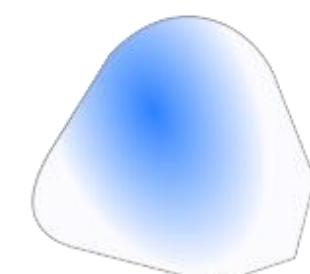
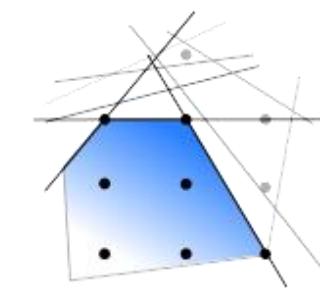
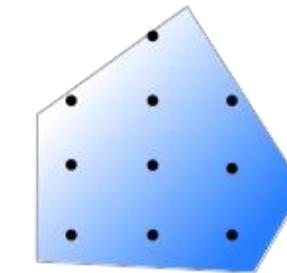
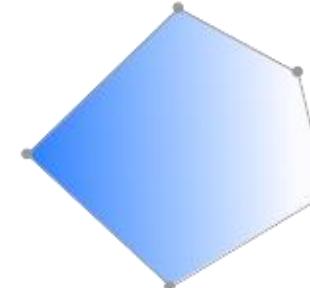
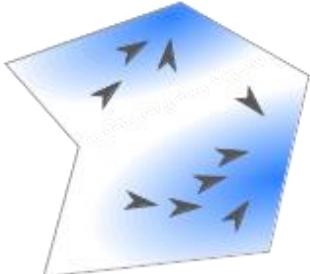
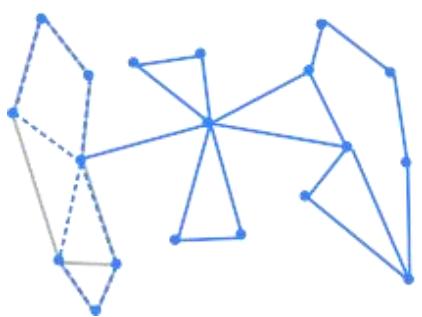
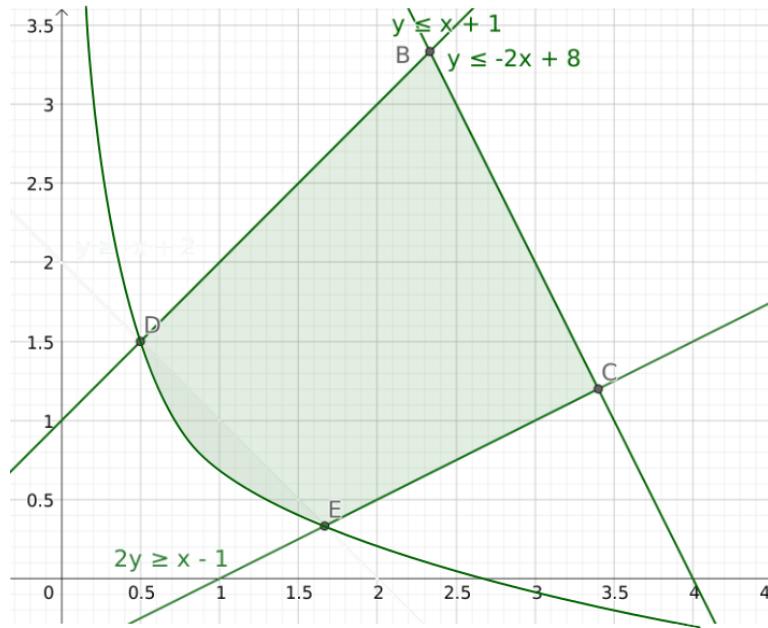
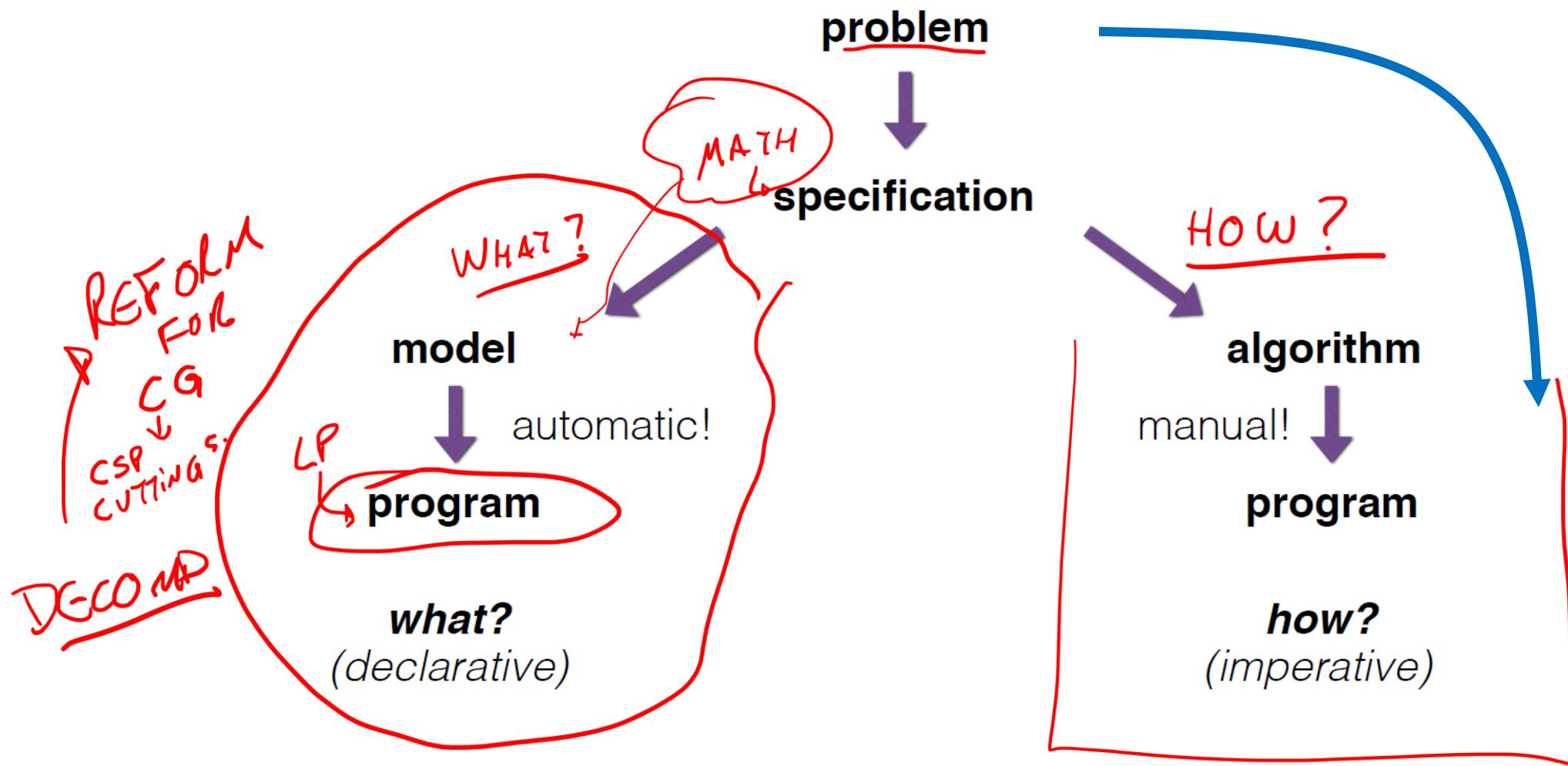


Construction Algorithms

COMP4691 / 8691



Previously on COMP4691(8691)



Previously on COMP4691(8691)

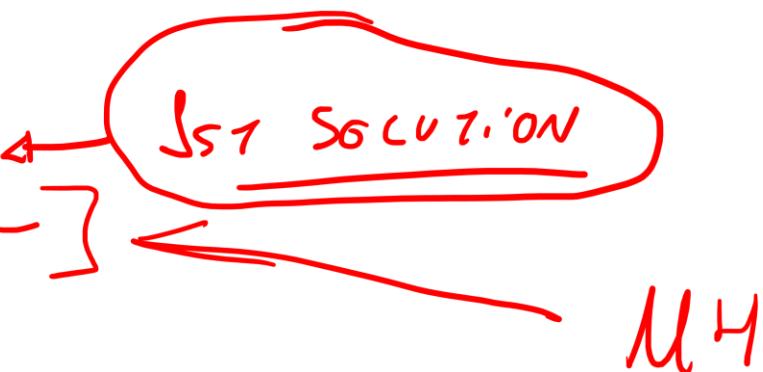
- Looked at a model-based approach to solving (MIPs)

Now

- Look at heuristic approaches

A classic

- Construct
- Improve



Today:

- Construct

SEARCHING
→ LNS
LARGE NEIGH SEARCH

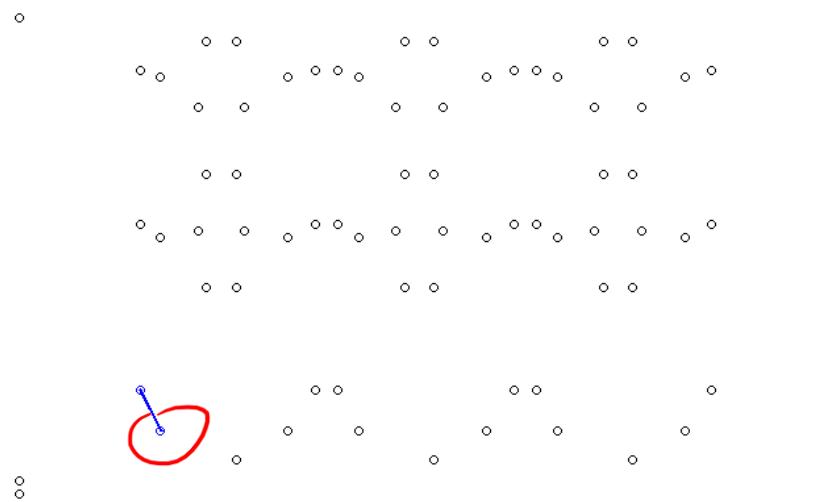
Construction algorithms

We need to start somewhere

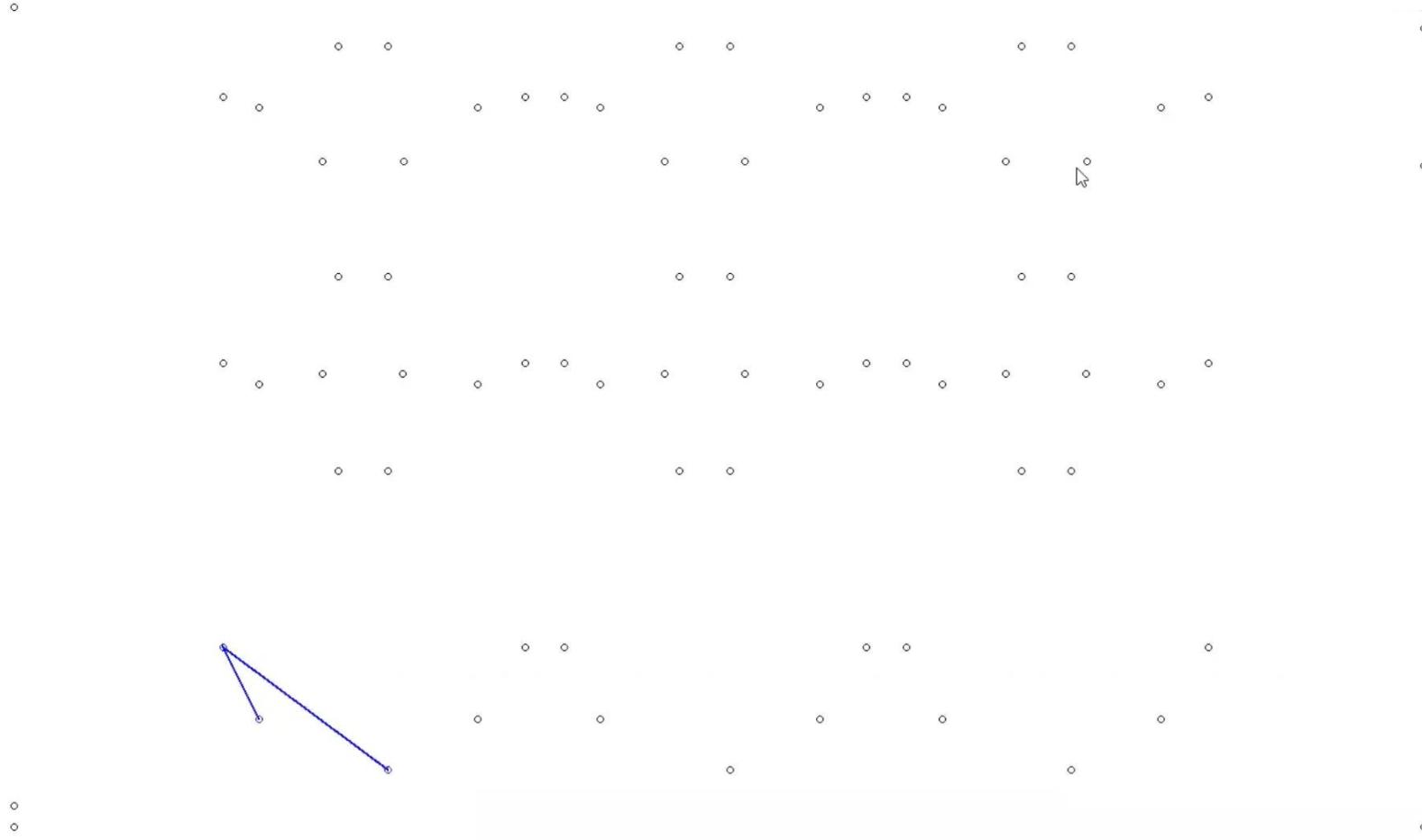
- When we first come to an instance we are looking at an empty page
 - In TSP, no cities are visited
 - In Knapsack, all the items are lying on the ground
- Initial solution often created by “constructing” one piece / element at a time

Construction algorithms – Nearest Neighbour

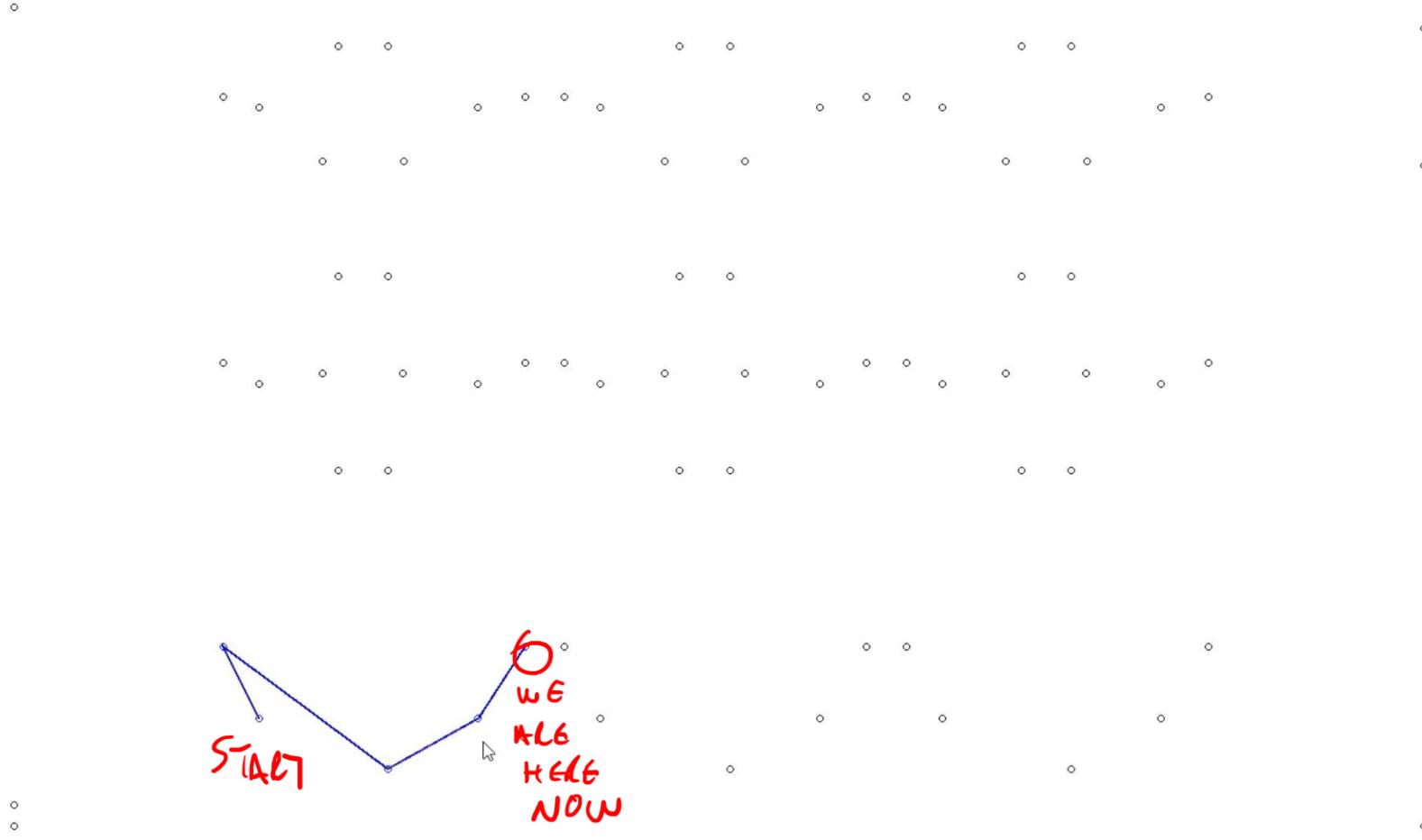
- For example: TSP
 - Starting with an empty tour, start in the “home” city
 - Repeat:
 - Add the city closest to the last-added city



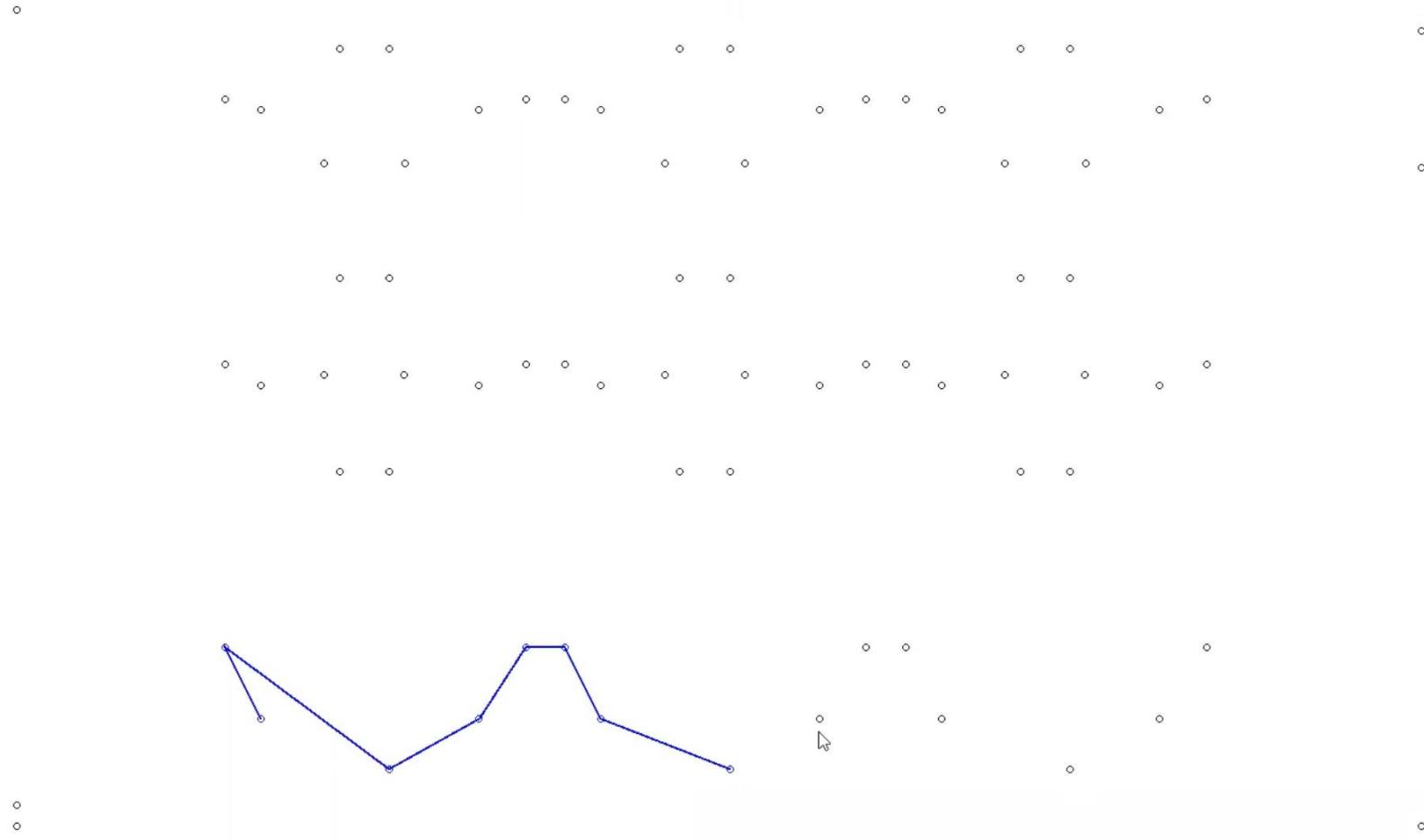
Construction algorithms – Greedy



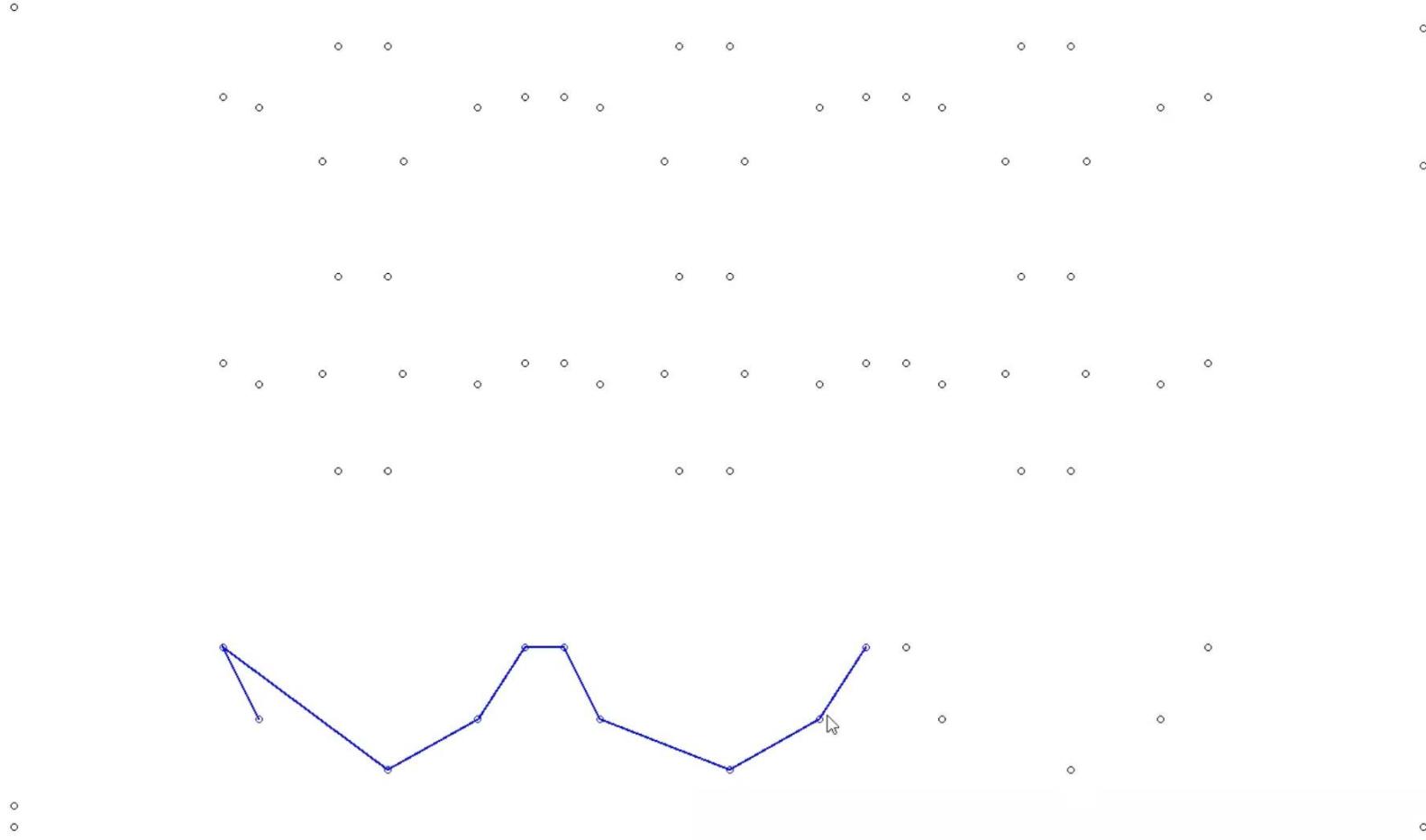
Construction algorithms – Greedy



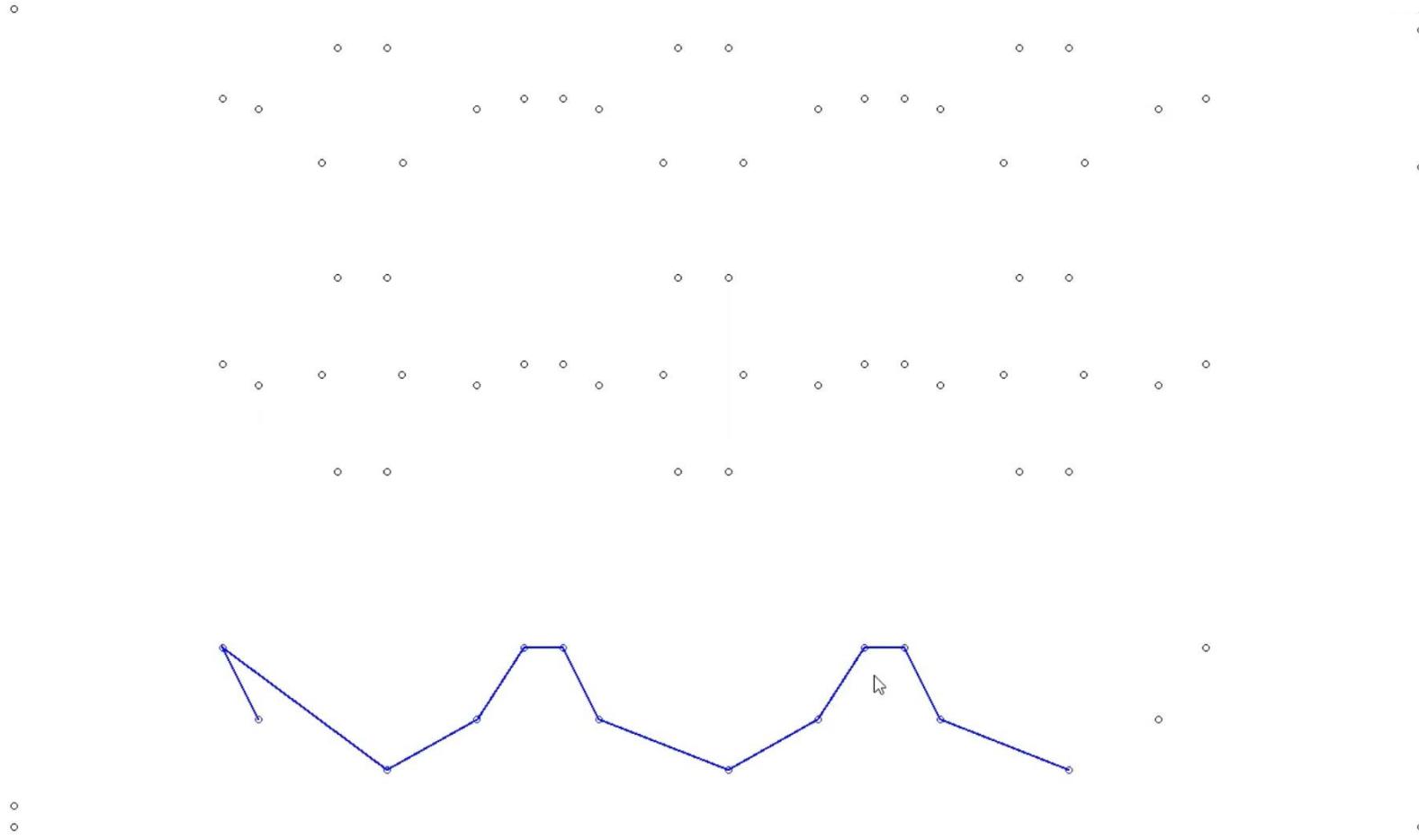
Construction algorithms – Greedy



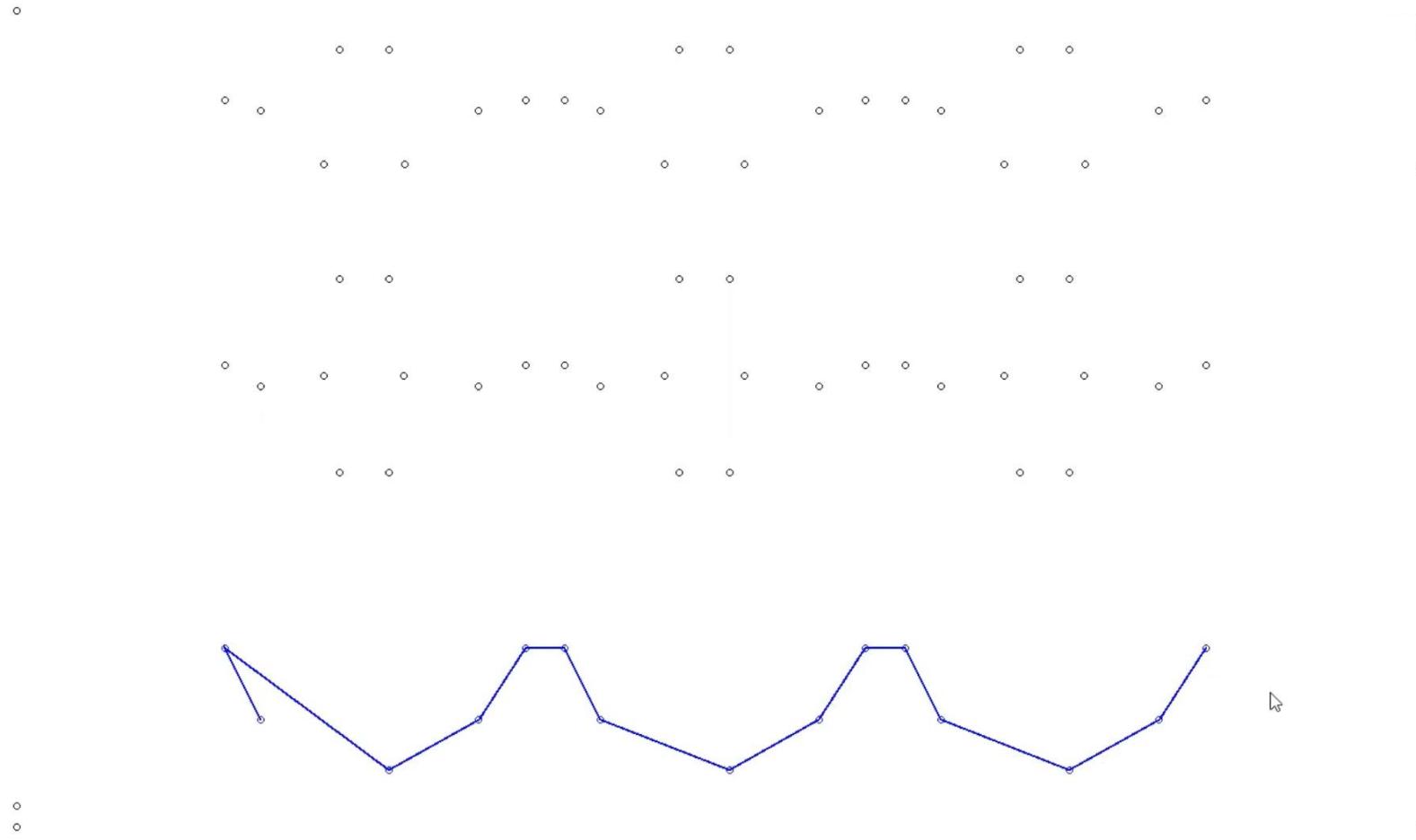
Construction algorithms – Greedy



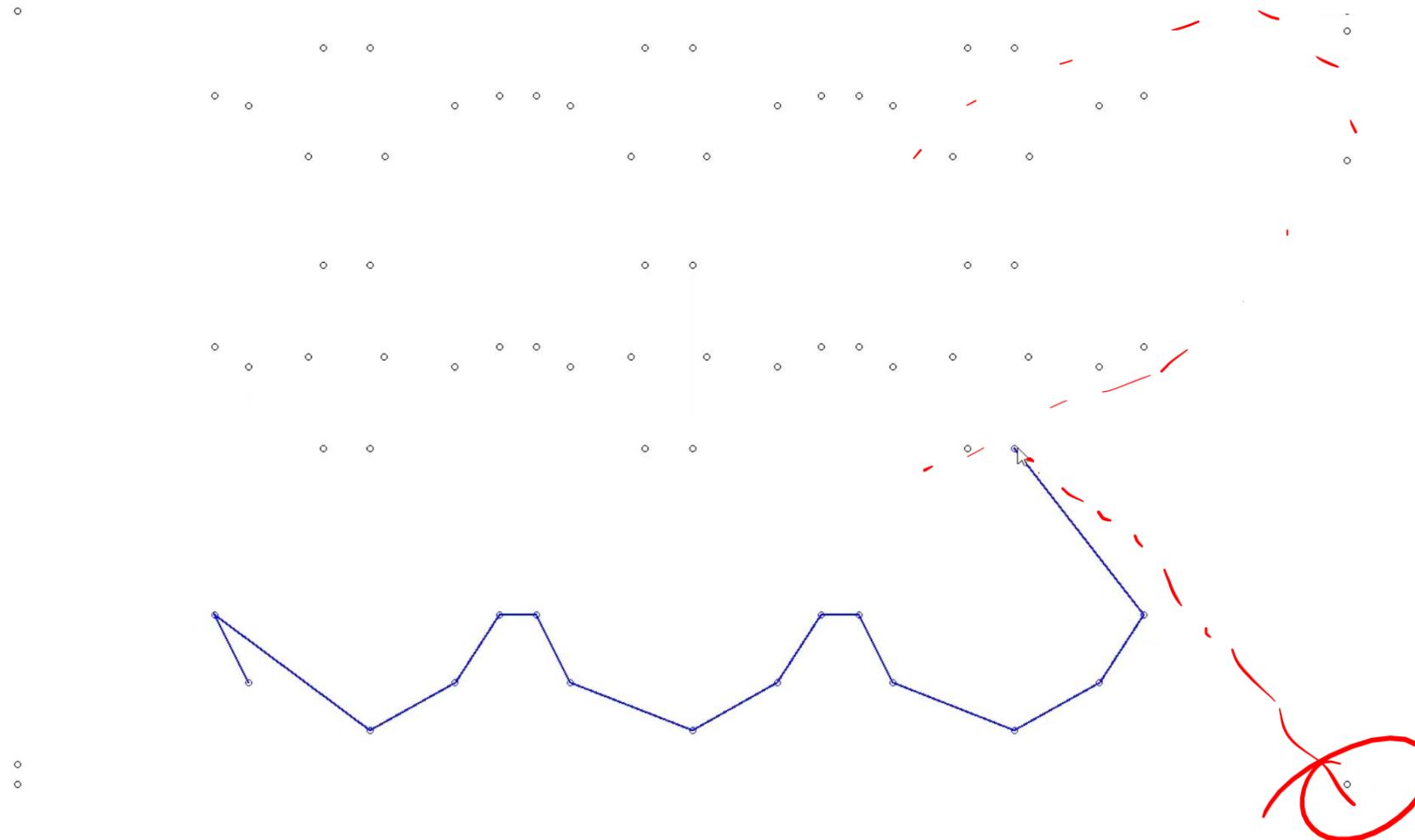
Construction algorithms – Greedy



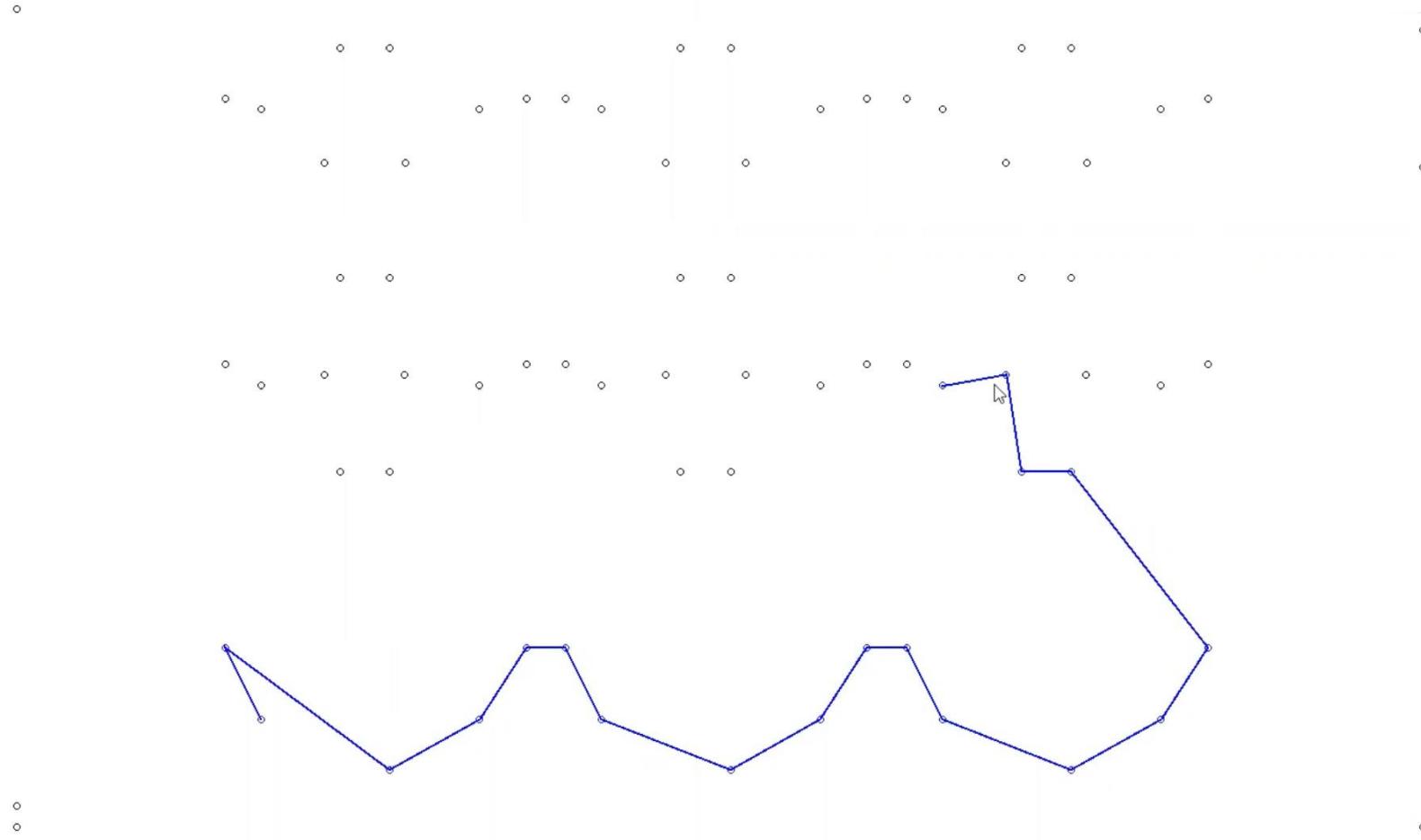
Construction algorithms – Greedy



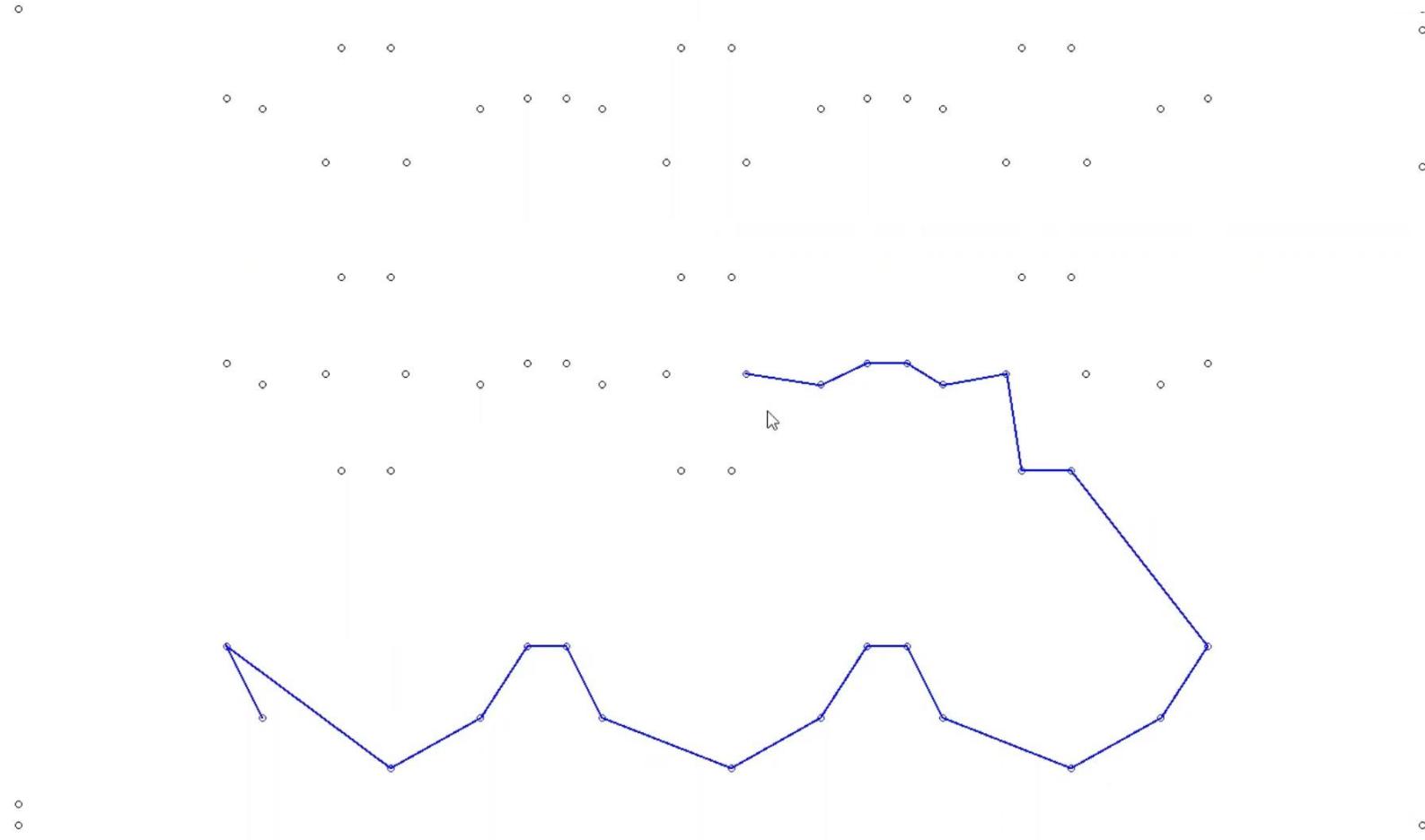
Construction algorithms – Greedy



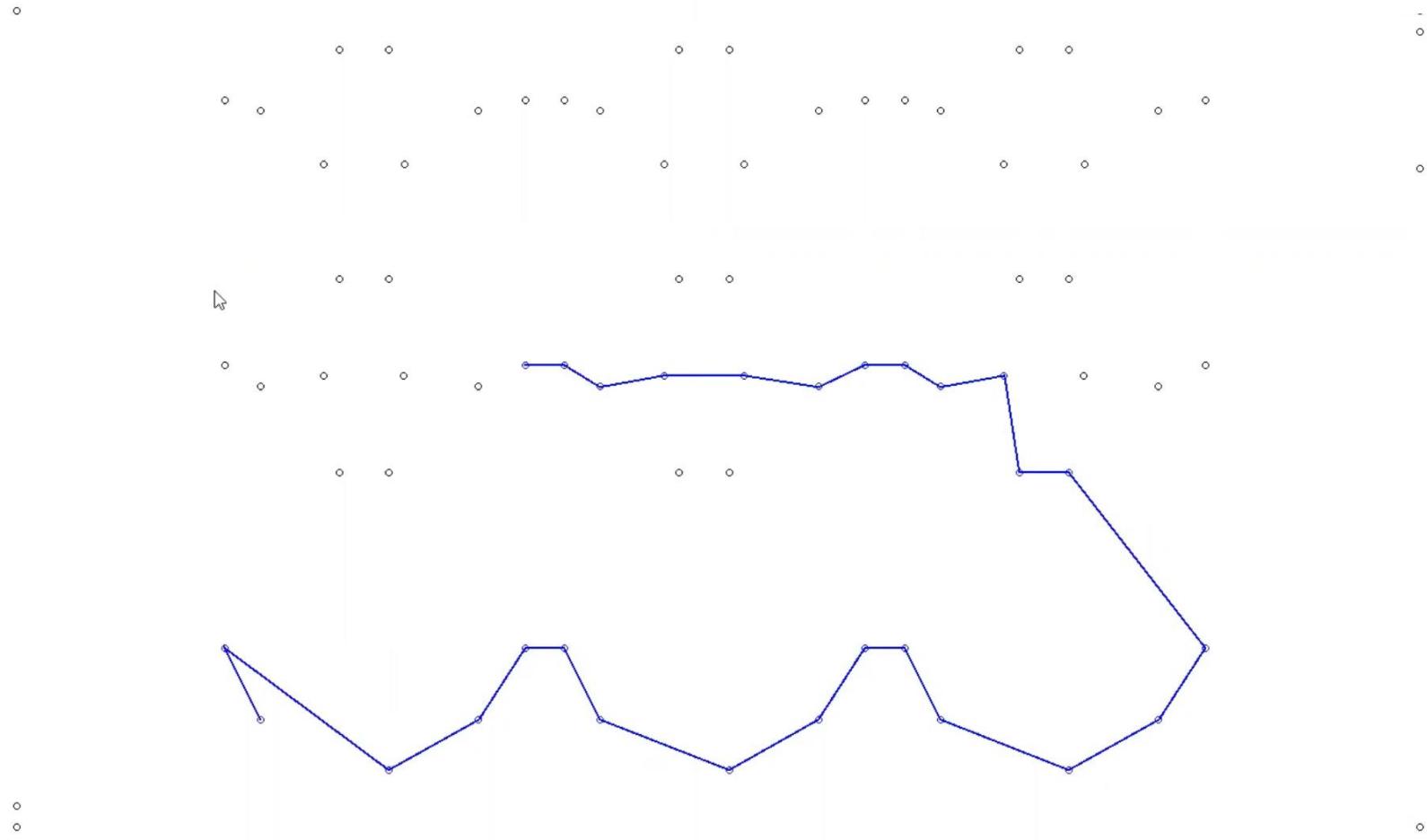
Construction algorithms – Greedy



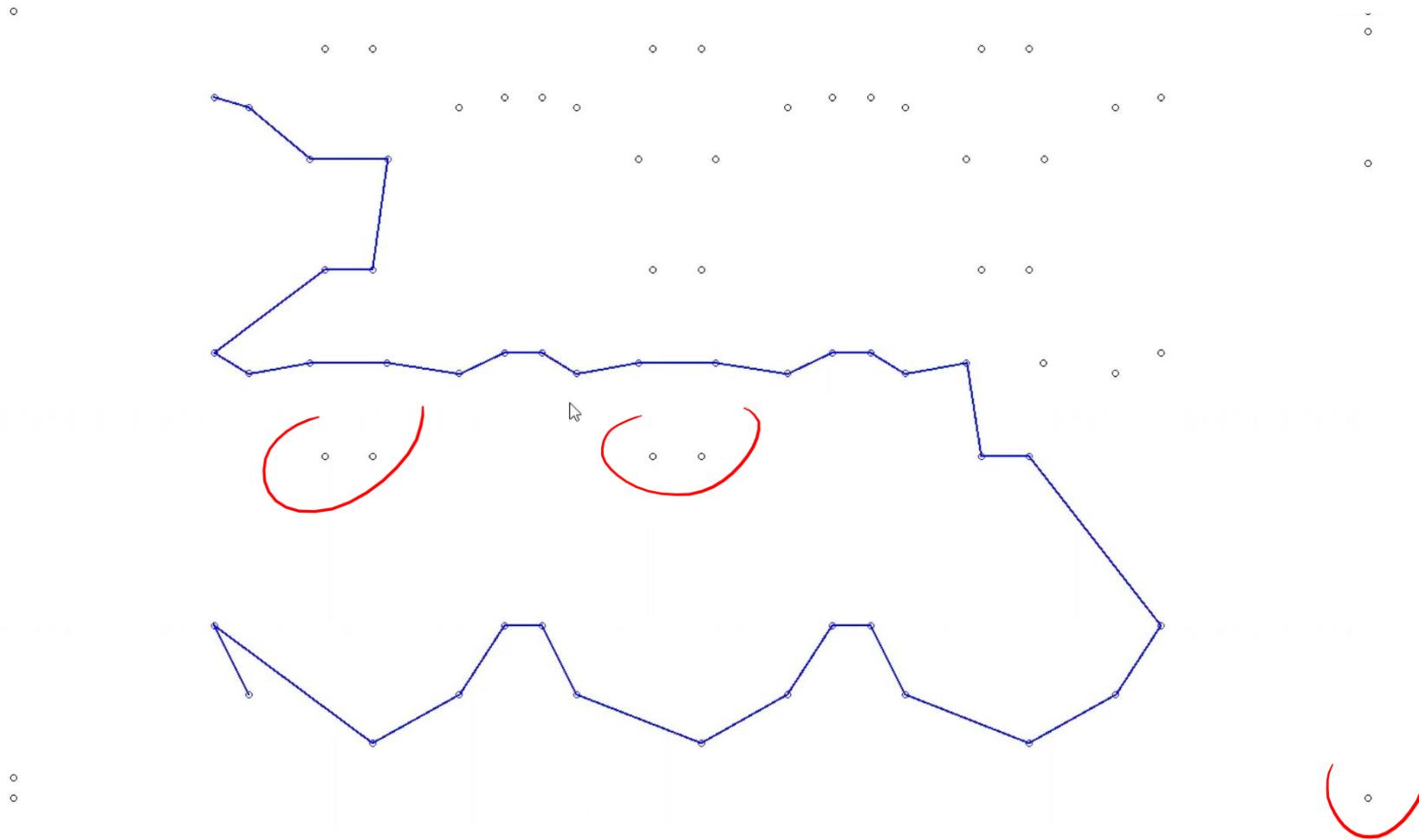
Construction algorithms – Greedy



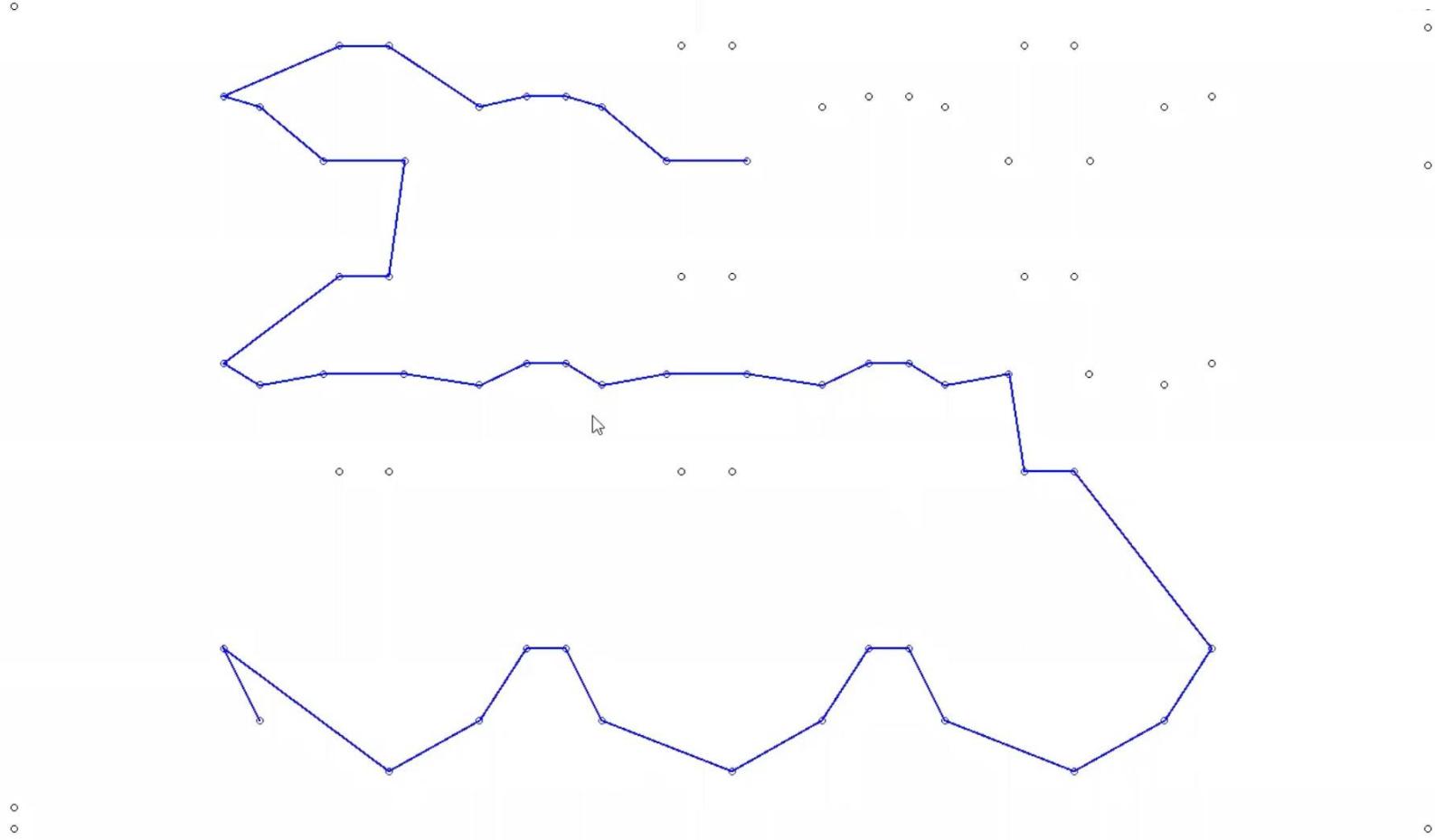
Construction algorithms – Greedy



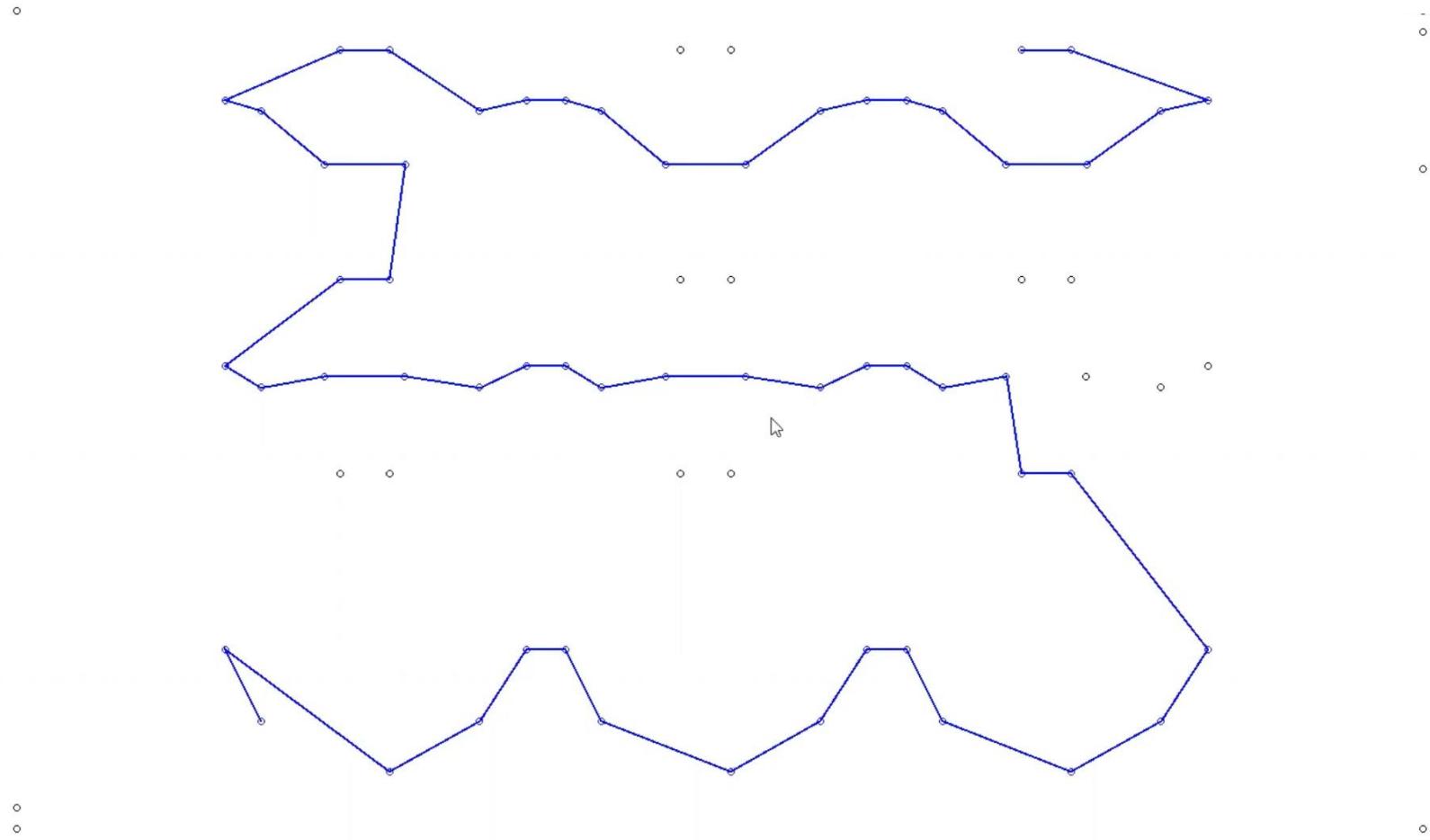
Construction algorithms – Greedy



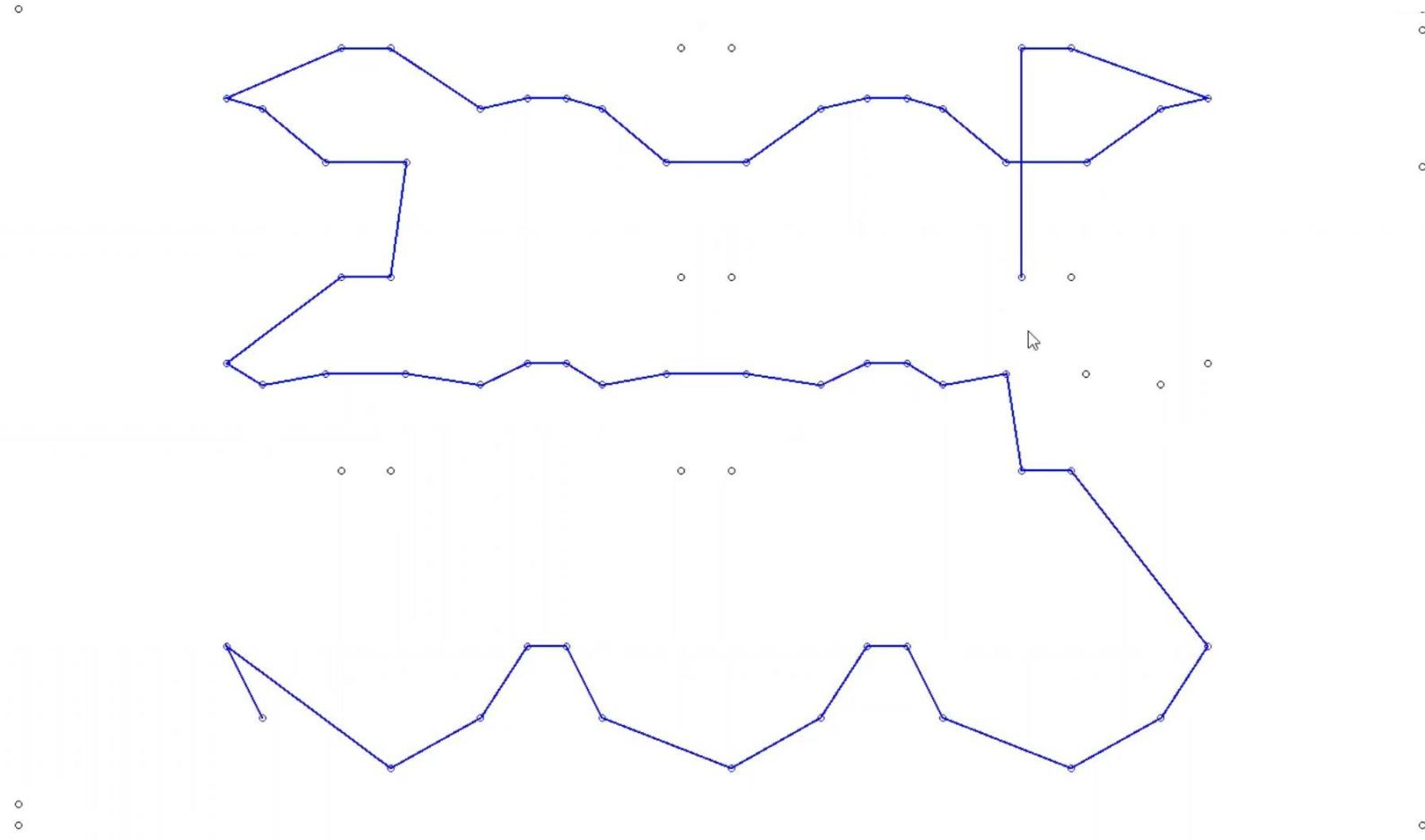
Construction algorithms – Greedy



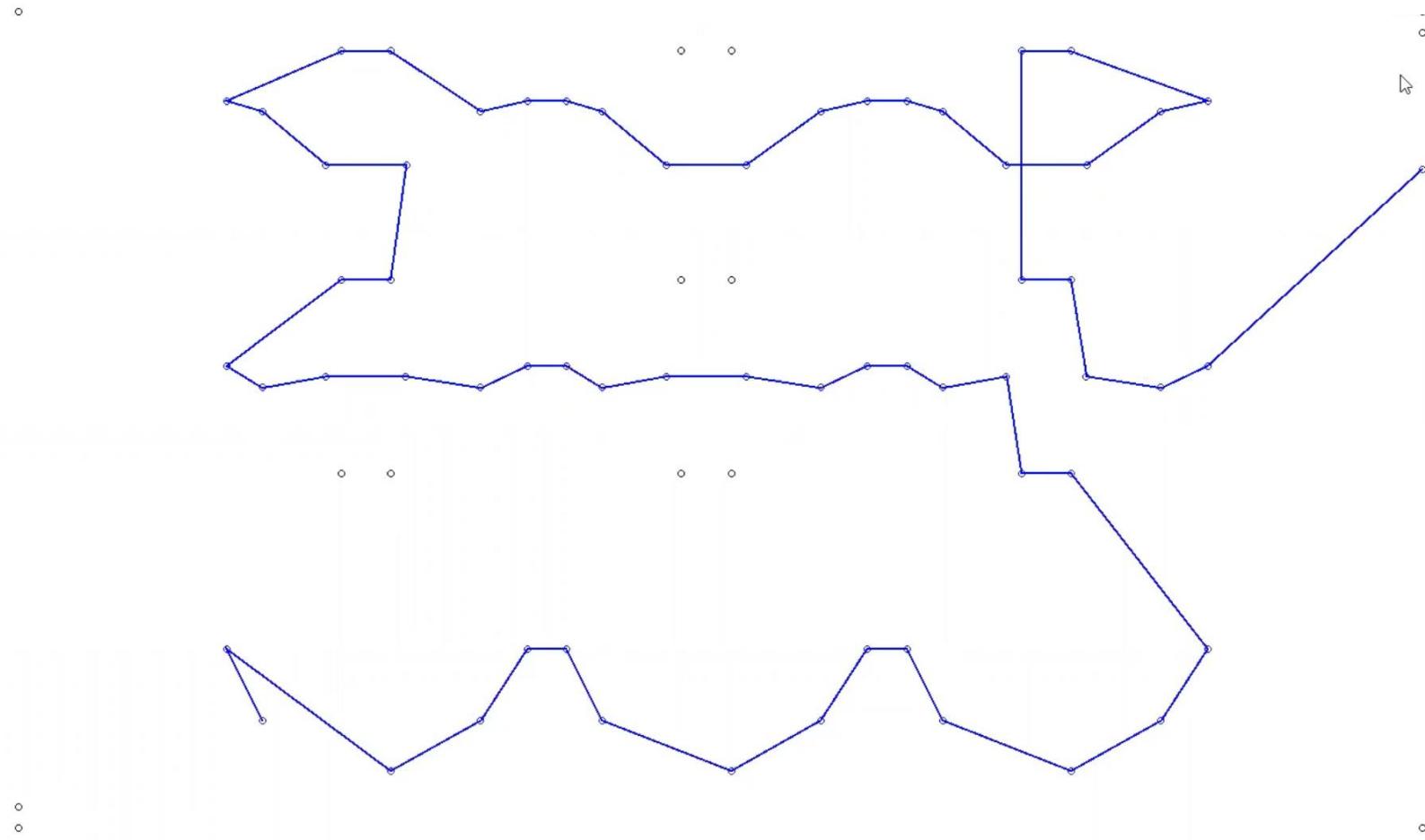
Construction algorithms – Greedy



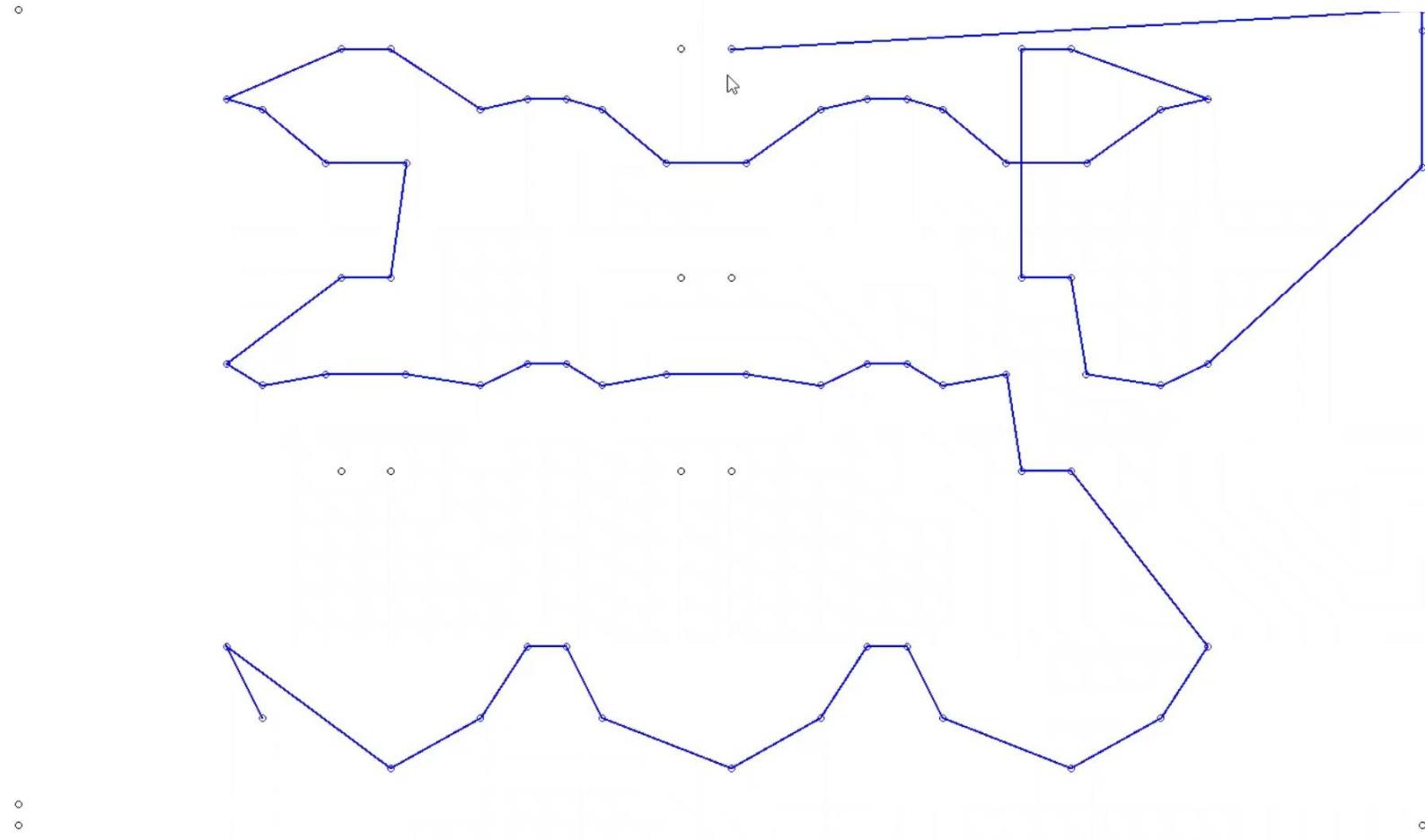
Construction algorithms – Greedy



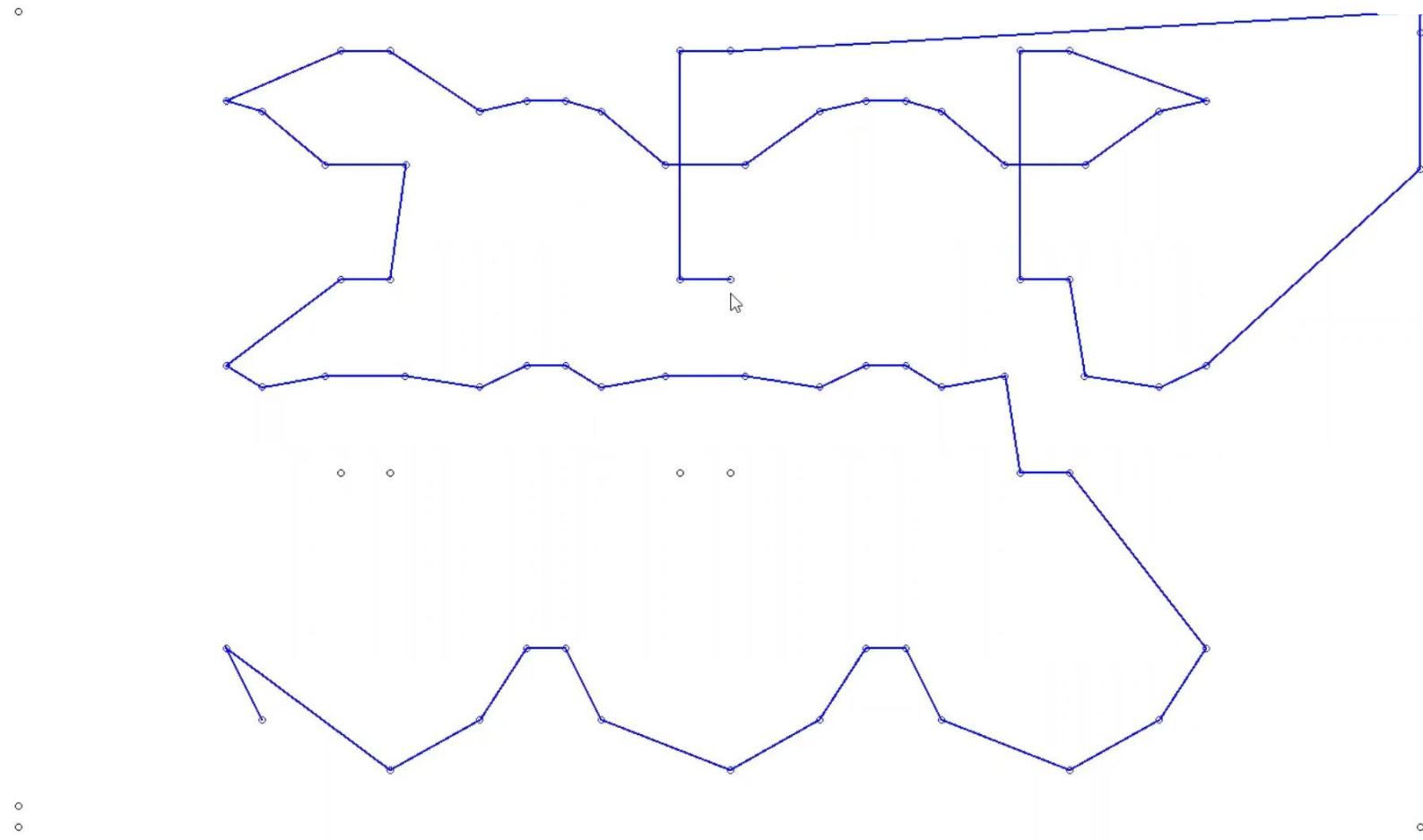
Construction algorithms – Greedy



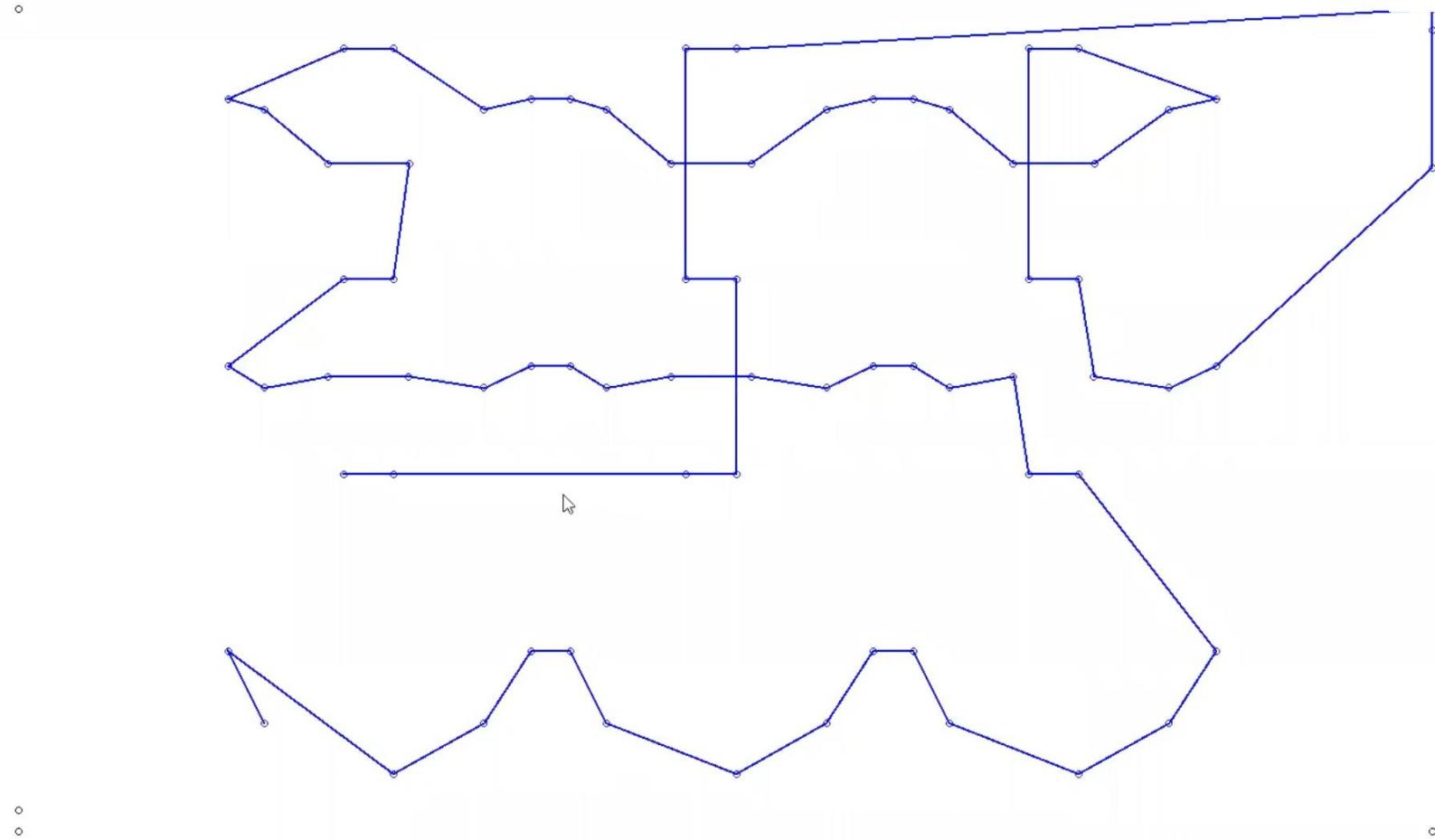
Construction algorithms – Greedy



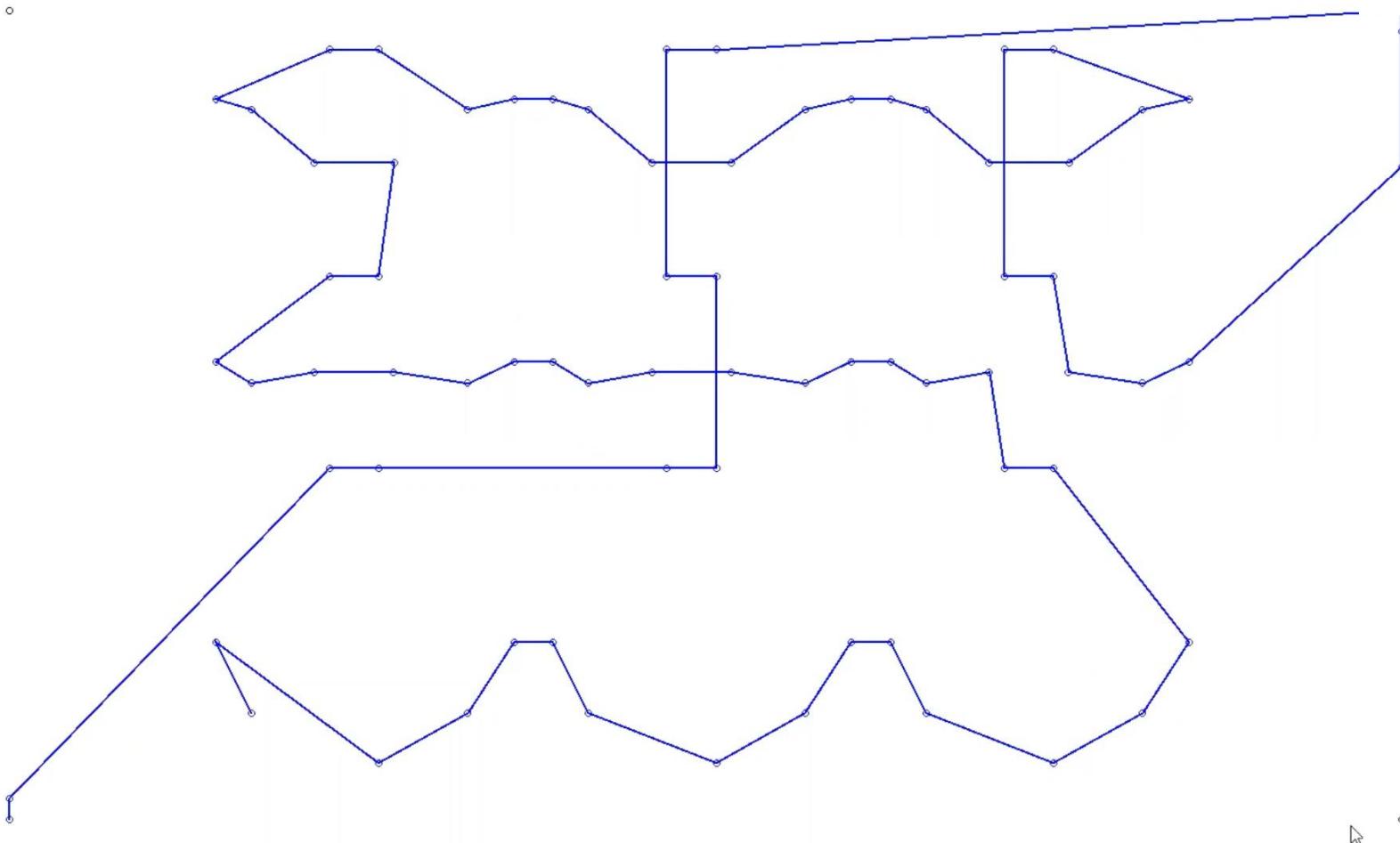
Construction algorithms – Greedy



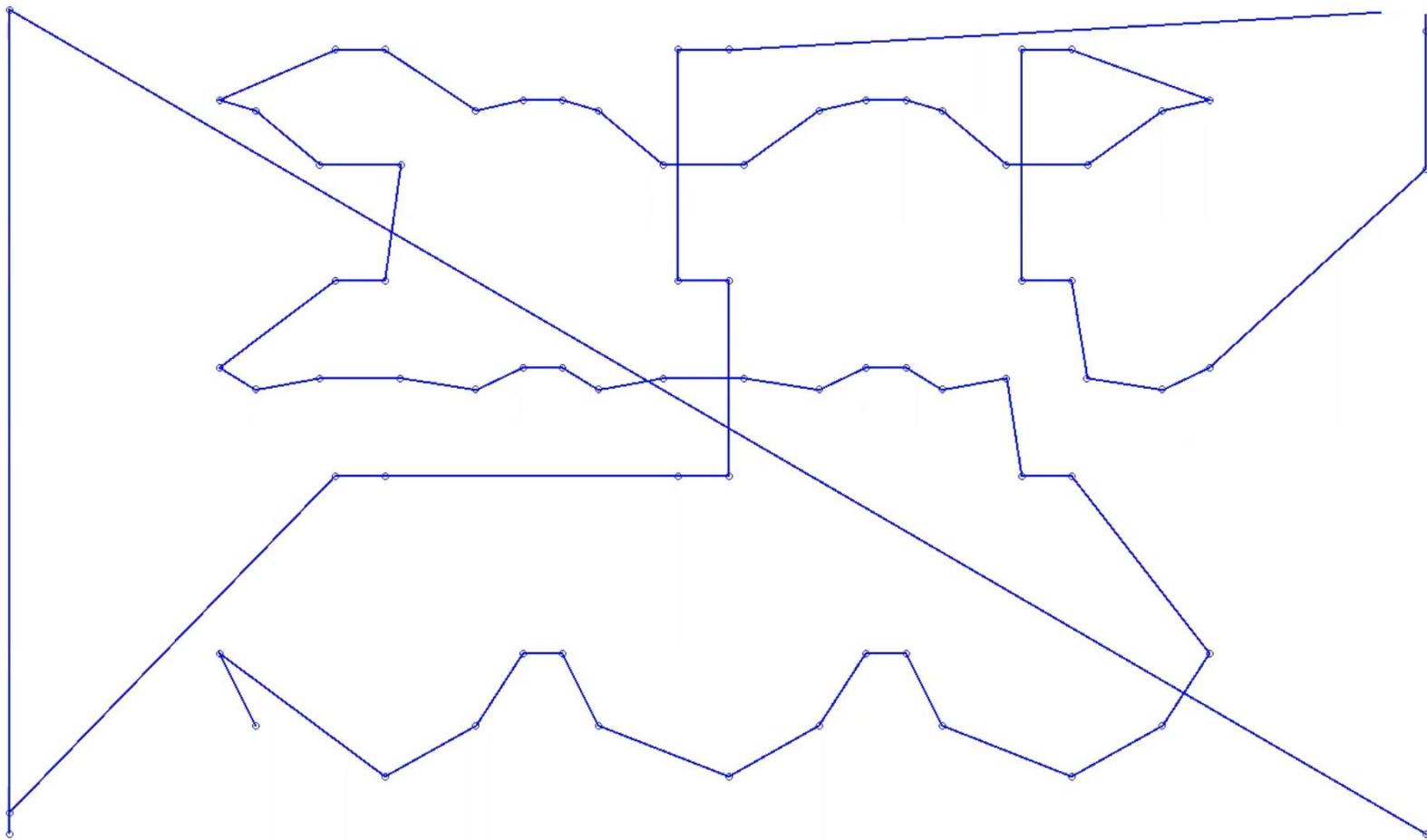
Construction algorithms – Greedy



Construction algorithms – Greedy



Construction algorithms – Greedy



Construction algorithms – Greedy

TSP

- “Nearest Neighbour” is a classic greedy algorithm
- Greedy algorithms perform locally-optimal moves
- However, they often make “structural” errors.
- Greedy algorithms exist for many problems:
 - Knapsack: Insert the item with the best value/weight ratio
 - Scheduling: Schedule the item with the earliest deadline

← REALLY
GOOD

OK

Construction algorithms – Minimum Insert Cost

A better greedy algorithm

- “Minimum Insert” is another greedy algorithm – but performs slightly better
- At each iteration
 - Find the customer that, when inserted, increases the cost by the least.
 - Insert it in the position that increases the cost by the least
 - Allows insert between any pair of customers – not just at the end

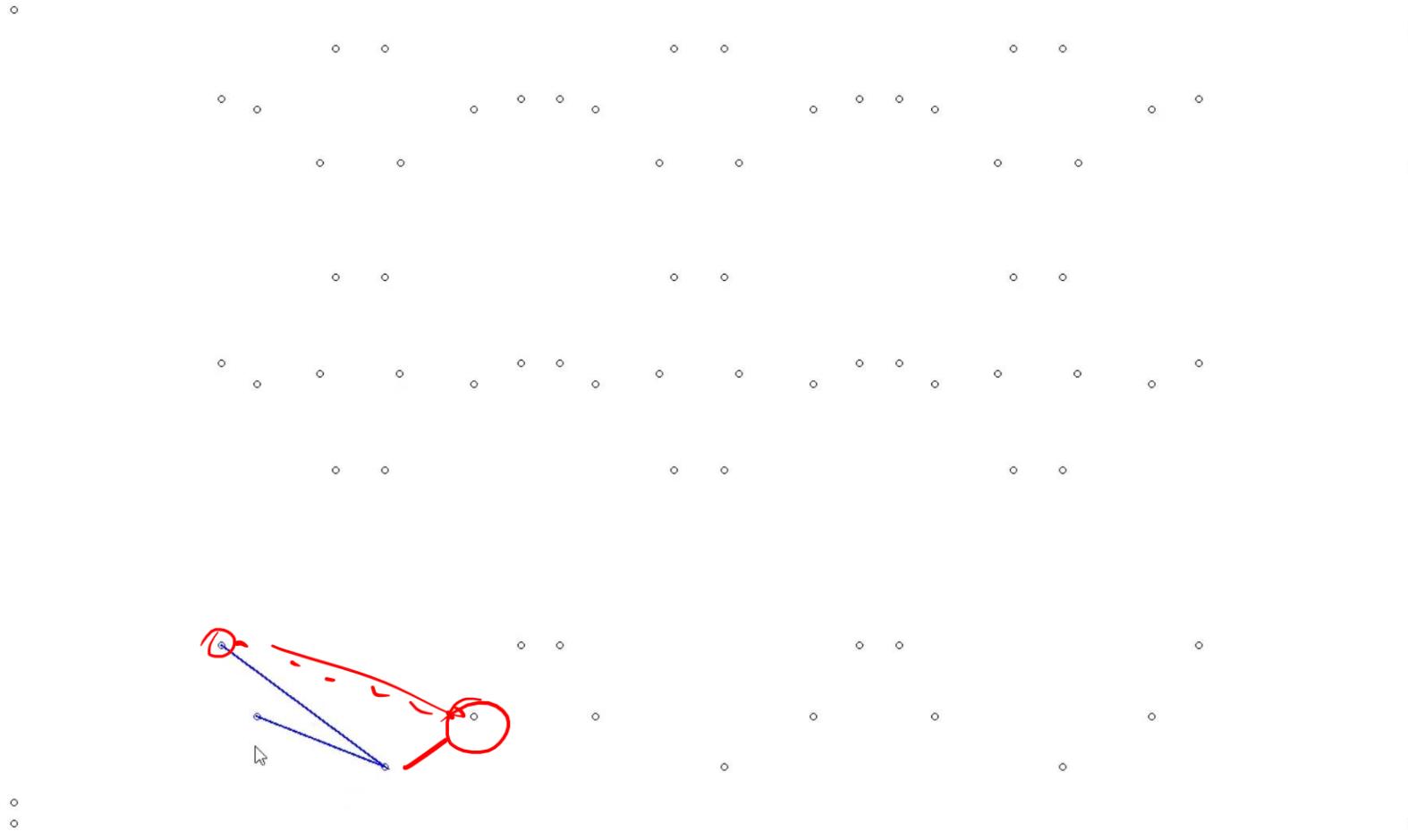
OPTIMAL TSP



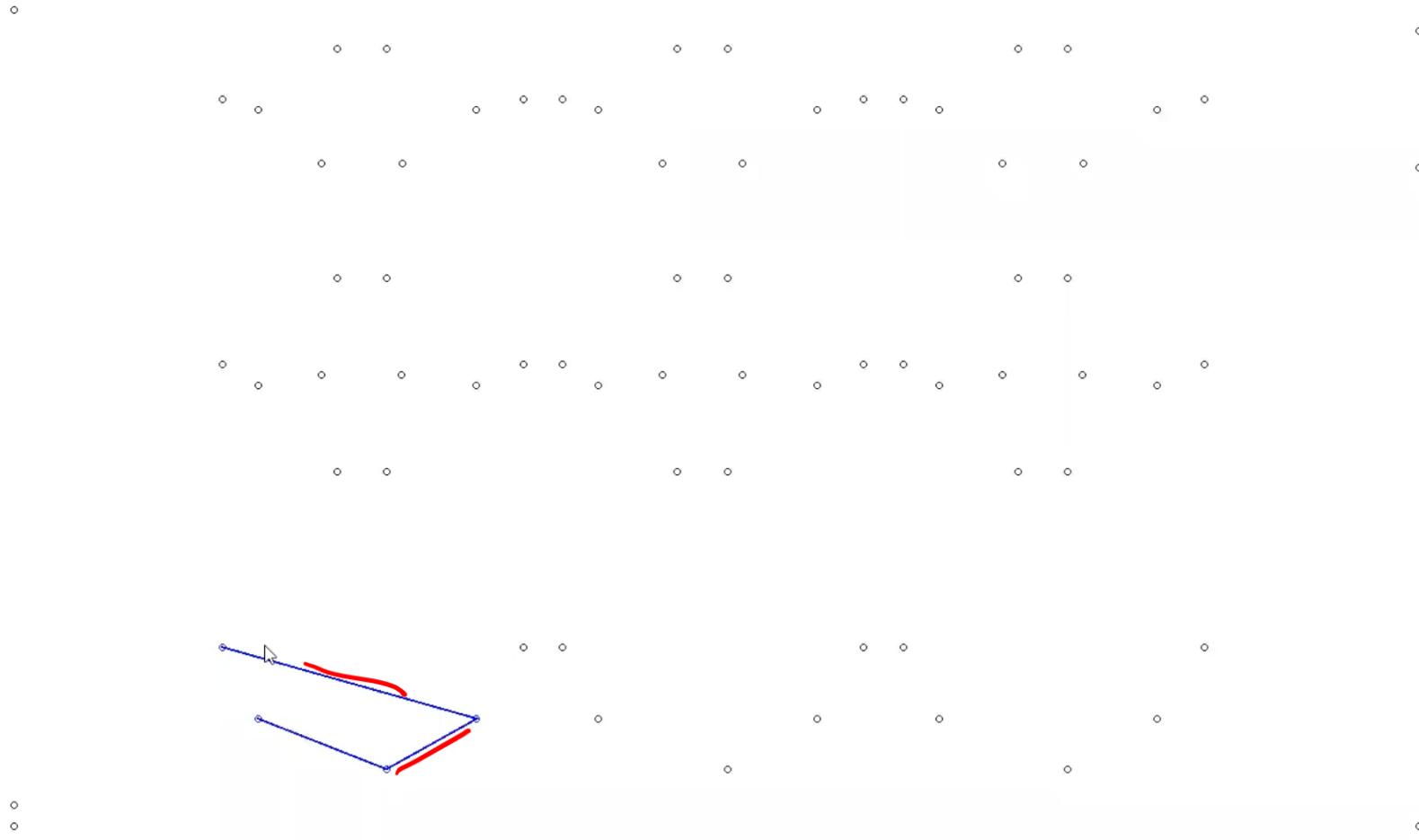
Construction algorithms – Minimum Insert Cost



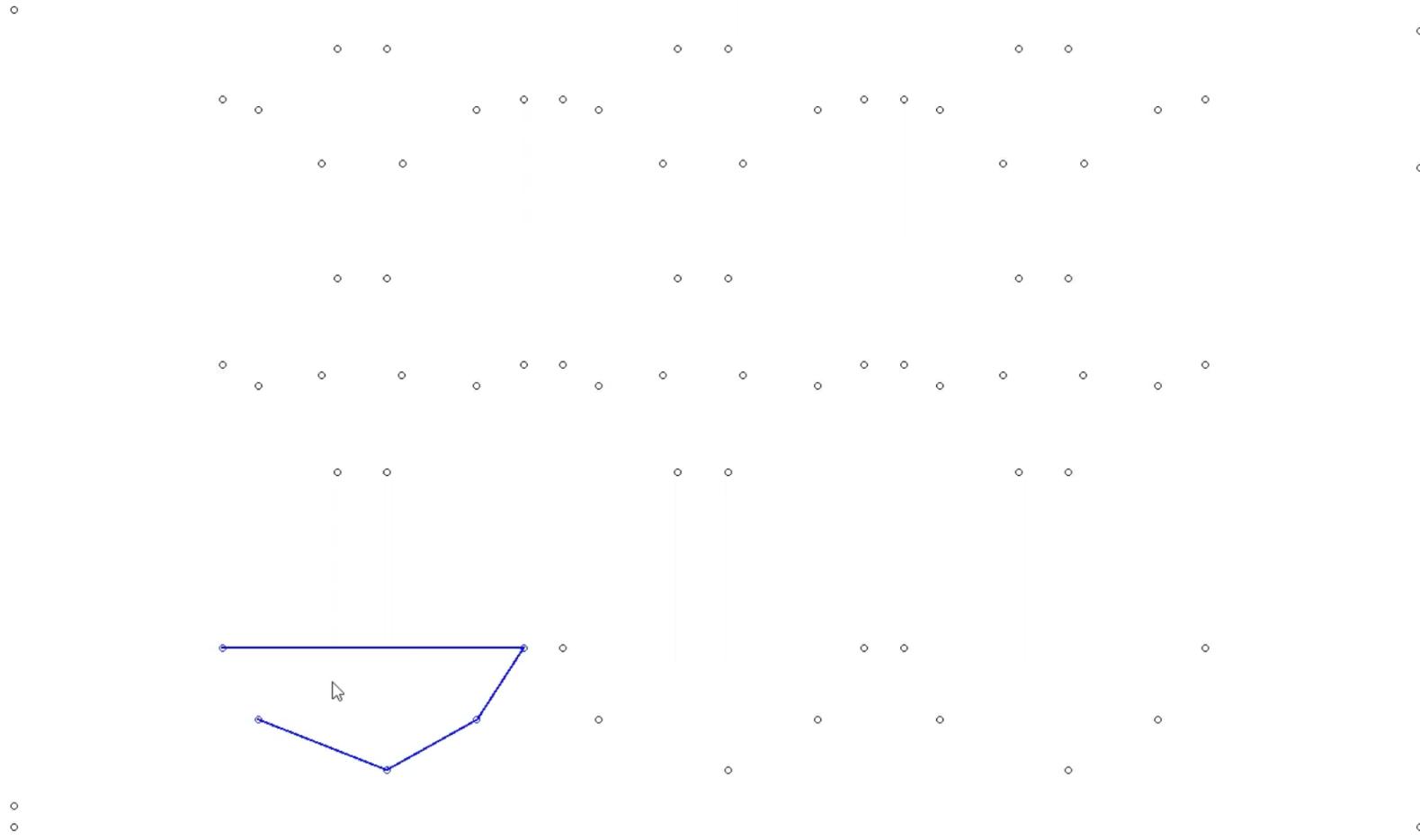
Construction algorithms – Minimum Insert Cost



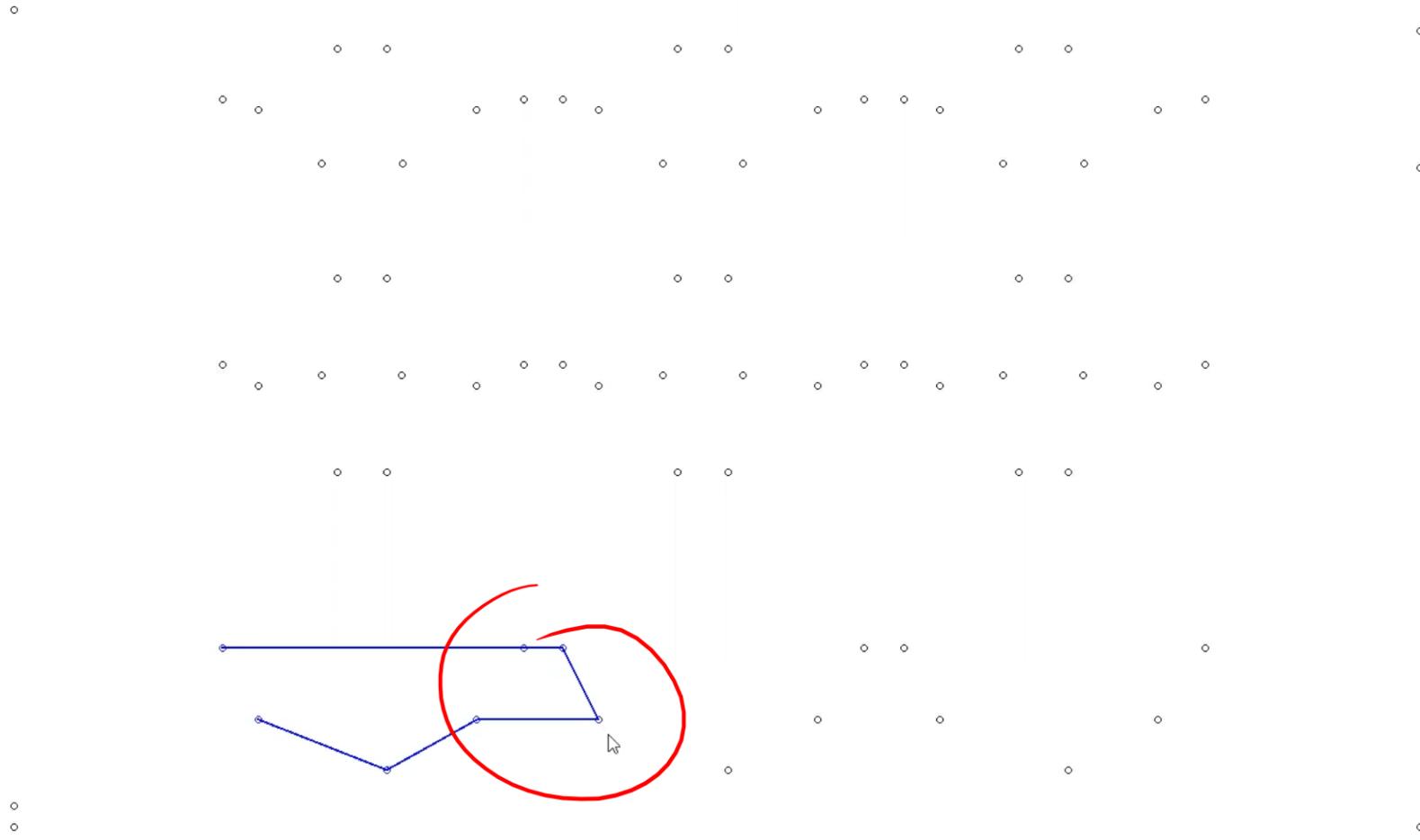
Construction algorithms – Minimum Insert Cost



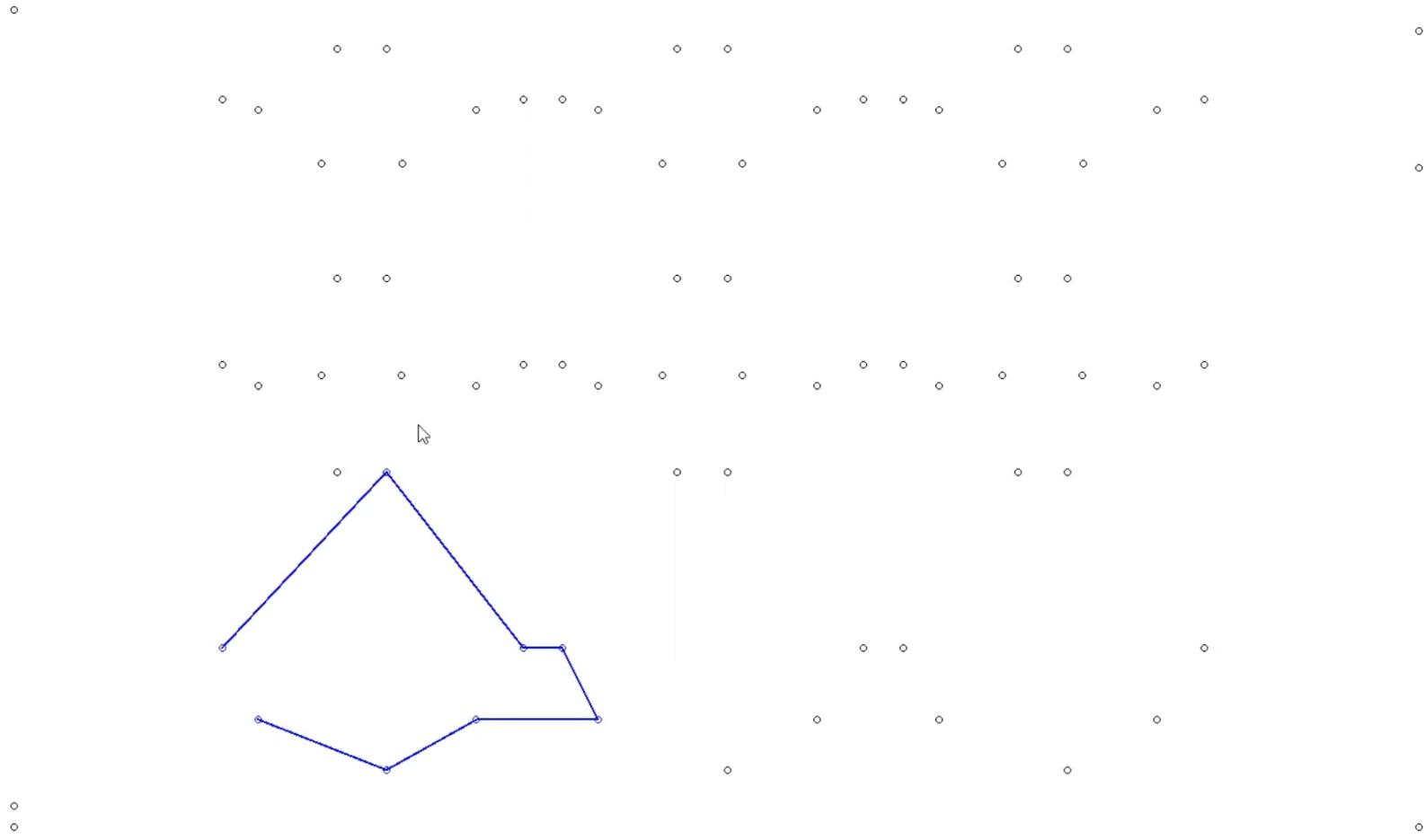
Construction algorithms – Minimum Insert Cost



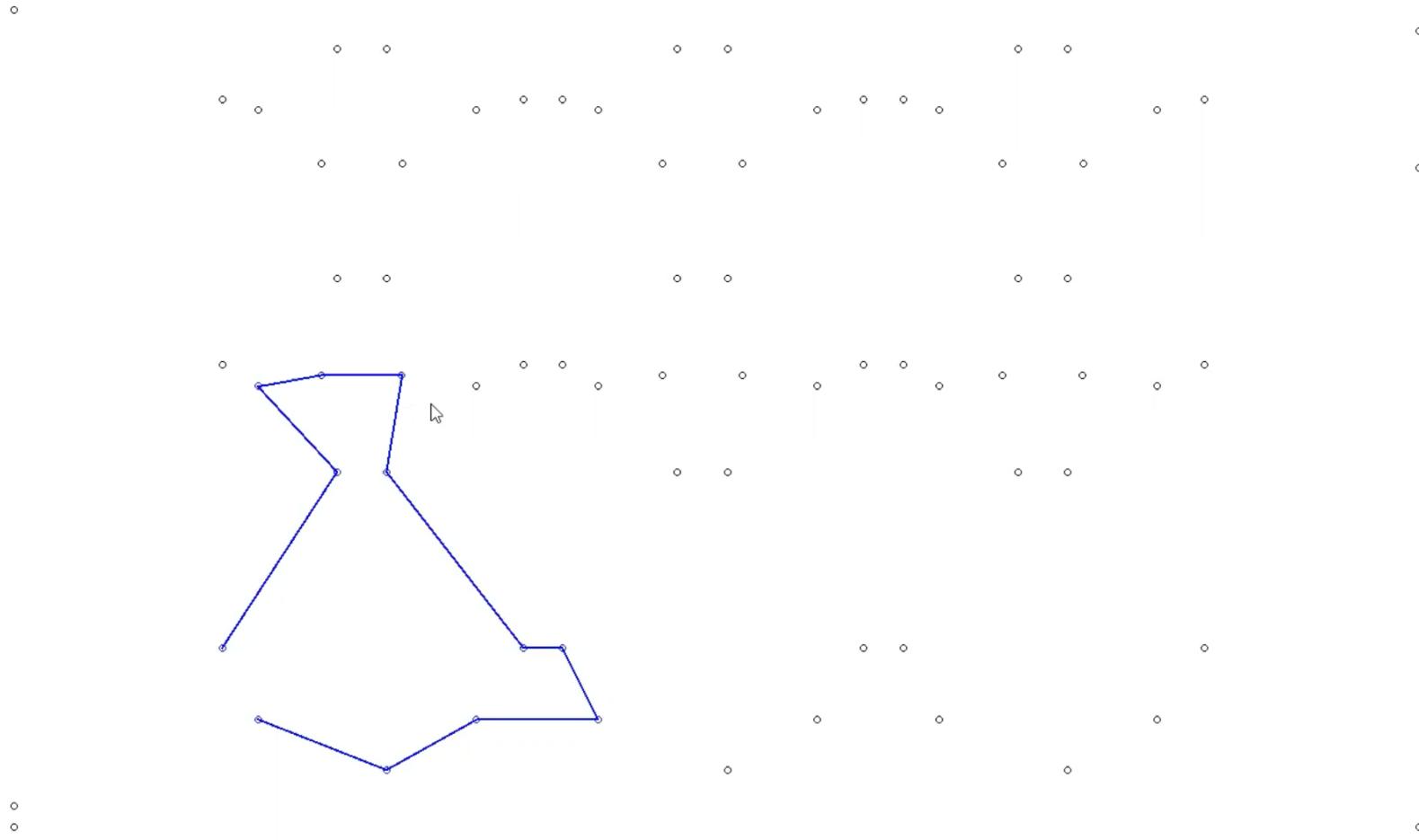
Construction algorithms – Minimum Insert Cost



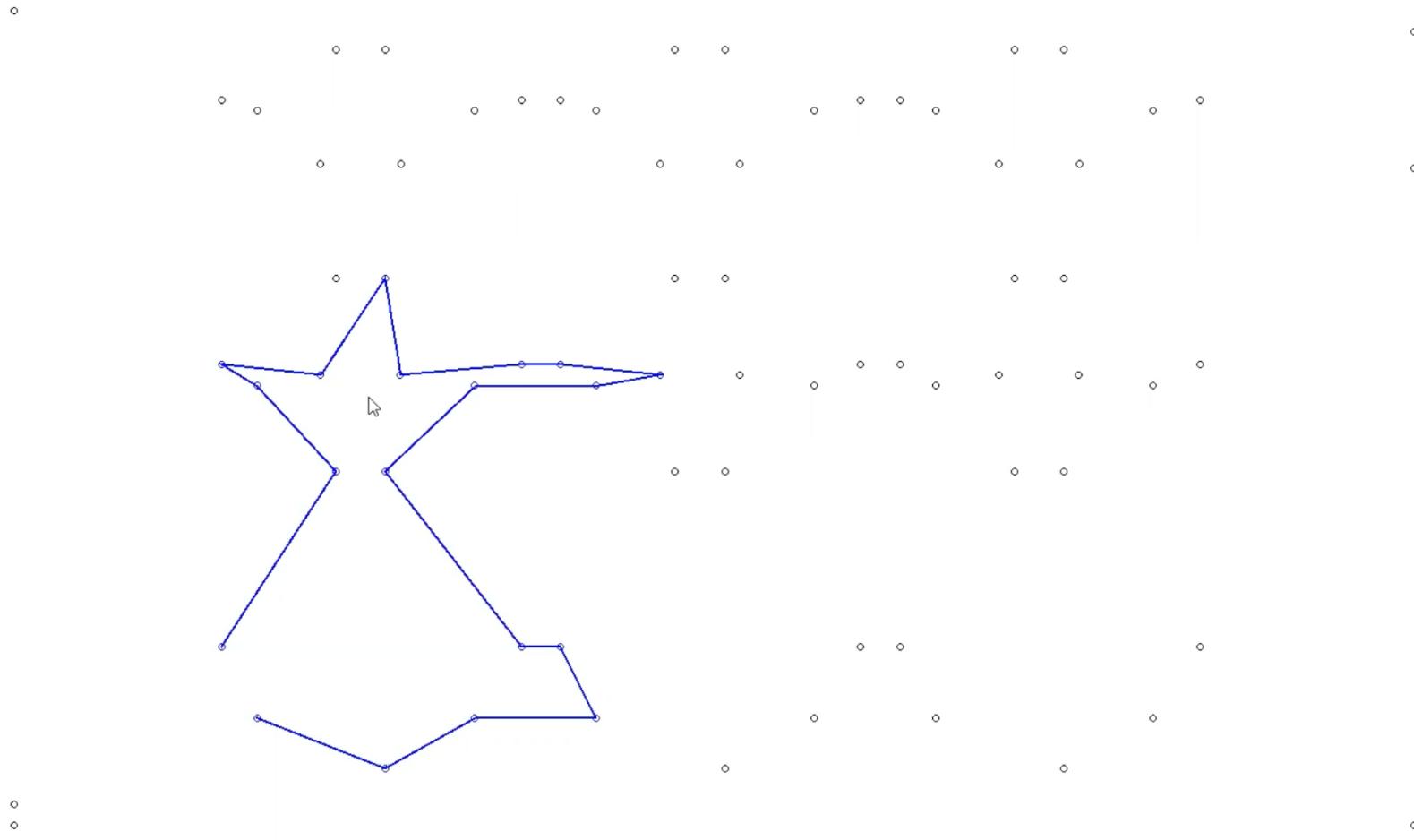
Construction algorithms – Minimum Insert Cost



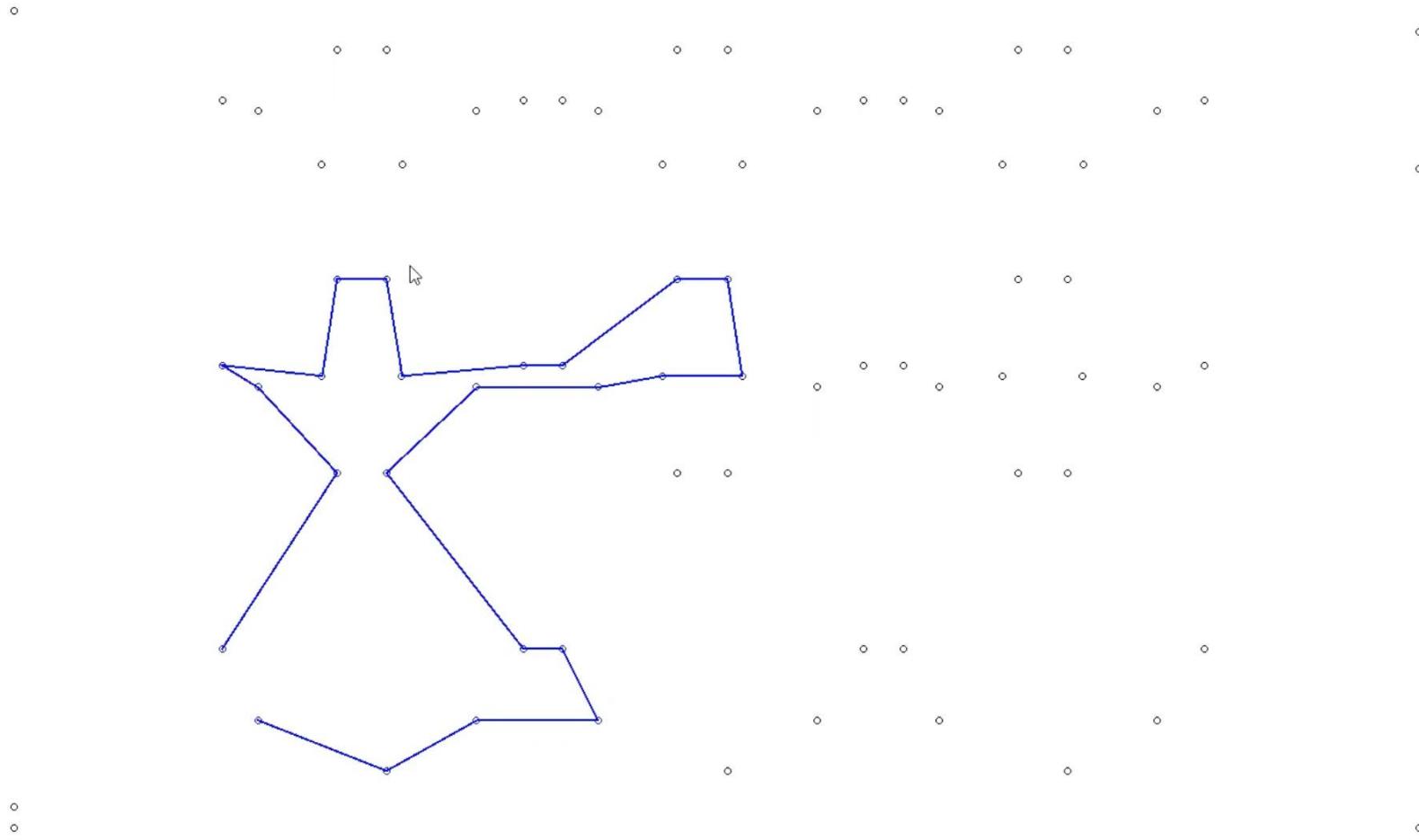
Construction algorithms – Minimum Insert Cost



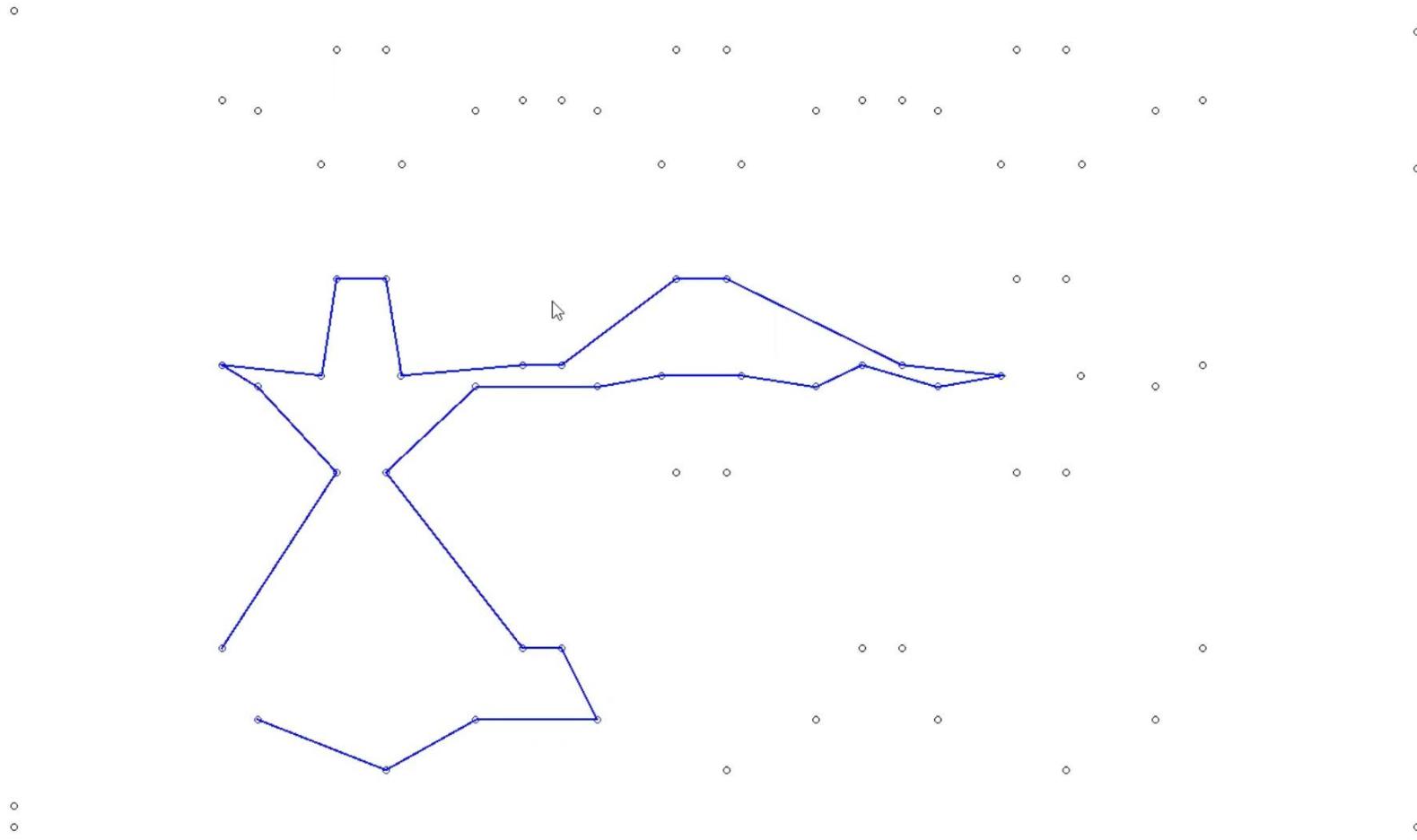
Construction algorithms – Minimum Insert Cost



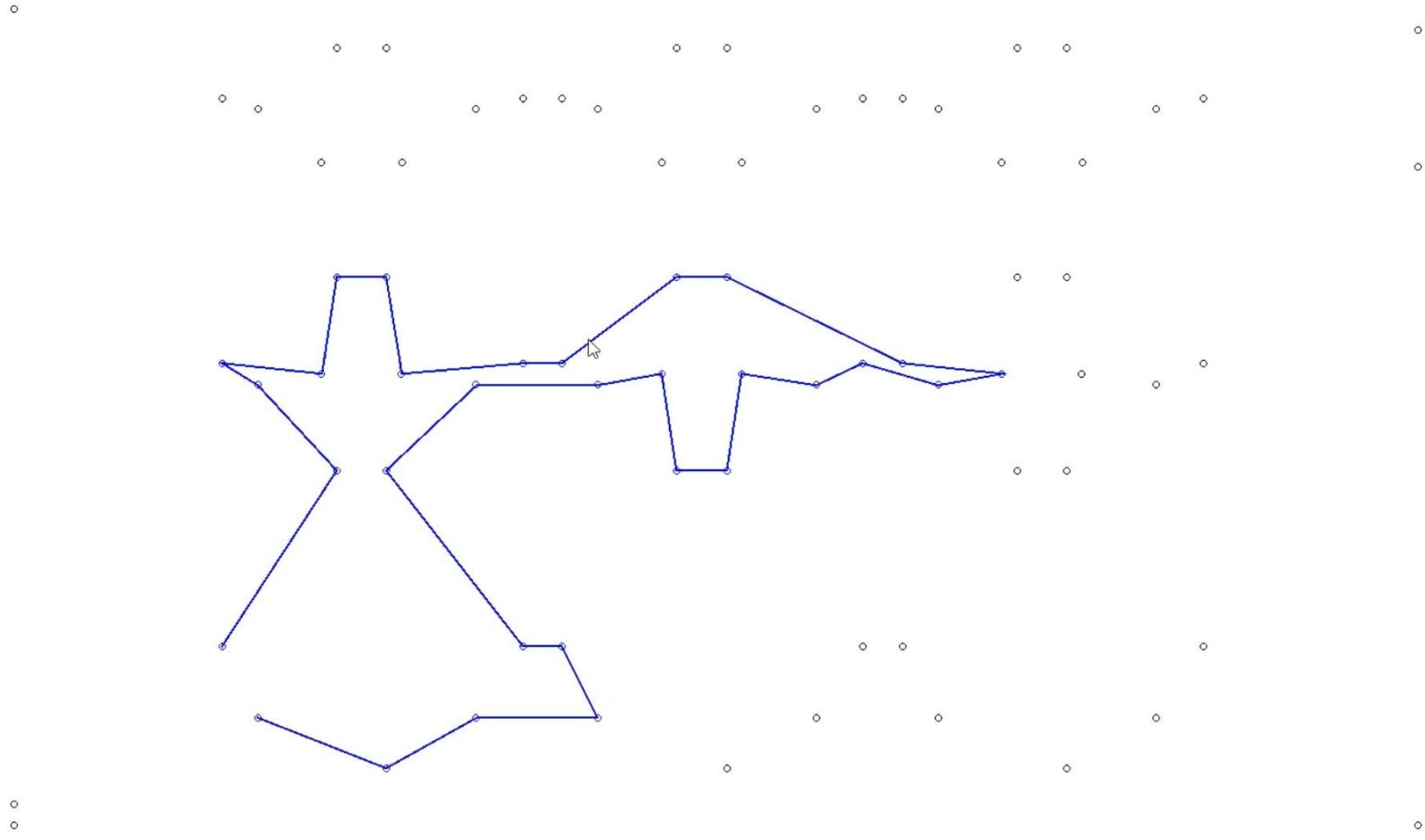
Construction algorithms – Minimum Insert Cost



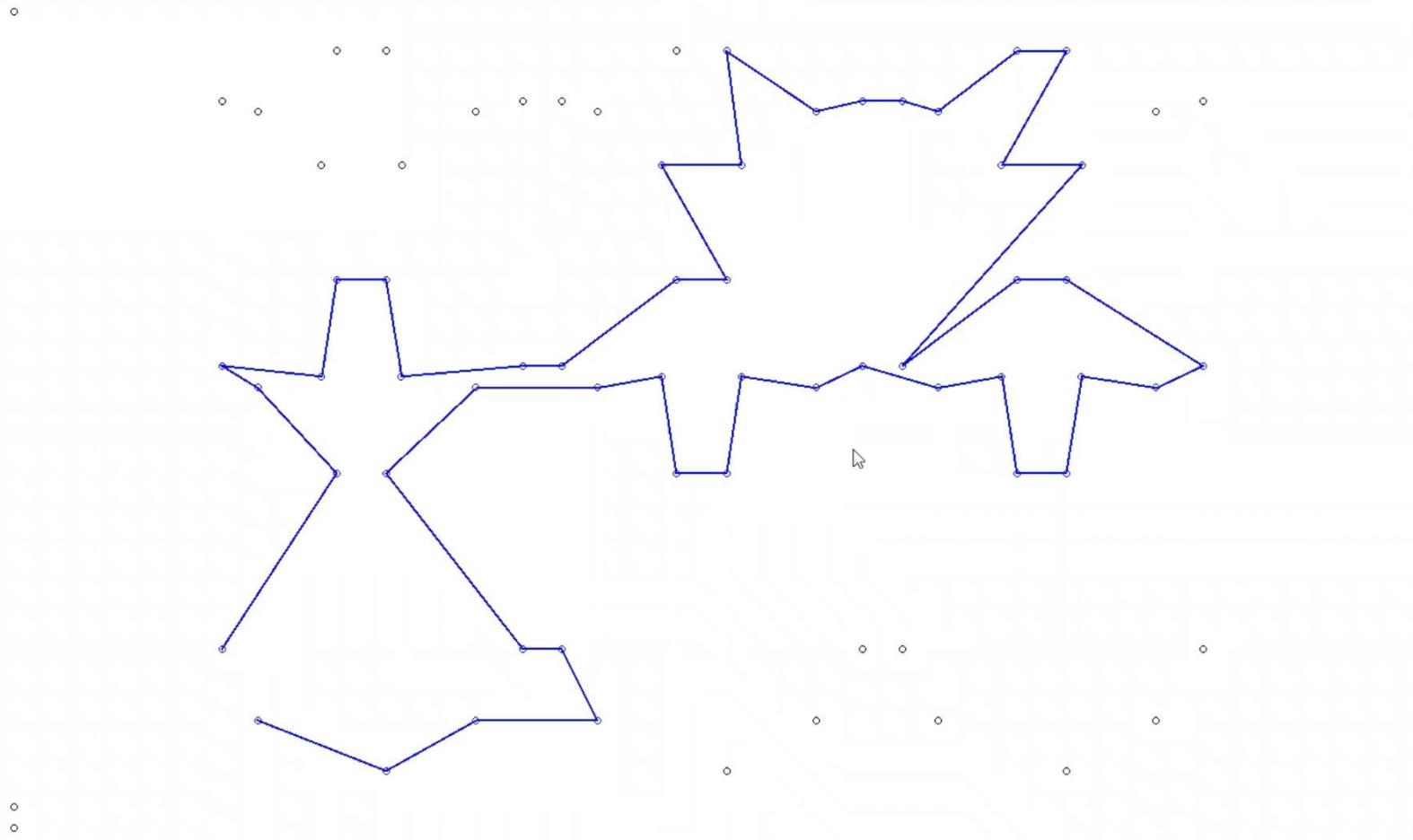
Construction algorithms – Minimum Insert Cost



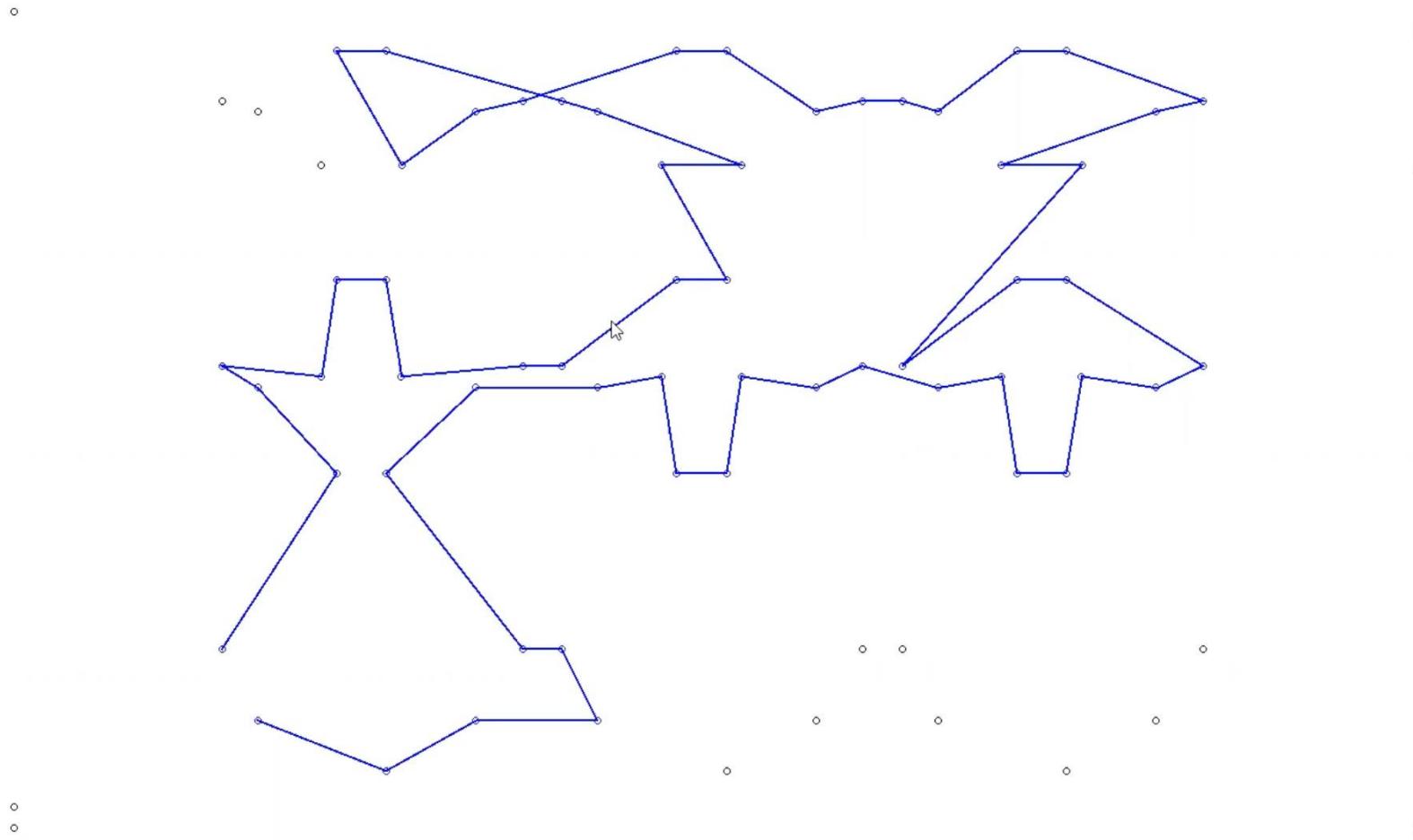
Construction algorithms – Minimum Insert Cost



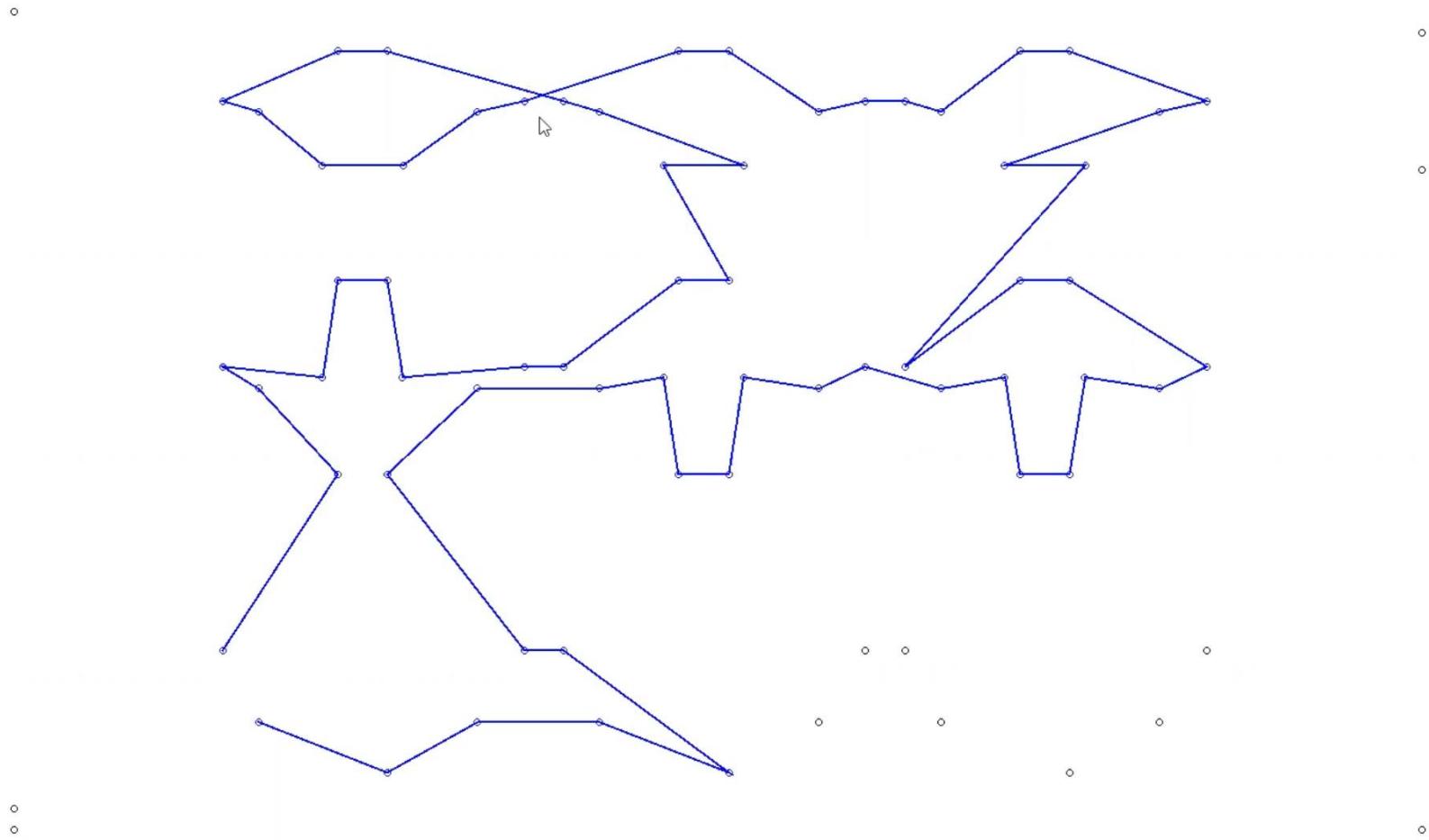
Construction algorithms – Minimum Insert Cost



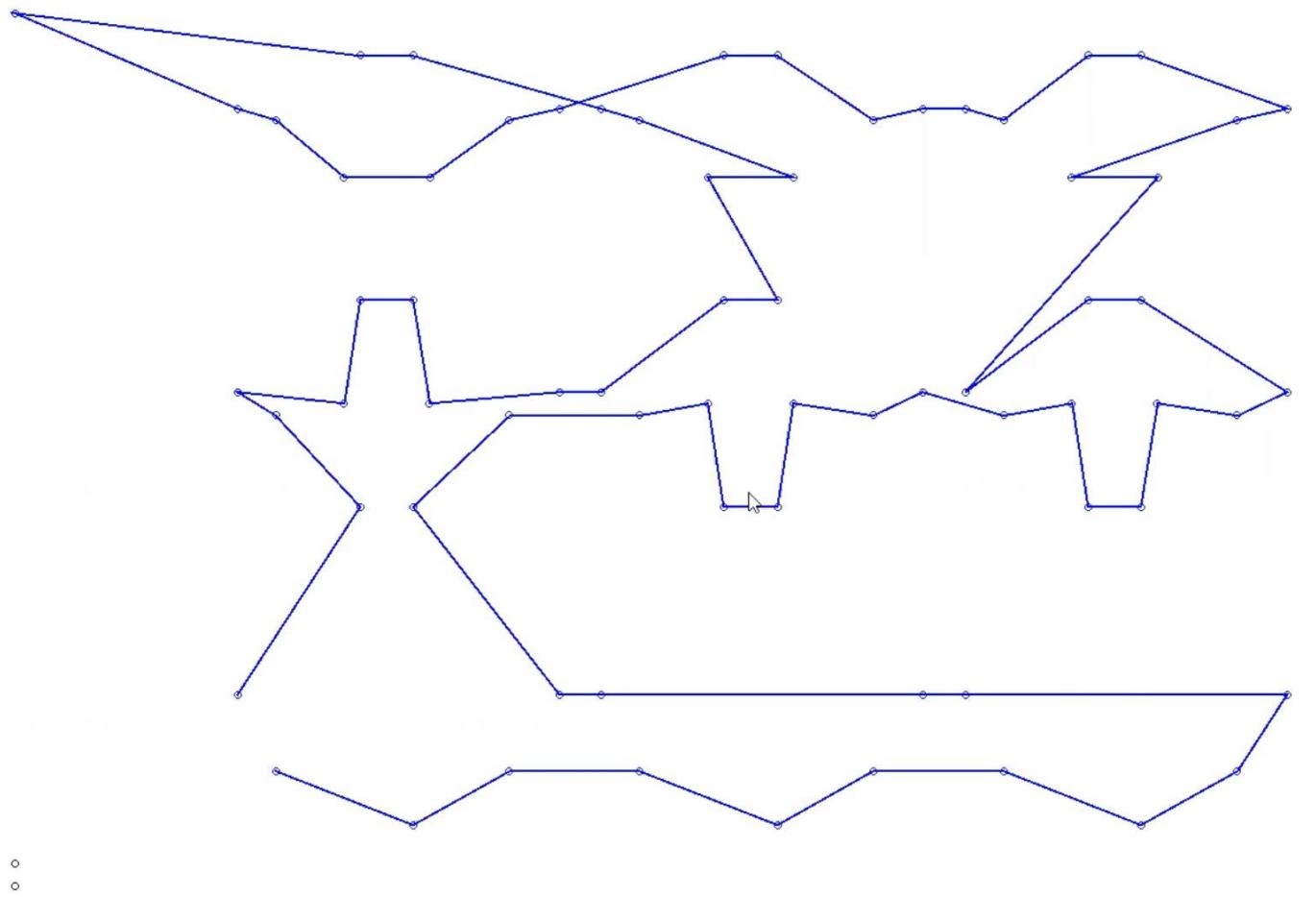
Construction algorithms – Minimum Insert Cost



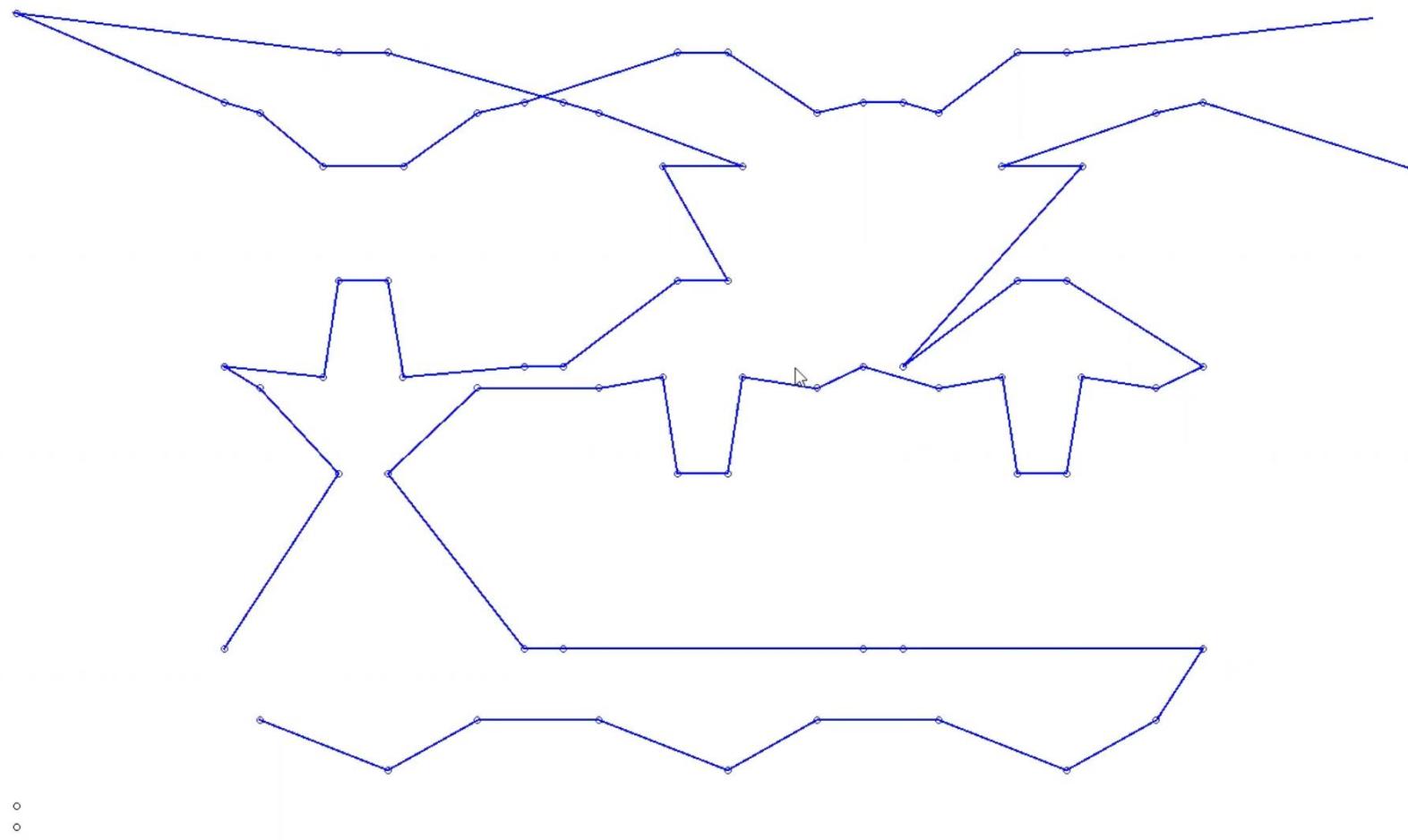
Construction algorithms – Minimum Insert Cost



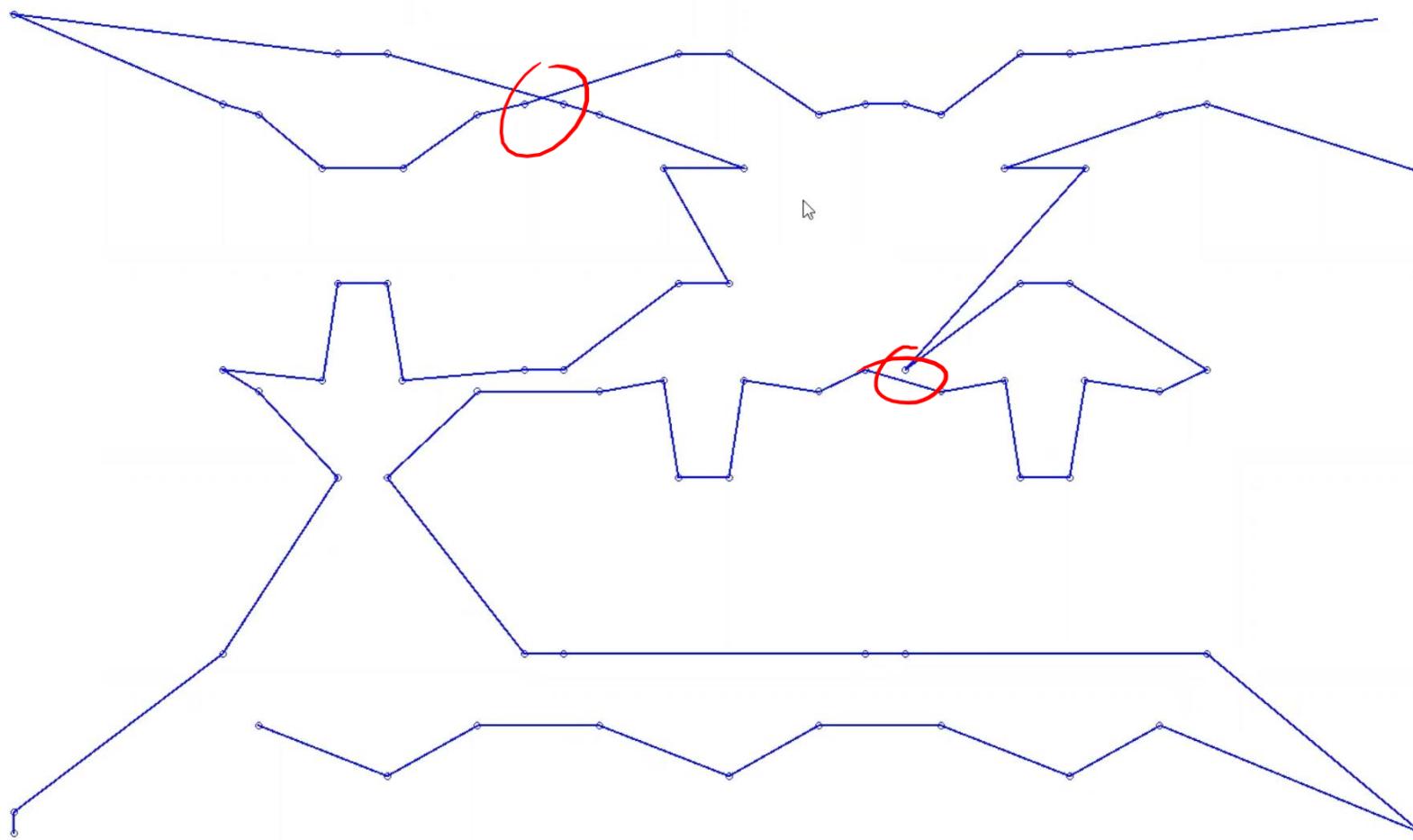
Construction algorithms – Minimum Insert Cost



Construction algorithms – Minimum Insert Cost



Construction algorithms – Minimum Insert Cost

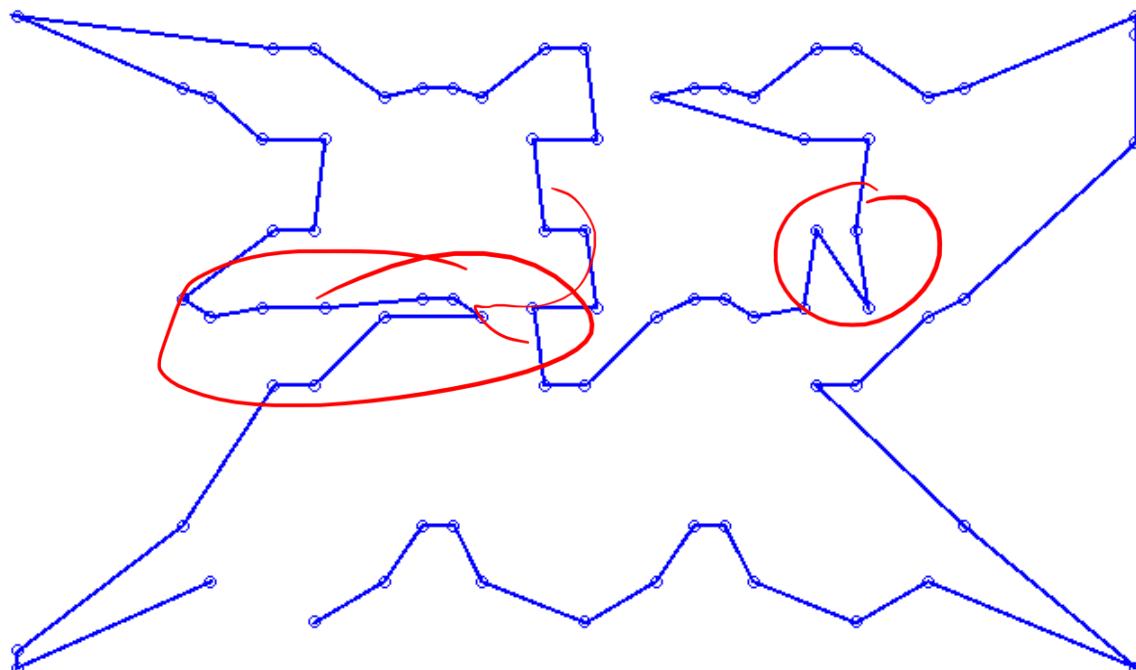


Construction algorithms – Randomisation

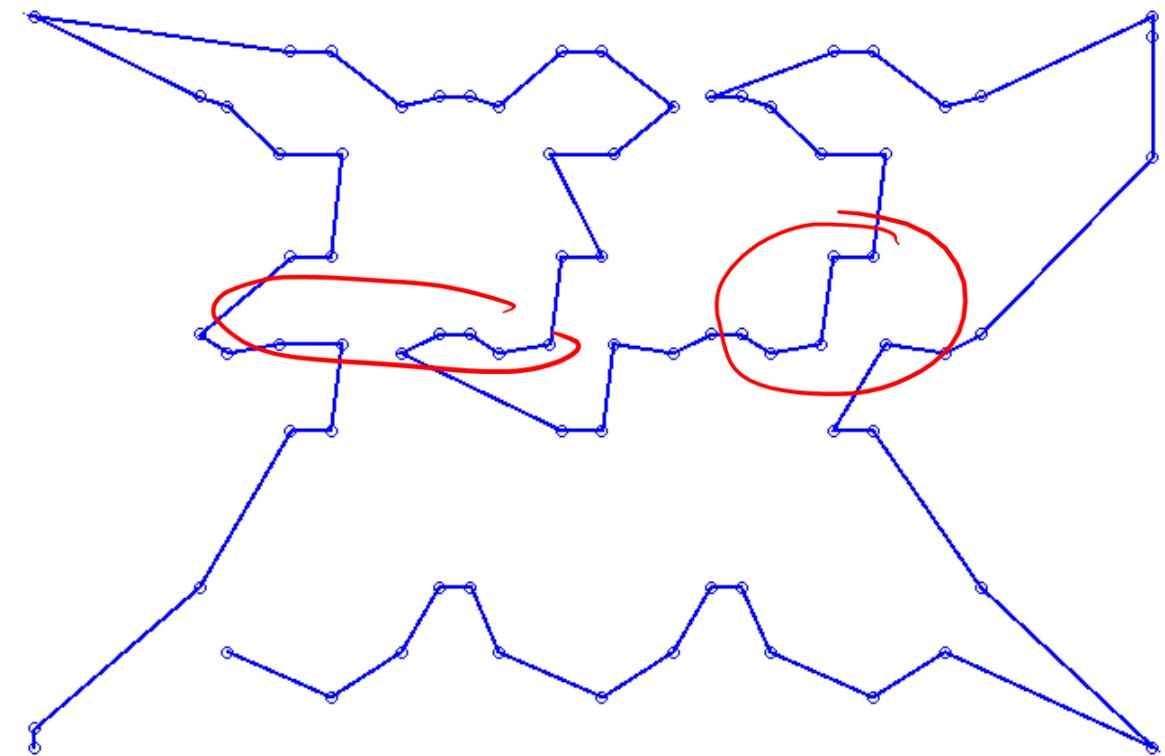
- At each iteration
 - Choose a **random customer**
 - Insert it in the position that increases the cost by the least
 - Allows insert between *any* pair of customers – not just at the end

Construction algorithms–Randomisation

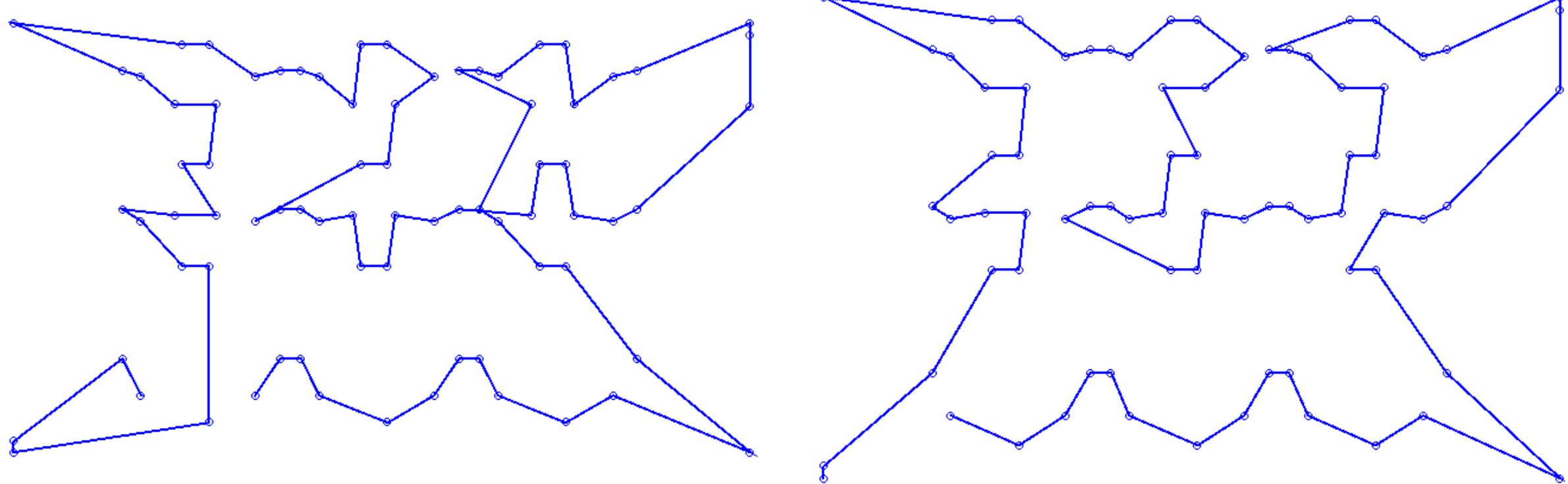
ONE RANDOM TOUR



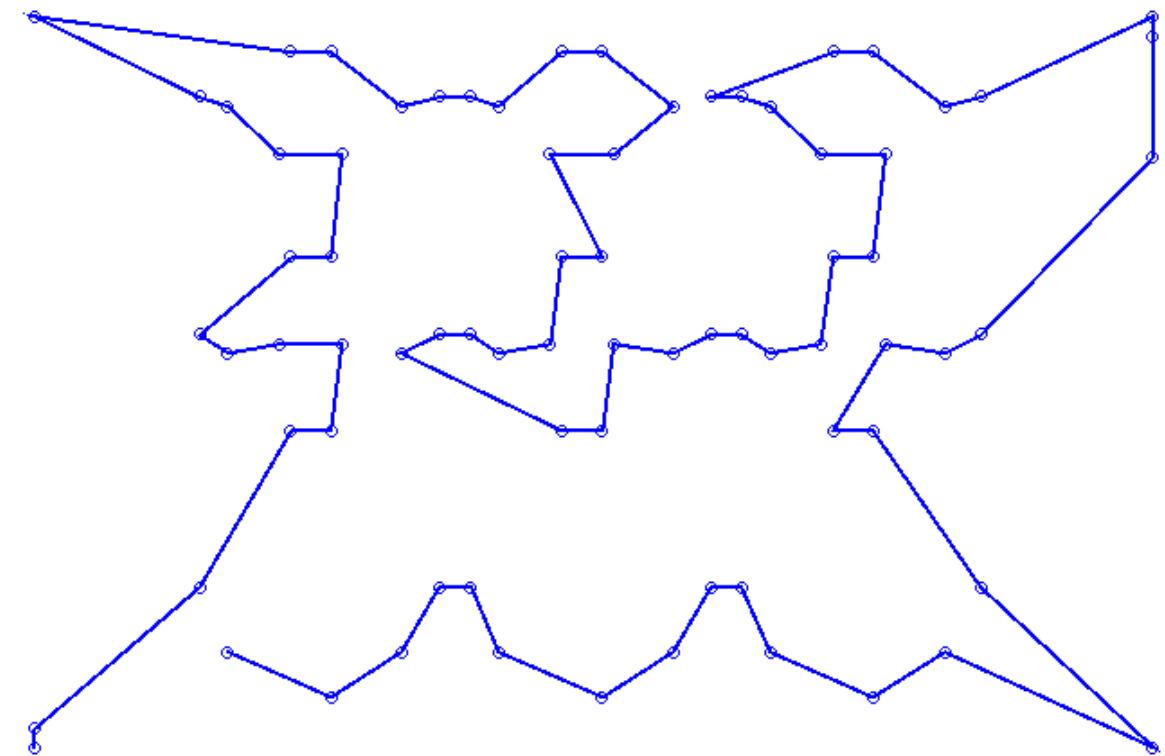
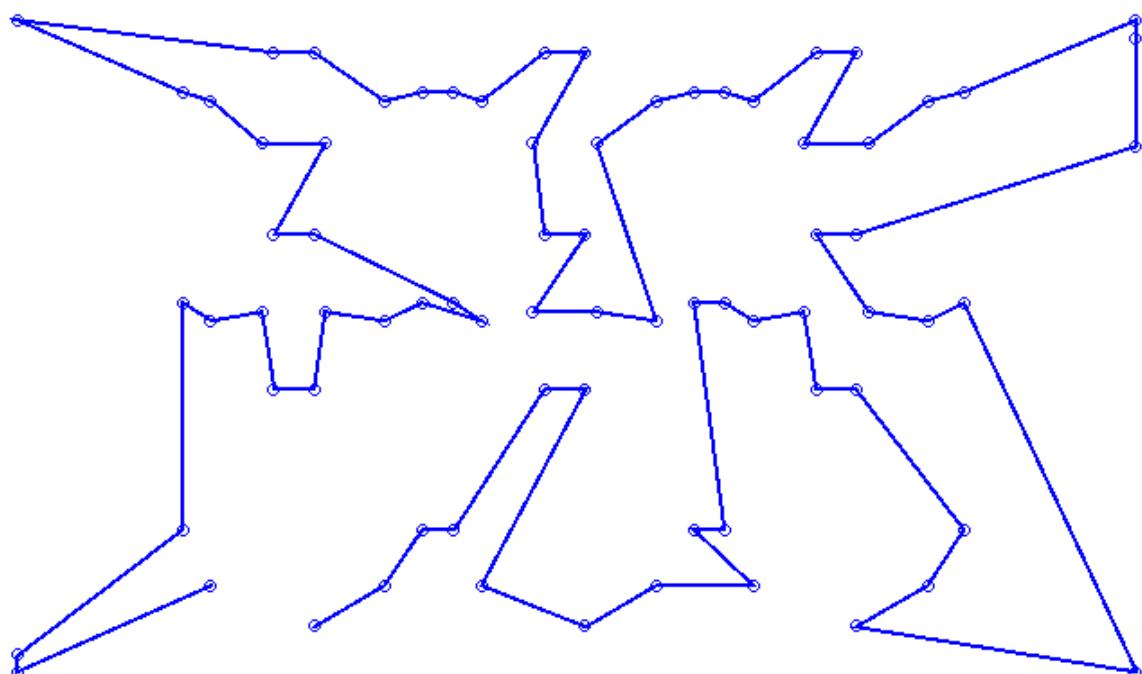
OPT. TOUR



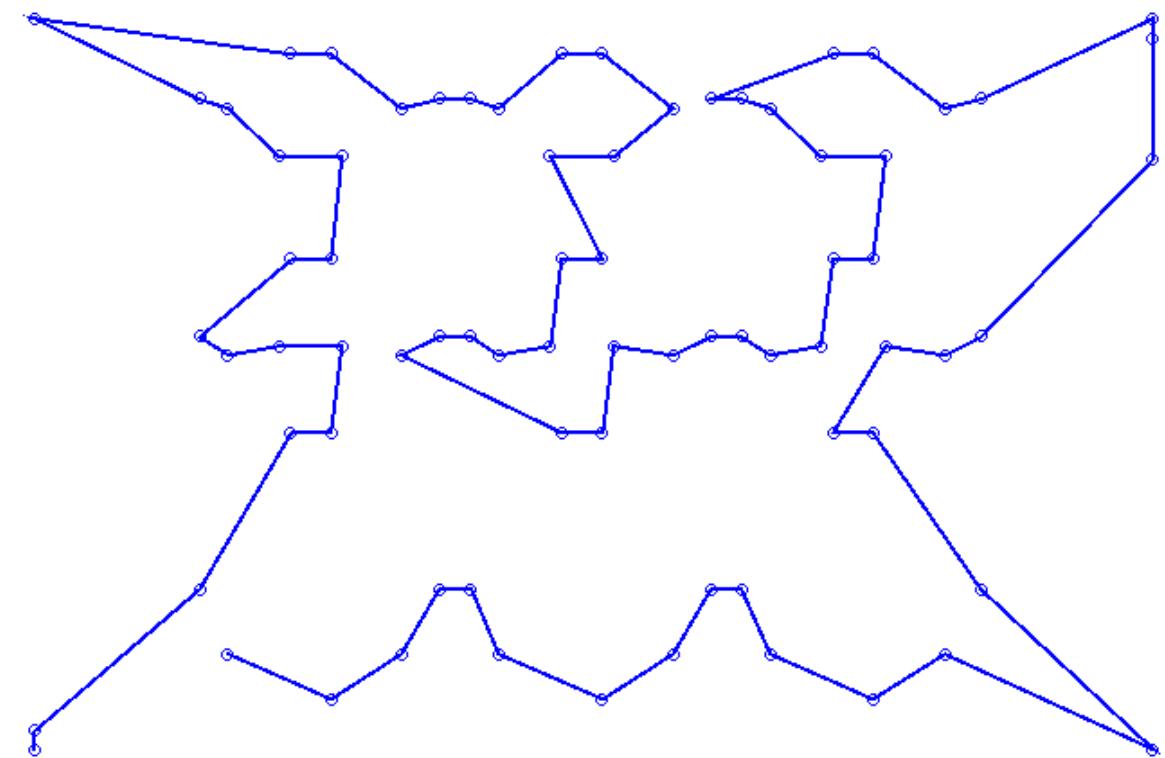
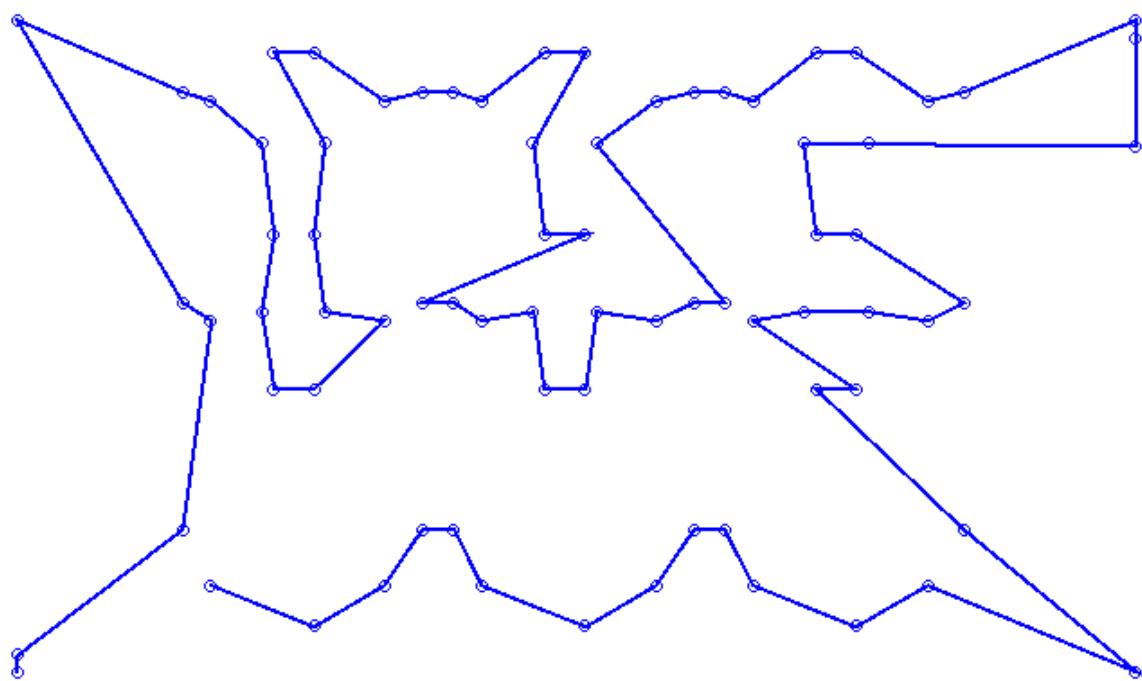
Construction algorithms–Randomisation



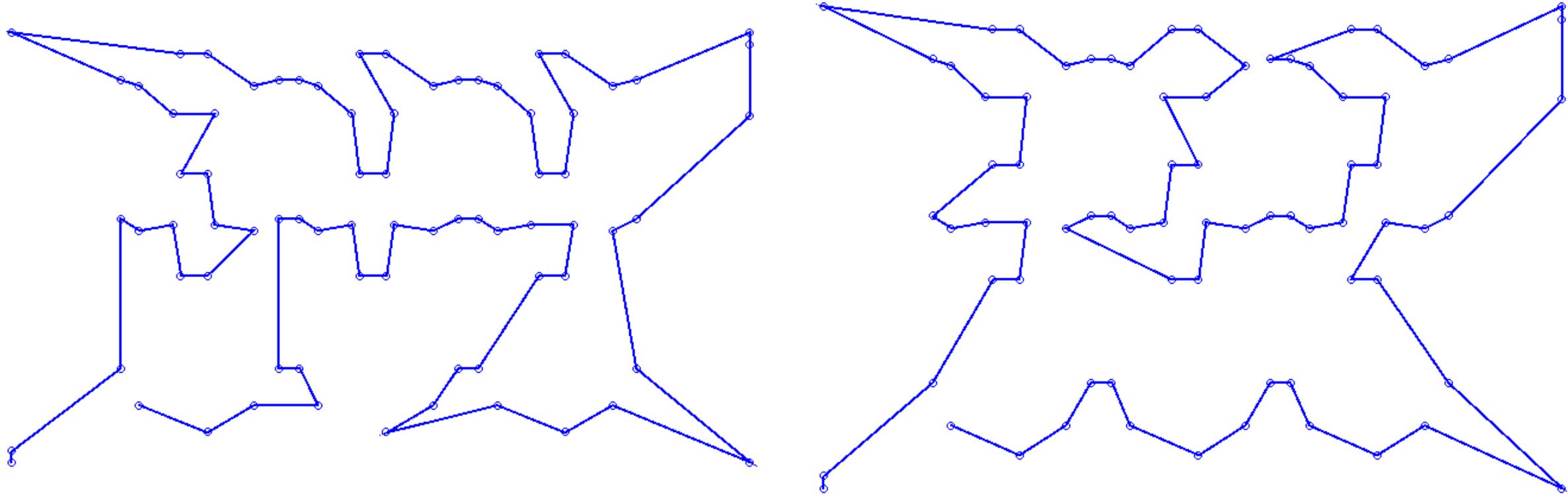
Construction algorithms–Randomisation



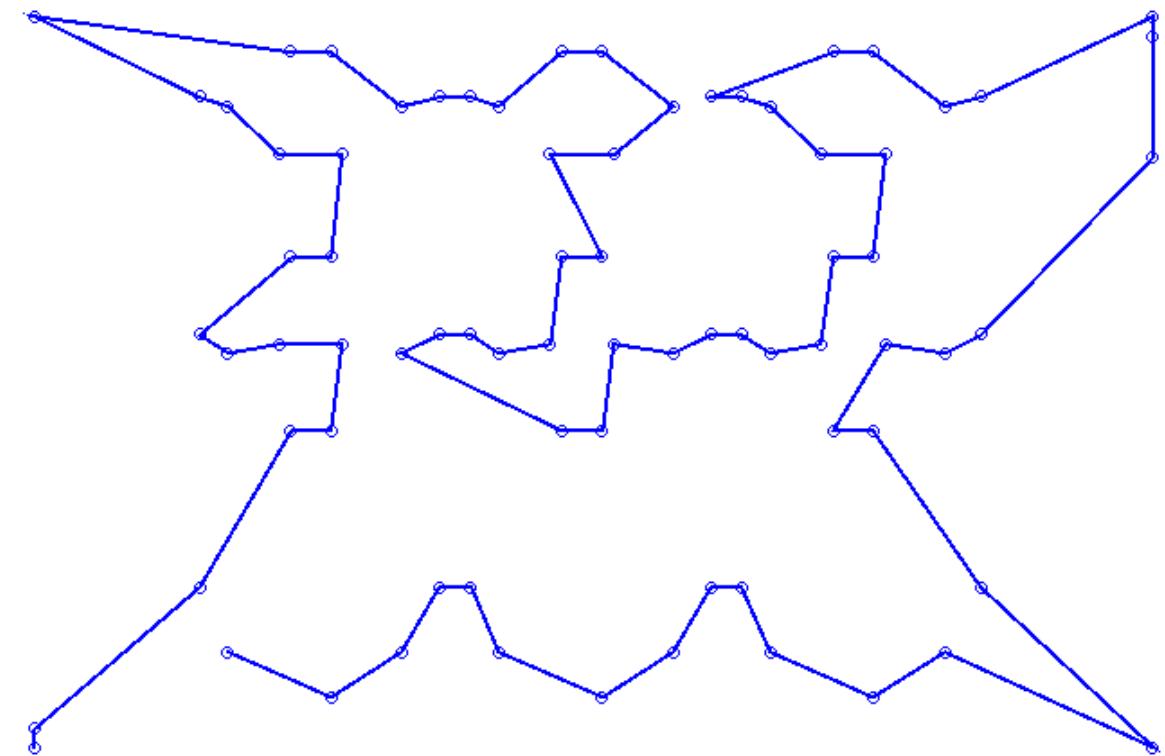
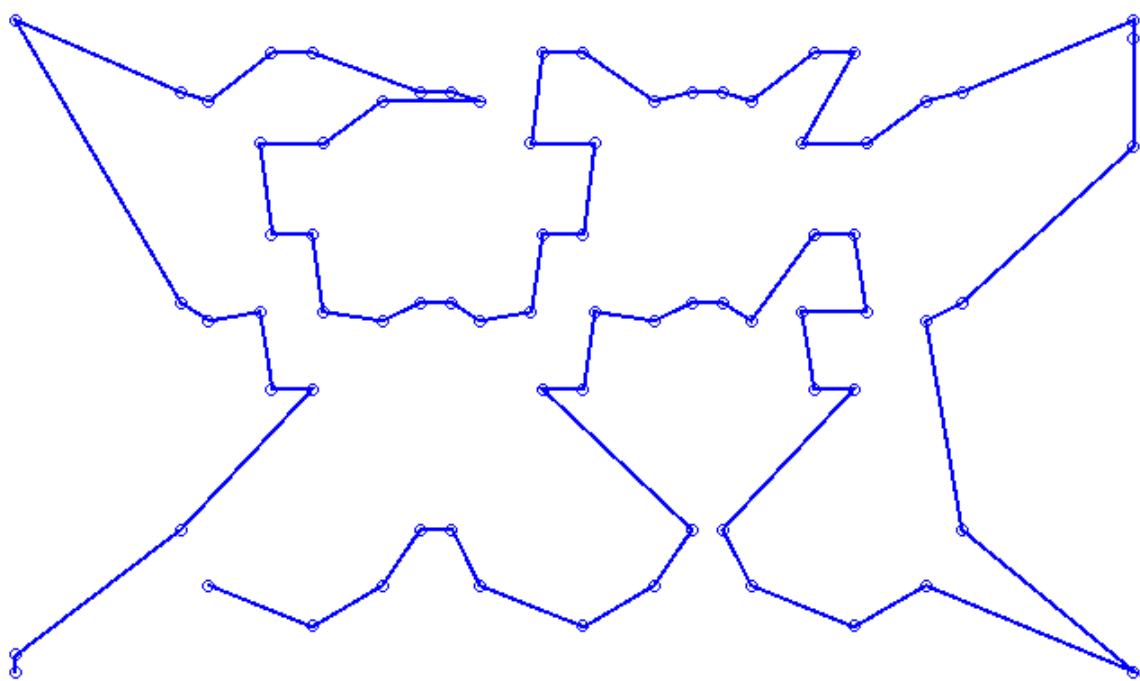
Construction algorithms–Randomisation



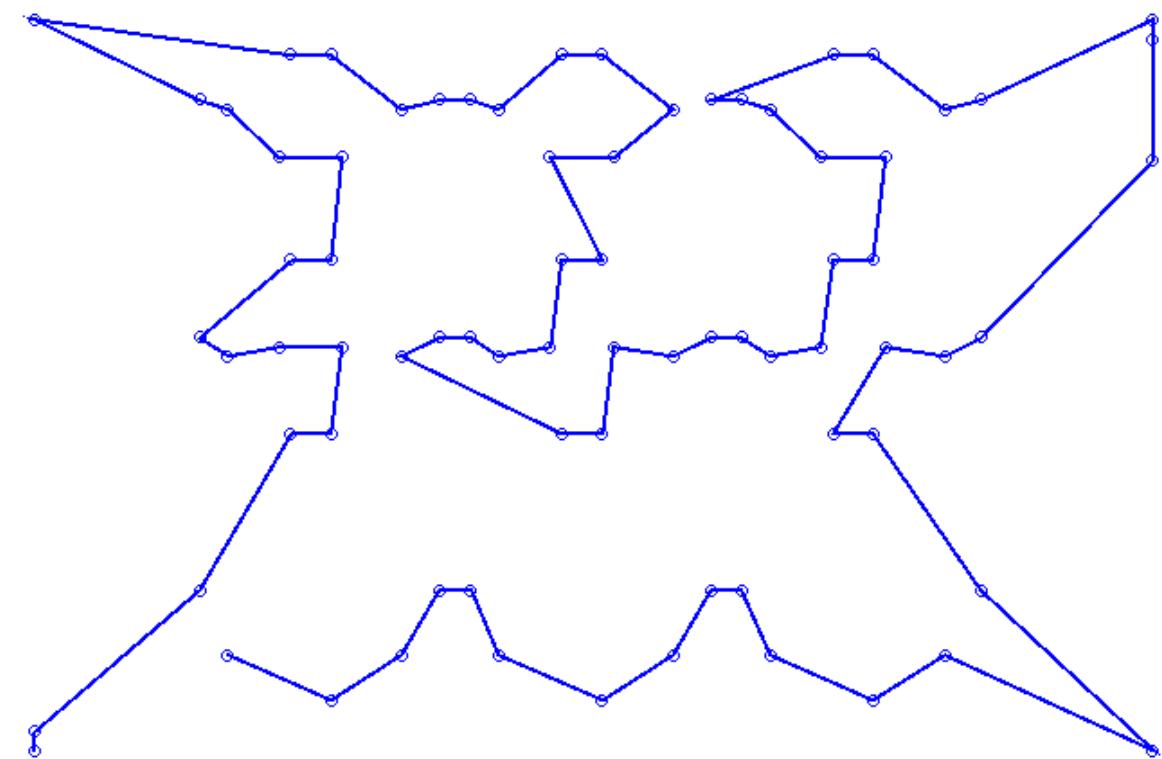
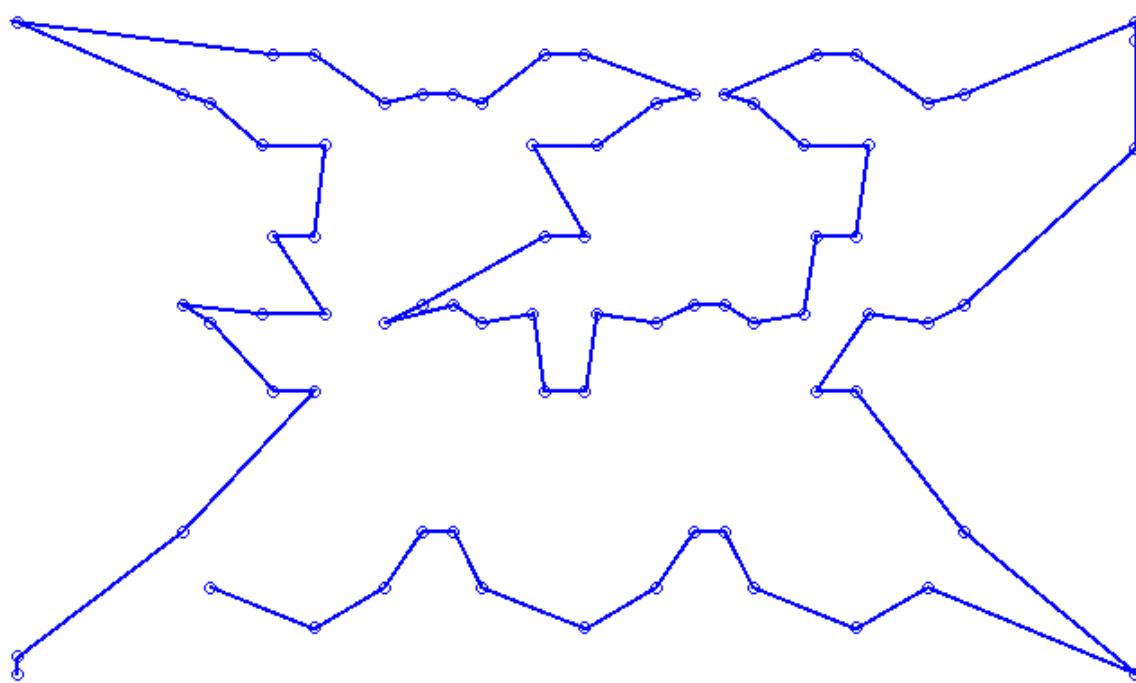
Construction algorithms–Randomisation



Construction algorithms–Randomisation

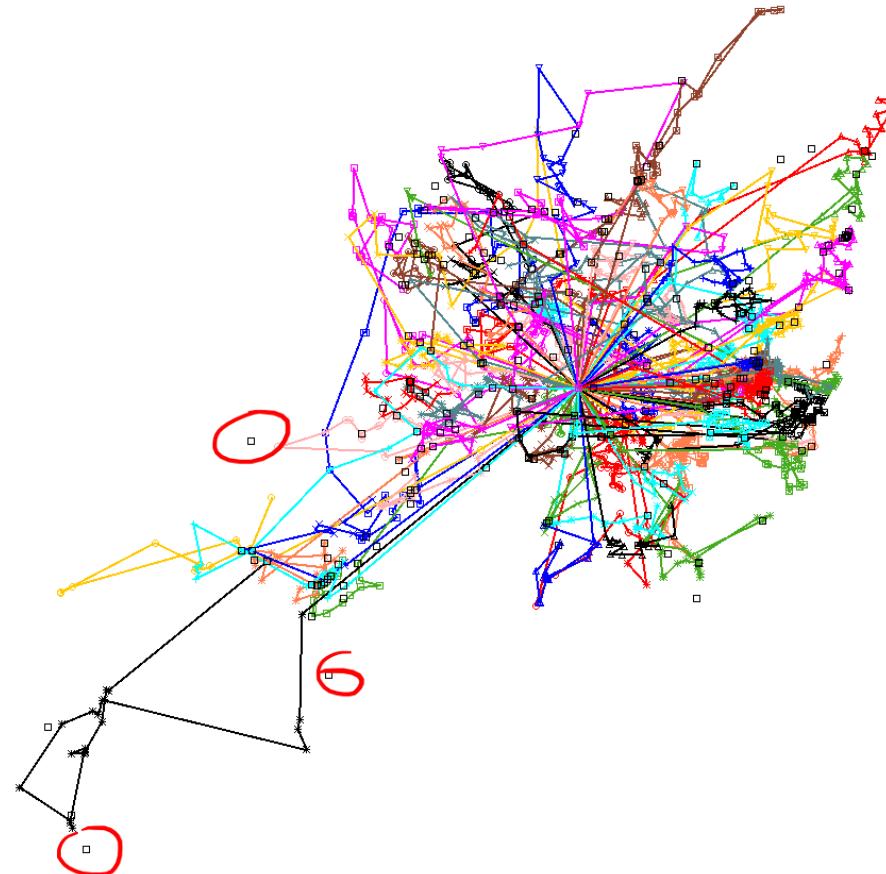
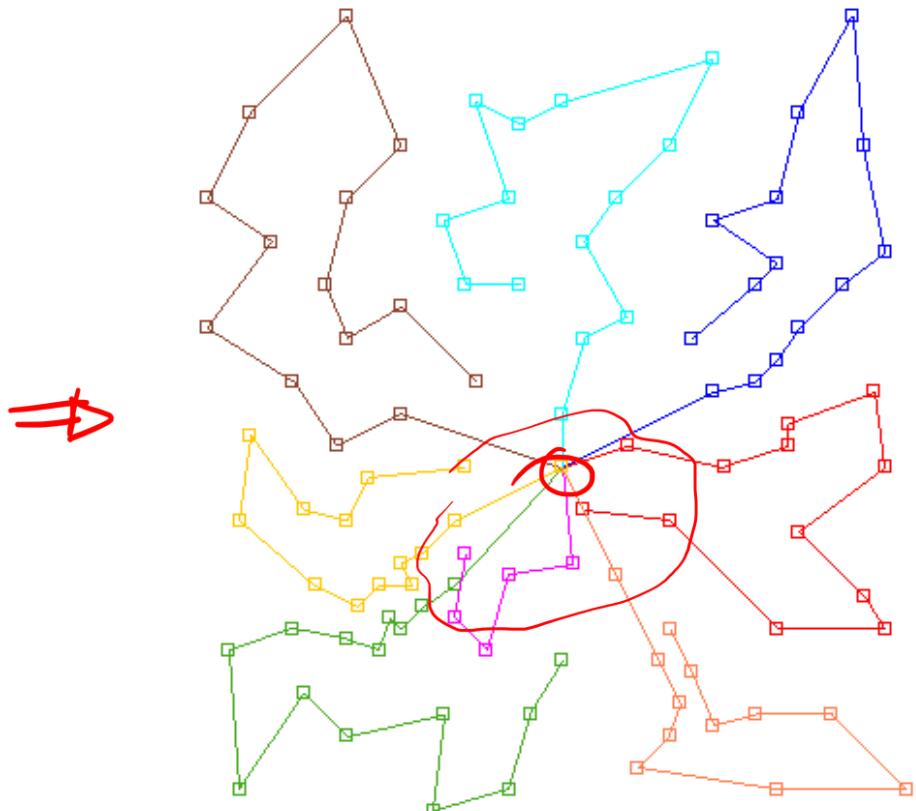


Construction algorithms–Randomisation

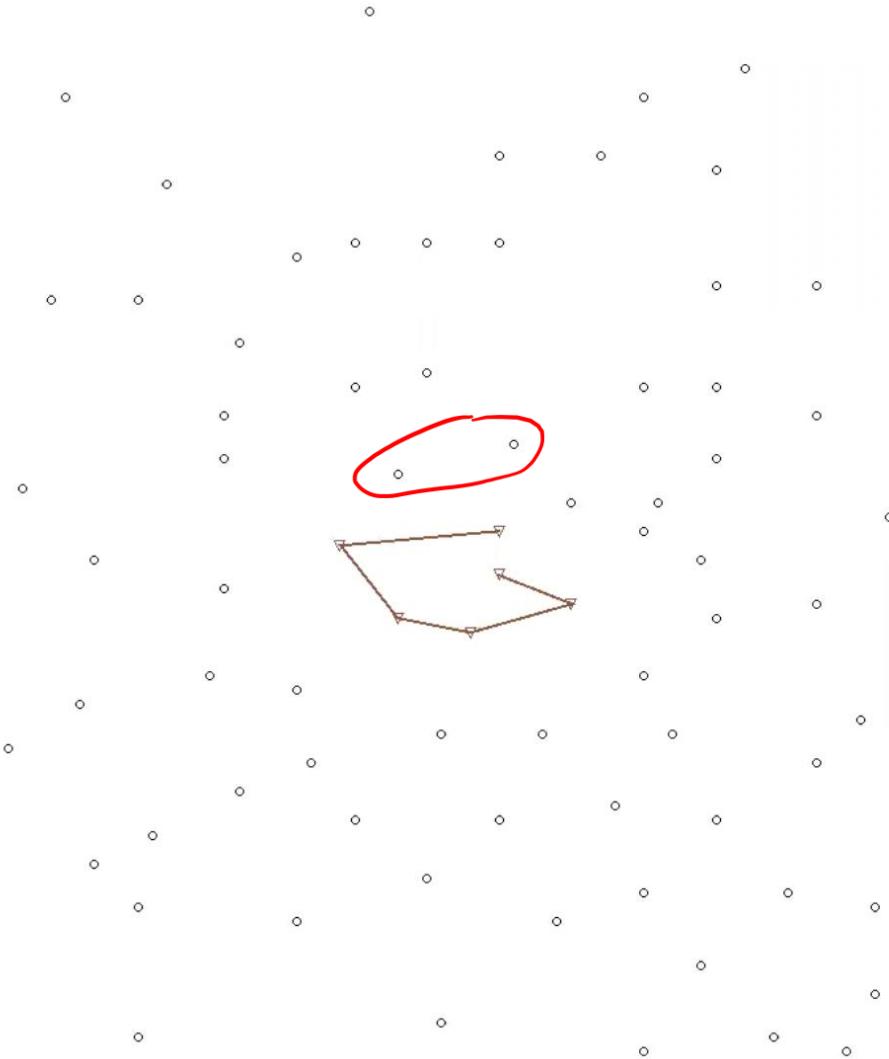


Vehicle Routing Problem (VRP)

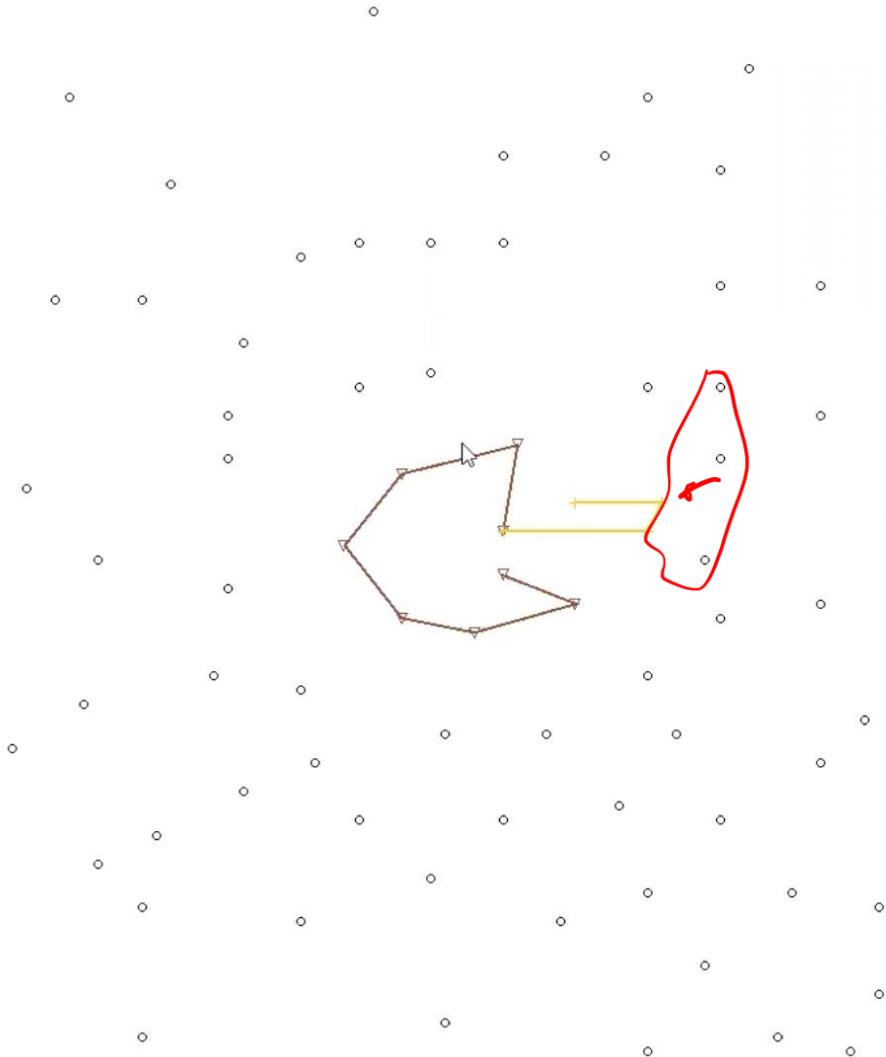
Given a set of customers, and a fleet of vehicles to make deliveries, find a set of routes that services all customers at minimum cost



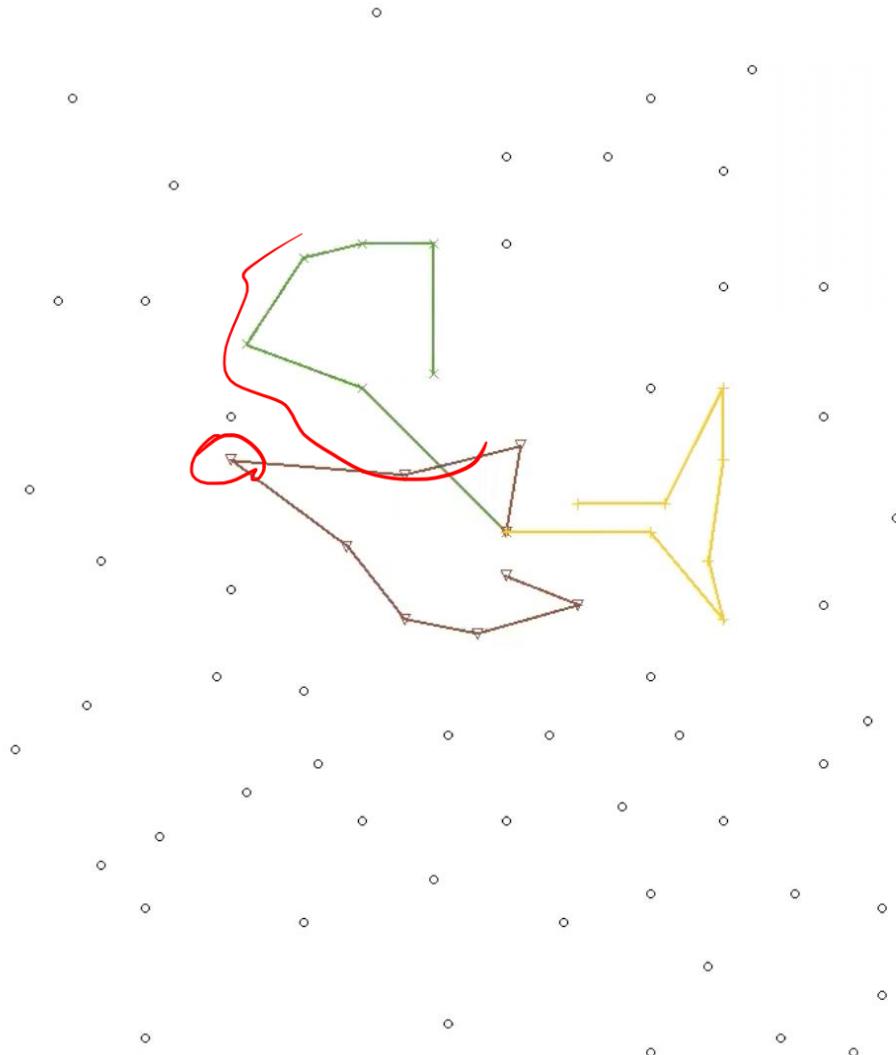
Minimum Insert Cost – VRP



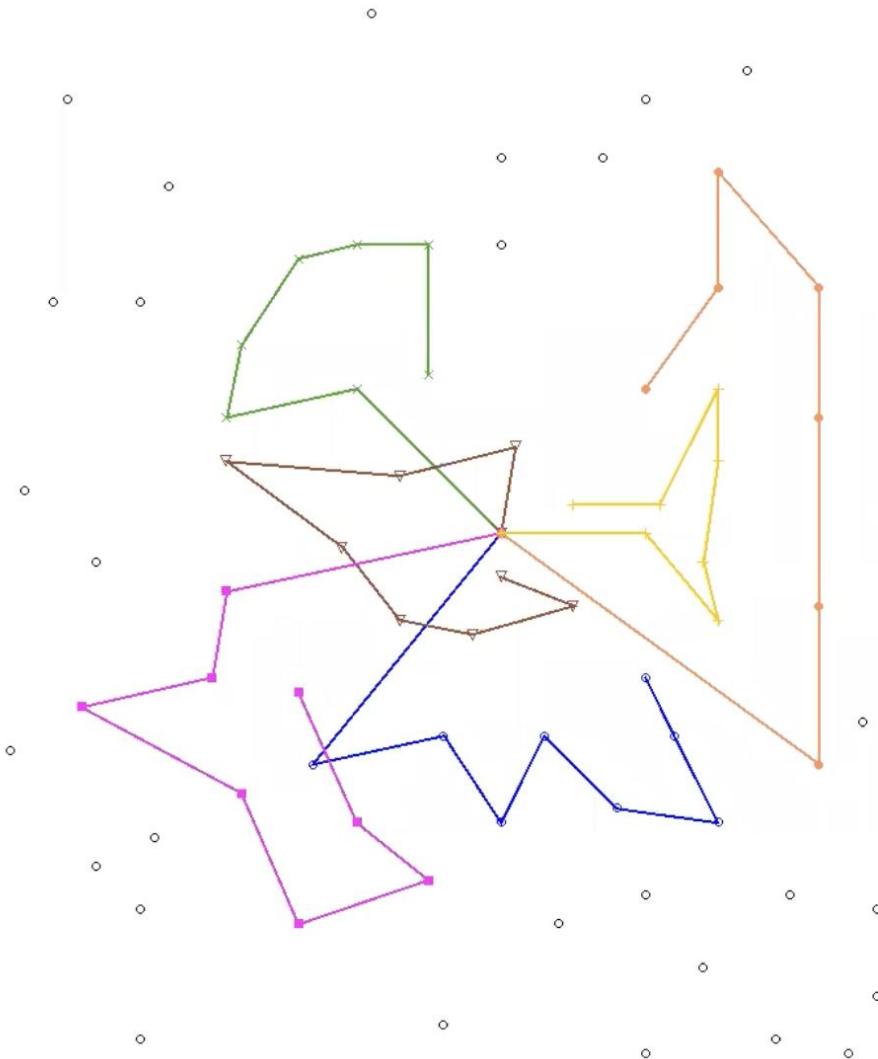
Minimum Insert Cost – VRP



Minimum Insert Cost – VRP

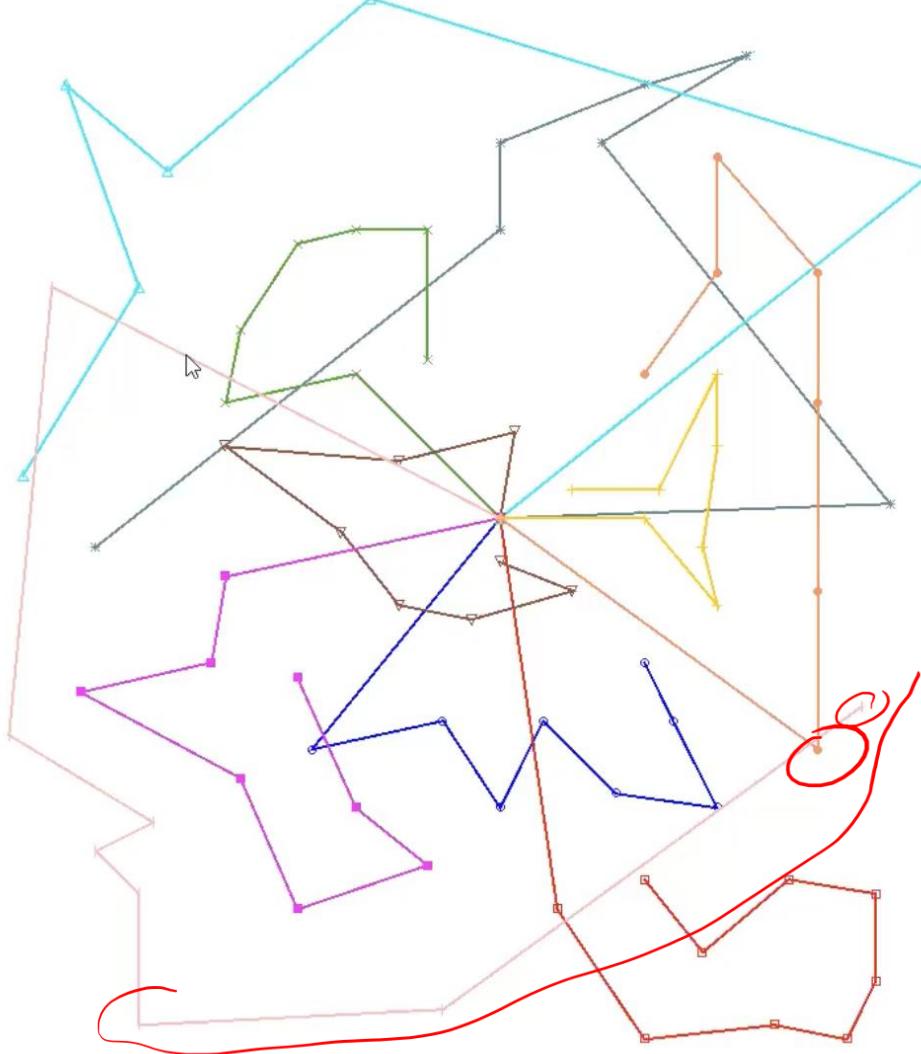


Minimum Insert Cost – VRP



Minimum Insert Cost – VRP

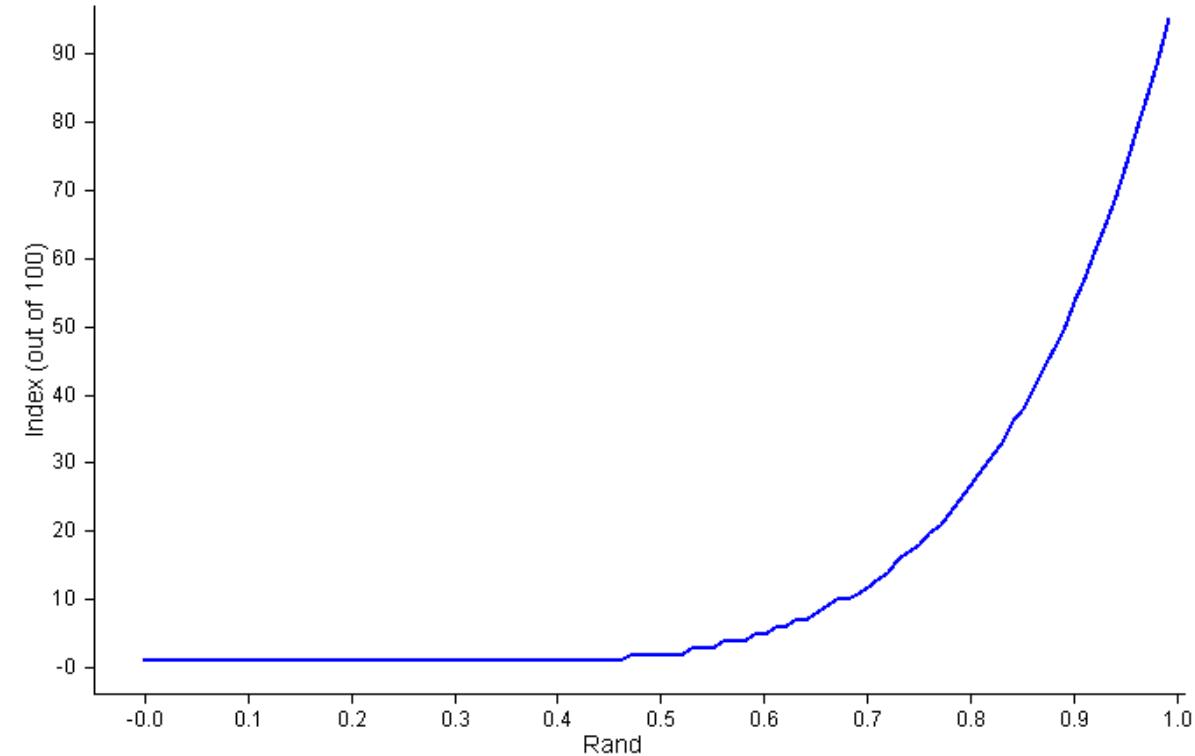
COMBINING
H/MY ← MATH OPT
AI: LNS ← NGx1 week
LARGE N. SEARCH



randomising the order
does not help 😞

Construction algorithms

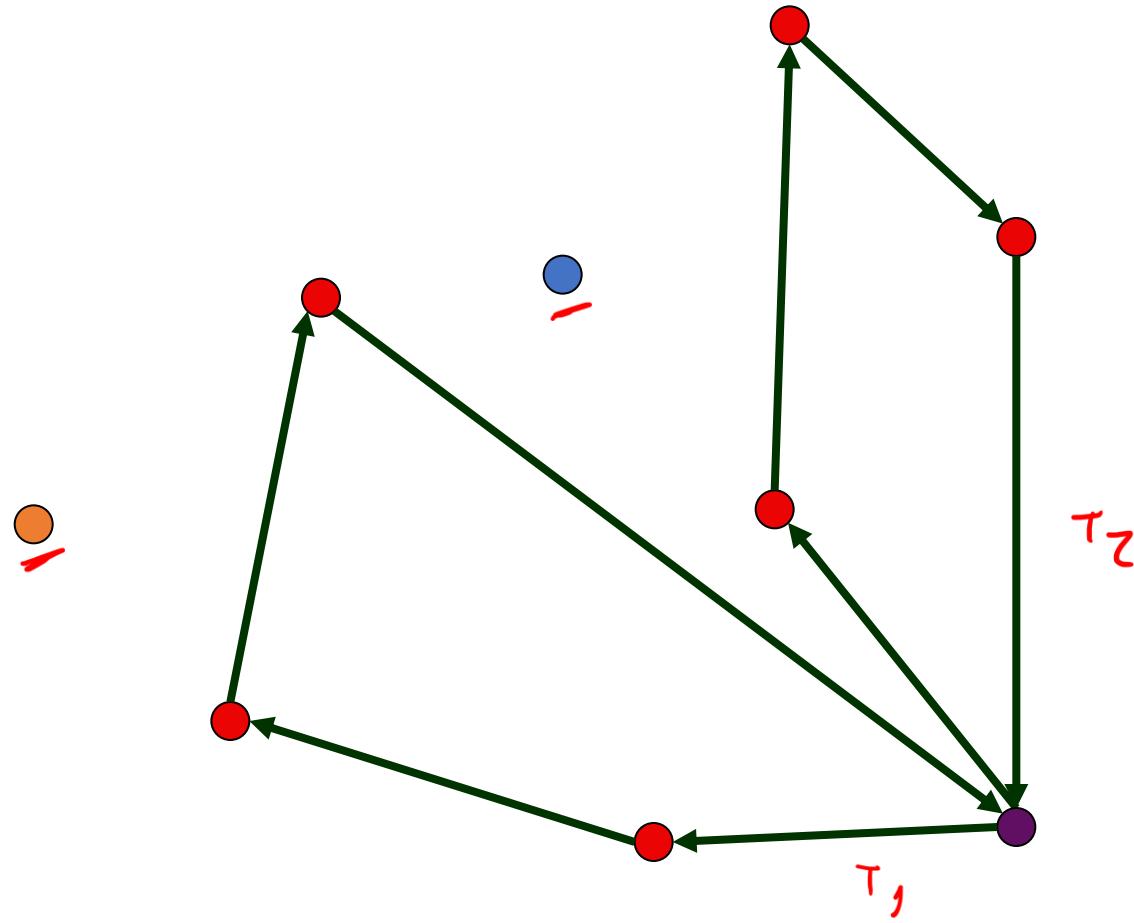
- Combine “greedy” and “random”
- Don’t always make the very “greediest” choice
- E.g. choose the r^{th} best of n choices
- $r = n * (\text{uniform}(0,1))^y$



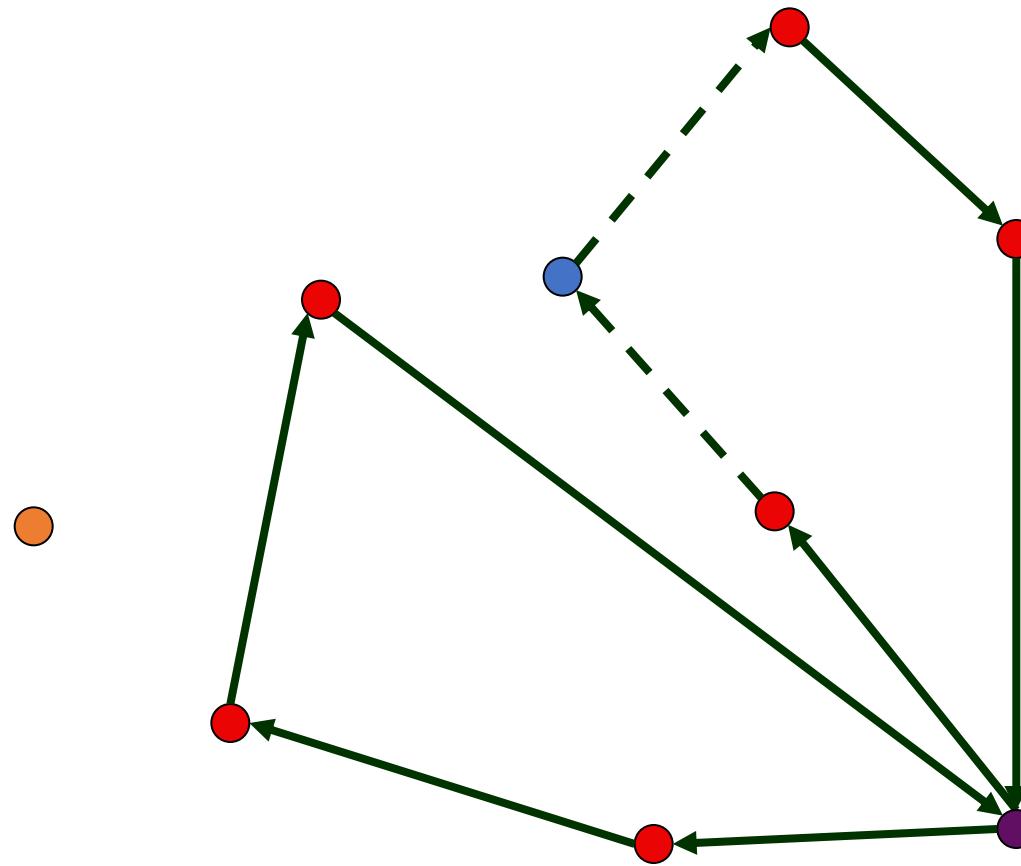
Construction algorithms

- We have seen that greedy algorithms lack “foresight”
- Next step: **Regret**

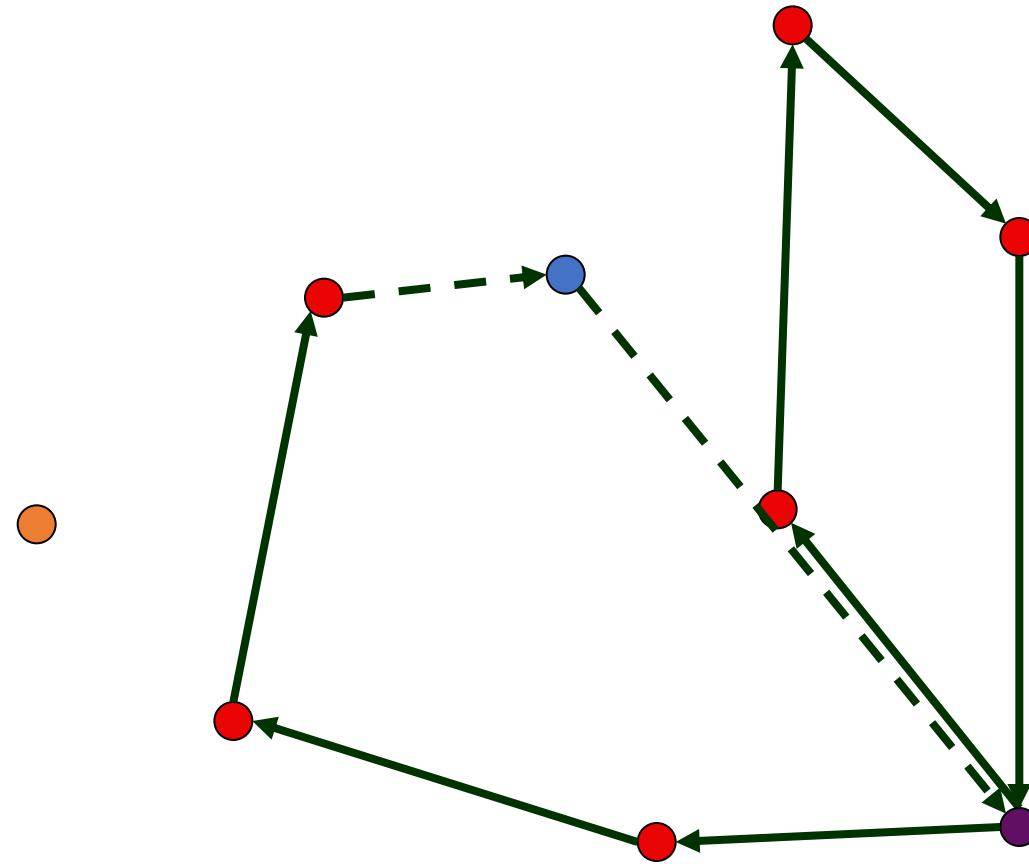
Regret – VRP



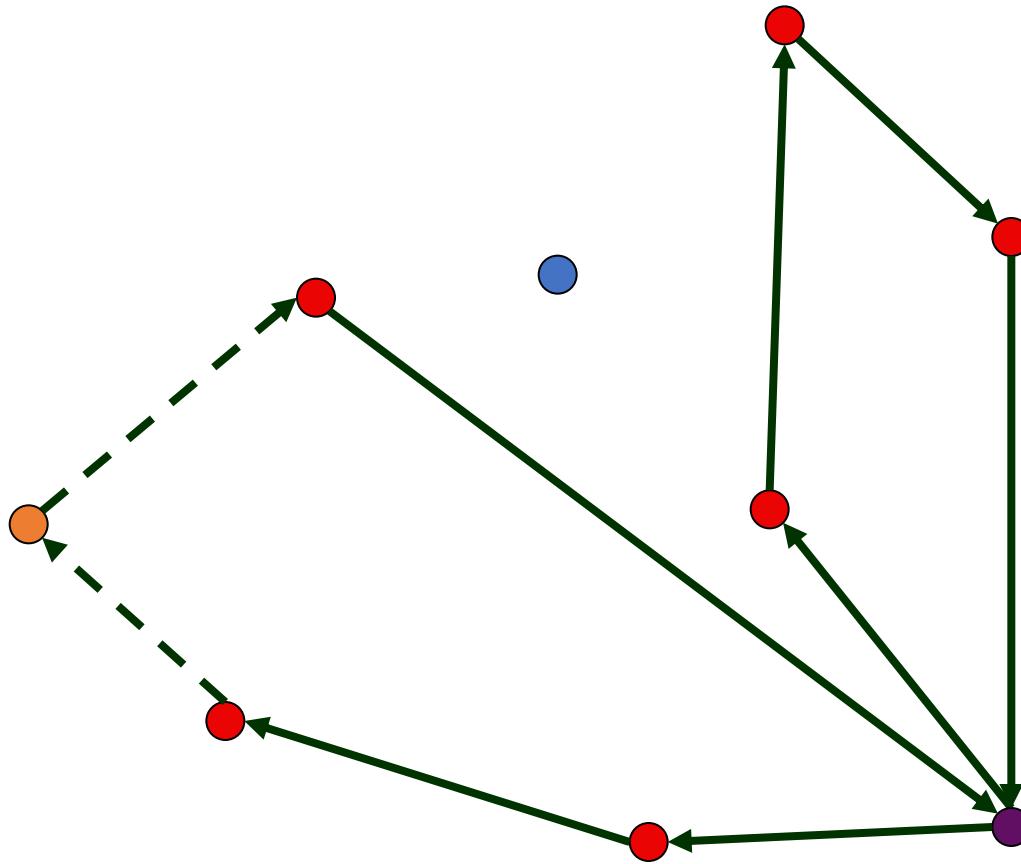
Regret – VRP



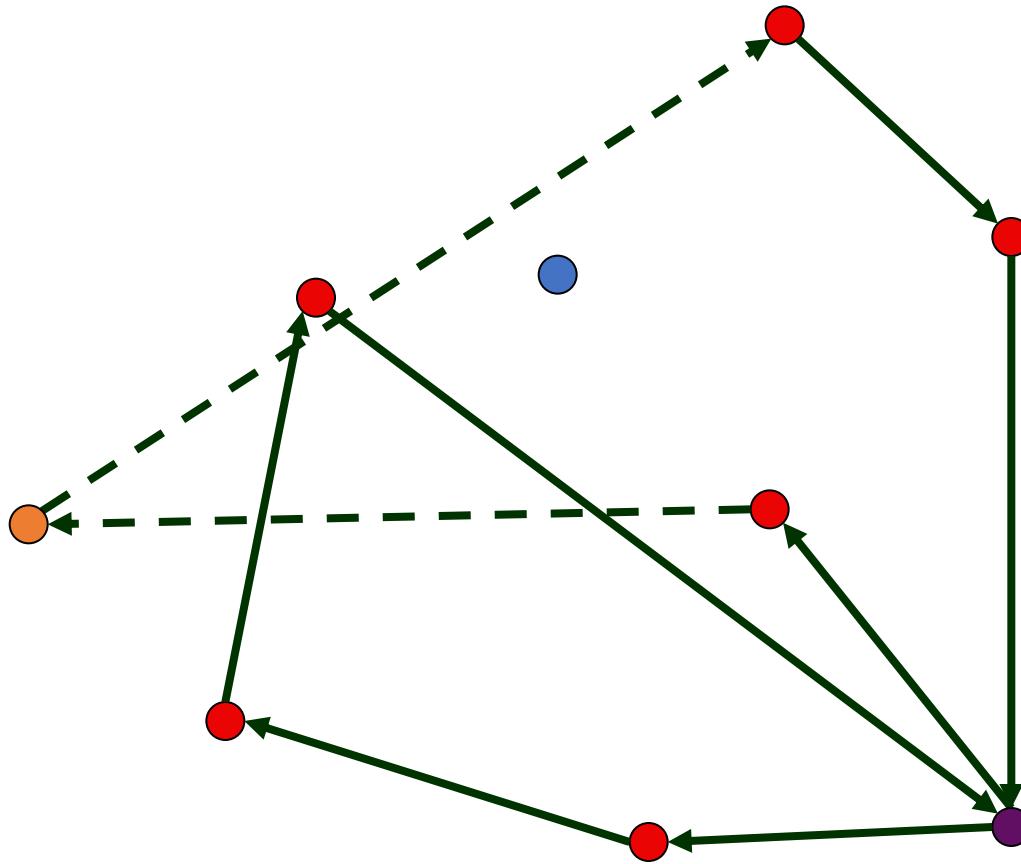
Regret – VRP



Regret – VRP



Regret – VRP



Regret

i-th



$$\begin{aligned}\text{regret}(i) &= C(\text{insert } i \text{ in } \underline{\text{2}^{\text{nd}}\text{-best route}}) - C(\text{insert } i \text{ in } \underline{\text{best route}}) \\ &= f(2, i) - f(1, i)\end{aligned}$$

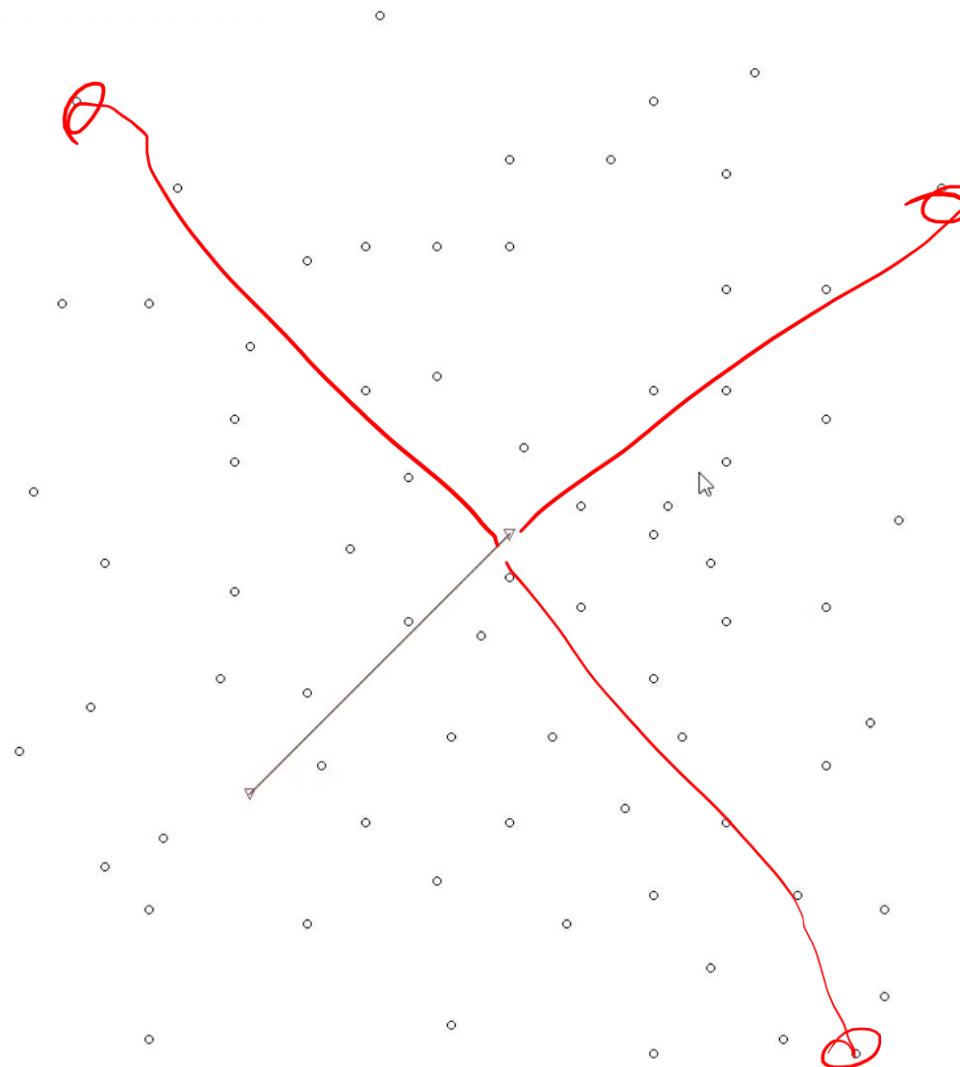
Select the choice i with largest regret

- VRP: next customer to be served is the one with largest regret

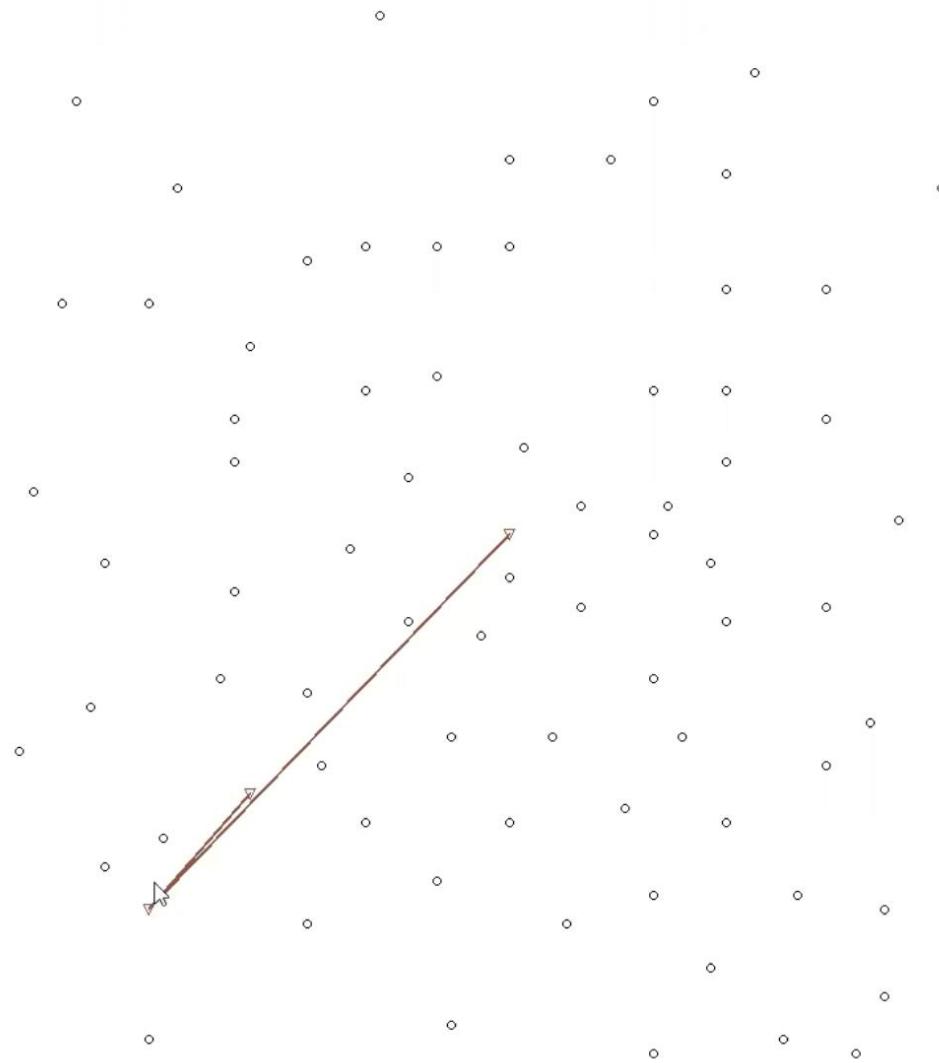
Can be extended to the k-th best routes:

$$\underline{k\text{-regret}}(i) = \sum_{j=2}^k (f(j, i) - f(1, i))$$

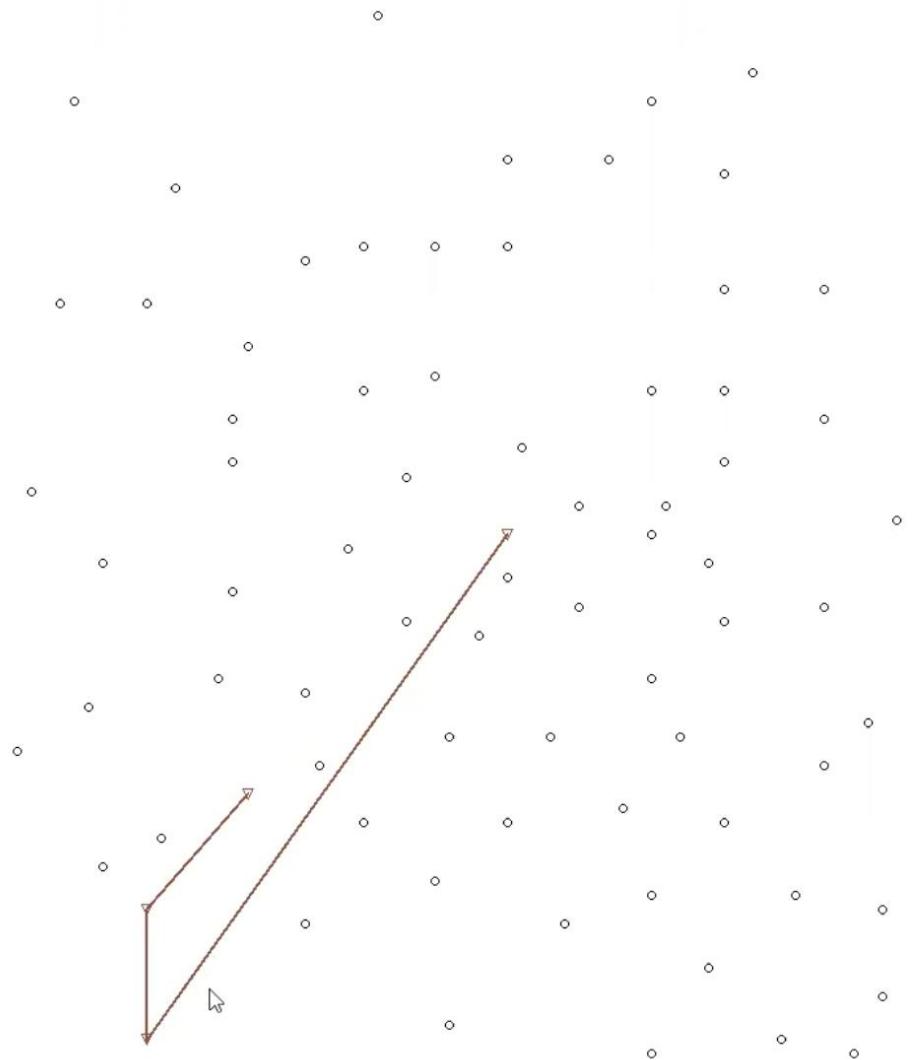
Regret



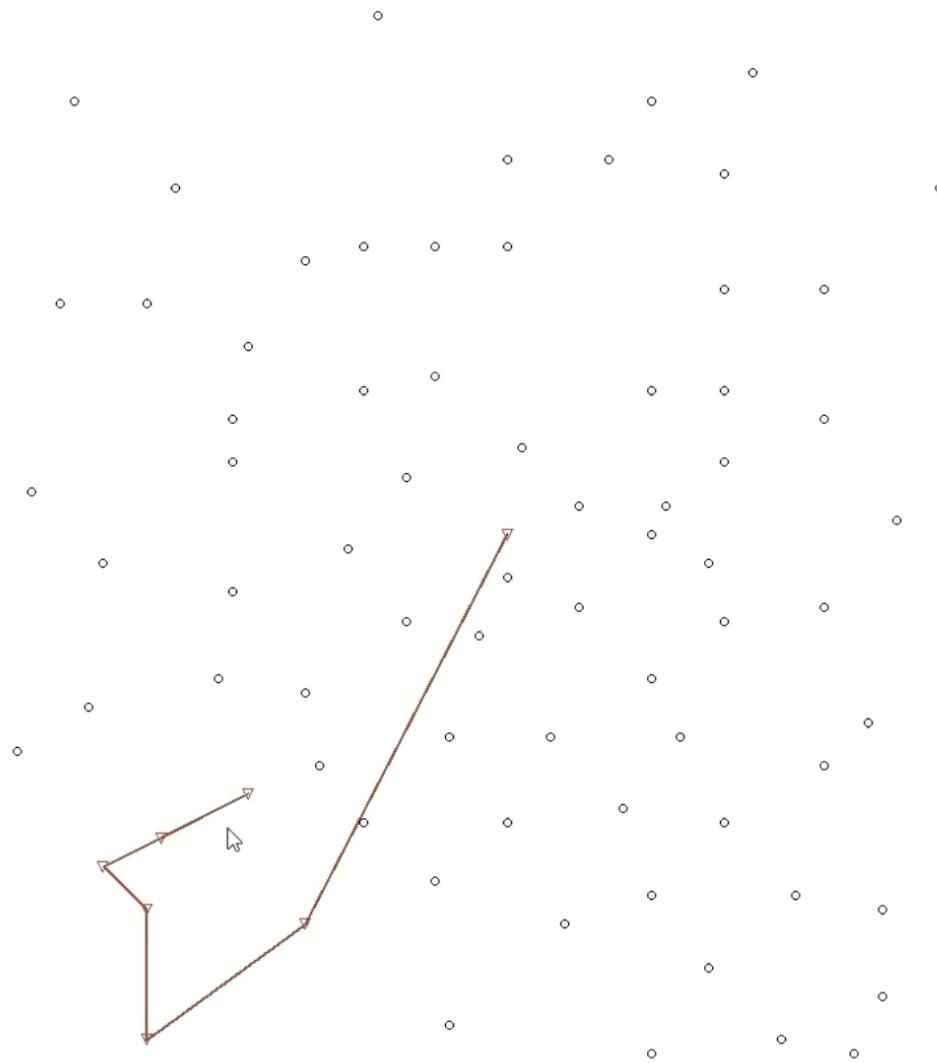
Regret



Regret



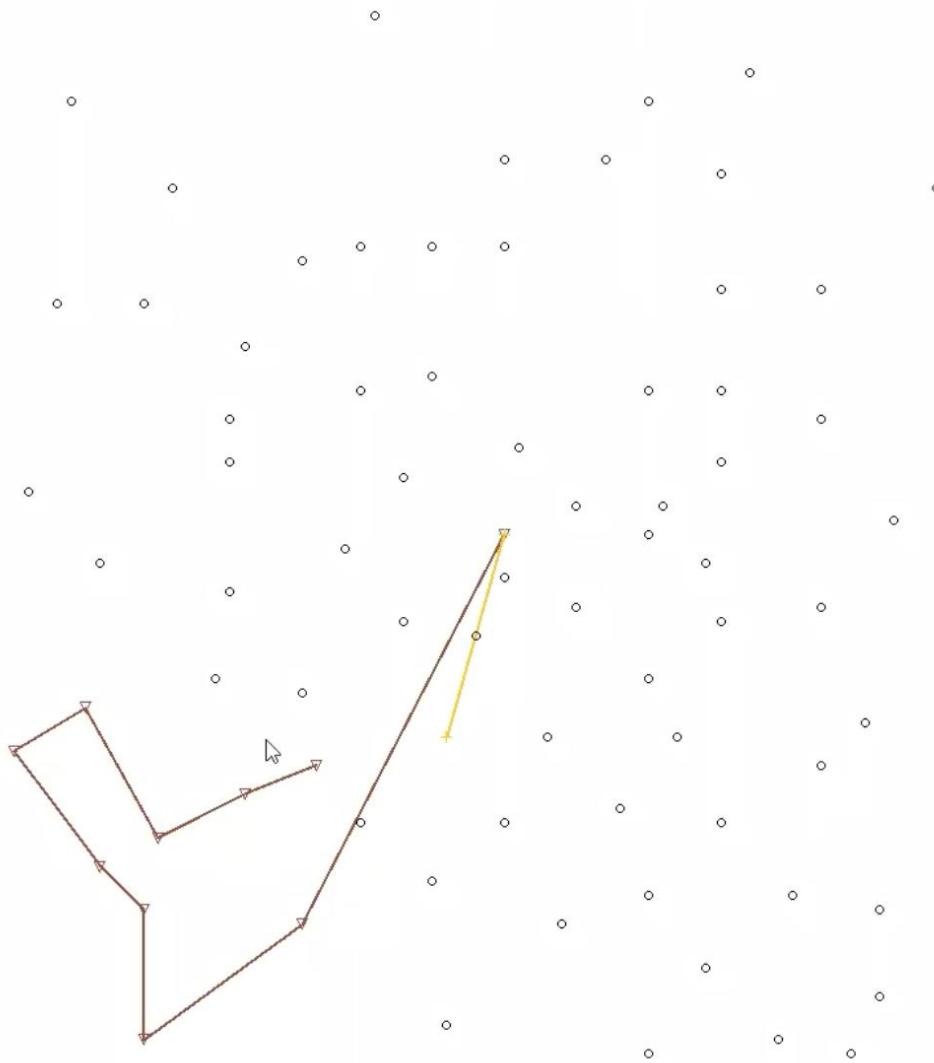
Regret



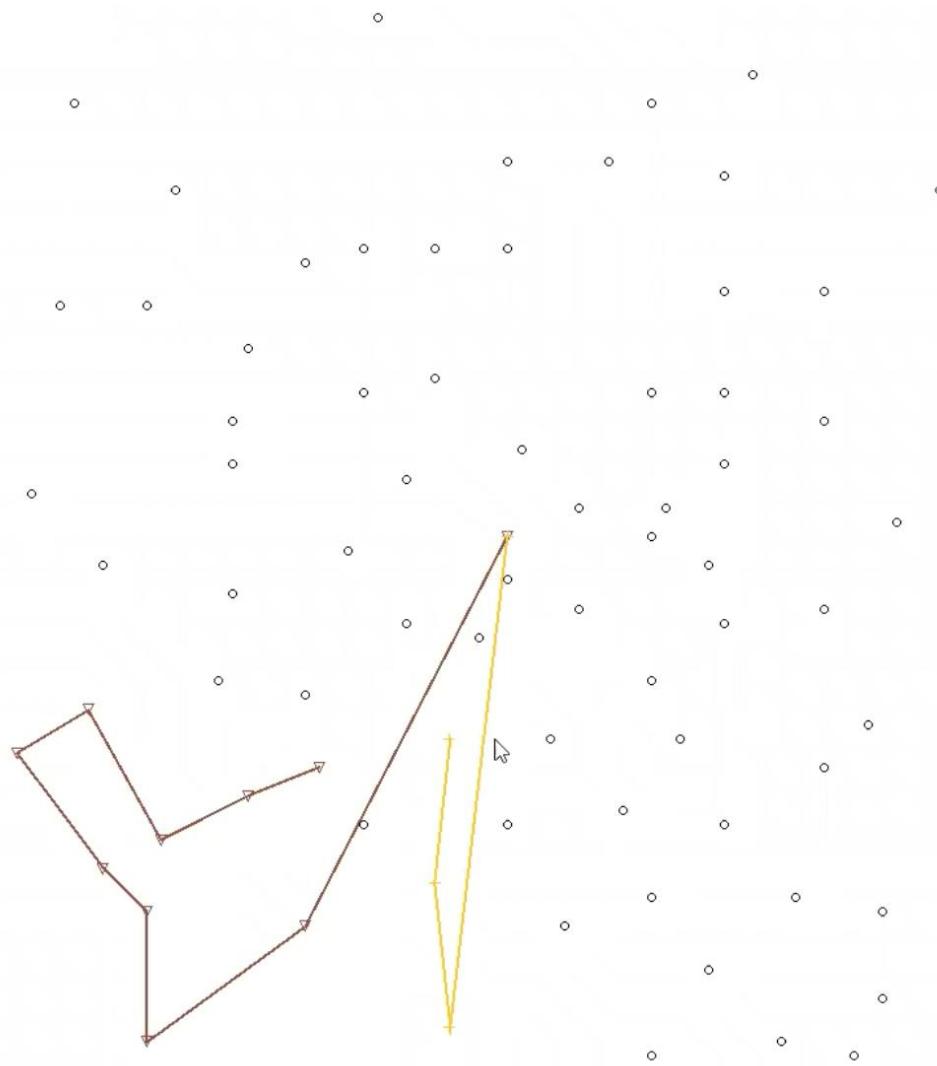
Regret



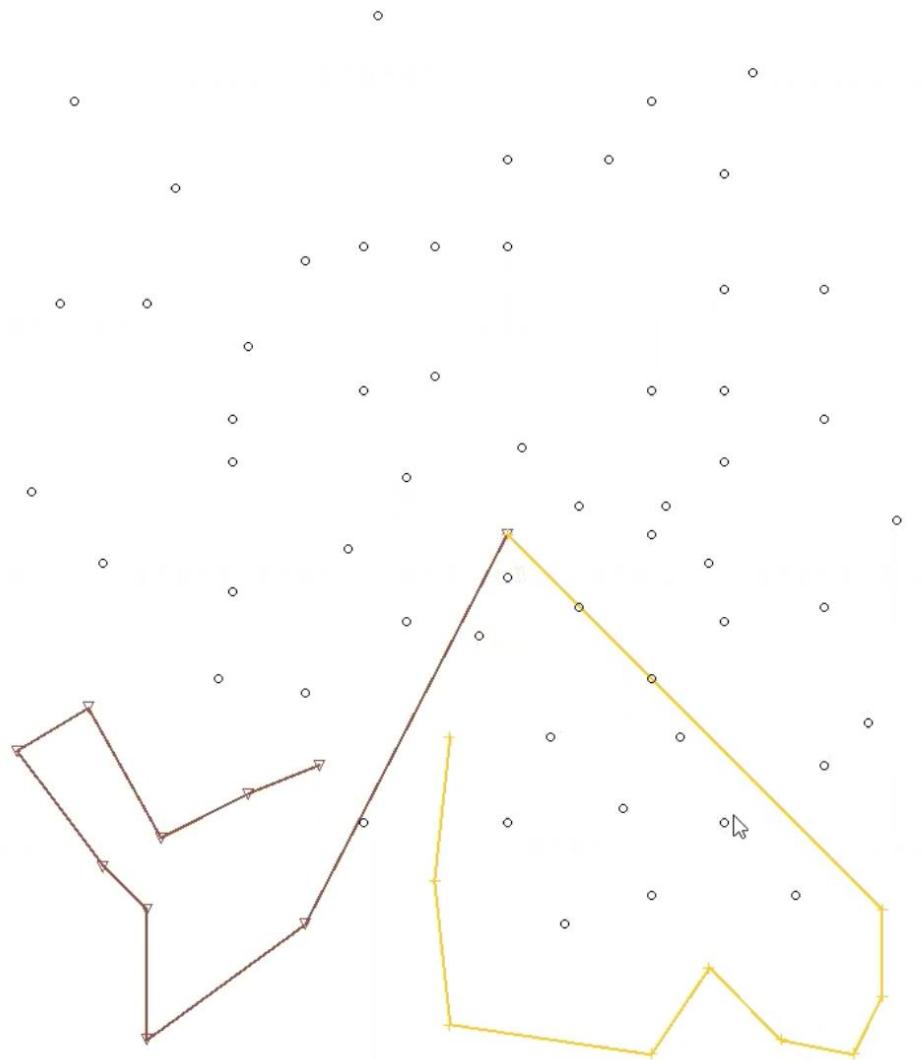
Regret



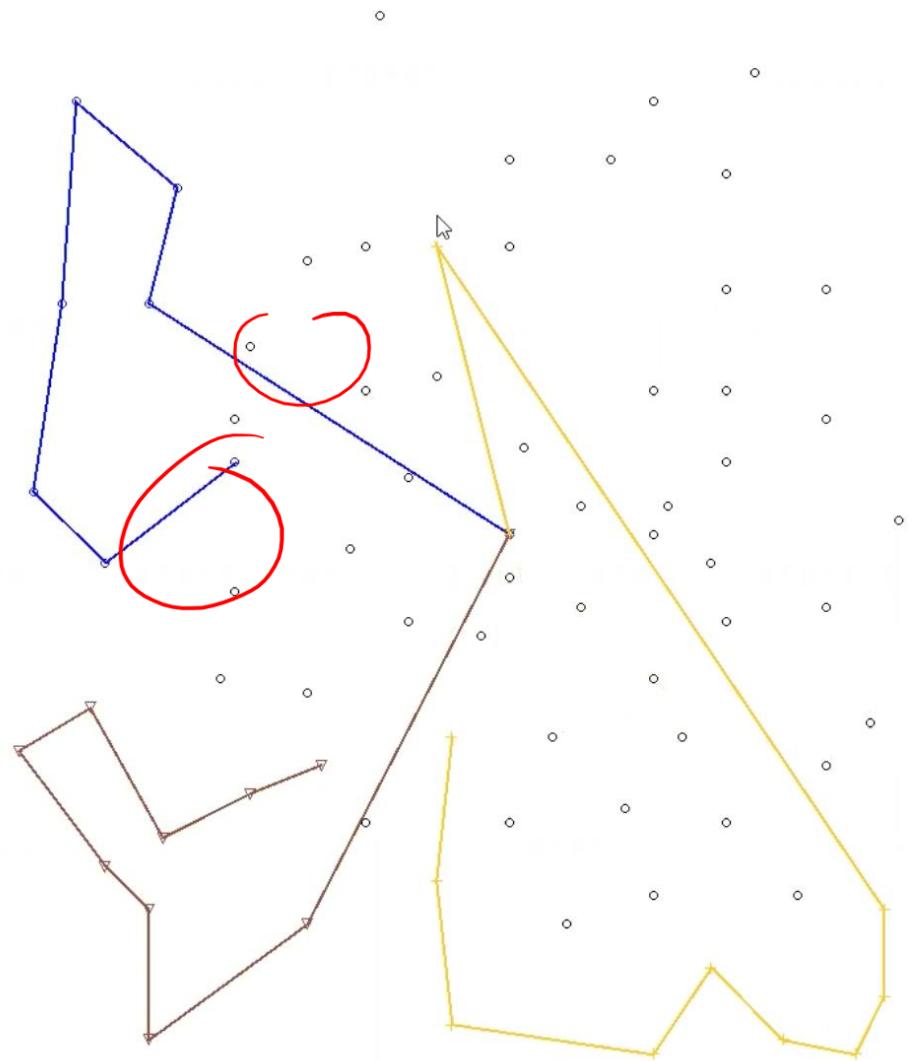
Regret



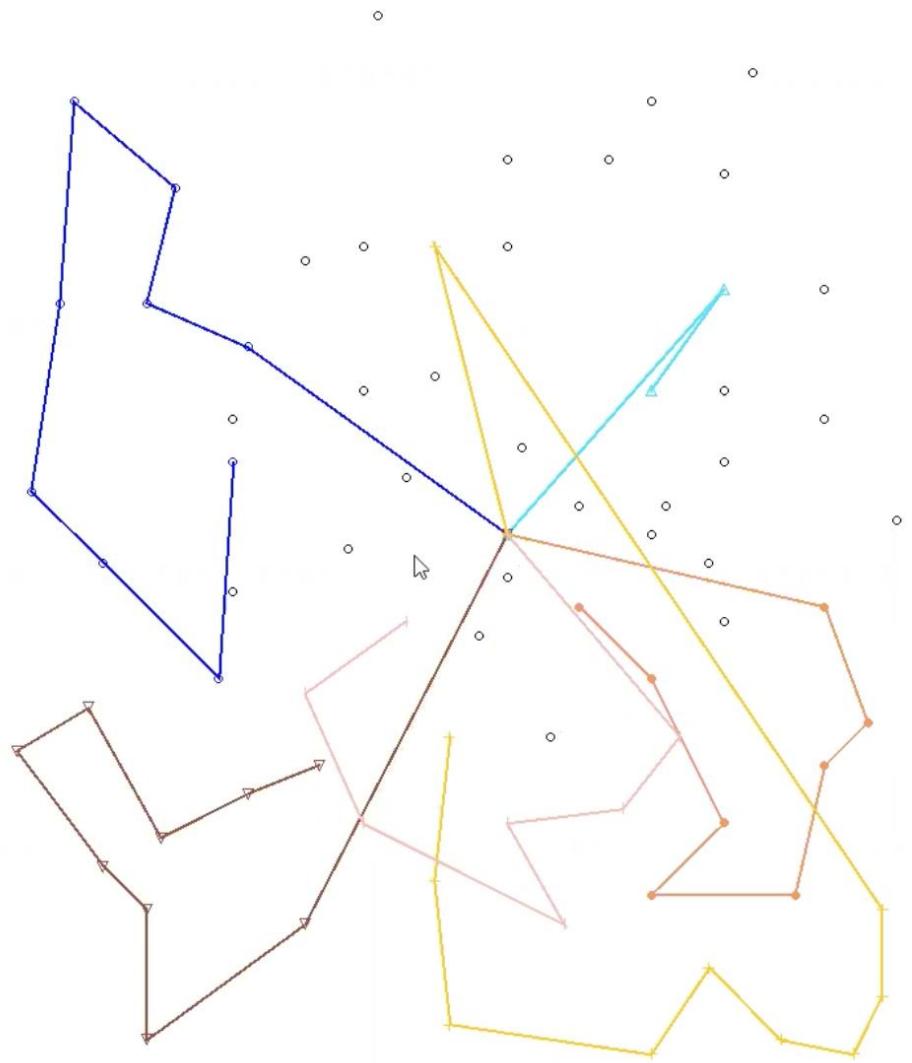
Regret



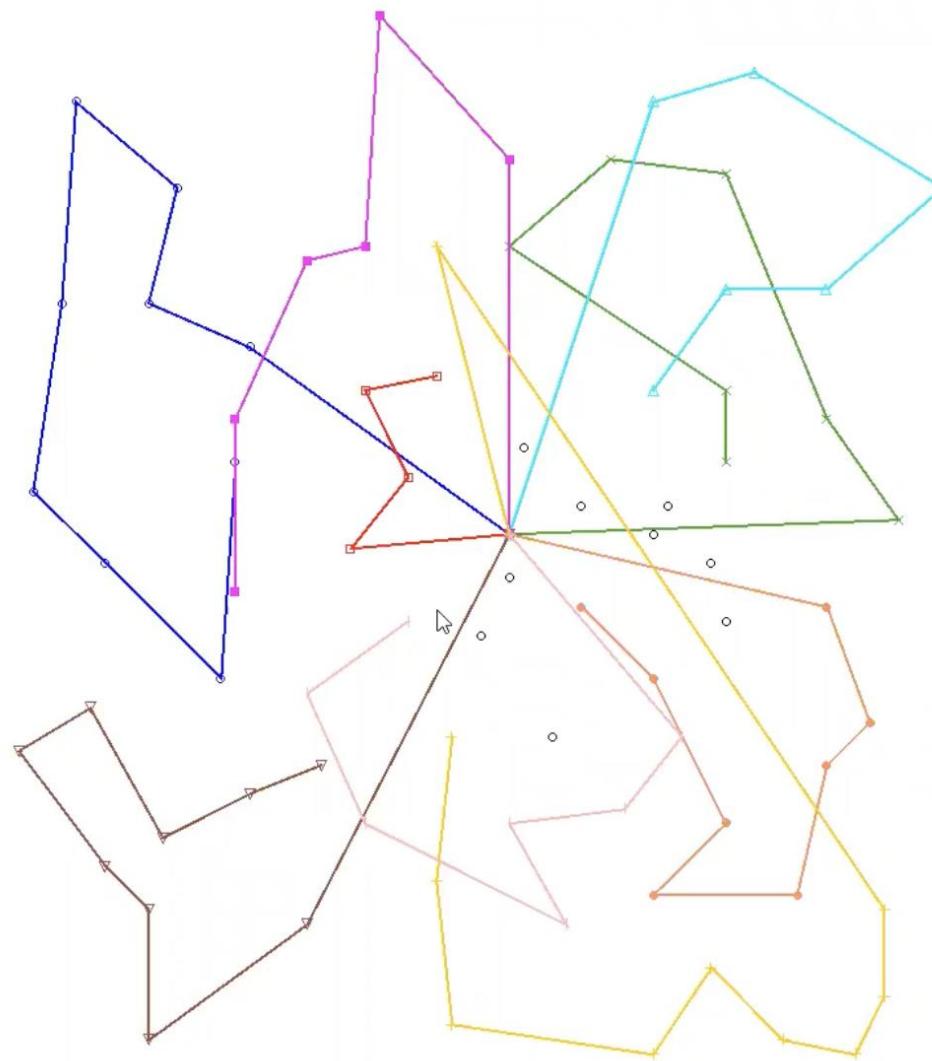
Regret



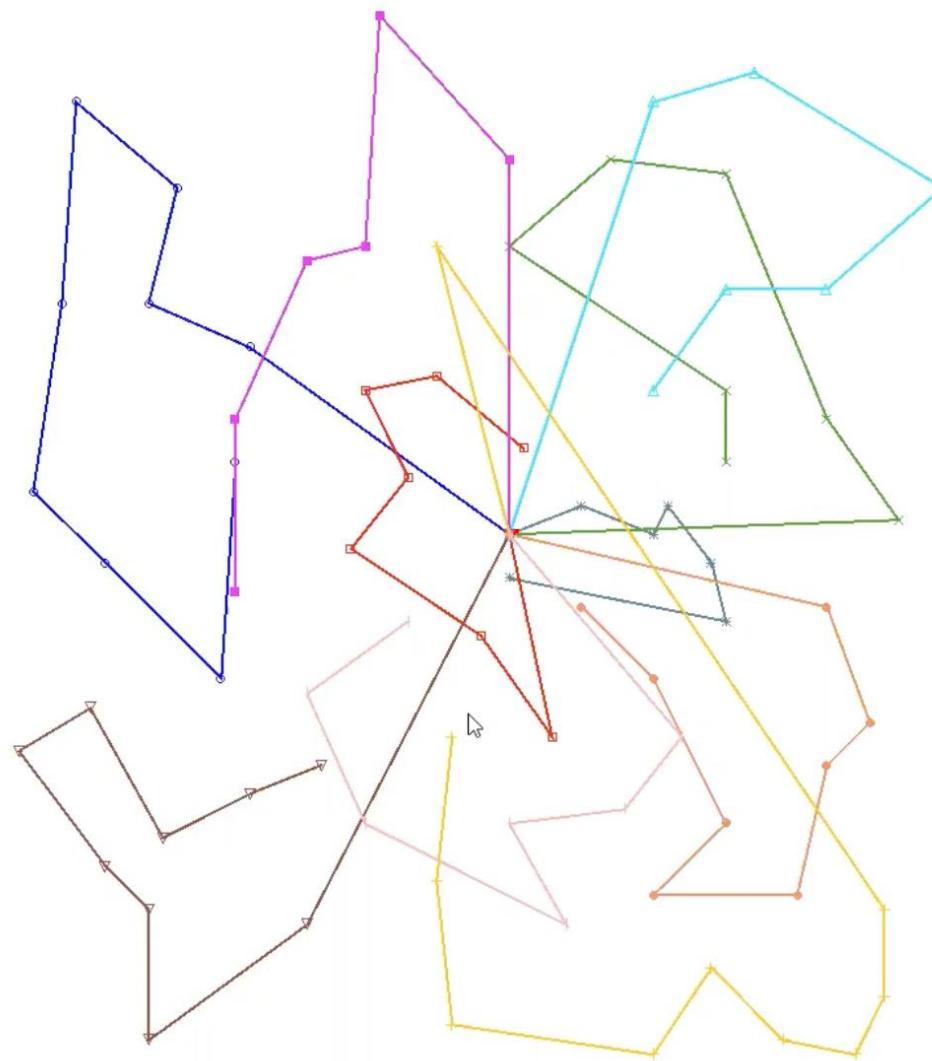
Regret



Regret



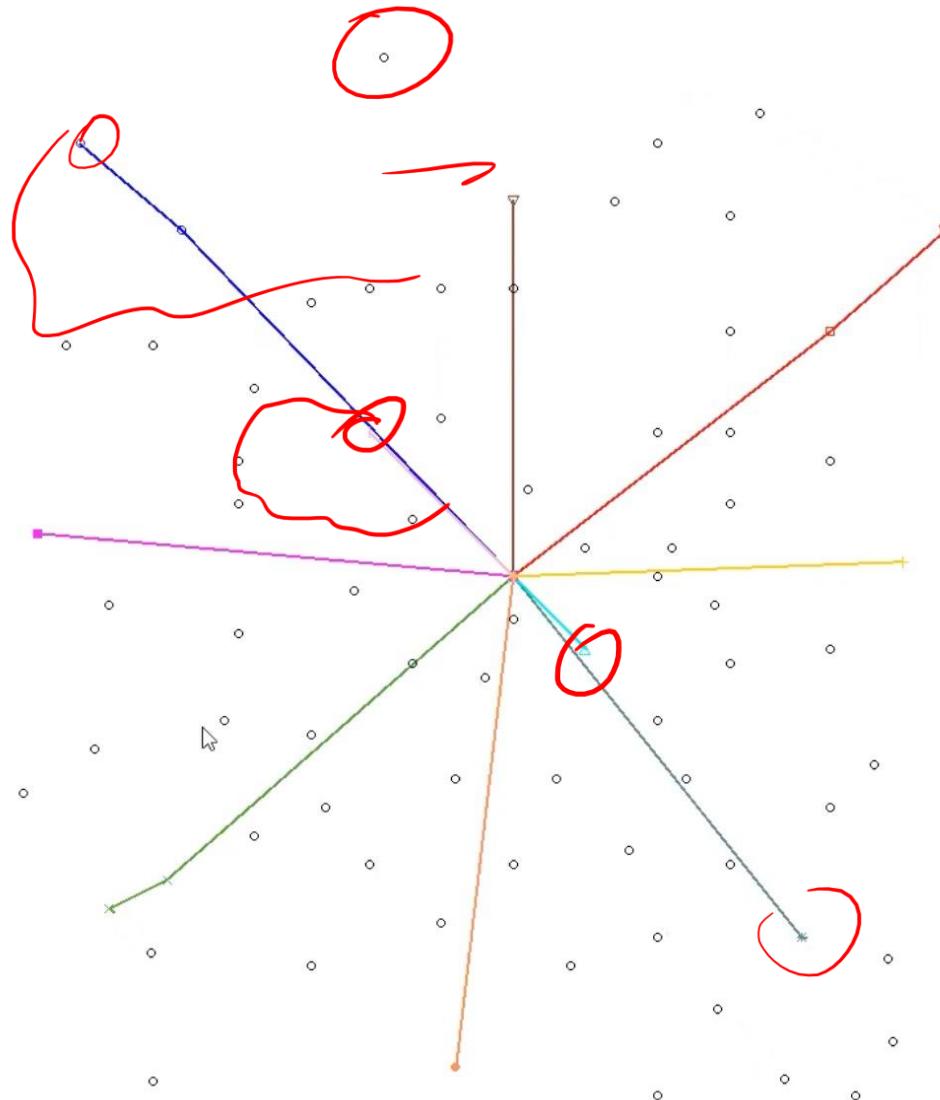
Regret



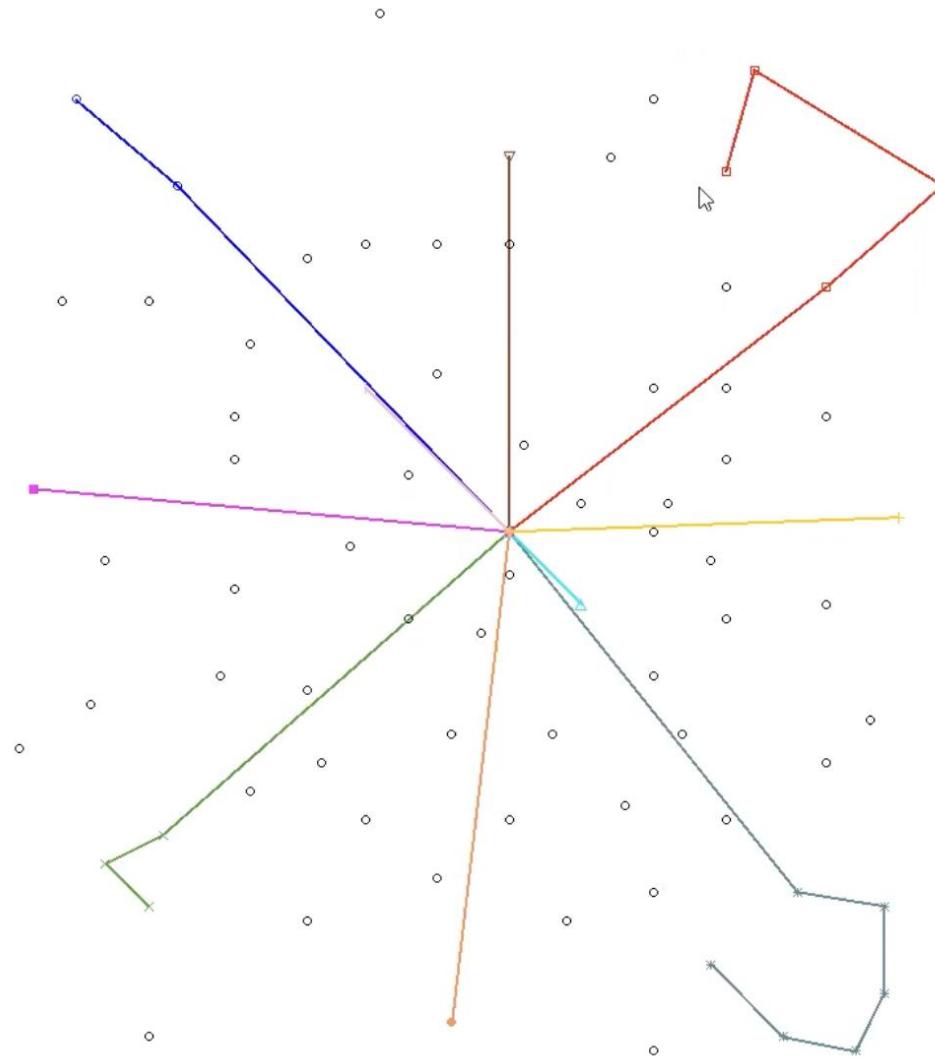
Regret

- And a little more structure...
- Regret works best when some structure is already present
- Use Seeds
 - Seed 1: The customer with max (dist to depot)
 - Seed 2: The customer with max (dist to Seed 1)
 - Seed 3: The customer with max (dist to Seed 1 + dist to Seed 2)
 - Seed 4: The customer with max (dist to Seed 1 + dist to Seed 2 + dist to Seed 3)
 - Etc

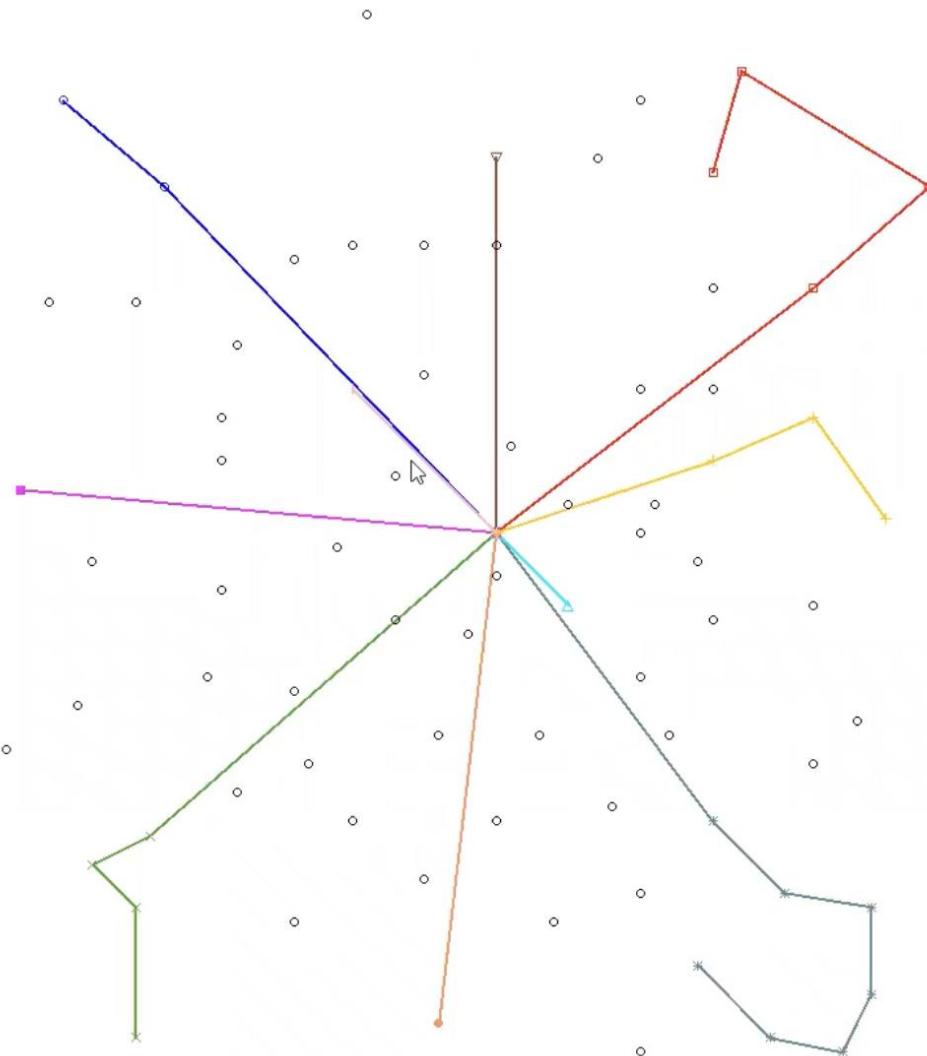
Regret with seeds



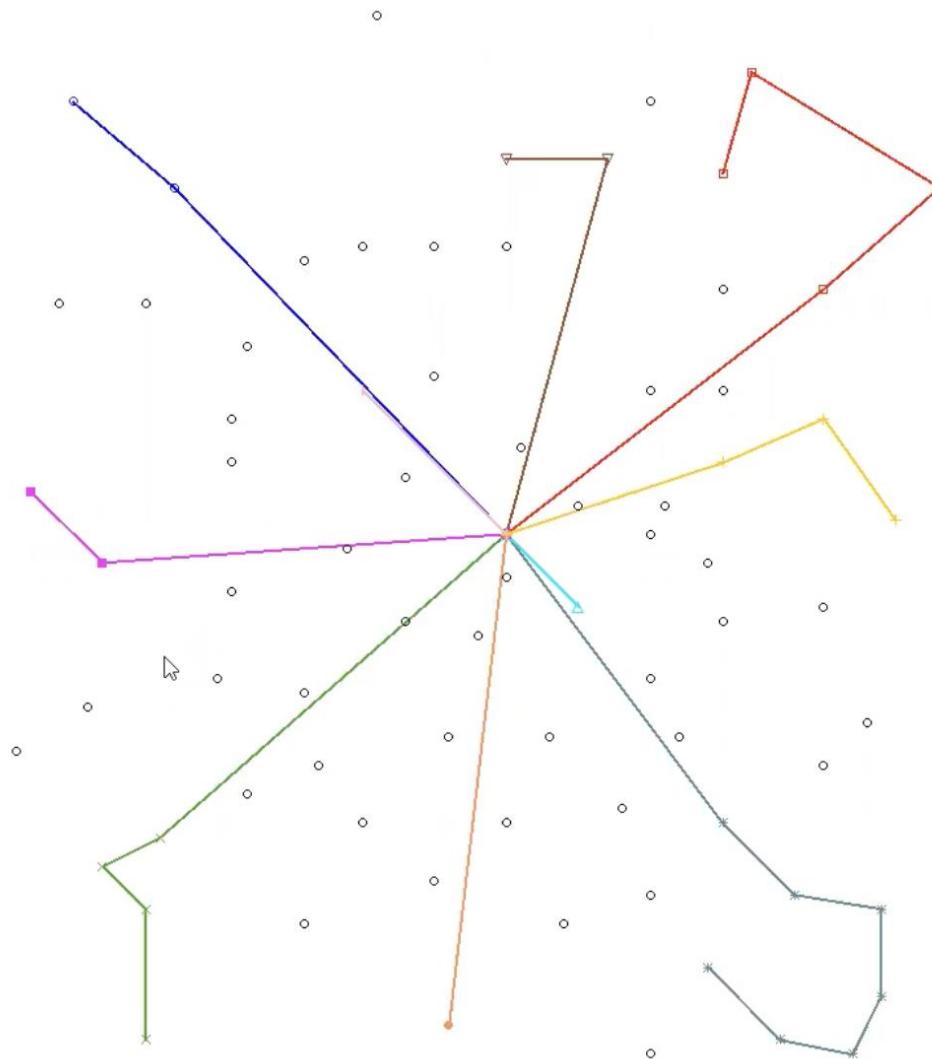
Regret with seeds



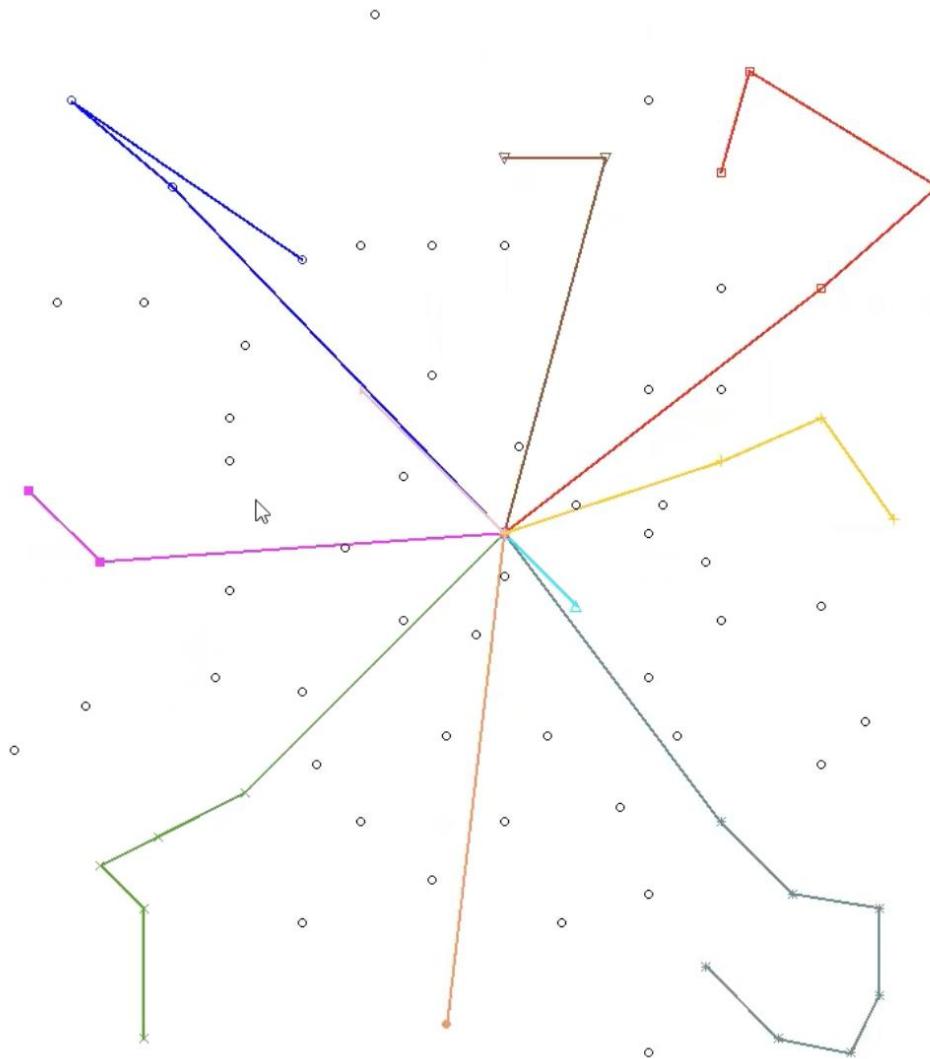
Regret with seeds



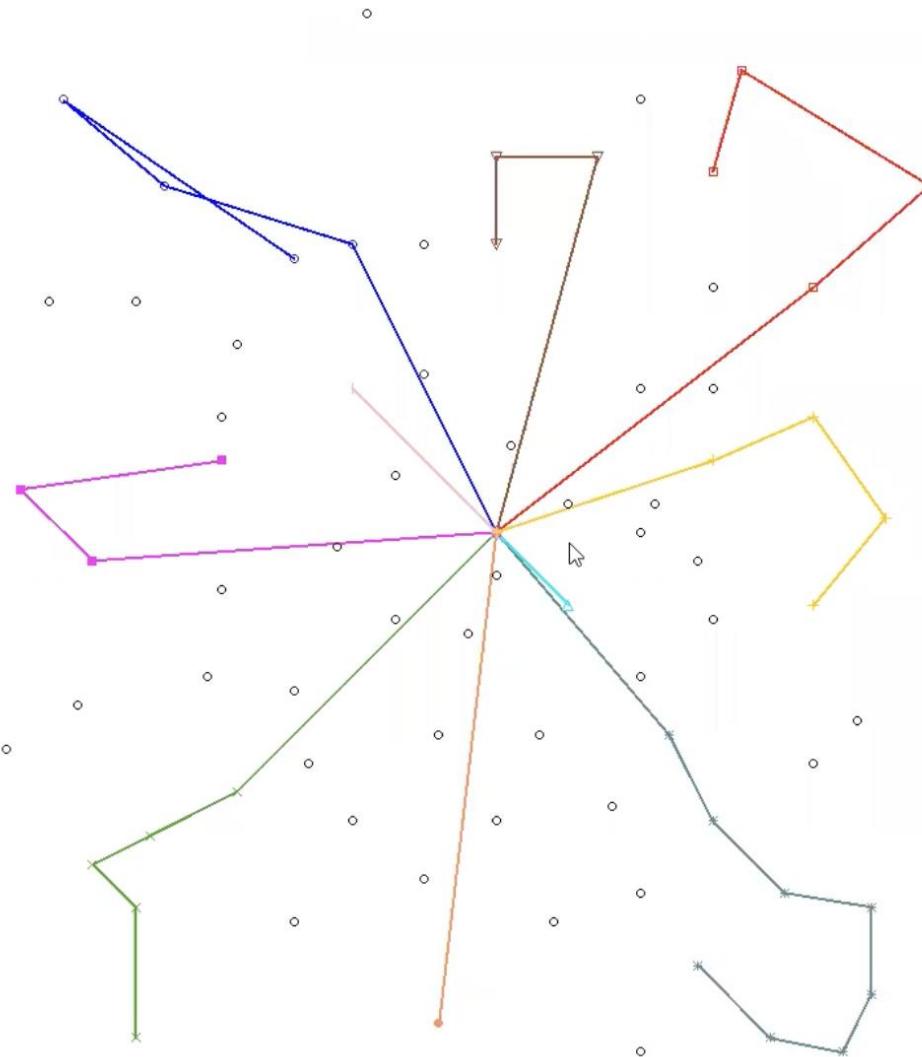
Regret with seeds



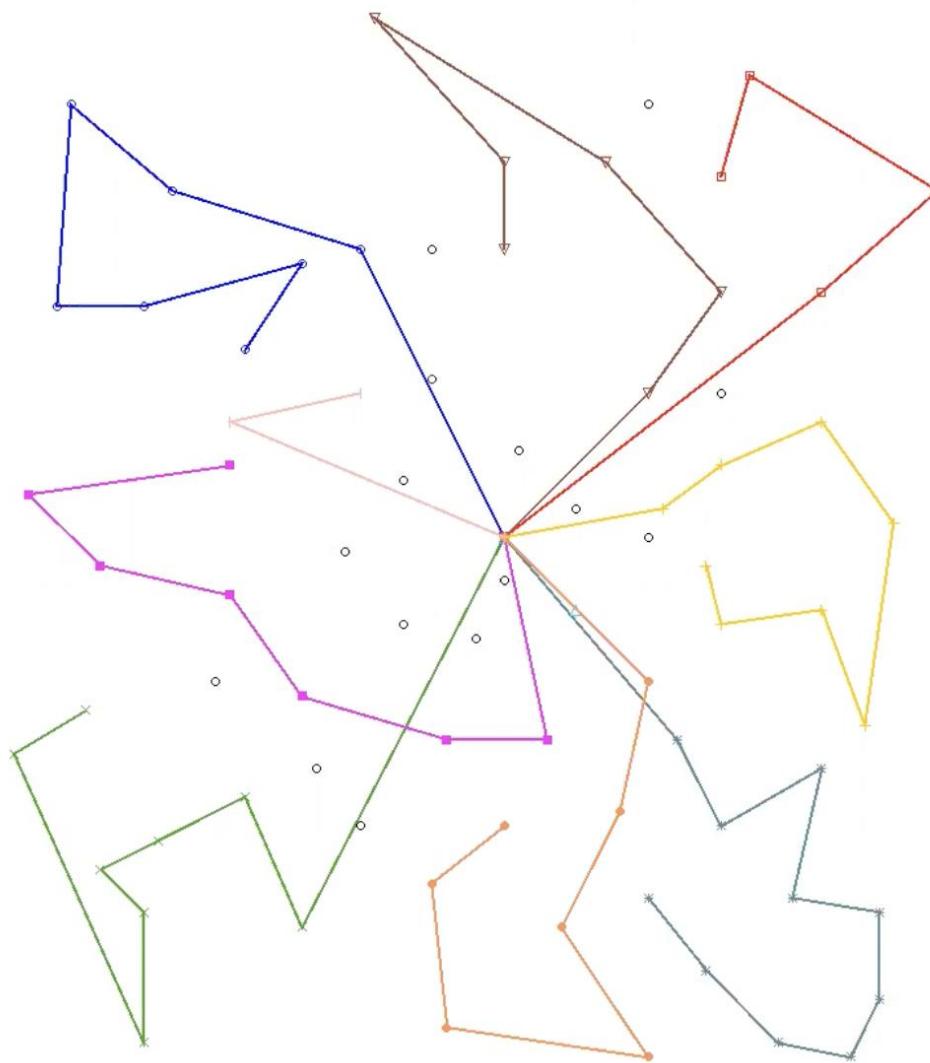
Regret with seeds



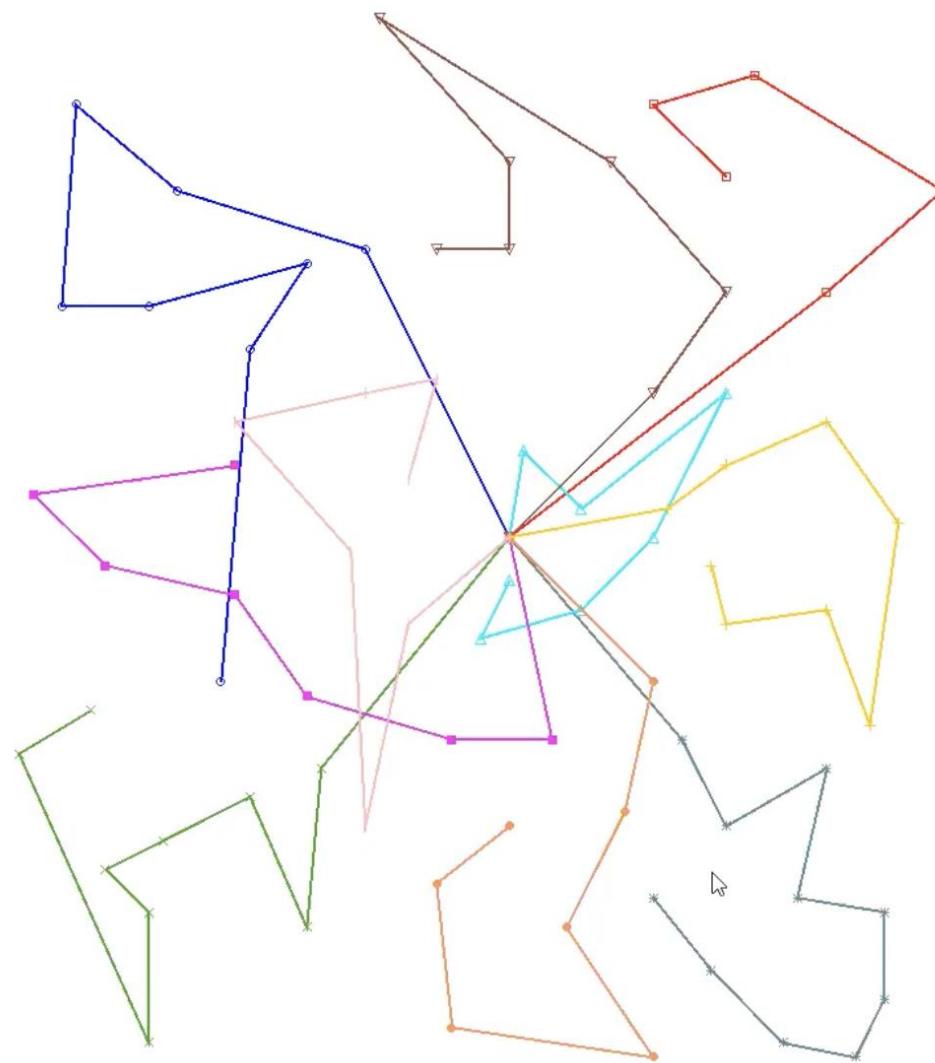
Regret with seeds



Regret with seeds

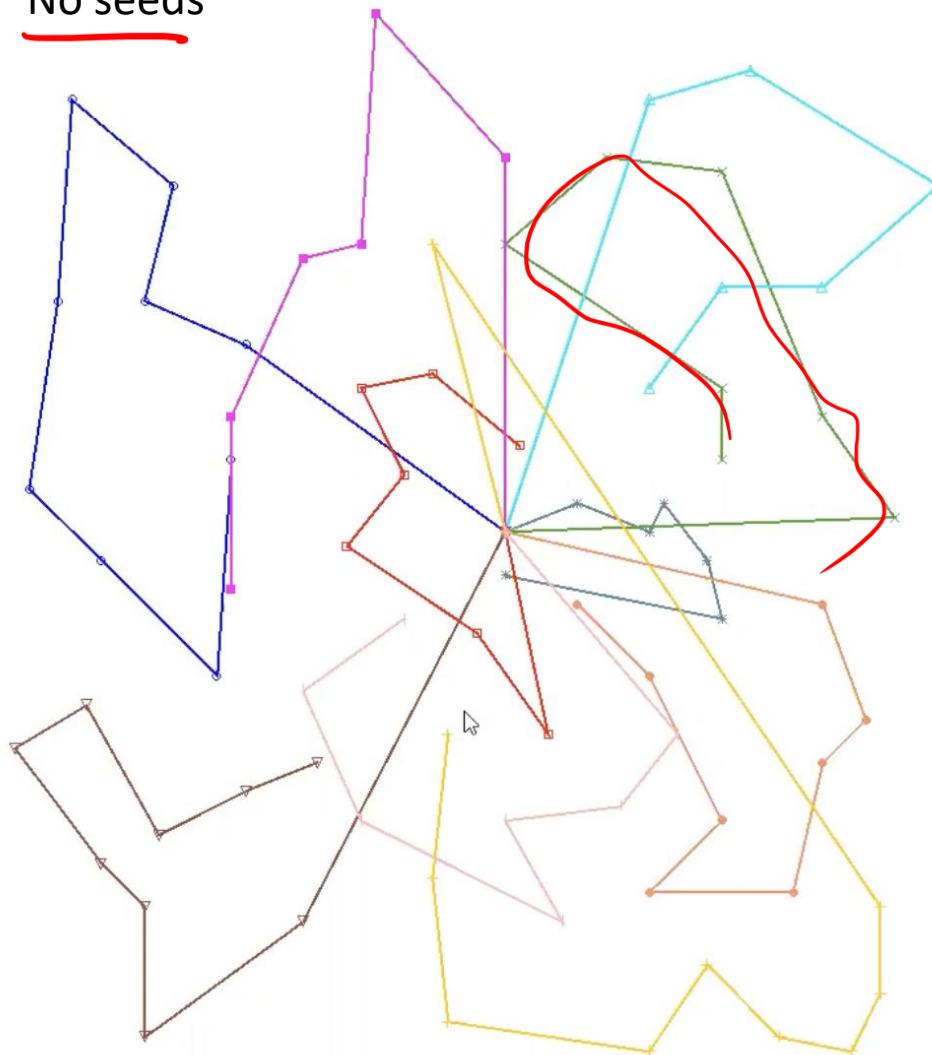


Regret with seeds

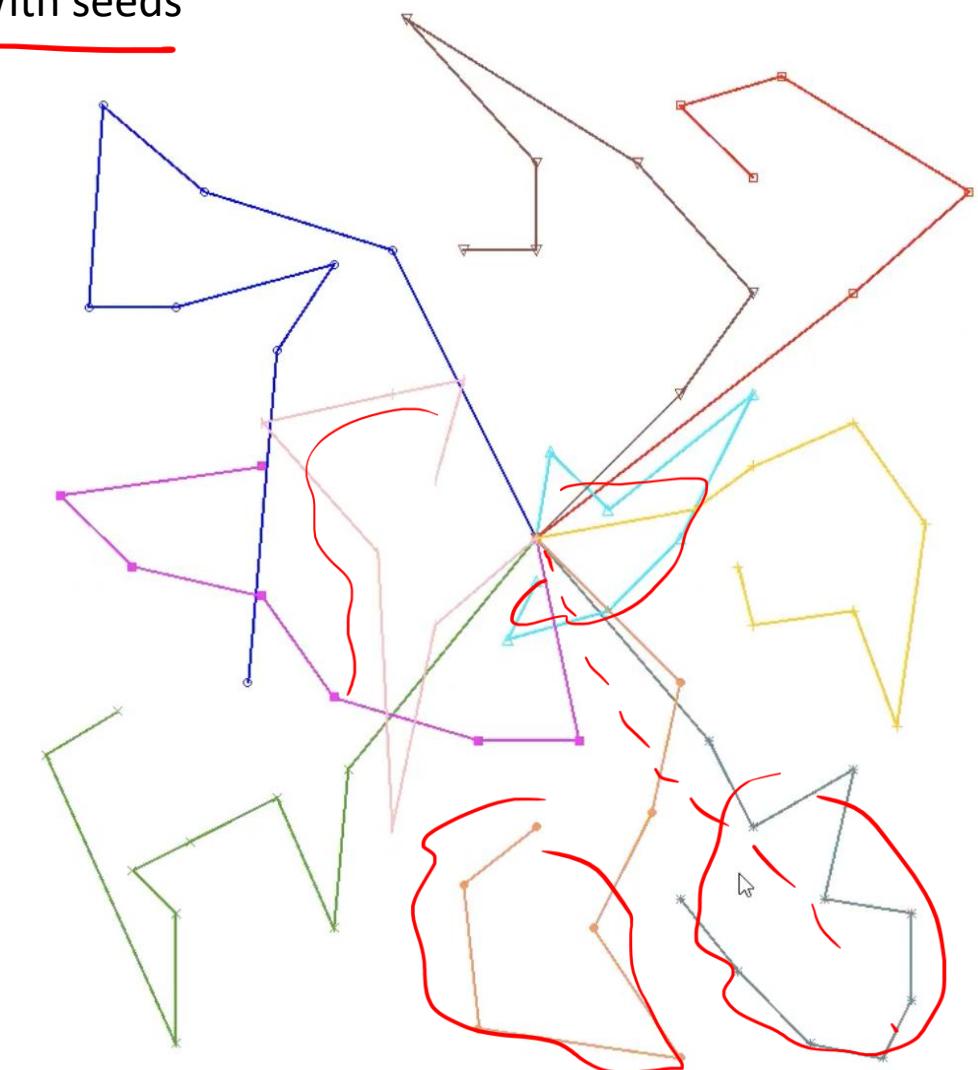


Regret with seeds

No seeds



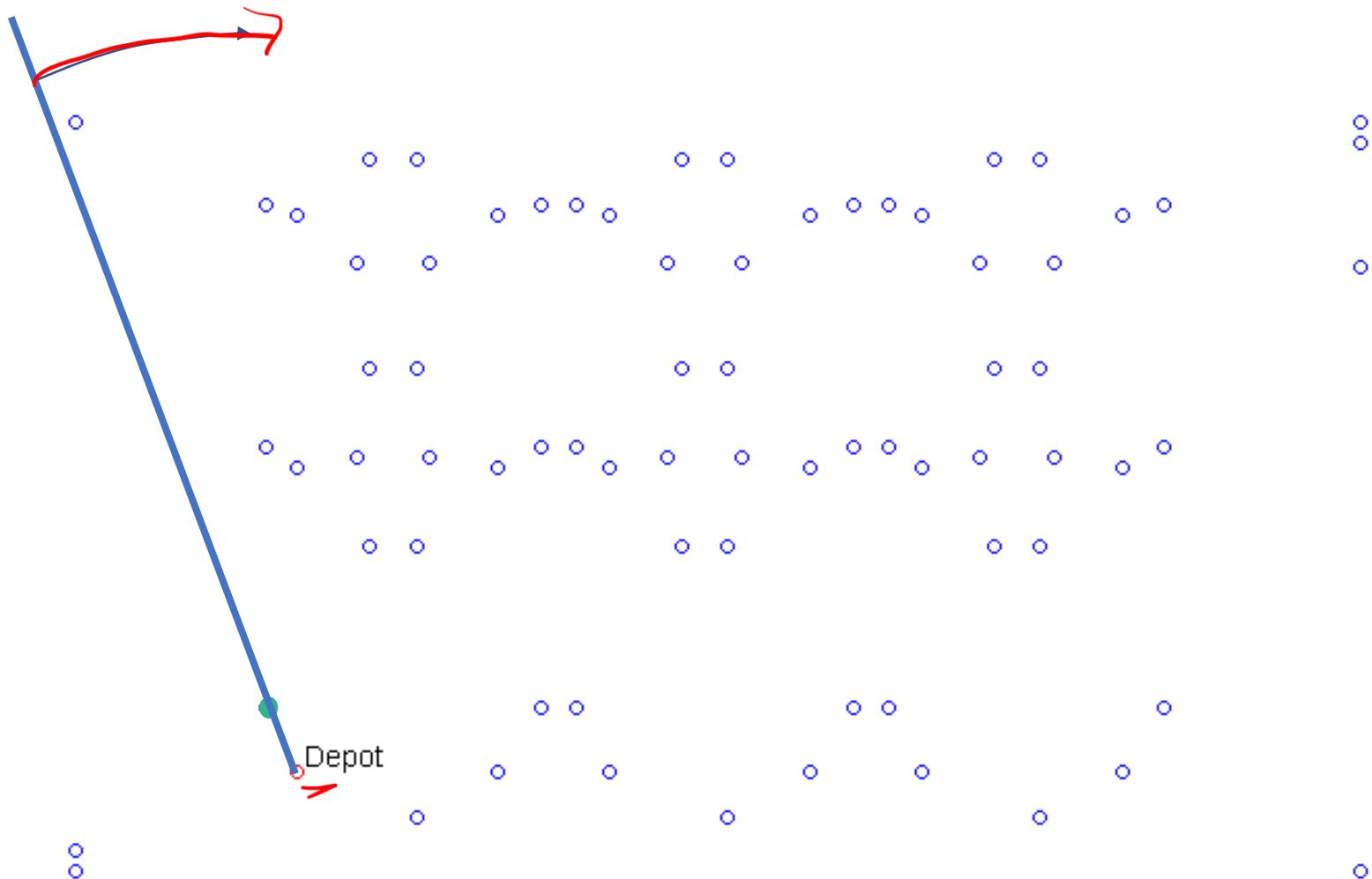
With seeds



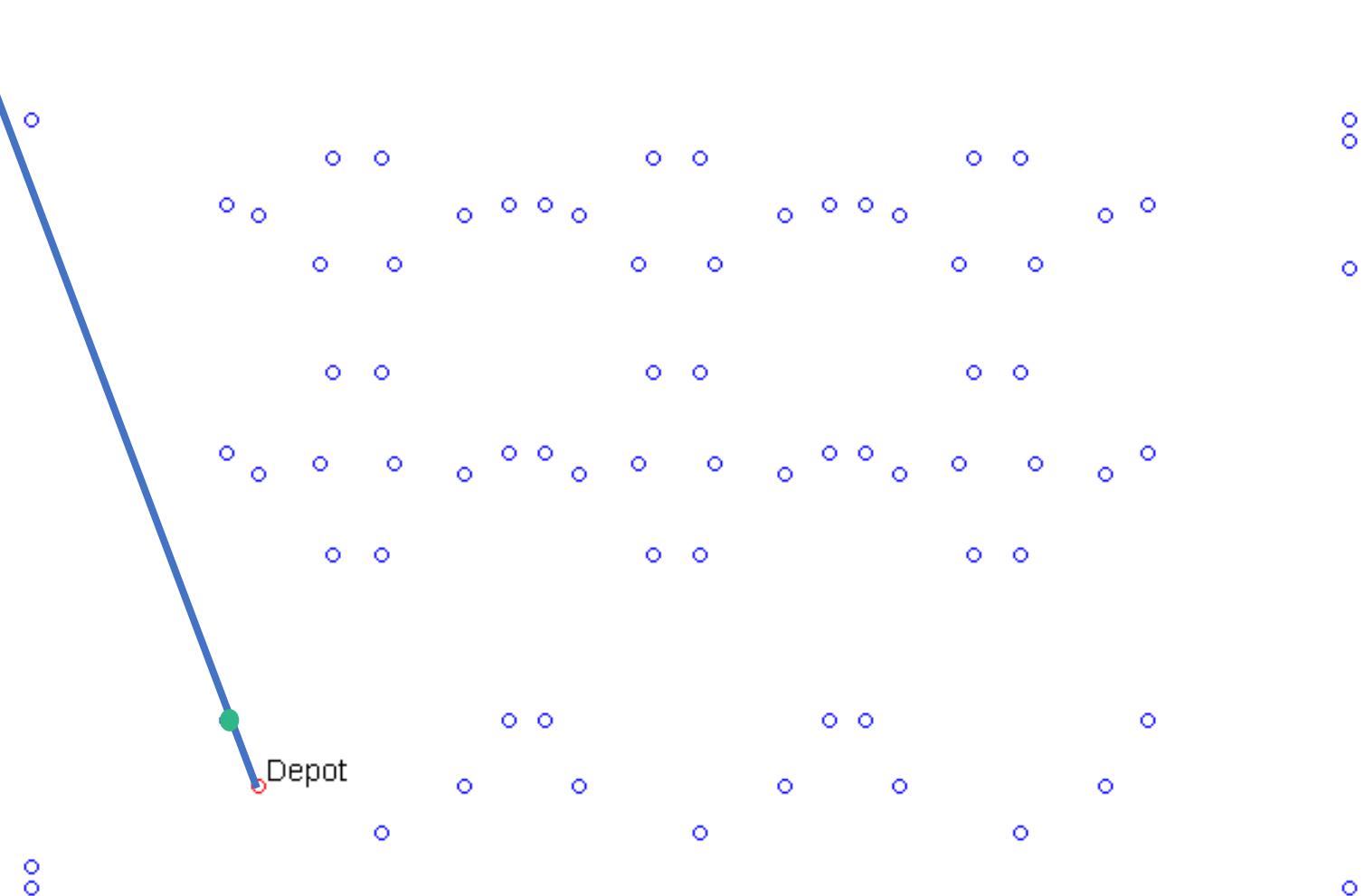
Other construction methods (for the VRP)

Bespoke methods

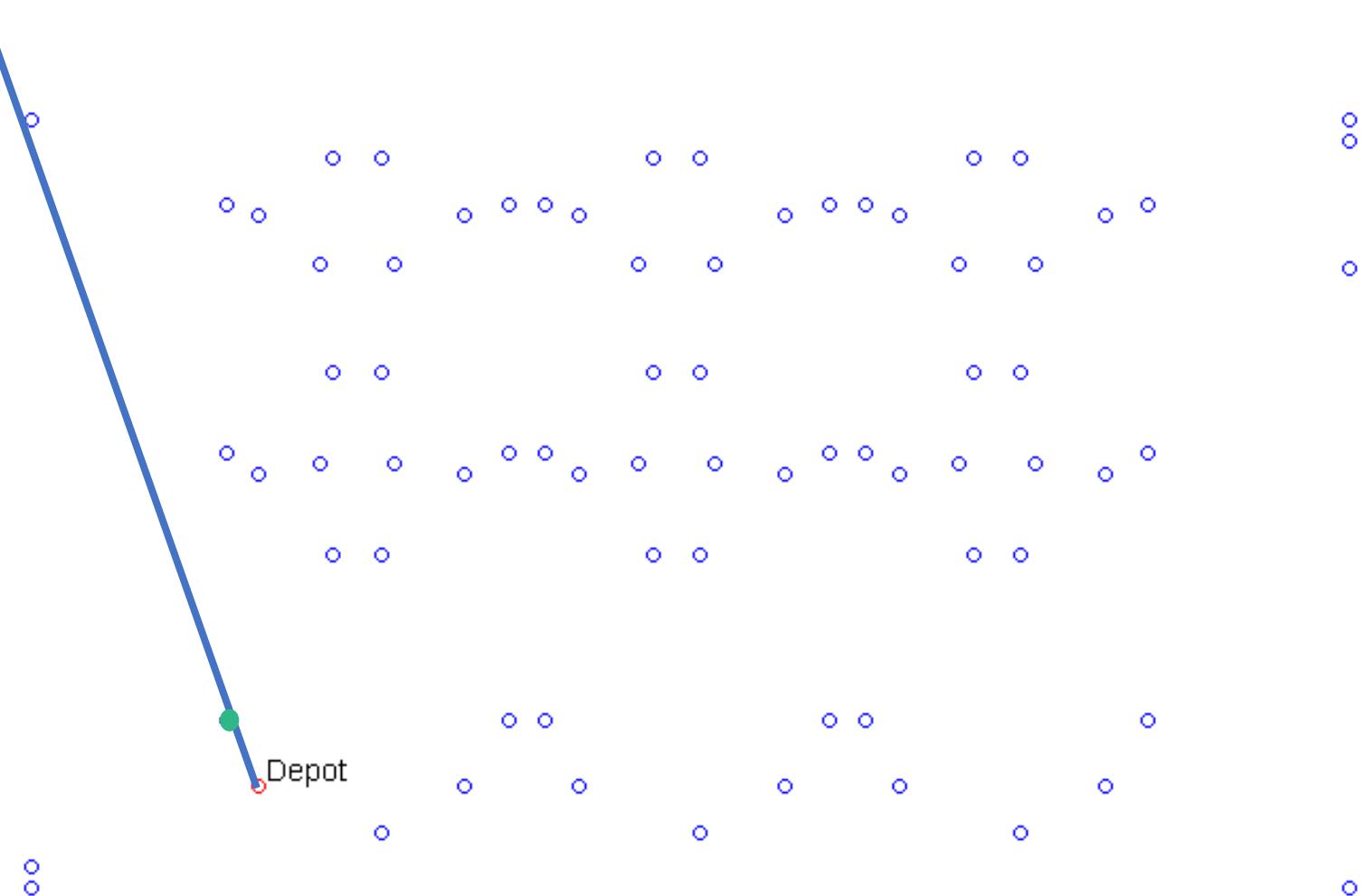
- Eg “Sweep”



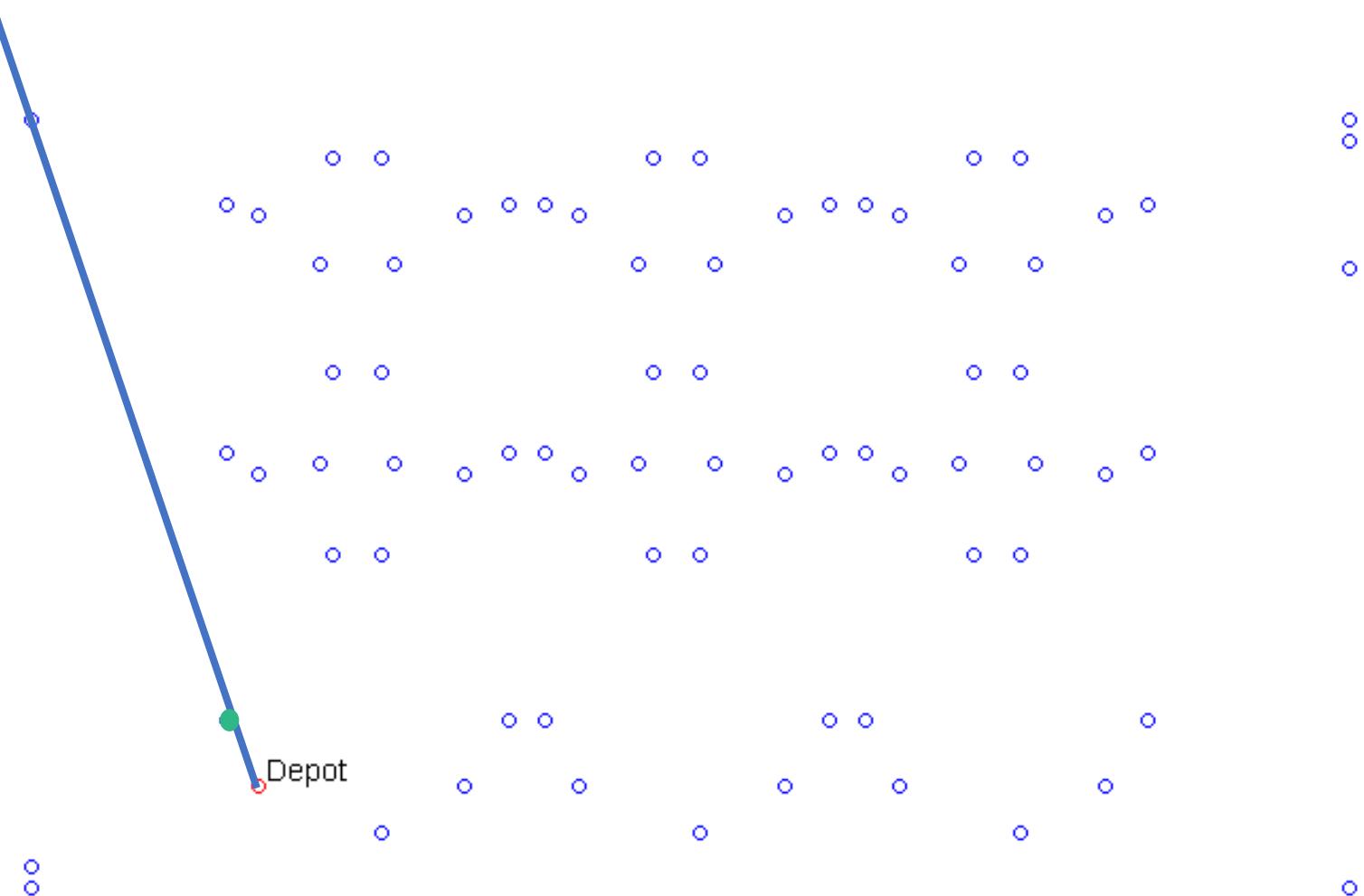
Other construction methods (for the VRP)



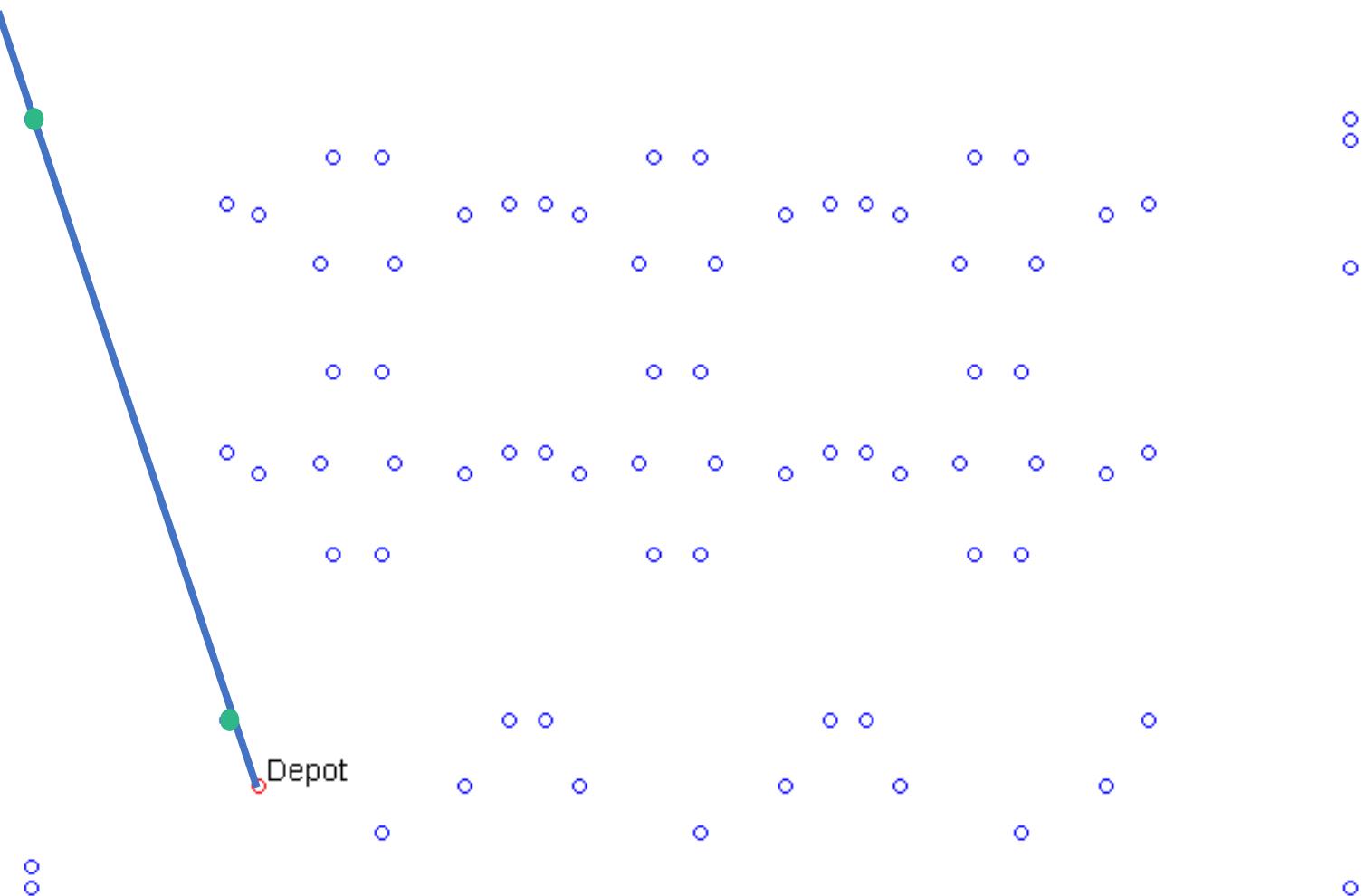
Other construction methods (for the VRP)



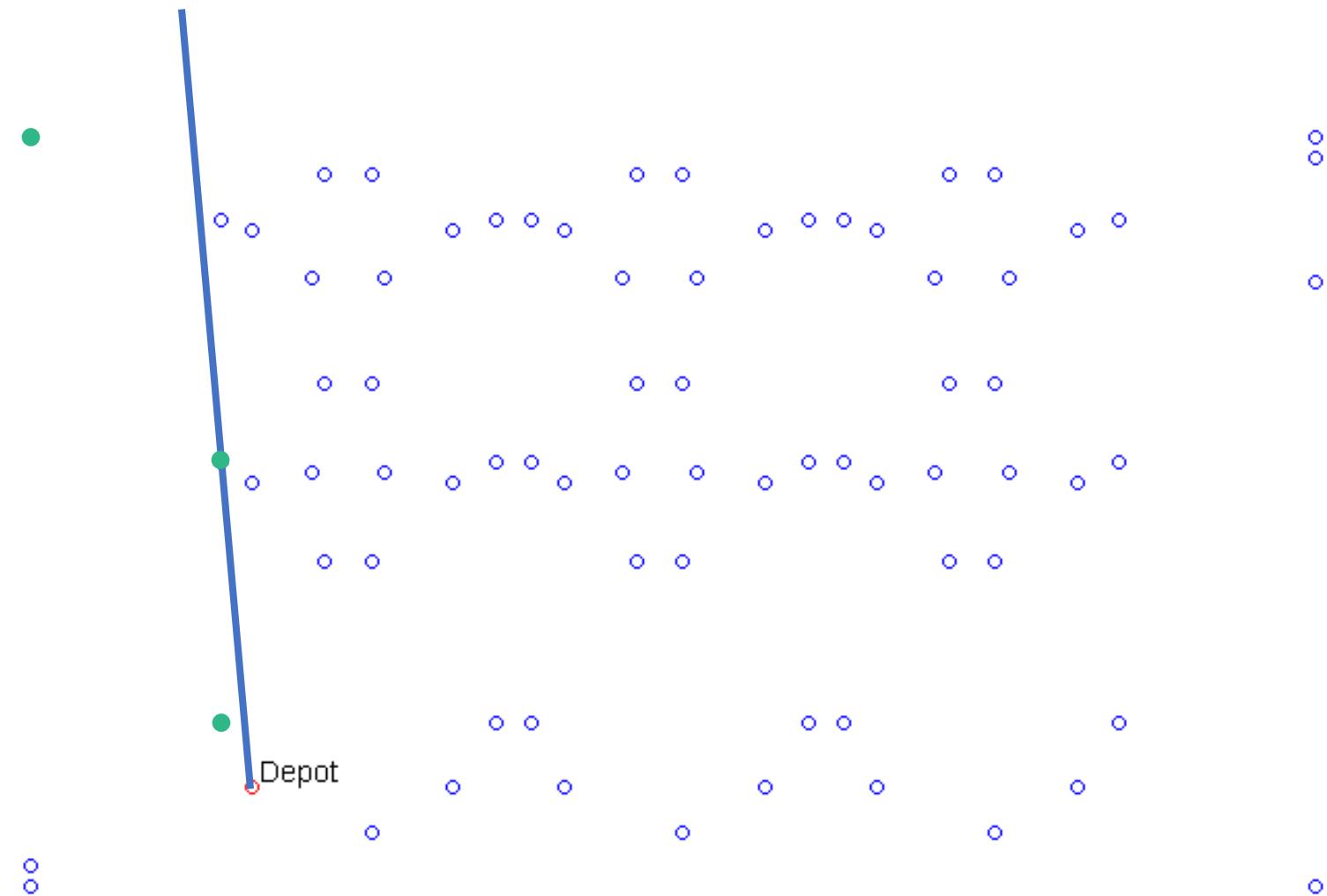
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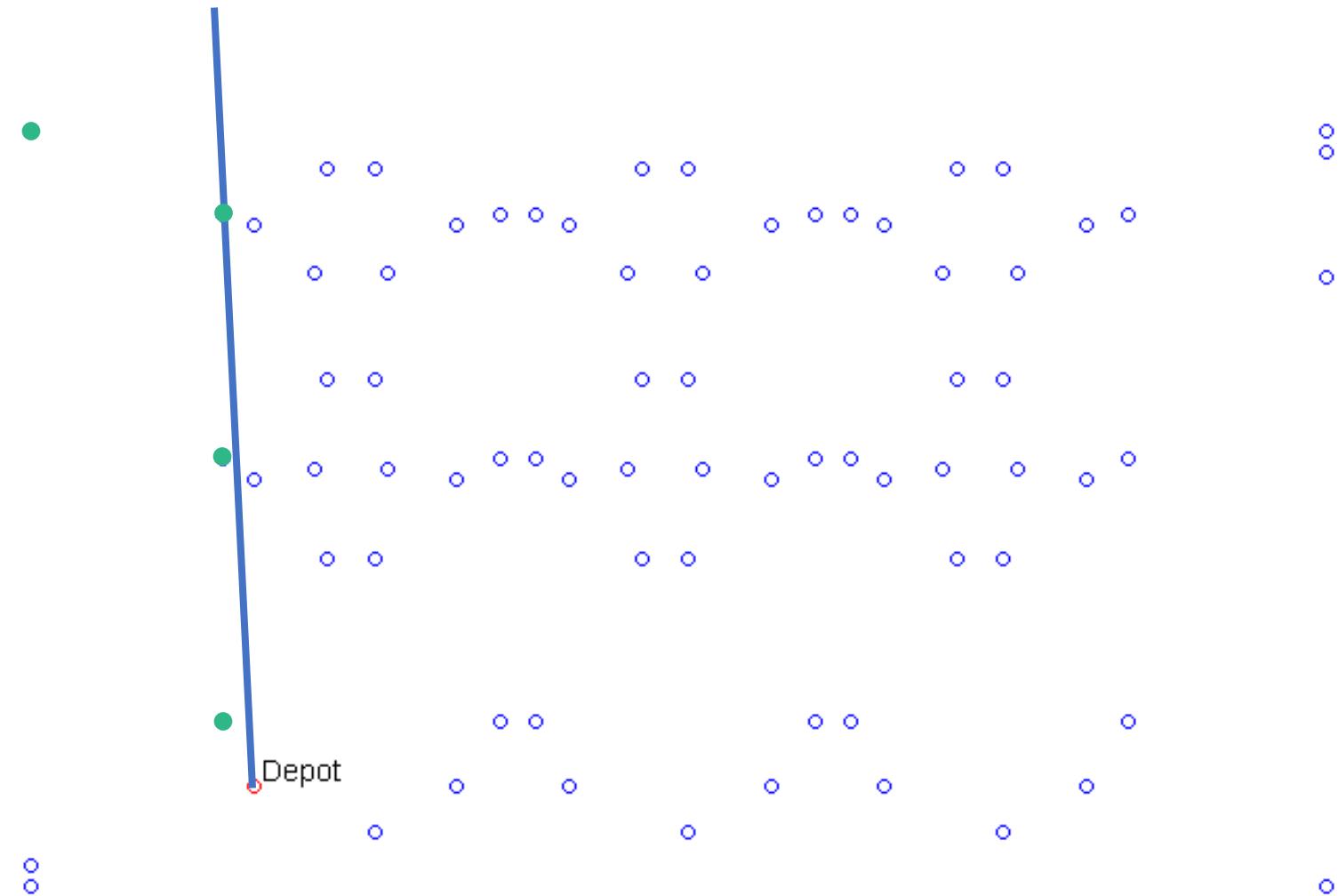
Other construction methods (for the VRP)



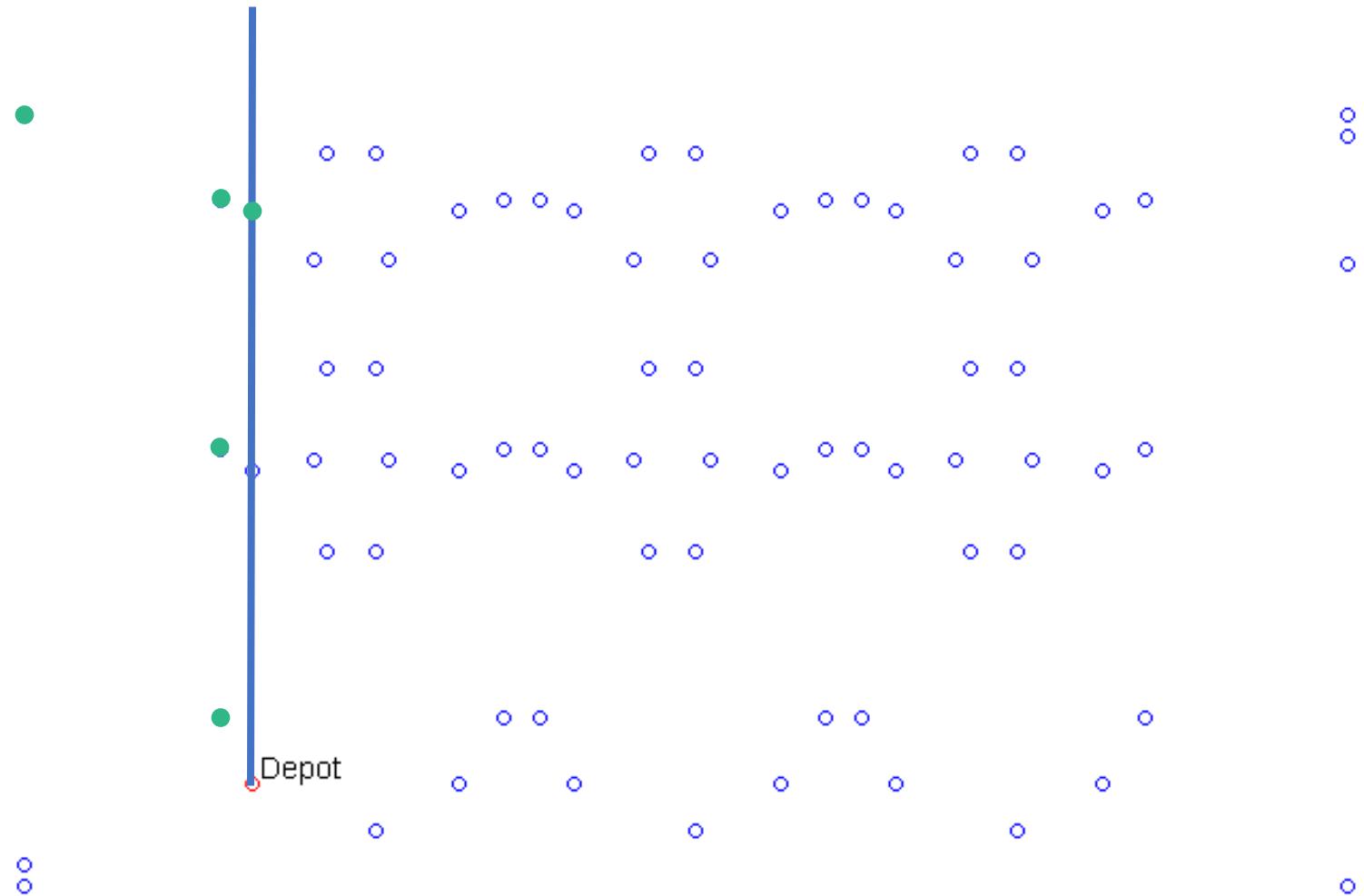
Other construction methods (for the VRP)



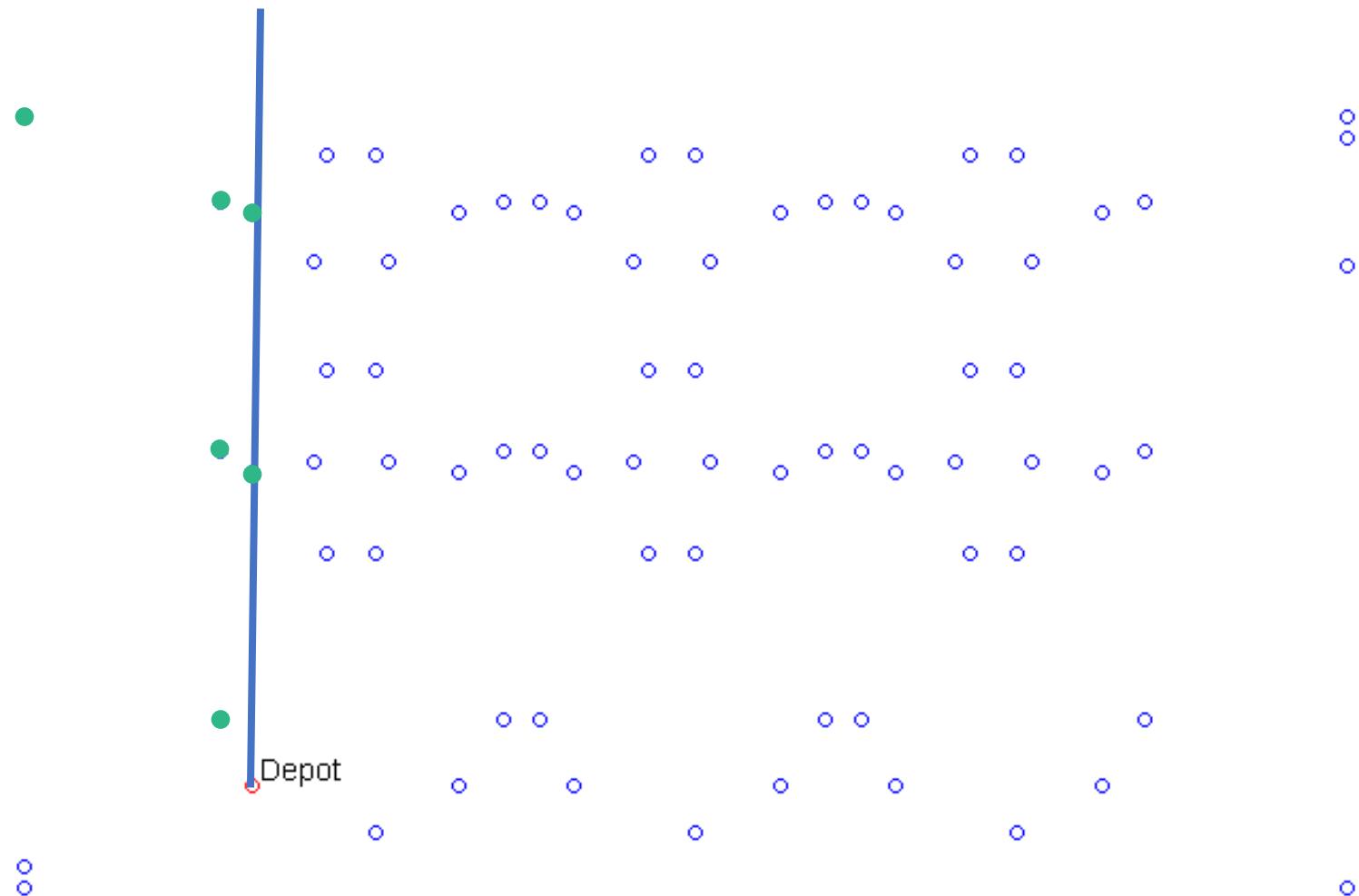
Other construction methods (for the VRP)



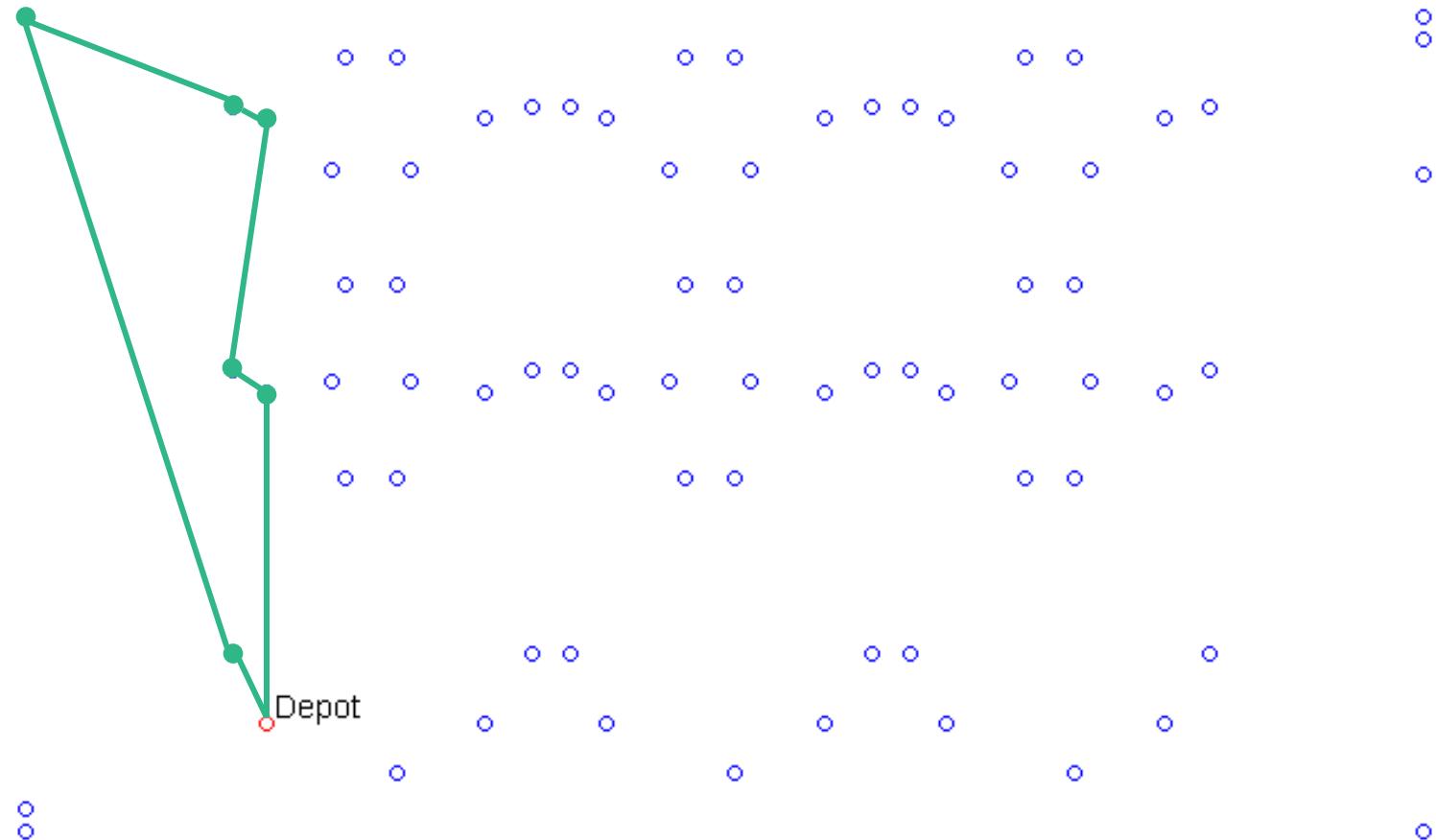
Other construction methods (for the VRP)



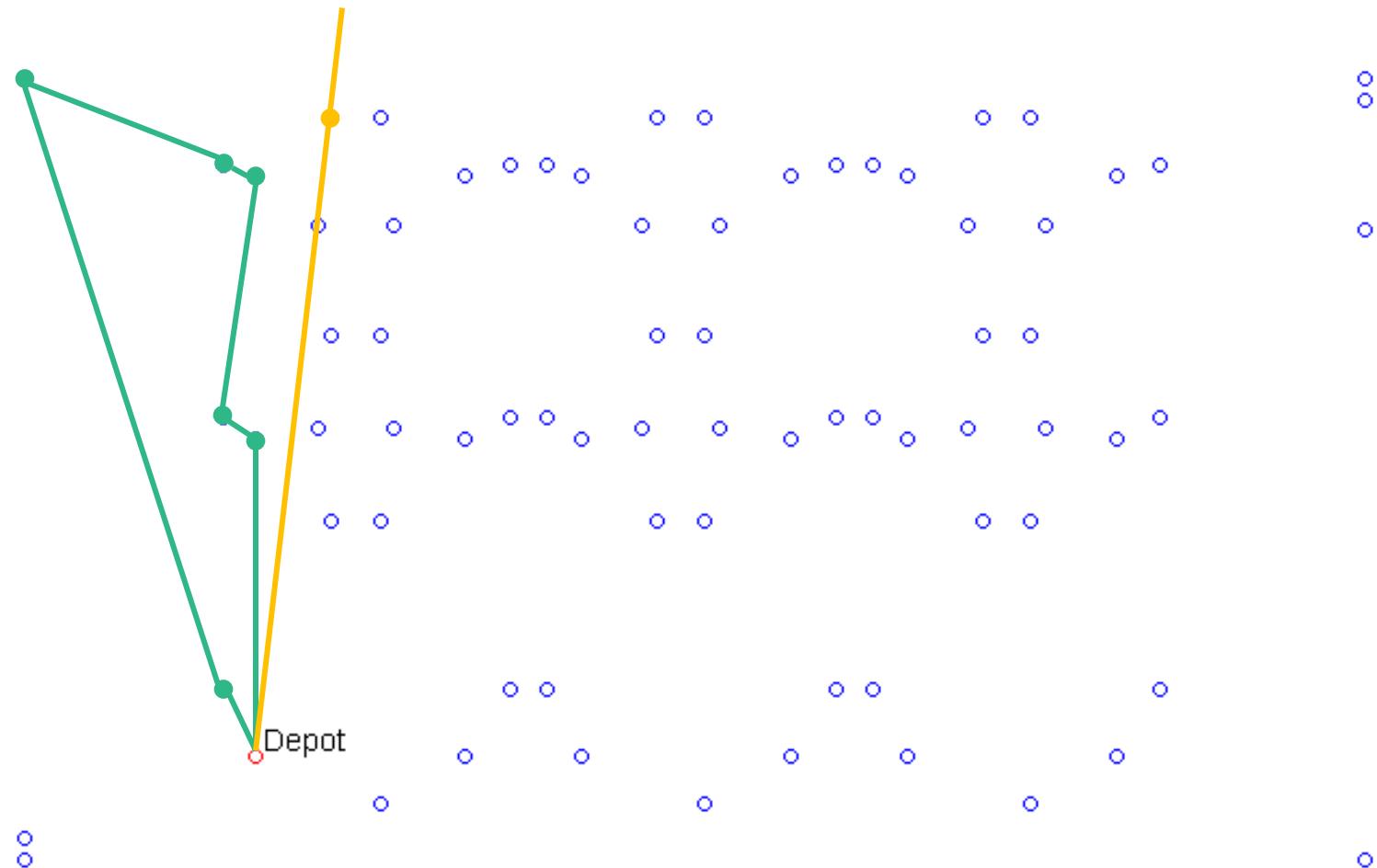
Other construction methods (for the VRP)



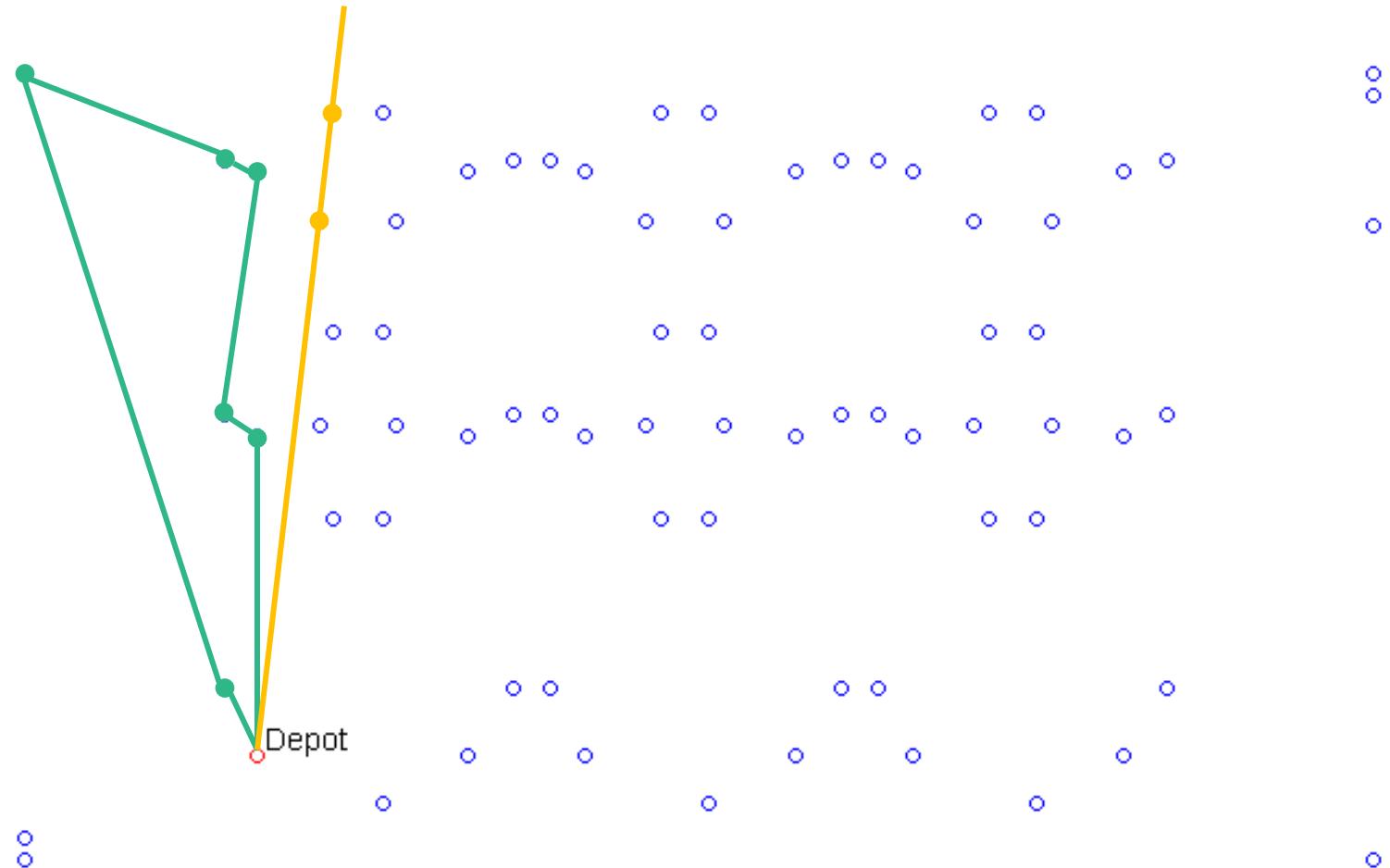
Other construction methods (for the VRP)



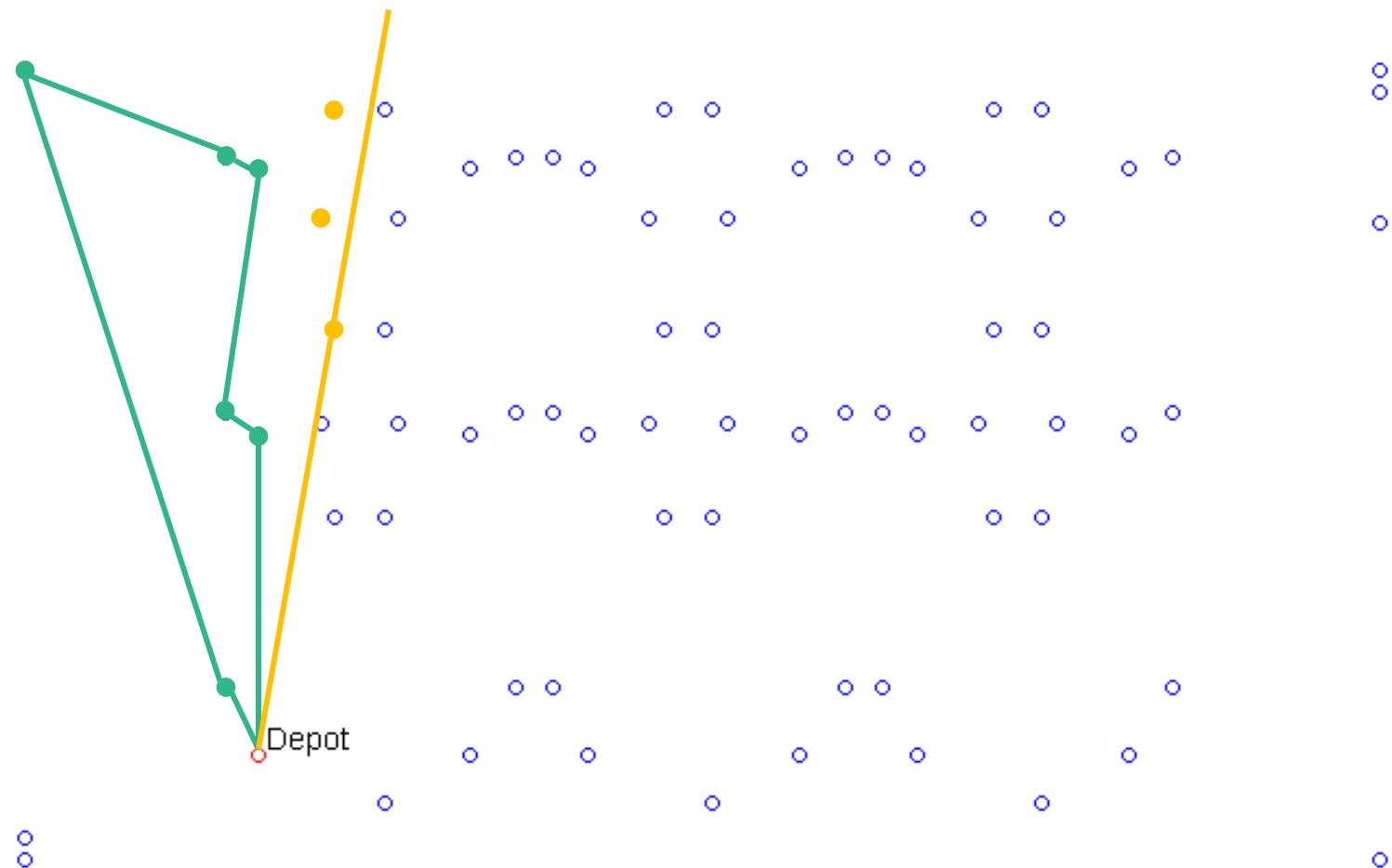
Other construction methods (for the VRP)



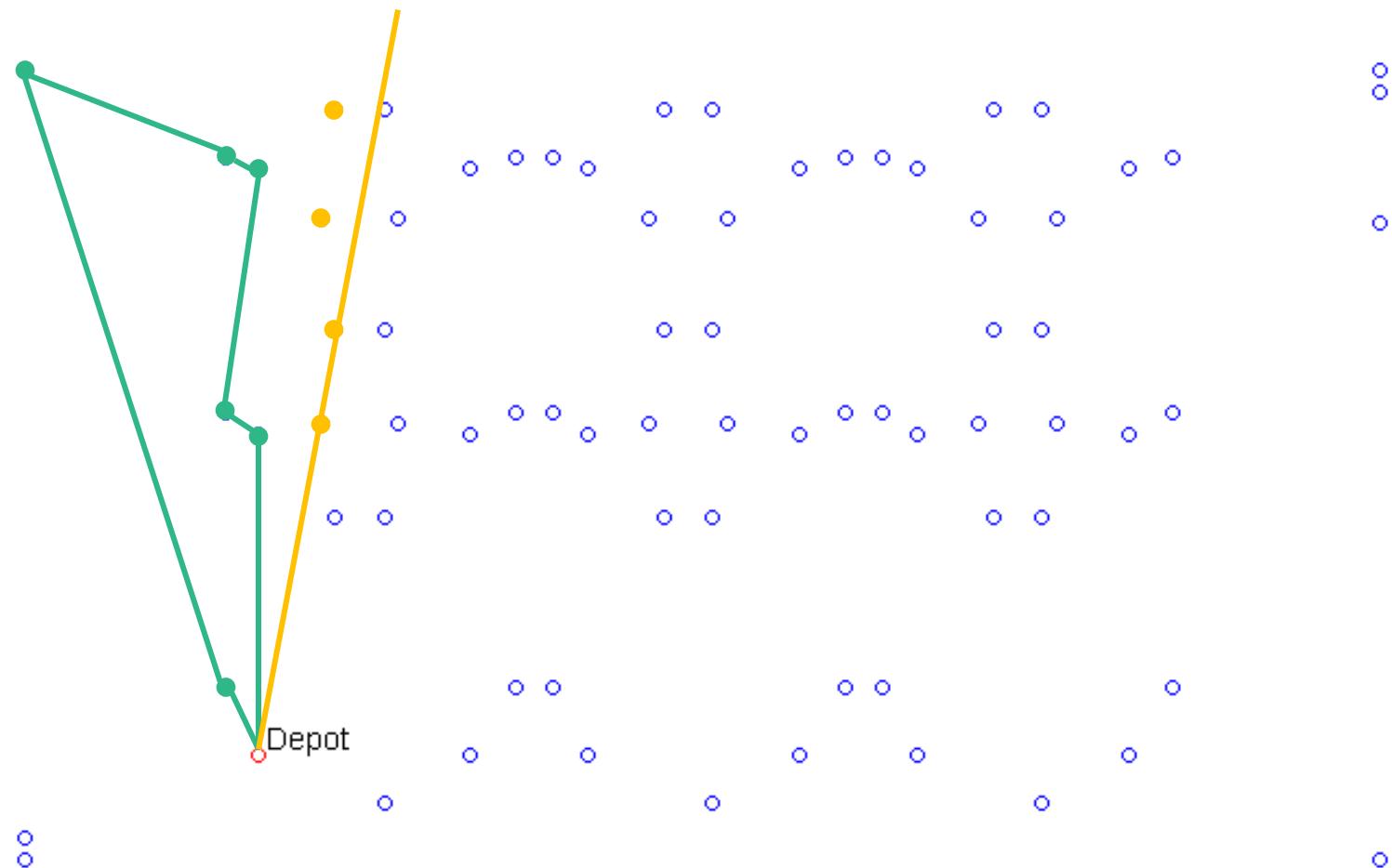
Other construction methods (for the VRP)



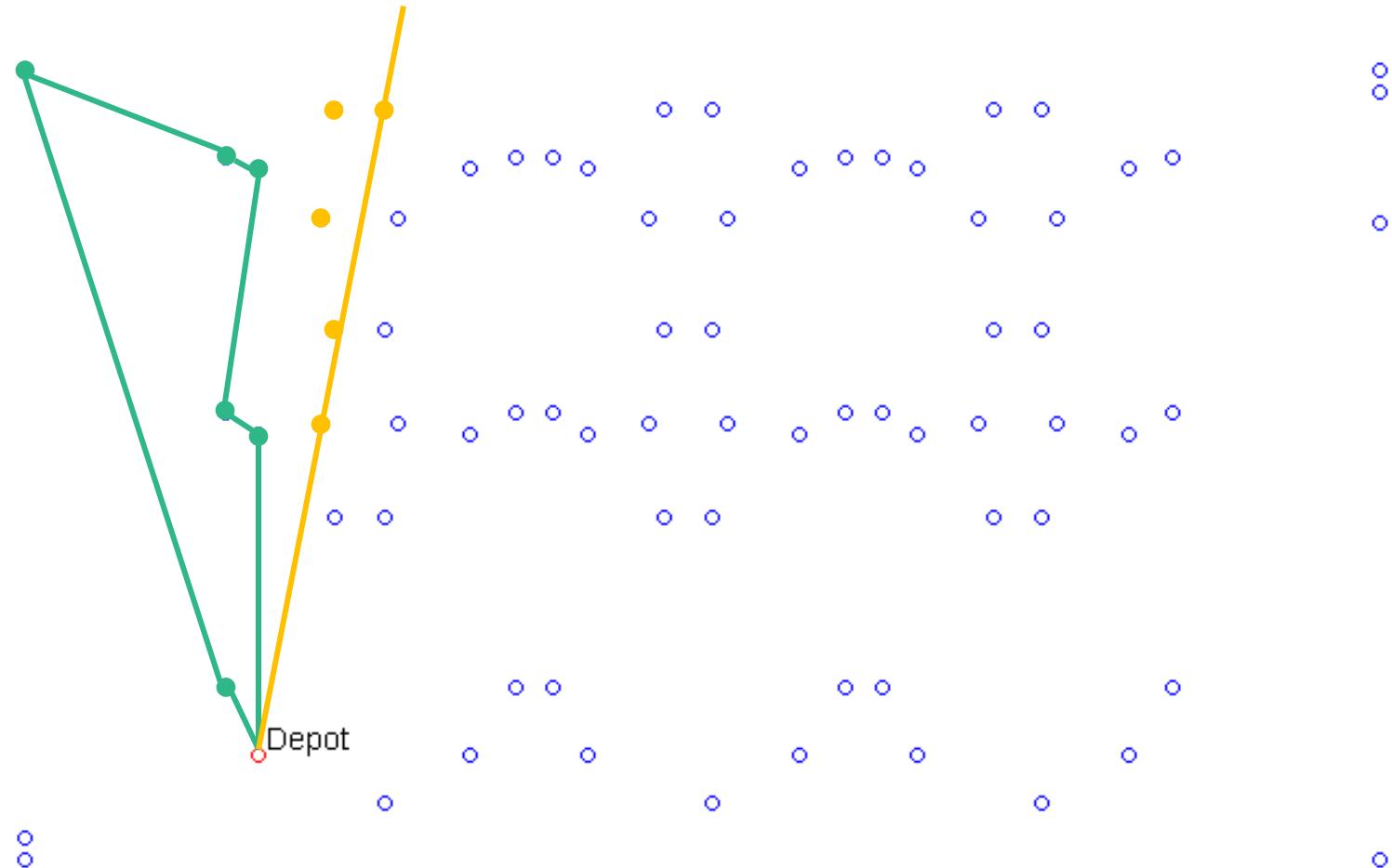
Other construction methods (for the VRP)



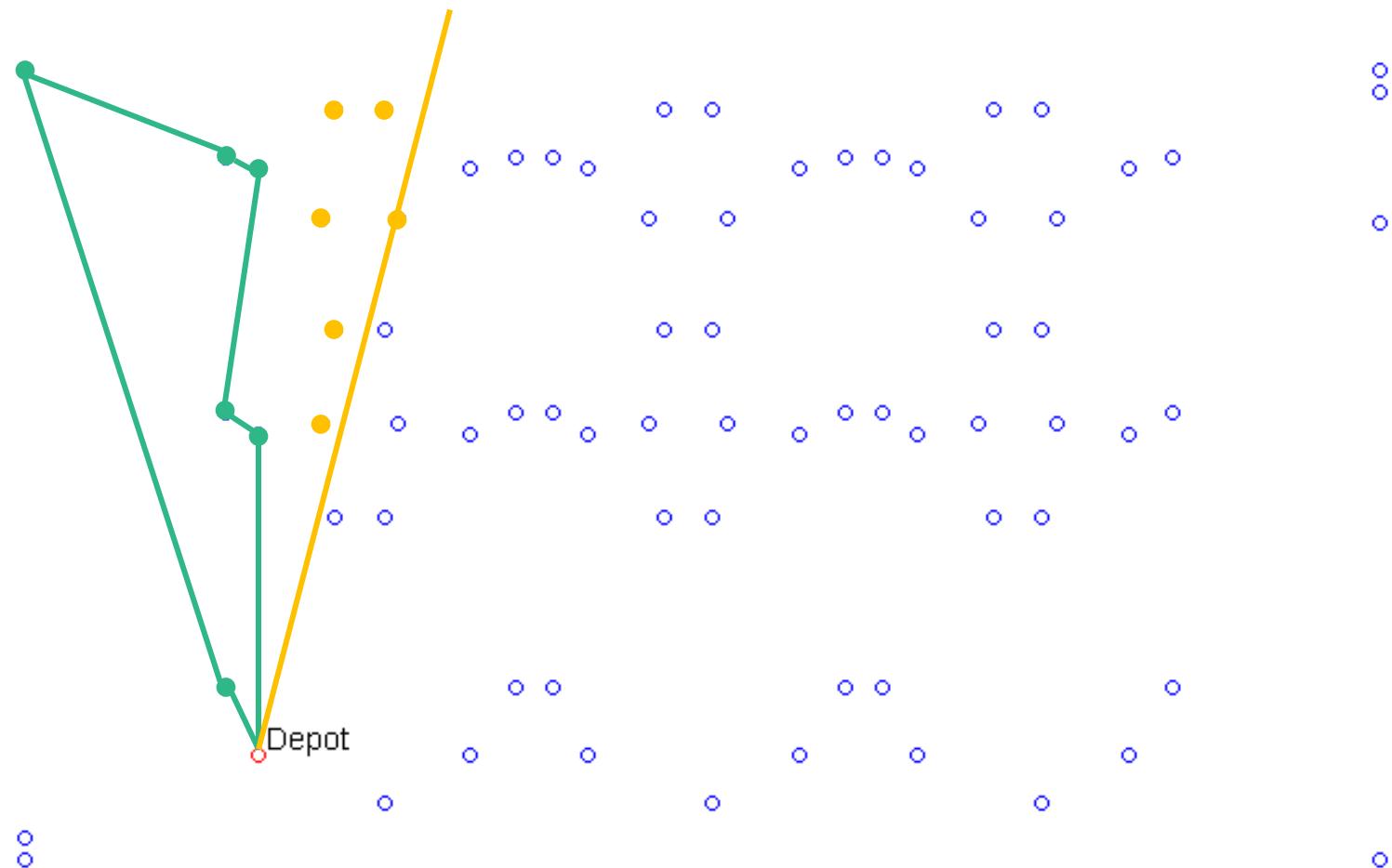
Other construction methods (for the VRP)



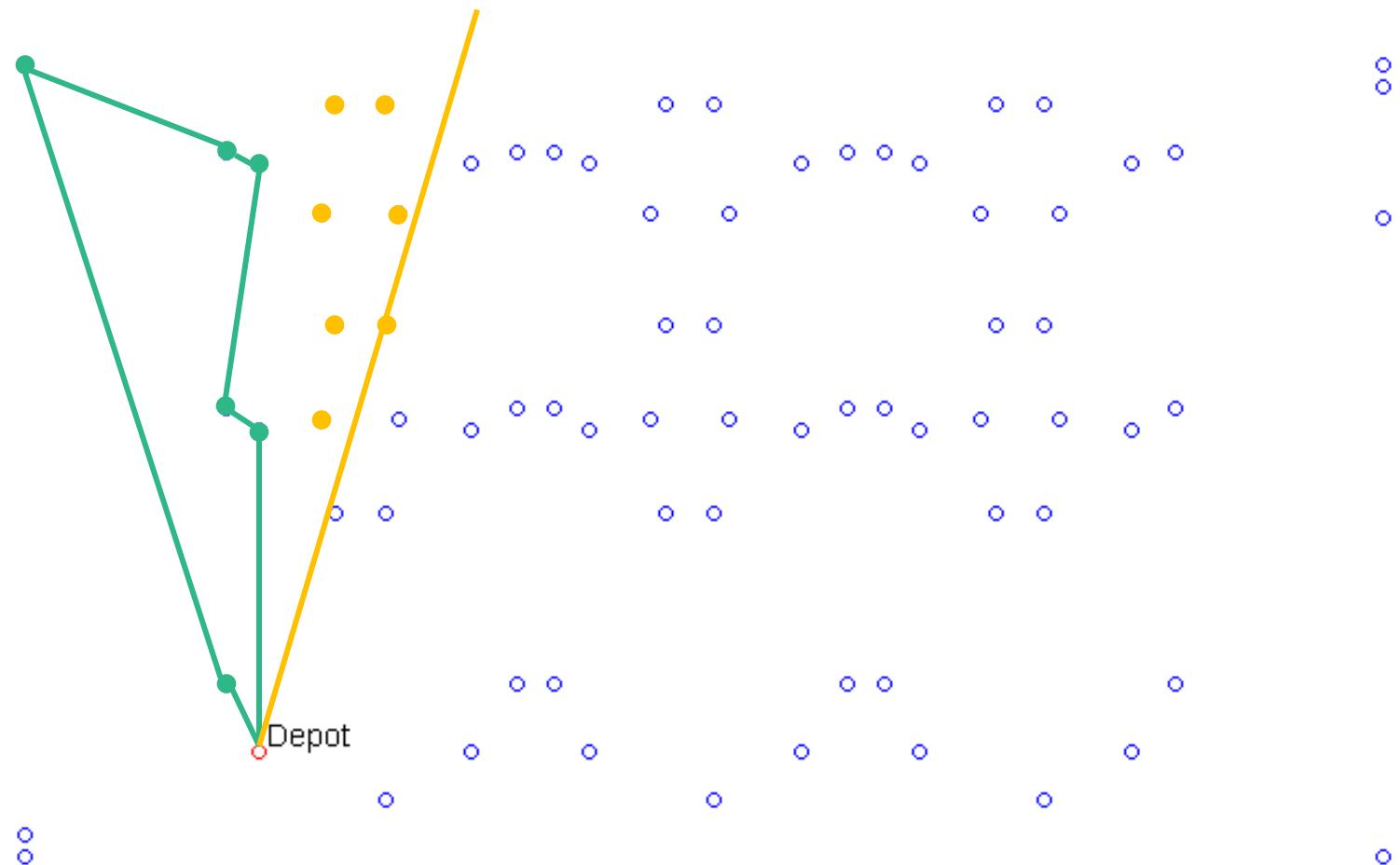
Other construction methods (for the VRP)



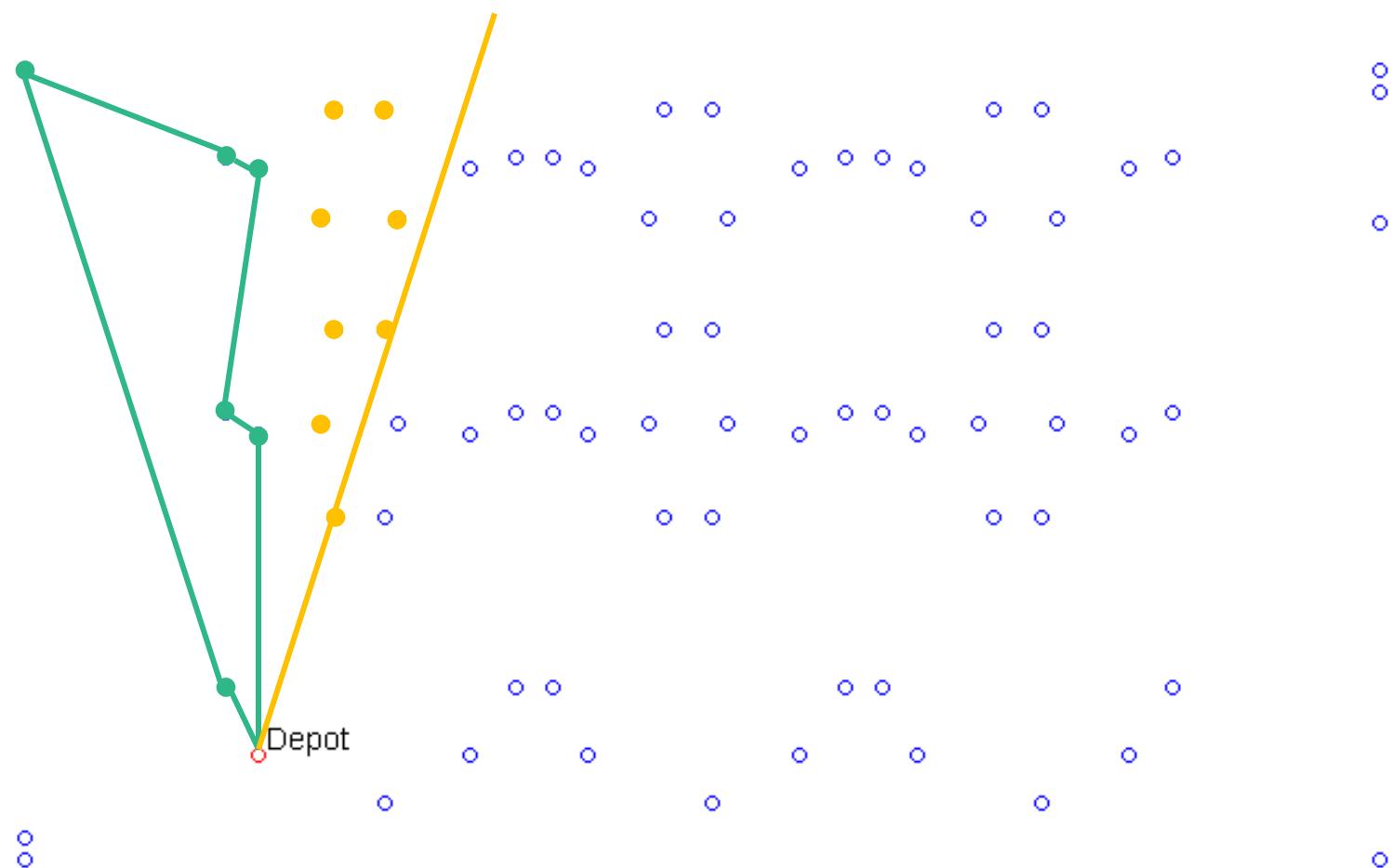
Other construction methods (for the VRP)



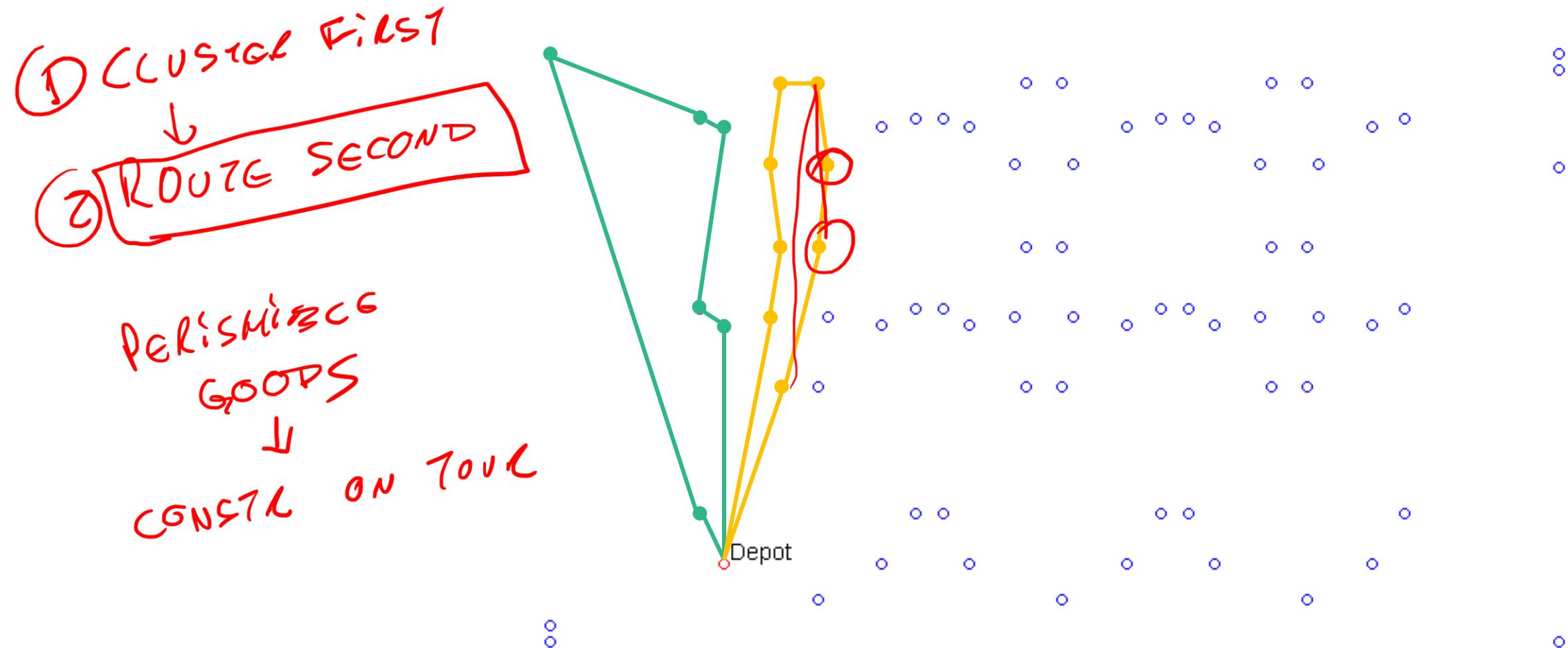
Other construction methods (for the VRP)



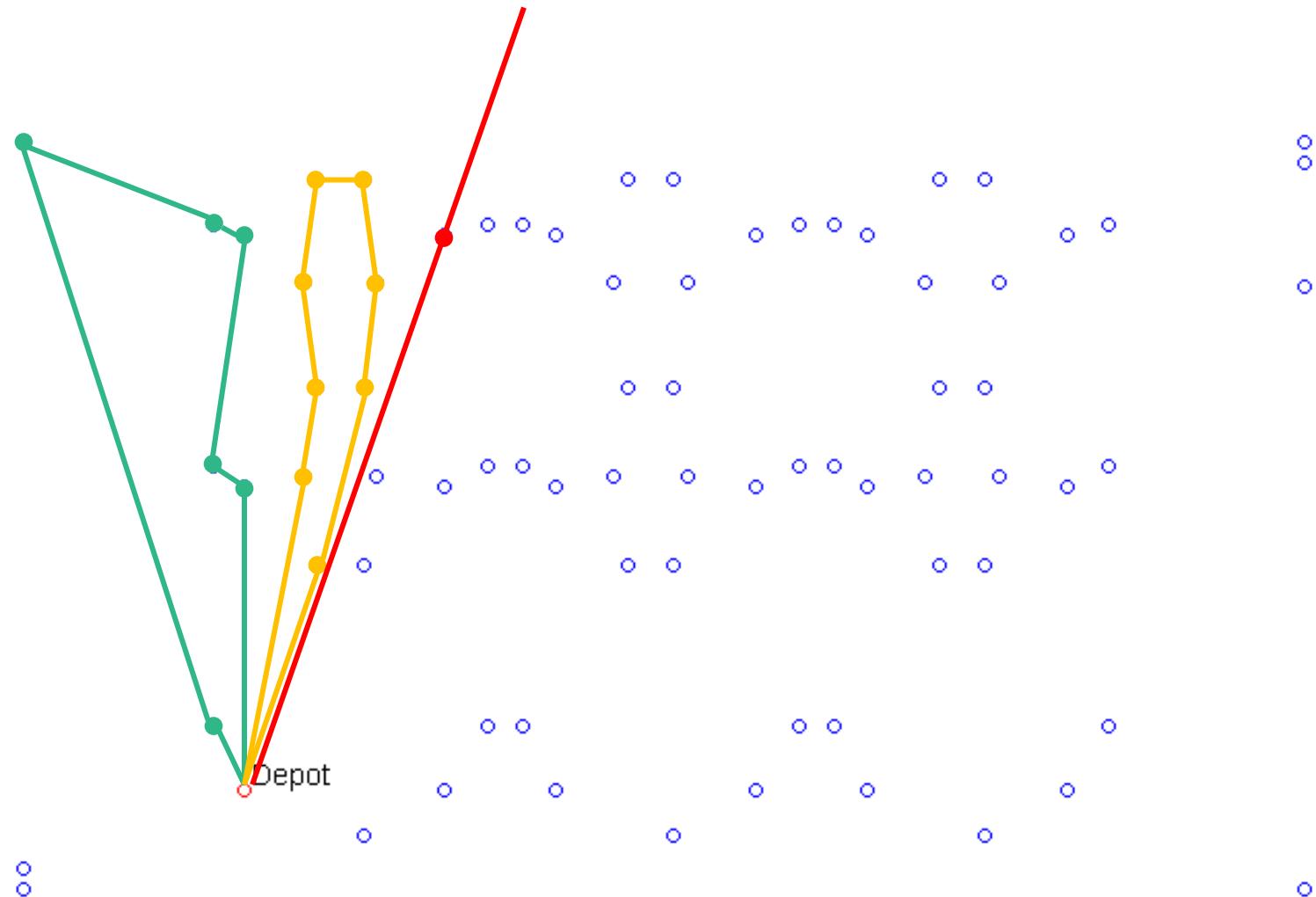
Other construction methods (for the VRP)



Other construction methods (for the VRP)



Other construction methods (for the VRP)



Other construction methods

3D Packing – Knapsack in 3D

- Working with a large mining co.
- They truck stuff from distributions centres near Perth to their mines
- Lots of stuff

Perth Woods
problem solved
Prisc Kiczy



3D Packing – Knapsack in 3D

- Working with a large mining co.
- They truck stuff from distributions centres near Perth to their mines
- Lots of stuff
- Lots and lots of stuff



The problem

- The company use 3PLs (Third-Party Logistics Providers) to move their stuff
- 3PLs are paid by the truck
- 3PLs say they pack their trucks well, but there is no incentive to do so



The problem

- The company *know* they are being ripped off
- They just don't know how much
- This is a “size of opportunity” project
- Aim: Answer the question

*If 3PLs packed their trucks really, really well,
how much could we save?*

- Answer: \$5M – Meh
- Answer: \$20M – Let's have a closer look



Problem statement

*Given a collection of items and a fleet of trucks,
fit as many items as possible into the trucks.*

Data:

- Each item has
 - dimensions WxHxD
 - a weight ↗
 - a due date ↗
 - a value ↗
- Each vehicle has
 - • multiple compartments with known dimensions WxHxD
 - a maximum load (weight) ↗
 - a departure date ↗

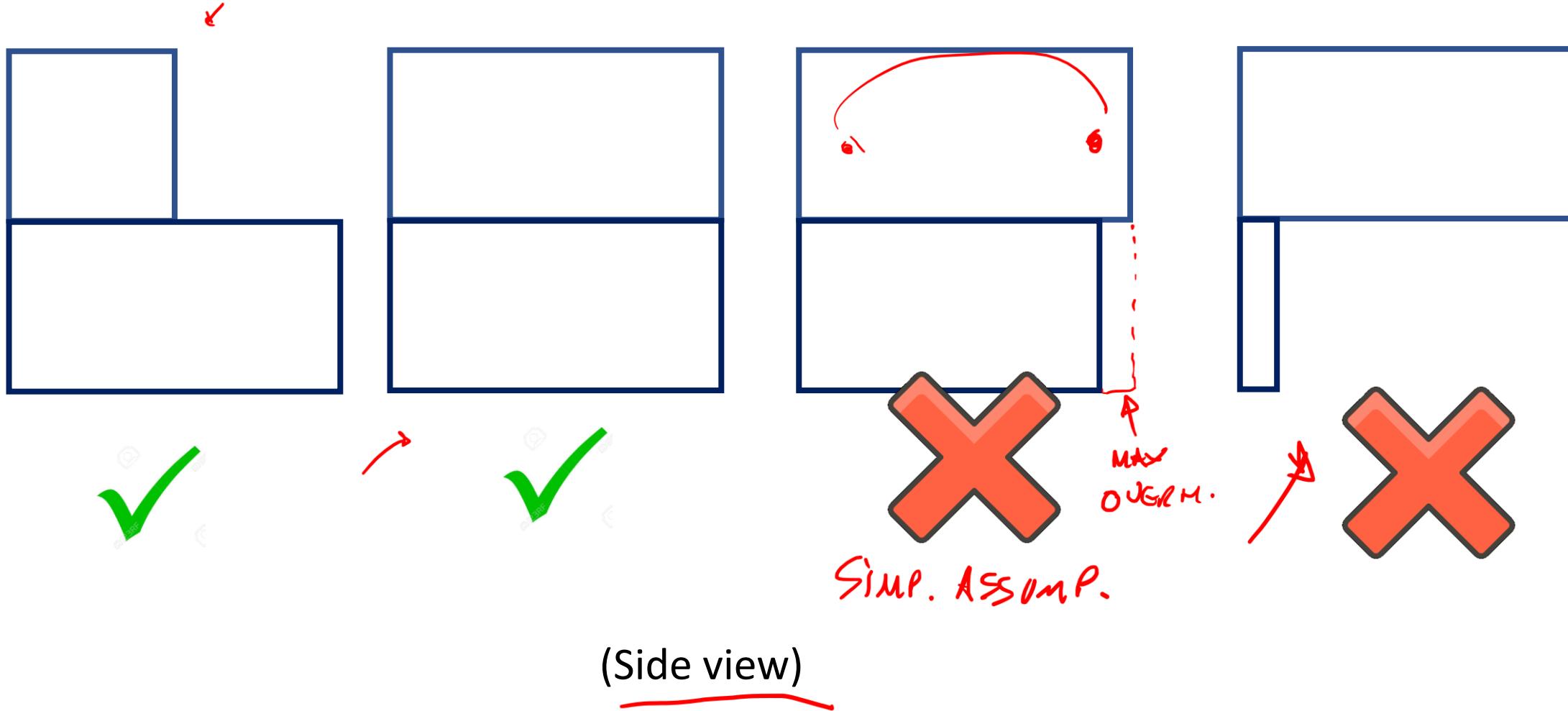


Problem statement

Constraints

- The vehicle's weight capacity constraint must be observed
 - Items cannot be placed on trucks departing after their due date
 - The items must be arranged so that they fit into the truck
 - They may be rotated the around the vertical axis, but can't be placed on their side
- Items may be placed on top of each other
 - Some may be marked “*top-load only*”
 - They must be supported from below with *no overhang*

No overhang



Problem statement

What we don't look at:

- Weight distribution / axle loading constraints
- Stacking constraints (e.g. item above cannot weigh more than item below)
- Tipping constraints (don't track centre of gravity)

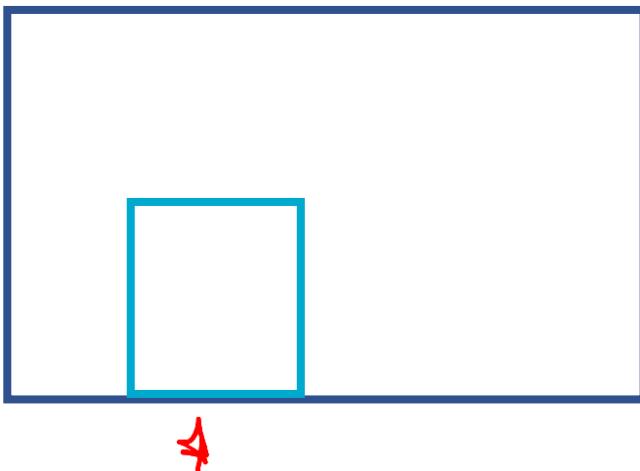
Solution

N S

- We will use a metaheuristic to solve the problem
- It will need two related subroutines:
 - Given a stack of items and a number of trucks, how can I distribute the items amongst the trucks
 - Given a truck and a stack of items, can I fit the items into the truck? ↗
- We know that the metaheuristic will call these subroutines many thousands of times.
 - So, packing must be FAST ↗

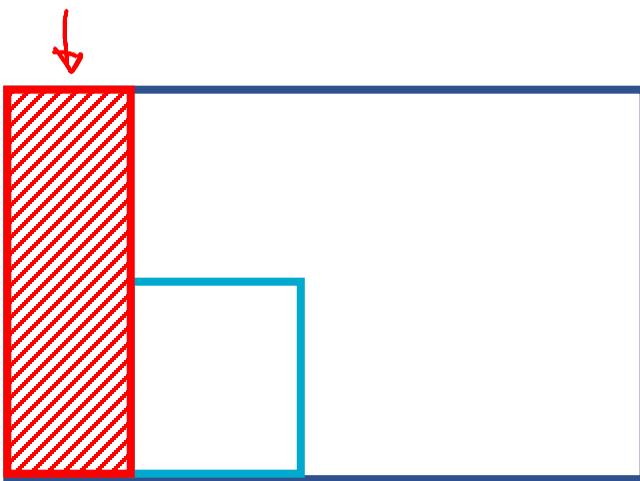
Solution

- The solution is based on a few observations.
- Observation 1: Might as well pack in corners



Solution

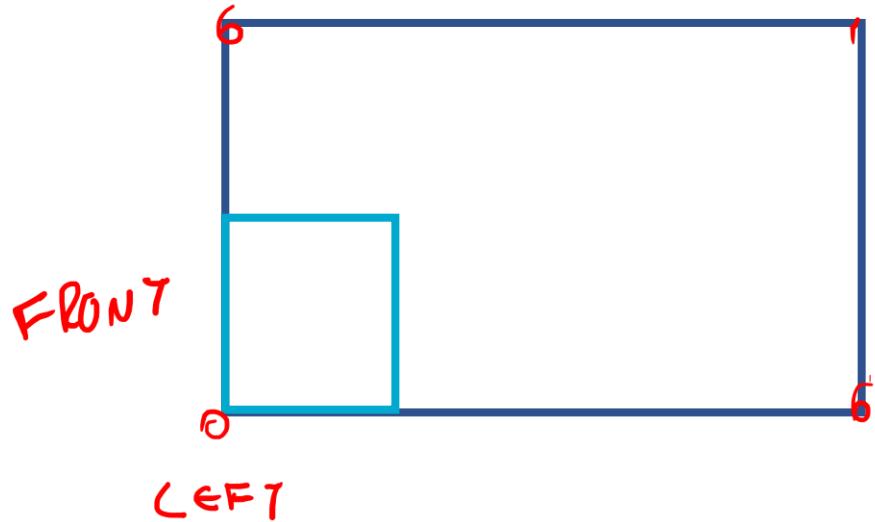
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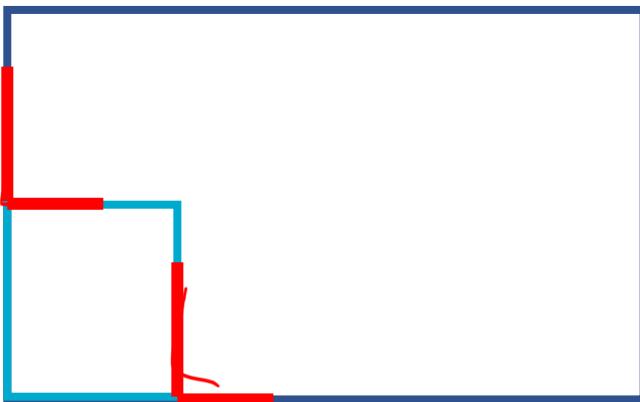
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TOP View



Solution

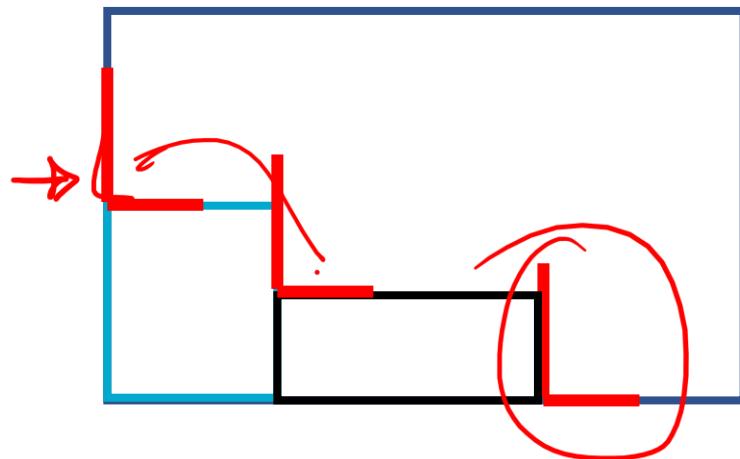
- The solution is based on a few observations.
- Observation 2: Corners are easy to track



Solution

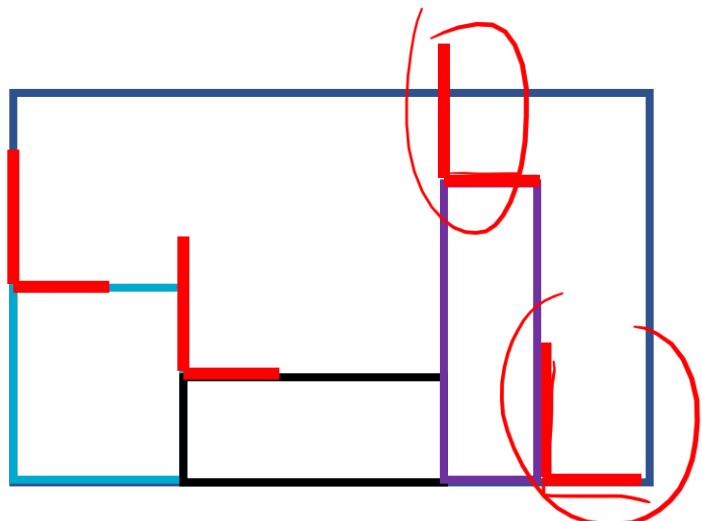
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Side view



Solution

- The solution is based on a few observations.
- Observation 2: Corners are easy to track



Solution

- The solution is based on a few observations.
- Observation 2': Corners are mostly easy to track



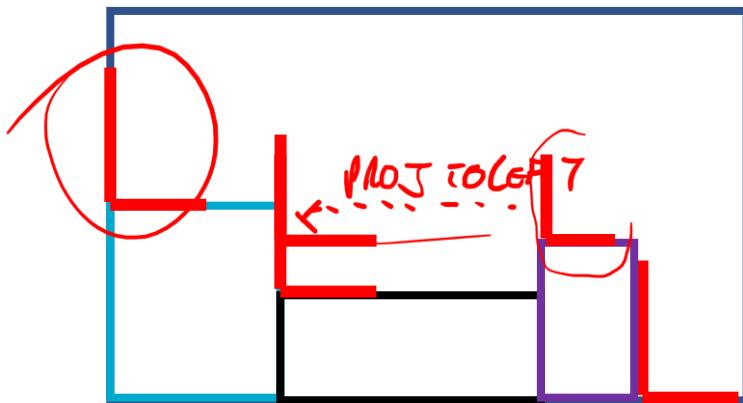
Recall that this is the view from the top:

- **Yellow is by the side of purple (not on top)**



Solution

- The solution is based on a few observations.
- Observation 2': Corners are mostly easy to track



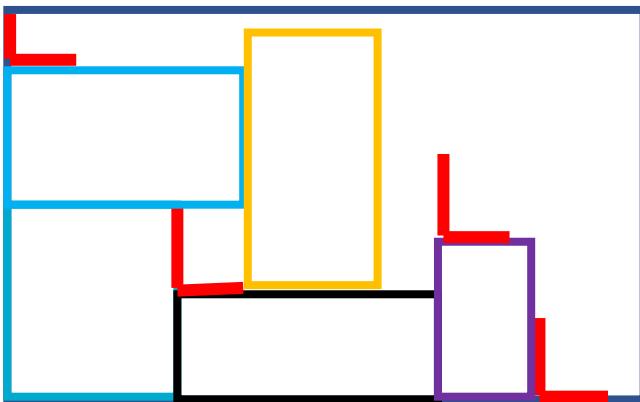
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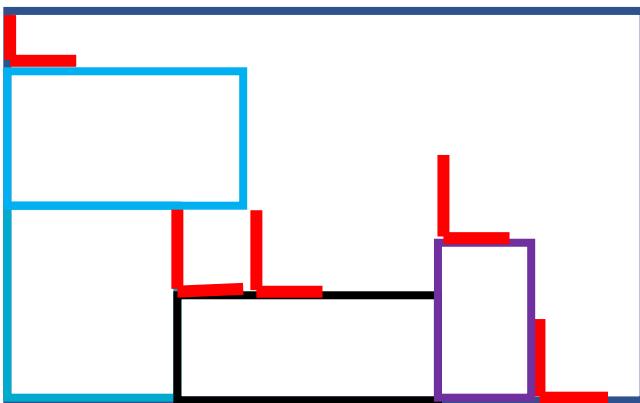
Solution

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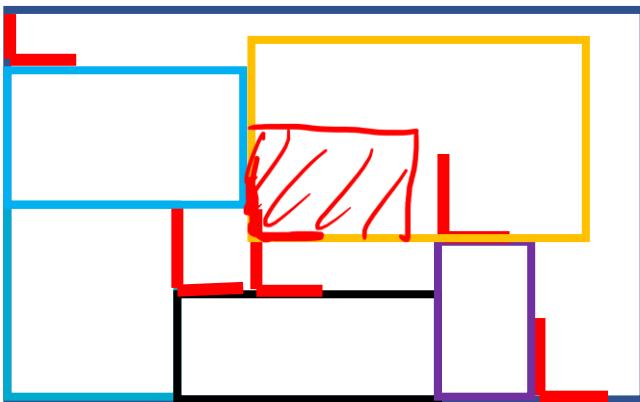
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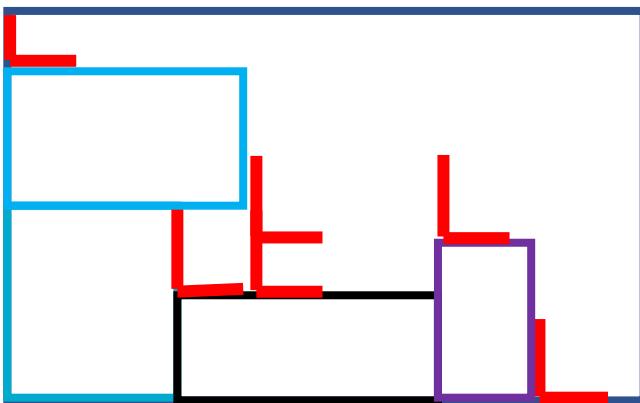
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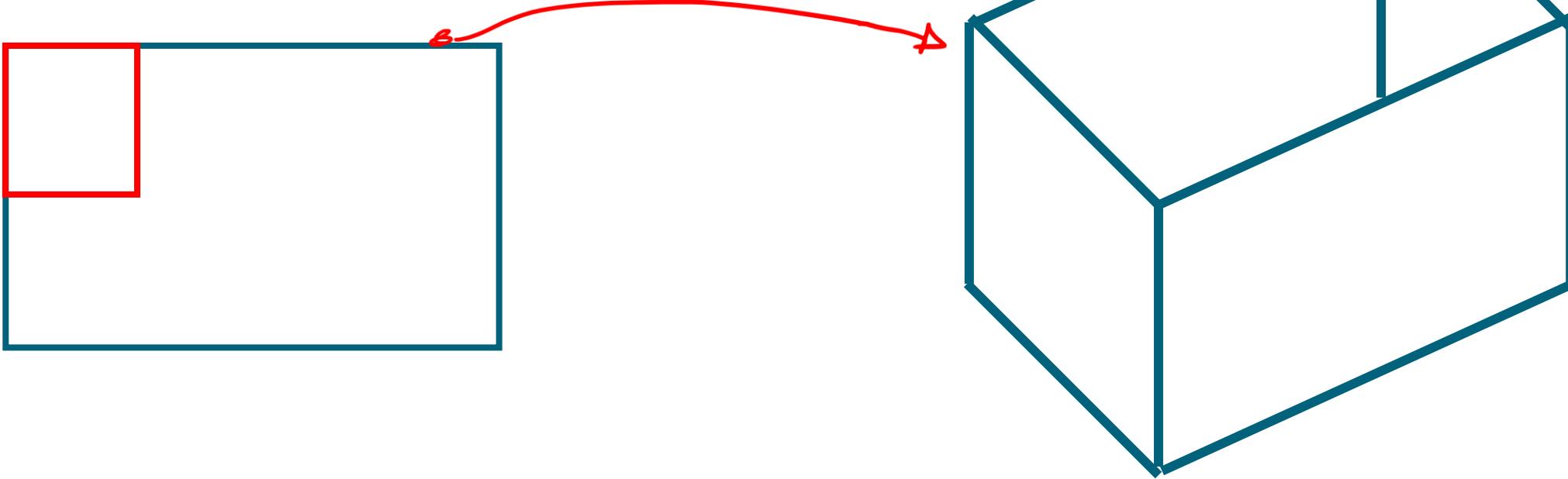
Solution

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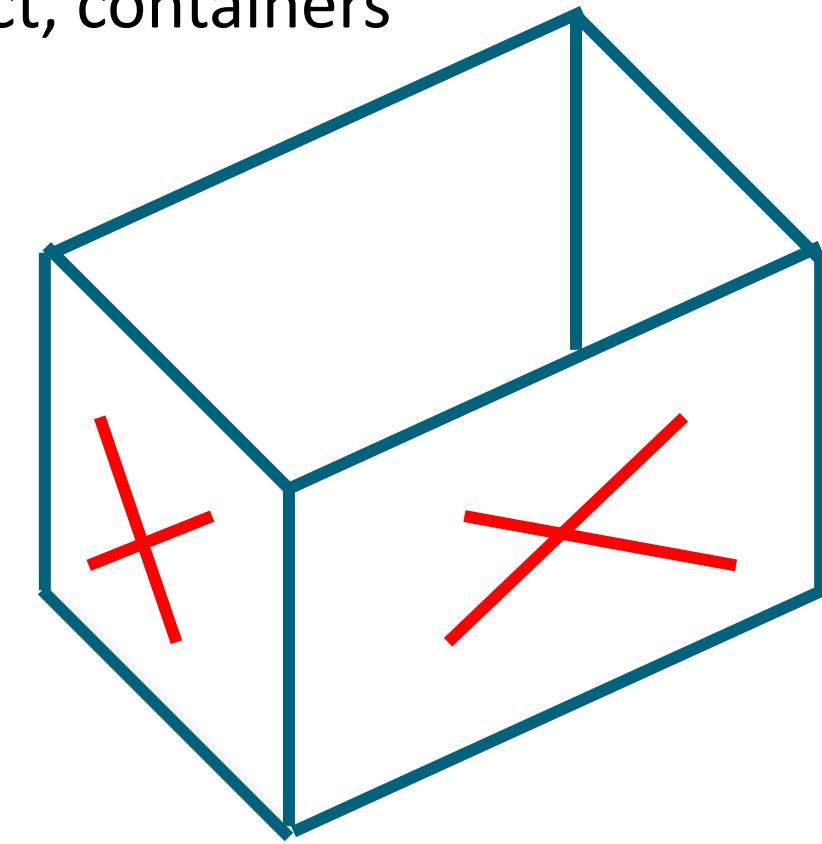
Solution

- The solution is based on a few observations.
- Observation 3: Tops of items are, in effect, containers



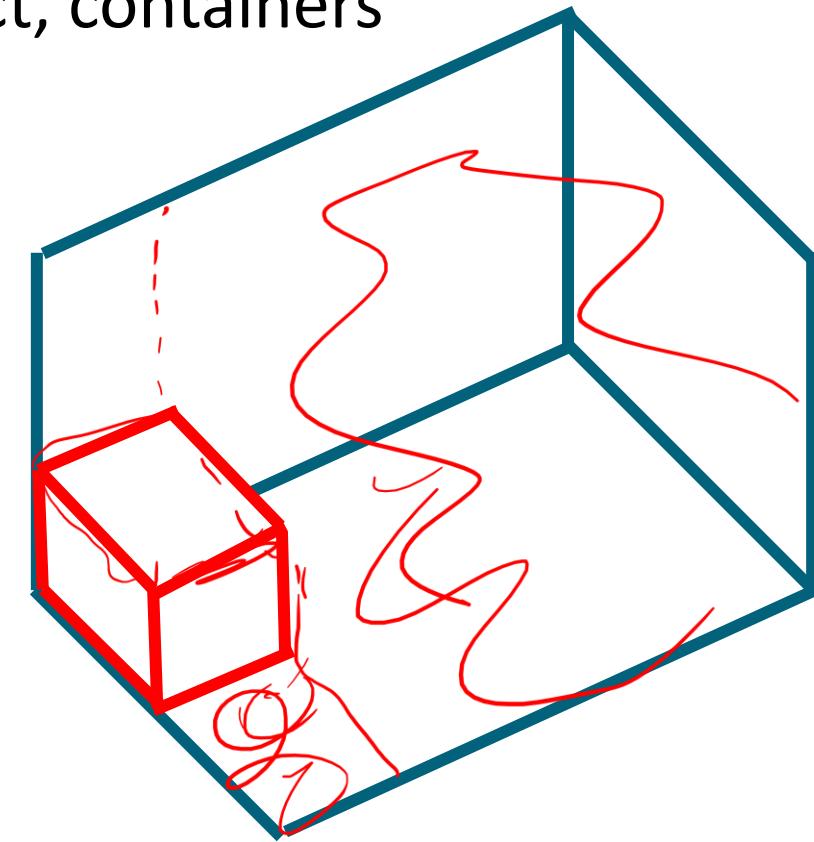
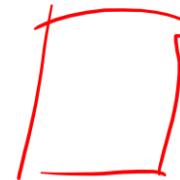
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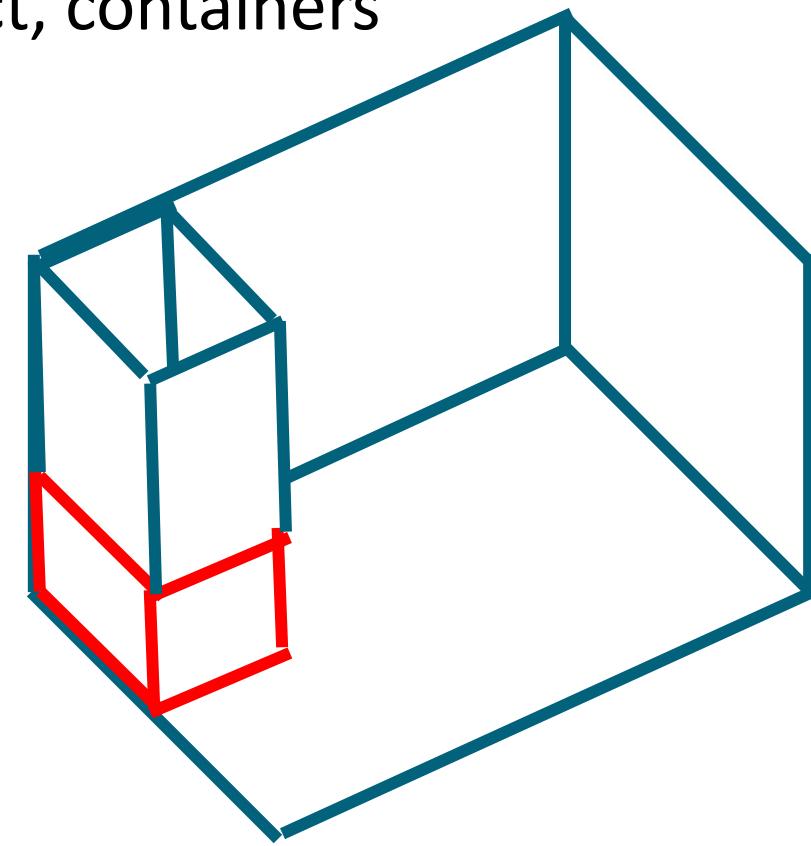
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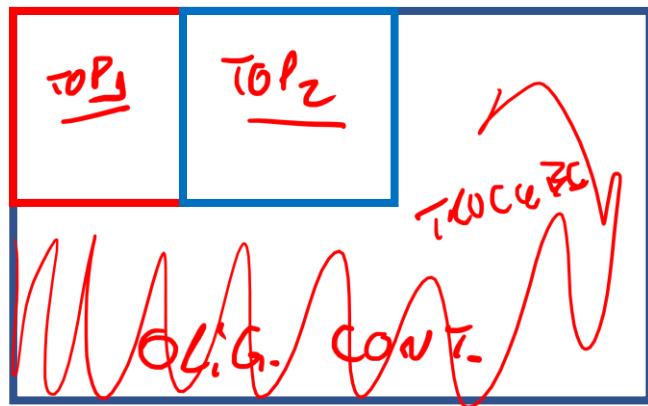
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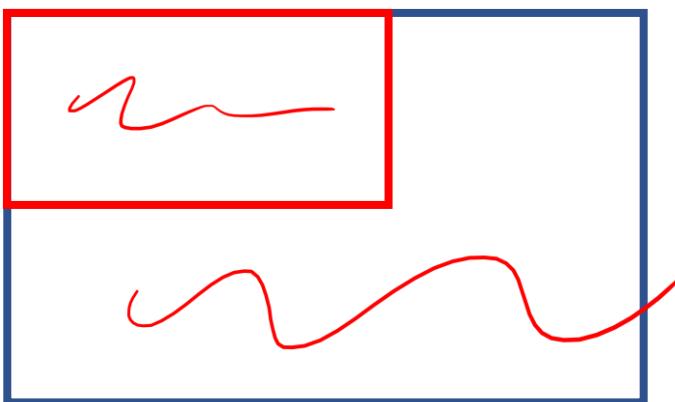
Solution

- The solution is based on a few observations.
- Observation 4: Adjoining tops (at the same height) can be joined



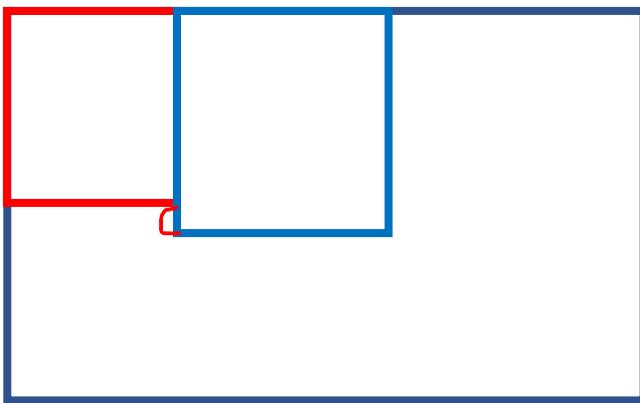
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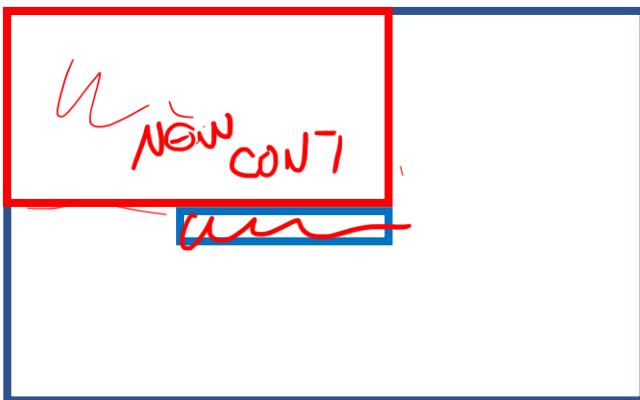
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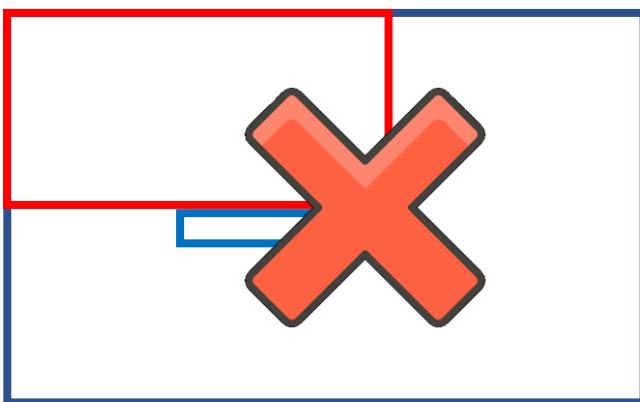
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Solution

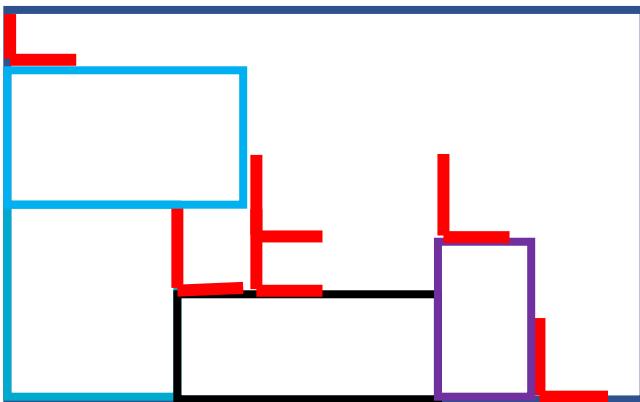
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Solution

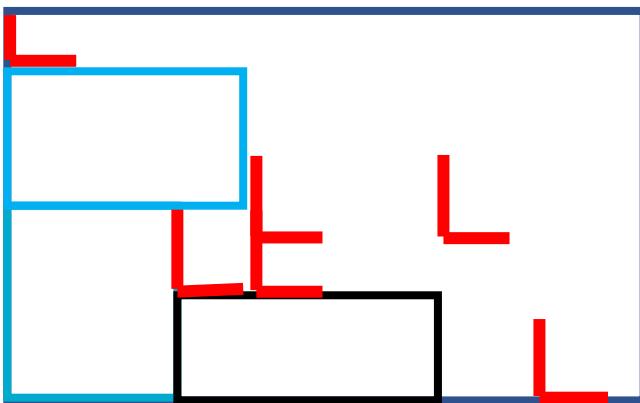
- The solution is based on a few observations.
- Observation 5: Taking things out of the middle is hard

EMPTY TIECON?



Solution

- The solution is based on a few observations.
- Observation 5: Taking things out of the middle is hard



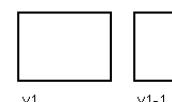
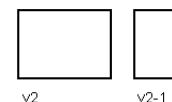
Solution

- The solution is based on a few observations.
- Observation 5: Taking things out of the middle is hard, so we just empty the container

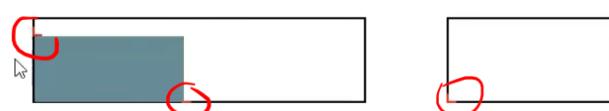


Solution

- Let's see it run ...



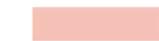
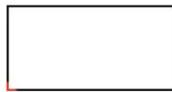
Solution



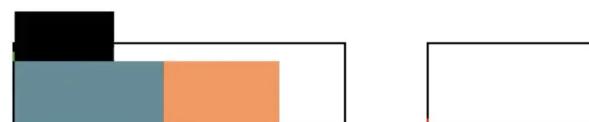
Solution



Solution



Solution



Solution

MAIN CLOCKED -LAICEC

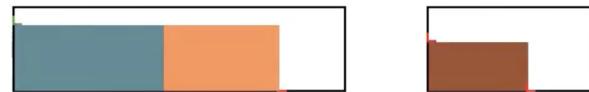
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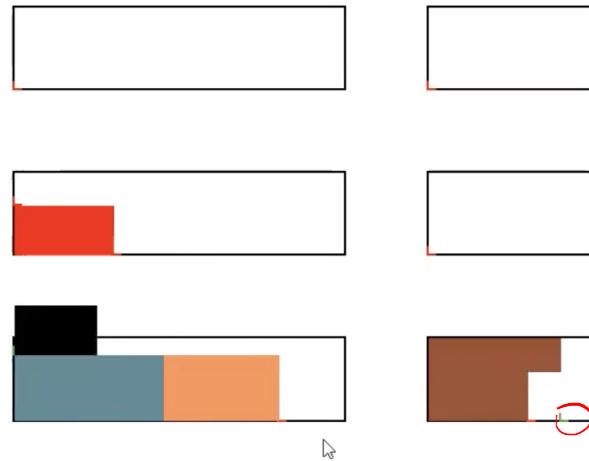
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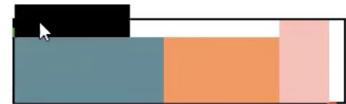
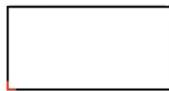
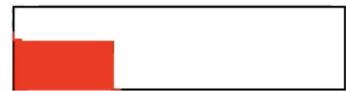
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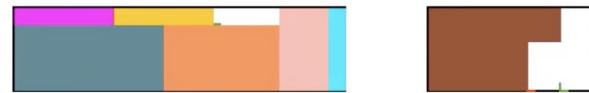
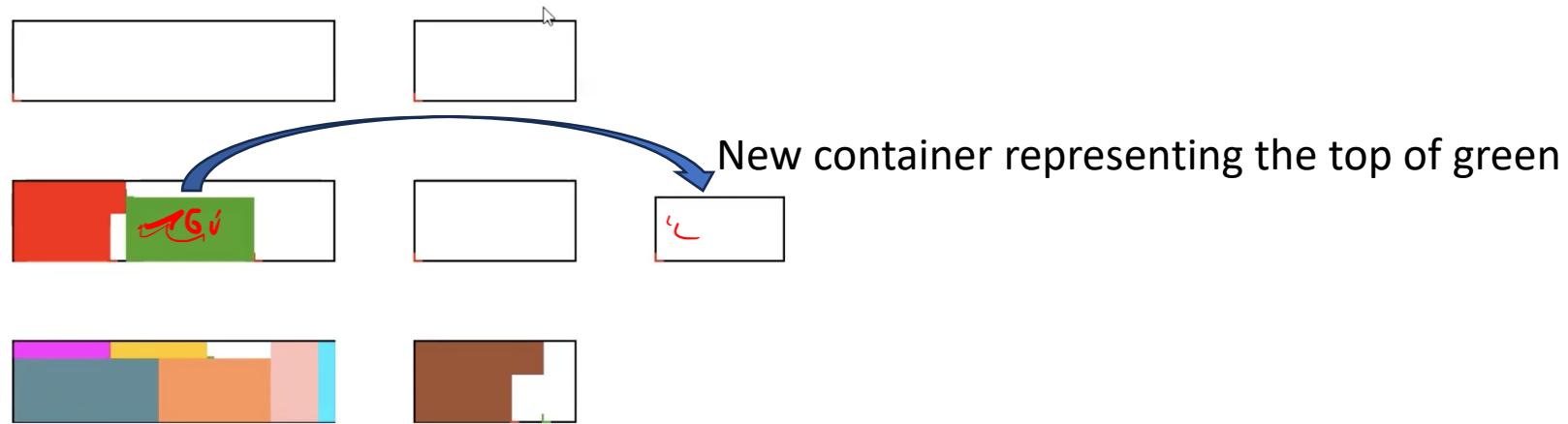
Solution



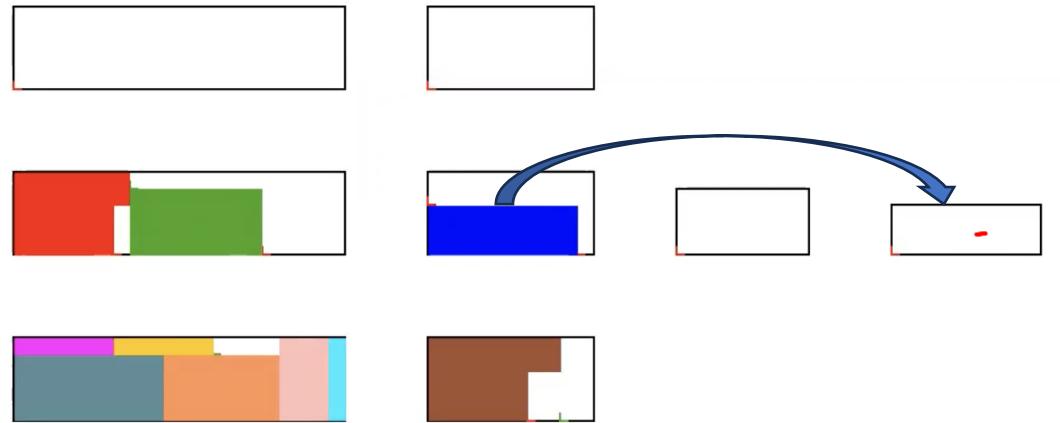
Solution



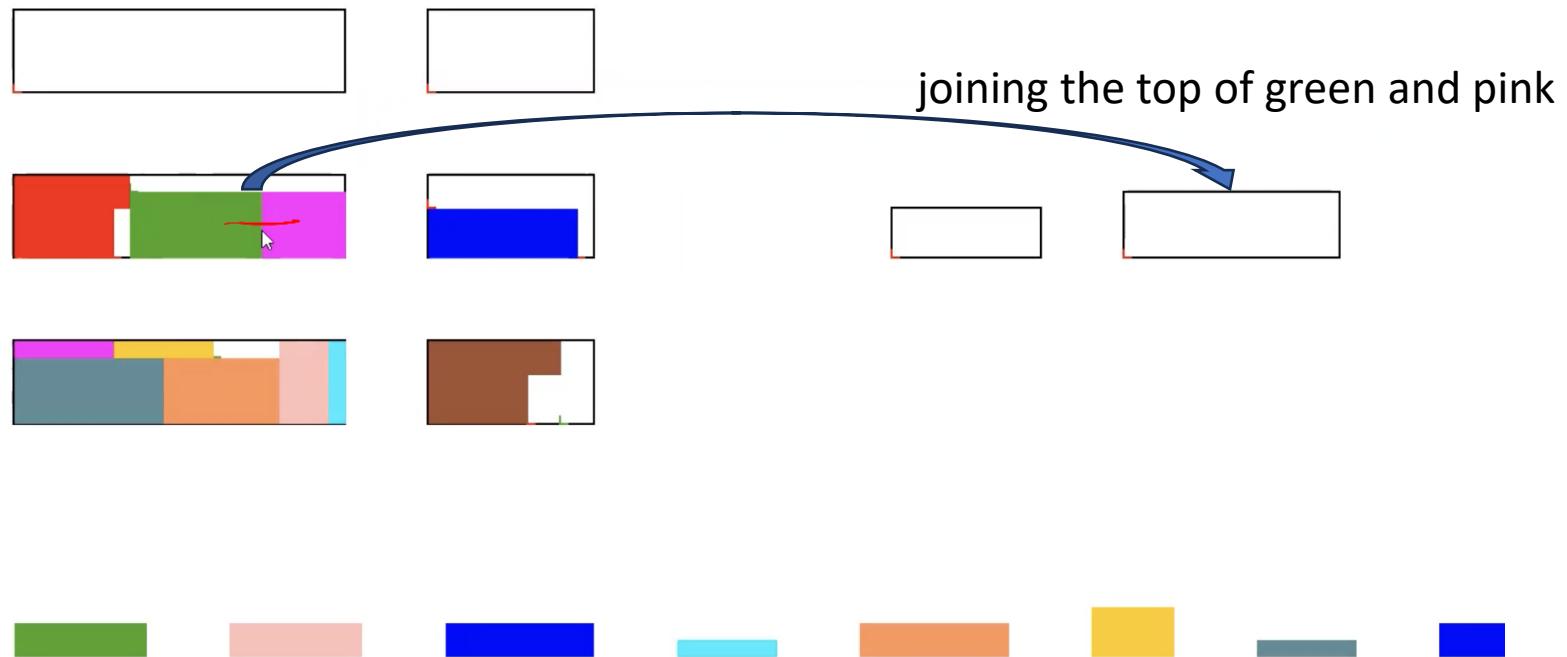
Solution



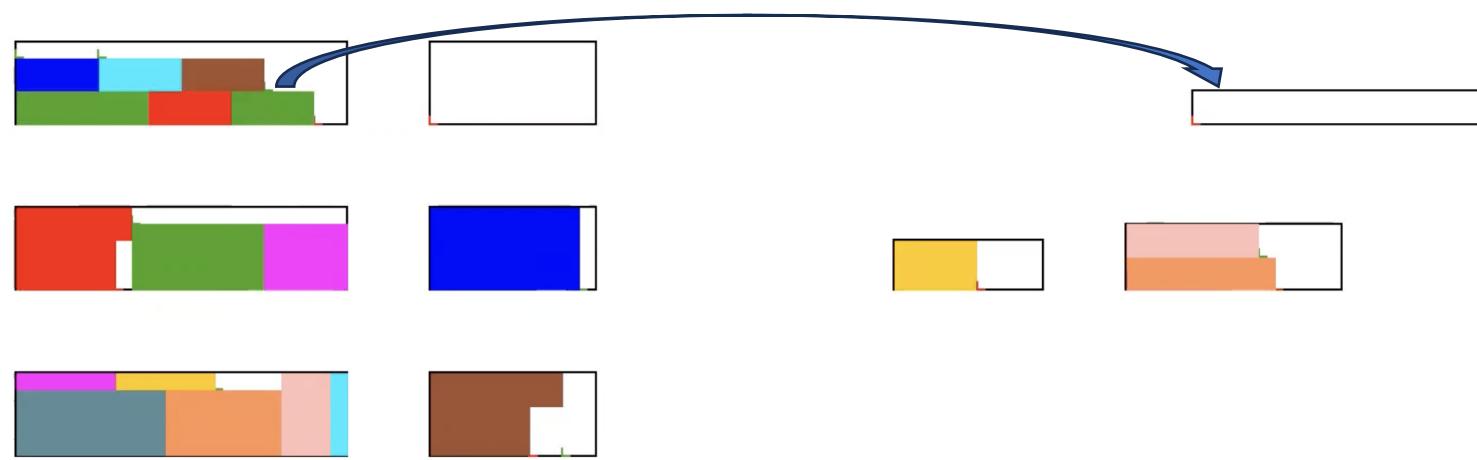
Solution



Solution



Solution



▷



Conclusion

- Constructing a solution is the first step in solving a combinatorial optimisation problem
- Building up a solution one element at a time is often a good place to start
- Greedy heuristics can give half-good solutions, but are often structurally unsound
- Some techniques can give a heuristic a bit of foresight, such as Regret

LOOK AHEAD

Next time

- Now we have our first solution ... can we improve it?

FRIDAY
QUIZ