## Assignment 2

Tasnim Sakib Apon

ID: 20241068

**CSE 490** 

## Assignment 2 Tagnin Salib Apon 20241068

! (a) Chiper text = plain tex (1) key stocker.

Co = 1. (1) So.

Both senter and occurer have same algoritm, so both In

Acice case, & Toudy knows the plain text and the chiper text,

then If he does X-oR operation with plaintent and

Chiper text, then he can get the key stocen easily.

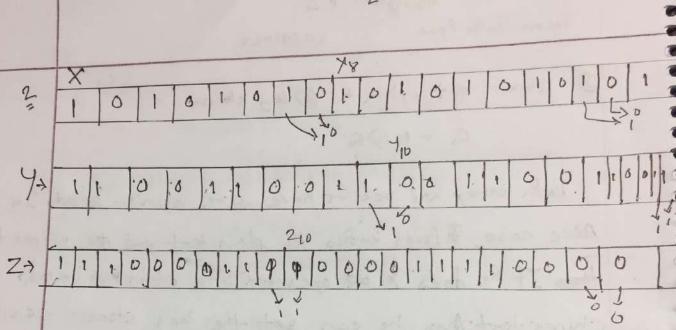
So-PoDCo, Si-P, DC, Sz=P2 DC2

Troudy cannot change the key stream but she can change plain text. I to t'. In order to get a new chiper text c'we can say that trudy will replace the equation of decreption.  $c = i\Phi$  s to  $c' = p\Phi$   $f\Phi$ s are know,  $c = c\Phi$ 

to get c'= p' DS.

We will replace the key and will get door

c'= p'DPDC



mg (x8, 410, Z10) = mes (1,0,1)=1.

The 1st key stream bit

X8=1=m-1:

so x has to be shifted

value will be placed in new to position.

Y0 = X13 @ X16 @ X12 @X18

Value in X0 = 0 0 1 0 00)

- 0

· Por Likewise

Zx =1, 20=0, Z1=0 Z2=0.

5° 2×072002210226 = 1

So Neu,

X = 01010101010101010

7 = 110011001100110011

Z = 1111 000001111 00006/111000

718 0 42, 0 222 = 1

ilst key stoean bit is 1.

For 2nd heystocan bit.

\*8=1, 710= 0, Z10=1.

mamas (1,0,1)=1

so register X steps. I does not & z steps.

Z13=1, X1, =0, X12=1, X18=0, X3+

×13 0 ×16 0 ×12 0 ×18 = 0

Z = 0Z 20= 6, Z 2 2 2 0 2 2 2 2 0

· 2 x 1 20 0 2, 1 2220

So 1010000

.. 2nd key to stoom bit = 0.

For this d key stream bit, 18=1, 410=0 210=1 mamas (1,0,1)=1

so register x steps, y doesn't and Z steps.
\*13=1, \*16=1, \*12=0, \*18=1

- : x130 x16 @ x120 x18= 1.

Zz=0, Z20=1, Z21=0, Z22=0

 $Z_{1} \oplus Z_{20} \oplus Z_{21} \oplus Z_{22} = 1$ 

NOW

X= 1001010101010101010.

7= 110011001100110011

Z= 10111100001111000011110.

×18 0 72, 0 722=1

- : 3 od & bit & treams 1.

4th: X8= 0, y10= 0, Z10=1

-", m = mag (0, 0,1)=0

Here x and s will step and z does not step.

\*13=1, \*16=0, ×12=1, ×18=0.

X13 @ X16 @ X12 @ X18=0

Y2021, Y2121,

So 720 D 421 = 0

Nov, X= 01001010101010101

4= 011001100110011001

Z= 101111000001111000011110.

\*18=1, 921=1 Z22=0

\$0 \*18 \$\P 721 \P Z22=0

... 44 is Key stocen bit = 0.

... Next 4 key stocen bit = 1010

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- €) Bits in each plain text > 69
- (6) Bits in each chiper text => 69
- co) Bits in the key => 56
- (d) Bits in each subley => 48
- @) Number of rounds => 16
- 43 Number of S-boxes => 8
- (9) Number of bits of input too 5-60x => 6
- (4) Numbers of bits of output of 5-60x =>9

In AES Add round key and byte substitution layer are used for confusion. The concept of confusion is making the relationship between the key and the chiper text. It increases the ambigity of chiper text. So adding round key and byte substitution in AES is considered as confusion.

The shift row layers and mix column layers are used for defusion. The aim of defusion is to hide the relationship between posses plain, and chiper text, so that attacher can easily find out the plain text. In AES, shift rows and mix column layers make the plain text complicated to understand.