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Topic 19: Community Detection

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Key points: .

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19.1 Stochastic Block Model (Abbe et al., 2015)

Consider an undirected graph *G*, with nodes *V* and edges *E*. Let

- *n* be a positive integer: the number of **vertices**
- *k* be a positive integer: the number of **communities**
- $p = (p_1, \dots, p_k)$ be a probability vector on $\{1, \dots, k\} := [k]$: the **prior** on the k communities
- **W** be a $k \times k$ symmetric matrix with entries $W_{ij} \in [0,1]$: the matrix of **connectivity probabilities**

then we have

Definition 19.1.1: Stochastic Block Model

The pair (\mathbf{X}, G) is drawn under $SBM(n, p, \mathbf{W})$ if \mathbf{X} is an n dimensional random vector with i.i.d. components distributed under p, and G is an n-vertex simple graph where vertices i and j are connected with probability W_{X_i,X_j} , **independently** of other pairs of vertices. And the **community** sets can be defined by

$$\Omega_i = \Omega_i(\mathbf{X}) := \{v \in [n] : X_v = i\}, i \in [k]$$

Immediately, we can define the symmetry of SBM as:

Definition 19.1.2: Symmetric SBM

An SBM is called symmetric if

- p is uniform
- W takes the same value on the diagonal and the same value off the diagonal

 (\mathbf{X}, G) is drawn under SSBM(n, k, A, B) if $p = \{1/k\}^k$ and \mathbf{W} takes avolue A on the diagonal and B off the diagonal.

19.1.1 Recovery

The goal of community detection is to recover the labels X by observing G, up to some level of accuracy. First, define **agreement** as

Definition 19.1.3: Agreement of Communities

The agreement between two community vectors \mathbf{x} , $\mathbf{y} \in [k]^n$ is obtained by maximizing the common components between \mathbf{x} and any relabelling of \mathbf{y} , that is

$$A(\mathbf{x}, \mathbf{y}) = \max_{\pi \in S_k} \frac{1}{n} \sum_{i=1}^{n} \mathbf{1} \left[x_i = \pi(y_i) \right]$$

where S_k is the group of permutations on [k].

The **relabelling** permutation is used to handle symmetric communities such as in SSBM, as it is impossible to recover the actual labels in this case. But it's possible to recover the **partition**. There are 2 types of partition recovery we consider

Exact Recovery First, consider the case of **exact recovery**:

Definition 19.1.4: Exact Recovery

Let $(\mathbf{X}, G) \sim SBM(n, p, W)$, the exact recovery is solved if there exists an algorithm that takes G as an input and outpus $\hat{\mathbf{X}} = \hat{\mathbf{X}}(G)$ such that $\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) = 1\right\} = 1 - o_p(1)$

In the SSBM case, algorithms that guarantee

$$A(\mathbf{X}, \hat{\mathbf{X}}) \to \frac{1}{k}$$

would be trivial.

Weak Recovery On the other hand, we the case of weak recovery defined as

Definition 19.1.5: Weak Recovery

Weak recovery or detection is solved SSBM(n,k,A,B) if for $(\mathbf{X},G) \sim SSBM(n,k,A,B)$, then $\exists \epsilon > 0$ and an algorithm that takes G as an input and outputs $\hat{\mathbf{X}}$ such that

$$\mathbb{P}\left\{A(\mathbf{X}, \hat{\mathbf{X}}) \ge \frac{1}{k} + \epsilon\right\} = 1 - o(1)$$

19.1.2 **Example:** SSBM(n,2)

Let's look at the example of $SSBM(n, 2, \alpha \frac{\log n}{n}, \beta \frac{\log n}{n})$, where

- *n*: number of vertices (assumed to be even for simplicity)
- for each $v \in [n]$, a binary label X_v is attached s.t.

$$|\{v \in [n] : X_v = 1\}| = n/2$$

• for each pair of distinct nodes $u, v \in [n]$, an edge is placed with probability

$$-\alpha \frac{\log n}{n} \text{ if } X_u = X_v$$

$$-\beta \frac{\log n}{n} \text{ if } X_u \neq X_v$$

where edges are placed independently conditionally on the vertex labels

• WLOG, $\alpha > \beta$

then we have the following theorem

Theorem 19.1.6: Exact Recovery in $SSBM(n, 2, \alpha \log(n)/n, \beta \log(n)/n)$

- Exact recovery in $SSBM(n, 2, \alpha \log(n)/n, \beta \log(n)/n)$ is solvable and efficiently so if $|\sqrt{\alpha} \sqrt{\beta}| > \sqrt{2}$ nad unsolvable if $|\sqrt{\alpha} \sqrt{\beta}| < \sqrt{2}$
- Exact recovery of the ground truth assignment of the partition (A, B) is also achieveable, that is: if

$$\frac{\alpha + \beta}{2} - \sqrt{\alpha \beta} > 1$$

i.e.

$$\alpha + \beta > 2$$
, $(\alpha - \beta)^2 > 4(\alpha + \beta) - 4$

the maximum likelihood estimator exactly recovers the communities (up to a global flip), with high probability.

See Abbe (2017) for the proof of this theorem.

In summary, for a graph structure G = (V, E) represented by adjacency matrix $\mathbf{X}_{n \times n}$, Stochastic Block Model (SBM)

- assumes that there is a symmetric matrix $\mathbf{P} = \{p_{ij}\} \in \mathbb{R}^{k \times k}$, for $k \ll n$ and a map $C : \{1, \dots, n\} \rightarrow \{1, \dots, k\}$, s.t. $\Pr(\mathbf{X}_{ij} = 1) = \mathbf{P}_{C(i), C(i)}$
- Define $\Pi = (\pi_1, \dots, \pi_n)' \in \mathbb{R}^{n \times k}$ where $\Pi_{ij} = 1$ if C(i) = j, and $\Pi_{ij} = 0$ otherwise
- Let $H = \mathbb{E}(X)$ be the probability matrix, then $H = \Pi P\Pi'$
- A variant of SBM is degree corrected SBM which incorporates the degree heterogeneity.
 - each node is assigned a parameter $\theta_i > 0$ such that $\Pr(\mathbf{X}_{ij} = 1) = \theta_i \theta_j \mathbf{P}_{C(i),C(j)}$
 - $\mathbf{H} = \mathbf{\Theta} \mathbf{\Pi} \mathbf{P} \mathbf{\Pi}' \mathbf{\Theta}$, where $\mathbf{\Theta} = \operatorname{diag}(\theta_1, \dots, \theta_n)$

19.2 SIMPLE Model (Fan et al., 2022)

In SBM, each $\pi_i \in \{e_1, \dots, e_K\}$ with e_k a one entry vector whose k-th component is one. But what if each node i can belong to K different communities? We generalize π_i to be a compositional vector, and interpret it as community membership profile for node i, then

$$\Pr\left(\mathbf{X}_{ij}=1\right) = \theta_i \theta_j \sum_{k=1}^K \sum_{l=1}^K \pi_i(k) \pi_j(l) p_{kl}$$

and $H = \Theta \Pi P \Pi' \Theta$

References

Emmanuel Abbe. Community detection and stochastic block models: recent developments. *The Journal of Machine Learning Research*, 18(1):6446–6531, 2017.