

# Lecture 16

Brute Force Algorithms & their Analysis: Finding Inversion in an Array, and Closest-Pair & Convex-Hull Problems.

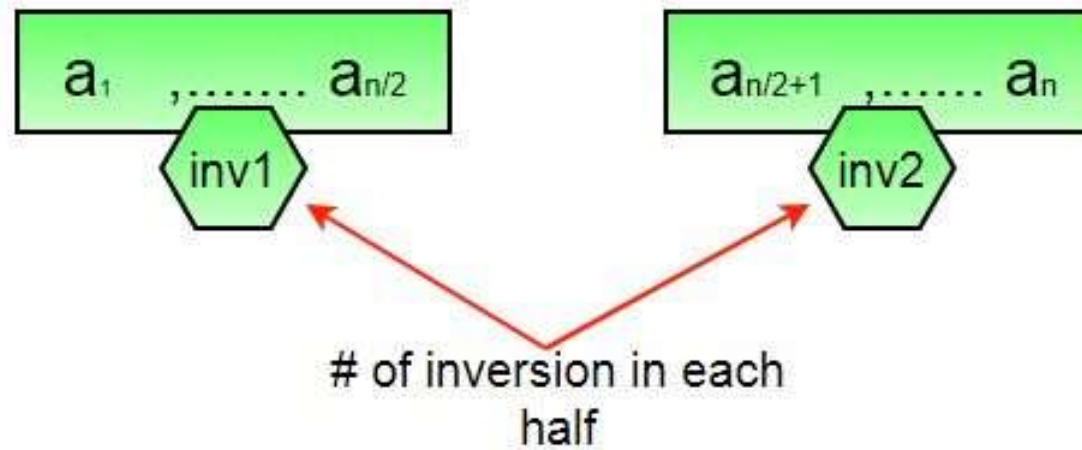




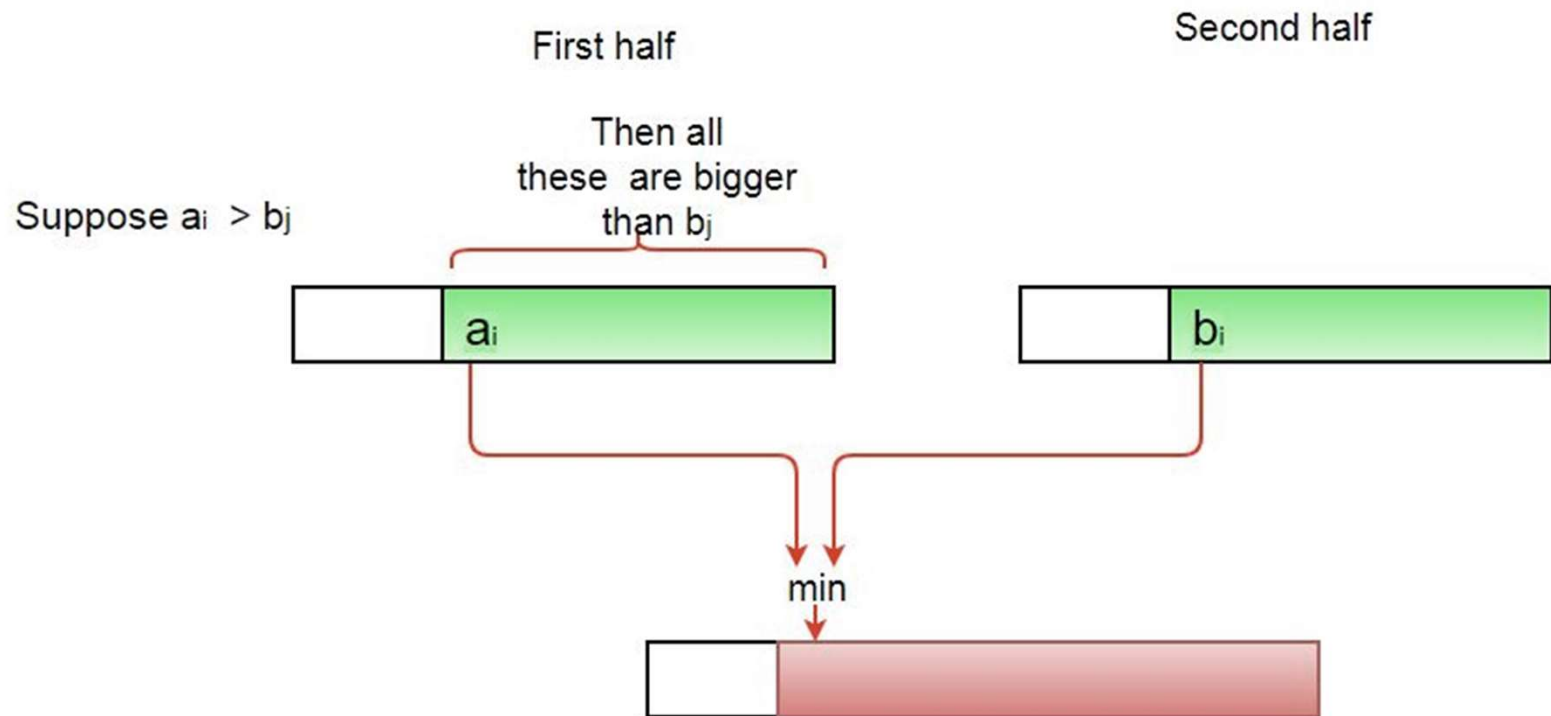
## Inversion Count in Array

- › **Inversion Count** for an array indicates – how far (or close) the array is from being sorted.
  - If the array is already sorted, then the inversion count is 0,
  - if the array is sorted in reverse order, the inversion count is the maximum.
- › Given an array **a[]**. The task is to find the inversion count of **a[]**. Where two elements  $a[i]$  and  $a[j]$  form an inversion if  $a[i] > a[j]$  and  $i < j$

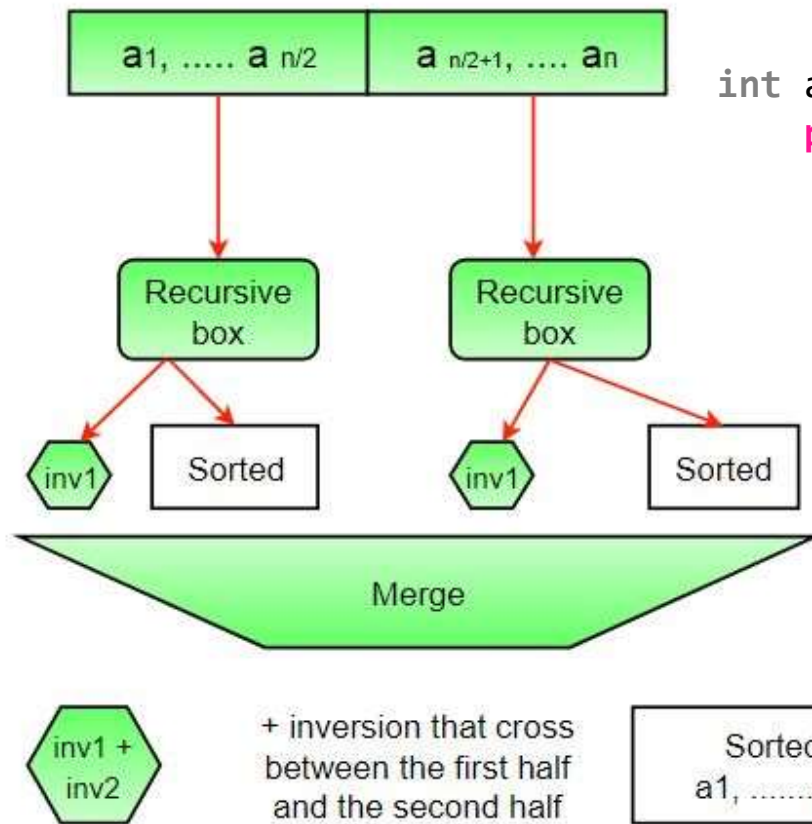
# Concept of Inversion



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```
int arr[] = { 1, 20, 6, 4, 5 };
printf(" Number of inversions are %d \n",
mergeSort(arr, 5));
```

## Time Complexity:

$O(n \log n)$ , The algorithm used is divide and conquer i.e. merge sort whose complexity is  $O(n \log n)$ .

## Auxiliary Space:

$O(n)$ , Temporary array.



# Inversions in Array: Steps {Self Execution}

**GetInvCount**(ar, n):

```
    inv_count = 0
    for i in range(n):
        for j in range(i + 1, n):
            if (ar[i] > ar[j]):
                inv_count += 1
    return inv_count
```

```
array = [1, 20, 6, 4, 5]
n = len(array)
print("Number of inversions are",
      GetInvCount(array, n))
```

**OUTPUT:**

Number of inversions are **5**

- **Time Complexity:**  $O(n^2)$ , Two nested loops are needed to traverse the array from start to end.
- **Auxiliary Space:**  $O(1)$ , No extra space is required.



## Inversion Count in Array: Examples

› **Input:**  $arr[] = \{8, 4, 2, 1\}$

**Output:** 6

**Explanation:** Given array has six inversions:

– (8, 4), (4, 2), (8, 2), (8, 1), (4, 1), (2, 1).

› **Input:**  $arr[] = \{1, 20, 6, 4, 5\}$

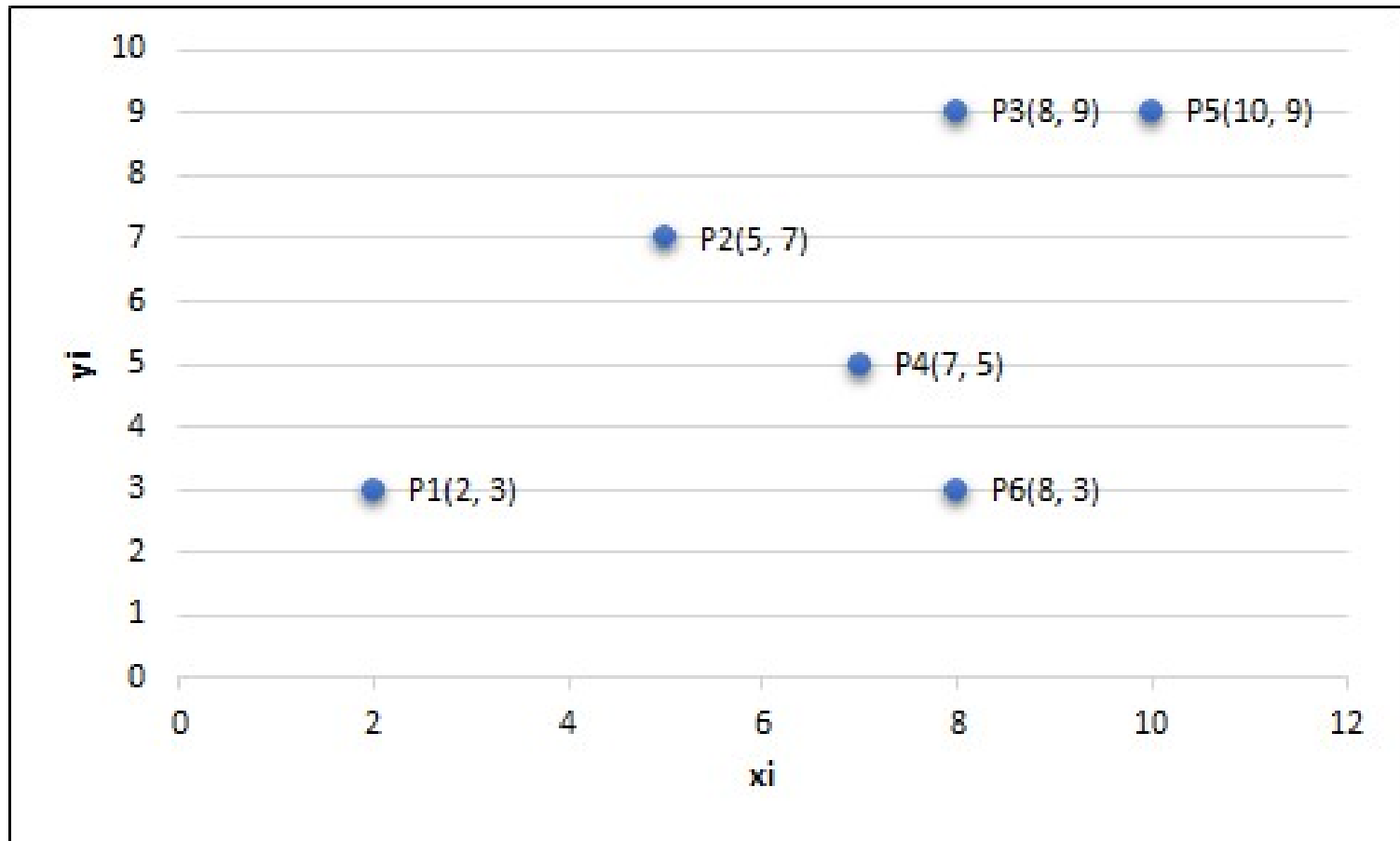
**Output:** 5

**Explanation:** Given array has five inversions:

– (20, 6), (20, 4), (20, 5), (6, 4), (6, 5).



## Brute Force: Closest Pair







# Brute Force: Closest Pair Problem

› Find distance between points

P1P2	P2P3	P3P4	P4P5	P5P6
P1P3	P2P4	P3P5	P4P6	
P1P4	P2P5	P3P6		
P1P5	P2P6			
P1P6				

› Formula to use: This is the basic distance formula

–  $d(p_i, p_j) = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$

– The least distance shows the closest pair



# Brute Force: Closest Pair Algorithm

## › *Algorithm BF – ClosestPair(p)*

›  $d_{min} = \infty$

› *FOR*  $i = 1$  *to*  $n - 1$

– *FOR*  $j = i + 1$  *to*  $n$

›  $d = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$

› *IF*  $d < d_{min}$  *THEN*

–  $d_{min} = d$

–  $index1 = i$

–  $index2 = j$

› *RETURN*  $index1, index2$

## › *Time Complexity*

–  $C(n) = \sum_{i=1}^{n-1} \sum_{j=i+1}^n 2$

–  $= 2 \sum_{i=1}^{n-1} (n - i - 1 + 1)$

–  $= 2 \sum_{i=1}^{n-1} (n - i)$

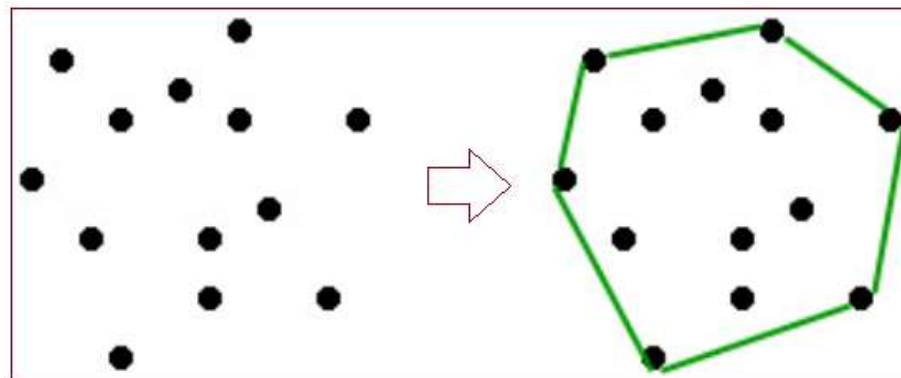
–  $= 2(n - 1) + (n - 2) \dots + 1$

–  $= 2 \cdot \frac{n(n-1)}{2} = n(n - 1) = n^2 - n$

–  $\Theta(n^2) \cong C(n)$

# Convex-Hull Problem : Divide & Conquer

- › Convex-Hull is the smallest convex polygon that contain all the points in the plane.
  - Given a set of  $X$  of point  $P_1(x_1, y_1), P_2(x_2, y_2) \dots, P_n(x_n, y_n)$  in plane. We want to find convex hull.
  - Divide and Conquer algorithm take  $O(n \log n)$  time to compute convex hull in clockwise order.
  - All the points are inside the polygon after connecting all outer points and called as Convex -Hull.
  - Input points must be pre-sorted by x-coordinates





## Divide & Conquer: Convex Hull Steps

1. Partition  $X$  into half  $X_1$  and  $X_2$  according to  $X$  coordinates.
2. If  $X_1$  is Empty, THEN Upper Half is simply line with end points  $P_1$  and  $P_2$
3. If  $X$  is not Empty, THEN Find  $P_{\max}$  in  $X_1$  which is farthest from linear  $P_1P_n$
4. If the TIE in  $P_{\max}$  THEN point the maximum angle  $P_{\max}, P_1, P_2$  can be
5. Now, Algorithm identifies all point if  $X_1$  that are left of line  $P_1P_{\max}$  – goto Step2

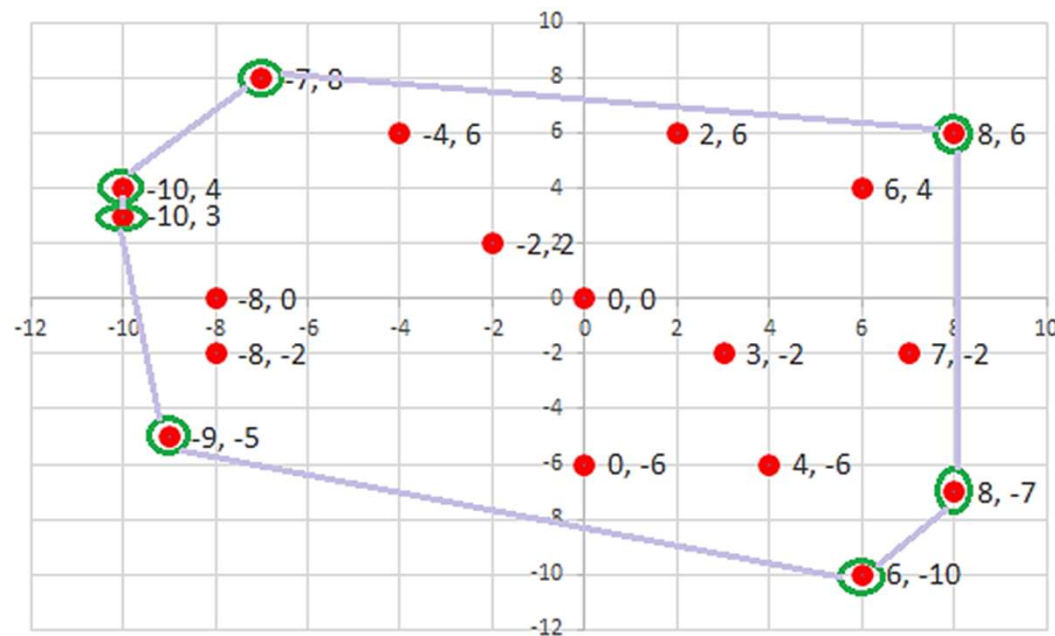


### **Divide and conquer : *Convex Hulls***

1. Divide the  $n$  points into two halves.
2. Find convex hull of each subset.
3. Combine the two hulls into overall convex hull.

# Convex-Hull Example

- › **Input** – Set of points:  $\{(-7,8), (-4,6), (2,6), (6,4), (8,6), (7,-2), (4,-6), (8,-7), (0,0), (3,-2), (6,-10), (0,-6), (-9,-5), (-8,-2), (-8,0), (-10,3), (-2,2), (-10,4)\}$
- › **Output** – Boundary points of convex hull are –
  - $(-9, -5) (6, -10) (8, -7) (8, 6) (-7, 8) (-10, 4) (-10, 3)$



[https://www.tutorialspoint.com/online\\_cpp\\_compiler.php](https://www.tutorialspoint.com/online_cpp_compiler.php)

# Thank You!!!

Have a good day

