

1 . [Recursion on Subsequences](#)

(**L6. Recursion on Subsequences | Printing Subsequences**)

Power-Set is The Most Efficient method for generating subsequences (

Power Set | Print all Subsequences)

Note : [Leetcode-pattern post](#)

```
1. import java.io.*;
2. import java.util.*;
3. class Subsequence {
4.     public static void main (String[] args) {
5.         int arr[]={1,2,3};
6.         int n=arr.length;
7.         helper(arr,n,0,new ArrayList<Integer>());
8.
9.     }
10.    static void helper(int arr[],int n,int i,ArrayList<Integer> list){
11.        if(i==n){
12.            System.out.println(list);
13.            return;
14.        }
15.        list.add(arr[i]);
16.        //Take the particular index into Subsequences
17.        helper(arr,n,i+1,list);
18.        list.remove(list.size()-1);
19.        //Do Not Take the particular index into Subsequences
20.        helper(arr,n,i+1,list);
21.    }
22. }
```

Questions On same pattern as Above Approach

2. Subsequence sum equal to k

(**L7. All Kind of Patterns in Recursion | Print All | Print one | Count**)

```
1. class Solution {
2.     public List<List<Integer>> SubsequenceSumEqualToK(int[] candidates, int target) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         helper(candidates,0,target,ans,new ArrayList<Integer>(),0);
5.         return ans;
6.     }
7.     void helper(int arr[],int i,int target,List<List<Integer>> ans,ArrayList<Integer>
temp,int sum){
8.         if(i==arr.length){
9.             if(sum==target) ans.add(new ArrayList<>(temp));
10.            return;
11.        }
12.        temp.add(arr[i]); sum+=arr[i];
13.        helper(arr,i+1,target,ans,temp,sum);
14.        int val=temp.remove(temp.size()-1); sum-=arr[i];
15.        helper(arr,i+1,target,ans,temp,sum);
16.    }
17. }
```

3. Print only 1 Subsequence Whose sum equal to K

( L7. All Kind of Patterns in Recursion | Print All | Print one | Count)

```
1. import java.util.*;
2. import java.io.*;
3. class Solution {
4.     static public void SubsequenceSumEqualToK(int[] candidates, int target) {
5.         helper(candidates,0,target,new ArrayList<Integer>(),0);
6.     }
7.     static boolean helper(int arr[],int i,int target,ArrayList<Integer> temp,int sum){
8.         if(i==arr.length){
9.             if(sum==target){
10.                 System.out.println(temp);
11.                 return true; //condition satisfied
12.             }
13.             return false; //condition does not satisfied
14.         }
15.         temp.add(arr[i]); sum+=arr[i];
16.         if(helper(arr,i+1,target,temp,sum)) return true;
17.         int val=temp.remove(temp.size()-1); sum-=arr[i];
18.         if(helper(arr,i+1,target,temp,sum)) return true;
19.         return false;
20.     }
21.     public static void main (String[] args) {
22.         SubsequenceSumEqualToK(new int[]{1,2,3,4,5},6);
23.     }
24. }
```

4.Count Subsequence sum equal to k

( L7. All Kind of Patterns in Recursion | Print All | Print one | Count)

```
1. import java.util.*;
2. import java.io.*;
3. class Solution {
4.     static public void SubsequenceSumEqualToK(int[] candidates, int target) {
5.
6.         System.out.println(helper(candidates,0,target,0));
7.     }
8.     static int helper(int arr[],int i,int target,int sum){
9.         if(i==arr.length){
10.             if(sum==target){
11.                 return 1; // if found return 1
12.             }
13.             return 0; //else return 0
14.         }
15.         sum+=arr[i];
16.         int l=helper(arr,i+1,target,sum); // answer from left call
17.         sum-=arr[i];
18.         int r=helper(arr,i+1,target,sum); // answer from right call
19.         return l+r; //return ans from left+right call
20.     }
21.     public static void main (String[] args) {
22.         SubsequenceSumEqualToK(new int[]{1,2,3,4,5},6);
23.     }
24. }
```

Some Approches on find combination sum equal to k

5. [Combination Sum](#) - You can pick an element any number of time :

(L8. Combination Sum | Recursion | Leetcode | C++ | Java)

```
1. class Solution {
2.     public List<List<Integer>> combinationSum(int[] candidates, int target) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         helper(candidates,0,target,ans,new ArrayList<Integer>());
5.         return ans;
6.     }
7.     void helper(int arr[],int i,int target,List<List<Integer>> res,List<Integer> ds){
8.         if(i==arr.length){
9.             if(target==0){
10.                 res.add(new ArrayList<>(ds));
11.             }
12.             return;
13.         }
14.         if(arr[i]<=target){ //only pick when the current element is less than or equal to target
15.             ds.add(arr[i]); // add current element
16.             //instead of adding current elemnt to a new variable we will try to
            reduce the current target to zero
17.             helper(arr,i,target-arr[i],res,ds); //pick, pick and pick the current element until ,
            the target become less than the current element
18.             ds.remove(ds.size()-1); //remove the current element
19.         }
20.         helper(arr,i+1,target,res,ds); // does not include current element
21.     }
22. }
```

- time complexity: $2^t(\text{target}) * k$
- Space complexity : $k(\text{Avg. Length}) * x(\text{combinations})$

6. [Combination Sum II](#)

(L9. Combination Sum II | Leetcode | Recursion | Java | C++)

```
1. class Solution {
2.     public List<List<Integer>> combinationSum2(int[] candidates, int target) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         Arrays.sort(candidates); //we are sorting since we need to give ans in sorted order and ignore duplicate
            combinations
5.         helper(candidates,0,target,ans,new ArrayList<Integer>());
6.         return ans;
7.     }
8.     void helper(int arr[],int ind,int target,List<List<Integer>> ans,ArrayList<Integer> ds){
9.         if(target==0){
10.             ans.add(new ArrayList<>(ds));
11.             return;
12.         }
13.         for(int i=ind; i<arr.length; i++){
14.             if(i>ind && arr[i]==arr[i-1]) continue; // ignore if same value occure for getting all unique
            combinations
15.             if(arr[i]>target) break; // this means we can not form combinations
16.             ds.add(arr[i]);
17.             helper(arr,i+1,target-arr[i],ans,ds);
18.             ds.remove(ds.size()-1);
19.         }
20.     }
21. }
```

- time complexity: $2^n(\text{for genrating all combination}) * k(\text{avg length of every combination,for adding combination to answers})$
- Space complexity : $k(\text{Avg. Length}) * x(\text{combinations})$

Approaches on subsets sum

[Subset-sum](#) : ([YouTube](#) L10. Subset Sum I | Recursion | C++ | Java)

```
1. import java.io.*;
2. import java.util.*;
3.
4. class Solution{
5.     ArrayList<Integer> subsetSums(ArrayList<Integer> arr, int N){
6.         ArrayList<Integer> ans=new ArrayList<>();
7.         helper(arr,0,0,ans);
8.         return ans;
9.     }
10.    void helper(ArrayList<Integer> arr,int idx,int sum,ArrayList<Integer> ans){
11.        if(idx==arr.size()){
12.            ans.add(sum);
13.            return;
14.        }
15.        sum+=arr.get(idx);
16.        helper(arr,idx+1,sum,ans);
17.        sum-=arr.get(idx);
18.        helper(arr,idx+1,sum,ans);
19.    }
20. }
```

- time complexity: 2^n
- Space complexity : 2^n (i did not understand why)
-

[Subset II](#) : ([YouTube](#) L11. Subset Sum II | Leetcode | Recursion)

```
1. class Solution {
2.     public List<List<Integer>> subsetsWithDup(int[] nums) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         Arrays.sort(nums);
5.         helper(nums,nums.length,0,new ArrayList<Integer>(),ans);
6.         return ans;
7.     }
8.     static void helper(int arr[],int n,int idx,ArrayList<Integer> list,List<List<Integer>> ans){
9.         // if(idx==n){
10.            ans.add(new ArrayList<>(list));
11.            // return;
12.            // }
13.            for(int i=idx; i<n; i++){
14.                if(i>idx && arr[i]==arr[i-1]) continue;
15.                list.add(arr[i]);
16.                //Take the particular index into Subsequences
17.                helper(arr,n,i+1,list,ans);
18.                list.remove(list.size()-1);
19.                //Do Not Take the particular index into Subsequences
20.            }
21.        }
22.    }
```

- time complexity: $2^n * n$ (assume that every subset is near about size n)
- Space complexity : $2^n * k$ (subsets avg length k)

Print all Permutations

Permutations :

(📺 L12. Print all Permutations of a String/Array | Recursion | Approach - 1)

```
1. class Solution {
2.     public List<List<Integer>> permute(int[] nums) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         helper(nums,nums.length,new ArrayList<Integer>(),ans,new boolean[nums.length]);
5.         return ans;
6.     }
7.     static void helper(int arr[],int n,ArrayList<Integer> list,List<List<Integer>> ans,boolean vis[]){
8.         if(list.size()==n){ //when the size of array data structure become equals to n this mean we have form out
           first permutation
9.             ans.add(new ArrayList<>(list));
10.            return;
11.        }
12.        for(int i=0; i<n; i++){
13.            if(vis[i]) continue; // if the current element is used before in this permutation do not use it
14.            list.add(arr[i]); vis[i]=true; //when use the current element mark it as true so that we can not use it
           again
15.            helper(arr,n,list,ans,vis);
16.            list.remove(list.size()-1); vis[i]=false; //while backtrack mark the used element as false so that we
           can use the same element in different permutation
17.        }
18.    }
19.
20. }
```

- time complexity: $n!$ (for permutations) * n
- Space complexity :if we ignore for ans operations then $O(n)$ (store ds) + $O(n)$ (map array)

Permutations Approach 2 :

(📺 L13. Print all Permutations of a String/Array | Recursion | Approach - 2)

```
1. class Solution {
2.     public List<List<Integer>> permute(int[] nums) {
3.         List<List<Integer>> ans=new ArrayList<>();
4.         helper(nums,nums.length,0,ans);
5.         return ans;
6.     }
7.     void helper(int arr[],int n,int idx,List<List<Integer>> ans){
8.         if(idx==n){ //if index is crossing the array boundary then push the current state of
           permuted numbers
9.             //in the array to ans
10.            ArrayList<Integer> ds=new ArrayList<>();
11.            for(int num:arr) ds.add(num);
12.            ans.add(new ArrayList<>(ds));
13.        }
14.        for(int i=idx; i<n; i++){ //try swapping this number with every number ahead
15.            swap(arr,i,idx); // swap each element till its nth element
16.            helper(arr,n,idx+1,ans); //recursively call the function from next index
17.            swap(arr,i,idx); //swap back to original state so that when recursion call returns
           to this level it can explore other possibilities
18.        }
19.    }
20.    void swap(int arr[],int i,int j){
21.        int temp=arr[i];
22.        arr[i]=arr[j];
23.        arr[j]=temp;
24.    }
25.
26. }
```

- time complexity: $n!$ (for permutations) * n
- Space complexity :if we ignore for ans operations . (recursion space $O(n)$, $O(n!)$ for returning answers)

N-Queens : (L14. N-Queens | Leetcode Hard | Backtracking)

```
1. class Solution {
2.     public List<List<String>> solveNQueens(int n) {
3.         char board[][]=new char[n][n];
4.         List<List<String>> ans=new ArrayList<>();
5.         for(int i=0; i<n; i++){
6.             for(int j=0; j<n; j++){
7.                 board[i][j]='.';
8.             }
9.         }
10.        helper(0,board,ans);
11.        return ans;
12.    }
13. void helper(int row,char board[][],List<List<String>> ans){
14.     if(row==board.length){
15.         ArrayList<String> res=new ArrayList<>();
16.         for(char ch[]:board){
17.             res.add(new String(ch));
18.         }
19.         ans.add(res);
20.         return;
21.     }
22.     for(int col=0; col<board.length; col++){
23.         if(board[row][col]=='.' && isQueenSafe(board,row,col)){ //check is the
queen postion is safe for placing or Not
24.             board[row][col]='Q'; //if safe then place thq queen
25.             helper(row+1,board,ans);
26.             board[row][col]='.'; //while backtrack remove the queen to explore
other posblties
27.         }
28.     }
29. }
30. //if(board[row][col]=='.' && left[col]==false && diagonal[row+col]==false &&
upper[(board.length-1)+(col-row)]==false) ← this hashing method is also can be used to
check whether the queen is placed at the right position or not
31.
32.     boolean isQueenSafe(char chess[][],int row,int col){ //function to check wheter
placing the queen at this postion is safe or not
33.         for(int i = row - 1, j = col - 1; i >= 0 && j >= 0; i--, j--)
34.             if(chess[i][j] == 'Q')
35.                 return false;
36.         for(int i = row - 1, j = col; i >= 0; i--)
37.             if(chess[i][j] == 'Q')
38.                 return false;
39.         for(int i = row - 1, j = col + 1; i >= 0 && j < chess.length; i--, j++){
40.             if(chess[i][j] == 'Q')
41.                 return false;
42.         }
43.         for(int i = row, j = col - 1; j >= 0; j--)
44.             if(chess[i][j] == 'Q')
45.                 return false;
46.         return true;
47. }
```

- time complexity:
- Space complexity :

Sudoku Solver : (L15. Sudoku Solver | Backtracking)

```
1. class Solution {
2.     public void solveSudoku(char[][] board) {
3.         helper(board);
4.     }
5.     boolean helper(char[][] board){
6.         for(int i=0; i<9; i++){
7.             for(int j=0; j<9; j++){
8.                 if(board[i][j]=='.'){ //traverse the matrix and find the empty
place
9. //                                once we find the empty place than we tried all the
numbers from 1 to 9 and check that it is a valid number or not by checking the
rules.
10.// and we find the correct number for that place than we find for the second
empty place in 9 X 9 matrix.for second empty place we repeat the same process
11.                 for(char ch='1'; ch<='9'; ch++){
12.                     if(isValid(board,i,j,ch)){
13.                         board[i][j]=ch;
14.                         if(helper(board)) return true; //and after all
recursive calls we got true, than we have to stop over there only and no need to
search for other solutions.
15.                     else board[i][j]='.'; //after getting the false from
solve(board) function we make all the places empty that we have filled . than
try for other member for first empty place
16.                 }
17.             }
18.             return false; //and if we doesn't get any number, so we
return false.
19.         }
20.     }
21. }
22. }
23. return true;
24. }
25. private boolean isValid(char[][] board, int row, int col, char c){
26.     for(int i = 0; i < 9; i++) {
27.         if(board[i][col] != '.' && board[i][col] == c) return false; //check
row
28.         if(board[row][i] != '.' && board[row][i] == c) return false; //check
column
29.         if(board[3 * (row / 3) + i / 3][3 * (col / 3) + i % 3] == c) return
false; //check 3*3 block
30.     }
31.     return true;
32. }
33. }
```


- time complexity: 9^m (m represents the number of blanks to be filled in), since each blank can have 9 choices.
- Space complexity :

M-Coloring Problem: (L16. M-Coloring Problem | Backtracking)

```
1. class solve {
2.     // Function to determine if graph can be coloured with at most M
   colours
3.     // such
4.     // that no two adjacent vertices of graph are coloured with same
   colour.
5.     public boolean graphColoring(boolean graph[][], int m, int n) {
6.         return helper(graph,new int[n],m,n,0);
7.     }
8.     boolean helper(boolean graph[][],int col[],int m,int n,int idx){
9.         if(idx==n) return true;
10.
11.         for(int i=1; i<=m; i++){ //trying for each colour
12.             if(isSafe(graph,col,i,idx)){ //is placing the colour
   isSafe(no adjcent node have same color) then we will color the node
   wich particlar color
13.                 col[idx]=i;
14.                 if(helper(graph,col,m,n,idx+1)) return true; //if
   this returns true this means we have sucessfully colord our graph
15.                 else col[idx]=0; //if it is not possible to color
   it with ith colour then we will get it backtrack to its original state
   to explore other possibilities
16.             }
17.         }
18.         return false;
19.     }
20.     boolean isSafe(boolean graph[][],int col[],int color,int idx){
21.         for(int i=0; i<graph.length; i++){ //checking adjacent
   nodes
22.             if(graph[idx][i]){
23.                 if(col[i]==color) return false;
24.             }
25.         }
26.         return true;
27.     }
28. }
```

- time complexity: $(M^N) * N(isSafe())$
- Space complexity : $N(color\ array) + N(recursion\ space)$

Palindrome Partitioning :

( L17. Palindrome Partitioning | Leetcode | Recursion | C++ | Java)

```
1. class Solution {
2.     public List<List<String>> partition(String s) {
3.         List<List<String>> ans=new ArrayList<>();
4.         helper(s,0,ans,new ArrayList<String>());
5.         return ans;
6.     }
7.     void helper(String s,int idx,List<List<String>>
ans,ArrayList<String> ds){
8.         if(idx==s.length()){
9.             ans.add(new ArrayList<>(ds));
10.            return;
11.        }
12.        for(int i=idx; i<s.length(); i++){
13.            if(isPalindrom(s,idx,i)){
14.                //1. check if the choosen substring is a palindrom
or not , if the current substring is a palindrom then add this into
our DS and check for further substring who are palindrom.
15.                //2. if the input is "aab", check if [0,0] "a" is
palindrome. then check [0,1] "aa", then [0,2] "aab".
16.                // While checking [0,0], the rest of string is "ab", use ab as
input to make a recursive call.
17.                ds.add(s.substring(idx,i+1)); //choose
18.                helper(s,i+1,ans,ds);
19.                ds.remove(ds.size()-1); //unchoose
20.            }
21.        }
22.    }
23.    boolean isPalindrom(String str,int l,int r){
24.        while(l<=r){
25.            if(str.charAt(l++)!=str.charAt(r--)) return false;
26.        }
27.        return true;
28.    }
29. }
```

- time complexity: $O(n \cdot (2^n))$
- Space complexity :

[Rat in a Maze](#) : ( L19. Rat in A Maze | Backtracking)

```
1. class Solution {
2.     public static ArrayList<String> findPath(int[][] m, int
    n) {
3.         ArrayList<String> ans=new ArrayList<>();
4.         helper(m,0,0,n,ans,new String());
5.         return ans;
6.     }
7.     static void helper(int m[][],int row,int col,int
    n,ArrayList<String> ans,String str){
8.         if(row<0 || col<0 || row>=n || col>=n ||
    m[row][col]==0) return;
9.         if(row==n-1 && col==n-1){
10.             ans.add(str);
11.             // System.out.println(str);
12.             return;
13.         }
14.         m[row][col]=0; //smart way of mark the direction as
    visited we make this cell 0 so in this qus zero treated as the
    object in the path sow we can not move again in this
    direction if we mark this node as 0
15.         helper(m,row+1,col,n,ans,str+"D");
16.         helper(m,row,col+1,n,ans,str+"R");
17.         helper(m,row-1,col,n,ans,str+"U");
18.         helper(m,row,col-1,n,ans,str+"L");
19.         m[row][col]=1; //while backtracking we revert back
    the maze to its original postion to explore diffrent paths
20.     }
21. }
```

- time complexity: $4^{(n*m)}$ (because we have to go into all four direction)
- Space complexity : $m*n$ (if we explore all the path for reaching $n-1$ and $m-1$)

Kth-Permutation Sequence : (L18. K-th Permutation Sequence | Leetcode)

```
1. class Solution {
2.     public String getPermutation(int n, int k) {
3.         //n=4, k=10
4.         int fact=1;
5.         ArrayList<Integer> list=new ArrayList<>();
6.         String ans="";
7.         for(int i=1; i<=n; i++){
8.             fact=fact*i;
9.             list.add(i);
10.        }
11.        //fact=24
12. fact=fact/list.size(); //fact=24/4=6 so-->
13. //we are doing this because there are fact/list.size() permutations sequence which will
14. // start from particular number
15. // say n = 4, you have {1, 2, 3, 4}
16. // If you were to list out all the permutations you have
17. // 1 + (permutations of 2, 3, 4) 6 permutation will start with 1
18. // 2 + (permutations of 1, 3, 4) 6 permutation will start with 2
19. // 3 + (permutations of 1, 2, 4) 6 permutation will start with 3
20. // 4 + (permutations of 1, 2, 3) 6 permutation will start with 4
21.
22. k=k-1; //since we are dealing with 0 based indexing k=10-1 => k=9
23. while(true) {
24.     ans+=String.valueOf(list.get(k/fact)); // (9/6=1) in zero based index we
25.     // come to know that out no
26.     // will start from 2 ((line 20) 2 + (permutations of 1, 3, 4) )
27.     // by k/fact we will know what will be our first no, it will tell us the
28.     // sequence in which our no lies so we will get our no
29.     list.remove(k/fact); // simply remove that no so list become 1 3 4
30.     if(list.size()==0) break;
31.     k=k%fact; // the 9%6=3 so the question break down find the 3rd permutation
32.     // sequence of 1 3 4
33.     fact=fact/list.size(); //6=6/3=2
34.     // so k=3 and fact=2 now repeat the above steps until list.size()!=0
35. }
36. return ans;
37. }
38. }
```

- time complexity: $O(n)$ (for looking for n numbers) * $O(n)$ (for everytime removing element from array it will take time so) = $O(n^2)$
- Space complexity : $O(n)$ (using array list and storing ans)