Effect handlers in OCaml

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Concurrency \neq Parallelism

- Concurrency
 - Programming technique
 - Overlapped execution of processes
- Parallelism
 - (*Extreme*) Performance hack
 - Simultaneous execution of computations

Concurrency ∩ Parallelism → Scalable Concurrency (Fibers) (Domains)

Schedulers

- Multiplexing fibers over domain(s)
 - Bake scheduler into the runtime system (GHC)
- Allow programmers to describe schedulers!
 - Parallel search —> LIFO work-stealing
 - Web-server —> FIFO runqueue
 - Data parallel —> Gang scheduling
- Algebraic Effects and Handlers

Algebraic Effects: Example

Dynamic wind

```
let dynamic_wind before_thunk thunk after_thunk =
  before_thunk ();
  let res =
    match thunk () with
    V -> V
    l exception e -> after_thunk (); raise e
    l effect e k ->
        after_thunk ();
        let res' = perform e in
        before_thunk ();
        continue k res'
  in
  after_thunk ();
  res
```

Effect systems and modularity

- Right now, we type effects like ML exceptions
 - (we're in the market for an effect system, if anyone has one lying around...)
- We need modularity, because effects can:
 - be local, dynamic and fresh
 - be abstracted, renamed and reuse
- We don't know statically whether two effects are the same

Scheduler Demo¹

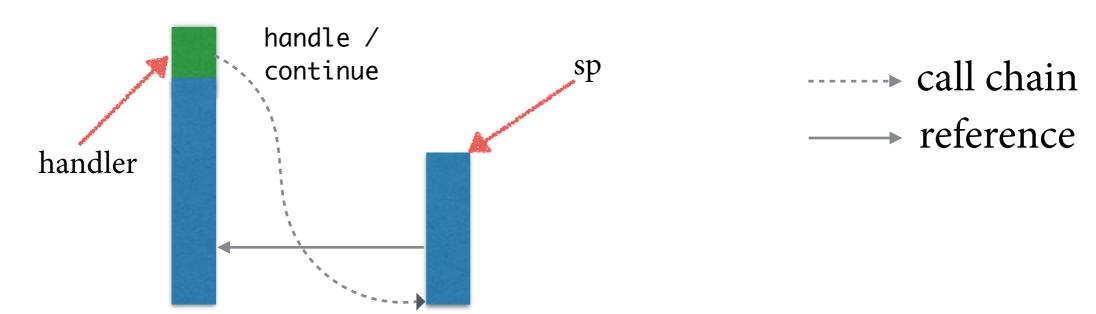
Generator from Iterator¹

```
type 'a t =
| Leaf
| Node of 'a t * 'a * 'a t
let rec iter f = function
  | Leaf -> ()
  Node (1, x, r) \rightarrow iter f l; f x; iter f r
(* val to_gen : 'a t -> (unit -> 'a option) *)
let to_gen (type a) (t : a t) =
  let module M = struct effect Next : a -> unit end in
  let open M in
  let step = ref (fun () -> assert false) in
  let first_step () =
    try
      iter (fun x -> perform (Next x)) t; None
    with effect (Next v) k ->
      step := continue k; Some v
  in
    step := first_step;
    fun () -> !step ()
```

[1] https://github.com/kayceesrk/ocaml15-eff/blob/master/generator.ml

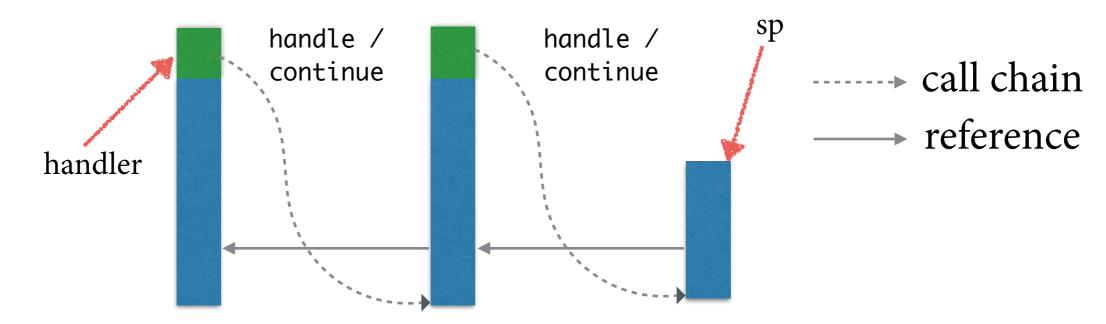
Implementation

- Fibers: Heap allocated, dynamically resized stacks
 - \sim 10s of bytes
 - No unnecessary closure allocation costs unlike CPS
- One-shot delimited continuations
 - Simplifies reasoning about resources sockets, locks, etc.
- Handlers —> Linked-list of fibers



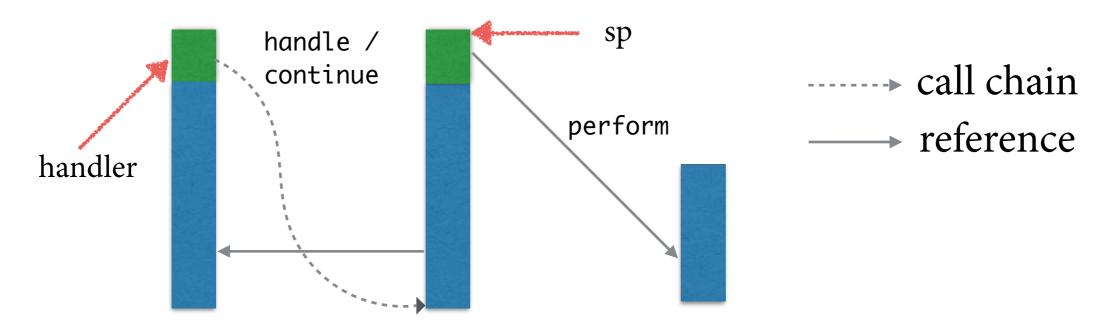
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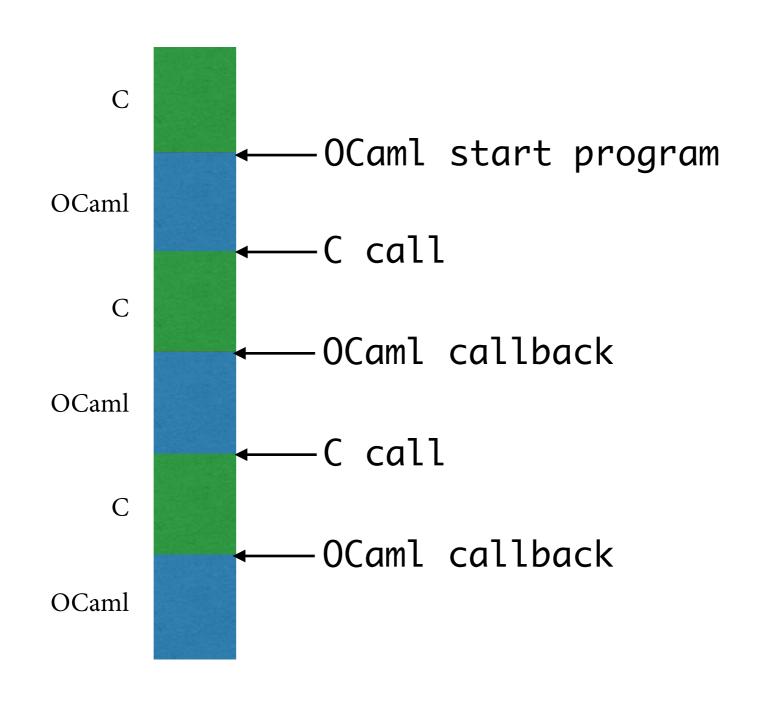


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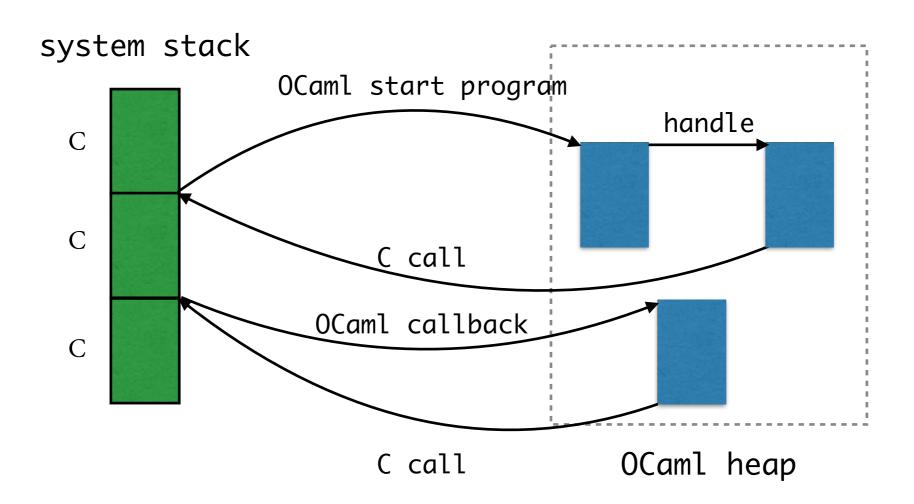
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Native-code fibers — Vanilla



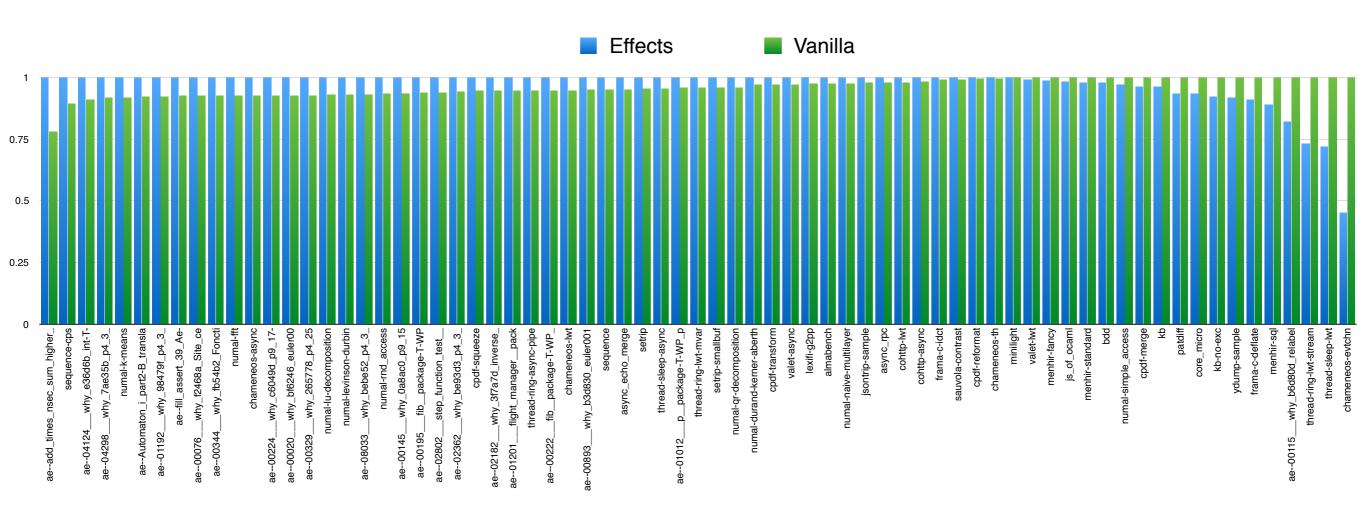
Native-code fibers — Effects



- 1. Stack overflow checks for OCaml functions
 - Simple static analysis eliminates many checks
- 2. FFI calls are more expensive due to stack switching
 - Specialise for calls which {allocate / pass arguments on stack / do neither}

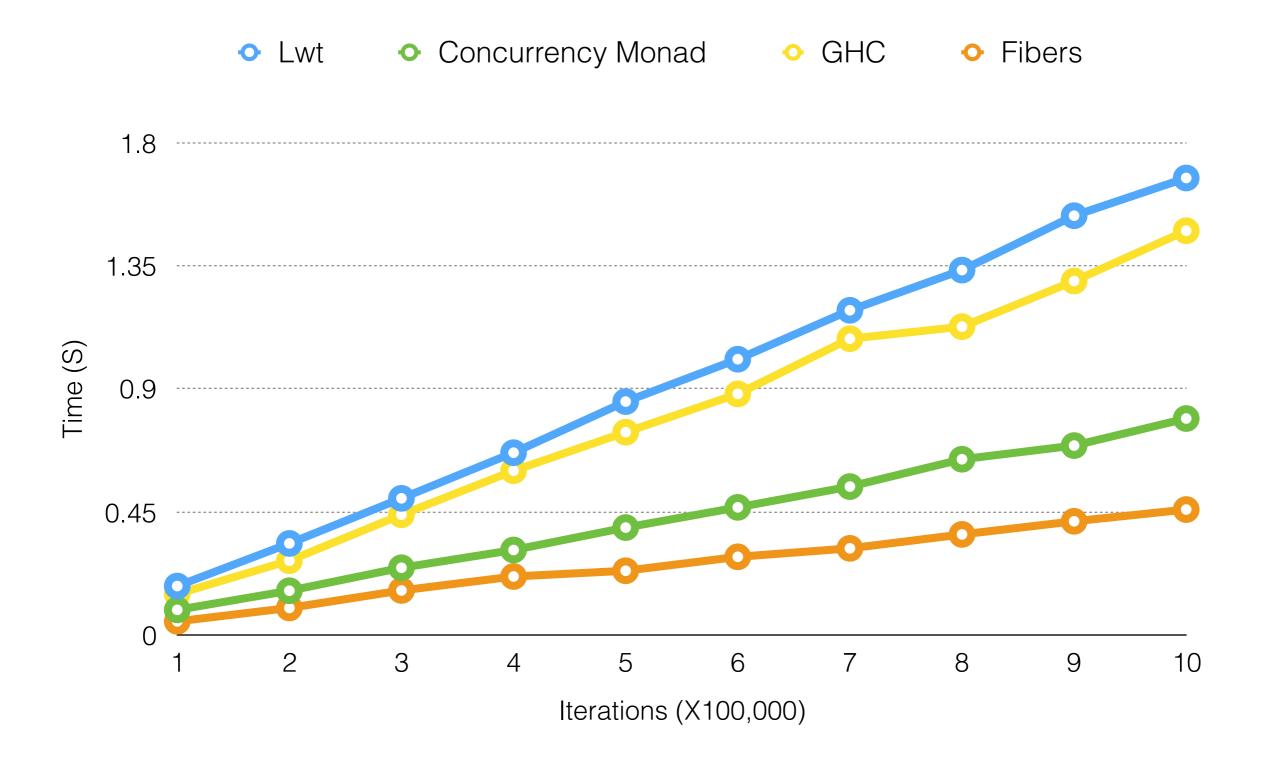
Performance: Vanilla OCaml

Normalised time (lower is better)

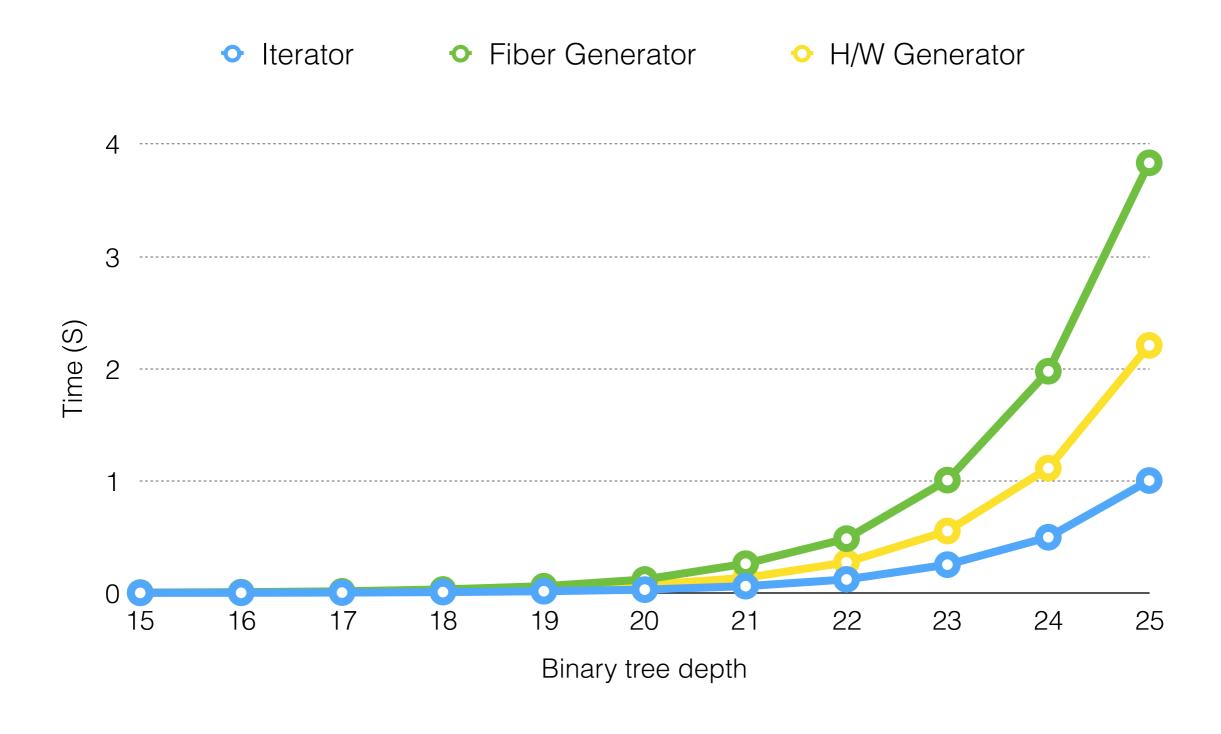


Effects ~0.9% slower

Performance: Chameneos-Redux



Performance: Generator



Javascript backend

- js_of_ocaml
 - OCaml bytecode —> Javascript
- js_of_ocaml compiler pass
 - Whole-program selective CPS transformation
- Work-in-progress!
 - Selective CPS translation

fin.

Multicore OCaml: https://github.com/ocamllabs/ocaml-multicore

Examples: https://github.com/kayceesrk/ocaml-eff-example