

Q1.

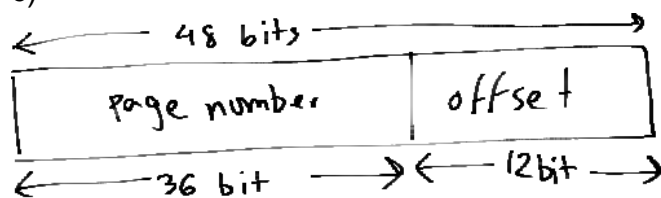
a)

$$\begin{aligned}\text{offset size} &= \log_2(\text{page size in bytes}) \\ &= \log_2(4\text{KB}) \\ &= \log_2(4 \times 1024) \\ &= 12 \text{ bits}\end{aligned}$$

b)

$$\begin{aligned}\text{number of virtual pages} &= \frac{\text{virtual memory size}}{\text{page size}} \\ &= \frac{2^{48} \text{ bytes}}{2^{12} \text{ bytes}} \\ &= 2^{36} \text{ virtual pages}\end{aligned}$$

c)

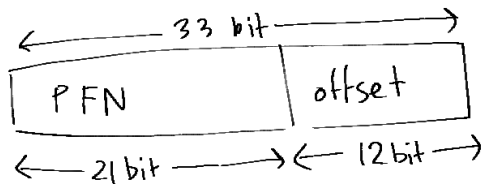


$$\begin{aligned}\text{bits for VPN} &= 48 - 12 \\ &= 36 \text{ bits}\end{aligned}$$

d)

$$\begin{aligned}
 \text{Physical address} &= \log_2 (\text{Physical memory size}) \\
 &= \log_2 (8 \text{ GB}) \\
 &= \log_2 (2^3 \times 2^{30}) \\
 &= 33 \text{ bits}
 \end{aligned}$$

$$\begin{aligned}
 \text{Number of bits for PFN} &= 32 - 12 \\
 &= 21 \text{ bits}
 \end{aligned}$$

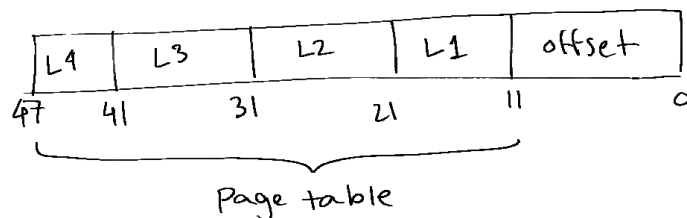


e)

$$\begin{aligned}
 \text{size of single page table} &= \text{No. of PTEs} \times \text{PTE size} \\
 &= \text{No. of virtual pages} \times 4 \text{ bytes} \\
 &= 2^{36} \times 2^2 \text{ bytes} \\
 &= 2^8 \times 2^{30} \text{ bytes} \\
 &= 256 \text{ GB}
 \end{aligned}$$

f)

$$\begin{aligned}
 \text{number of bits for each level page table} &= \log_2 \left(\frac{\text{page size}}{\text{PTE size}} \right) \\
 &= \log_2 \left(\frac{4 \times 1024 \text{ bytes}}{4 \text{ bytes}} \right) \\
 &= \log_2 (1024) \\
 &= 10 \text{ bits}
 \end{aligned}$$



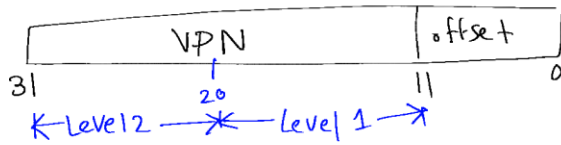
Q2.

$$\begin{aligned}\text{Memory required for all processes} &= 1024 \times \text{Memory per process} \\ &= 1024 \times 4\text{MB} = 4\text{GB}\end{aligned}$$

$$\text{Memory per process} = 4\text{MB}$$

$$\begin{aligned}\text{No. of pages} &= \frac{\text{Virtual Address}}{\text{Page Size}} \\ &= \frac{2^{32}}{2^{12}} = 2^{20}\end{aligned}$$

$$\begin{aligned}\text{bits for VPN} &= \log_2(2^{20}) \\ &= 20 \text{ bits}\end{aligned}$$



$$\begin{aligned}L2 &= \text{entries} \times \text{entry size} \\ &= 2^{10} \times 4 \text{ bytes} \\ &= 4\text{KB}\end{aligned}$$

$$\begin{aligned}L1 &= 2^{10} \times 4\text{KB} \\ &= 4\text{MB}\end{aligned}$$

$$\text{Memory per process} = L1 + L2 \approx 4\text{MB}$$

$$\begin{aligned}\text{Physical Address} &= \log_2(\text{Physical memory size}) \\ &= \log_2(2^{33}) \\ &= 33 \text{ bit}\end{aligned}$$

$$\begin{aligned}\text{bits for PFN} &= 33 - 12 \\ &= 21 \text{ bits}\end{aligned}$$

$$\begin{aligned}\text{PTE size} &= 21 + 10 = 31 \text{ bits} \\ &\Rightarrow 4 \text{ bytes (rounded up)}\end{aligned}$$

$$\begin{aligned}\text{No. of PTE} &= \frac{\text{Page size}}{\text{Entry size}} \\ &= \frac{2^{12}}{4} = 2^{10}\end{aligned}$$

Q3

6 bit virtual address space

16 byte pages/frames

Mapping

VPN	PFN
0	8
1	3
2	11
3	1

Translate virtual to physical (numbers in decimal)

a) 20

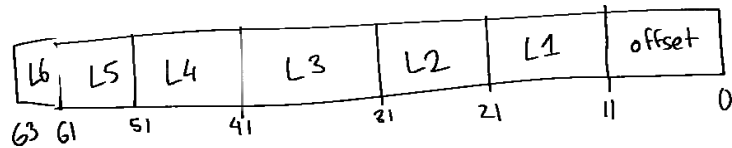
b) 40

don't know how to do this

Q4

a) 6 levels

b)



c)

Maximum number of pages = $2^{44} + 2^{34} + 2^{24} + 2^{14} + 2^4 + 1$

Q5.

Virtual Address Space: 24 bits

Physical Address Space: 32 bits

Page Size: 4 KB (2^{12} bytes) -> 12 bits

$VPN = 24 - 12 = 12$ bits of virtual address

a) 0000 0000 0011 0000 1010 1111

VPN: 0000 0000 0011 = 0x003 -> 0x032A0 (valid)

Offset: 0000 1010 1111 = 0x0AF

Physical address: 0x032A00AF

b) 0000 0000 0101 0000 0000 0101

VPN: 0000 0000 0101 = 0x005 -> 0x0121A (not valid)

Offset: 0000 0000 0101 = 0x005

Not valid.

c) 0000 0000 0111 0001 1111 0111

VPN: 0000 0000 0111 = 0x007 -> 0x000AB (valid)

Offset: 0001 1111 0111 = 0x1F7

Physical address: 0x000AB1F7