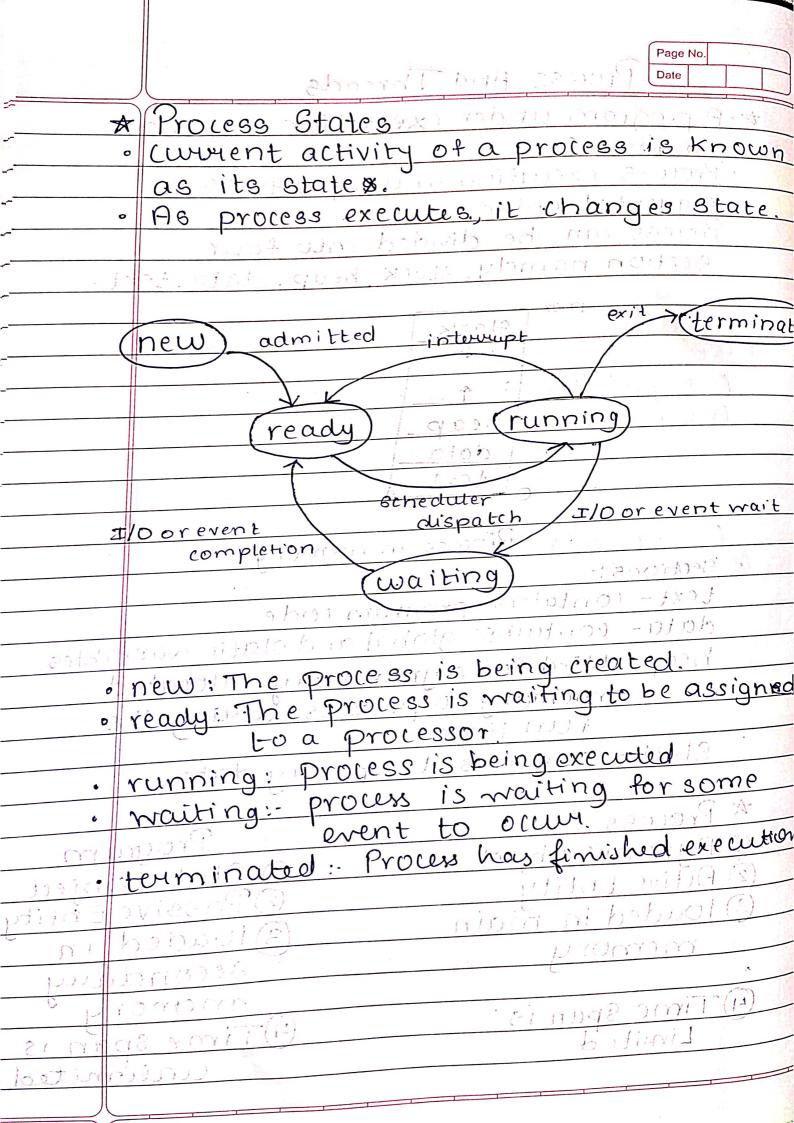
## Page No. Process And Threads A program under execution is called as Process. Process execution must progress in sequential order. Process can be divided into four section namely, stack, heap, data, text. Stack heap data text Process in memory Sections: text- contains program code data- contains global and static variables. heap-used for dynamically allocated memory to process during its run time Stack- contains temporary data Program \* Process Process 1 Dynamic Object De Static Object 2) Passive Entity 2 Active Entity 3 10aded in main 3 loaded in Becondary memory memory 4) Time span is 4 Time Span is unlimited Limited



	Page No.				
	Date				
-	* Process Control Block (PCB)				
	14 man 1 10 1 10 > 0 pm 11 04 11				
	(10100x +01				
	Process State				
	Process Number				
	Program Counter				
ted	registers (no him)				
	Millianton La memory limits				
	l'integralies listion openities				
	000				
	(0 hia) 21,0019				
11/23	1. 7				
	Process state: - Specifies curvient state of process				
	Process Number: Each process has some				
	unique identification number				
	called as Process Indentification				
	number (PID)				
	Program counter: Pointer indicates the				
	address of next instruction				
	to be executed for process.				
	Registers: - Registers Like stack pointer, index				
	Registers: Registers Like stack pointer, index oregisters, general-purpose regustar et where process need to be				
	ete where process need to be				
	stoned for execution for				
	ounning state.				
	memory limits: Memory used by the operating system				
-	File management: - Handles lists of open files.				
	<del></del>				

Millian ....

6T>MT>LT				
Page No.				
Date				
ypes of Schedulers				
ongterm Scheduler				
selects a job or process from jobpool				
from secondary memory and loads				
nto main memory.				
Executes less frequently				
controls degree of multiprogramming.				
Secondary				
(new) (ready)				
memory obposed A				
maille of through a separate A.				
dillor about the property of				
Short-term Scheduler				
selects a job from ready queue and				
submit to CPU.				
As it selects only one job at atime, it is				
invoked very frequently.				
Main objective is to increase system performance.				
ww Gramm 15 10 01 1 44 bill May 4617.				
(ready) (running)				
Texecution (13) (17) (10) (10)				
232593 1 17 19 19 19 19 19 19 19 19 19 19 19 19 19				
Mid-term Scheduler de sort of sorting				
Mid-term Scheduler is part of smapping				
so it is known as smapper/smap scheduler.				
gramid-termischen process, and to				
the process from main-memory to				
secondary-memory and vice-versa.				
· As it removes the process from memory it reduces the degree of multiprogramming				
14 security the argine of				
(mim) (sm)				

() L

* Context Switch	4-1 -171111
	27.70
requires serving the state of	of the
old process and loading t	he saved
State of the new proces	38. 1NIS
requires serving the state of old process and loading to state of the new process to stack is known as contact the multiple process to a stack is multiple process.	ext switch
· Enghles multiple process to 1	ise single
CPU.	
Chip 20	(4/2)
* Threads	
. A throad is a flow of ex	ecution
through the process couce	, // (/)
its mun program co care	C) 0 1 7 1 () (4)
that keeps track of inst	ructions
to execute next, system which hold its convent venicobles and stacks	siegisters
which hold its convent	working
venicibles and stacks	which
hold execution history	41. mom
· Light-weight Process & LWF	9 mm
· Babic unit of CPU whilize	ation,
· Subset of process.	
- Run in shared memory &	spaces.
· Showes tile description?	mul-bim (
shares tile system conte	xt:him.+
· Shares signal handling	21 11 0,0
Threads are dependent	· Winn X
of Fireman-aidin mont 200	17401401
Office of Air Francisco	1000112
The state of the s	or sieA

the device of walthough

			Page No.  Date		
	*	User Level Threads	Kernel Level Threads		
	0	ULT are faster to	OKLT are stower		
		create and manage			
	2	Threads implemented	2) Threads implemented		
		at user revel known	at Kennel level		
•		as ULT	known as KLT.		
	3	can run on any	3 Specitic to		
		0.8	11000		
		Benefits			
		Efficient Communicati	on:- lommunication		
		between multiple .thr	between multiple threads is easier		
		as the throad a ah	outes common		
	-	address space:			
	(2)	Computational speedup: On single			
4		can be executed speedily by creating multiple threads in the			
I.o.		can be executed speedily by			
		creating multiple threads in the			
1		process> contd[nextpg]			
	11	•			

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		Date	
X	Benefits	choord lovel vest	
(3	Responsiveness:	Allows a program to	
in the R	run oven it pa	rt Bf. it is blocked	
100 100	Using multithreading		
į	of leavily to	awant in the war to	
4	Resoure Sharing:	- All threads of aprocess	
	Shares resour	ces such as memory,	
	douta, files etc.		
	·		
		Persisting & Persisting	
0.0	to between it empli	Common de l'amment	
	Harton J. Blanca	it significan accounted to	