

## Matrix Decomposition

✓ LU

✓ QR

→ Spectral (square matrices)

SVD

$\underline{A}\underline{x} = \lambda\underline{x}$  ← Eigenvector  $\underline{x}$  is scaled by  $\lambda$   
when premultiplied by  $\underline{A}$

$\underline{A}$  must be square (column & row space must be equal)

To solve analytically, find all  $\lambda$ , such that

$$|\underline{A} - \lambda\underline{I}| = 0 \quad \text{Determinant search}$$

Issue: No closed form solutions for polynomials of size  $\geq 5$

One could do numerical root finding, but that is typically not stable, and one would still need the eigenvectors  $\underline{x} \Rightarrow$  Need iterative solvers for the eigenproblem

Two classes of solvers:

\* 1) Finding largest (or smallest)  $\lambda$

2) Finding the entire spectrum (or a portion of it)

### Largest Eigenvalue

Restrict to real, symmetric A

### Rayleigh Quotient

Let  $\underline{x}$  be an eigenvector of  $\underline{A}$ , then

$$\underline{A}\underline{x} = \lambda \underline{x}$$

and

$$\underline{x}^T \underline{A} \underline{x} = \lambda \underline{x}^T \underline{x} \quad (\text{scalar equation})$$

$$\therefore \lambda = \frac{\underline{x}^T \underline{A} \underline{x}}{\underline{x}^T \underline{x}} \quad \leftarrow \text{Given } \underline{x} \text{ \& } \underline{A}, \text{ find } \lambda$$

### Power Iteration

Let  $\underline{v}_0$  be any vector such that

$\|\underline{v}_0\| = 1$  and  $\underline{v}_0$  is not an eigenvector

Let  $\underline{q}_1, \underline{q}_2, \dots, \underline{q}_n$  be the orthonormal set of eigenvectors,

then

$$\underline{v}_0 = a_1 \underline{q}_1 + a_2 \underline{q}_2 + \dots + a_n \underline{q}_n$$

Consider  $\underline{A} \underline{v}_0$

$$\underline{A} \underline{v}_0 = \underline{A} (a_1 \underline{q}_1 + a_2 \underline{q}_2 + \dots + a_n \underline{q}_n)$$

$$= a_1 \underline{A} \underline{q}_1 + a_2 \underline{A} \underline{q}_2 + \dots + a_n \underline{A} \underline{q}_n$$

$$= a_1 \lambda_1 \underline{q}_1 + a_2 \lambda_2 \underline{q}_2 + \dots + a_n \lambda_n \underline{q}_n$$

$$= \lambda_1 (a_1 \underline{q}_1 + a_2 \frac{\lambda_2}{\lambda_1} \underline{q}_2 + \dots + a_n \frac{\lambda_n}{\lambda_1} \underline{q}_n)$$

$$\underline{A}^2 \underline{v}_0 = \underline{A} (\underline{A} \underline{v}_0) = \underline{A} \lambda_1 (a_1 \underline{q}_1 + \dots + a_n \frac{\lambda_n}{\lambda_1} \underline{q}_n)$$

$$= \lambda_1 (a_1 \underline{A} \underline{q}_1 + a_2 \underline{A} \frac{\lambda_2}{\lambda_1} \underline{q}_2 + \dots + a_n \frac{\lambda_n}{\lambda_1} \underline{A} \underline{q}_n)$$

$$= \lambda_1 (a_1 \lambda_1 \underline{q}_1 + a_2 \frac{\lambda_2^2}{\lambda_1} \underline{q}_2 + \dots + a_n \frac{\lambda_n^2}{\lambda_1} \underline{q}_n)$$

$$= \lambda_1^2 (a_1 \underline{q}_1 + a_2 \left(\frac{\lambda_2}{\lambda_1}\right)^2 \underline{q}_2 + \dots + a_n \left(\frac{\lambda_n}{\lambda_1}\right)^2 \underline{q}_n)$$

$$\lambda_1^2 \dots = \lambda_1^2 / \lambda_1^2 = \dots = (\lambda_2 / \lambda_1)^2 \dots (\lambda_n / \lambda_1)^2$$

$$\underline{A}^p \underline{v}_0 = \lambda_1^p \left( a_1 \underline{q}_1 + a_2 \left( \frac{\lambda_2}{\lambda_1} \right)^p \underline{q}_2 + \dots + a_n \left( \frac{\lambda_n}{\lambda_1} \right)^p \underline{q}_n \right)$$

Let  $|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$  (order eigenvalues)

Then

$$\lim_{p \rightarrow \infty} \left( \frac{\lambda_j}{\lambda_1} \right)^p = 0 \text{ for } j \neq 1$$

$$\Rightarrow \lim_{p \rightarrow \infty} \frac{\underline{A}^p \underline{v}_0}{\lambda_1^p} = a_1 \underline{q}_1 \text{ with } a_1 = \underline{q}_1^T \underline{v}_0$$

Combine this result with the Rayleigh Quotient

Algorithm: Power Algorithm

$\underline{v}_0 \Rightarrow$  Some vector with  $\|\underline{v}_0\| = 1$

for  $k=1, 2, \dots$

$$\underline{w} = \underline{A} \underline{v}_{k-1}$$

$$\underline{v}_k = \underline{w} / \|\underline{w}\|$$

$$\lambda_{(k)} = \underline{v}_k^T \underline{A} \underline{v}_k$$

$$\underline{A}^k \underline{v}_0 = \underline{A} \underline{A} \underline{A} \left( \underline{A} \underline{v}_0 \right)$$

$\underline{v}_2$   
 $\underline{v}_1$

$$\rightarrow \underline{v}_k^T \underline{v}_k = 1$$

$\rightarrow$  Rayleigh quotient

end

This converges at a rate of

$$\| \underline{v}_k - (\pm \underline{q}_1) \| = O \left( \left| \frac{\lambda_2}{\lambda_1} \right|^k \right)$$

$$| \lambda_{(k)} - \lambda_1 | = O \left( \left| \frac{\lambda_2}{\lambda_1} \right|^{2k} \right)$$

This causes an issue if  $\lambda_1 \sim \lambda_2$

In this case, try an inverse iteration w/shift

### Inverse Iteration with Shift

Let  $\mu \in \mathbb{R}$ , such that  $\mu$  is not an eigenvalue of  $\underline{A}$ . Then,  $(\underline{A} - \mu \underline{I})$  has the same eigenvectors as  $\underline{A}$  with eigenvalues  $\lambda_j - \mu = \hat{\lambda}_j$

$$\text{Note: } [(\underline{A} - \mu \underline{I}) - \hat{\lambda} \underline{I}] \underline{q} = \underline{0} \Rightarrow \lambda = \mu + \hat{\lambda}$$

Extension: Eigenvectors of  $(\underline{A} - \mu \underline{I})^{-1}$  are the same as those for  $\underline{A}$ , and the eigenvalues

for  $(\underline{A} - \mu \underline{I})^{-1}$  are  $(\lambda_j - \mu)^{-1}$

Let  $\mu$  be close to  $\lambda_1$ , then  $|\lambda_1 - \mu|^{-1}$  will be much larger than  $|\lambda_j - \mu|^{-1}$  for  $j > 1$

Algorithm: Inverse iteration with shift

Let  $\underline{v}_0$  be some vector with  $\|\underline{v}_0\| = 1$ ,

choose  $\mu > 0$

for  $k = 1, 2, \dots$

Solve  $(\underline{A} - \mu \underline{I}) \underline{w} = \underline{v}_{k-1}$  for  $\underline{w}$

$\underline{v}_k = \underline{w} / \|\underline{w}\|$       Normalize  $\underline{v}_k$

$\lambda_{(k)} = \underline{v}_k^T \underline{A} \underline{v}_k$       Rayleigh Quotient

end

Convergence order of

$$\|\underline{v}_k - (\pm \underline{q}_1)\| = \mathcal{O}\left(\left|\frac{\mu - \lambda_1}{\mu - \lambda_2}\right|^k\right)$$

$$|\lambda_{(k)} - \lambda_1| = \mathcal{O}\left(\left|\frac{\mu - \lambda_1}{\mu - \lambda_2}\right|^{2k}\right)$$

Now combine above ideas

### Algorithm: Rayleigh Quotient Iteration

$\underline{v}_0$  is some vector w/  $\|\underline{v}_0\| = 1$

$$\lambda_{(0)} = \underline{v}_0^T \underline{A} \underline{v}_0$$



Do not need

to select a shift  $\mu$

for  $k=1, 2, \dots$

Solve  $(\underline{A} - \lambda_{(k-1)} \underline{I}) \underline{w} = \underline{v}_{k-1}$  for  $\underline{w}$

$$\underline{v}_k = \underline{w} / \|\underline{w}\|$$

$$\lambda_{(k)} = \underline{v}_k^T \underline{A} \underline{v}_k$$

end

This method has a convergence of

$$\|\underline{v}_{k+1} - (\pm \underline{q}_j)\| = \mathcal{O}(\|\underline{v}_k - (\pm \underline{q}_j)\|^3)$$

$$|\lambda_{(k+1)} - \lambda_j| = \mathcal{O}(|\lambda_{(k)} - \lambda_j|^3)$$



Cubic order convergence of the eigenvector  $\underline{q}_j$   
closest to  $\underline{v}_0$

See Lecture 27 of Trefethan & Bau

$\underline{A}\underline{x} = \lambda\underline{x}$   $\leftarrow$  Eigenvector  $\underline{x}$  is scaled by  $\lambda$   
when premultiplied by  $\underline{A}$

Two classes of solvers:

- 1) Finding largest (or smallest)  $\lambda$
- \* 2) Finding the entire spectrum (or a portion of it)

### Spectrum Calculations

Try to find all or a subset of the eigenvalue spectrum

Recall that any square matrix has the  
Schur Decomposition

$$\underline{A} = \underline{Q} \underline{T} \underline{Q}^T, \text{ where } \underline{T} \text{ is upper triangular}$$

Eigenvalue computations can try to find  
this decomposition, in which  $\underline{A}$  &  $\underline{T}$  are similar

this decomposition, in which  $\underline{A}$  &  $\underline{I}$  are similar

Recall, eigenvalues of  $\underline{A}$  appear on the diagonal of  $\underline{I}$

Note: If A is symmetric and real, then

$$\underline{A} = \underline{Q} \underline{I} \underline{Q}^T = \underline{S} \underline{\Lambda} \underline{S}^{-1}$$

↑  
diagonal

The above looks similar to QR decomposition,

where  $\underline{A} = \underline{Q} \underline{R}$   $\nwarrow$  upper triangular

## Recall Householder reflections

$$\begin{bmatrix} x & x & x \\ x & x & x \\ x & x & x \end{bmatrix} \xrightarrow{Q_1^T} \begin{bmatrix} x & x & x \\ 0 & x & x \\ 0 & x & x \end{bmatrix}$$

For the eigenproblem, we need  $\underline{Q}_1^T \underline{A} \underline{Q}_1$

Then

$$\underbrace{Q_n^T Q_{n-1}^T \dots Q_1^T}_{\underline{Q}^T} \underline{A} \underbrace{Q_1 Q_2 \dots Q_n}_{\underline{Q}} = \underline{I}$$

and then  $\underline{A} = \underline{Q} \underline{I} \underline{Q}^T$

Consider  $\underline{Q}_1^T \underline{A} \underline{Q}_1$

$$\begin{bmatrix} x & x & x \\ 0 & x & x \\ 0 & x & x \end{bmatrix} \rightarrow \begin{bmatrix} x & x & x \\ x & x & x \\ x & x & x \end{bmatrix}$$

$\underline{Q}_1^T \underline{A}$ 
 $\underline{Q}_1^T \underline{A} \underline{Q}_1$

Fill-in of  
the zeros

$\therefore$  The original Householder approach  
will not work

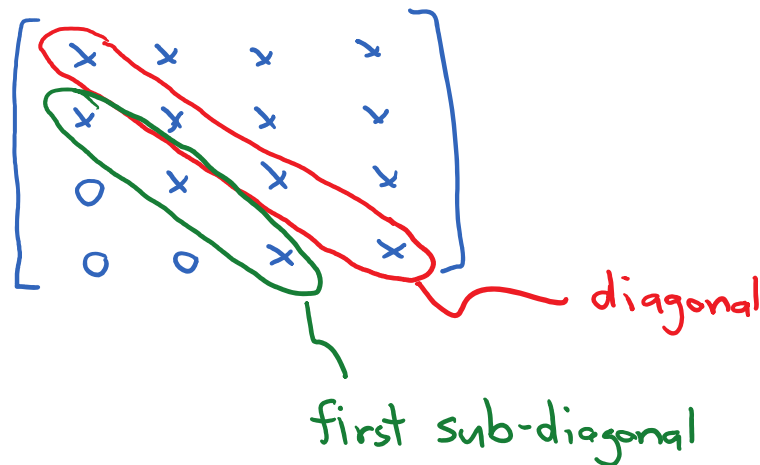
$\Rightarrow$  Not possible to get a Schur Decomposition  
directly

Instead, two steps are needed:

- 1) Reduce to upper Hessenberg form, which is nearly upper triangular
- 2) Iterate until upper triangular is obtained

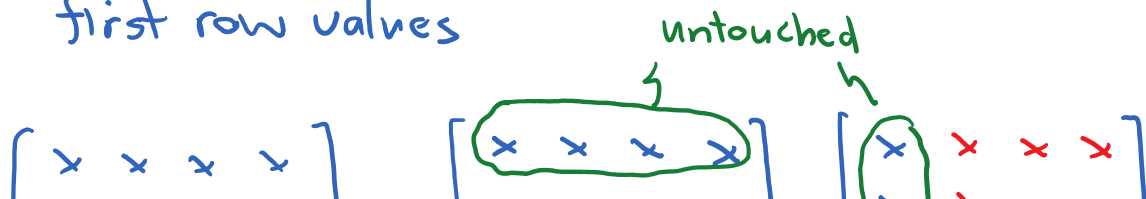
Details of these two steps:

- 1) Upper Hessenberg Matrix: A matrix with zeros below the first sub-diagonal



Let  $\underline{Q}_1^T$  be a unitary matrix ( $\underline{Q}_1^T \underline{Q}_1 = \underline{I}$ )

that zeros out values below the first subdiagonal of the first column, but does not touch the first row values



$$\begin{array}{ccc}
 \begin{bmatrix} \times & \times & \times & \times \\ \times & \times & \times & \times \\ \times & \times & \times & \times \\ \times & \times & \times & \times \end{bmatrix} & \rightarrow & \begin{bmatrix} \times & \times & \times & \times \\ \times & \times & \times & \times \\ 0 & \times & \times & \times \\ 0 & \times & \times & \times \end{bmatrix} \rightarrow \begin{bmatrix} \times & \times & \times & \times \\ \times & \times & \times & \times \\ 0 & \times & \times & \times \\ 0 & \times & \times & \times \end{bmatrix} \\
 \underline{A} & & \underline{Q}_1^T \underline{A} \quad \underline{Q}_1^T \underline{A} \underline{Q}_1
 \end{array}$$

Use a Householder reflector to assure ortho-normality

$$\underline{Q} = \begin{bmatrix} \underline{I} & \underline{0} \\ \underline{0} & \underline{F} \end{bmatrix}$$

Algorithm: Householder Reduction to Upper Hessenberg Form

for  $k = 1$  to  $m-2$  with  $\underline{A} \in \underline{M}_{mm}$

$$\underline{x} = \underline{A}(k+1:m, k)$$

$$\underline{v}_k = \text{sign}(x_1) \|\underline{x}\|_2 \underline{e}_1 + \underline{x}$$

$$\underline{A}(k+1:m, k:m) = \underline{A}(k+1:m, k:m)$$

$\swarrow \underline{Q}^T \underline{A}$

$$-2 \underline{v}_k (\underline{v}_k^T \underline{A}(k+1:m, k:m))$$

$$\underline{A}(1:m, k+1:m) = \underline{A}(1:m, k+1:m)$$

$$-2 (\underline{A}(1:m, k+1:m) \underline{v}_k) \underline{v}_k^T$$

end

$$\hookrightarrow \underline{Q}^T \underline{A} \underline{Q}$$

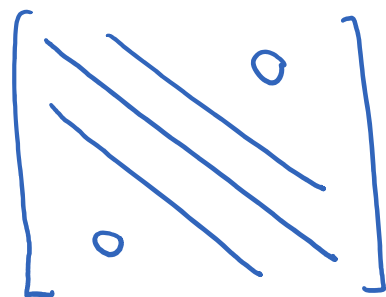
$\Rightarrow \underline{A}$  then converts to upper Hessenberg

Note:  $\underline{Q}$  is never formed

$$\text{Cost: } \mathcal{O}\left(\frac{10}{3} m^3\right)$$

If  $\underline{A}$  is symmetric, then the cost reduces to  $\mathcal{O}\left(\frac{4}{3} m^3\right)$

and the result is tri-diagonal



Why?

2) Iterate to Upper Triangular Form

Focus on real, symmetric matrices

Turn to  $\underline{A} = \underline{Q} \underline{T} \underline{Q}^T$  (Schur decomposition)

$\underline{A}$  could be any matrix or the result of part 1) (i.e., upper Hessenberg)

$$\underline{T} = \underline{Q}^T \underline{A} \underline{Q}$$

Make this an iteration

Given  $\underline{A}_k$ , let  $\underline{A}_{k+1} = \underline{Q}_k^T \underline{A}_k \underline{Q}_k$

Now, let  $\underline{A}_k = \underline{Q}_k \underline{R}_k$  be the QR decomposition of  $\underline{A}_k$

Then,

$$\begin{aligned} \underline{A}_{k+1} &= \underline{Q}_k^T \underline{A}_k \underline{Q}_k = \underline{Q}_k^T (\underline{Q}_k \underline{R}_k) \underline{Q}_k \\ &= \underline{I} \underline{R}_k \underline{Q}_k = \underline{R}_k \underline{Q}_k \end{aligned}$$

Given  $\underline{A}_k$ , find  $\underline{Q}_k \underline{R}_k$ , then  $\underline{A}_{k+1} = \underline{R}_k \underline{Q}_k$



⇒ This is the QR Algorithm for eigenproblems

Algorithm: QR for Eigenproblems

$$\text{Let } \underline{A}_0 = \underline{A}$$

for  $k=1, 2, \dots$

$$\underline{Q}_k \underline{R}_k = \underline{A}_{k-1} \quad (\text{QR of } \underline{A}_{k-1})$$

$$\underline{A}_k = \underline{R}_k \underline{Q}_k \quad \text{Recombination in reverse}$$

end

Converge to some tolerance,

result will be upper triangular matrix  $\underline{I}$

$$\underline{A} \underline{x} = \lambda \underline{x} \quad \text{for square } \underline{A}$$

$\lambda$ : eigenvalue

$\underline{x}$ : eigenvector

$$(\underline{A} - \lambda \underline{I}) \underline{x} = \underline{0}$$

For nontrivial solutions  $\det(\underline{A} - \lambda \underline{I}) = 0$

Solution approaches:

- Characteristic equation  $\rightarrow$  find roots  $\lambda$
- Iteration methods

Power iteration method (largest  $|\lambda|$ )

Inverse power iteration method (smallest  $|\lambda|$ )

Inverse iteration with shifts

## Rayleigh quotient iteration

- Spectrum calculations

### Schur decomposition

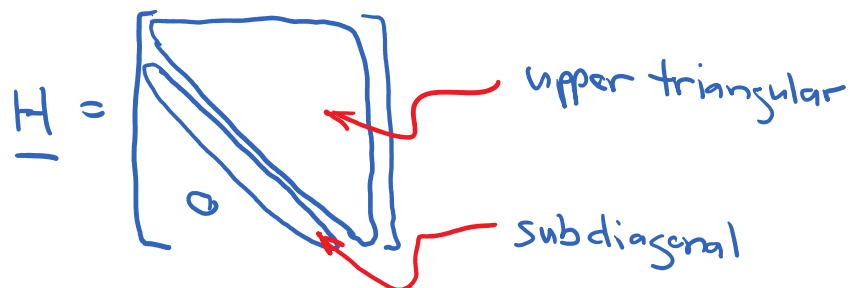
$$\underline{A} = \underline{Q} \underline{T} \underline{Q}^T \quad \text{with } \underline{T}: \text{Upper triangular} \\ \underline{Q}: \text{Unitary } (\underline{Q}^T \underline{Q} = \underline{I})$$

$\underline{A}$  &  $\underline{T}$  are similar & consequently have same eigenvalues

Diagonals of  $\underline{T}$  are its eigenvalues

Multi-step process to obtain eigenvalues:

- ① Compute upper Hessenberg form (upper triangle plus subdiagonal) by using Householder reflections  $\rightarrow \underline{H} = \underline{Q}_1^T \underline{A} \underline{Q}_1$



- ② Iterate from upper Hessenberg form to

upper triangular  $\underline{T} = \underline{Q}_2^T \underline{H} \underline{Q}_2$

③ Eigenvalues of  $\underline{A}$  are on diagonal of  $\underline{T}$

Other approaches (beyond current scope)

- Givens rotations replacing Householder reflections (both have unitary  $\underline{Q}$ )
- Subspace iteration

Lanczos } Krylov subspaces  
Arnoldi }