



DATA STRUCTURES AND ITS APPLICATIONS

Basic Concept and Definitions: Trees

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Basic Concept and Definitions: Trees

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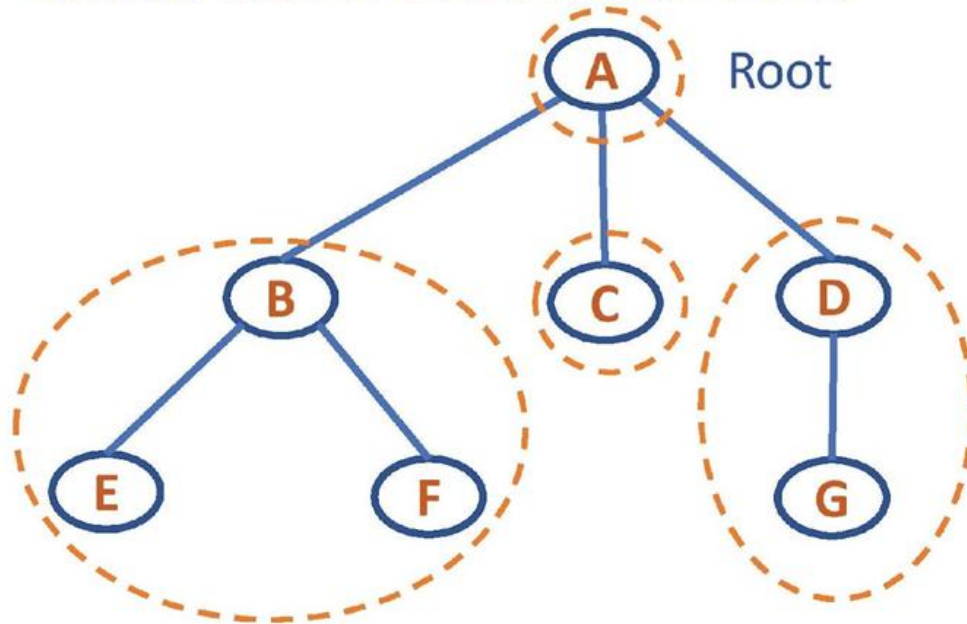
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DATA STRUCTURES AND ITS APPLICATIONS

Trees



- Non Linear Data Structure
- Finite nonempty set of elements
 - One element is the root
 - Remaining elements are partitioned into $m \geq 0$ disjoint subsets each of which is itself a tree

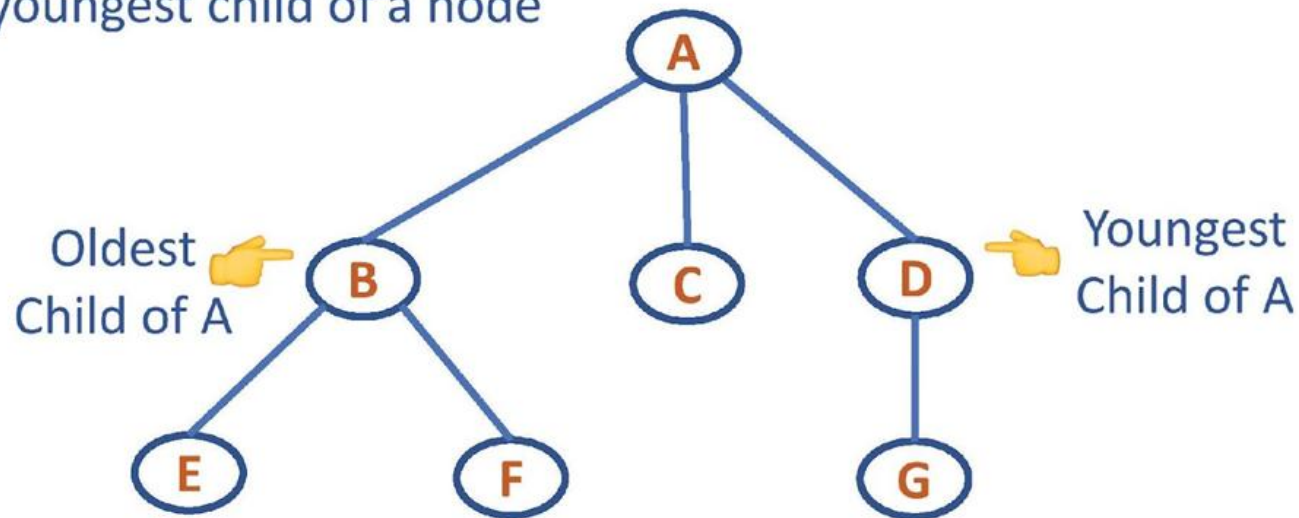


DATA STRUCTURES AND ITS APPLICATIONS

Trees



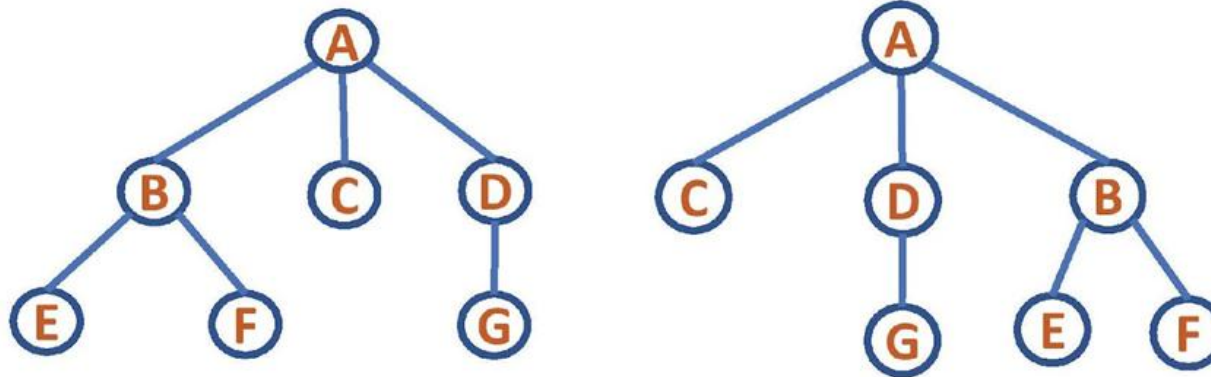
- Ordered Tree: a tree in which subtrees of each node forms an ordered set
- In such a tree we define first, second, ..., Last child of a particular node
- First child is called the oldest child and the last child the youngest child of a node



DATA STRUCTURES AND ITS APPLICATIONS

Trees

As unordered trees the below figures are equivalent but as ordered trees, they are different



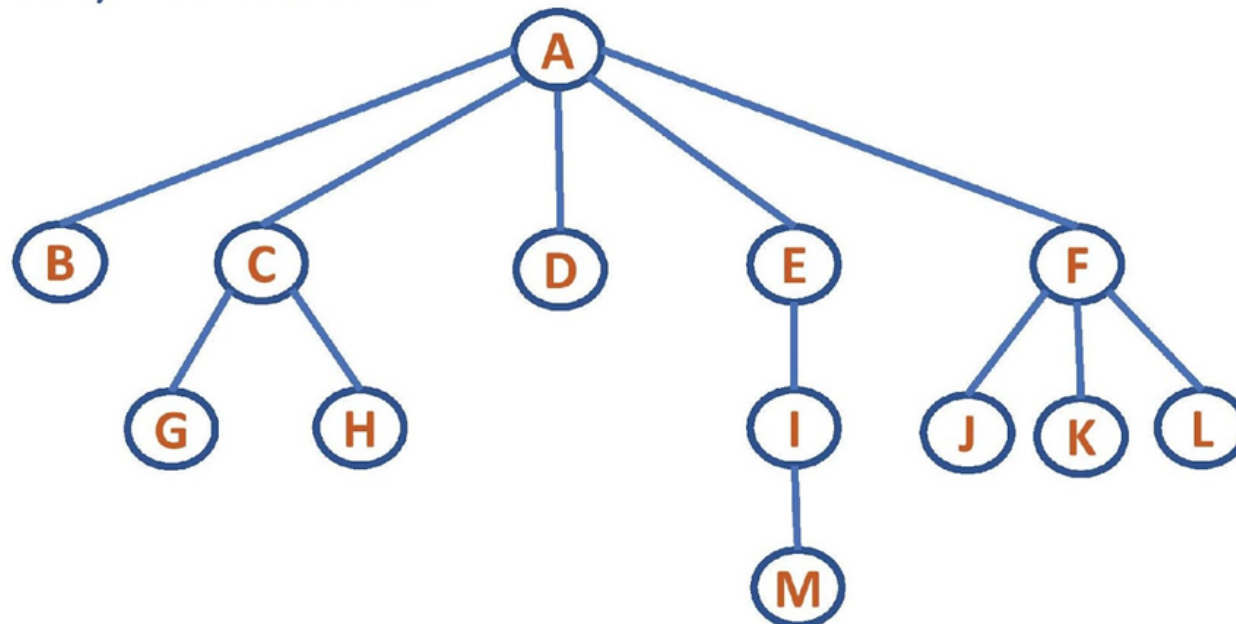
DATA STRUCTURES AND ITS APPLICATIONS

N-ary Tree & Forest



n-ary tree: A rooted tree in which each node has no more than **n** children

- A binary tree is an n-ary tree with $n=2$
- n-ary tree with $n=5$



- Forest: is an ordered set of ordered trees

DATA STRUCTURES AND ITS APPLICATIONS

Tree



Representation of trees:

Tree node options:

```
struct treenode{  
    int info;  
    struct treenode *child[MAX];  
};
```

where MAX is a constant

Restrictions with the above implementation:

- A node cannot have more than MAX children. Therefore cannot expand the tree

DATA STRUCTURES AND ITS APPLICATIONS

Tree



2nd implementation:

- All the children of a given node are linked and only the oldest child is linked to the parent
- A node has link to first child and a link to immediate sibling

```
struct treenode{  
    int info;  
    struct treenode *child;  
    struct treenode *sibling;  
};
```


DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree



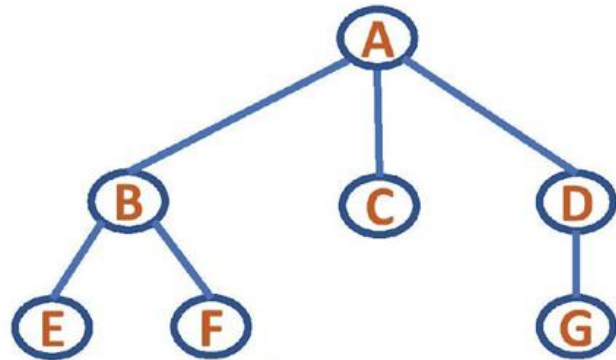
Left Child – Right Sibling Representation

- Link all the siblings of a node
- Delete all links from a node to its children except for the link to its leftmost child
- The left child in binary tree is the node which is the oldest child of the given node in an n-ary tree, and the right child is the node to the immediate right of the given node on the same horizontal line. Such a binary tree will not have a right sub tree
- The node structure corresponds to that of

Data	
Left Child	Right Sibling

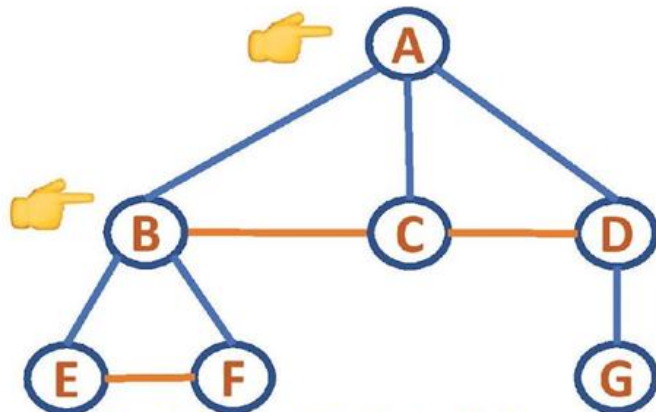
DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree



3-ary tree

Link all siblings of a Node



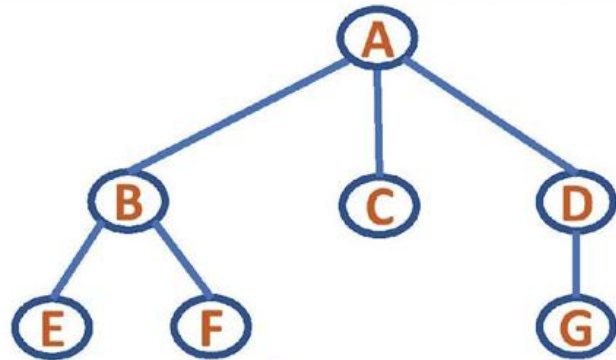
Delete all links from a Node to its children except for the link to its leftmost child

Left child – Right sibling

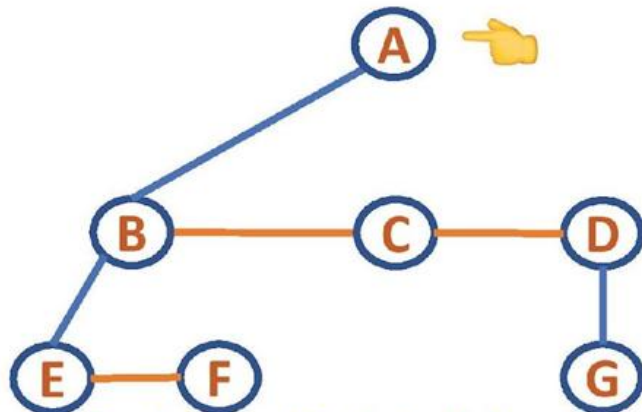
Representation of the above 3-ary tree

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree

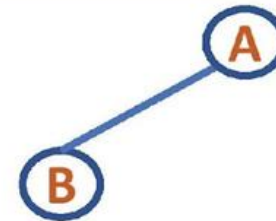


3-ary tree



Left child – Right sibling

Representation of the above 3-ary tree

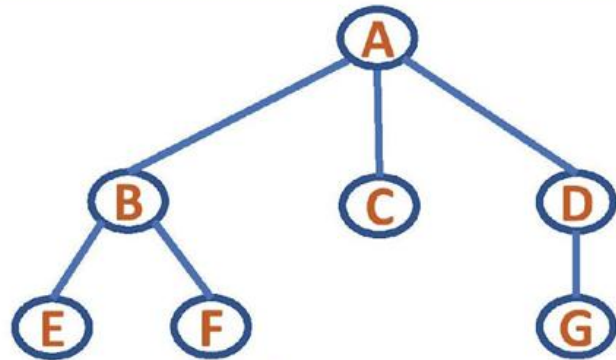


Node Below is B – Left child
No Right sibling so – no Right child

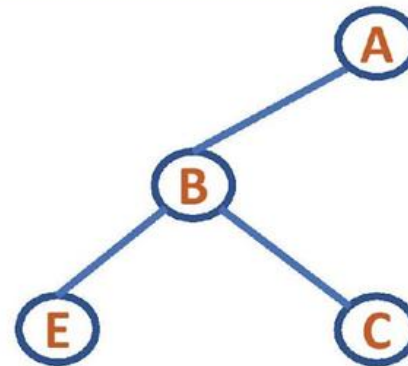
Corresponding
binary tree

DATA STRUCTURES AND ITS APPLICATIONS

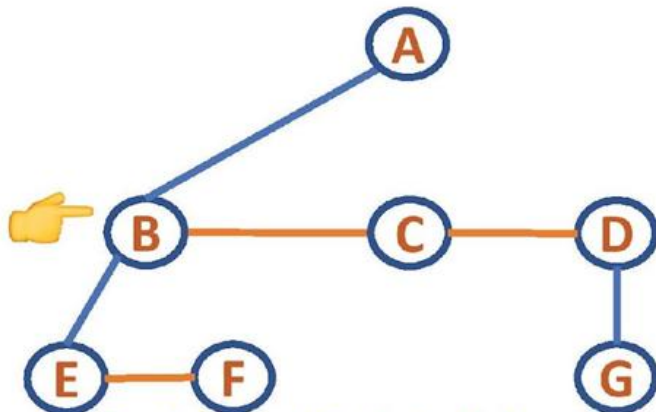
Conversion of an n-ary Tree to a Binary Tree



3-ary tree



Node Below is E – Left child
Next Right sibling is C – Right child



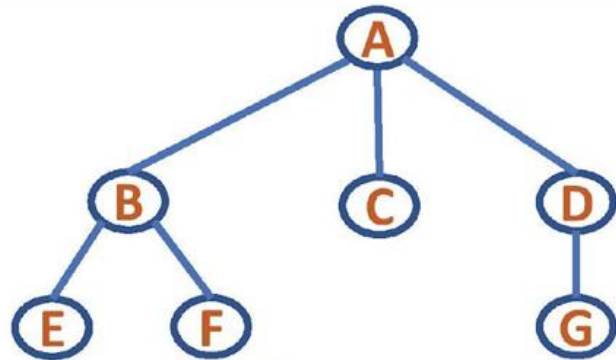
Left child – Right sibling

Representation of the above 3-ary tree

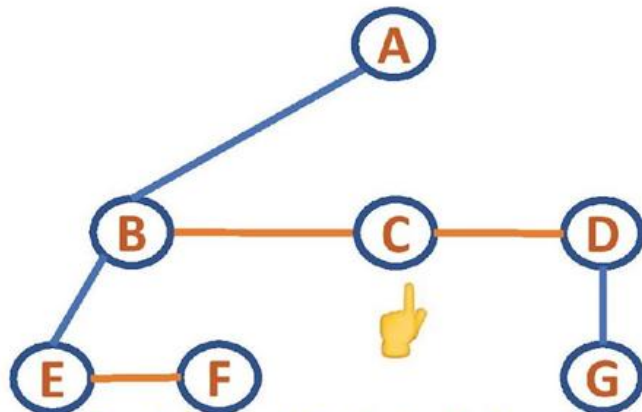
Corresponding
binary tree

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree

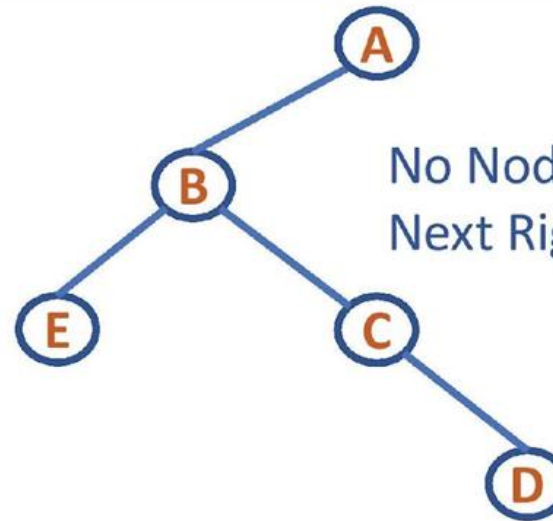


3-ary tree



Left child – Right sibling

Representation of the above 3-ary tree

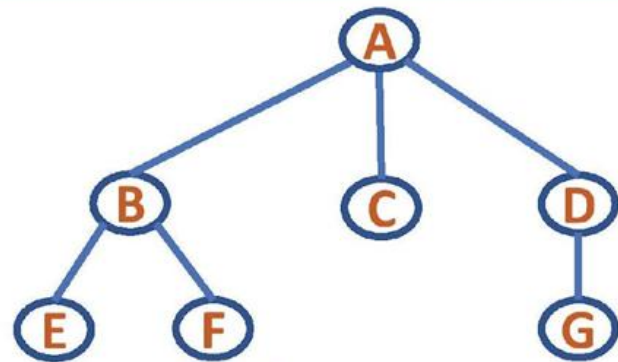


No Node Below so – no Left child
Next Right sibling is D – Right child

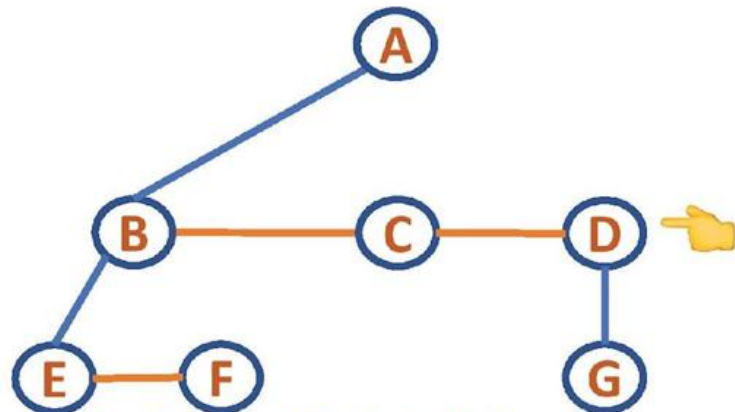
Corresponding
binary tree

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree

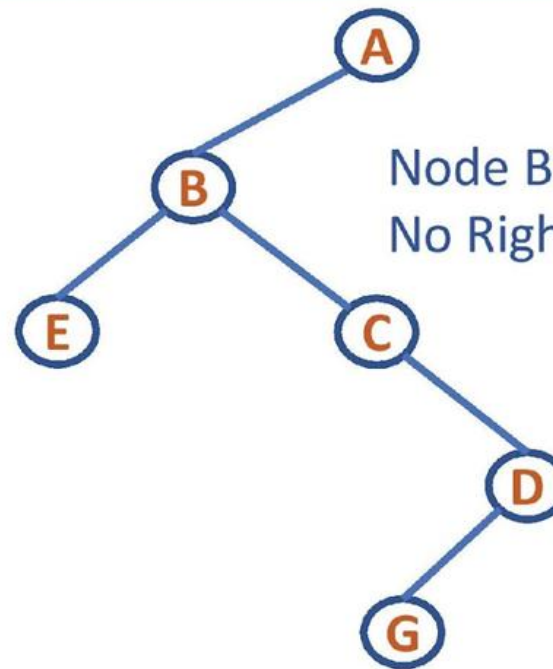


3-ary tree



Left child – Right sibling

Representation of the above 3-ary tree

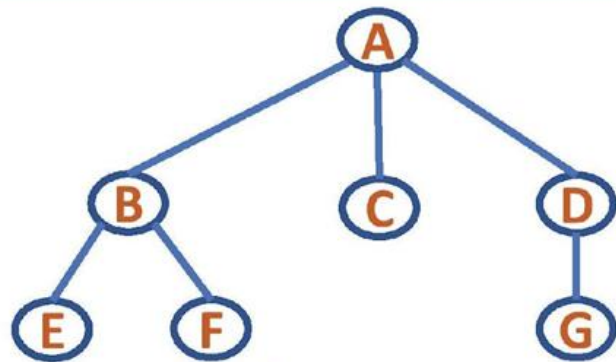


Corresponding
binary tree

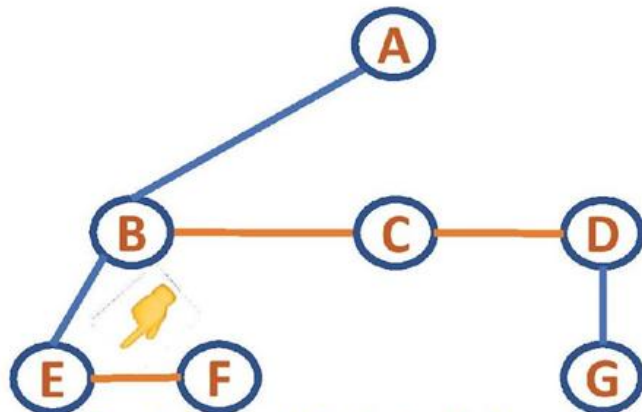
Node Below is G – Left child
No Right sibling so – no Right child

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree

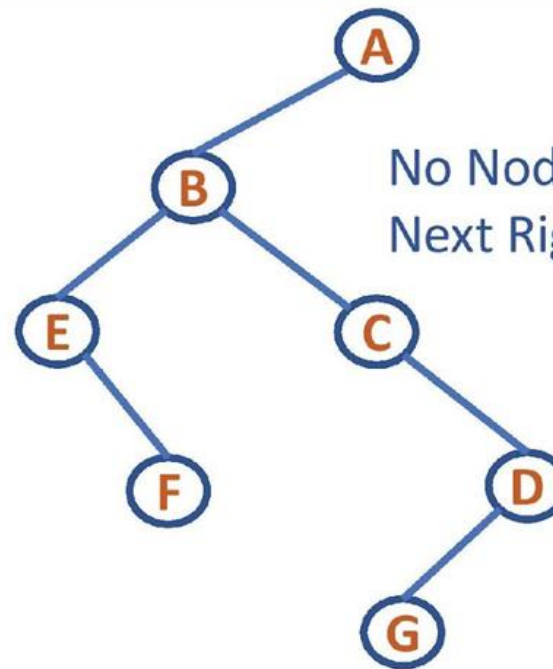


3-ary tree



Left child – Right sibling

Representation of the above 3-ary tree

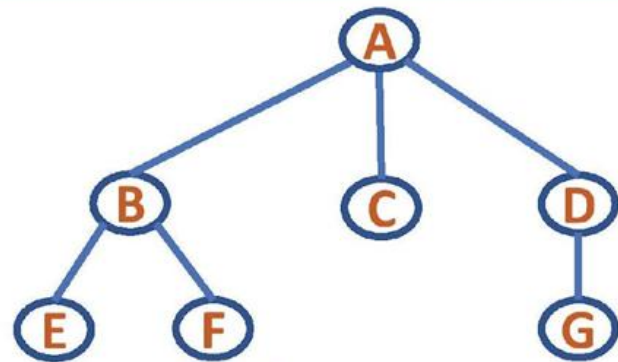


Corresponding
binary tree

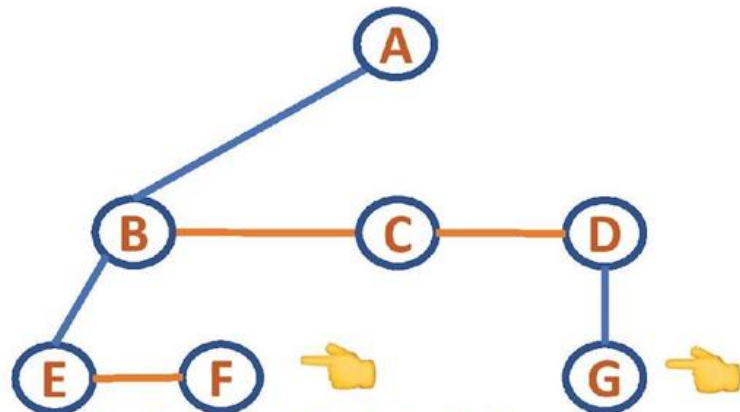
No Node Below so – no Left child
Next Right sibling is F – Right child

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of an n-ary Tree to a Binary Tree

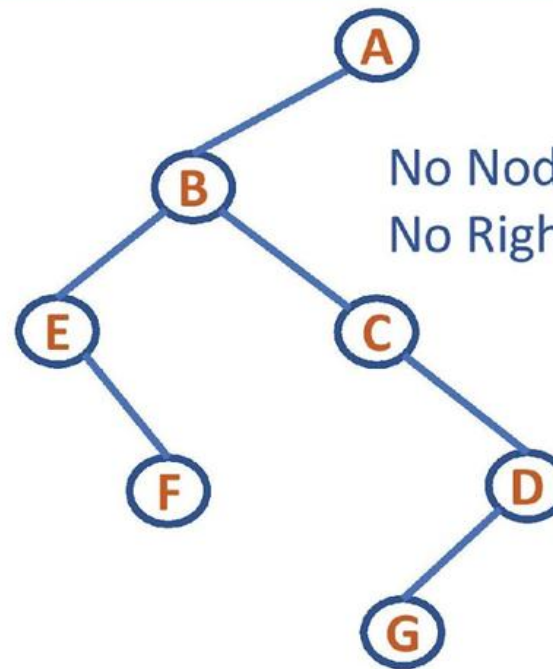


3-ary tree



Left child – Right sibling

Representation of the above 3-ary tree



Corresponding
binary tree

No Node Below so – no Left child
No Right sibling so – no Right child

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of a Forest to a Binary Tree



- Right Child of the root node of every resulting binary tree will be empty. This is because the root of the tree we are transforming has no siblings.
- On the other hand, if we have a forest then these can all be transformed into a single binary tree as follows:
 - First obtain the binary tree representation of each of the trees in the forest
 - Link all the binary trees together through the right sibling field of the root nodes

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of a Forest to a Binary Tree



Conversion of a Forest to a Binary Tree can be formally defined as follows:

- If T_1, \dots, T_n is a forest of n trees, then the binary tree corresponding to this forest, denoted by $B(T_1, \dots, T_n)$:

- is empty if $n = 0$
- has root equal to root (T_1)
- has left subtree equal to $B(T_{11}, T_{12}, \dots, T_{1m})$

where T_{11}, \dots, T_{1m} are the subtrees of root(T_1)

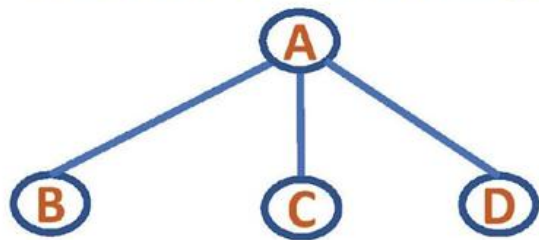
- has right subtree $B(T_2, \dots, T_n)$

DATA STRUCTURES AND ITS APPLICATIONS

Conversion of a Forest to a Binary Tree



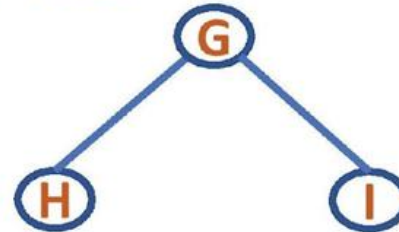
Consider the following Forest with three Trees



T1

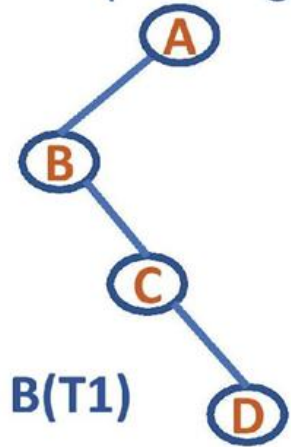


T2

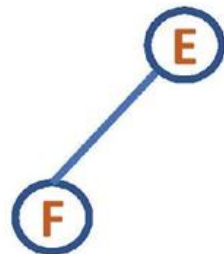


T3

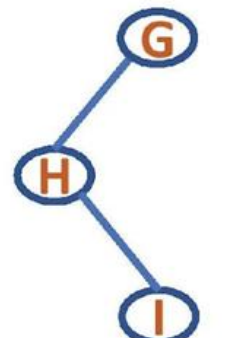
Corresponding Binary Trees



B(T1)



B(T2)



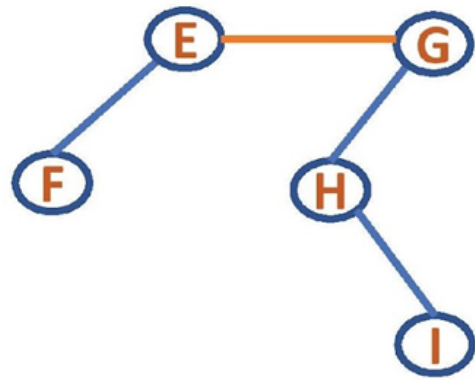
B(T3)

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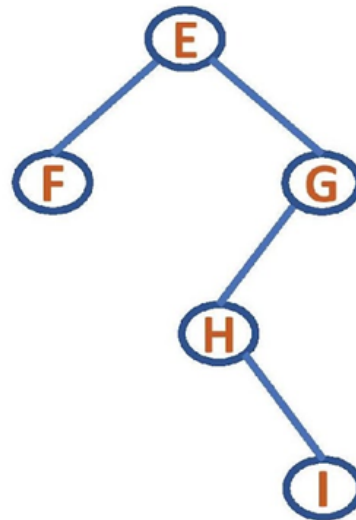
Conversion of a Forest to a Binary Tree



Link $B(T_2)$ and $B(T_3)$



G becomes Right Child of E



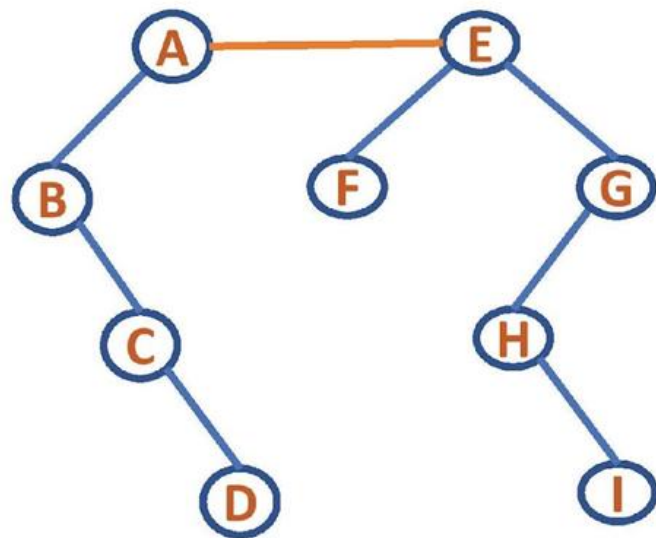
Corresponding Binary Tree
 $B(T_2, T_3)$

DATA STRUCTURES AND ITS APPLICATIONS

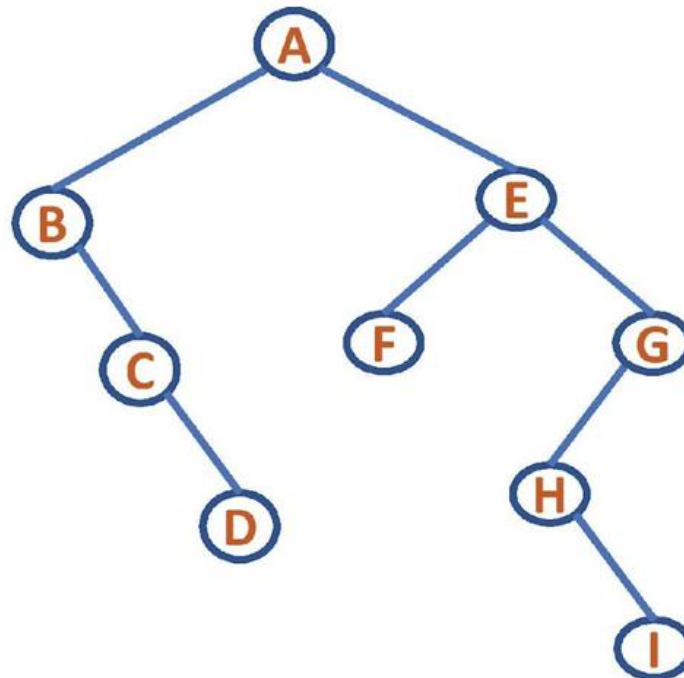
Conversion of a Forest to a Binary Tree



Link B(T1) and B(T2, T3)



E becomes Right Child of A



Corresponding Binary Tree B(T1, T2, T3)

1. A forest can be defined as:

- a) A set of disconnected graphs
- b) A collection of disjoint trees
- c) A binary tree with multiple roots
- d) A weighted tree

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2. If an N-ary tree has n nodes, the maximum number of edges is:

- a) n
- b) $n - 1$
- c) $n + 1$
- d) n^2

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3. Which of the following is true about the relationship between a tree and a forest?

- a) A tree can be converted into a forest by removing its root
- b) A forest can be converted into a tree by adding a super root
- c) Both (a) and (b)
- d) Neither (a) nor (b)

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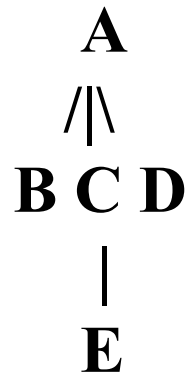
4. In an N-ary tree, the maximum number of children any node can have is:

- a) N
- b) $N - 1$
- c) $\log(N)$
- d) $2N$

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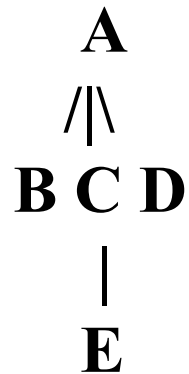
5. Consider the tree below:



If we convert this tree into a forest by removing the root A, the resulting forest will contain:

- a) 2 trees: (B) and (C–D–E)
- b) 3 trees: (B), (C–E), (D)
- c) 3 trees: (B), (C), (D–E)
- d) 4 trees: (A), (B), (C), (D–E)

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If we convert this tree into a forest by removing the root A, the resulting forest will contain:

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- c) 3 trees: (B), (C), (D–E)
- d) 4 trees: (A), (B), (C), (D–E)



THANK YOU

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