

# Probability in Computing

## Assignment - 4

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1) The probability distribution vector after 4 time steps is  
 $[0.1938, 0.1772, 0.629]$

2)  $T_{ij}(\text{expected time after } i \text{ to } j) = j^n - i^n$   
 $(i \leq j)$

3) expected number of steps to reach is 10.

4) Stationary distribution for the given random walk is:  
 $\left[\frac{2}{7}, \frac{1}{7}, \frac{3}{14}, \frac{1}{7}, \frac{1}{7}, \frac{1}{14}\right]$

5a) we use constructive proof to show that there exists a colouring with two colours of graph s.t. no triangle is monochromatic.

Given a graph with 3-colouring using red, blue and green. Replace all nodes' colour having blue with green. Now, in the graph, every ~~triangle~~ triangle previously had 3-coloured nodes but since one of them is changed, every triangle must have 2 colours (~~one~~ red which is unchanged and green). There can't be a triangle with all green as that would imply red wasn't there previously and hence would violate 3 colouring.

Hence, proved.



5b) we can use Schöningh's algorithm (variant).

1) Pick a random assignment  $(x_1, x_2, \dots, x_n) = (a_1, a_2, \dots, a_n)$ .  
//  $x_i$  are nodes and  $a_i \in \{R, B\}$ .

2) Repeat  $3n$  times, terminating if a satisfying assignment is found

- If there exists a monochromatic triangle, do:

- Pick a random vertex (uniformly) from the triangle, and flip the colour of node

- Re-evaluate the graph (check if ~~2-colourable~~ 2 colouring satisfies)

3) If assignment doesn't satisfy (there exists a monochromatic triangle), GOTO step 1 (restart)

let  $X_t = \frac{(\text{num of})}{(\text{nodes})}$  matching the current assignment,  $(A_1, A_2, \dots, A_n)$   
where  $t$  is number of time steps.

Consider a random walk on  $\{0, 1, \dots, n\}$  with  $X_t$ .  
Need to calculate expected time for  $X_t = n$ .

consider a triangle with vertices  $(v_1, v_2, v_3)$ .

without loss of generality, let the colouring be  $\{R, R, R\}$ .

~~Let  $A_i$  be the correct assignment from the solution.~~  
The current assignment can be of the form:

$(B, R, R), (R, B, R), (R, R, B)$

or  $(R, B, B), (B, R, B), (B, B, R)$

→ If so with  $\text{Prob} = \frac{1}{3}$ , we get a monochromatic  $\Delta$  & then flip a node to increase  $X_t$  value by 1

↓  
Similar to previous case but with  $\text{Prob} = \frac{2}{3}$  can increase  $X_t$  value by 1.

hence, with prob = atleast  $\frac{1}{3}$ ,  $x_t$  increases by 1 every time. This is similar to the case of Schöningh's algo.

Thus, using that algorithm ~~as~~ result, the expected run time for  $x_t = n$  is  $O(1.33^n)$