	Synchronization Page No. Date						
4	Sunchronization or a						
	is the task of co-ordination						
2 10	is the task of co-ordinating the execution of the process in such a way that no						
	two processes can have access to Bame						
11	shared data and resources.						
1							
	types [on the basis of synchronization]						
	(1) Independent Process -> does not affect other praces						
4 '	(I) (n-onerative Prace of the prace						
理力	1 Co-operative Process -> affect the execution of other process						
	OF Other process						
-	Dana in the same i						
N	Race Conditionis a situation where						
	multiple processes access and manipulate						
7	the same data concurrently and						
	Outcome of the execution depends on						
	the order in which the instructions						
	execute a compagno of the company of						
	El El Distra mala has conflict la manin						
*	* Critical section, is a code segment						
	that can be accessed by only one						
(a processigationatione of partition for						
	· Critical sections contains shared						
1	variables which need to be synchroniz						
	to maintain consistency of the data						
2	variables, barred A : parting be being the						
C.	Jeens Jick of Barriel Par Victorio						
	Lundo, of the of hourselle was						
4	work & word of hot for a mind						
1 %	Entry Section in Thom						
	critical section						
	Exit Section						
	remainder section						
	while (True);						

Page No.

DENtry section: Each process needs permiss
to enter critical section, each process
goes to critical section through entry sec
i) Exit section: Each critical section
is terminated by exit section.
ii) Remainder section: The code after
exit section is remainder section.

A Solution to Critical Section must satisfy Pollowing three condition requirement

in vietical sugion section and other the no other process is allowed to execute in the vietical section.

Progress: In no process is executing in cuitical section and other process is waiting outside the critical region then only thoese process that one not executing in remainder section can participate in deciding which will enter next in cuitical section.

Bounded waiting: A bound must exists on number of times a process es are allowed to enter witical section, after a process has made a suguest to enter critical section.

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1.	Type	11	-				

* Semaphore is a vaniable used control access to a commor by multiple processes and ow critical section problem. est is an integer variable, which can be accessed only through two operations wait() and signal().

Advantages

allows

Only one process into withcal section

1 There is no resource wastage.

· Types of semaphores

D Binary semaphore

1) Counting semaphore.

* Bounded-Buffer Problem. also known as

· Produces consumer problem is

a classical synchronizational problem.

. We can solve this by using semaphores. · by meating two counting semaphores "full" and "empty" to keep track

of the current number of fulland

empty buffers respectively.

· Priodicers produce a product and consumers consumes the product, but both use one of the container

each time.

· Producer: - O Creates item and adds to buffer m po not want to overflow buffer

· consumer: - O Removes items from buffor.

1 Do not want to get ahead of produces.

