$$\frac{3}{3} + \frac{3}{5}$$
 $\frac{2}{5}$
 $\frac{2}{5}$
 $\frac{2}{5}$
 $\frac{2}{5}$
 $\frac{2}{5}$
 $\frac{2}{5}$

Let e be the root node.

$$y_e = 0$$
.

$$ye-y_c=3 \Rightarrow y_c=-3$$

$$y_e - y_b = 7 \Rightarrow y_b = -7$$

$$y_b - y_a = -4 \Rightarrow y_a = -3$$
.

so the dual slades are

$$2ac = -3 + 8 - (-3) = 8$$

$$z_{bc} = -7 + 6 - (-3) = 2$$

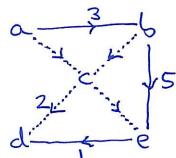
$$\frac{3}{2}$$
cd = $-3+6-13=-10$

$$z_{da} = 13 + 1 - (-3) = 17$$

Zed is infeasible so the current basis is not optimal. Let (c,d) enter the basis. This forms the cycle

so t=2 and (c,e) leaves the basis.

The new basis is



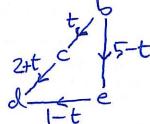
with those arcs whose dual slacks require updating indicated.

$$\tilde{Z}_{bc} = 2 + (-10) = -8$$

$$\tilde{z}_{ce} = 0 - (-10) = 10$$

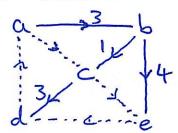
For the next iteration, choose arc (b,c) to enter the basis.

This forms the cycle



so t=1 and are (e,d) leaves the basis.

The new solution is



The new dual slacks are $\tilde{z}_{ac} = -2 - (-8) = 6$

$$\hat{Z}_{da} = 17 + (-8) = 9$$
 $\hat{Z}_{ed} = 0 - (-8) = 8$

So the current solution is dual and primal feasible, and hence optimal.

The optimal value is $3 \cdot (-4) + 1 \cdot 6 + 4 \cdot 7 + 3 \cdot 6 = 40$.

- 2. (a.) F. There are only finitely many basic solutions but as some as there are 2 optimal solutions there are infinitely many because any point on the edge connecting these two solutions will also be optimal.
 - (b.) F. For example, maximise -z, subject to x,≤0.
 - (c.) F. The example on p.114 shows that if c, is reduced by more than /z the solution becomes nonoptimal.
 - (d.) F. We need the yi's at the outset in order to compute the dual slack variables. Thereafter we do not need them.
 - (e.) F. We may include a new demand node to represent the surplus, with an arc running to it from each supply node with unit cost of O to represent leaving alone any material which flows along that arc in the solution.