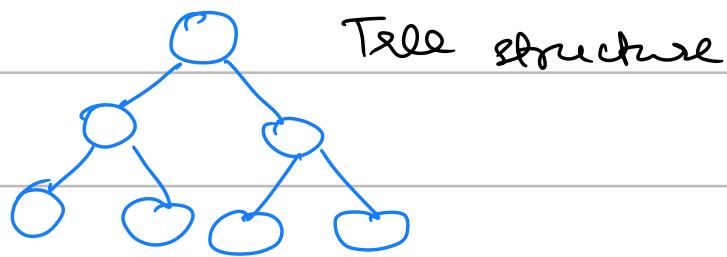
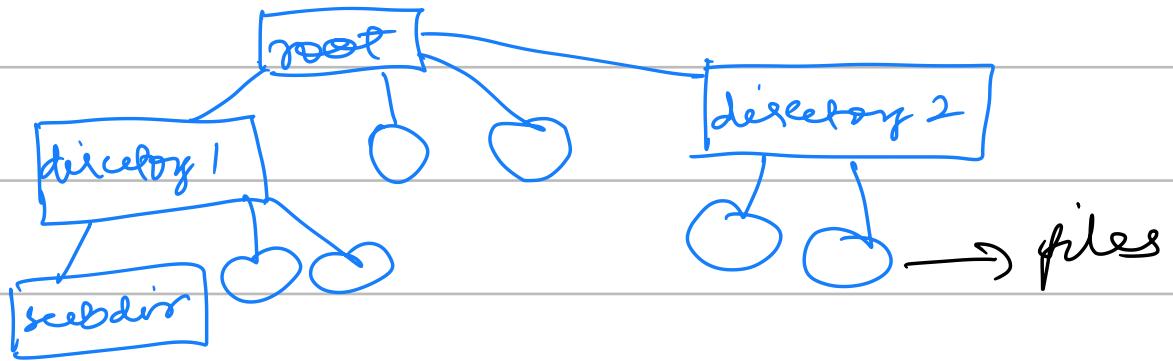


Trees in Python



Tree structure

Mostly we see use of tree in OOPS file structure



Another example is a companies hierarchy

What exactly is a tree structure?

There exist a node , root node main node, this node can have more nodes as its children.

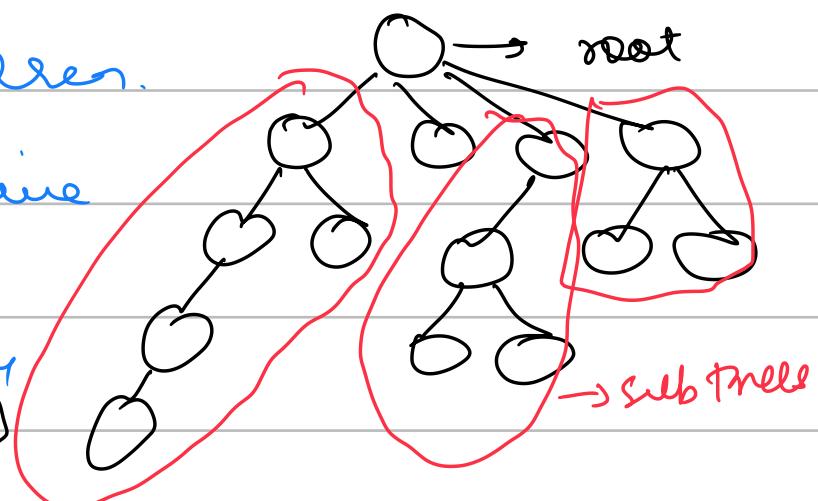
Now each of these children can further have more children.

Trees a highly recursive

structure and we

mostly use "recursion"

To solve trees.



* Binary trees

In binary trees each node can only have either 0 or 1 or 2 nodes.

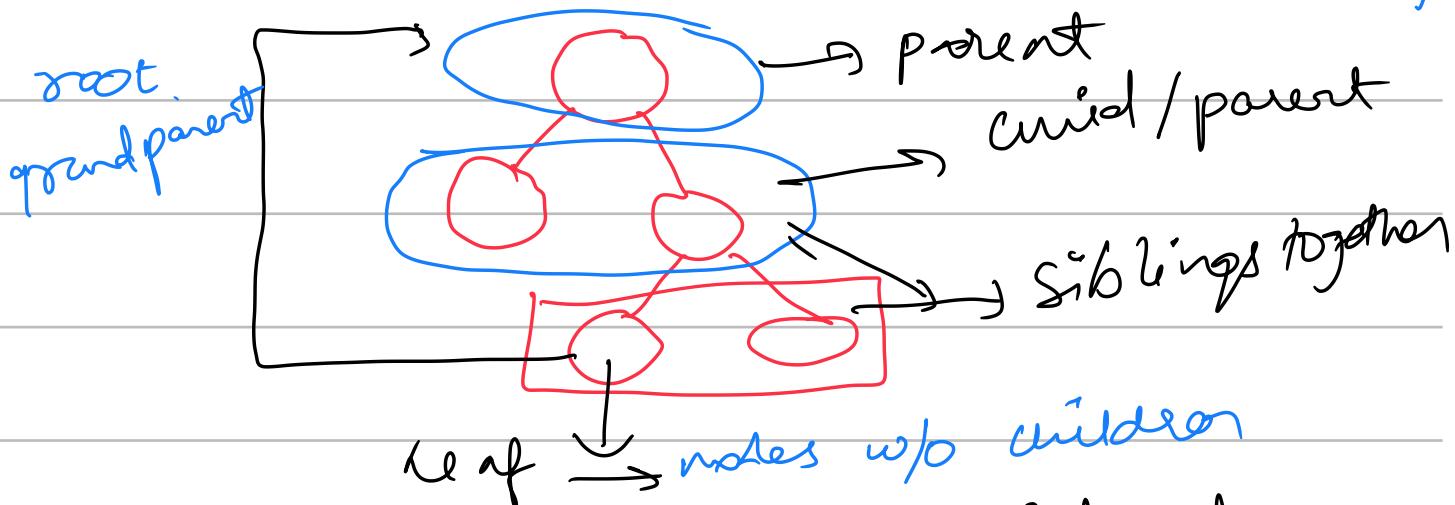


* Terminology of trees

Node → individual data point in tree

Root → top most node is called root.

It's like "head" in Linkedlist. We can search the whole tree or traverse the whole tree through



Important terms → node, root, parent, child, and leaf.

* So what will a node store

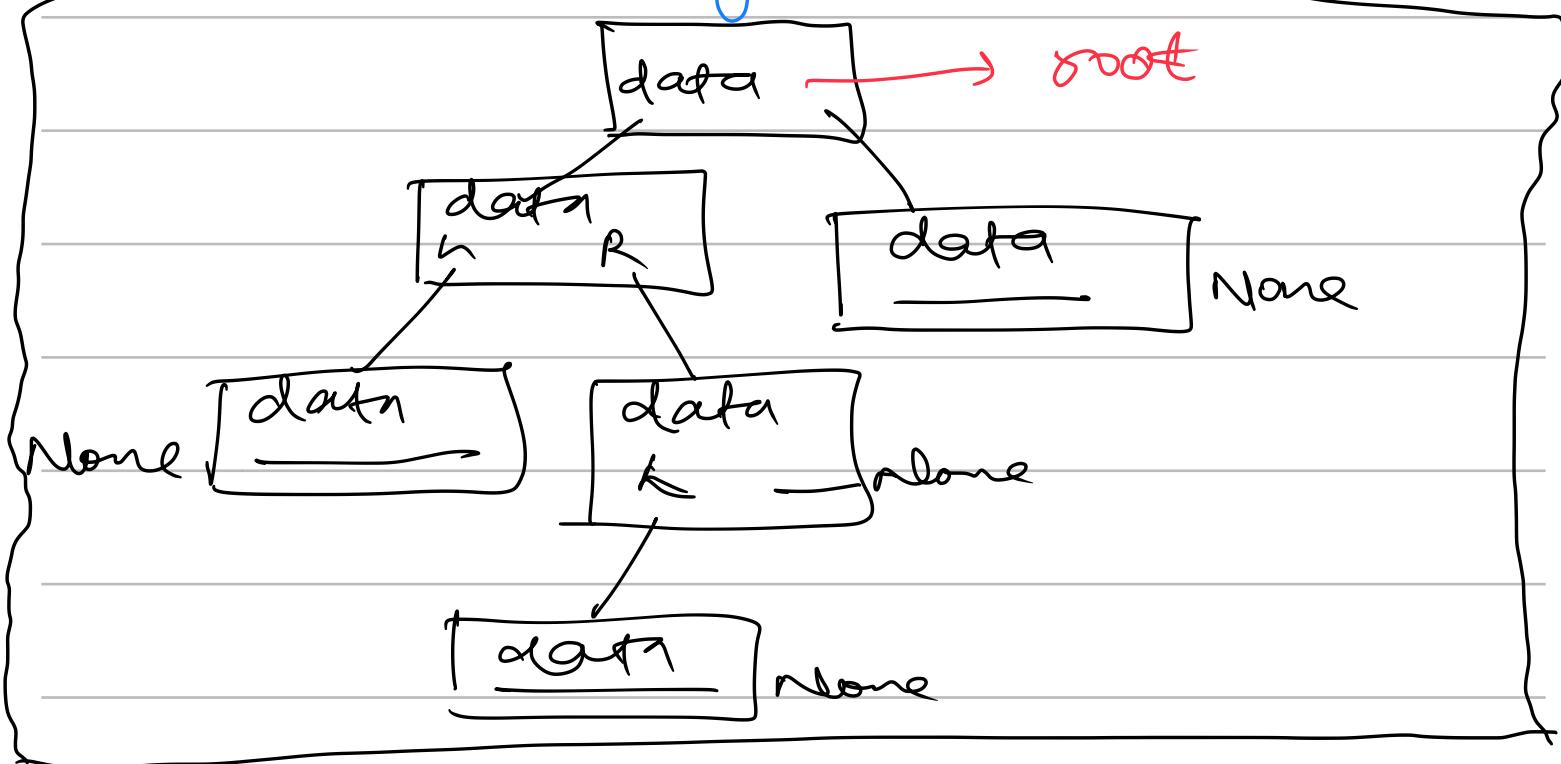


This way Node will have 3 attributes

- Address of Child 1 Address of Child 2
- 1.) Data
 - 2.) Address of Child 1
 - 3.) Address of Child 2

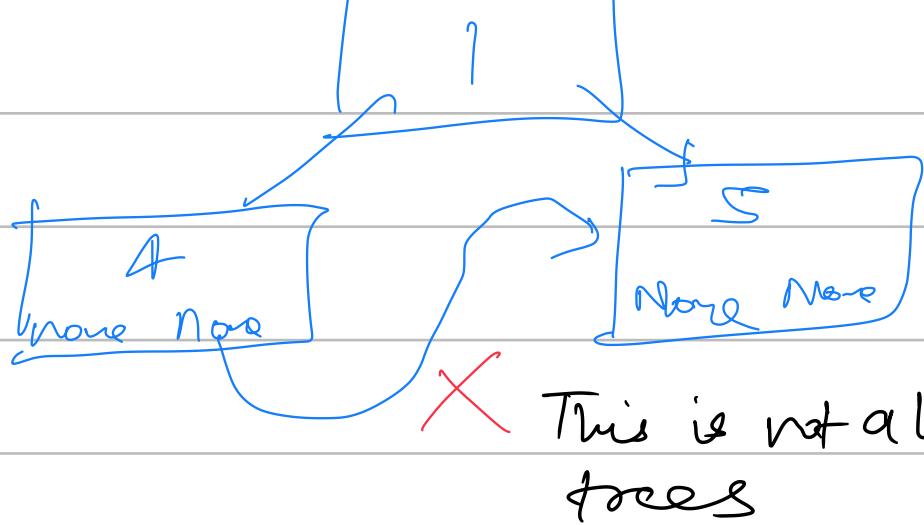
In case of binary trees these child are called left & right.

keep Address of the particular node's children. The **addresses** can also be **None** meaning that particular node doesn't have any children.



Note - Check github for code for BT.

Note → Trees doesn't have cycles



```

class BinaryTreeNode:

    def __init__(self,data):
        self.data = data #our node should have some data
        self.left = None
        self.right = None #two address holders for left node and right node

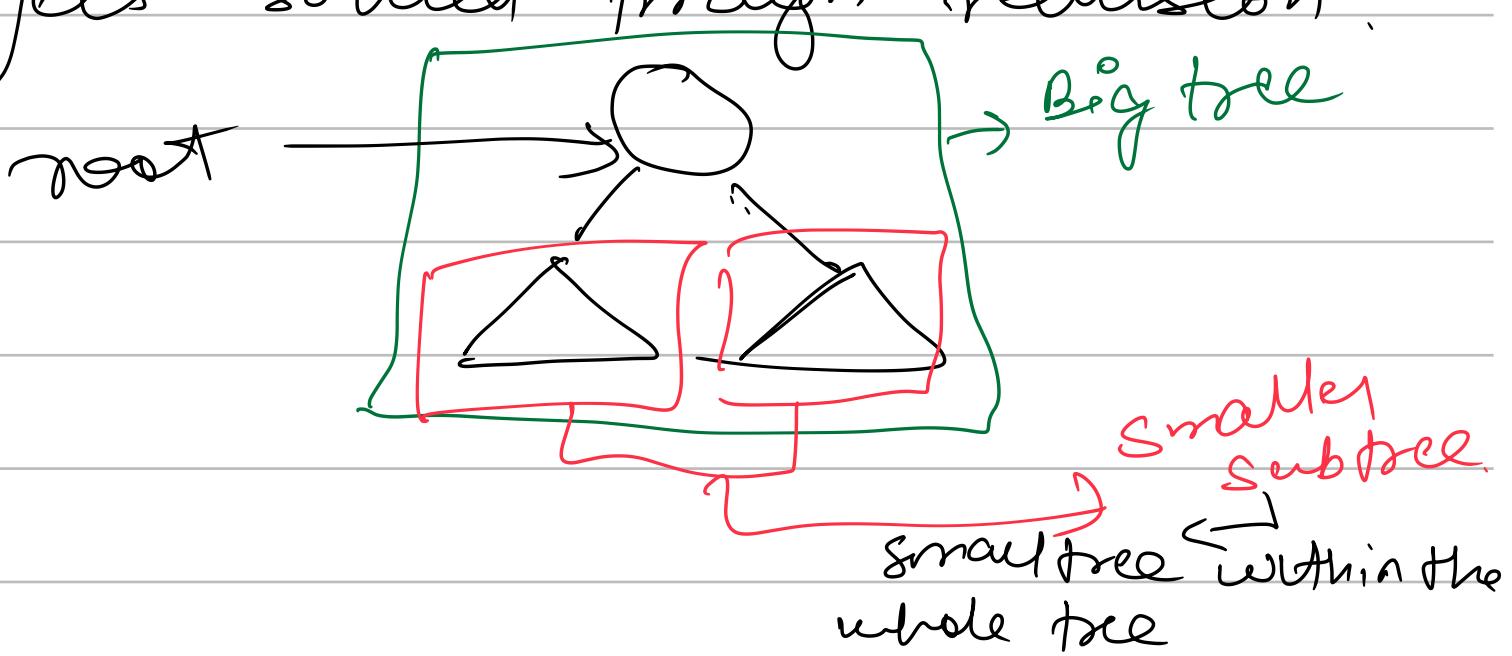
    rootNode = BinaryTreeNode(2)
    node1 = BinaryTreeNode(4)
    node2 = BinaryTreeNode(8)

    #connected children to parent node
    rootNode.left = node1
    rootNode.right = node2

    print(rootNode.data)
    print(rootNode.left.data)
    print(rootNode.right.data)

```

Note → most of the tree problems
gets solved through recursion.



we will print root's data and then
use recursion on the remaining left &
right subtree

What what will be the base case?

① If root a leaf node, then there
is no point on calling on the left
node & right node.

Now there are two problems

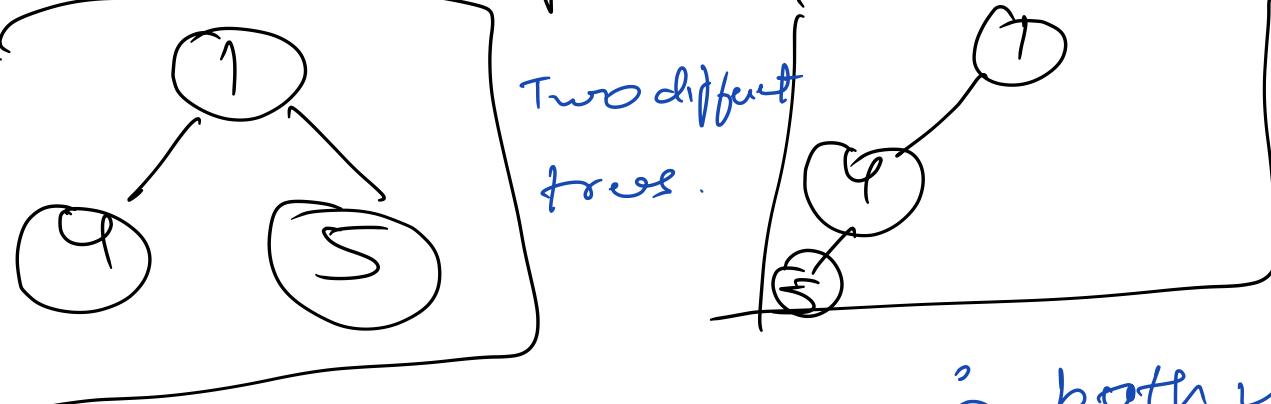
with this approach

I) what if root is "None" i.e.
tree is empty.

II) we need have a overhead of
checking if left is None then don't
call, if right is "None" then don't call.

So in most of our cases we take
the base case that If root is None

But this still does not tell us the
complete picture



Using the above scenario both will print 1,4,5. But there is confusion b/w the parent child relationship.

So what to do?

→ need to print parent child relation.
we need to print in a way like
root: L, R or 1 : L4, R5

so to print left nodes data we
can do root.left.data → this print
left child's data.

"Now this can only be root.lft
not None"

or else we will get an error.

Taking user input for Binary Tree?
we take input recursively as well
user gives root and take input for

left subtree and right subtree.
All this except for left & right subtree
will be through recursion.

Note :- while taking input the
whole first subtree whether be at
left subtree or right subtree has to
be completed before moving to the
other subtree.

given not find number of nodes in
tree ?

Base Case if root is None
return 0. Otherwise call
recursion on left tree, pass on the
right. Add 1 + leftcount + right count
and you are done.

Believe in Recursion, just think
About Root

Traversals on Binary Tree

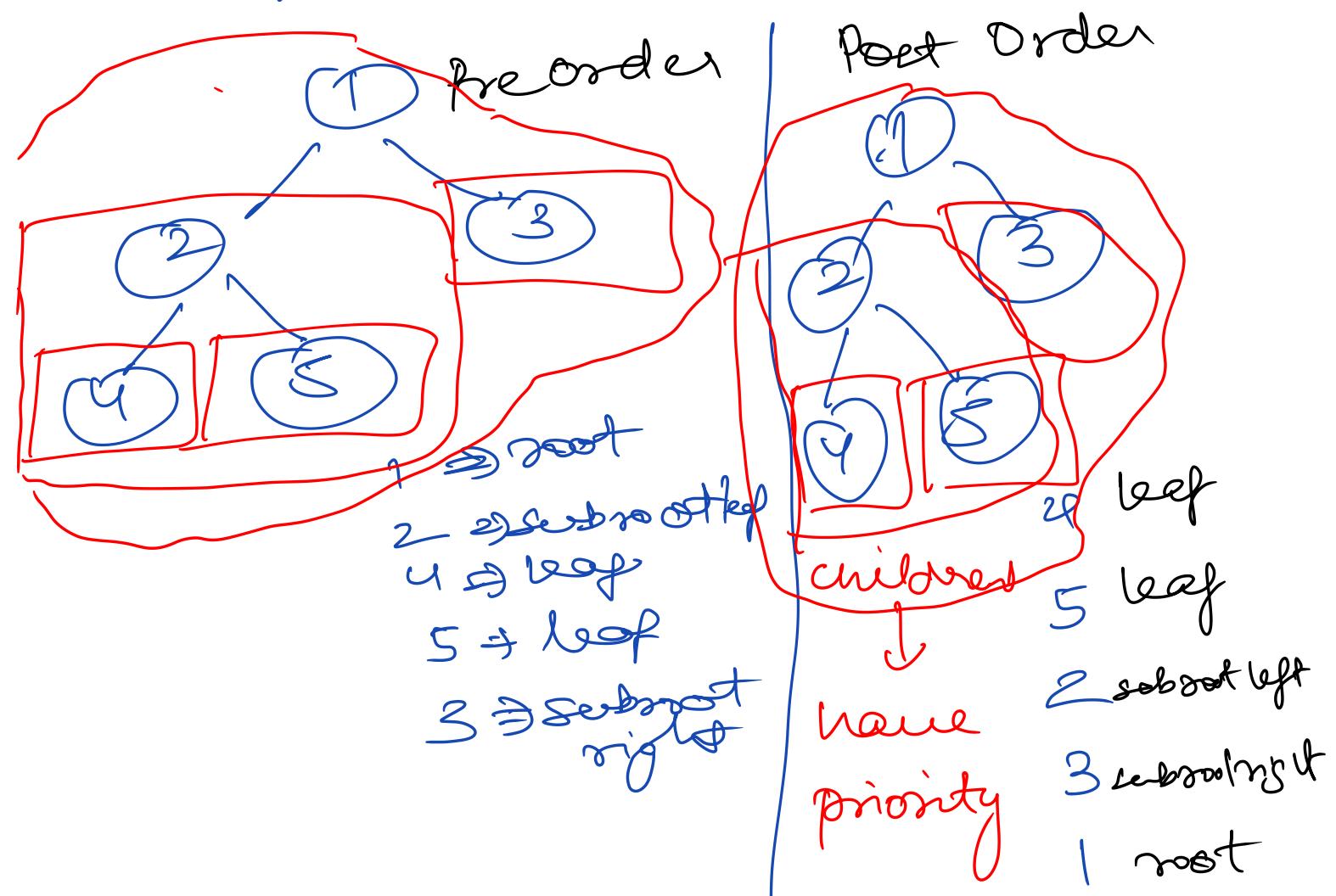
a) Pre order

Root left right

never root takes for himself
before children

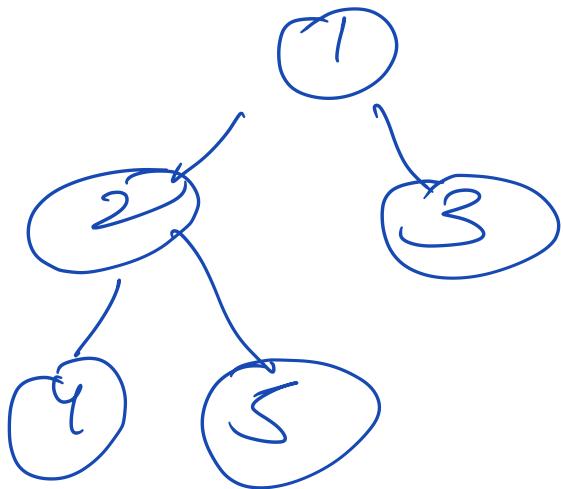
b.) Post Order

First children are taken care
of. before the root



In order Traversal

left \Rightarrow root \Rightarrow right



4 \rightarrow left leaf

2 \rightarrow subroot left

5 \Rightarrow right leaf

1 \Rightarrow root

3 \Rightarrow subroot right

Time Complexity of Trees.

def numnodes(π):
 base case

a = numnode(left)

b = numnodes(right)

return $1 + a + b$

