

Trees

Arrays

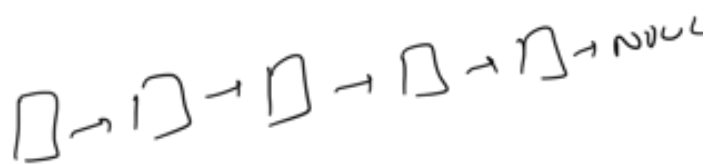
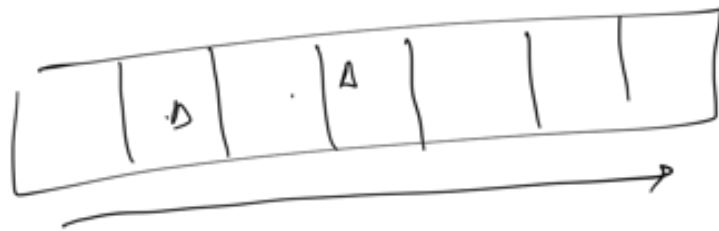
LL

Queue

Stack

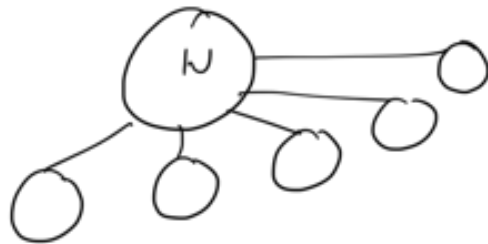
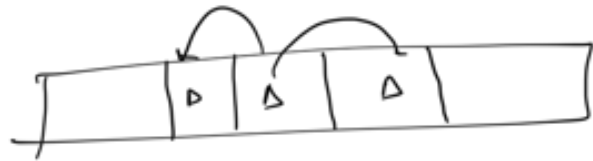
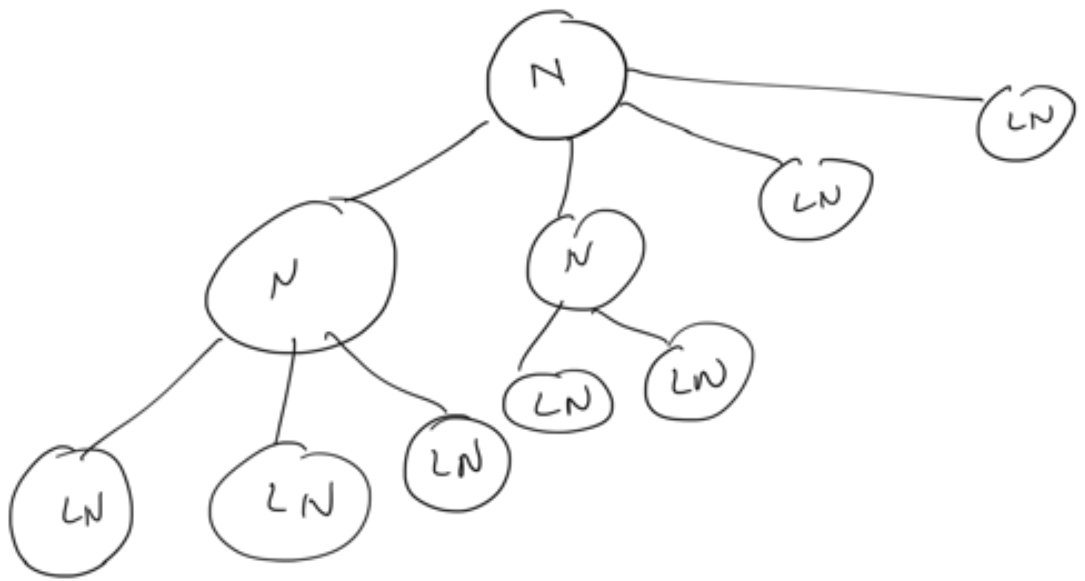
Hash Map

Hash Set



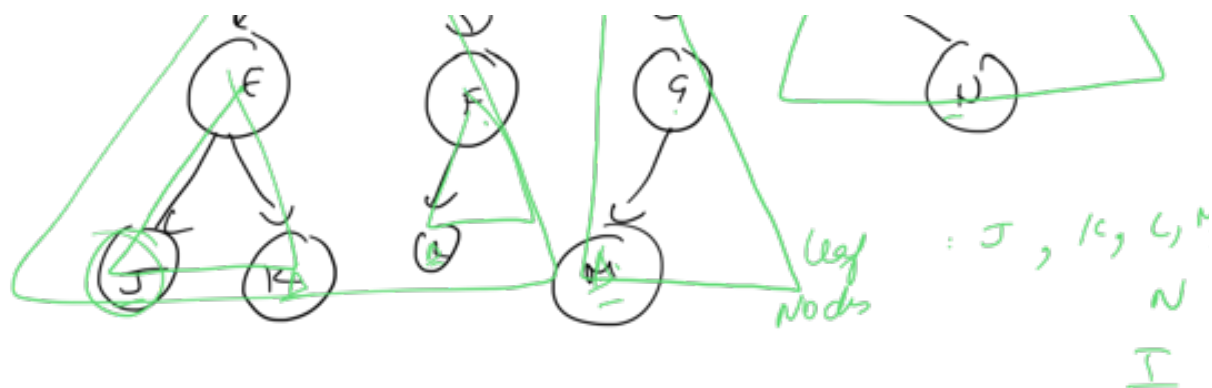
Hierarchical DS

Tree



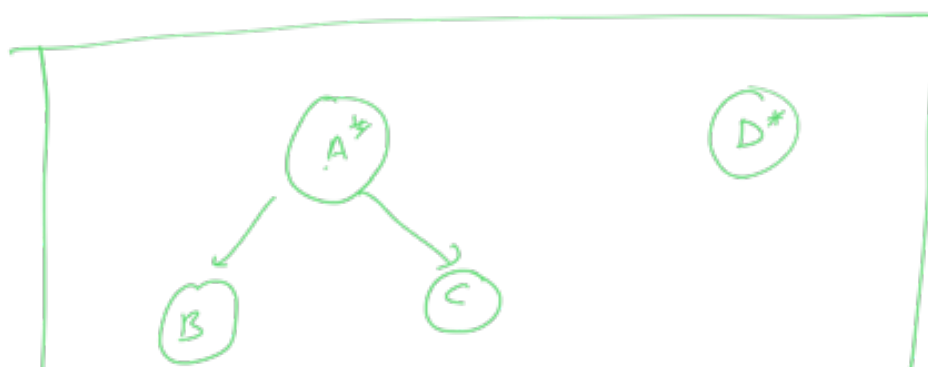
Root $\equiv A$





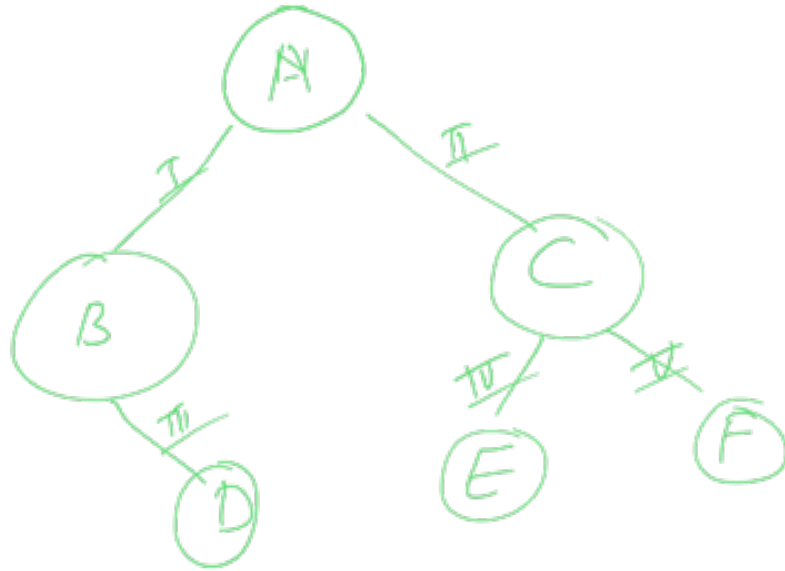
Tree \equiv Recursive

Tree : a collection of nodes, N
 in which all of the nodes are connected
 and $\# \text{ edges} = N - 1$



5 Nodes

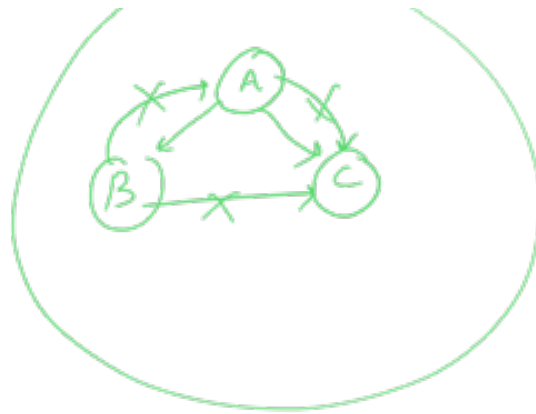
5 edge



Tree \equiv collection of Nodes

A'

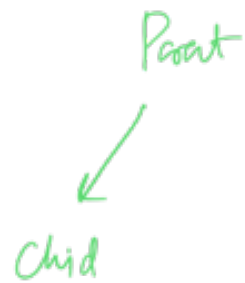
N nodes



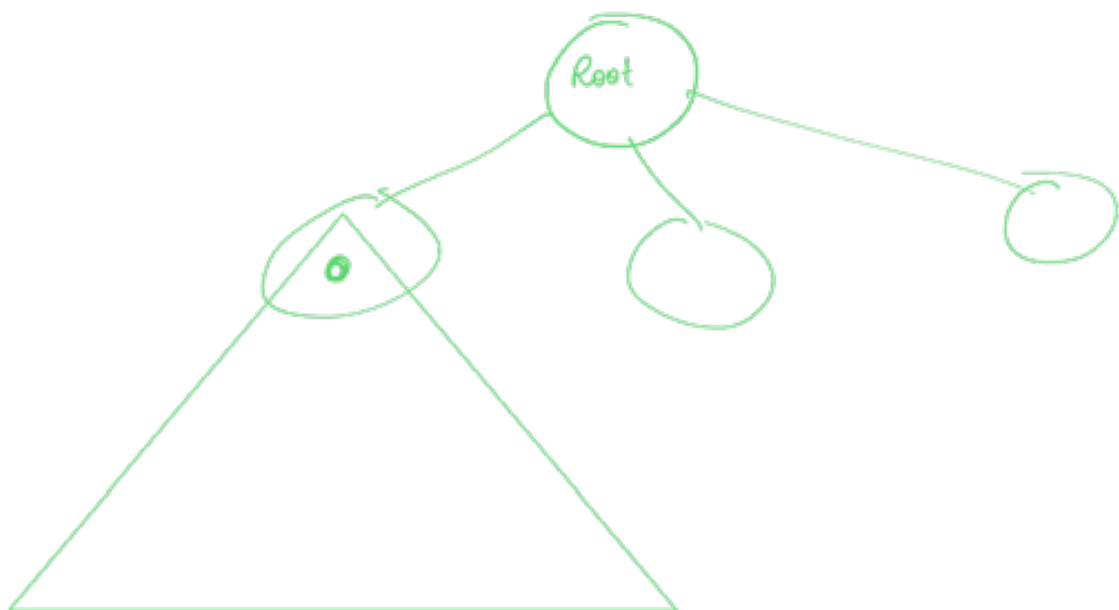
0 edge 1st ans
1 edge 2nd ans
1 edge 3rd

N nodes

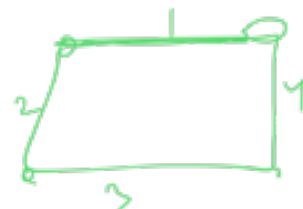
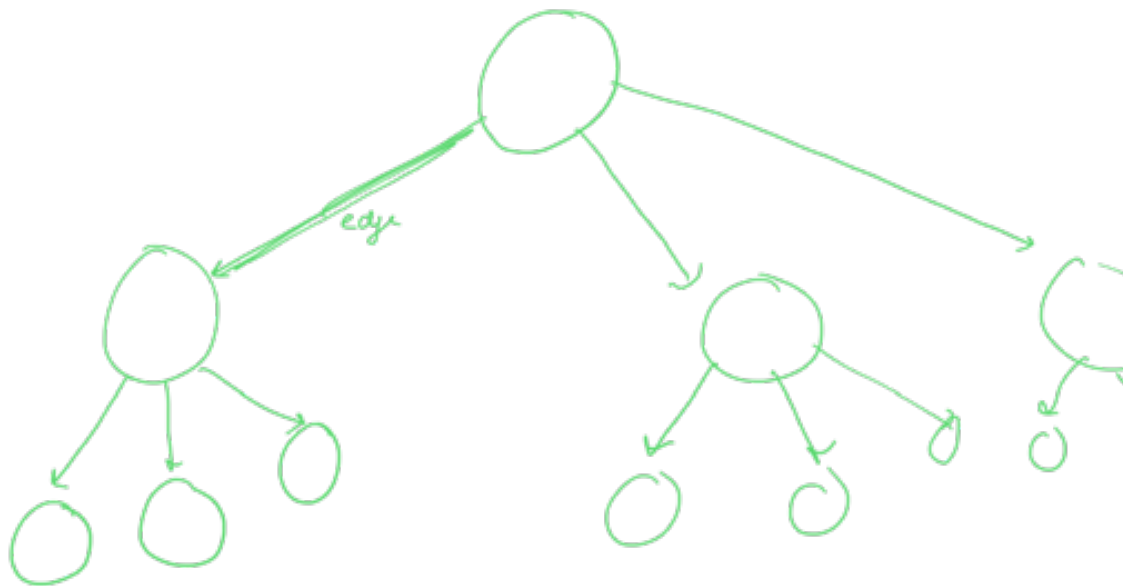
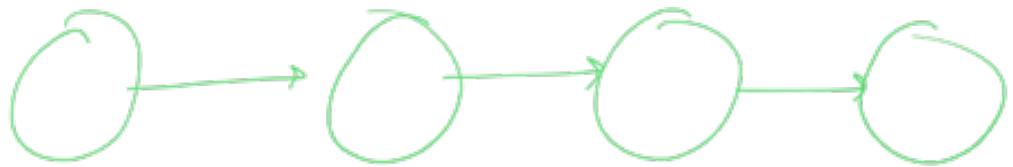
atls $N-1$ edges



N nodes $\equiv N-1$ edges



L.L

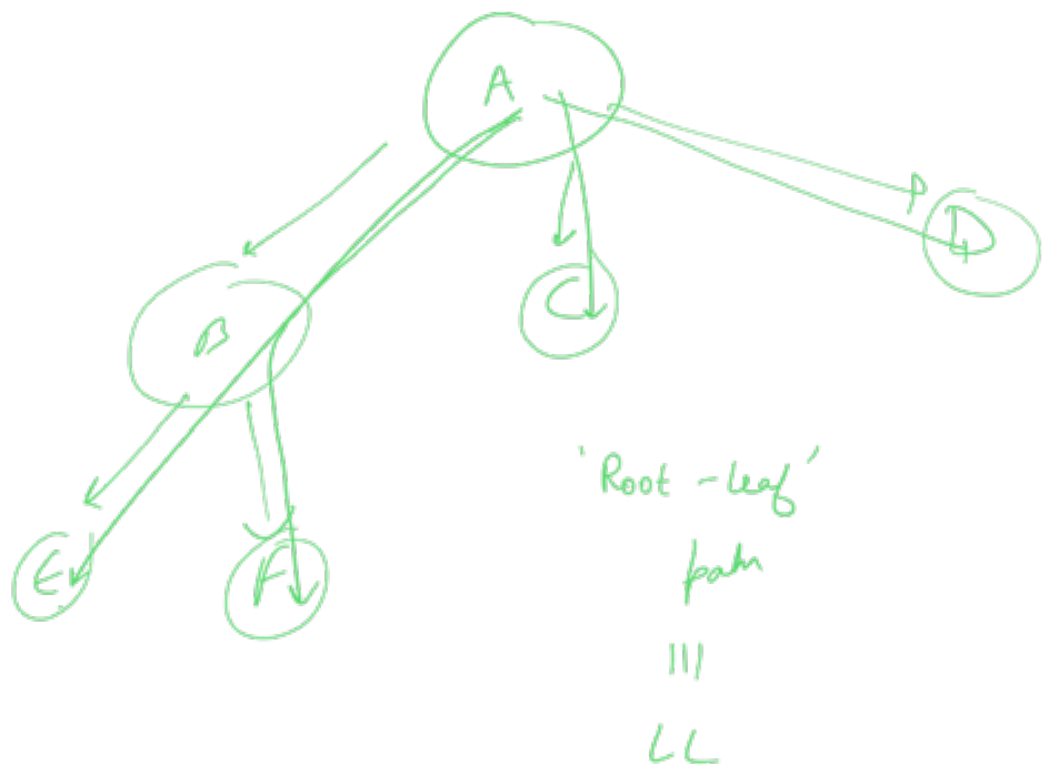
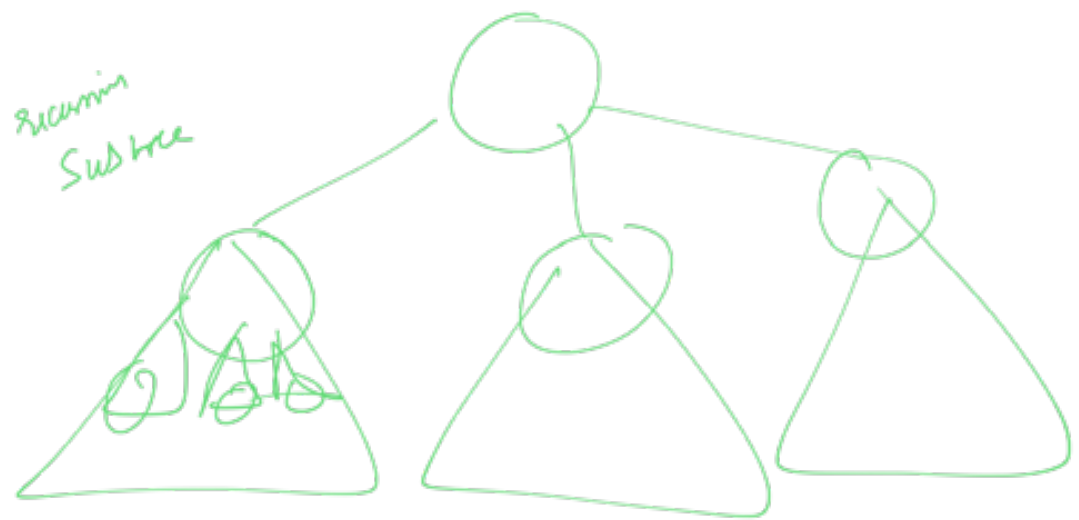


Tree

nodes \equiv data points

Edges

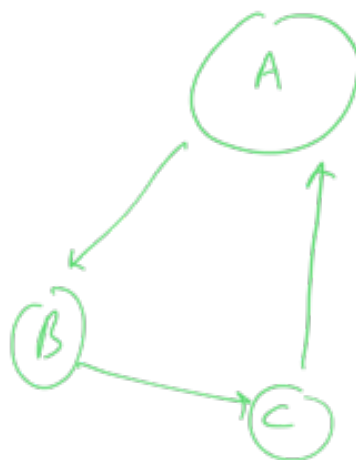
Structure



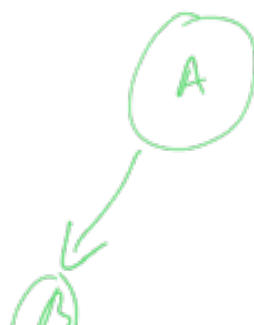


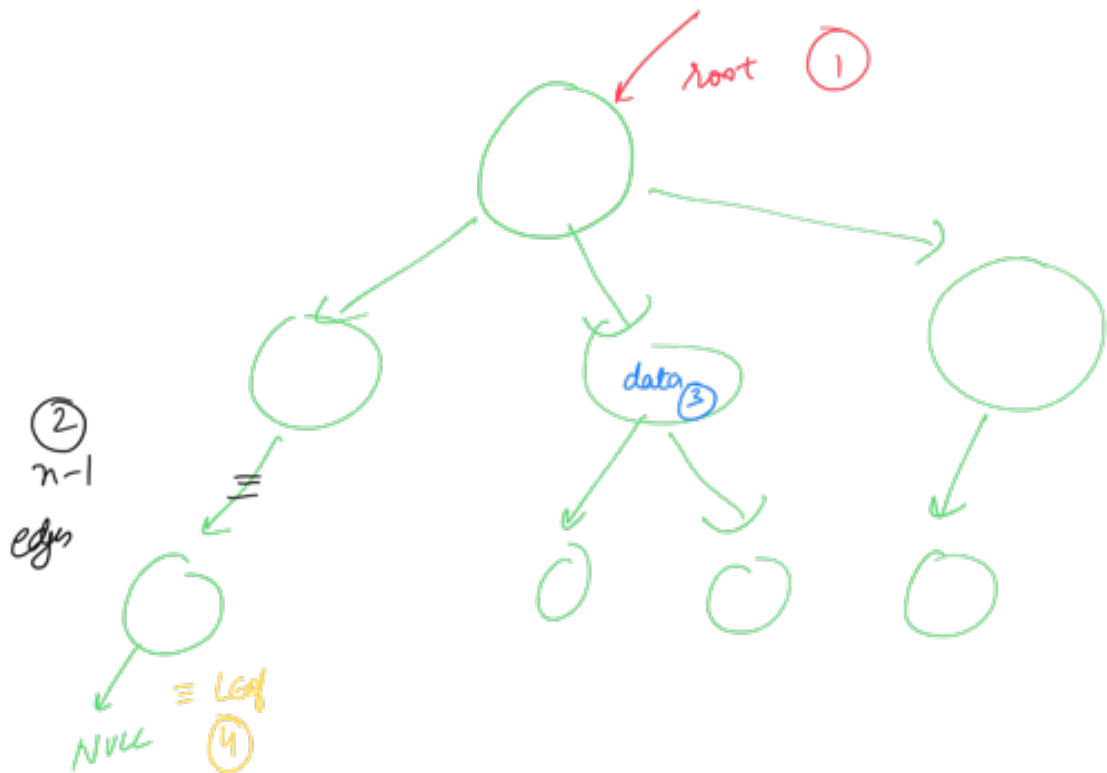
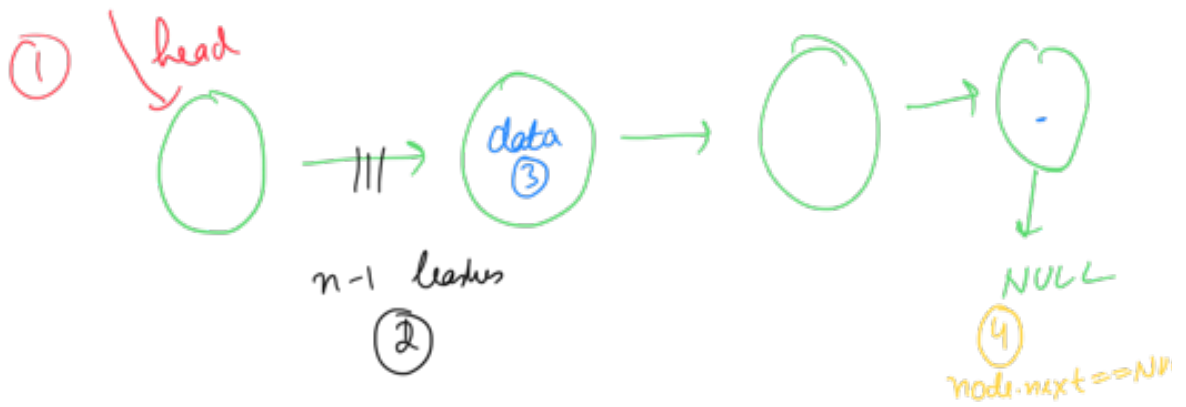
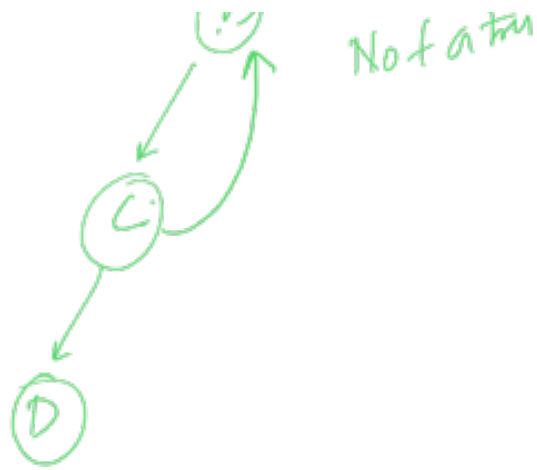
NOT A TREE

every node must have 1 parent
Or has
for
root



NOT a tree
loop X
N edges X



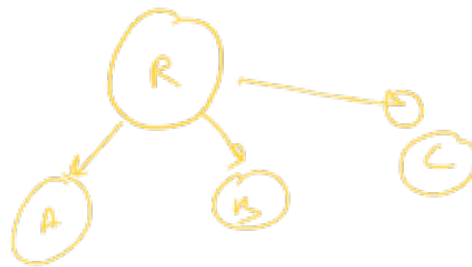


Tree is basically hierarchical L.C



L.C is basically a special
type of a Tree

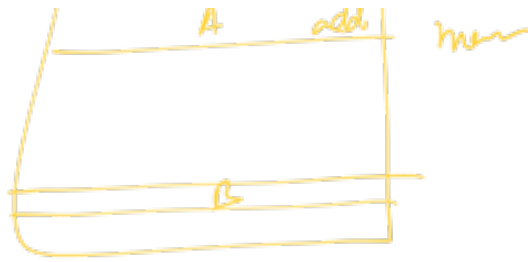
\equiv Tree is a more generic d.s
as compared to a L.C



Path root to Leaf \equiv L.C 😊

L.C is a Tree with any node \neq 1 child





```

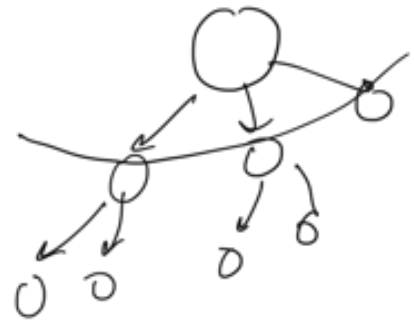
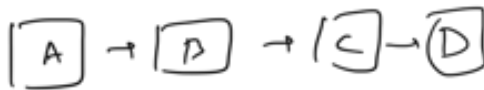
class LinkedListNode
{
    int data ✓
    LinkedListNode next
}

```

```

class TreeNode
{
    int data; ✓
    TreeNode [ ] children;
}

```



Binary Tree

k-ary Tree : max number of children one node
can have
= k

n-ary Tree with $n = 15$

Binary Tree
 $k = 2$

every node in the binary tree will have at most 2 children 😊

✓
0 child or 1 child or 2 children

class TreeNode

{

int data;

TreeNode[] children;

}

class BinaryTreeNode {
int data;

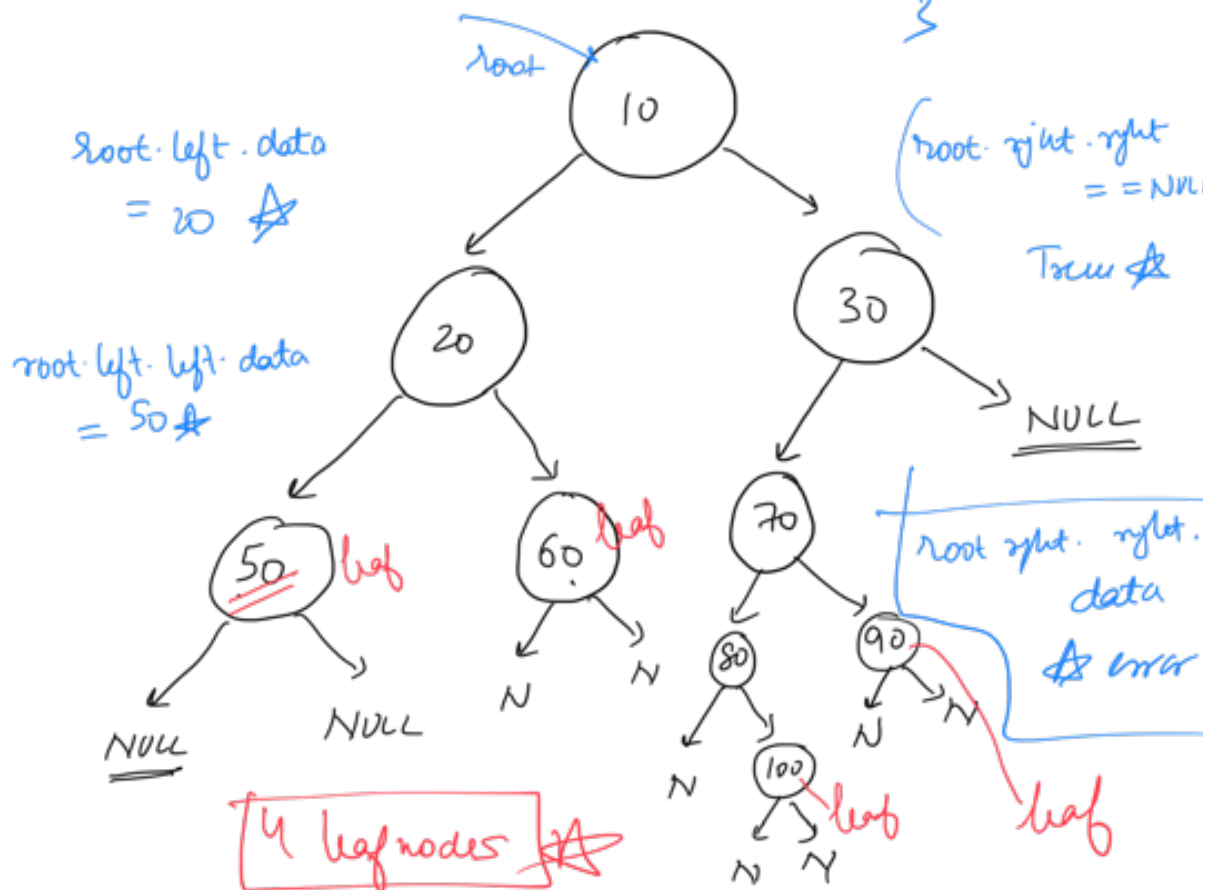
TreeNode createNode(i)

{
// memory alloc
TreeNode temp

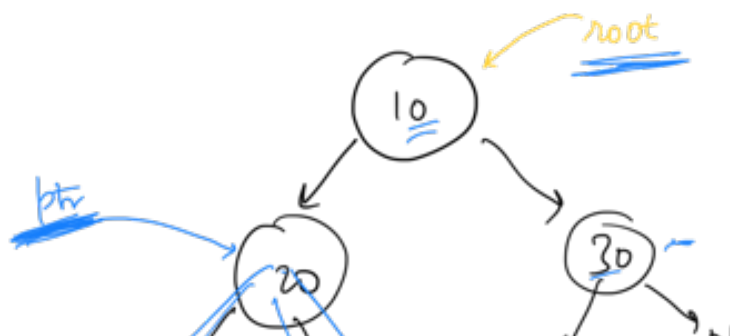


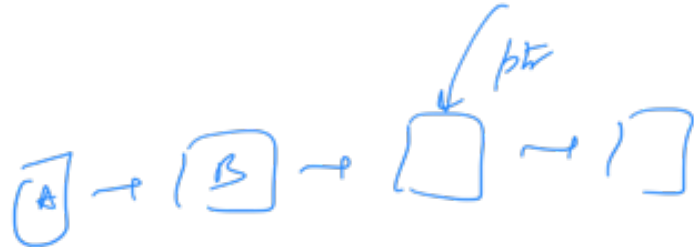
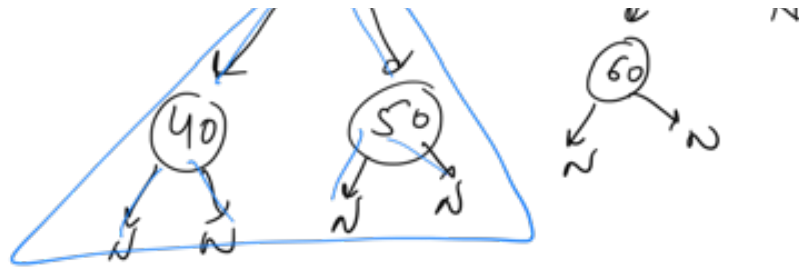
TreeNode leftChild
TreeNode rightChild

= new TreeNode
temp.data = v;
temp.leftChild = NULL;
temp.rightChild = NULL;
return temp

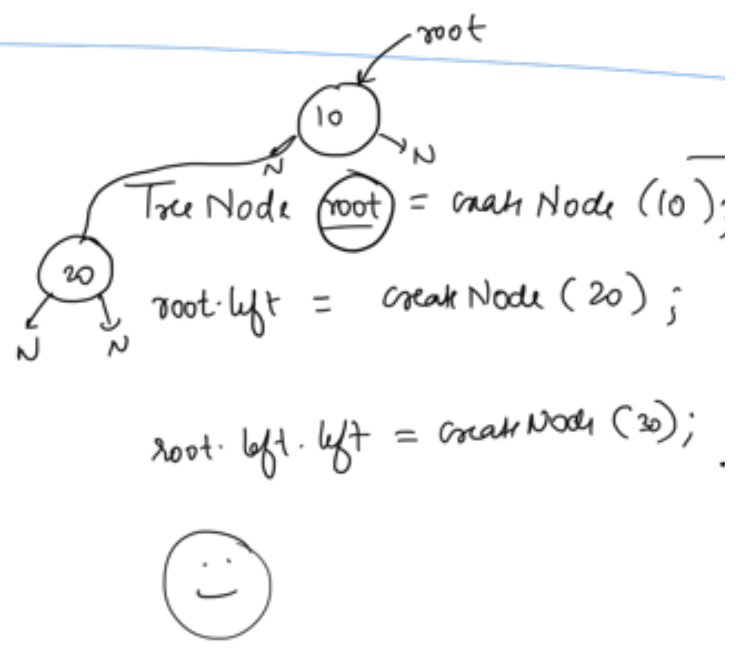
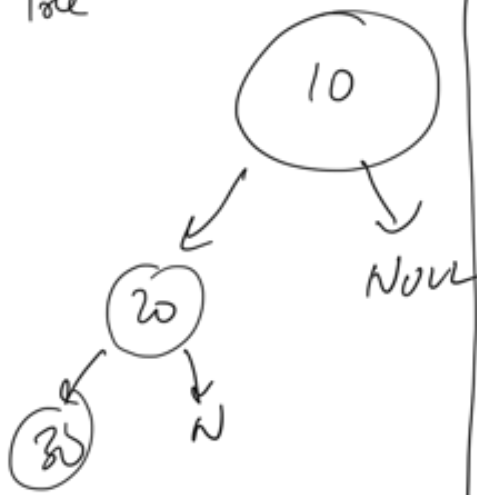


Leaf = node with 0 children





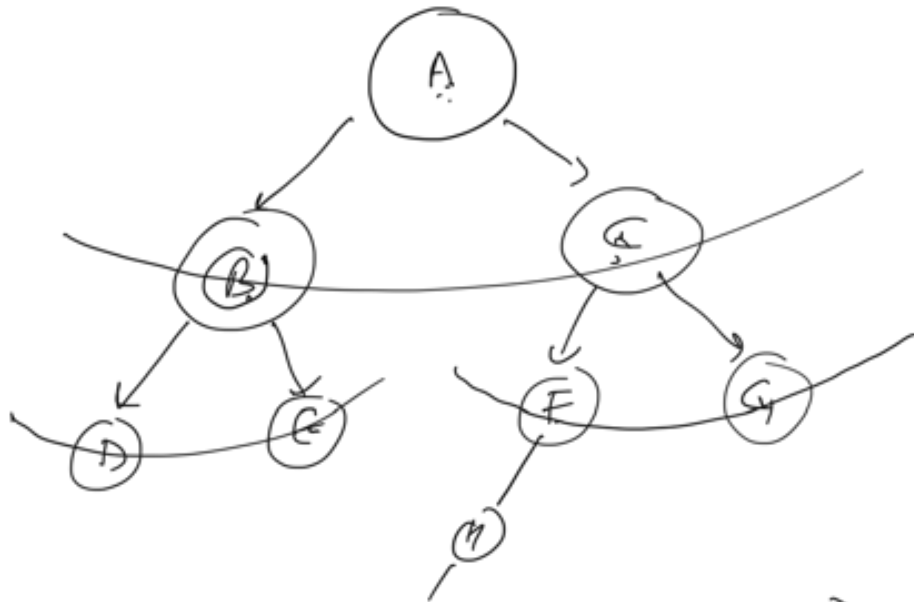
Tree



```

void traverseLL (LLNode head)
{
    LLNode temp = head;
    while (temp.next != NULL)
    {
        print(temp.data);
        temp = temp.next;
    }
}

```



Void recurTree (TreeNode root)

{ if (root == NULL) //sanity

return;

if (root->left == NULL AND
root->right == NULL)

return

Base Case

// recur

I

→ recurTree (root->leftChild)

II

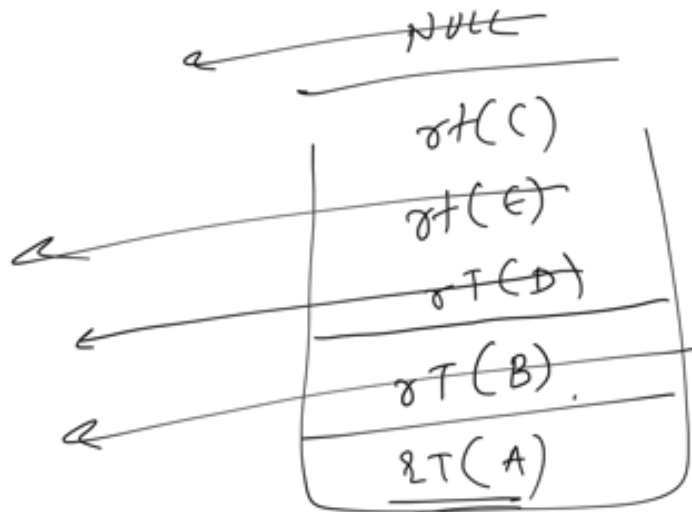
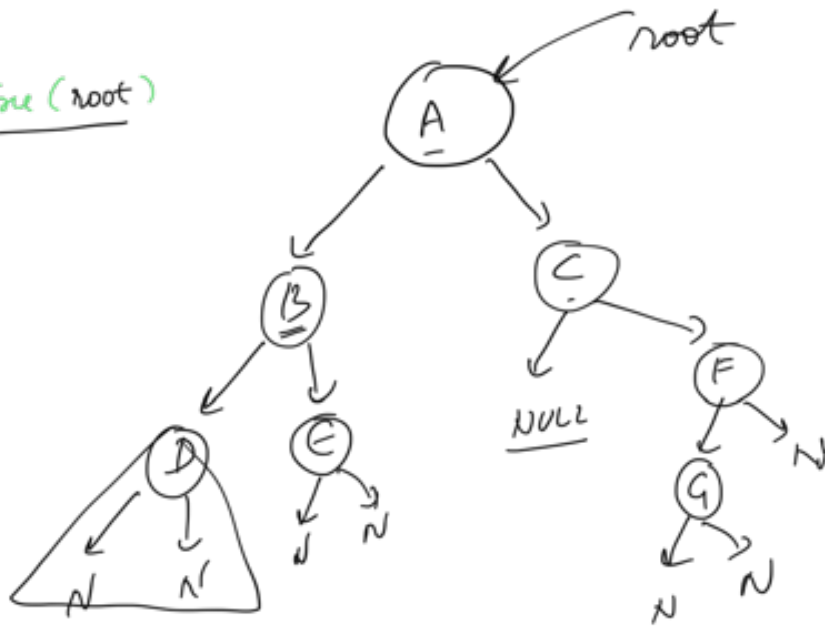
→ recurTree (root->rightChild)

III

print (root->data)

}

recursion (root)

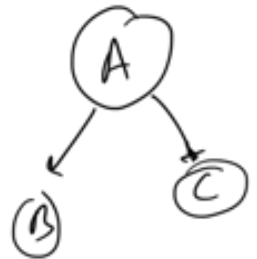


Reaction (A)

print(root.data) $\equiv A = \odot$

recurse (root, left, right),

near the (young child) -


$$\int A$$

new func

II

```

101 // recur tree
102 print (root.data) == A
103 recurT (root.right child)

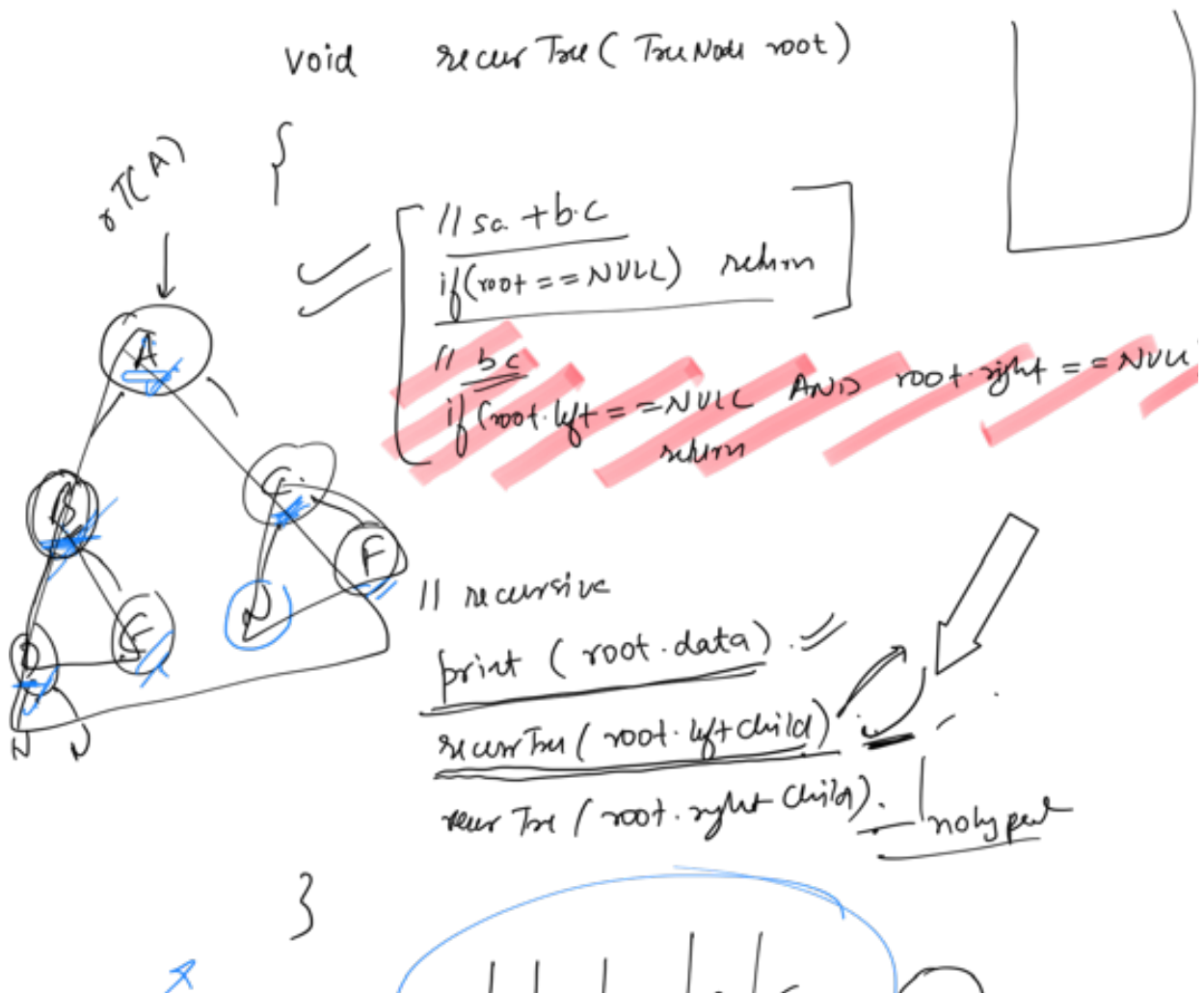
```

III

```

recur Tree (root left child)
recur T (root.right child)
print (root.data); == A

```

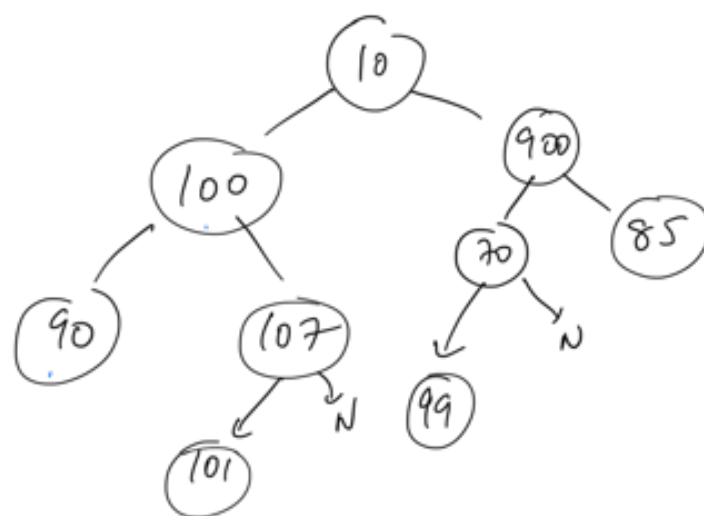
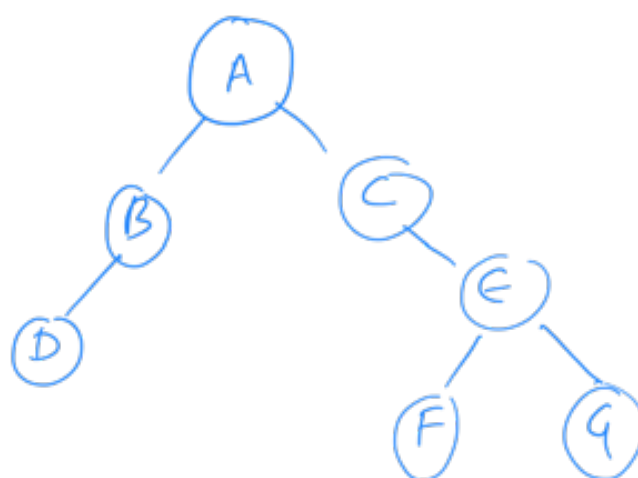


(A | B | D | E | C | F)

Preorder Traversal ☺

(recursively put parent first)

A
B
D
C
E
F
G



10

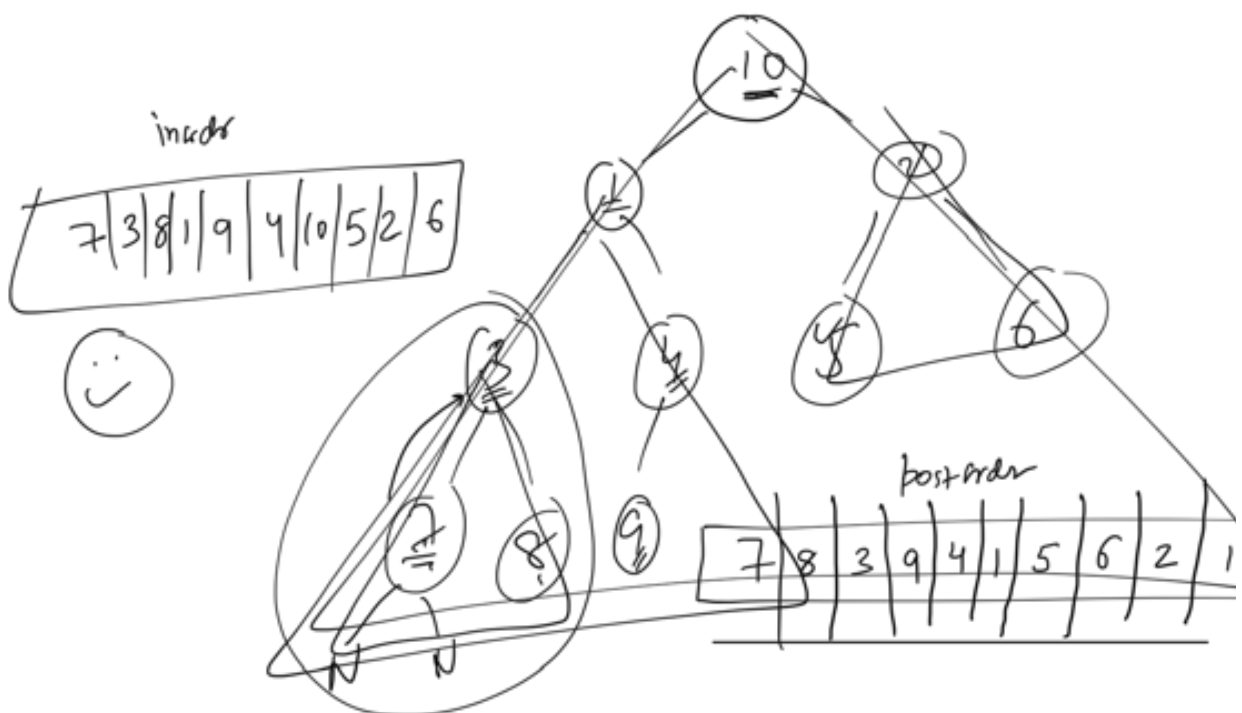
100

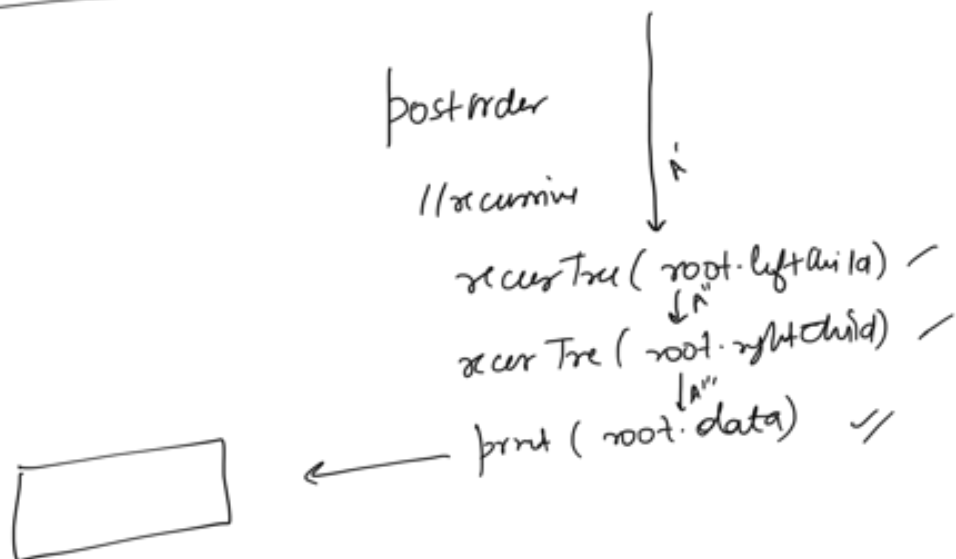
90
107
101
900
70
99
85

Inorder

// recursion

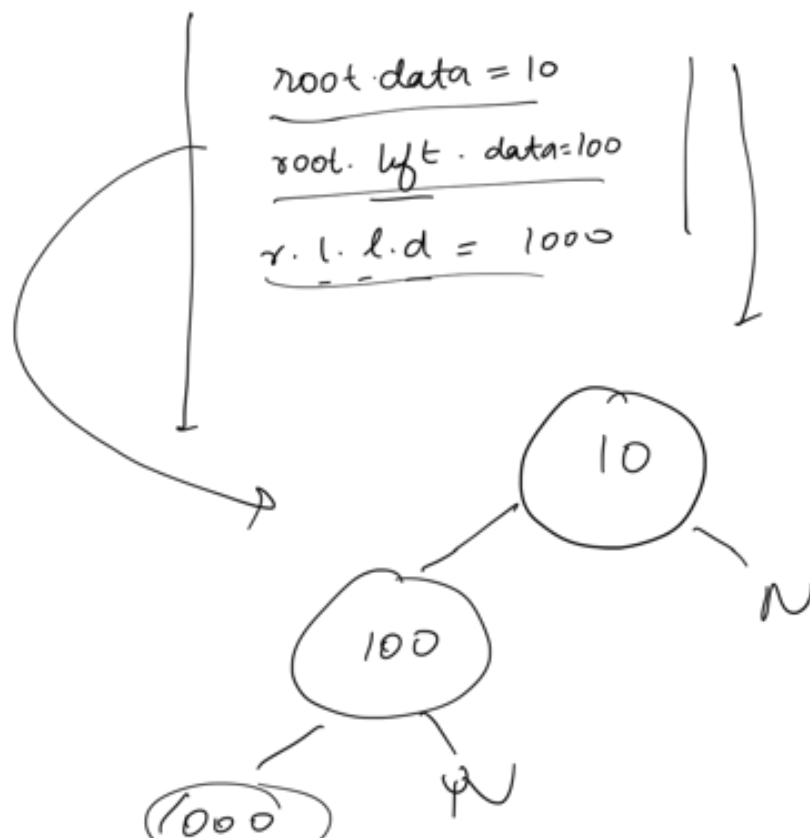
101 recurTree (root, leftChild) ✓
102 print (root.data) → A
103 recurTree (root, rightChild) →

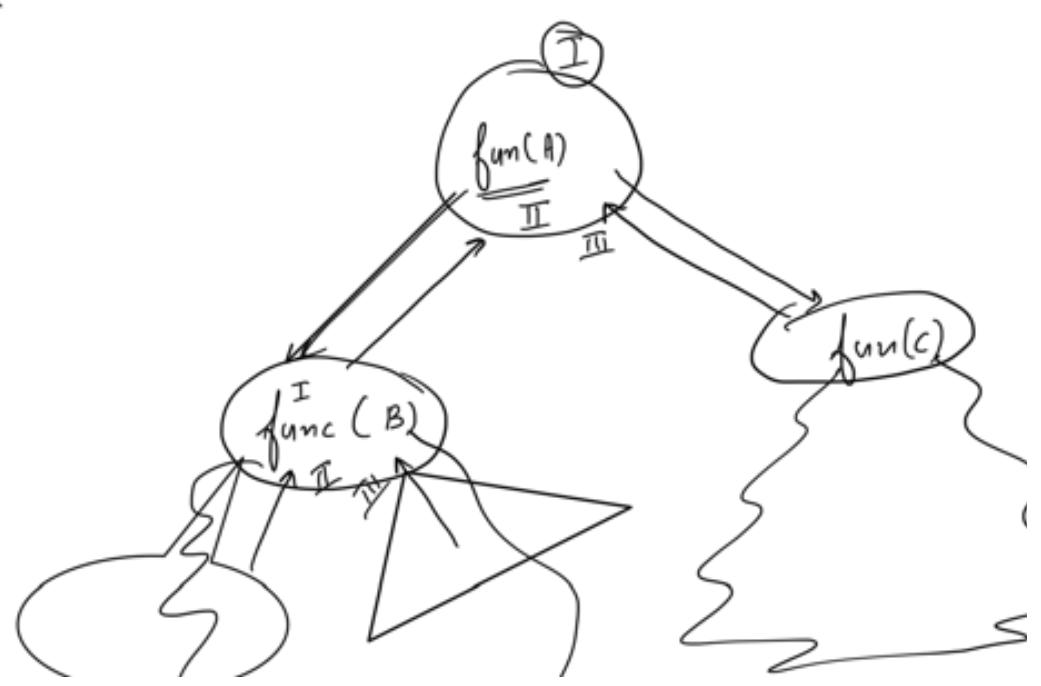
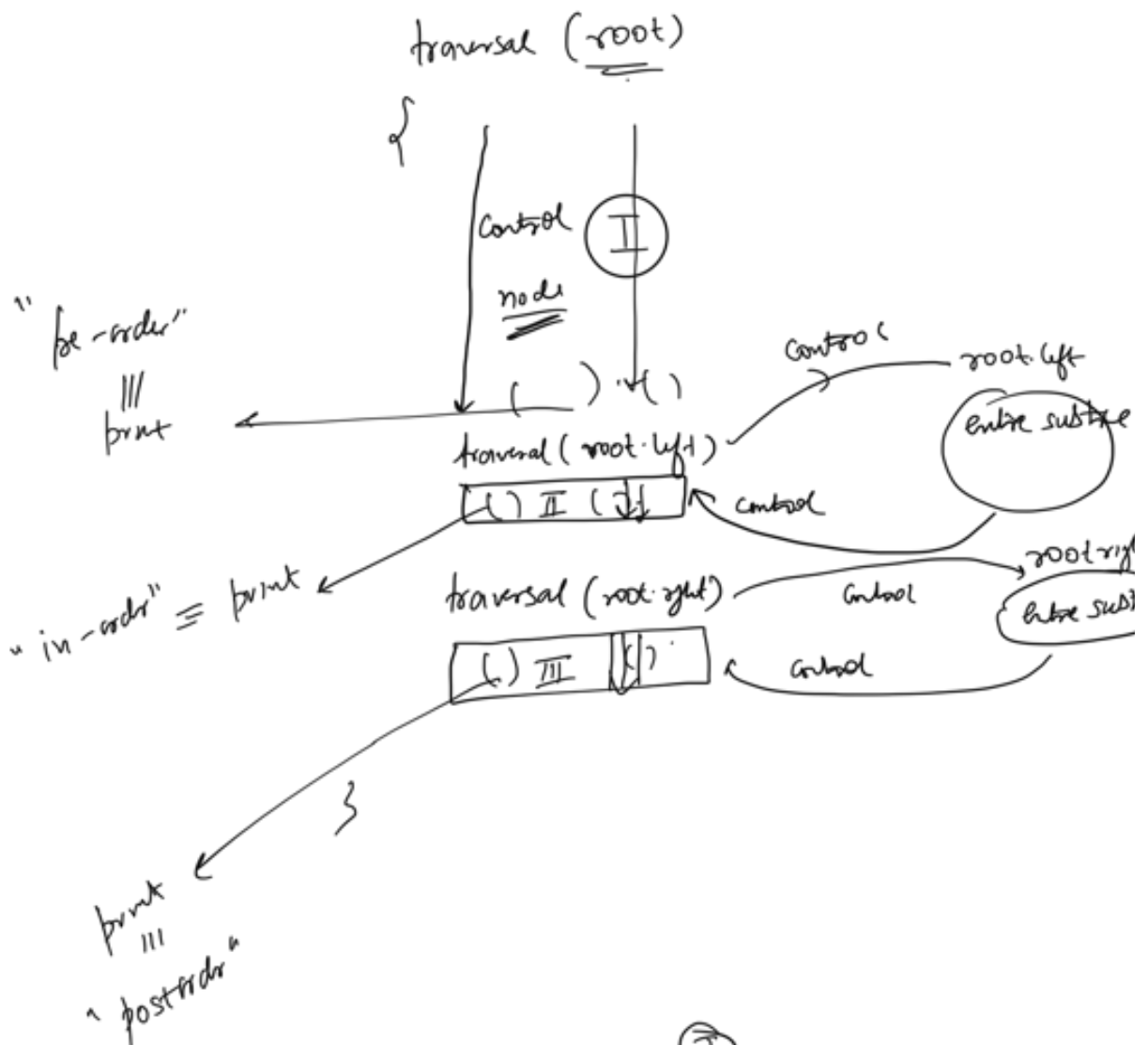


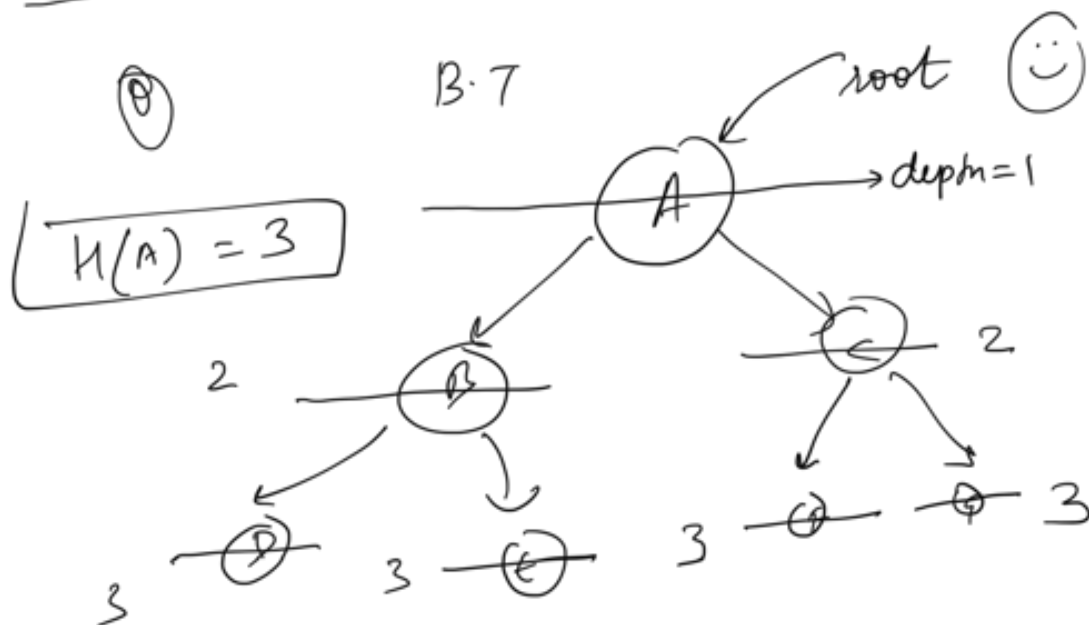
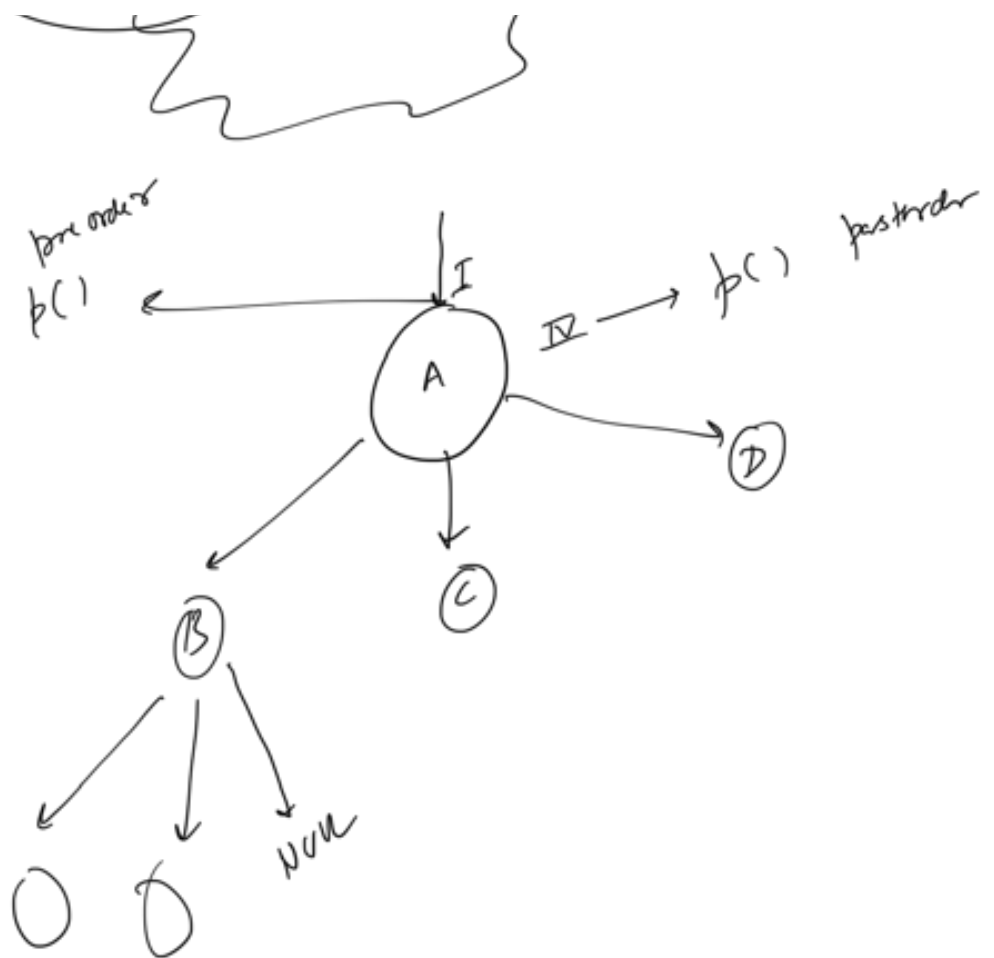


$$TC = O(N)$$

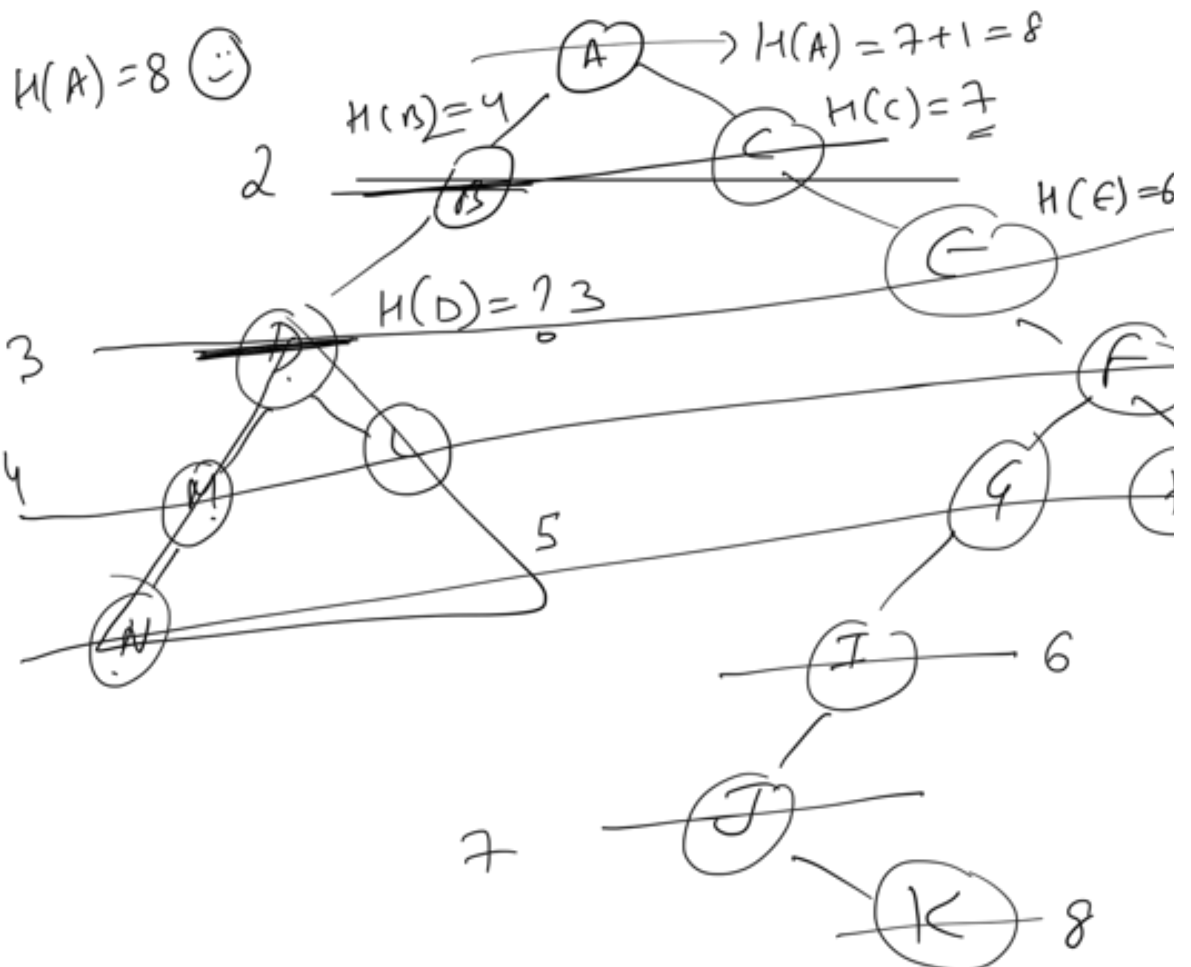
Sc S.C = (depth of Tree)
 = length of path from root to leaf







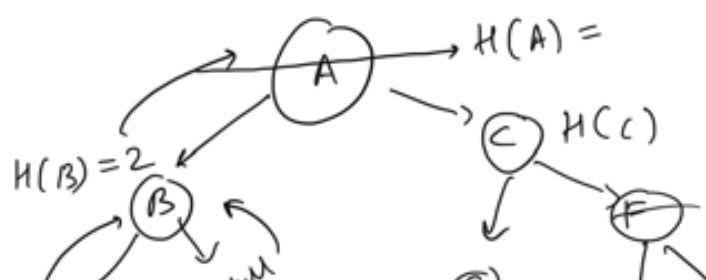
What is height of the Tree?

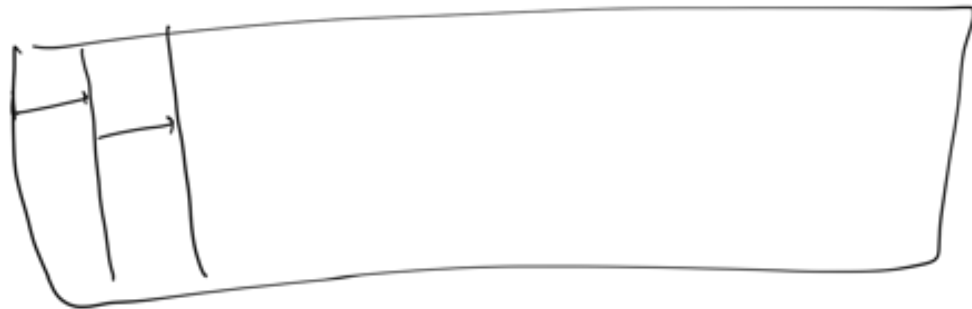
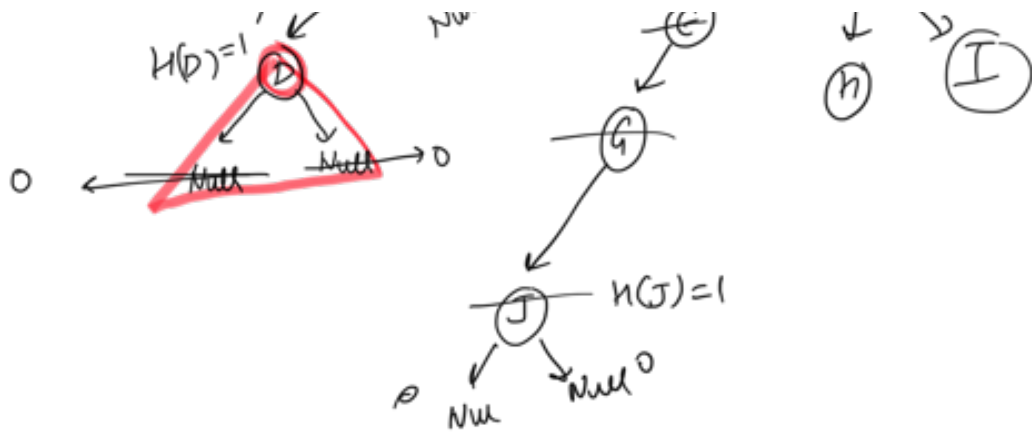


😊

$\text{Height} = \# \text{ edges in path root} \rightarrow \text{leaf node} + 1$

$H(\text{root}) = 1$





int getHeight (TreeNode root)

{

// sanity cum base

if (root == NULL)
return 0

// Base Case

OR

Option B

if (root.left == NULL
 ||
 root.right == NULL)
return 1;



Option A

Base \equiv Sanity

// recursive calls

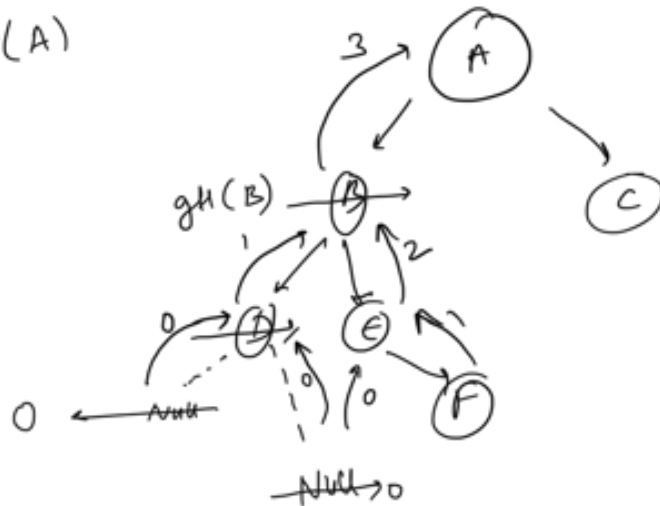
→ int height = getHeight (root.left) <=



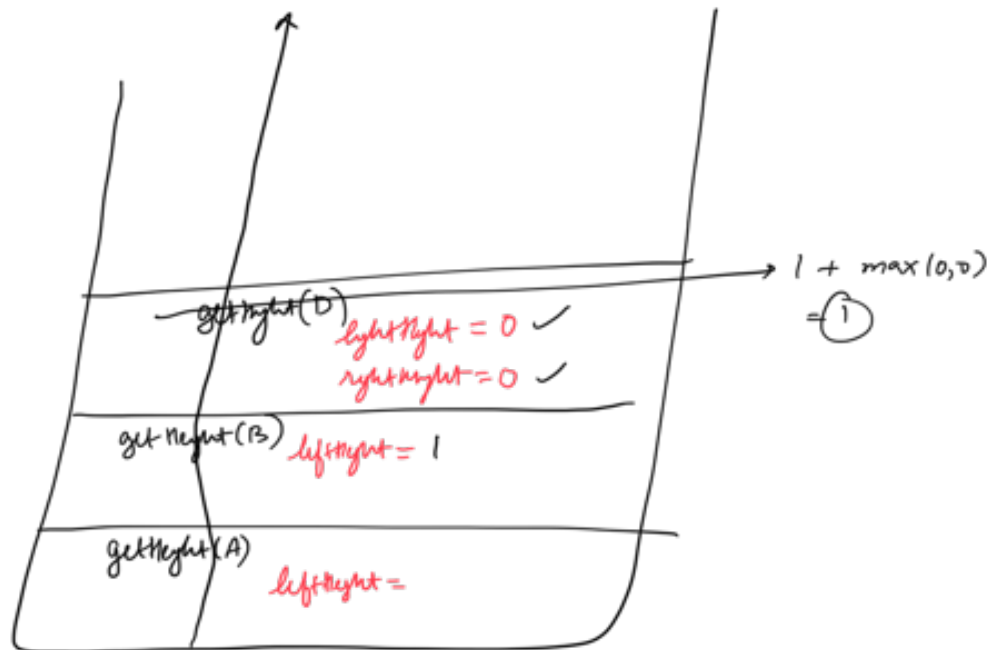
\rightarrow int rheight = getheight (root.right) =
return = 1 + max(lheight, rheight)

}

getheight(A)



recursion
stack



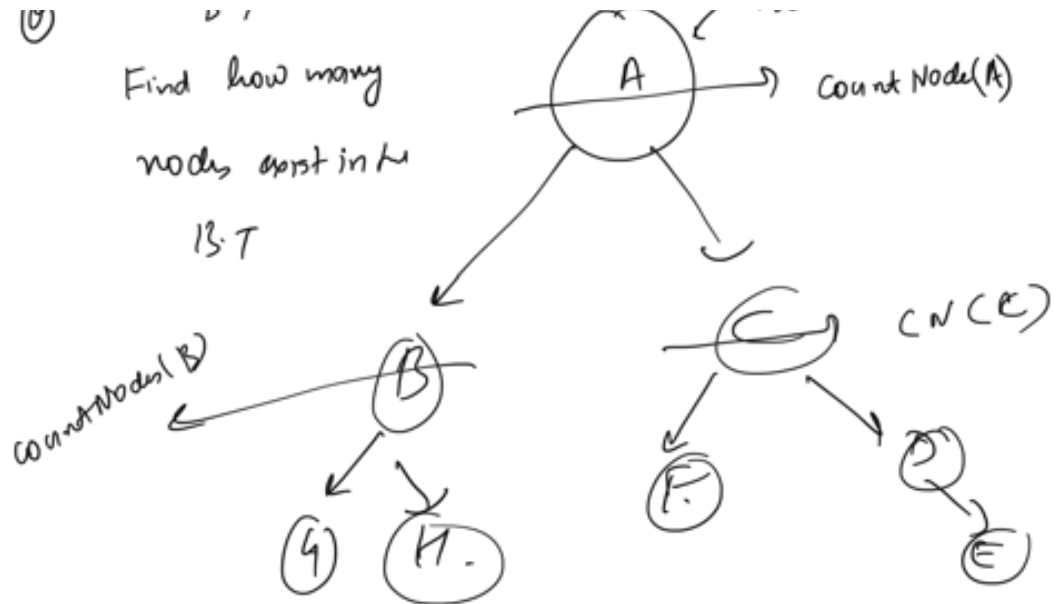
A

Given
R.T

root

Q

Find how many
nodes exist in
B.T



```

int countNode(TreeNode root)
{

```

// sanity check cum Base

```

if (root == NULL)
    return 0;

```

// Base case

```

if (root->left == NULL && root->right == NULL)
    return 1;

```

// recur

```

int lC = countNode(root->left);
int rC = countNode(root->right);
return 1 + lC + rC;

```

T.C = #state
generated
x
Time stack

$O(N)$
x

$O(1)$

$T.C = O(N)$

S.C
ans b
RS = $O(\text{depth Tree})$

$= O(N)$



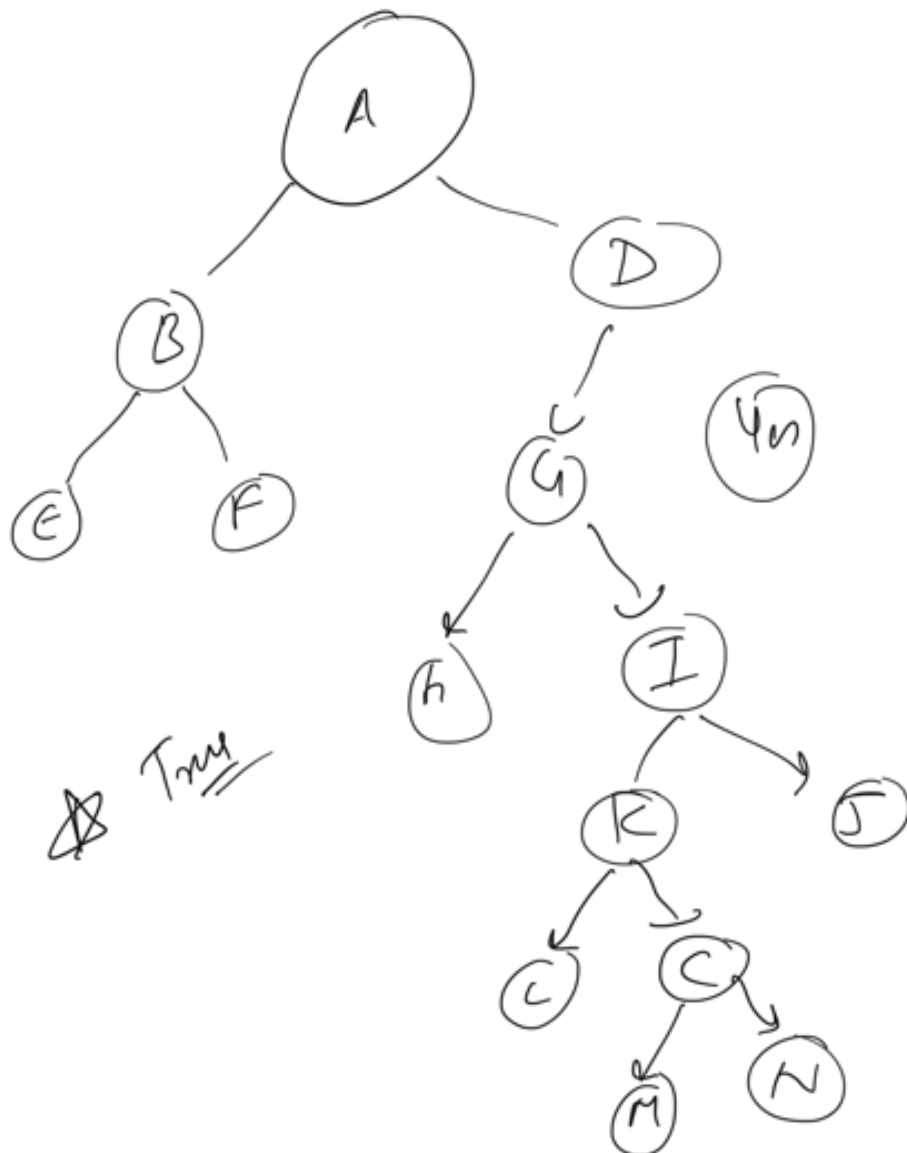
Q

B.T

you want to find

if there exists a node

whose data value = 'c'



~~Star~~ Tree

Pre order / Post order / In order \equiv T.C = $O(N)$

boolean checkExistence (TreeNode root)

```
{  
    // sanity check // cum  
    if (root == NULL) Negative B.C  
        return false;  
}
```

// Positive Base Case

```
if (root.val == 'c')  
    return true;
```

Optional

```
// Neg B.C  
if (root.left == NULL || root.right == NULL)  
    return false;
```

// rec

```
return (  
    checkExistence (root.left)  
    ||  
    checkExistence (root.right)  
);
```



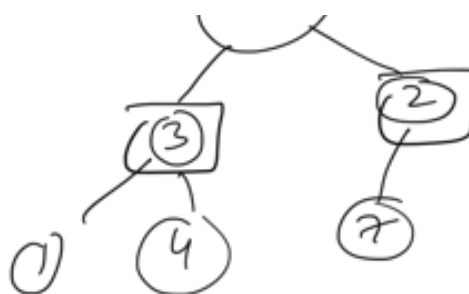
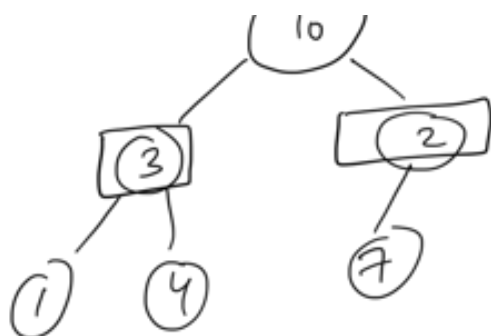
}

Q

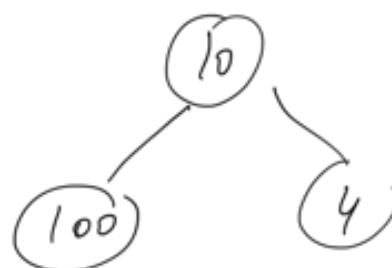
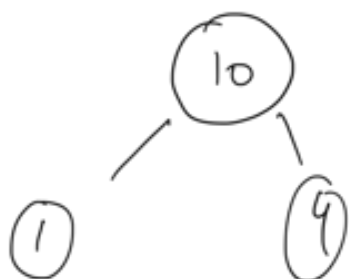
Given two binary trees

root A

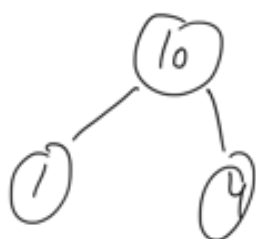
root B
10



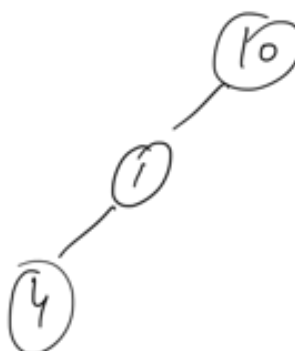
Yes.



No.

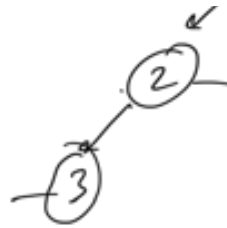


No.



* Traversal order is a boiled info in which
 * the structure detail get lost
 *



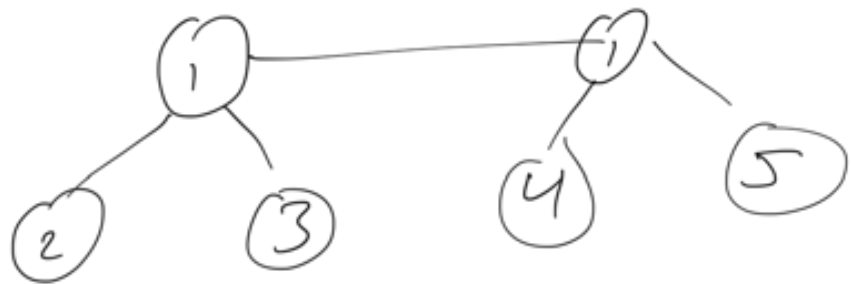


Pre order:

1, 2, 3

1, 2, 3

Recursion 😊



bool isIdentical(TreeNode A, TreeNode B)

{

// sanity check cum B.C

if (A == NULL and B == NULL) ==
return True;

\rightarrow if ($rA == NULL$ || $rB == NULL$)
 return False

// Base Case (Negative Base Case)
 if ($rA.data \neq rB.data$)
 return False;

~~// Base Case (Positive B.C)~~
~~if ($rA.data == rB.data$)~~
~~return True~~

// recur call

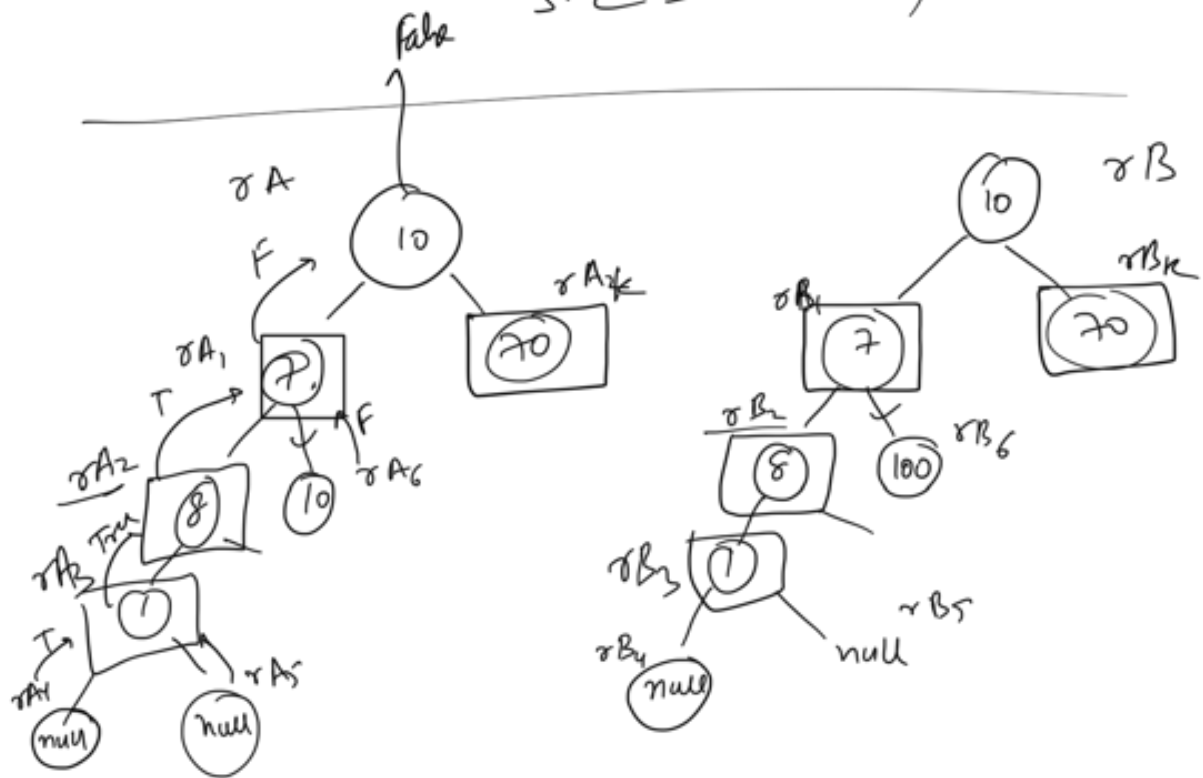


return (isIdentical ($rA.left$, $rB.left$))
 and
 (isIdentical ($rA.right$, $rB.right$))
 R.S.T

}

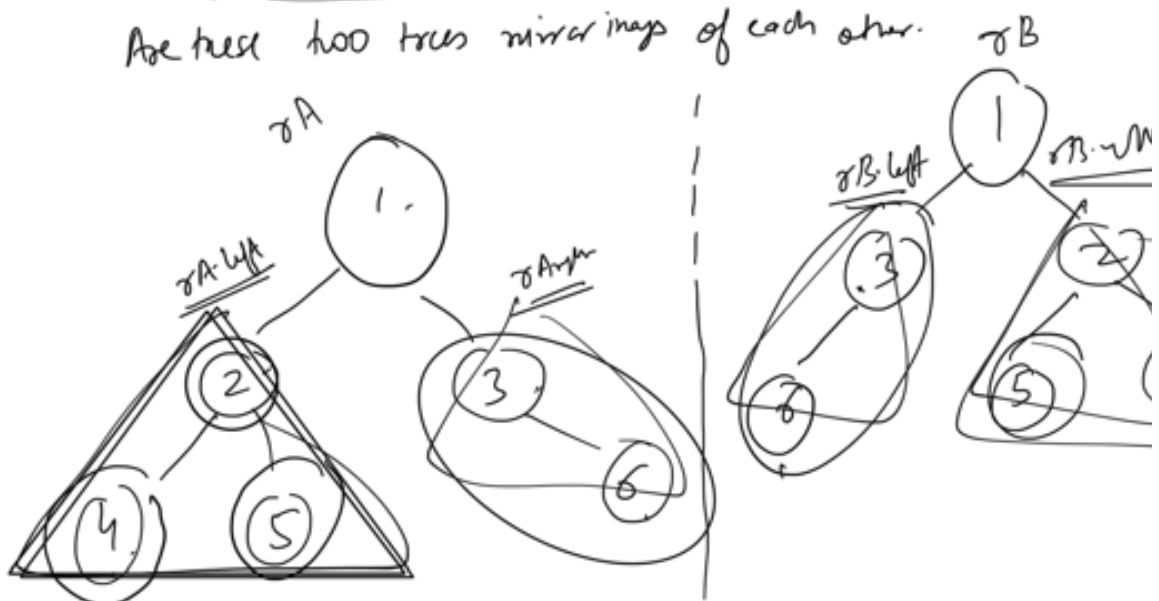
$$\begin{aligned}
 T.C &= \# \text{ static cpy bce} \times O(1) \\
 &= O(N) \\
 &= O(N)
 \end{aligned}$$

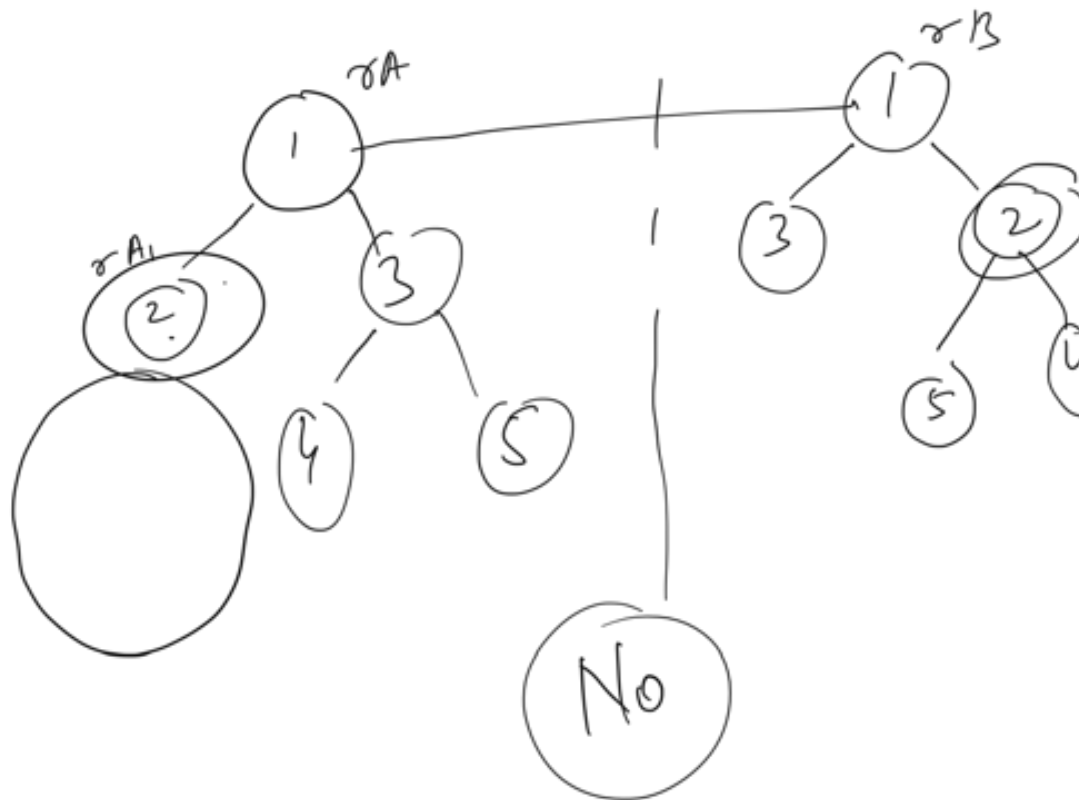
$$S.C = O(H)$$



① Given 2 Trees B.T

Are these two trees mirror images of each other.





```

bool isMirror (TreeNode rootA,   TreeNode rootB)
{
    😊
    // sanity check
    if (rootA == null and rootB == null)
        return true;

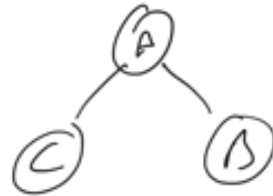
    if (rootA == null || rootB == null)
        return false

    if (rootA.val != rootB.val)
        return false
}

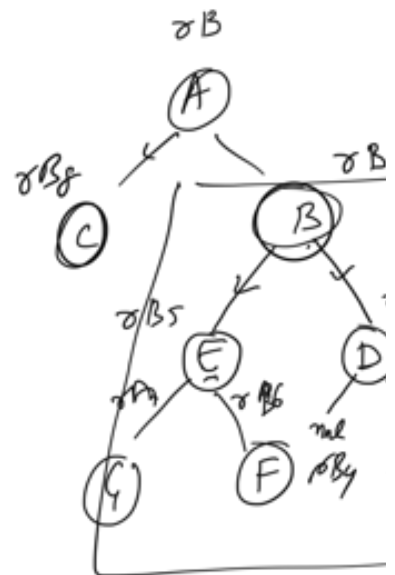
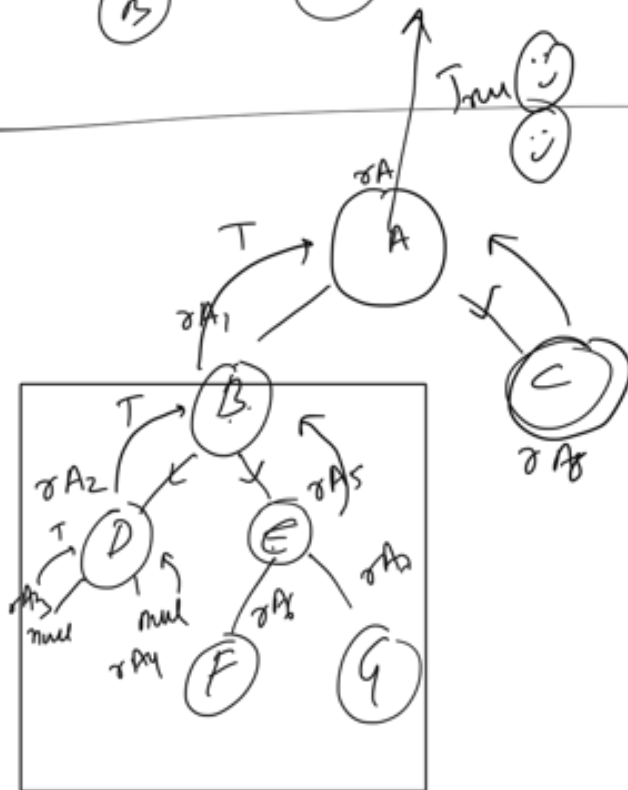
```

// Recurs
return

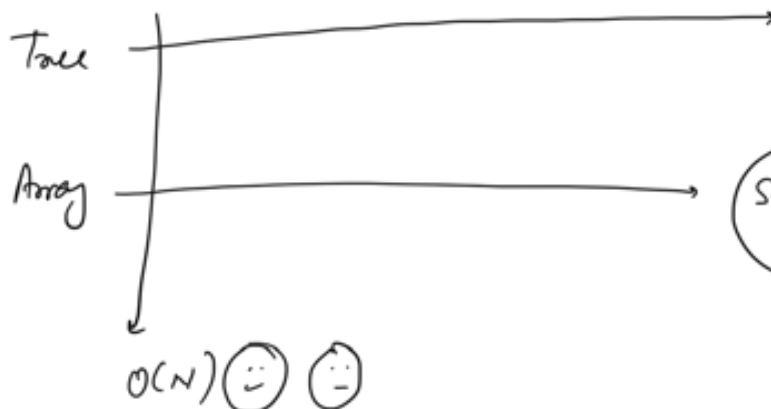
isMirror(rootA.left, rootB.right)
and
isMirror(rootA.right, rootB.left)



True 😊
😊



Check Extra



Binary Search Tr

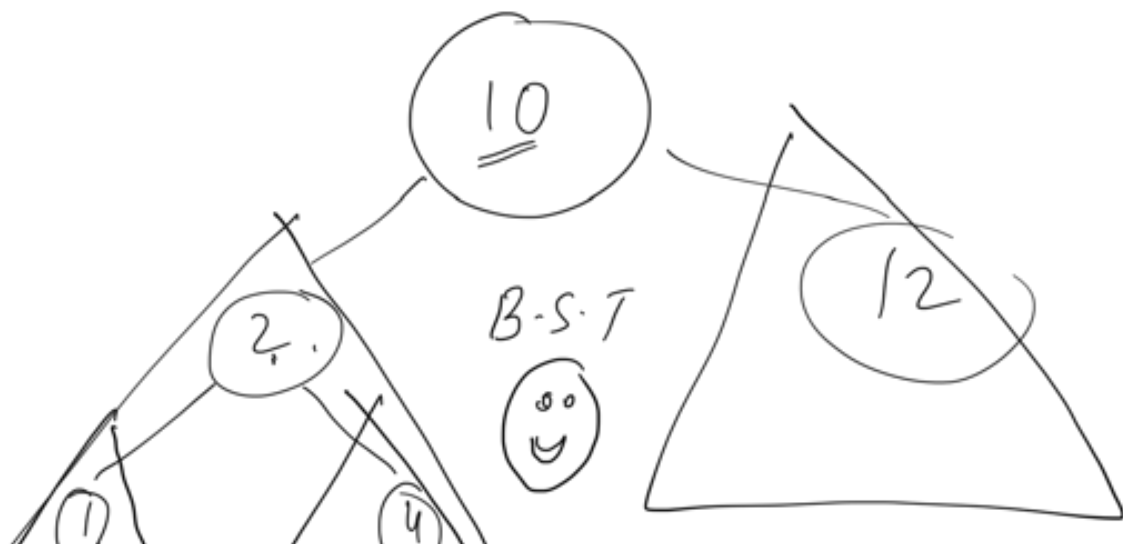
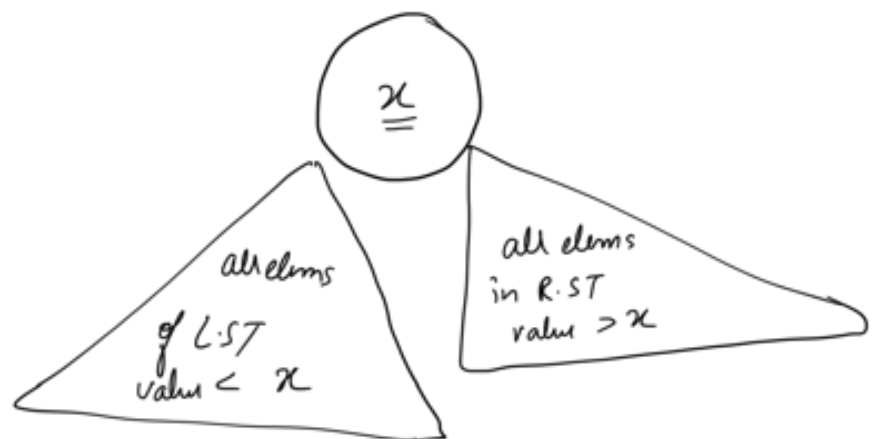
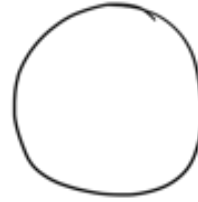
Sorted Array
Binary Search

$O(\log N)$



Binary Search Trees

- (i) \rightarrow B.T
- (ii) helps in searching
- (iii) recursively





Recursion

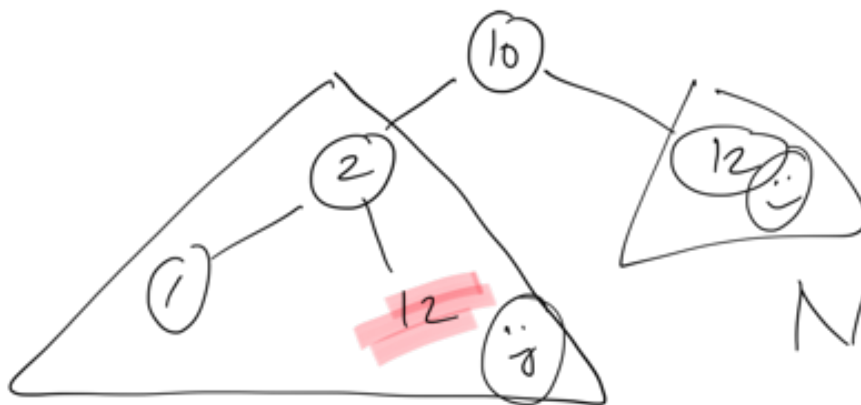
every elem in your L.S.T. should be

smaller than your root
or equal to

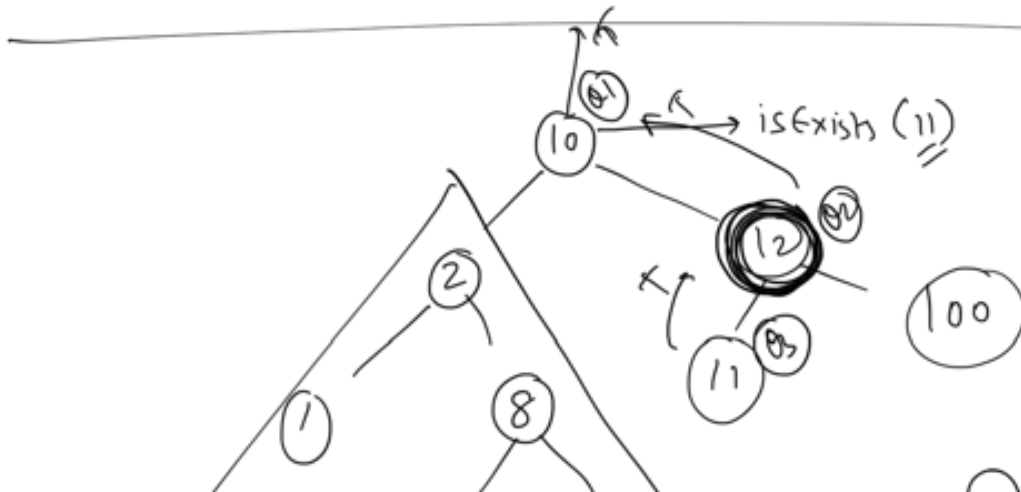
AND

every elem in your R.S.T. should be

bigger than your root



NOT
a B.S.T





45 (✓)

$$\begin{aligned}
 T.C &= \# \text{ stack exps} \times \text{Time for} \\
 &= O(n) \times O(1) \\
 &= O(n)
 \end{aligned}$$

```

bool findInBST (TreeNode root, int val)
{
    if (val == null)
        return false;
    if (root == null)
        return false;

```

```

    if (root.data == val) ✓
        return true; ✓

```

// recur

```

    if (val > root.data)
        return findInBST(root.right);

```

```

    if (val < root.data)
        return findInBST(root.left);

```



