that computes a value for a state (but does not depend on any path to that state) that may be helpful in guiding the search.

**Estimated Distance to Goal** – this type of heuristic function depends on the state and the goal. The classic example is straight-line distance used as an estimate for actual distance in a road network. This type of information can help increase the efficiency of the search

# Implementing the Search Strategies

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## Depth first

- Pick first element of Q
- Add path extensions to front of Q

## Breadth first

- Pick first element of Q
- Add path extensions to end of Q

## Best first (Informed) or greedy *uses heuristics*

*Not guaranteed of finding the best or optimal path ☹*

- Pick “best” (measured by heuristic value of state) element of Q
- Add path extensions anywhere in Q (It may be more efficient to keep the Q ordered in some way so as to make it easier to find the “best” element

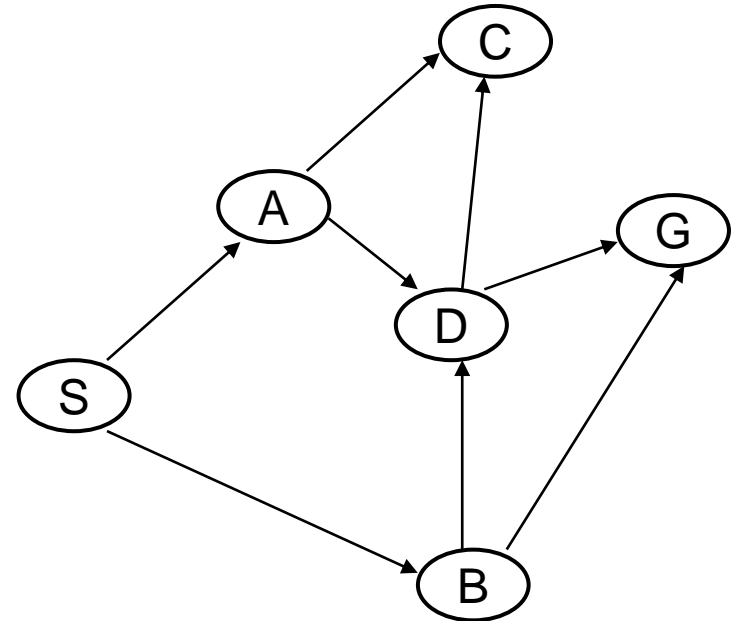
**In worst case it would examine all the same paths that DFS and BFS would examine – obviously, order will not be the same and so do the path**

# Depth First Search



Pick First element of Q; Add path extensions to front of Q

	Q	Visited
1		
2		
3		
4		
5		



Added paths in **blue**

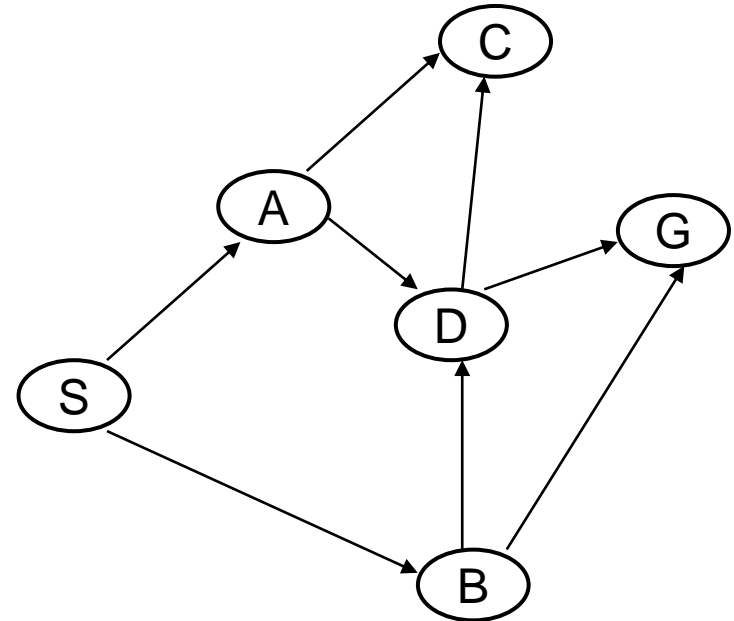
Path is shown in reversed order, the node's state is the first entry

# Depth First Search



Pick First element of Q; Add path extensions to front of Q

	Q	Visited
1	(S)	S
2		
3		
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5		



Added paths in blue

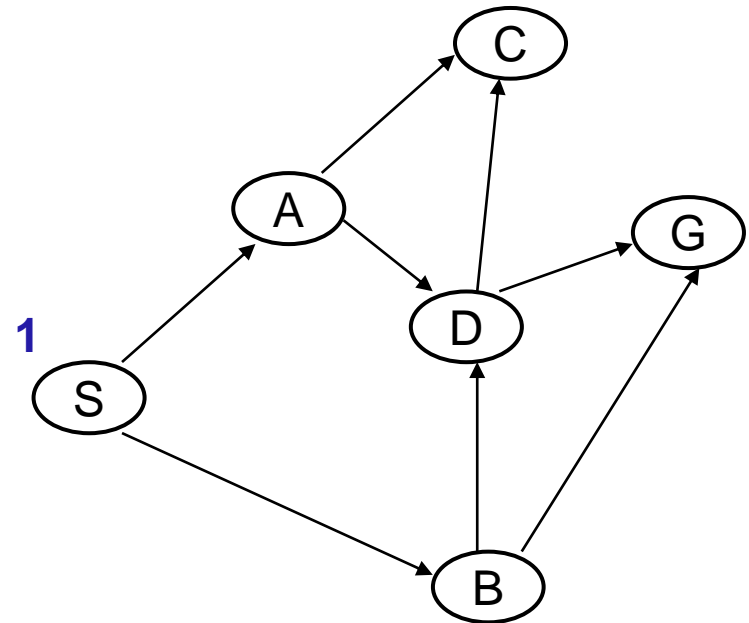
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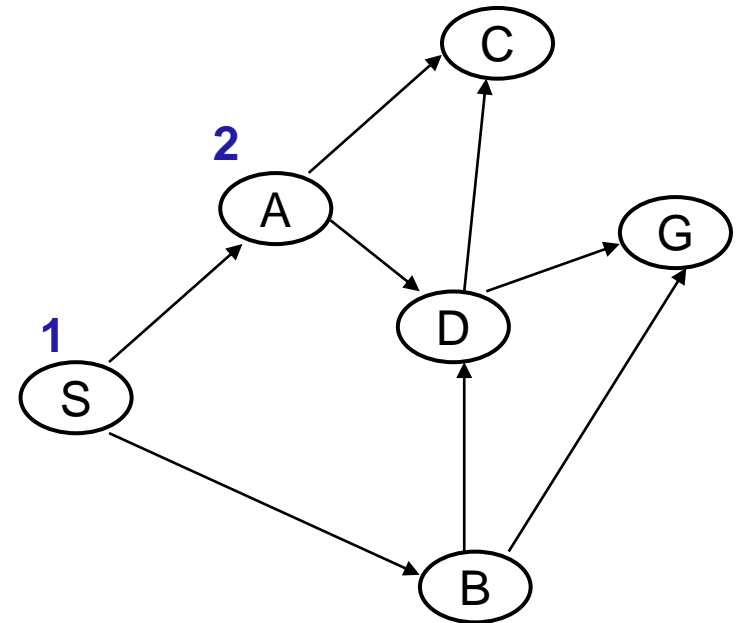


# Depth First Search



Pick First element of Q; Add path extensions to front of Q

	Q	Visited
1	(S)	S
2	(A S) (B S)	A, B, S
3		
4		
5		



Added paths in blue

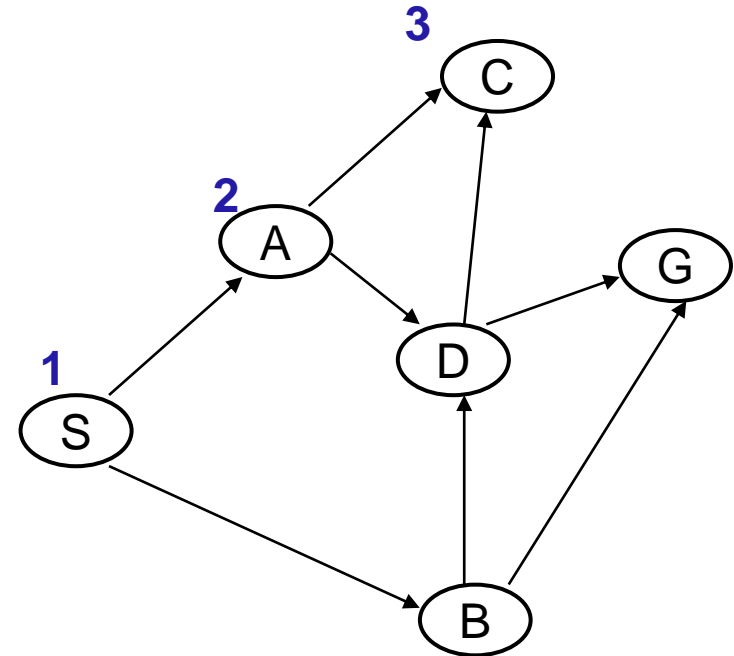
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# Depth First Search



Pick First element of Q; Add path extensions to front of Q

	Q	Visited
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3	(C A S) (D A S) (B S)	C, D, B, A, S
4		
5		



Added paths in **blue**

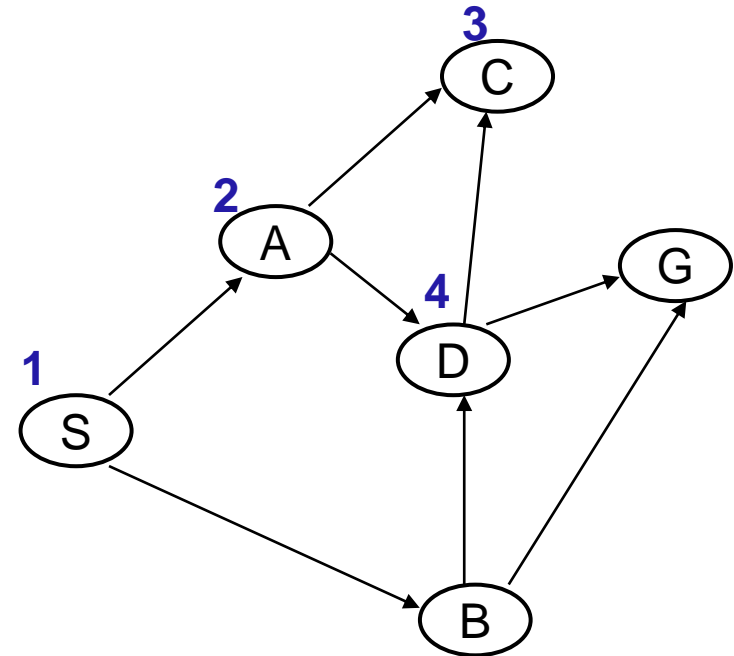
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Pick First element of Q; Add path extensions to front of Q

	Q	Visited
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4	(D A S) (B S)	C, D, B, A, S
5		



Added paths in **blue**

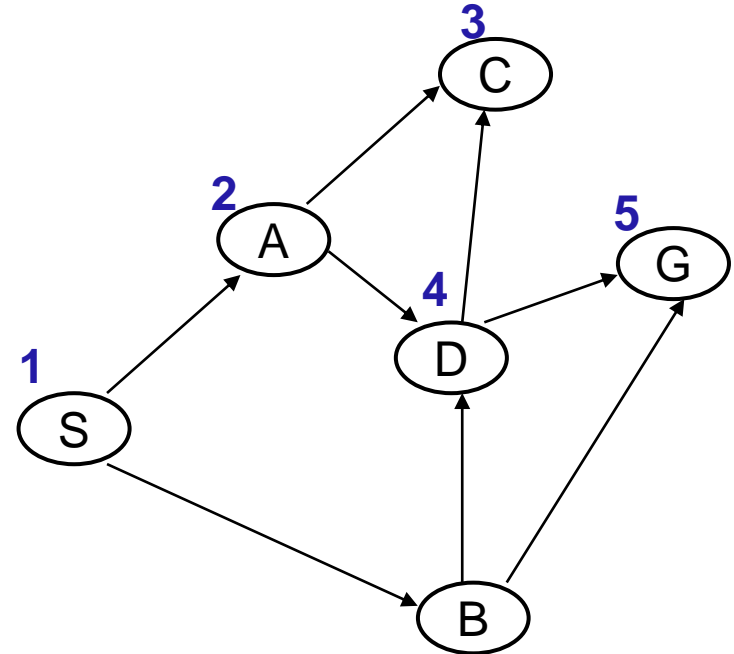
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# Depth First Search



Pick First element of Q; Add path extensions to front of Q

	Q	Visited
1	(S)	S
2	(A S) (B S)	A, B, S
3	(C A S) (D A S) (B S)	C, D, B, A, S
4	(D A S) (B S)	C, D, B, A, S
5	(G D A S) (B S)	G, C, D, B, A, S



Added paths in **blue**

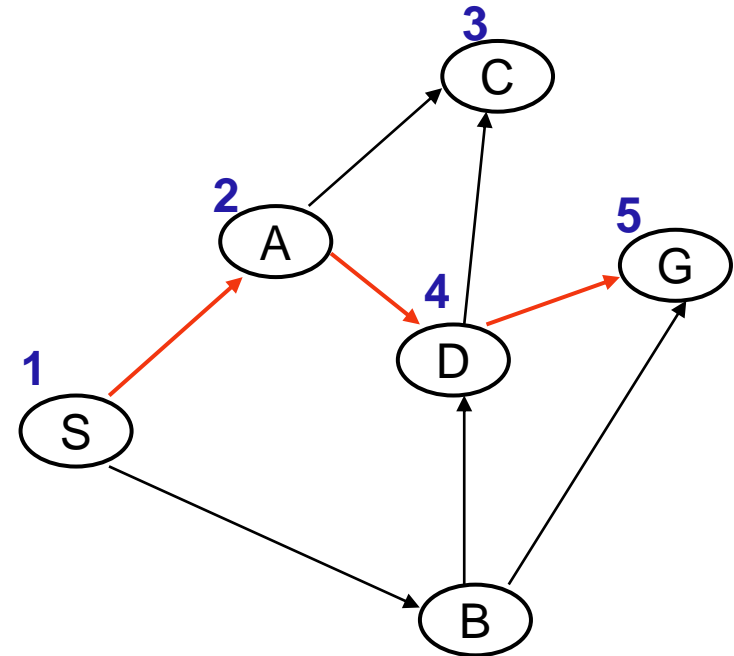
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# Depth First Search



Pick First element of Q; Add path extensions to front of Q

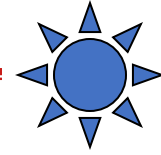
	Q	Visited
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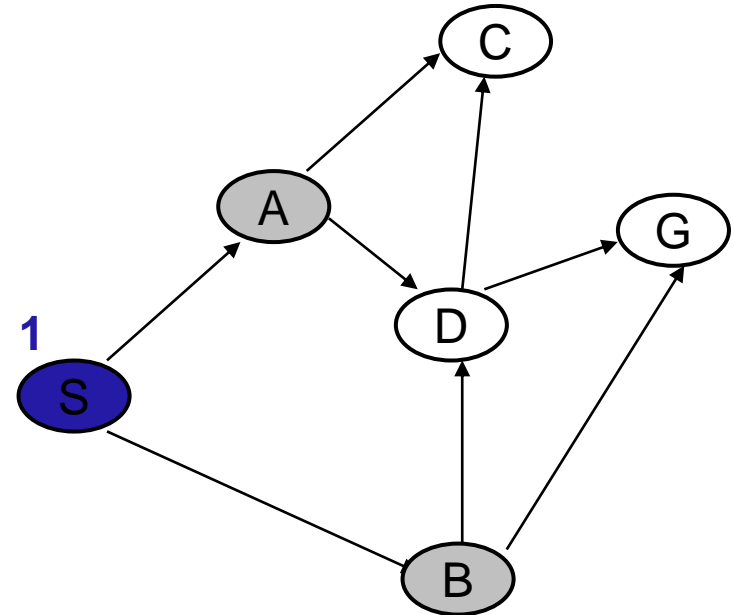
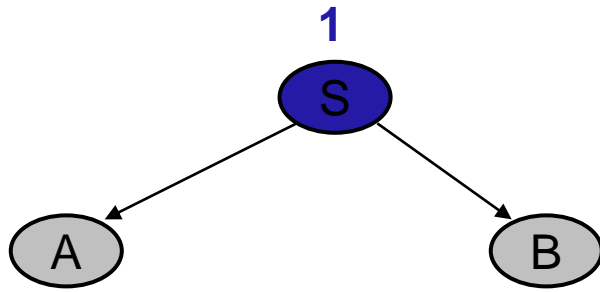
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# Depth First Search



Also we have look at this as,

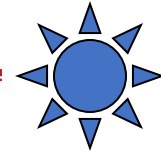


**Numbers indicate order pulled off of Q (expanded)**

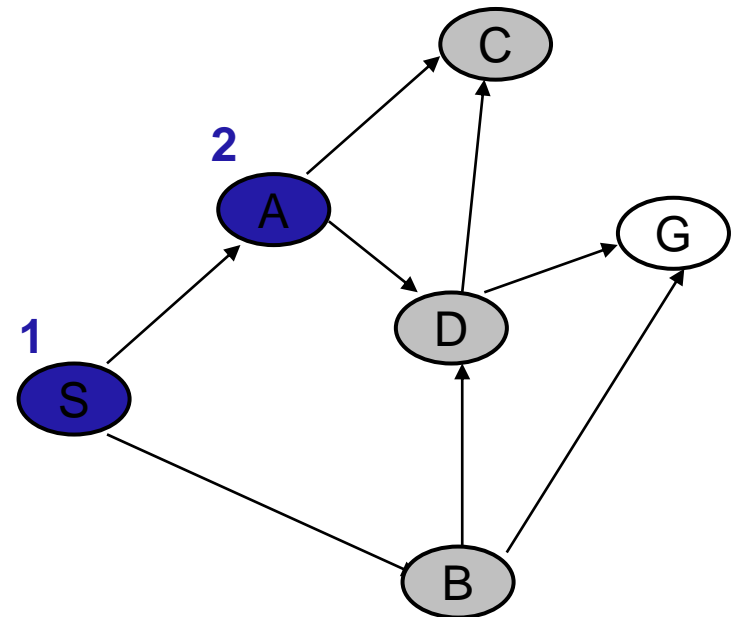
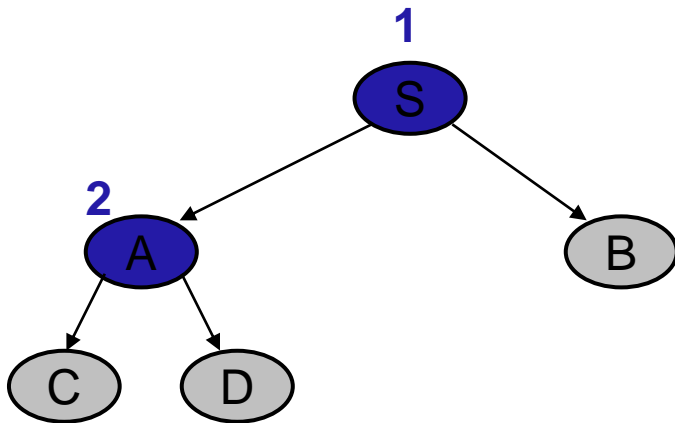
**Dark blue fill = Visited & Expanded**

**Light grey fill = Visited**

# Depth First Search



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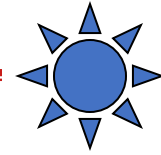


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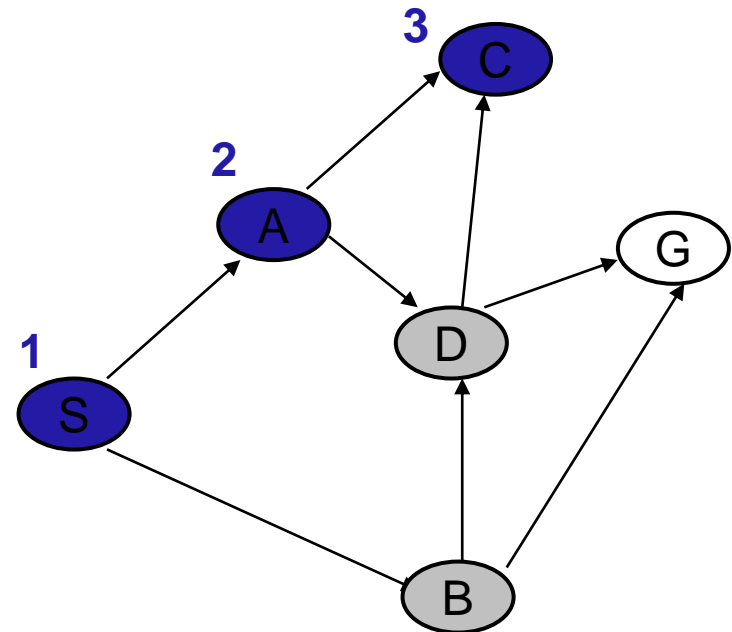
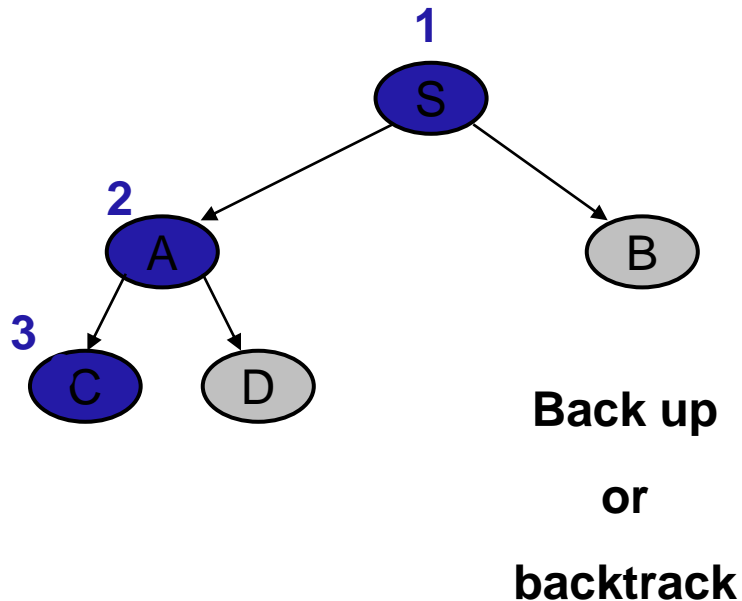
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**Light grey fill = Visited**

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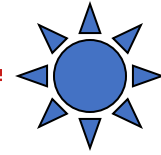
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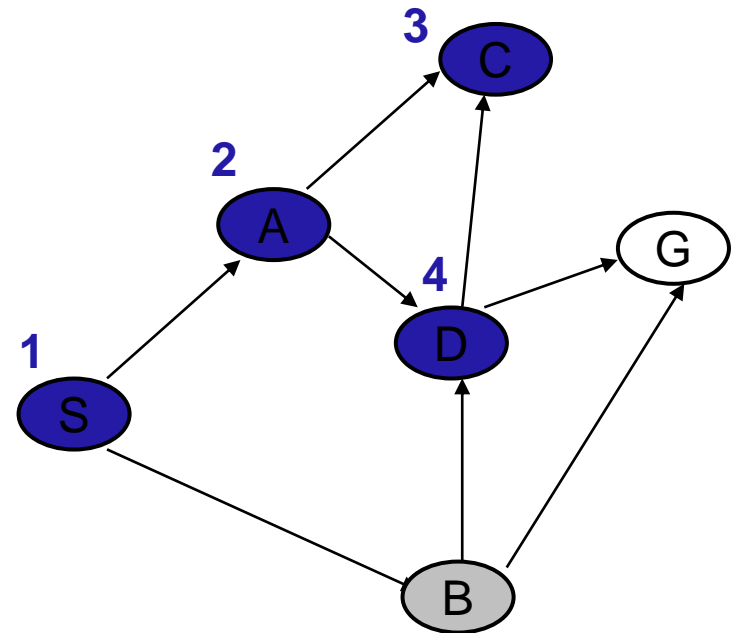
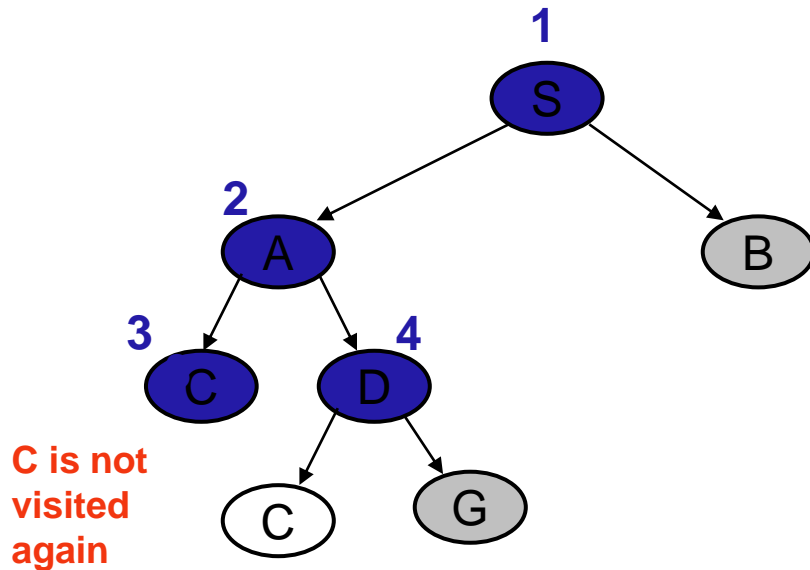
Light grey fill = Visited



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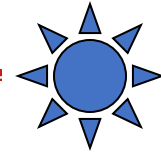


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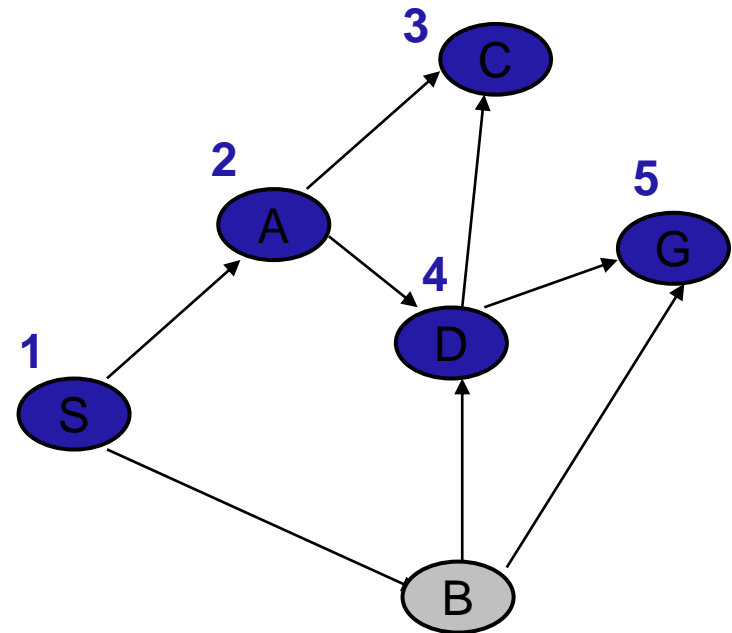
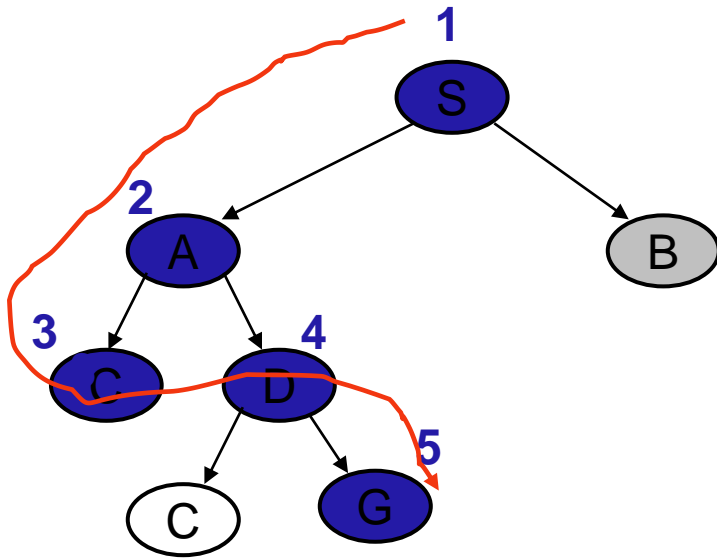
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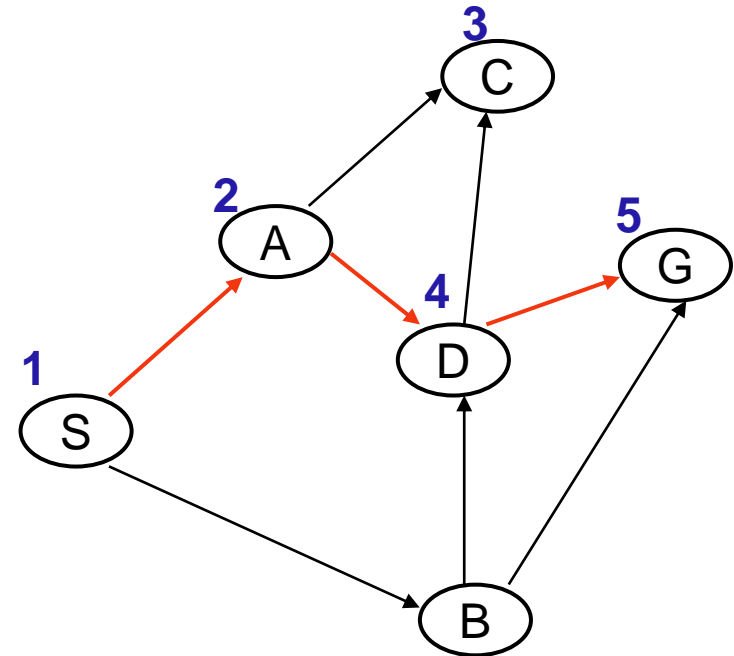
Light grey fill = Visited

# Depth First Search (without Visited List)



Pick First element of Q; Add path extensions to front of Q

	Q
1	(S)
2	(A S) (B S)
3	(C A S) (D A S) (B S)
4	(D A S) (B S)
5	(C D A S) (G D A S) (B S)
6	(G D A S) (B S)



Added paths in **blue**

Path is shown in reversed order, the node's state is the first entry

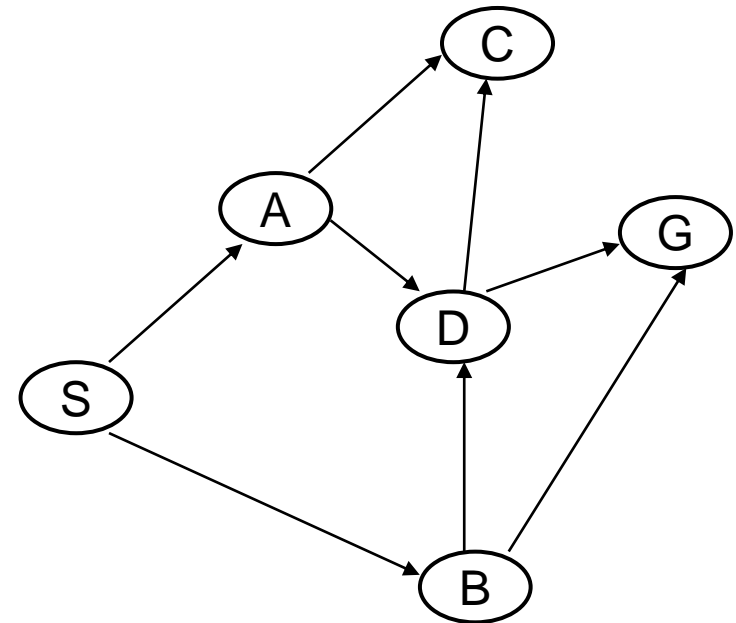
**Do not extend a path to a state if the resulting path would have a loop**

# Breadth First Search



Pick First element of Q; Add path extensions to end of Q

	Q	Visited
1	(S)	S
2		
3		
4		
5		



Added paths in blue

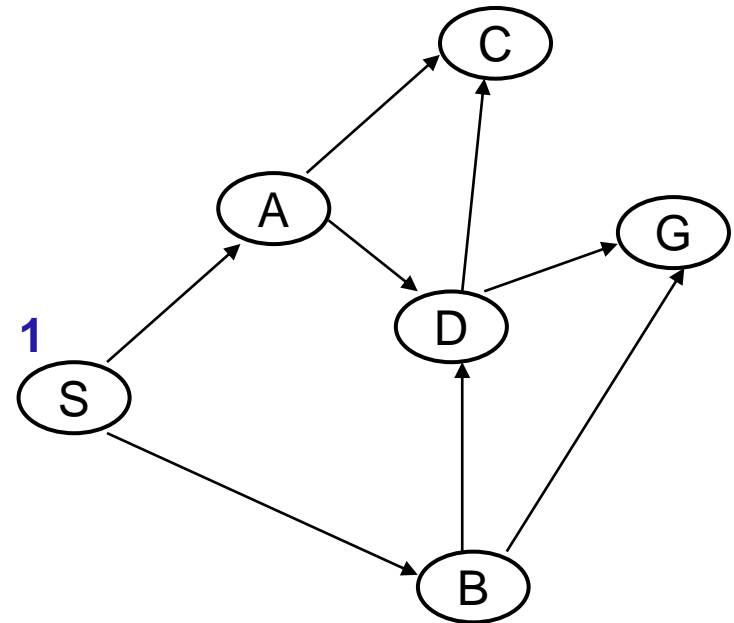
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Pick First element of Q; Add path extensions to end of Q

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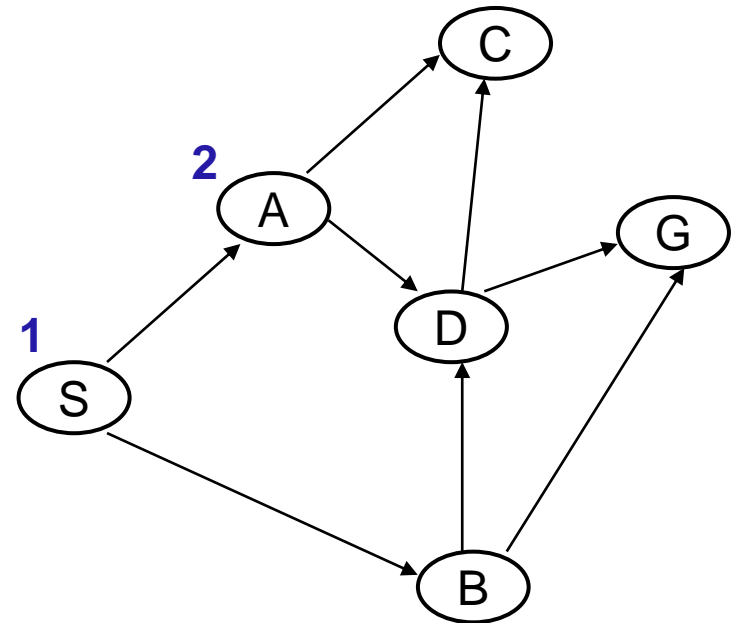
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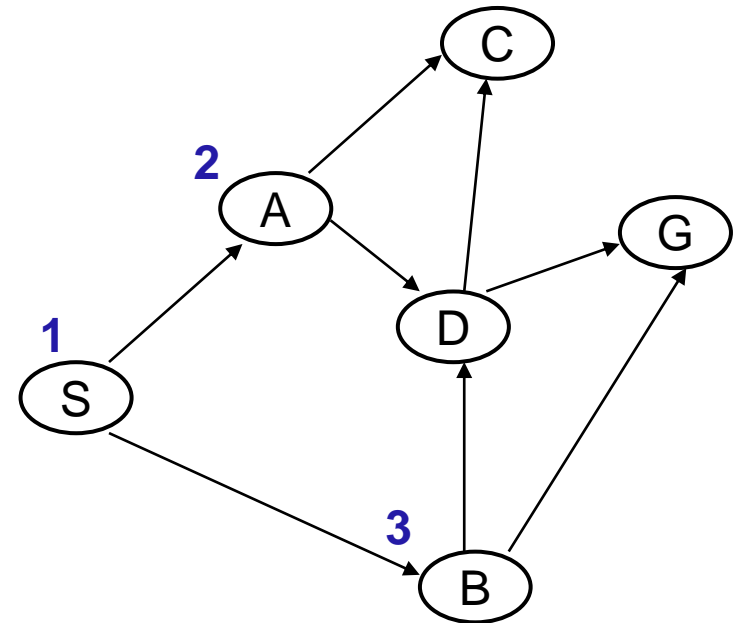
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6		



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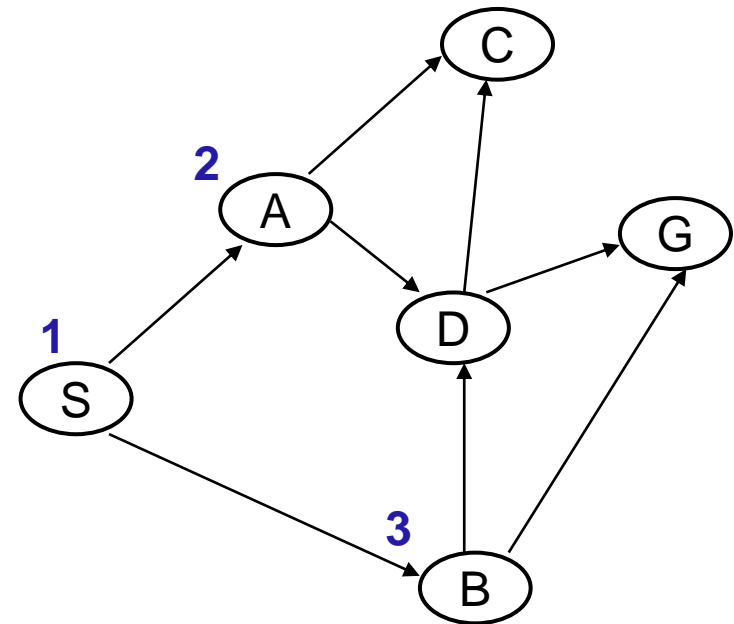
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4	(C A S) (D A S) (G B S)*	G,C,D,B,A,S
5		
6		



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**\*We could have stopped here, when the first path to the goal was generated**

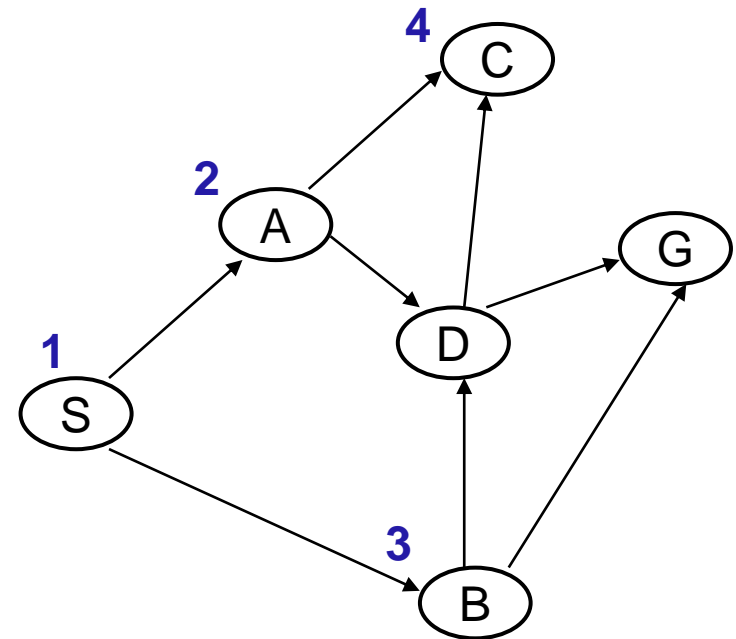


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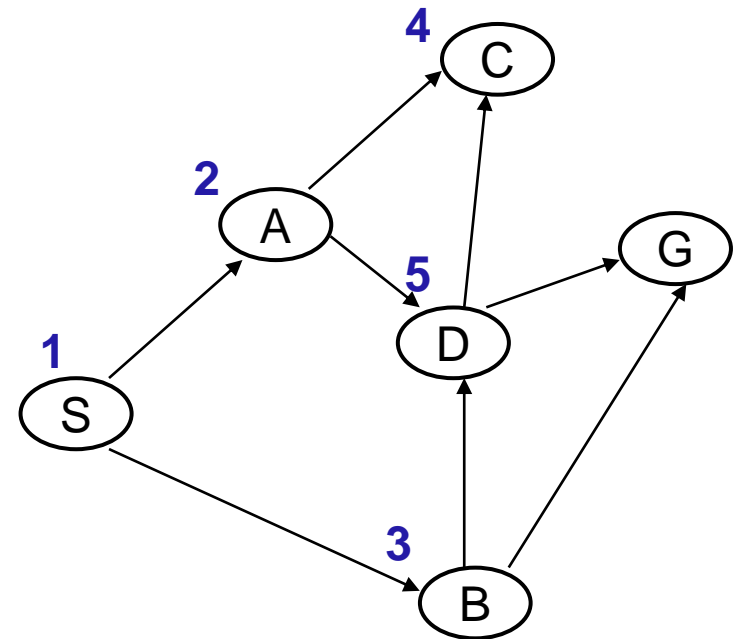
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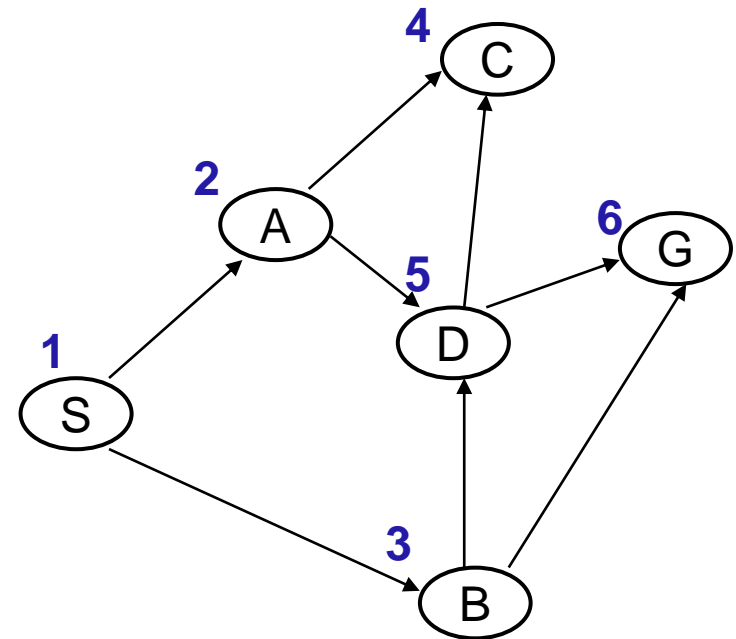
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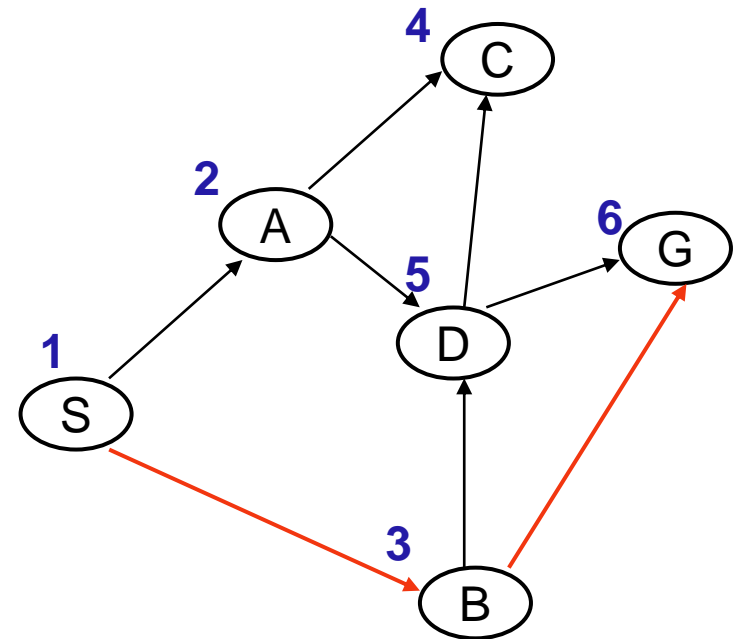
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Added paths in **blue**

Path is shown in reversed order, the node's state is the first entry

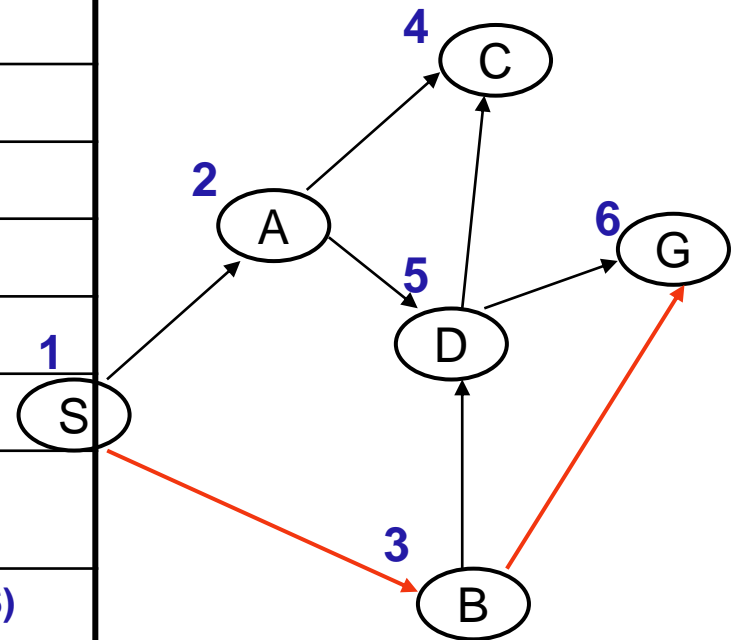
**\*We could have stopped here, when the first path to the goal was generated**

# Breadth First (without Visited list)



Pick First element of Q; Add path extensions to end of Q

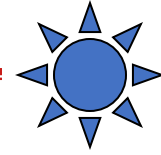
	Q
1	(S)
2	(A S) (B S)
3	(B S) (C A S) (D A S)
4	(C A S) (D A S) (D B S) (G B S)*
5	(D A S) (D B S) (G B S)
6	(D B S) (G B S) (C D A S) (G D A S)
7	(G B S) (C D A S) (G D A S) (C D B S) (G D B S)



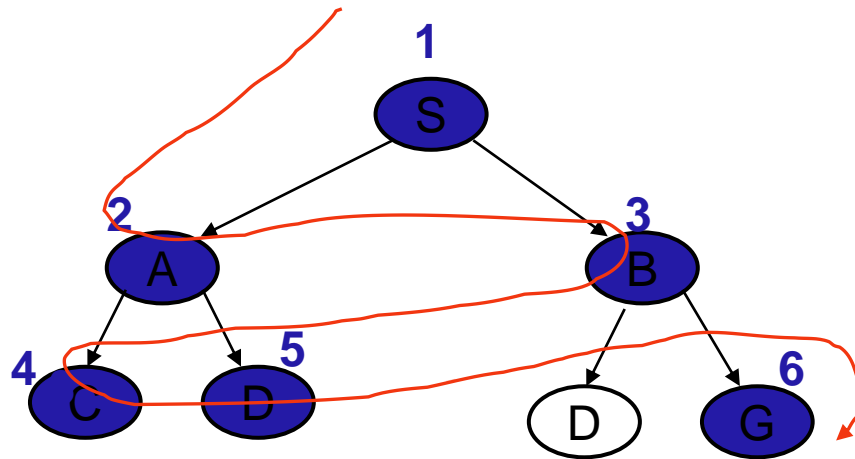
Added paths in **blue**

Path is shown in reversed order, the node's state is the first entry

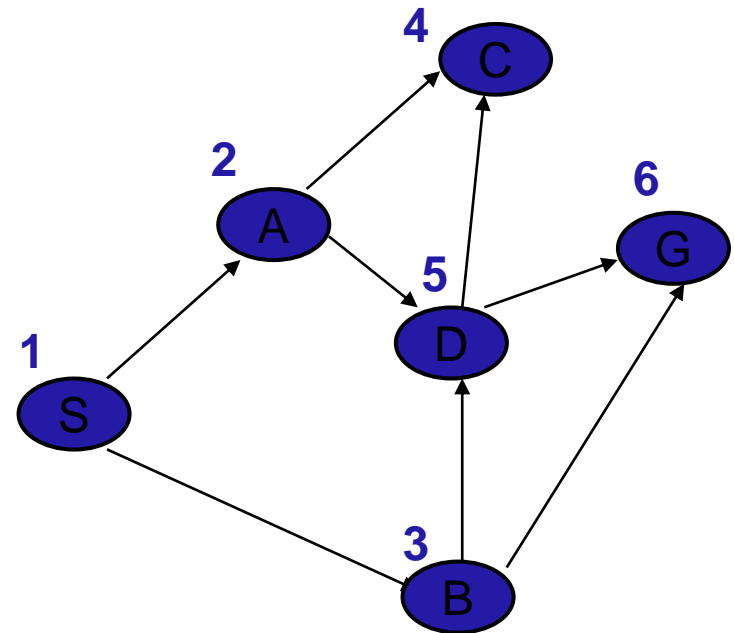
# Breadth First Search



Also we have look at this as,



**D is not visited again**



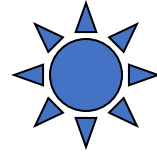
**Numbers indicate order pulled off of Q (expanded)**

**Dark blue fill = Visited & Expanded**

**Light grey fill = Visited**

## **Implementing Issues: Finding the best node**

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**There are many possible approaches to finding the best node in Q.**

**Scanning Q to find lowest values**

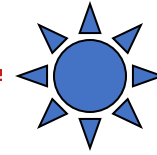
**Sorting Q and picking the first element**

**Keeping the Q sorted by doing “sorted” insertions**

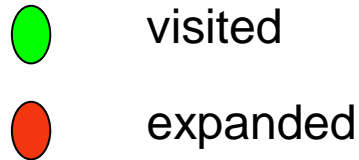
**Keeping Q as a priority queue** – like in data structures

**Which of these is best will depend among other thing on how many children nodes have on average**

# Worst Case Space

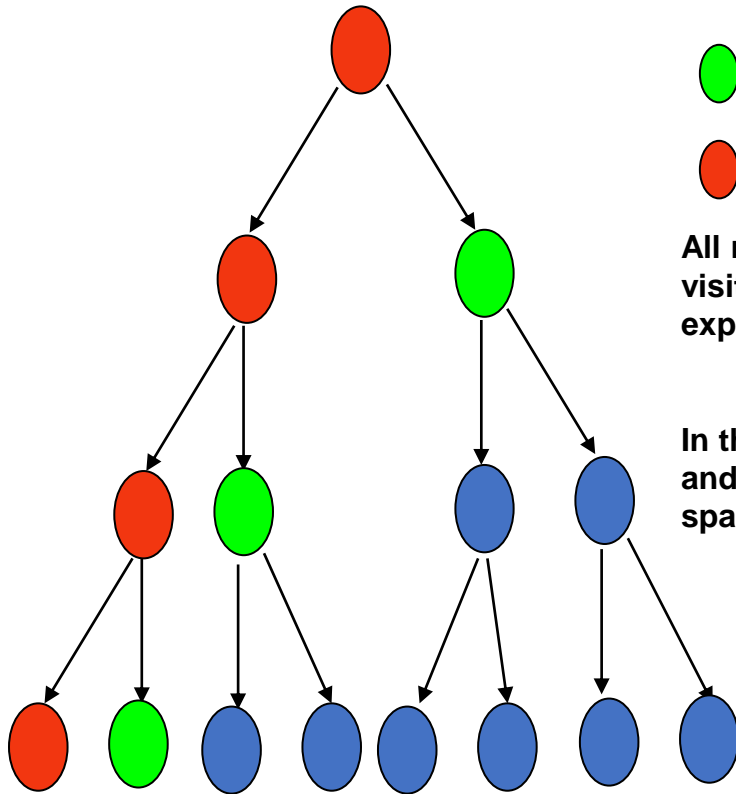


$$\text{Max Q size} = \text{Max}(\# \text{Visited} - \# \text{Expanded})$$



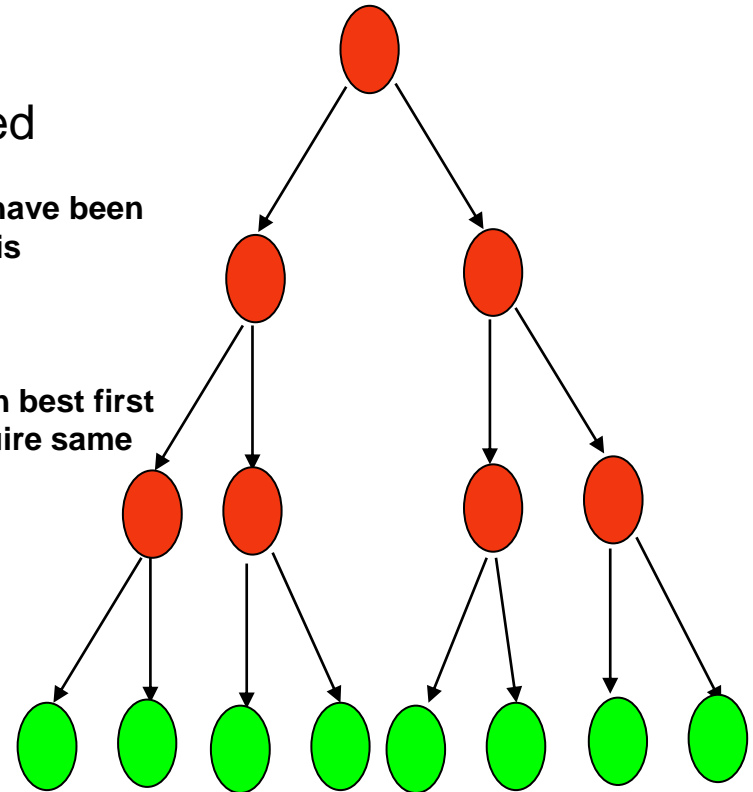
All nodes at depth  $d$  have been visited but none of it is expanded

In the worst case both best first and breadth first require same space



Depth First max Q size

$$(b-1)d \approx bd$$



Breadth First max Q size

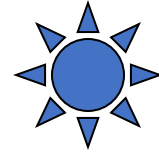
$$b^d$$



# Worst Case Space

**Max Q size = Max(#Visited - #Expanded)**

## In a tree structured search



### Q holds unexpanded siblings of the node

## Path length cannot be greater than d

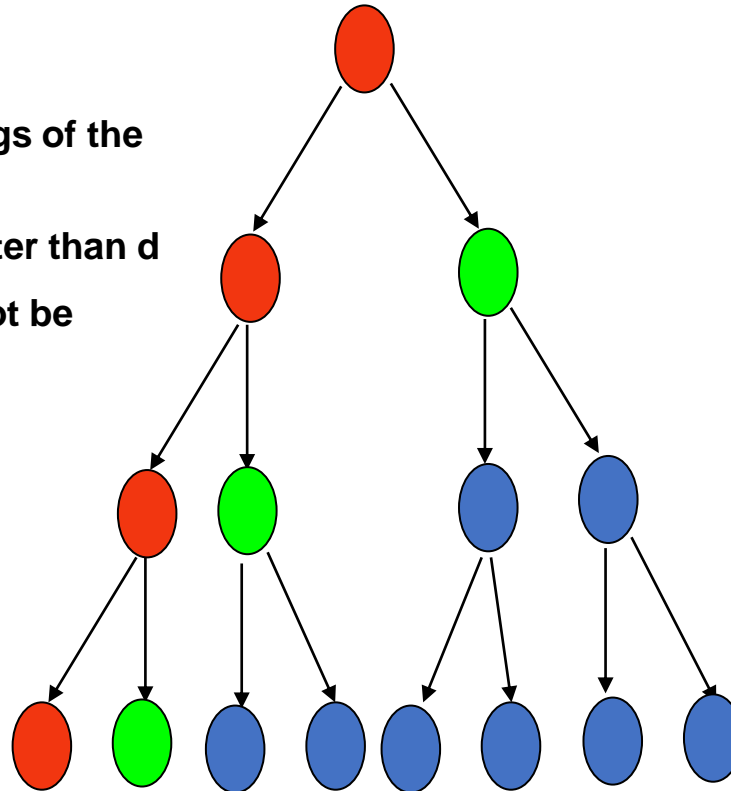
## Unexpanded siblings cannot be greater than $b-1$



visited



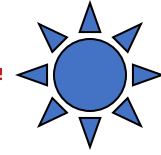
expanded



## Depth First max Q size

$$(b-1)d \approx bd$$

# Worst Case Running Time



Max time  $\propto$  Max #Visited

The number of states in the search space may be exponential in some “depth” parameter, e.g. number of actions in a plan, number of moves in a game

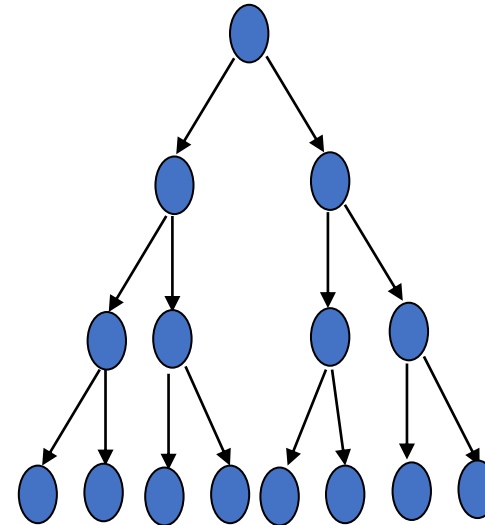
All the searches, with or without visited list, may have to visit each state at least once, in the worst case

So, all searches will have worst case running times that are at least proportional to the total number of states and therefore exponential in the “depth” parameter.

d=0

d=1

d=2



b=2

Or is Uniform

Number of choices  
in a move

d is depth

b is branching factor

$$b^d < \frac{(b^{d+1} - 1)}{(b - 1)} < b^{d+1}$$

states in tree

# Cost and Performance of Any-Path Methods

Searching a tree with branching factor  $b$  and depth  $d$   
(without using a Visited list – tree search)



Search Method	Worst Time	Worst Space	Fewest States?	Guaranteed to find Path?
Depth-first	$b^{d+1}$	$bd$	No	Yes*
Breadth-first	$b^{d+1}$	$b^d$	Yes	Yes
Best-first	$b^{d+1}^{**}$	$b^d$	No	Yes*

*\*If there are no infinitely long paths in the search space*

*\*\*Best-first needs more time to locate the best node in Q*

*Not very precise estimates but just tradeoffs*

**Worst case time is proportional to number of nodes added to Q**

**Worst case space is proportional to maximal length of Q**

**DFS requires less space i.e. linear**

# Cost and Performance of Any-Path Methods

Searching a tree with branching factor  $b$  and depth  $d$   
(using a Visited list)



Search Method	Worst Time	Worst Space	Fewest States?	Guaranteed to find Path?
Depth-first	$b^{d+1}$	<del><math>b^d</math></del> $b^{d+1}$	No	Yes*
Breadth-first	$b^{d+1}$	<del><math>b^d</math></del> $b^{d+1}$	Yes	Yes
Best-first	$b^{d+1}^{**}$	<del><math>b^d</math></del> $b^{d+1}$	No	Yes*

*\*If there are no infinitely long paths in the search space*

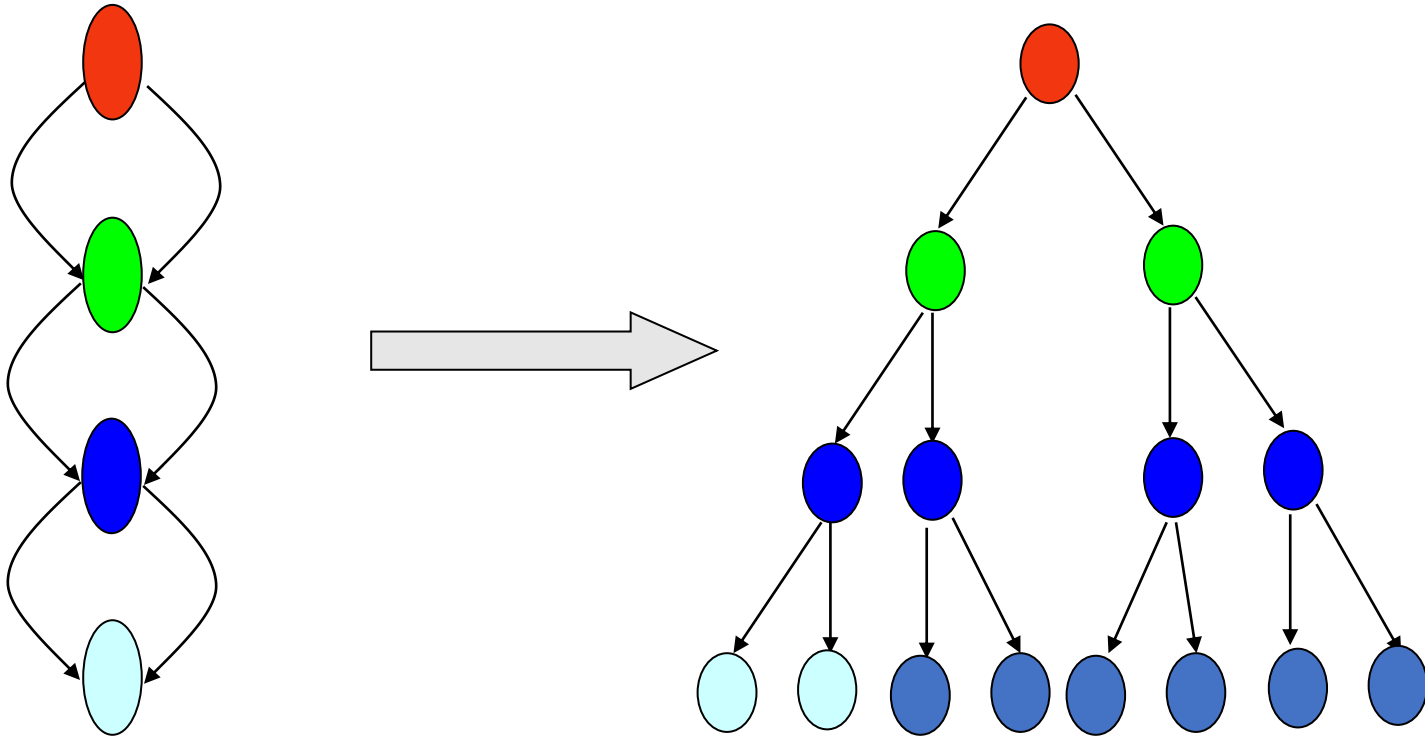
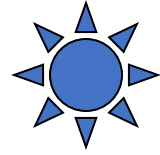
*\*\*Best-first needs more time to locate the best node in Q*

**Worst case time is proportional to number of nodes added to Q**

**Worst case space is proportional to maximal length of Q (and Visited list)**

# States Vs Paths

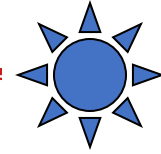
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Visited list can save lot of time as shown by this example.....there could be exponentially larger number of path than states

# Space

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In large search problem, memory is often the limiting factor

Imagine search a tree with branching factor 8 and depth 10. Assume a node requires just 8 bytes of storage. Then breadth-first search might require up to

$$(2^3)^{10} \times 2^3 = 2^{33} \text{ bytes} = 8,000 \text{ Mbytes} = 8 \text{ Gbytes}$$

One strategy is to trade time for memory. For example, we can emulate breadth-first search by repeated applications of depth-first, each up to a preset depth limit. This is called progressive deepening search (PDS)

- 1-  $C=1$
- 2- Do DFS to max depth  $C$ . If path found, return it.
- 3- Otherwise, increment  $C$  and go to 2

# **PDS – Best of Both Worlds**

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**Depth-First Search (DFS) has small space requirements (linear in depth) but has major problems:**

- **DFS can run forever in search spaces with infinite length paths**
- **DFS does not guarantee finding shallowest goal**

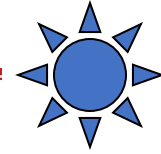
**Breadth first Search (BSF) guarantees finding shallowest goal, even in the presence of infinite paths, but it has one great problem:**

- **BFS requires a great deal of space (exponential in depth)**

**Progressive Deepening Search (PDS) has the advantages of DFS and BFS**

- **PDS has small space requirements (linear in depth)**
- **PDS guarantees finding shallowest goal**

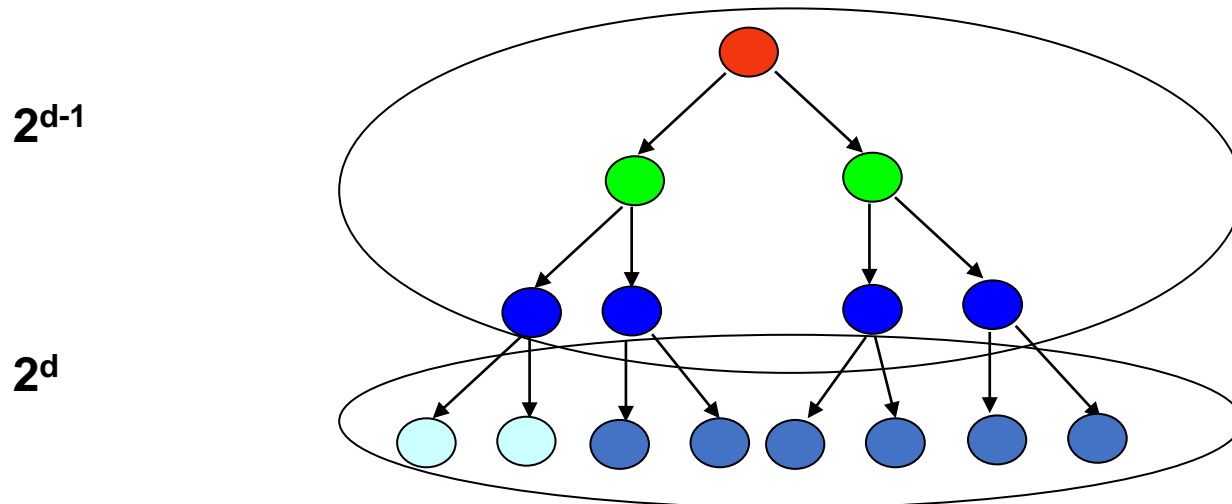
# Progressive Deepening Search



**Isn't PDS too expensive?** You look at nodes again and again so wastage of time but.....only in small searches or graphs

**In exponential trees, time is dominated by deepest search**

**For example, if branching factor is 2, then the number of nodes at depth  $d$  is  $2^d$  while the total number of nodes in all previous levels is  $2^d - 1$ , so the difference between looking at whole tree versus only the deepest level is at worst a factor of 2 in performance**





# Progressive Deepening Search

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Compare the ratio of average time spent on PDS with average time spent on a single DFS with the full depth tree:

$$(\text{Avg time for PDS})/(\text{Avg time for DFS}) \approx (b+1)/(b-1)$$

b	ratio
2	3
3	2
5	1.5
25	1.08
100	1.02

**Progressive deepening is an effective strategy for difficult searches.**

# Questions

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