Relational Algebra: Gives an evaluation plan for executing SQL queries

1. Selection

Q: Find the details of employees

SQL: select

*

From emp;

RA: σ (emp)

2. Projection

Q: Find the salary of employees

SQL: select

Empno, salary

From emp;

RA: $\pi_{empno,salary}$ (emp)

3. Selection with a condition

Q: Find the details of employees whose salary is greater than 2000

SQL: select

*

From emp

Where salary>2000;

RA:

σ_{salary>2000} (emp)

4. Selection and projection with a condition

Q: Find the hiredate of employees who are salesmen

SQL: select

Empno, hiredate

From emp

Where job='Salesmen';

RA1: $\pi_{empno, hiredate} (\sigma_{job='Salesmen'} (emp))$

RA2: $\sigma_{job='Salesmen'}$ ($\pi_{empno, hiredate}$ (emp))

5. Cartesian Product

SQL: select

*

From emp, dept

RA: σ (emp X dept)

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6. Join (Join is a subset to Cartesian Product)
  SQL: select
        *
        From emp, dept
        Where emp.deptno=dept.deptno
                        o (emp ⋈ dept)
  RA1:
             \sigma_{\text{emp.deptno=dept.deptno}} ( emp X dept )
  RA2:
7. Join with aliases
  SQL: select
        From emp e, dept d
        Where e.deptno=d.deptno
         \sigma_{e.deptno=d.deptno}(\rho_e(emp) \times \rho_d(dept))
  RA:
8. Join with projection
  SQL: select
        Empno, e.deptno, dname
        From emp e, dept d
        Where e.deptno=d.deptno
        \pi_{empno, e.deptno, dname} (\sigma_{e.deptno=d.deptno} (\rho_{e}(emp) \times \rho_{d}(dept)))
  RA:
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9. Nested Query
  Q: Find the names of departments of managers
  SQL:
  Query a: select
             deptno
             From emp
             Where job='Manager'
                  select
  Query b:
                  dname
                  From dept
                  Where deptno in (query a)
  Query c: substitute query a in query b
  RA:
                 r \leftarrow \pi_{deptno} \left( \sigma_{job='Manager'} \left( emp \right) \right)
                                \pi_{dname} (\sigma_{deptno in r} (dept))
10.
       Left outer join
             select
  SQL:
             *
             From emp e, dept d
             Where e.deptno=d.deptno(+);
                     \sigma_{e.deptno=d.deptno(+)} ( \rho_e(emp) \times \rho_d(dept) )
  RA:
                     \sigma ( \rho_e(emp) \triangleright \rho_d(dept))
```

11. Right outer join

SQL: select

*

From emp e, dept d

Where e.deptno(+)=d.deptno;

RA: $\sigma_{e.deptno(+)=d.deptno}$ ($\rho_e(emp) \times \rho_d(dept)$)

 σ (ρ_e (emp) ρ_d (dept))

12. Full outer join

SQL: select

*

From emp e, dept d

Where e.deptno(+)=d.deptno(+);

RA: $\sigma_{e.deptno(+)=d.deptno(+)}$ ($\rho_e(emp) \times \rho_d(dept)$)

 σ (ρ_e (emp) ρ_d (dept))

13. Group functions

SQL:

Select

Count(*)

From emp;

RA: $\pi_{count(*)}$ (emp)

14. Group function with group by clause SQL:

Select

Deptno, max(salary)

From emp

Group by deptno;

RA: $\pi_{deptno, count(*)}$ (\mathcal{G}_{deptno} (emp))

15. Group function with group by and having clausesQ: Find the deptno with atleast one managerSQL:

Select
Deptno, count(*)
From emp
Where job='Manager'
Group by deptno
Having count(*)>=1

RA: $\pi_{deptno, count(*)}(\sigma_{count(*)}) = 1$ ($G_{deptno}(\sigma_{job='Manager'}(emp)))$)

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Set Operations
16. Union
 SQL:
          select prof1 from programmer
          union
          select prof2 from programmer
       \pi_{prof1} (programmer) \bigcup \pi_{prof2} (programmer)
 RA:
17.
   Intersection
 SQL:
      Select pname from programmer
      Intersect
      Select pname from software
 RA: \pi_{pname} (programmer) \cap \pi_{pname} (software)
18.
      Difference
 Q: Find the programmer who haven't done any project
      SQL:
          Select pname from programmer
          minus
          Select pname from software
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RA: π_{pname} (programmer) - π_{pname} (software)

Update operations

19. Insert

SQL:

Insert into sailor values('S1','Hari','Vskp',7)

RA: sailor \leftarrow sailor \cup ('S1','Hari','Vskp',7)

20. Delete

SQL:

Delete sailor where sid='S10'

RA: sailor \leftarrow sailor $-\sigma_{\text{sid}='S10'}$ (sailor)

21. Update

SQL:

Update sailor set city='Guntur' where sid='S2'

RA: sailor $\leftarrow \pi_{city='Guntur'}$ ($\sigma_{sid='S2'}$ (sailor))

 π , σ , ρ , \bowtie , X, \bowtie , \bowtie , X, \bowtie , \bowtie , \bowtie , \bowtie , \bowtie ,