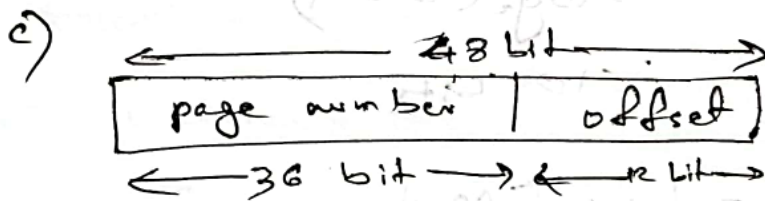


01: (a)

$$\begin{aligned}
 \text{offset size} &= \log_2 (\text{page size in bytes}) \\
 &= \log_2 (4 \text{ KB}) \\
 &= \log_2 (4 \times 1024) \\
 &= 12 \text{ bits}
 \end{aligned}$$

b) number of virtual pages =  $\frac{\text{Virtual memory size}}{\text{page size}}$

$$\begin{aligned}
 &= \frac{2^{48} \text{ bytes}}{2^{12} \text{ bytes}} \\
 &= 2^{36} \text{ bytes}
 \end{aligned}$$

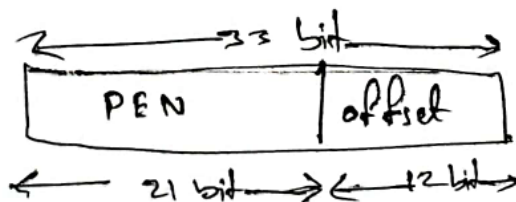


bits for VPN =  $48 - 12 = 36 \text{ bit}$

d) Physical address =  $\log_2 (\text{Physical memory size})$

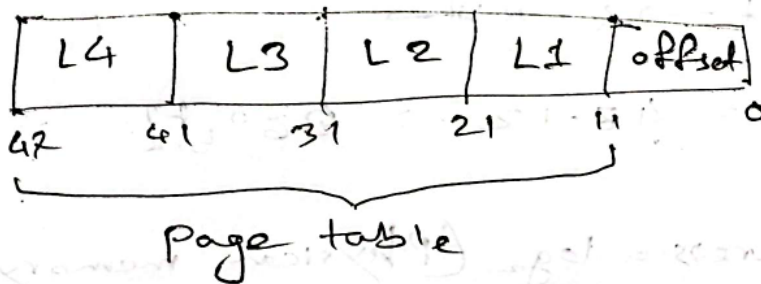
$$\begin{aligned}
 &= \log_2 (8 \text{ GB}) \\
 &= \log_2 (2^3 \times 2^{30}) \\
 &= 33 \text{ bits}
 \end{aligned}$$

Number of bits for PEN =  $33 - 12 = 21 \text{ bits}$



e) size of single page table = No of PTEs  $\times$  PTE size  
 $= \text{No of virtual pages} \times 4 \text{ bytes}$   
 $= 2^{36} \times 2^2 \text{ bytes}$   
 $= 2^8 \times 2^{30} \text{ bytes}$   
 $= 256 \text{ GB}$

f) Number of bits for  
each level page table =  $\log_2 \left( \frac{\text{Page size}}{\text{PTE size}} \right)$   
 $= \log_2 \left( \frac{4 \times 1024 \text{ bytes}}{4 \text{ bytes}} \right)$   
 $= \log_2 (1024)$   
 $= 10 \text{ bits}$



2. Memory required for =  $1024 \times \text{Memory per process}$

$$= 1024 \times 4 \text{ MB} = 4 \text{ GB}$$

Memory per process = 4 MB

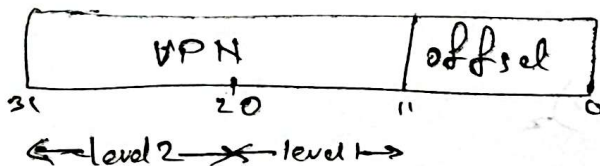
$$\text{No. of pages} = \frac{\text{Virtual Address}}{\text{Page size}}$$

$$= \frac{2^{32}}{2^{12}} = 2^{20}$$

$$\text{bits for VPN} = \log_2(2^{20})$$

$$\text{Physical Address} = \log_2(\text{Physical memory size})$$

$$= 33 \text{ bit}$$



$$\text{Bits for PPN} = 33 - 12$$

$$= 21 \text{ bits}$$

$$L2 = \text{entries} \times \text{entry size}$$

$$= 2^{10} \times 4 \text{ bytes}$$

$$= 4 \text{ KB}$$

$$\text{PTE size} = 21 + 10 = 31 \text{ bits}$$

$$\Rightarrow 4 \text{ bytes (rounded up)}$$

$$\text{No. of PTE} = \frac{\text{Page size}}{\text{Entry size}}$$

$$L_1 = 2^{10} \times 4 \text{ KB}$$

$$= 4 \text{ MB}$$

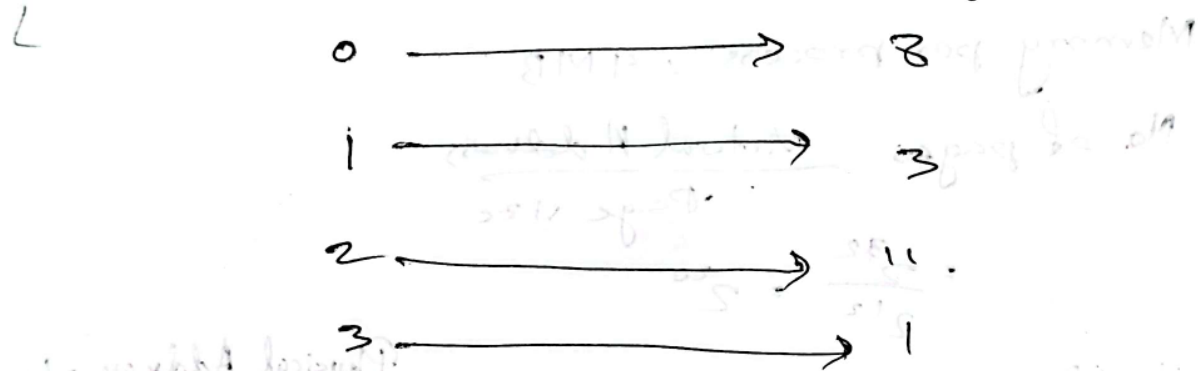
$$= \frac{2^{12}}{4} = 2^{10}$$

$$\text{Memory per process} = L_1 + L_2$$

$$\approx 4 \text{ MB}$$

3. (a) given

Virtual Page Number      Physical Frame



So,

Virtual address =  $(20)_{10}$

=  $(010100)_2$

VPN =  $(01)_2$   
=  $(1)_{10}$

offset =  $(0100)_2$   
=  $(4)_{10}$

Physical address =  $\text{PFN} \times \text{Page space} + \text{offset}$   
=  $3 \times 16 + 4 = 52 \quad (\text{Ans})$

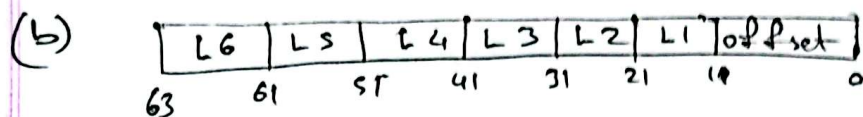
(b) VA =  $(40)_{10} = (101000)_2$

VPN =  $(10)_2 = (2)_{10}$

offset =  $(1000)_2 = (8)_{10}$

Physical Address =  $11 \times 16 + 8 = 184 \quad (\text{Ans})$

4. (a) 6 levels



(c) Maximum number of pages =  $2^{44} + 2^{34} + 2^{24} + 2^{14} + 2^{4+1}$

5. Virtual Address Space: 24 bits

Physical Address Space: 32 bits

Page Size: 4 KB ( $2^{12}$  bytes)  $\rightarrow$  12 bits

VPN =  $24 - 12 = 12$  bits of virtual address,

a) 0000 0000 0011 0000 1010 1111

VPN: 0000 0000 0011 =  $0x0003 \rightarrow 0x032A0$  (valid)

offset: 0000 1010 1111 =  $0x0AF$

Physical address:  $0x032A00AF$

b) 0000 0000 0101 0000 0000 0101

VPN: 0000 0000 0101 =  $0x0005 \rightarrow 0x0121A$  (Not Valid)

offset: 0000 0000 0101 =  $0x0005$

Not Valid

c) 0000 0000 0111 0001 1111 0111

VPN: 0000 0000 0111 =  $0x0007 \rightarrow 0x000AB$  (Valid)

offset: 0001 1111 0111 =  $0x1F7$

Physical address:  $0x000AB1F7$