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# **Topic 19: Community Detection**

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### Key points: .

**Disclaimer**: The note is built on Prof. Jinchi Lv's lectures of the course at USC, DSO 607, High-Dimensional Statistics and Big Data Problems.

# 19.1 Stochastic Block Model (Abbe et al., 2015)

Consider an undirected graph G, with nodes V and edges E. Let

- *n* be a positive integer: the number of **vertices**
- *k* be a positive integer: the number of **communities**
- $p = (p_1, \dots, p_k)$  be a probability vector on  $\{1, \dots, k\} := [k]$ : the **prior** on the k communities
- **W** be a  $k \times k$  symmetric matrix with entries  $W_{ij} \in [0,1]$ : the matrix of **connectivity probabilities**

then we have

# **Definition 19.1.1: Stochastic Block Model**

The pair  $(\mathbf{X}, G)$  is drawn under  $SBM(n, p, \mathbf{W})$  if  $\mathbf{X}$  is an n dimensional random vector with i.i.d. components distributed under p, and G is an n-vertex simple graph where vertices i and j are connected with probability  $W_{X_i,X_j}$ , **independently** of other pairs of vertices. And the **community** sets can be defined by

$$\Omega_i = \Omega_i(\mathbf{X}) := \{v \in [n] : X_v = i\}, i \in [k]$$

Immediately, we can define the symmetry of SBM as:

### **Definition 19.1.2: Symmetric SBM**

An SBM is called symmetric if

- p is uniform
- W takes the same value on the diagonal and the same value off the diagonal

 $(\mathbf{X}, G)$  is drawn under SSBM(n, k, A, B) if  $p = \{1/k\}^k$  and  $\mathbf{W}$  takes avolue A on the diagonal and B off the diagonal.

## 19.1.1 Recovery

The goal of community detection is to recover the labels X by observing G, up to some level of accuracy. First, define **agreement** as

### **Definition 19.1.3: Agreement of Communities**

The agreement between two community vectors  $\mathbf{x}$ ,  $\mathbf{y} \in [k]^n$  is obtained by maximizing the common components between  $\mathbf{x}$  and any relabelling of  $\mathbf{y}$ , that is

$$A(\mathbf{x}, \mathbf{y}) = \max_{\pi \in S_k} \frac{1}{n} \sum_{i=1}^n \mathbf{1} [x_i = \pi(y_i)]$$

where  $S_k$  is the group of permutations on [k].

The **relabelling** permutation is used to handle symmetric communities such as in SSBM, as it is impossible to recover the actual labels in this case. But it's possible to recover the **partition**. There are 2 types of partition recovery we consider

**Exact Recovery** First, consider the case of **exact recovery**:

### **Definition 19.1.4: Exact Recovery**

Let  $(\mathbf{X}, G) \sim SBM(n, p, W)$ , the exact recovery is solved if there exists an algorithm that takes G as an input and outpus  $\hat{\mathbf{X}} = \hat{\mathbf{X}}(G)$  such that  $\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) = 1\right\} = 1 - o_p(1)$ 

In the SSBM case, algorithms that guarantee

$$A(\mathbf{X}, \hat{\mathbf{X}}) \to \frac{1}{k}$$

would be trivial.

Weak Recovery On the other hand, we the case of weak recovery defined as

#### **Definition 19.1.5: Weak Recovery**

Weak recovery or detection is solved SSBM(n,k,A,B) if for  $(\mathbf{X},G) \sim SSBM(n,k,A,B)$ , then  $\exists \epsilon > 0$  and an algorithm that takes G as an input and outputs  $\hat{\mathbf{X}}$  such that

$$\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) \ge \frac{1}{k} + \epsilon\right\} = 1 - o(1)$$

### 19.1.2 **Example:** SSBM(n,2)

Let's look at the example of  $SSBM(n, 2, \alpha \frac{\log n}{n}, \beta \frac{\log n}{n})$ , where

- *n*: number of vertices (assumed to be even for simplicity)
- for each  $v \in [n]$ , a binary label  $X_v$  is attached s.t.

$$|\{v \in [n] : X_v = 1\}| = n/2$$

• for each pair of distinct nodes  $u, v \in [n]$ , an edge is placed with probability

$$-\alpha \frac{\log n}{n} \text{ if } X_u = X_v$$

$$-\beta \frac{\log n}{n} \text{ if } X_u \neq X_v$$

where edges are placed independently conditionally on the vertex labels

• WLOG,  $\alpha > \beta$ 

then we have the following theorem

# **Theorem 19.1.6: Exact Recovery in** $SSBM(n, 2, \alpha \log(n)/n, \beta \log(n)/n)$

- Exact recovery in  $SSBM(n, 2, \alpha \log(n)/n, \beta \log(n)/n)$  is solvable and efficiently so if  $|\sqrt{\alpha} \sqrt{\beta}| > \sqrt{2}$  nad unsolvable if  $|\sqrt{\alpha} \sqrt{\beta}| < \sqrt{2}$
- Exact recovery of the ground truth assignment of the partition (A, B) is also achieveable, that is: if

$$\frac{\alpha + \beta}{2} - \sqrt{\alpha \beta} > 1$$

i.e.

$$\alpha + \beta > 2$$
,  $(\alpha - \beta)^2 > 4(\alpha + \beta) - 4$ 

the maximum likelihood estimator exactly recovers the communities (up to a global flip), with high probability.

See Abbe (2017) for the proof of this theorem.

In summary, for a graph structure G = (V, E) represented by adjacency matrix  $\mathbf{X}_{n \times n}$ , Stochastic Block Model (SBM)

- assumes that there is a symmetric matrix  $\mathbf{P} = \{p_{ij}\} \in \mathbb{R}^{k \times k}$ , for  $k \ll n$  and a map  $C : \{1, \dots, n\} \rightarrow \{1, \dots, k\}$ , s.t.  $\Pr(\mathbf{X}_{ij} = 1) = \mathbf{P}_{C(i), C(i)}$
- Define  $\Pi = (\pi_1, \dots, \pi_n)' \in \mathbb{R}^{n \times k}$  where  $\Pi_{ij} = 1$  if C(i) = j, and  $\Pi_{ij} = 0$  otherwise
- Let  $\mathbf{H} = \mathbb{E}(\mathbf{X})$  be the probability matrix, then  $\mathbf{H} = \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}'$
- A variant of SBM is degree corrected SBM which incorporates the degree heterogeneity.
  - each node is assigned a parameter  $\theta_i > 0$  such that  $\Pr(\mathbf{X}_{ij} = 1) = \theta_i \theta_j \mathbf{P}_{C(i),C(j)}$
  - $\mathbf{H} = \mathbf{\Theta} \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}' \mathbf{\Theta}$ , where  $\mathbf{\Theta} = \text{diag}(\theta_1, \dots, \theta_n)$

# **19.2** SIMPLE Model (Fan et al., 2022)

In SBM, each  $\pi_i \in \{e_1, \dots, e_K\}$  with  $e_k$  a one entry vector whose k-th component is one. But what if each node i can belong to K different communities? We generalize  $\pi_i$  to be a compositional vector, and interpret it as community membership profile for node i, then

$$\Pr\left(\mathbf{X}_{ij}=1\right) = \theta_i \theta_j \sum_{k=1}^K \sum_{l=1}^K \pi_i(k) \pi_j(l) p_{kl}$$

and  $\mathbf{H} = \mathbf{\Theta} \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}' \mathbf{\Theta}$ . Now, consider a new statistical tests for testing whether any given pair of nodes share the same membership profiles, and providing the associated p-values.

# 19.2.1 Problem Setting

For an undirected graph G = (V, E) with n nodes, let  $\mathbf{X} = \{x_{ij}\} \in \mathbb{R}^{n \times n}$  be the **symmetric** adjacency matrix. Under a probabilistic model, assume  $x_{ij}$  is an independent realization from a Bernoulli random variable for all upper triangular entries of random matrix  $\mathbf{X}$ . Consider the adjacency matrix with the deterministic-random latent structure

$$X = H + W$$

where

- $\mathbf{H} = \{h_{ij}\} \in \mathbb{R}^{n \times n}$  is the deterministic mean matrix of low rank  $K \ge 1$
- $\mathbf{W} = \{w_{ij}\} \in \mathbb{R}^{n \times n}$  is a symmetric random matrix with zero mean and independent entries on and above the diagonal

Assume *V* is decomposed into *K* disjoint latent communities

$$C_1, \cdots, C_K$$

where each node i is associated with the community membership probability vector

$$\boldsymbol{\pi}_i = (\pi_i(1), \cdots, \pi_i(K))' \in \mathbb{R}^K$$

s.t.

$$\Pr(i \in C_k) = \pi_i(k), \ k = 1, \dots, K$$

here, K is unknown but bounded away from  $\infty$ .

# 19.2.2 Hypothesis Testing

For any given pair of nodes  $i \neq j \in V$ , the goal is to infer whether they share the same community identity with quantified uncertainty level based on adjacency matrix X, the hypothesis is

$$H_0: \pi_i = \pi_j \qquad \qquad H_1: \pi_i \neq \pi_j$$

More explicitly, consider the DCMM (Degree Corrected Mixed Membership) model as the underlying network model, s.t. the probability of a link between nodes i and j can be written as

$$\Pr(\mathbf{X}_{ij} = 1) = \theta_i \theta_j \sum_{k=1}^K \sum_{l=1}^K \pi_i(k) \pi_j(l) p_{kl}$$

and

$$H = \Theta \Pi P \Pi' \Theta$$

in matrix form, where  $\Pi = (\pi_1, \dots, \pi_n)' \in \mathbb{R}^{n \times k}$  and  $\Theta = \text{diag}(\theta_1, \dots, \theta_n)$ . Consider

- No degree homogeneity:  $\mathbf{\Theta} = \sqrt{\theta}\mathbf{I}_n$ , then  $\mathbf{H} = \theta \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}'$ . If we eigen-decompose  $\mathbf{H} = \mathbf{V} \mathbf{D} \mathbf{V}'$  where  $\mathbf{D} = \operatorname{diag}(d_1, \cdots, d_K)$  with  $|d_1| \geq |d_2| \geq \cdots |d_K| > 0$  is the matrix of all K non-zero eigenvalues and  $V = (v_1, \cdots, v_K) \in \mathbb{R}^{n \times K}$  is the eigenvectors.
  - the column space spanned by  $\Pi$  is the same as the eigenspace spanned by the top K eigenvectors of matrix  $\mathbf{H}$
  - mean matrix **H** is **not** observable: replace it with adjacency matrix **X** and conduct eigen-decomposition to get eigenvalues  $\hat{d}_1, \dots, \hat{d}_n$  and eigenvectors  $\hat{v}_1, \dots, \hat{v}_n$ . We assume that

$$\left|\hat{d}_1\right| \ge \left|\hat{d}_2\right| \ge \dots \ge \left|\hat{d}_n\right|$$

and let 
$$\hat{\mathbf{V}} = (\hat{v}_1, \cdots, \hat{v}_K) \in \mathbf{R}^{n \times K}$$
.

**Without degree heterogeneity** first, consider the case where  $\mathbf{\Theta} = \sqrt{\theta} \mathbf{I}_n$  and  $\mathbb{E}(\mathbf{X}) = \mathbf{H} = \theta \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}'$ . If  $\pi_i = \pi_j$ , then nodes i and j are exchangeable and  $\mathbf{V}(i) = \mathbf{V}(j)$ . The test statistic for membership information of node i and j is given as

$$T_{ij} = \left[\hat{\mathbf{V}}\right]$$

# References

Emmanuel Abbe. Community detection and stochastic block models: recent developments. *The Journal of Machine Learning Research*, 18(1):6446–6531, 2017.