



# Data Structures & Its Applications

---

**Kusuma K V**

Department of Computer Science & Engineering

# DATA STRUCTURES & ITS APPLICATIONS

---

## UNIT 1: Skip List

**Kusuma K V**

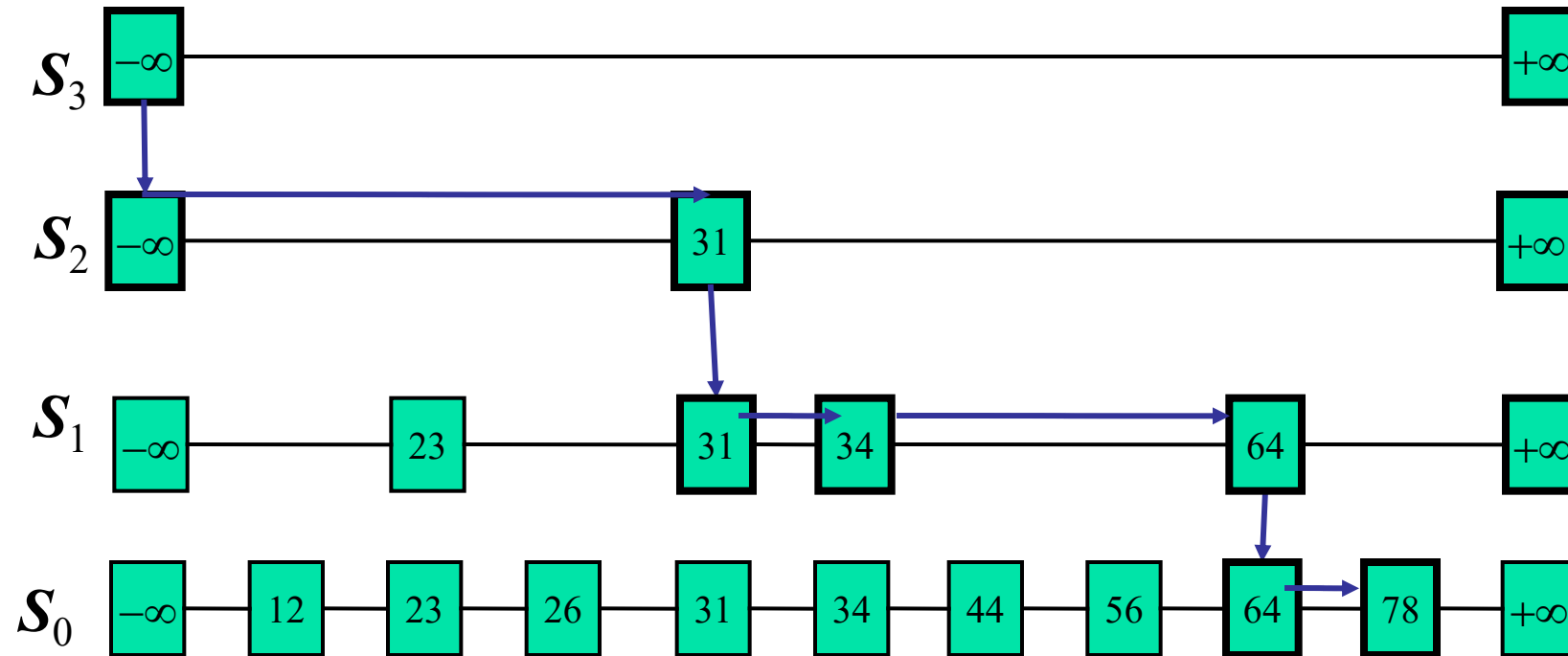
Department of Computer Science & Engineering

### Search

We search for a key  $x$  in a skip list as follows:

- We start at the first position of the top list
- At the current position  $p$ , we compare  $x$  with  $y$  i.e.,  $\text{key}(\text{after}(p))$ 
  - $x = y$ : we return  $\text{element}(\text{after}(p))$
  - $x > y$ : we “scan forward”
  - $x < y$ : we “drop down”
- If we try to drop down past the bottom list, we return `NO_SUCH_KEY`
- Example: search for 78

### Searching in Skip List Example



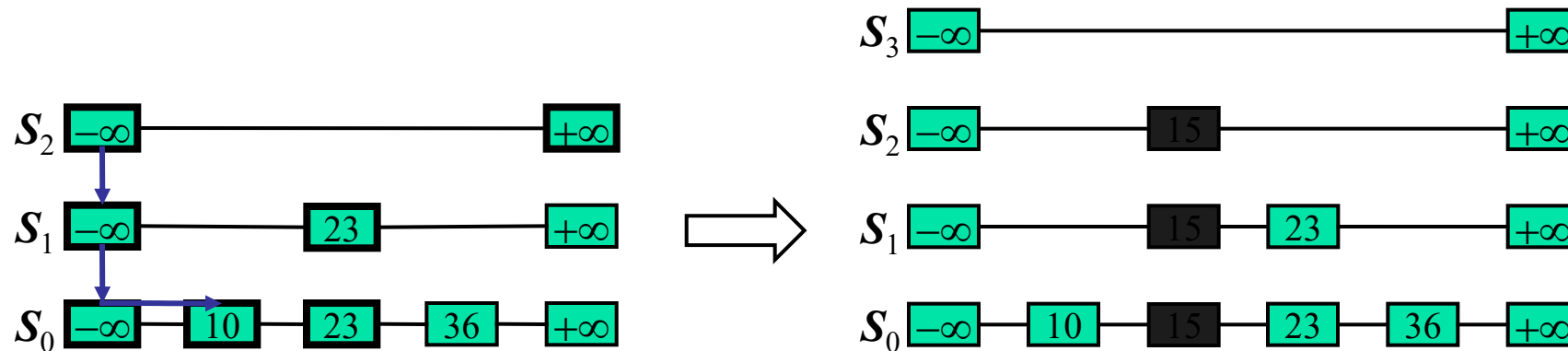
### Insertion

- The insertion algorithm for skip lists uses **randomization** to decide how many references to the new item should be added to the skip list
- We then insert the new item in this bottom-level list at its appropriate position. After inserting the new item at this level we “flip a coin”
  - If the flip comes up tails, then we stop right there.
  - If the flip comes up heads, we move to next higher level and insert in this level at the appropriate position

- A randomized algorithm performs coin tosses (i.e., uses random bits) to control its execution
- Its running time depends on the outcomes of the coin tosses
- We analyze the expected running time of a randomized algorithm under the following assumptions
  - the coins are unbiased, and
  - the coin tosses are independent
- The worst-case running time of a randomized algorithm is large but has very low probability (e.g., it occurs when all the coin tosses give “heads”)

### Insertion in Skip List Example

- Suppose we want to insert 15
- Do a search, and find the spot between 10 and 23
- Suppose the coin comes up “head” two times



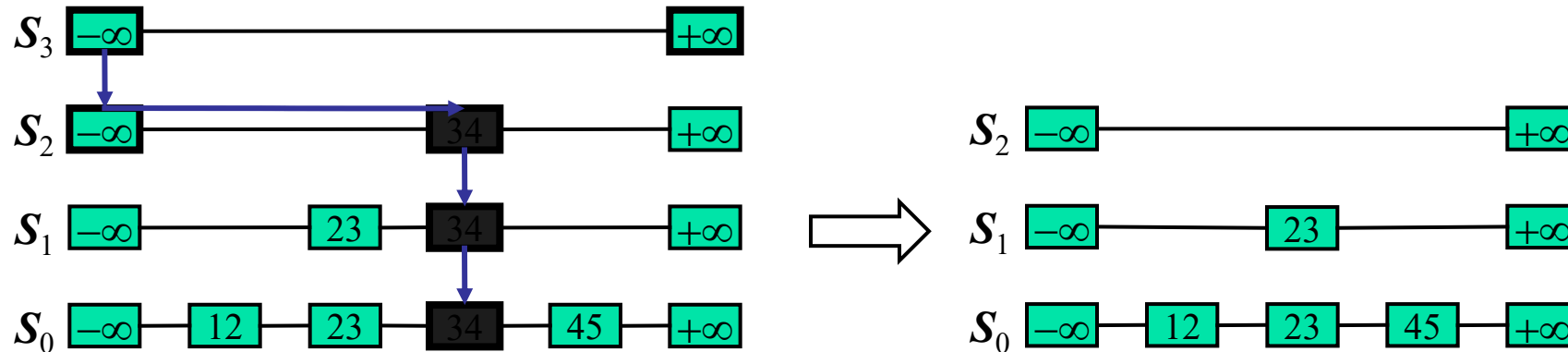
### Deletion

- We begin by performing a search for the given key  $k$ .
- If a position  $p$  with key  $k$  is not found, then we return the NO\_SUCH\_KEY element
- Otherwise, if a position  $p$  with key  $k$  is found (it would be found on the bottom level), then we remove all the position above  $p$
- If more than one upper level is empty, remove it



### Deletion in Skip List Example

- Suppose we want to delete 34
- Do a search, find the spot between 23 and 45
- Remove all the position above p

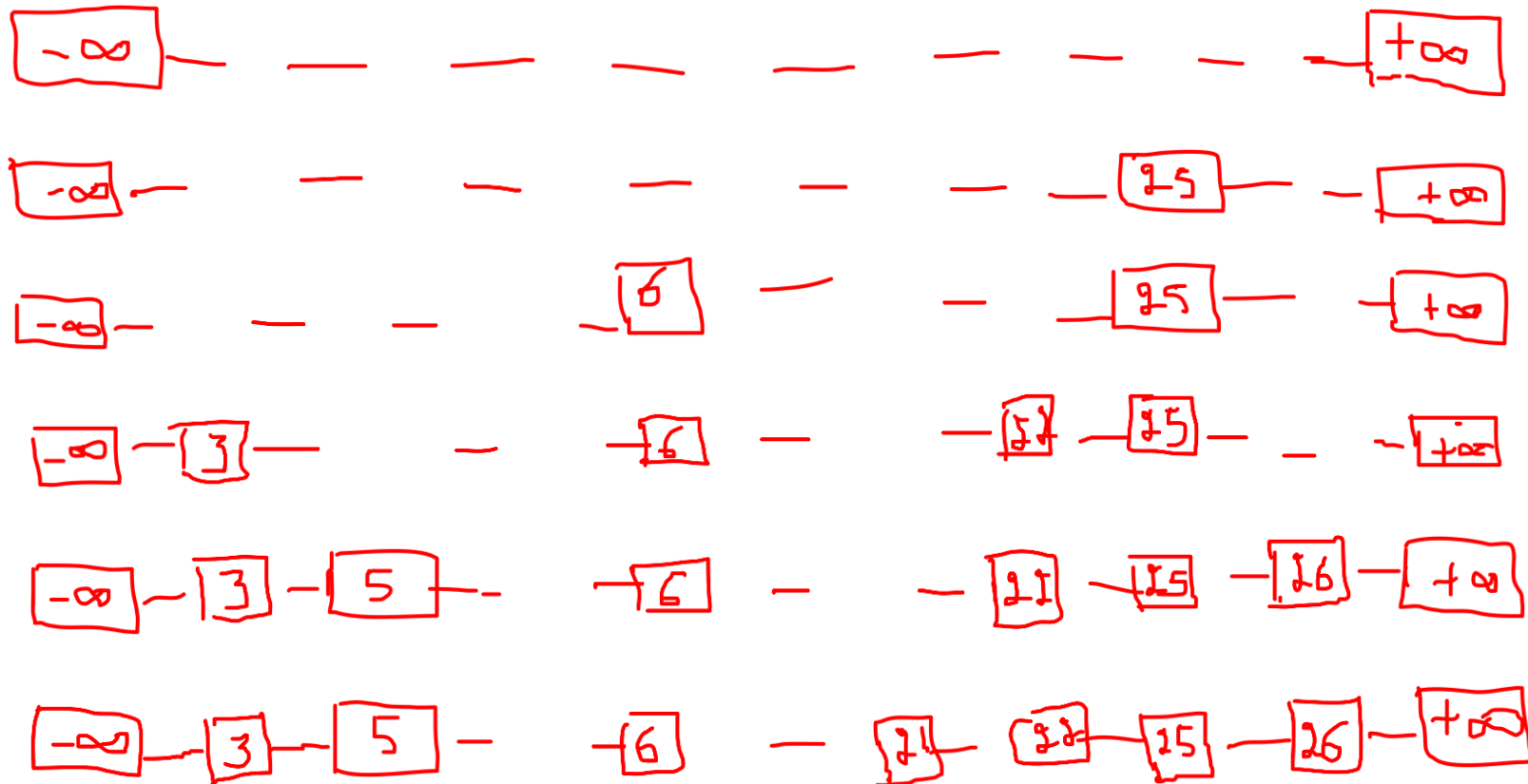


Starting with an empty skip list, insert these keys, with these “randomly generated” levels

- 5, with level 1
- 26, with level 1
- 25, with level 4
- 6, with level 3
- 21, with level 0
- 3, with level 2
- 22, with level 2

Note that the ordering of the keys in the skip list is determined by the value of the keys only; the levels of the nodes are determined by the random number generator only

5 with level 1, 26 with level 1, 25 with level 4, 6 with level 3  
21 with level 0, 3 with level 2, 22 with level 2



- Presentation Slides: Advanced Data Structures  
Dr. RamaMoorthy Srinath, Professor, CSE, PESU

**1. While searching for key K in a skip list, what is the correct rule for horizontal movement at the current level?**

- a) Move right while  $\text{next.key} \leq K$ .
- b) Move right while  $\text{next.key} < K$ ; drop down when  $\text{next.key} \geq K$  or  $\text{next} == \text{NULL}$ .
- c) Always move right exactly once per level.
- d) Drop down first, then move right to find K.

1. While searching for key K in a skip list, what is the correct rule for horizontal movement at the current level?
  - a) Move right while `next.key <= K`.
  - b) Move right while `next.key < K`; drop down when `next.key >= K` or `next == NULL`.
  - c) Always move right exactly once per level.
  - d) Drop down first, then move right to find K.

### 2. Which sequence best describes inserting a new key K?

- a) Random level → Allocate node → Search path → Link node in found level only.
- b) Search path (collect update[]) → Random level for node → Allocate node → Splice node into all levels  $\leq$  nodeLevel.
- c) Allocate node → Splice at level 0 → Randomly add higher links later.
- d) Search path → Allocate node → Splice only at top level.

### 2. Which sequence best describes inserting a new key K?

- a) Random level → Allocate node → Search path → Link node in found level only.
- b) Search path (collect update[]) → Random level for node → Allocate node → Splice node into all levels  $\leq$  nodeLevel.
- c) Allocate node → Splice at level 0 → Randomly add higher links later.
- d) Search path → Allocate node → Splice only at top level.



**3. How is the level (height) of a newly inserted node typically determined?**

- a) Fixed at 1 for all nodes.
- b) Count set bits in key value.
- c) Repeatedly “flip a coin” with probability  $p$  of going up a level until fail or max level reached.
- d) Proportional to current list size ( $\lceil \log_2 n \rceil$ ).

**3. How is the level (height) of a newly inserted node typically determined?**

- a) Fixed at 1 for all nodes.
- b) Count set bits in key value.
- c) Repeatedly “flip a coin” with probability  $p$  of going up a level until fail or max level reached.
- d) Proportional to current list size ( $\lceil \log_2 n \rceil$ ).

**4. To delete key K, which approach is correct?**

- a) Search only level 0; unlink node; done.
- b) Search from top without storing predecessors; when found, drop to level 0 and delete.
- c) Search building update[] of last nodes  $< K$  at each level; if found, relink  $\text{update}[i] \rightarrow \text{forward}[i]$  to skip target at all involved levels; free node.
- d) Mark node deleted; let garbage collector collapse levels.

**4. To delete key K, which approach is correct?**

- a) Search only level 0; unlink node; done.
- b) Search from top without storing predecessors; when found, drop to level 0 and delete.
- c) Search building update[] of last nodes  $< K$  at each level; if found, relink  $update[i] \rightarrow forward[i]$  to skip target at all involved levels; free node.
- d) Mark node deleted; let garbage collector collapse levels.

**5. After deleting a node, when (if ever) should the current highest level of the skip list be reduced?**

- a) Never; levels are fixed after first insertion.
- b) Whenever any node is deleted.
- c) When the header's forward pointer at the top level becomes NULL (no elements at that level).
- d) When  $\text{current size} < \text{previous size}/2$ .

**5. After deleting a node, when (if ever) should the current highest level of the skip list be reduced?**

- a) Never; levels are fixed after first insertion.
- b) Whenever any node is deleted.
- c) When the header's forward pointer at the top level becomes NULL (no elements at that level).
- d) When  $\text{current size} < \text{previous size}/2$ .



**THANK YOU**

---

**Kusuma K V**

Department of Computer Science & Engineering

**[kusumakv@pes.edu](mailto:kusumakv@pes.edu)**