Econometrics August 24, 2023

# Topic 19: Community Detection

by Sai Zhang

### Key points: .

**Disclaimer**: The note is built on Prof. Jinchi Lv's lectures of the course at USC, DSO 607, High-Dimensional Statistics and Big Data Problems.

### 19.1 Stochastic Block Model

Consider an undirected graph G, with nodes V and edges E. Let

- *n* be a positive integer: the number of **vertices**
- *k* be a positive integer: the number of **communities**
- $p = (p_1, \dots, p_k)$  be a probability vector on  $\{1, \dots, k\} := [k]$ : the **prior** on the k communities
- **W** be a  $k \times k$  symmetric matrix with entries  $W_{ij} \in [0,1]$ : the matrix of **connectivity probabilities**

then we have

### **Definition 19.1.1: Stochastic Block Model**

The pair  $(\mathbf{X}, G)$  is drawn under  $SBM(n, p, \mathbf{W})$  if  $\mathbf{X}$  is an n dimensional random vector with i.i.d. components distributed under p, and G is an n-vertex simple graph where vertices i and j are connected with probability  $W_{X_i,X_j}$ , **independently** of other pairs of vertices. And the **community** sets can be defined by

$$\Omega_i = \Omega_i(\mathbf{X}) := \{v \in [n] : X_v = i\}, i \in [k]$$

Immediately, we can define the symmetry of SBM as:

#### **Definition 19.1.2: Symmetric SBM**

An SBM is called symmetric if

- p is uniform
- W takes the same value on the diagonal and the same value off the diagonal

 $(\mathbf{X}, G)$  is drawn under SSBM(n, k, A, B) if  $p = \{1/k\}^k$  and  $\mathbf{W}$  takes avolue A on the diagonal and B off the diagonal.

## 19.1.1 Recovery

The goal of community detection is to recover the labels X by observing G, up to some level of accuracy. First, define **agreement** as

#### **Definition 19.1.3: Agreement of Communities**

The agreement between two community vectors  $\mathbf{x}$ ,  $\mathbf{y} \in [k]^n$  is obtained by maximizing the common components between  $\mathbf{x}$  and any relabelling of  $\mathbf{y}$ , that is

$$A(\mathbf{x}, \mathbf{y}) = \max_{\pi \in S_k} \frac{1}{n} \sum_{i=1}^n \mathbf{1} \left[ x_i = \pi(y_i) \right]$$

where  $S_k$  is the group of permutations on [k].

The **relabelling** permutation is used to handle symmetric communities such as in SSBM, as it is impossible to recover the actual labels in this case. But it's possible to recover the **partition**. There are 2 types of partition recovery we consider

**Exact Recovery** First, consider the case of **exact recovery**:

### **Definition 19.1.4: Exact Recovery**

Let  $(\mathbf{X}, G) \sim SBM(n, p, W)$ , the exact recovery is solved if there exists an algorithm that takes G as an input and outpus  $\hat{\mathbf{X}} = \hat{\mathbf{X}}(G)$  such that  $\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) = 1\right\} = 1 - o_p(1)$ 

In the SSBM case, algorithms that guarantee

$$A(\mathbf{X}, \hat{\mathbf{X}}) \to \frac{1}{k}$$

would be trivial.

**Weak Recovery** On the other hand, we the case of **weak recovery** defined as

#### **Definition 19.1.5: Weak Recovery**

Weak recovery or detection is solved SSBM(n,k,A,B) if for  $(\mathbf{X},G) \sim SSBM(n,k,A,B)$ , then  $\exists \epsilon > 0$  and an algorithm that takes G as an input and outputs  $\hat{\mathbf{X}}$  such that

$$\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) \ge \frac{1}{k} + \epsilon\right\} = 1 - o(1)$$

## 19.1.2 **Example:** SSBM(n,2)

Let's look at the example of  $SSBM(n, 2, \alpha \frac{\log n}{n}, \beta \frac{\log n}{n})$