

BRICS – Efficient Techniques for Estimating the Farness-Centrality in Parallel

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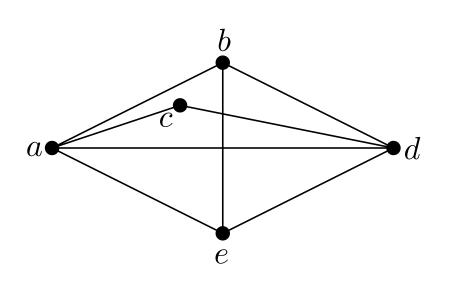
Why Graphs?



For a node v in graph G(u,v), it is the sum of the shortest distances from node v to all other nodes in the graph.

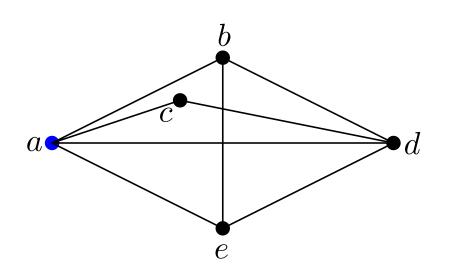
$$farness[u] = \sum_{v \in V} d(u, v)$$





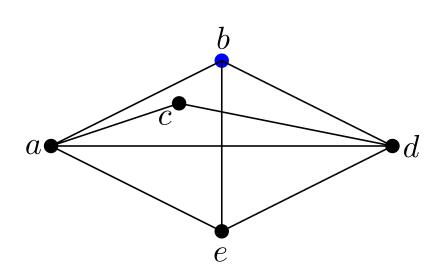
	a	b	c	d	e	FC
a						
b						
c						
$\mid d \mid$						
$\mid e \mid$						





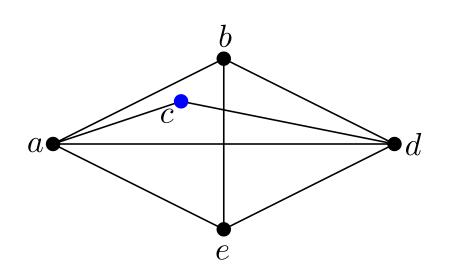
	a	b	c	d	e	FC
a	0	1	1	1	1	4
b						
c						
d						
$\mid e \mid$						





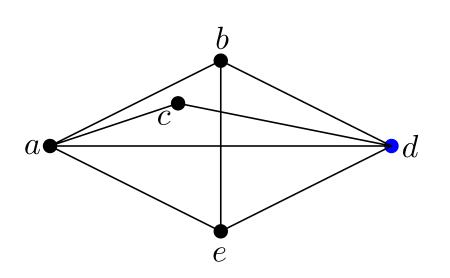
	a	b	c	d	e	FC
a	0	1	1	1	1	4
b	1	0	2	1	1	5
c						
d						
e						





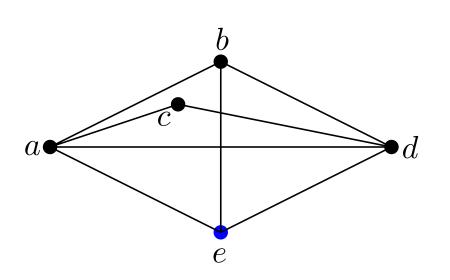
	a	b	c	d	e	FC
a	0	1	1	1	1	4
b	1	0	2	1	1	5
c	1	2	0	1	2	6
d						
$\mid e \mid$						





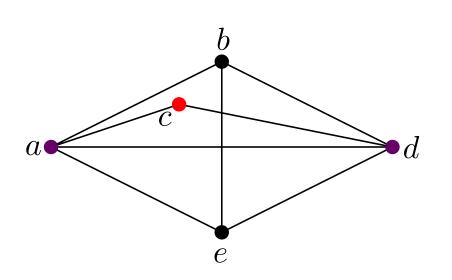
	a	b	c	d	e	FC
a	0	1	1	1	1	4
b	1	0	2	1	1	5
c	1	2	0	1	2	6
$\mid d \mid$	1	1	1	0	1	4
e						





	a	b	c	d	e	FC
a	0	1	1	1	1	4
b	1	0	2	1	1	5
c	1	2	0	1	2	6
d	1	1	1	0	1	4
e	1	1	2	1	0	5





	a	b	c	d	e	FC
a	0	1	1	1	1	4
b	1	0	2	1	1	5
c	1	2	0	1	2	6
d	1	1	1	0	1	4
e	1	1	2	1	0	5

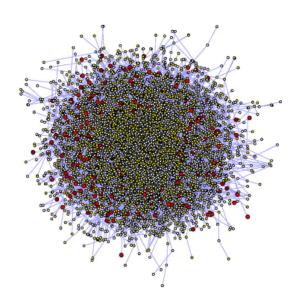
Why Approximate FC?

• Time complexity:

- FC: O(nm)

- Approx FC: O(km)

10M nodes graph



- FC: 2 - 3 days

- Approx FC: 3 - 5 Hrs

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Related Work

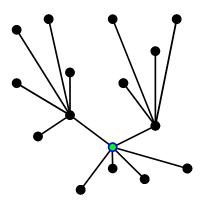
• Eppstein et al: Random-K, first randomized Approximation of Centrality.(1999)

 Cohen et al: Hybrid method of sampling and pivoting to estimate the closeness centrality.

 Sariyuce et al: BADIOS framework for computing exact betweenness centrality.



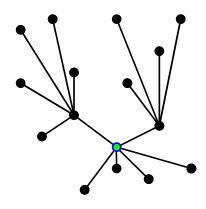
Approaches

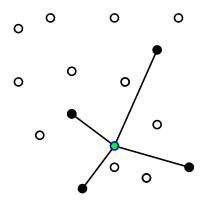


Exact FC



Approaches

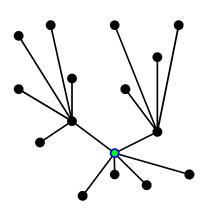


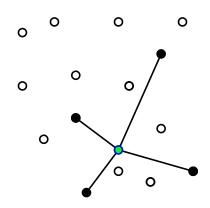


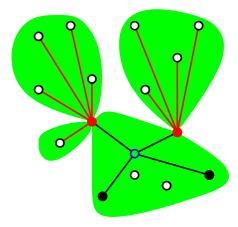
Exact FC Approx FC(Random-K)



Approaches







Exact FC

Approx FC(Random-K)

Our Approach(BRICS)



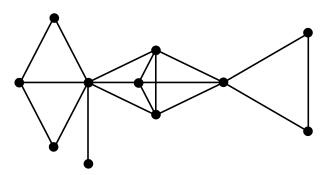
BRICS-FrameWork

- BCT
- Redundant Nodes 3D & 4D
- Identical Nodes
- Chaining



BCC based approach

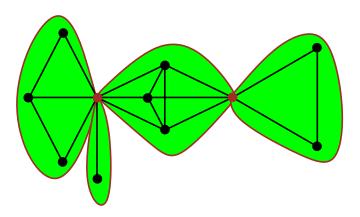
 $\mathsf{Graph}\ G$





BCC based approach

BCC(G)

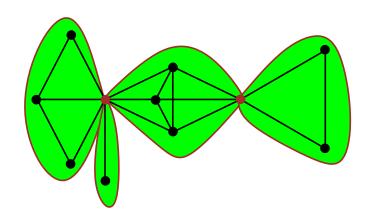


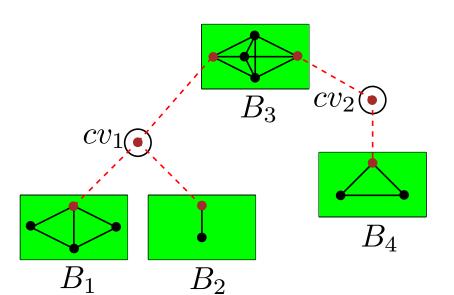


BCC based approach

BCC(G)

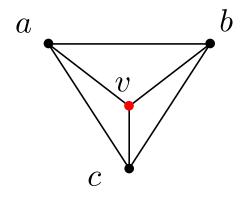
BCT(G)

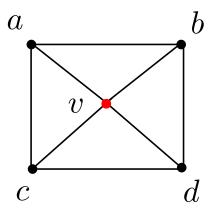






Redundant Nodes

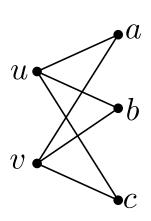


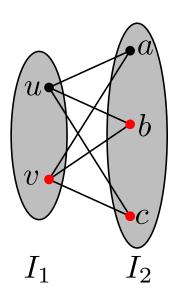


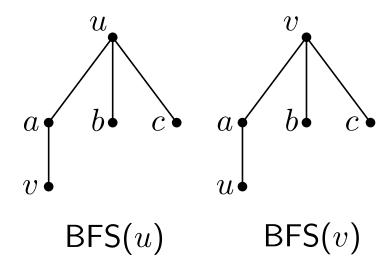
No shortest path pass through v, unless v is source/destination in R3, similarly in R4.



Identical Nodes









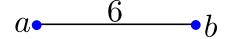
Chain Nodes





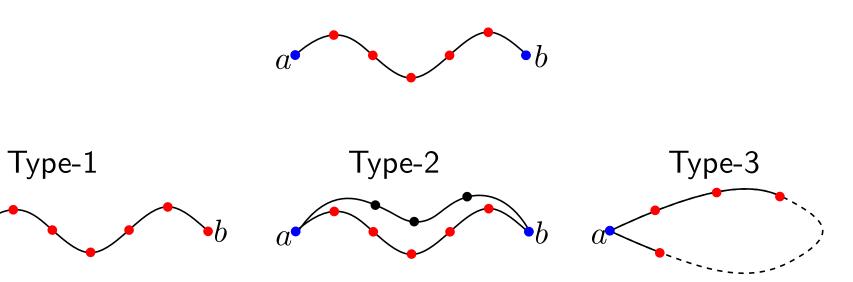
Chain Nodes





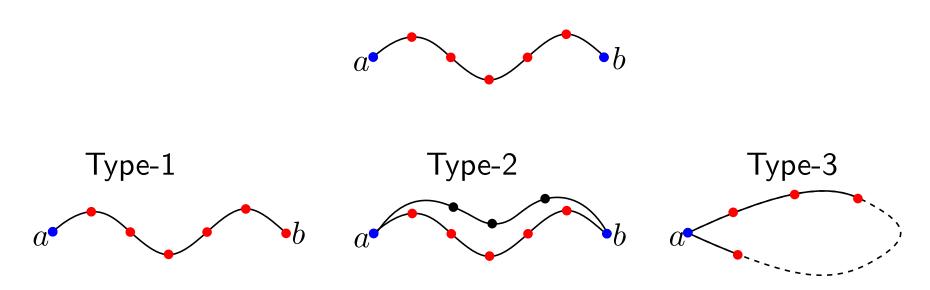


Chain Nodes





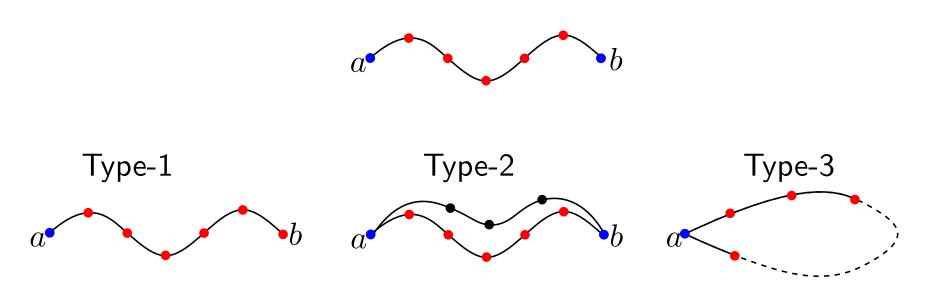
Chain Nodes



All the redundant nodes can be removed directly from Graph G.



Chain Nodes

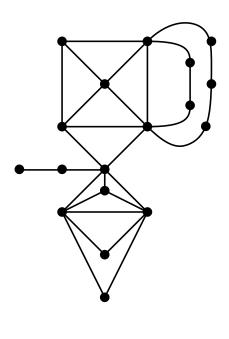


All the redundant nodes can be removed directly from Graph G.

Upto 60% of total chains are of Type-2 and 80% are redundant nodes.

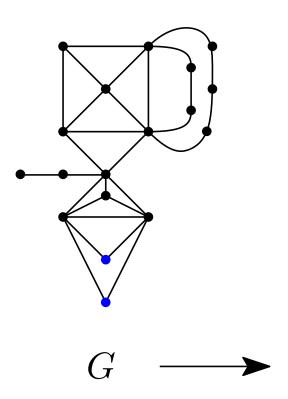


Procedure

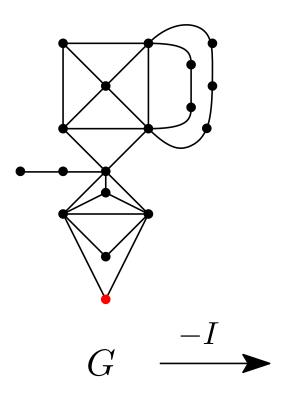


G

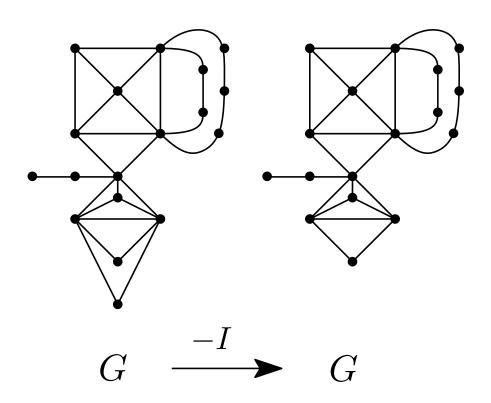




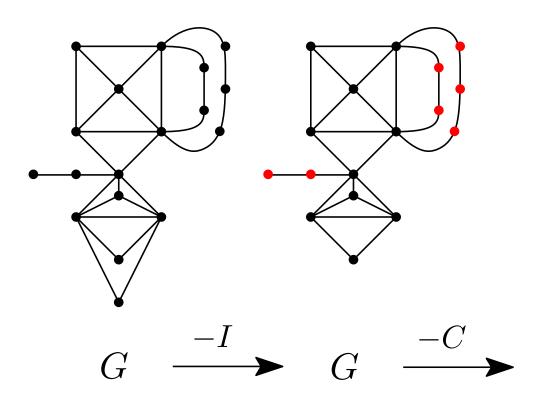




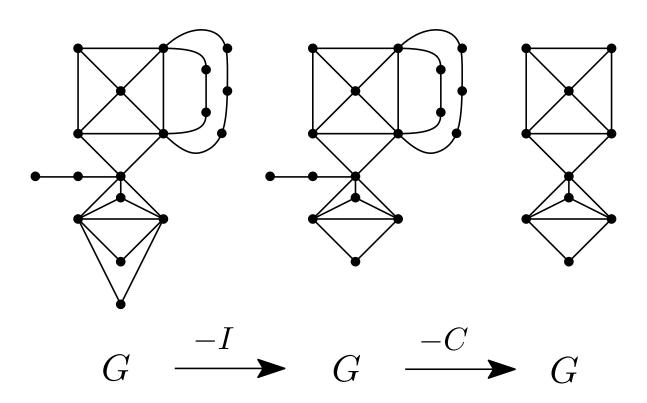




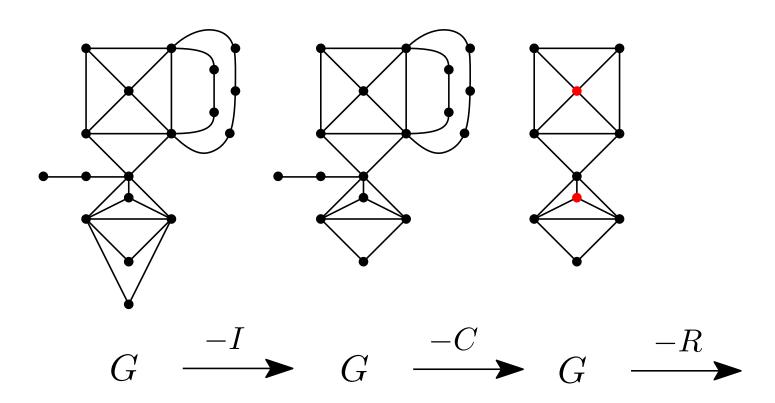




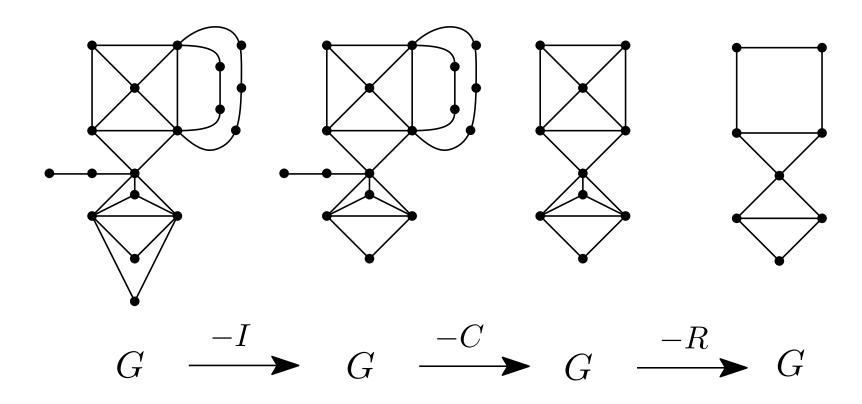




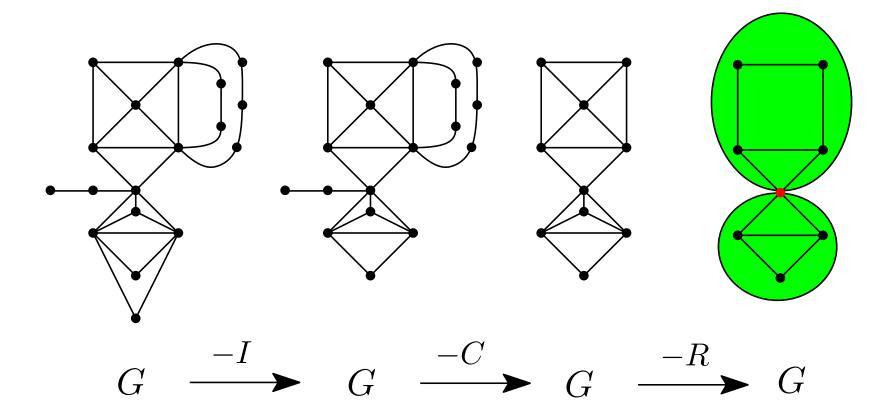




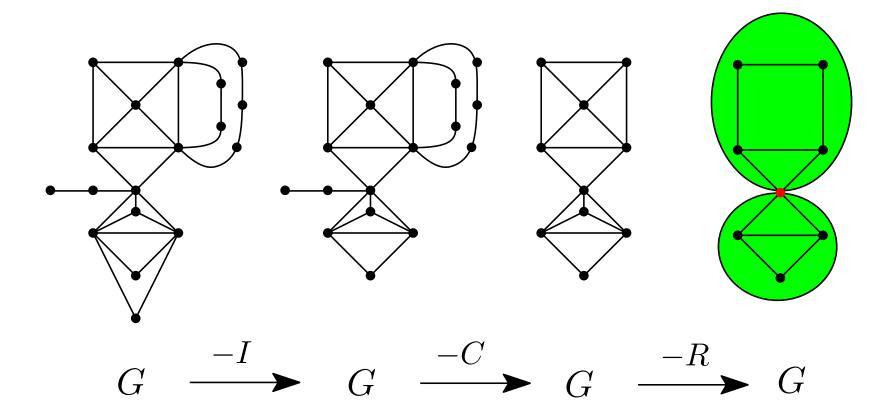




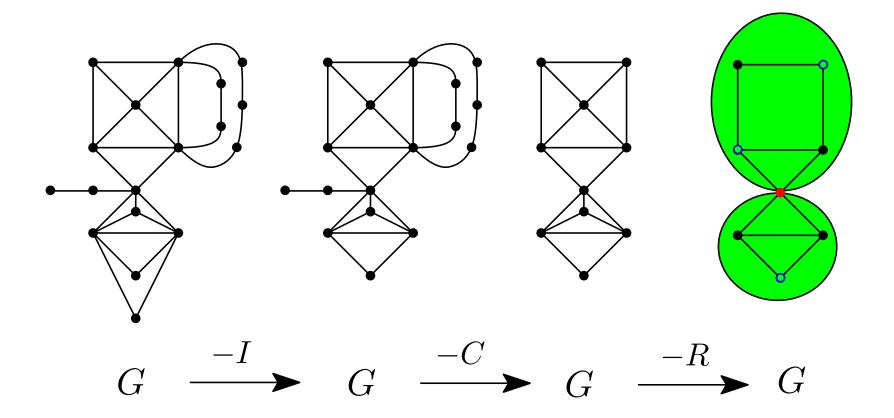






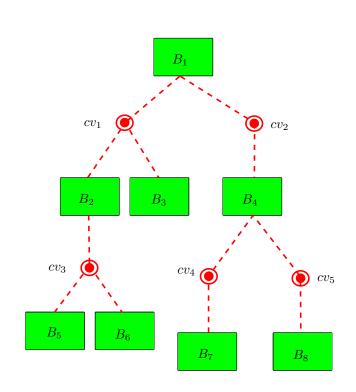






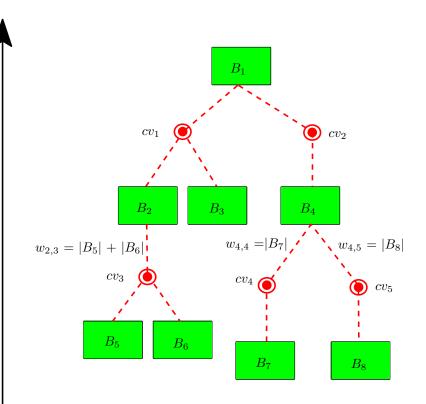




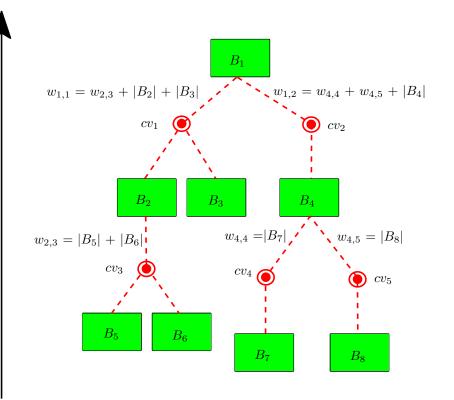




Top Down Approach



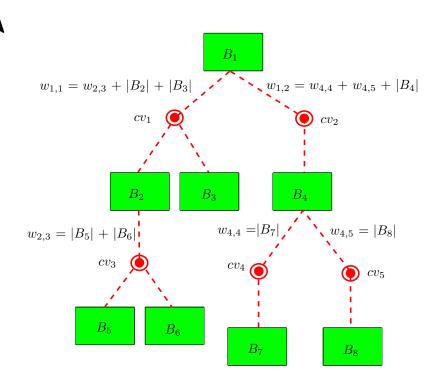


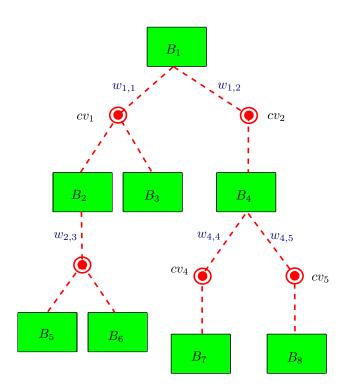






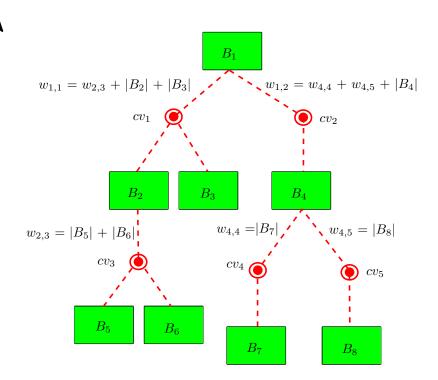
Procedure

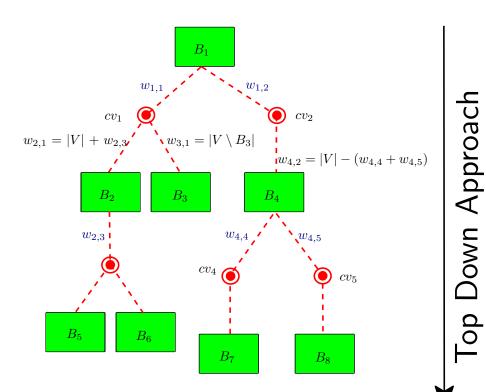




Top Down Approach

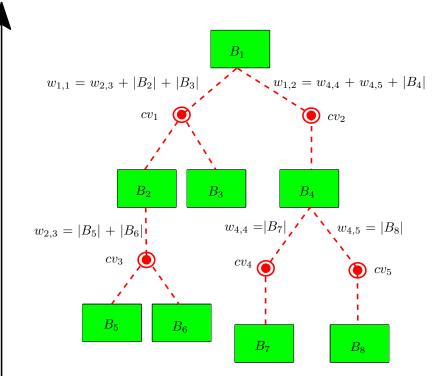


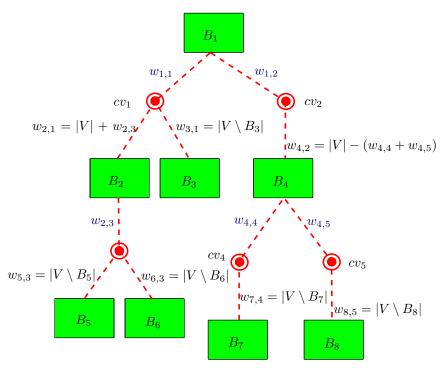






Procedure



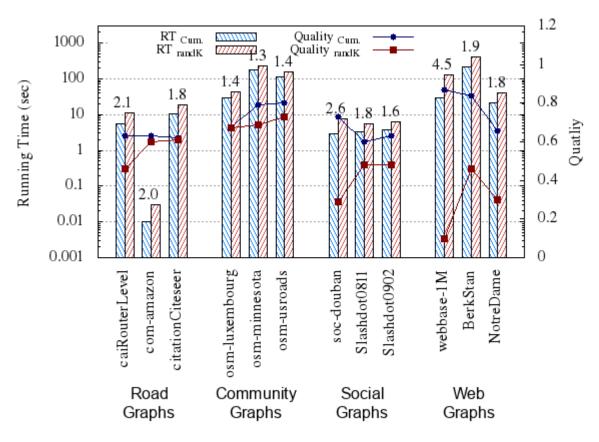


Top Down Approach



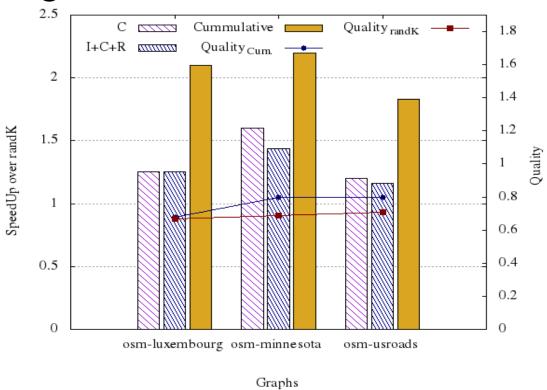
Speed up Vs Quality

Quality =
$$\frac{\sum_{v \in V} AR(v)}{n}$$
 where $AR(v) = \frac{farness_estimated(v)}{farness_actual(v)}$





Profiling for Road Networks

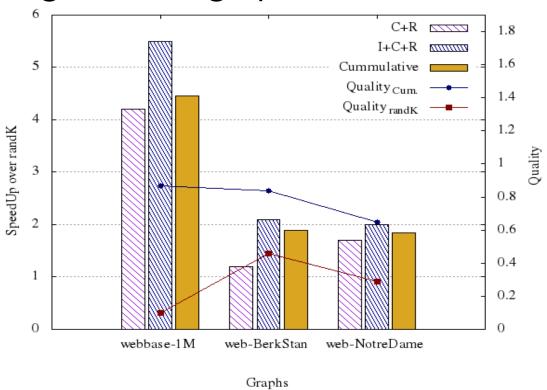


I & R < 1% C: 70 - 85% and Largest BiCC 90%

Distribution of BiCC: Skewed, Slightly better quality



Profiling for Web graphs

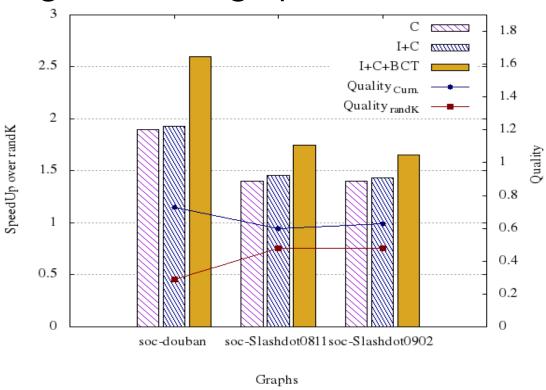


I: 44 %, C: 54%, R: 2.4 % and Large BiCC 27.3 % Nodes

Distribution of BiCC: nearly balanced, better quality



Profiling for Social graphs

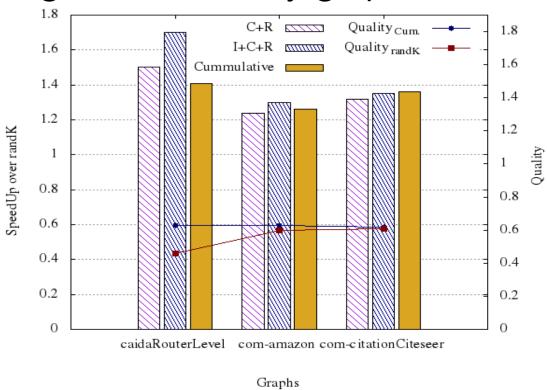


I & C \leq 40 % but R < 1% and Large BiCC: 72 %

Distribution of BiCC: Skewed, Slightly better quality



Profiling for Community graphs



I, C & R \leq 40 % and Large BiCC: 80 %

Distribution of BiCC: Skewed, Slightly better quality

Thank You

Questions???