

NPTEL MOOC

PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 6, Lecture 1

Madhavan Mukund, Chennai Mathematical Institute

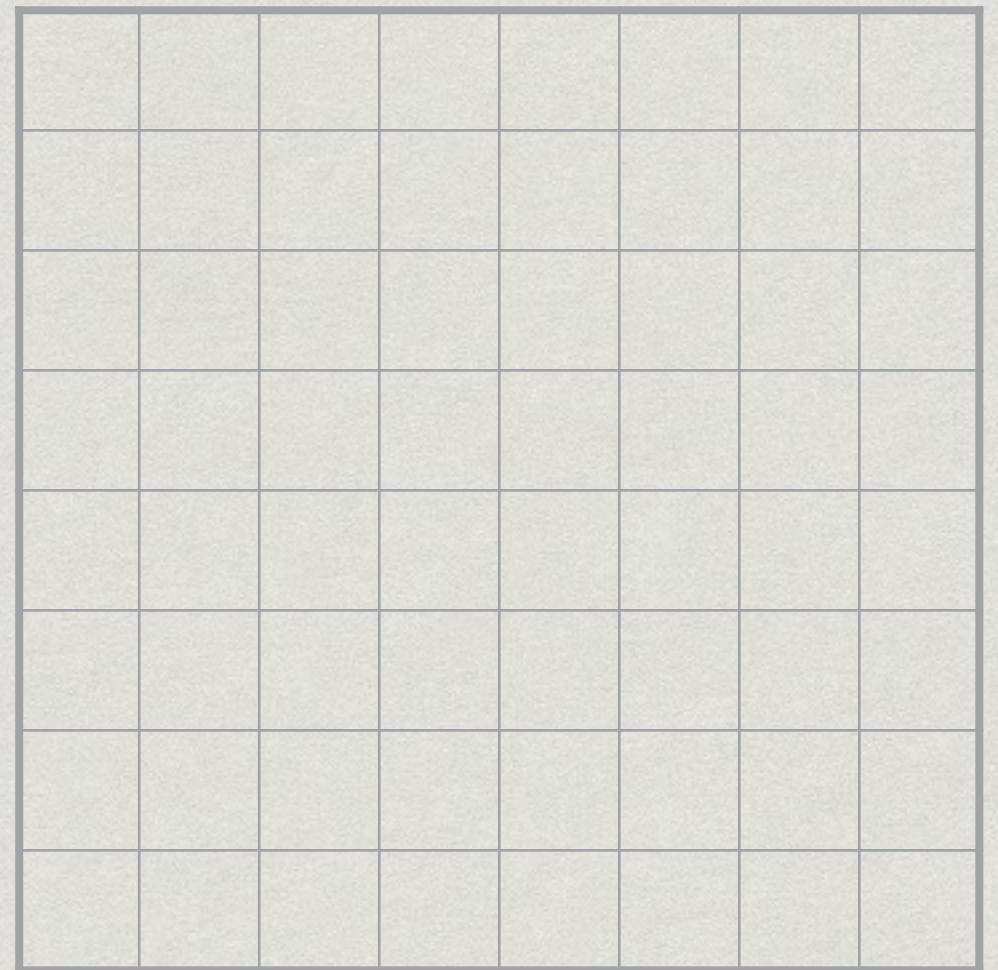
<http://www.cmi.ac.in/~madhavan>

Backtracking

- * Systematically search for a solution
- * Build the solution one step at a time
- * If we hit a dead-end
 - * Undo the last step
 - * Try the next option

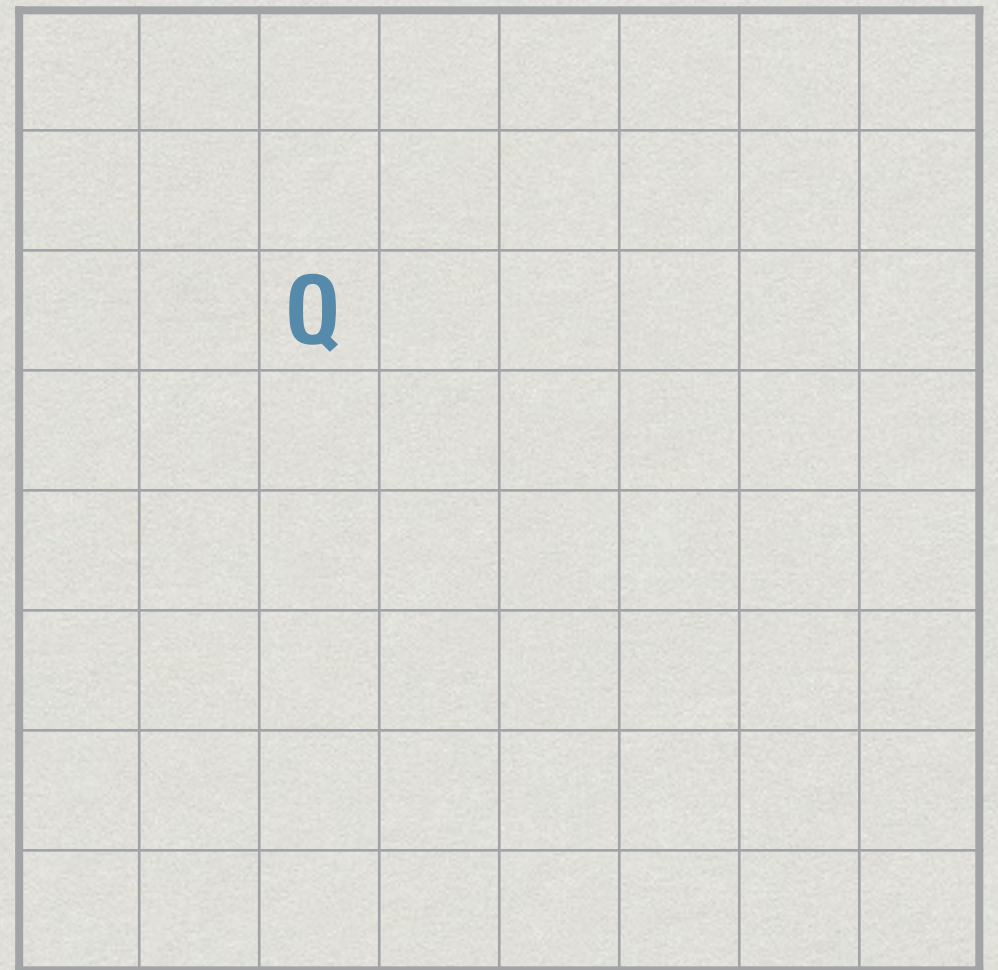
Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal



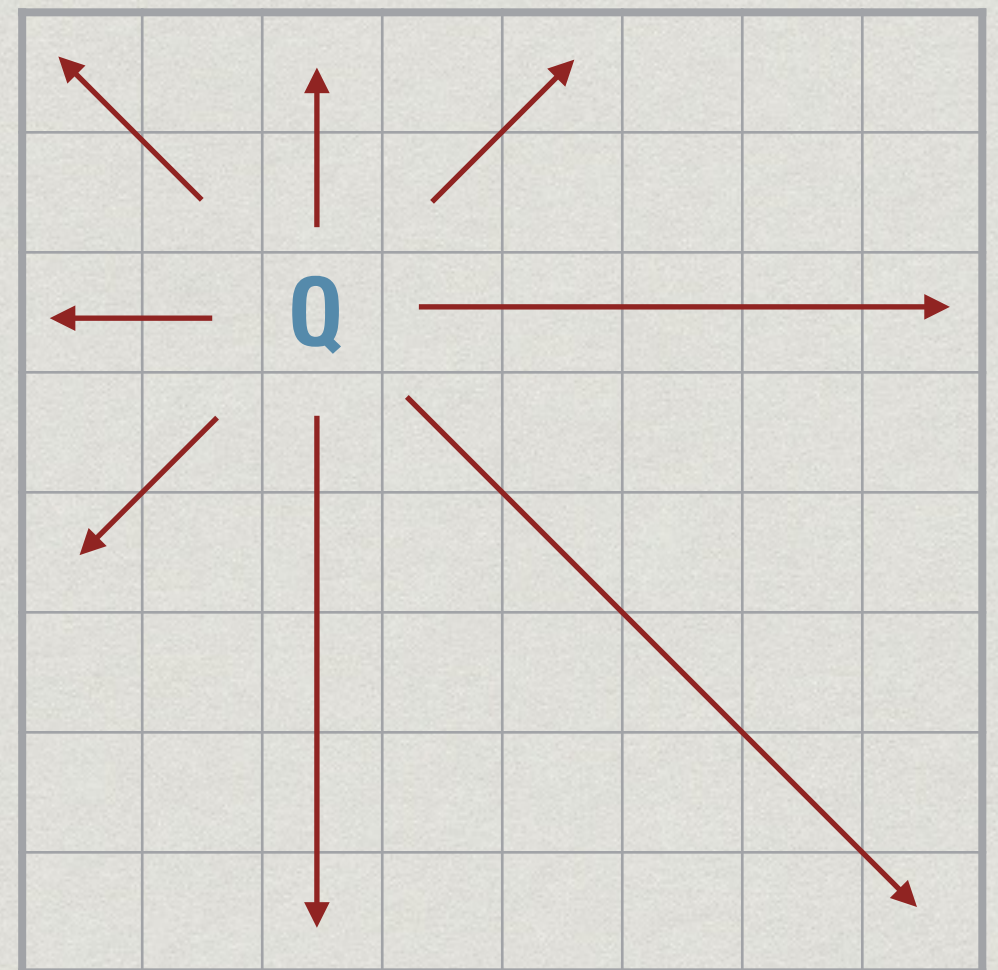
Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal



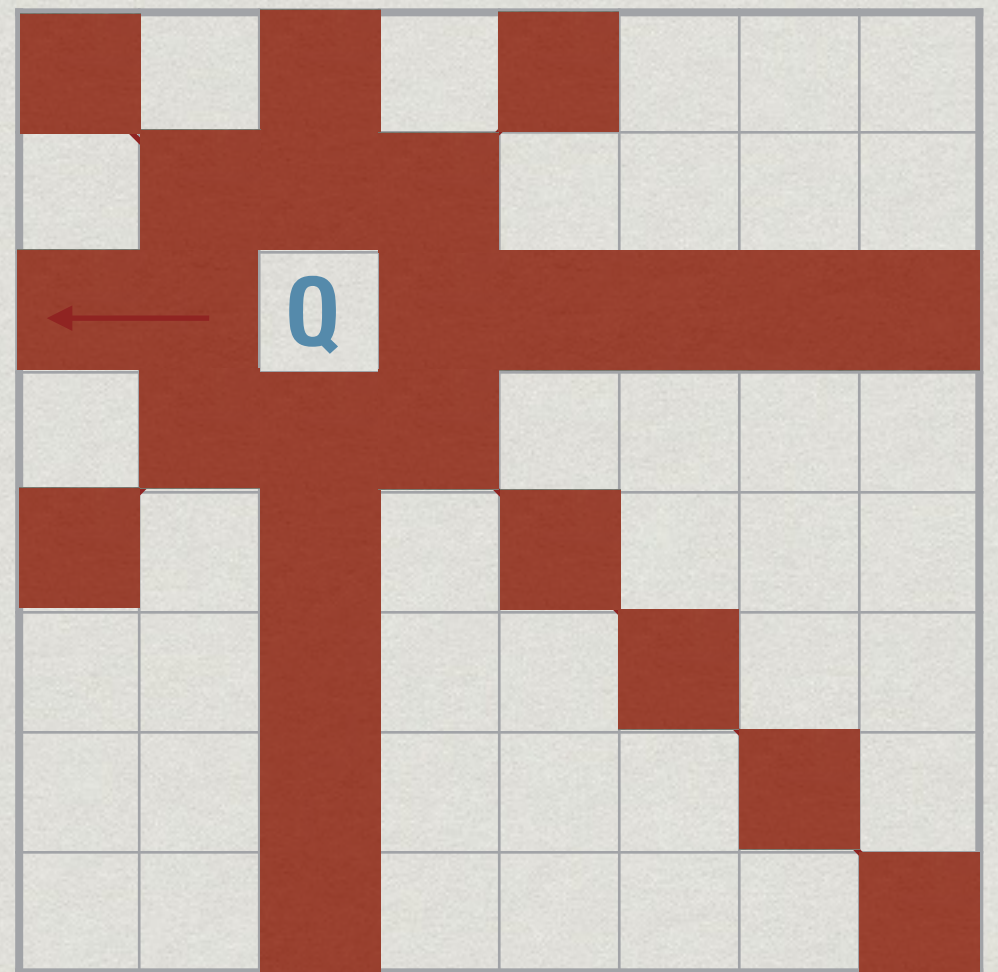
Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal



Eight queens

- * Place 8 queens on a chess board so that none of them attack each other
- * In chess, a queen can move any number of squares along a row column or diagonal

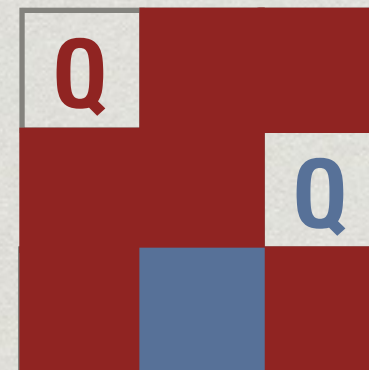


N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible

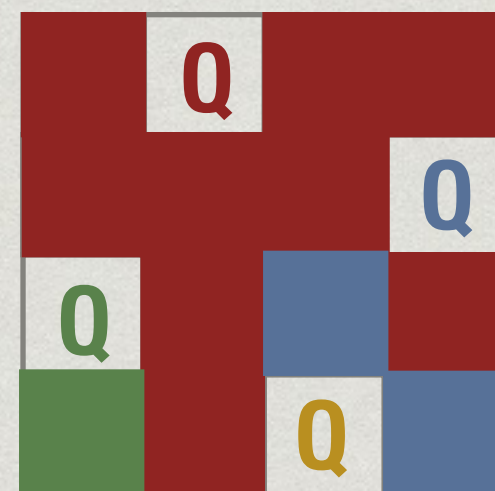
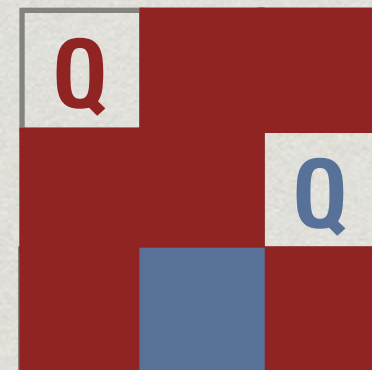
N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible



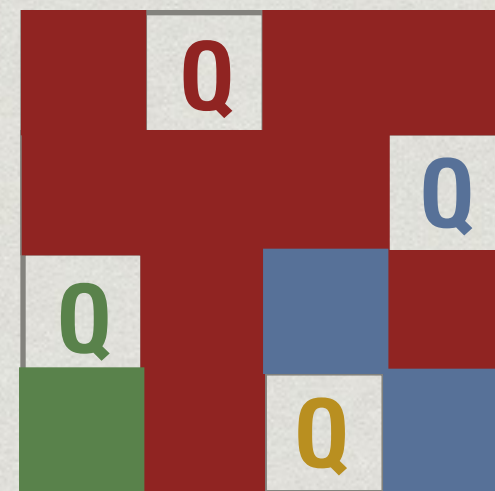
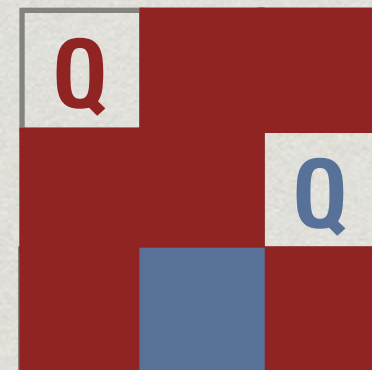
N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible
- * N = 4 is possible



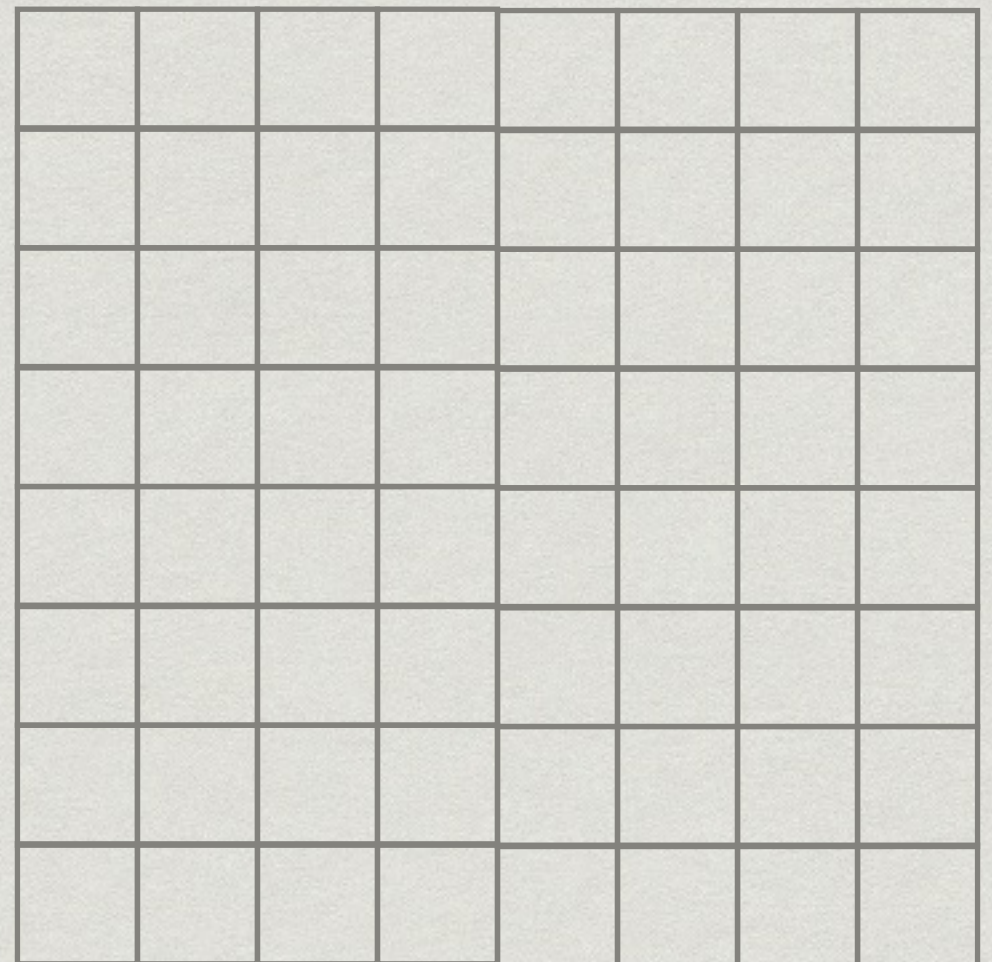
N queens

- * Place N queens on an N x N chess board so that none attack each other
- * N = 2, 3 impossible
- * N = 4 is possible
- * And all bigger N as well



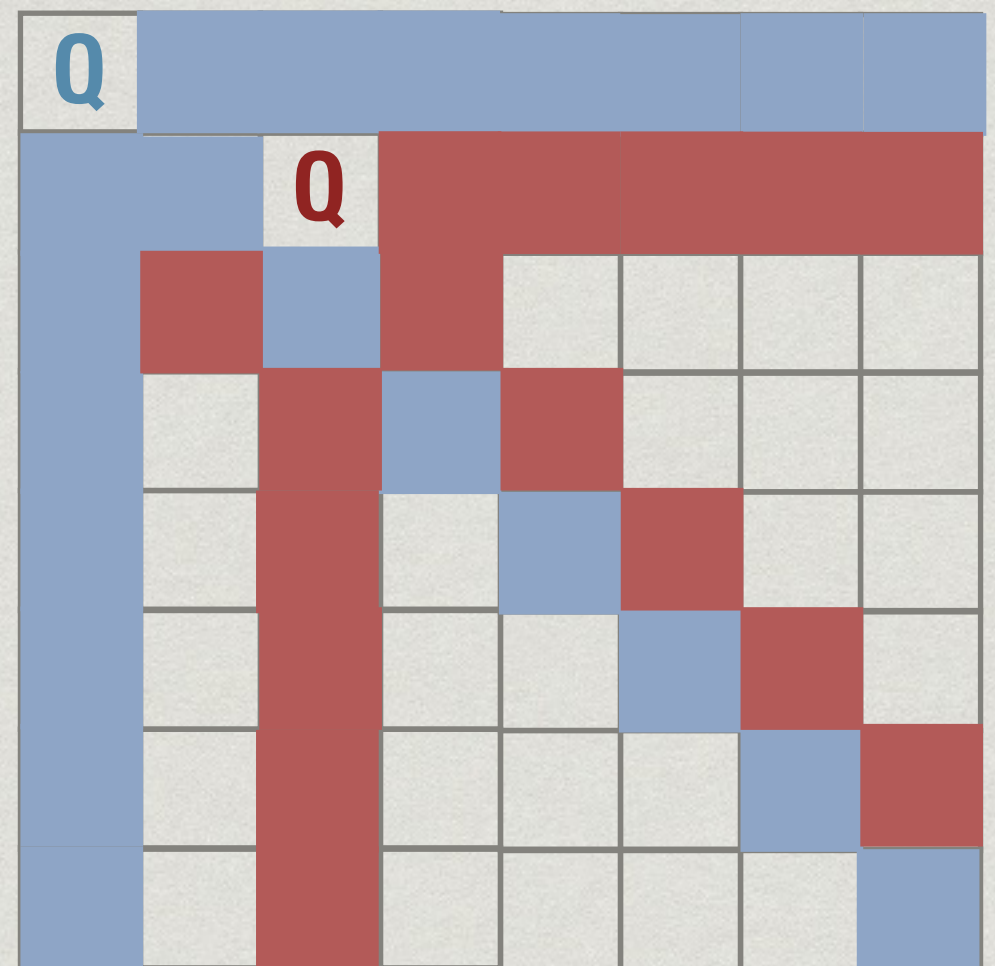
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



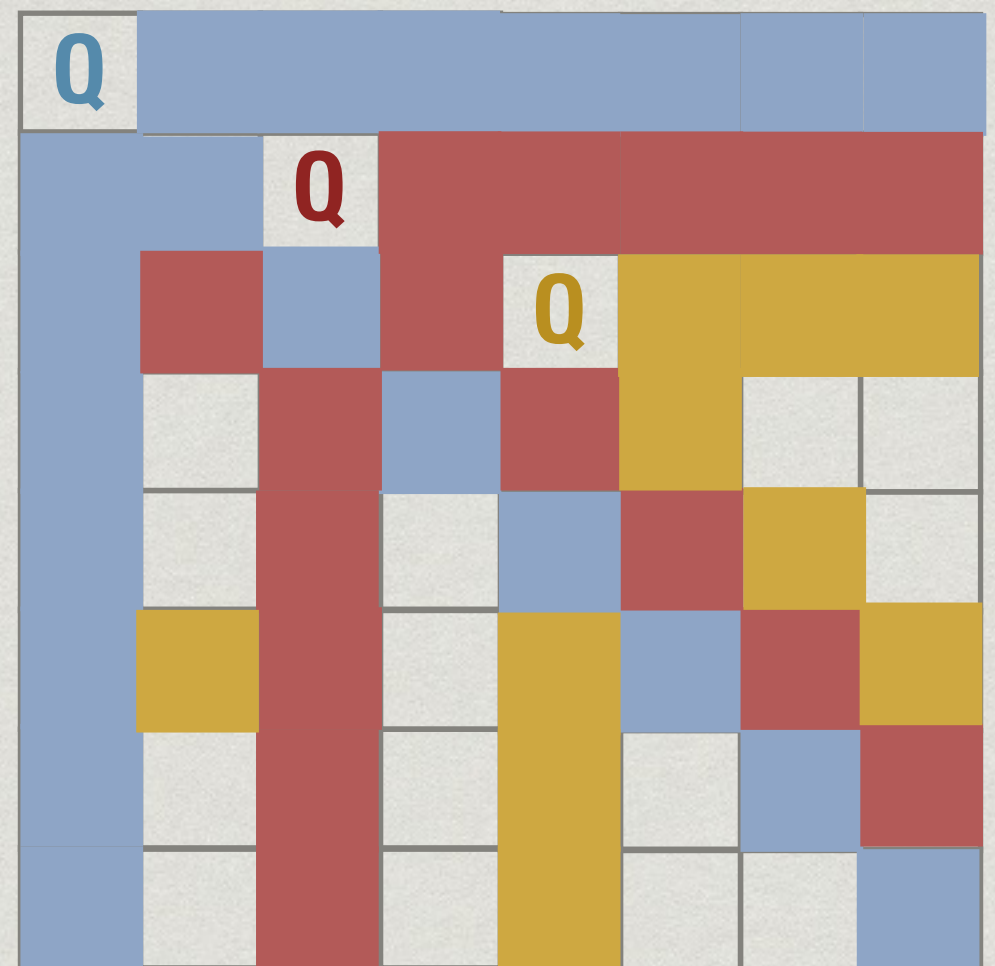
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



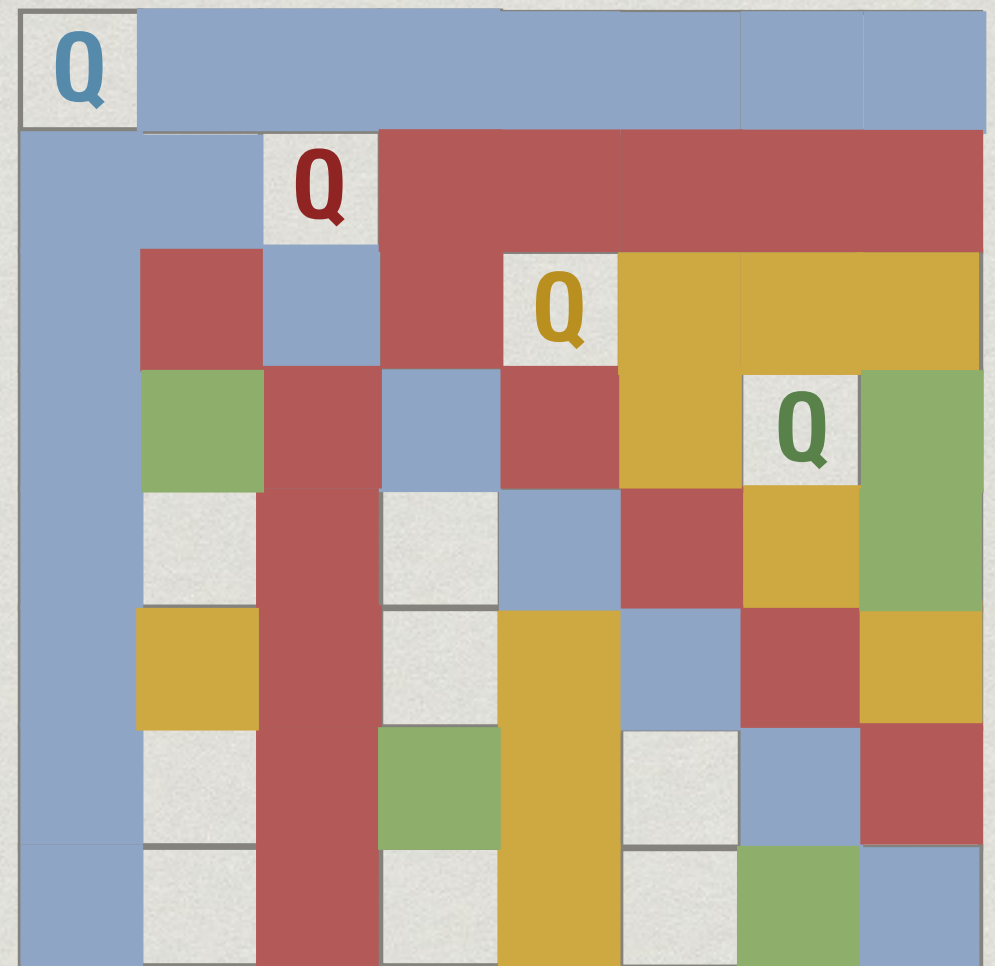
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



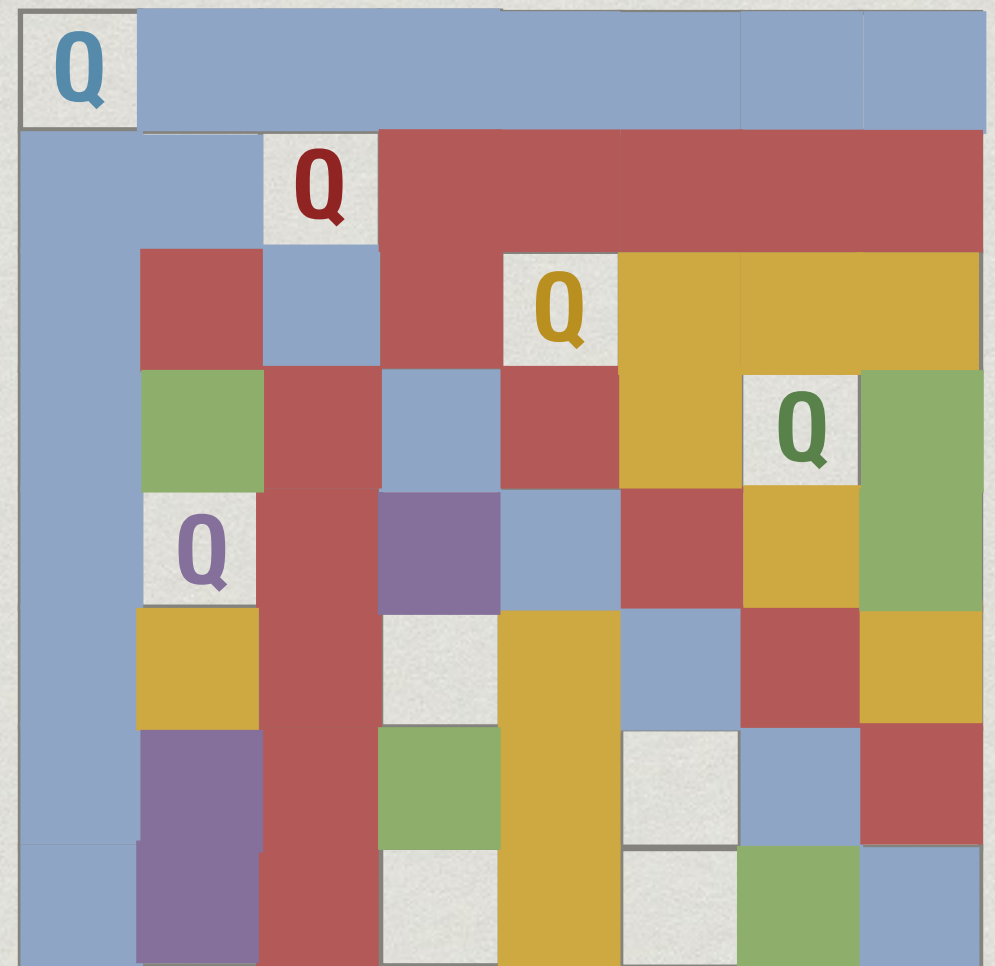
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



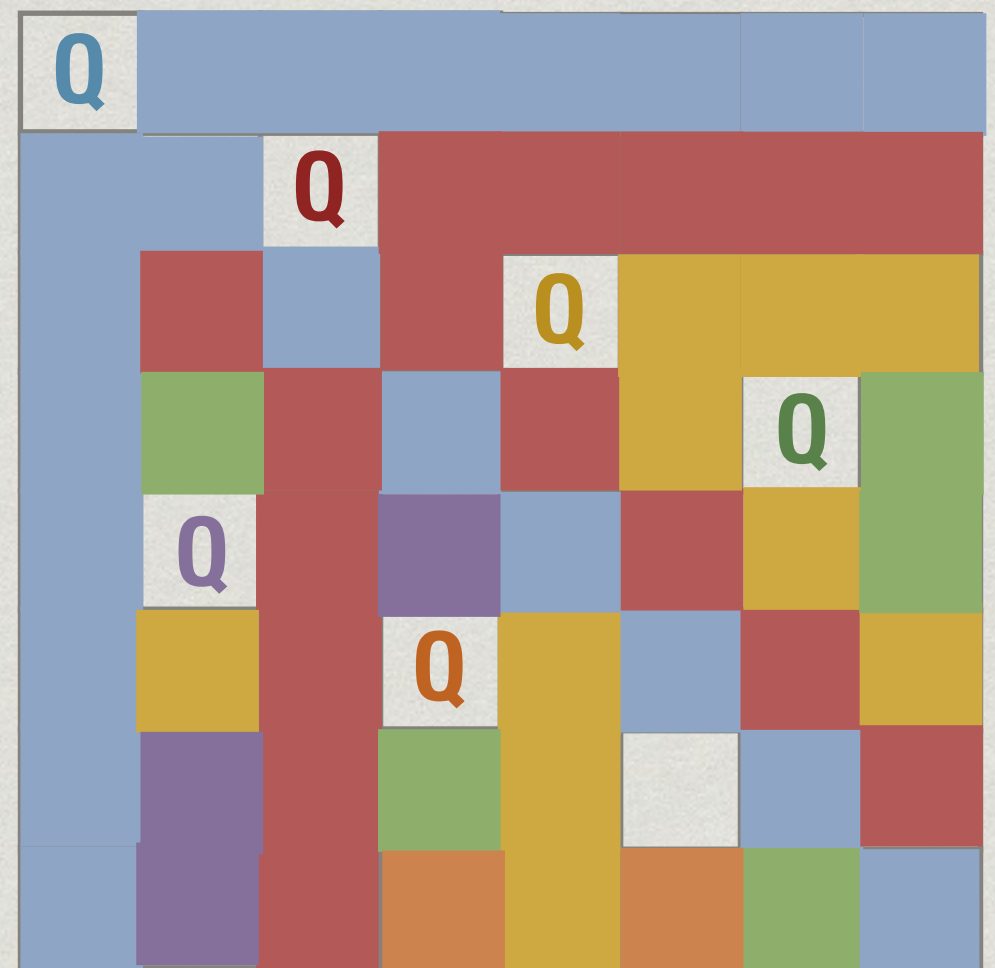
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



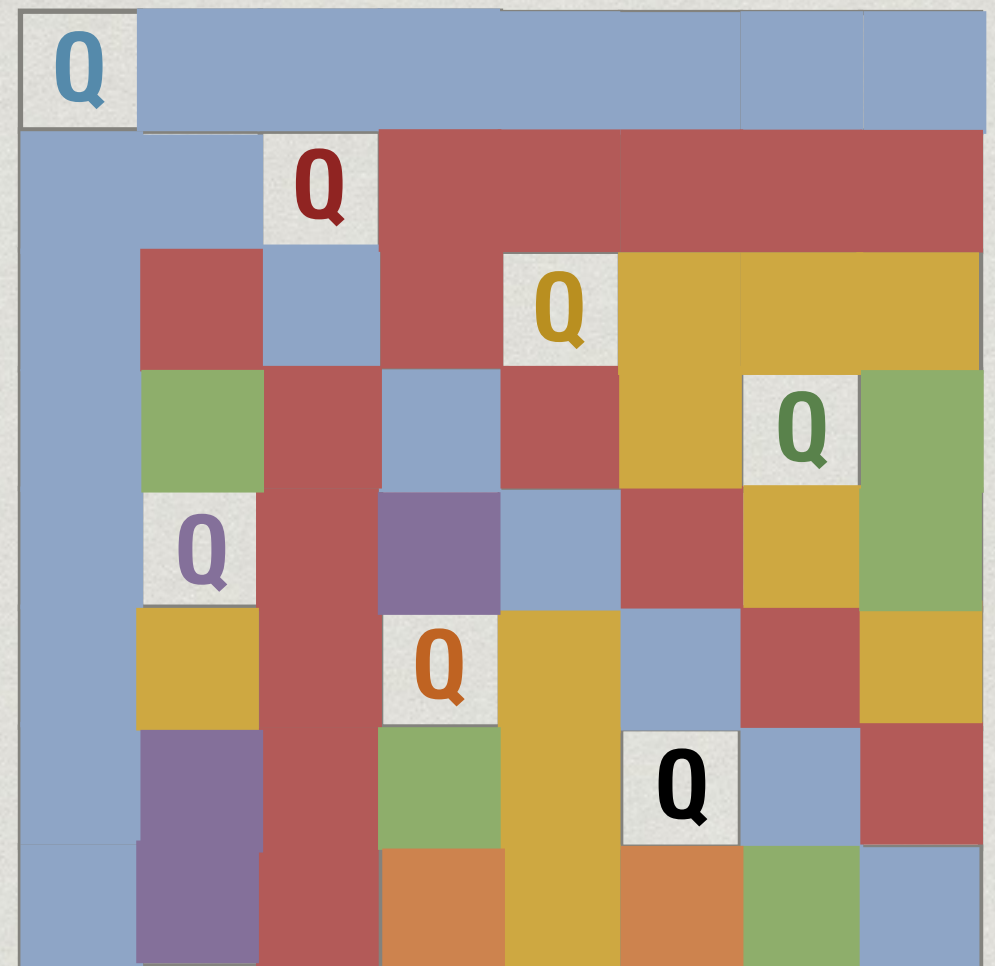
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



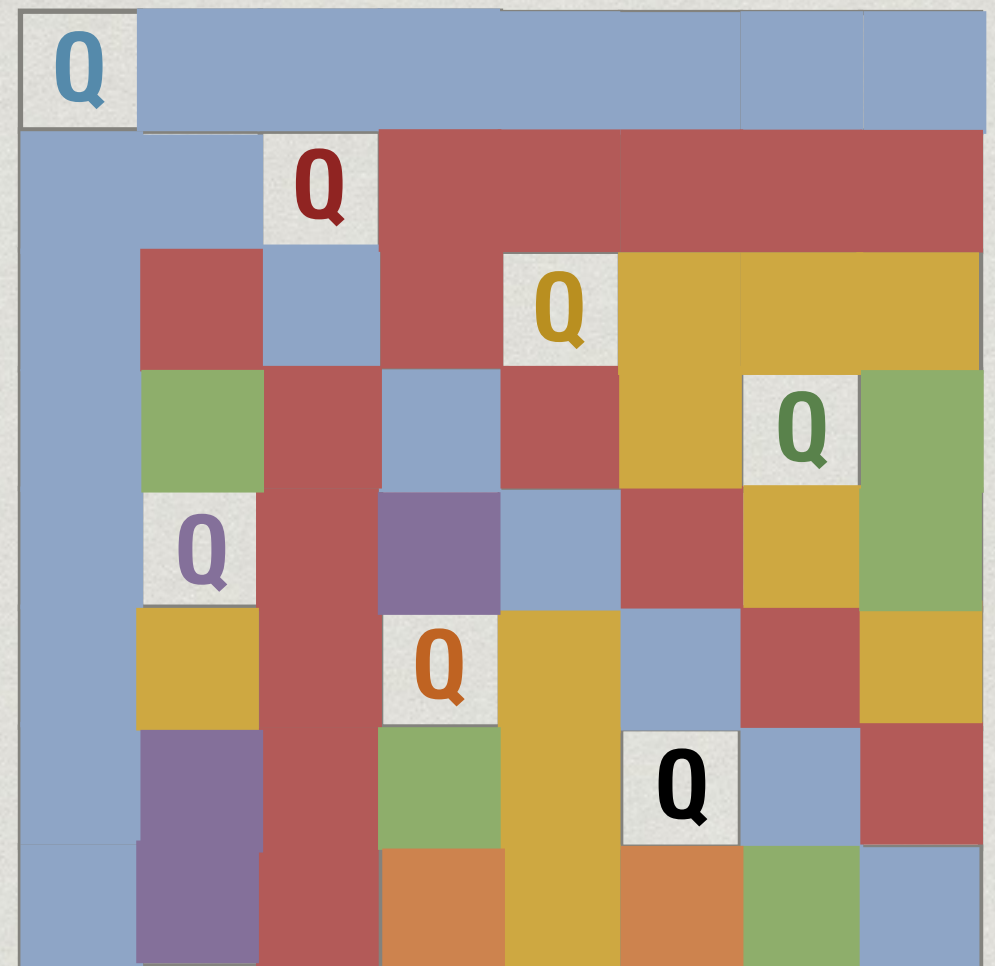
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column



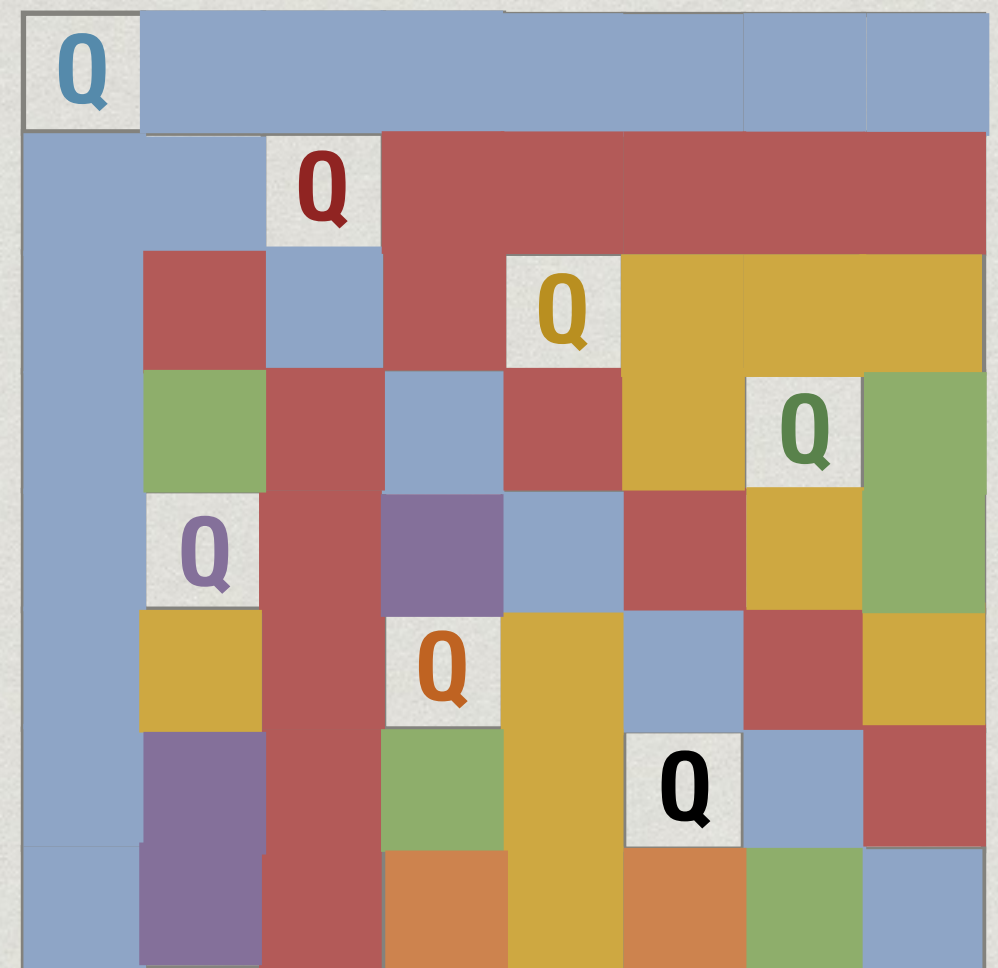
8 queens

- * Clearly, exactly one queen in each row, column
- * Place queens row by row
- * In each row, place a queen in the first available column
- * Can't place a queen in the 8th row!



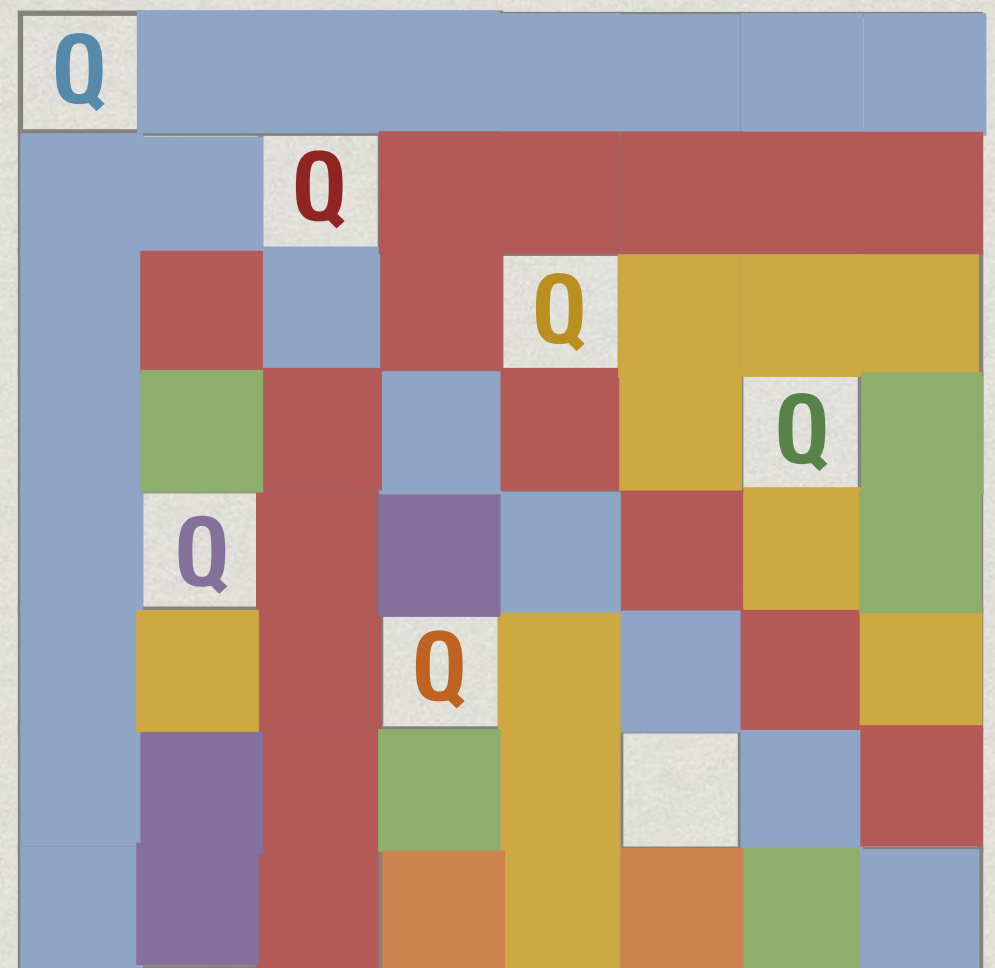
8 queens

- * Can't place the a queen in the 8th row!



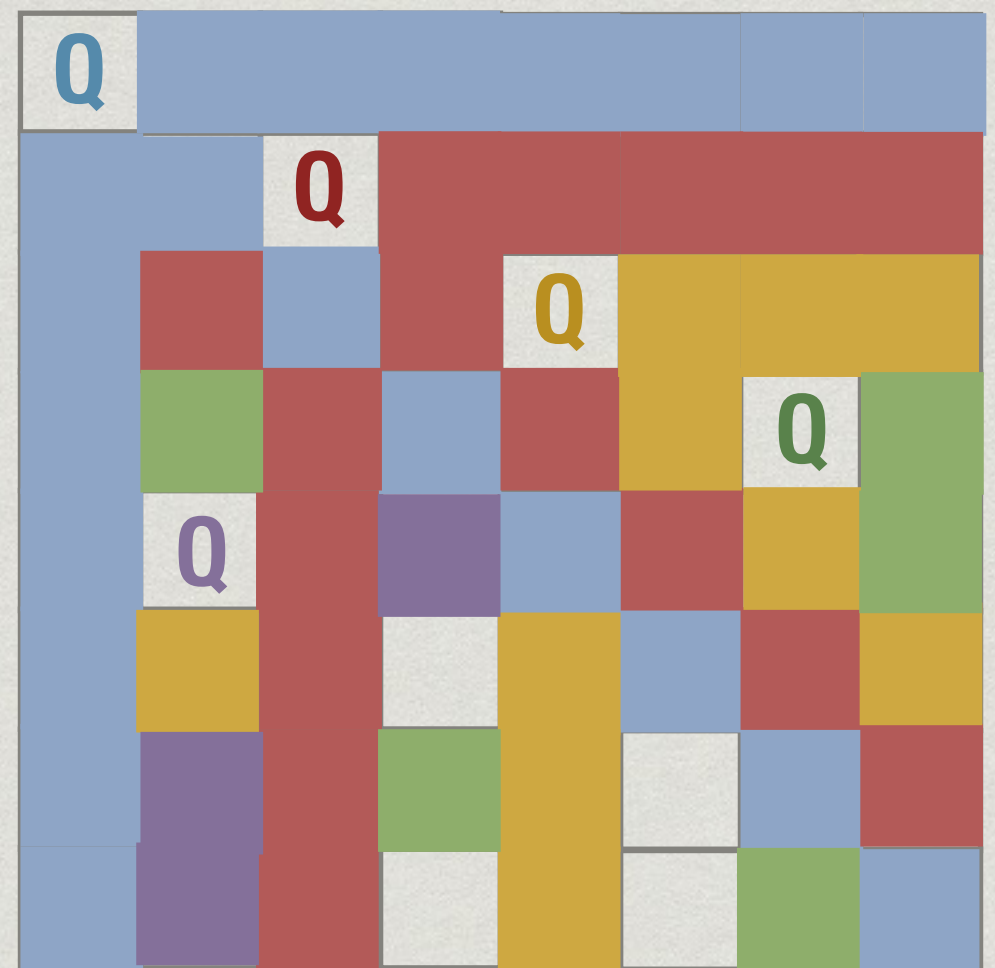
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice



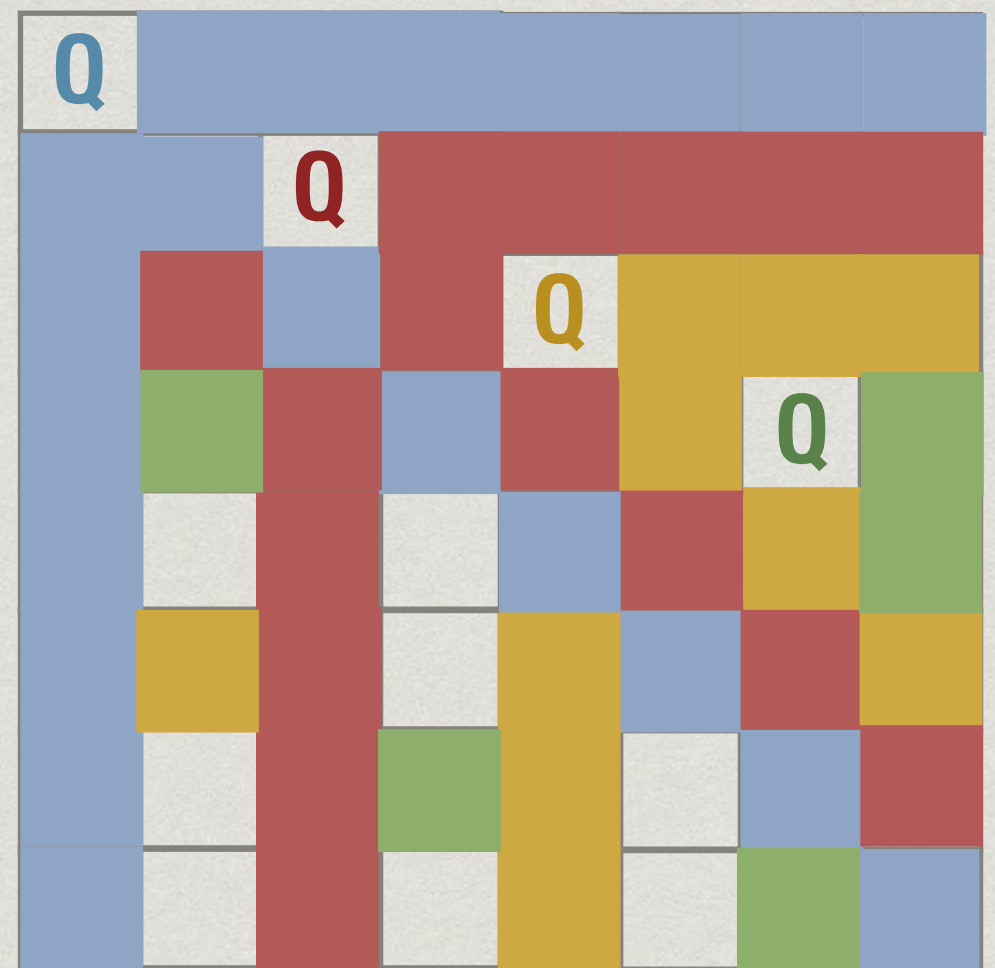
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice



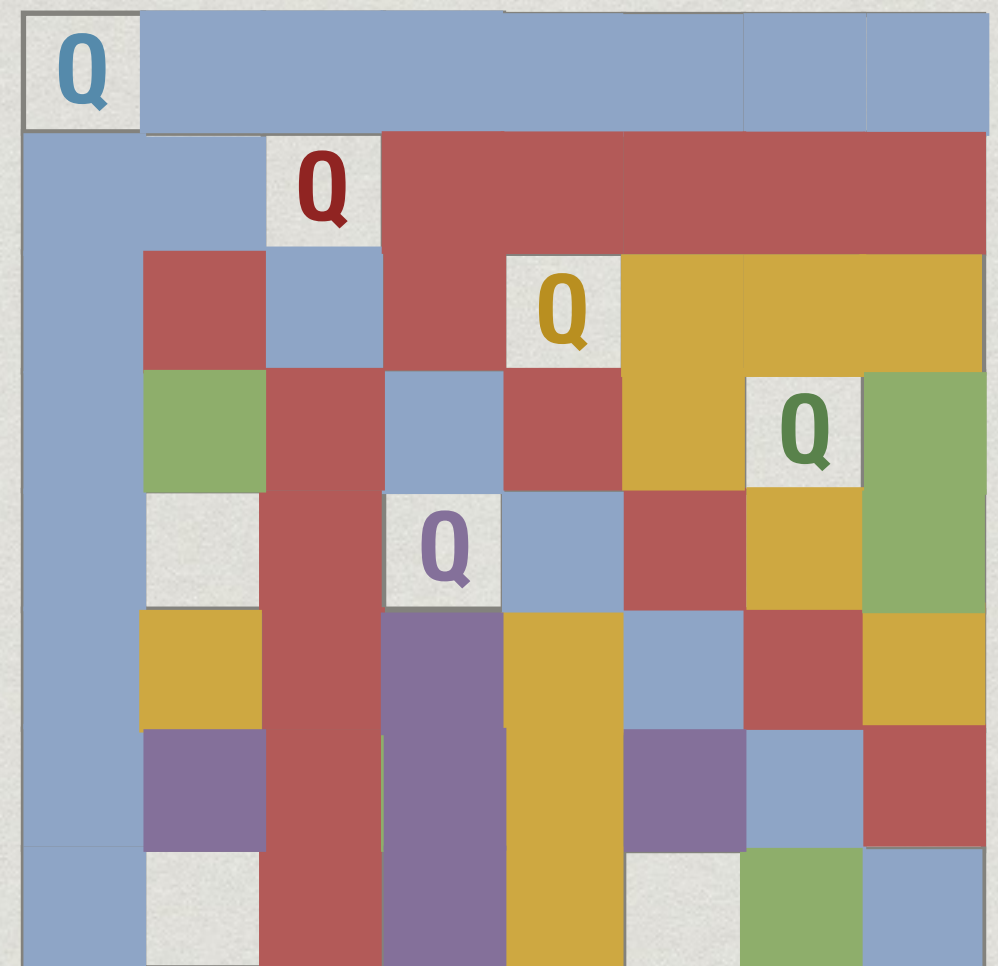
8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



8 queens

- * Can't place the a queen in the 8th row!
- * Undo 7th queen, no other choice
- * Undo 6th queen, no other choice
- * Undo 5th queen, try next



Backtracking

- * Keep trying to extend the next solution
- * If we cannot, undo previous move and try again
- * Exhaustively search through all possibilities
- * ... but systematically!

Coding the solution

- * How do we represent the board?
- * $n \times n$ grid, number rows and columns from 0 to $n-1$
 - * `board[i][j] == 1` indicates queen at (i, j)
 - * `board[i][j] == 0` indicates no queen
- * We know there is only one queen per row
- * Single list `board` of length n with entries 0 to $n-1$
 - * `board[i] == j` : queen in row i , column j , i.e. (i, j)

Overall structure

```
def placequeen(i,board): # Trying row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
    else:
        return(False) # Row i failed
```


Updating the board

- * Our 1-D and 2-D representations keep track of the queens
- * Need an efficient way to compute which squares are free to place the next queen
- * $n \times n$ **attack** grid
 - * **attack**[i][j] == 1 if (i, j) is attacked by a queen
 - * **attack**[i][j] == 0 if (i, j) is currently available
- * How do we undo the effect of placing a queen?
 - * Which **attack**[i][j] should be reset to 0?

Updating the board

- * Queens are added row by row
- * Number the queens 0 to $n-1$
- * Record earliest queen that attacks each square
 - * $\text{attack}[i][j] == k$ if (i, j) was first attacked by queen k
 - * $\text{attack}[i][j] == -1$ if (i, j) is free
- * Remove queen k — reset $\text{attack}[i][j] == k$ to -1
 - * All other squares still attacked by earlier queens

Updating the board

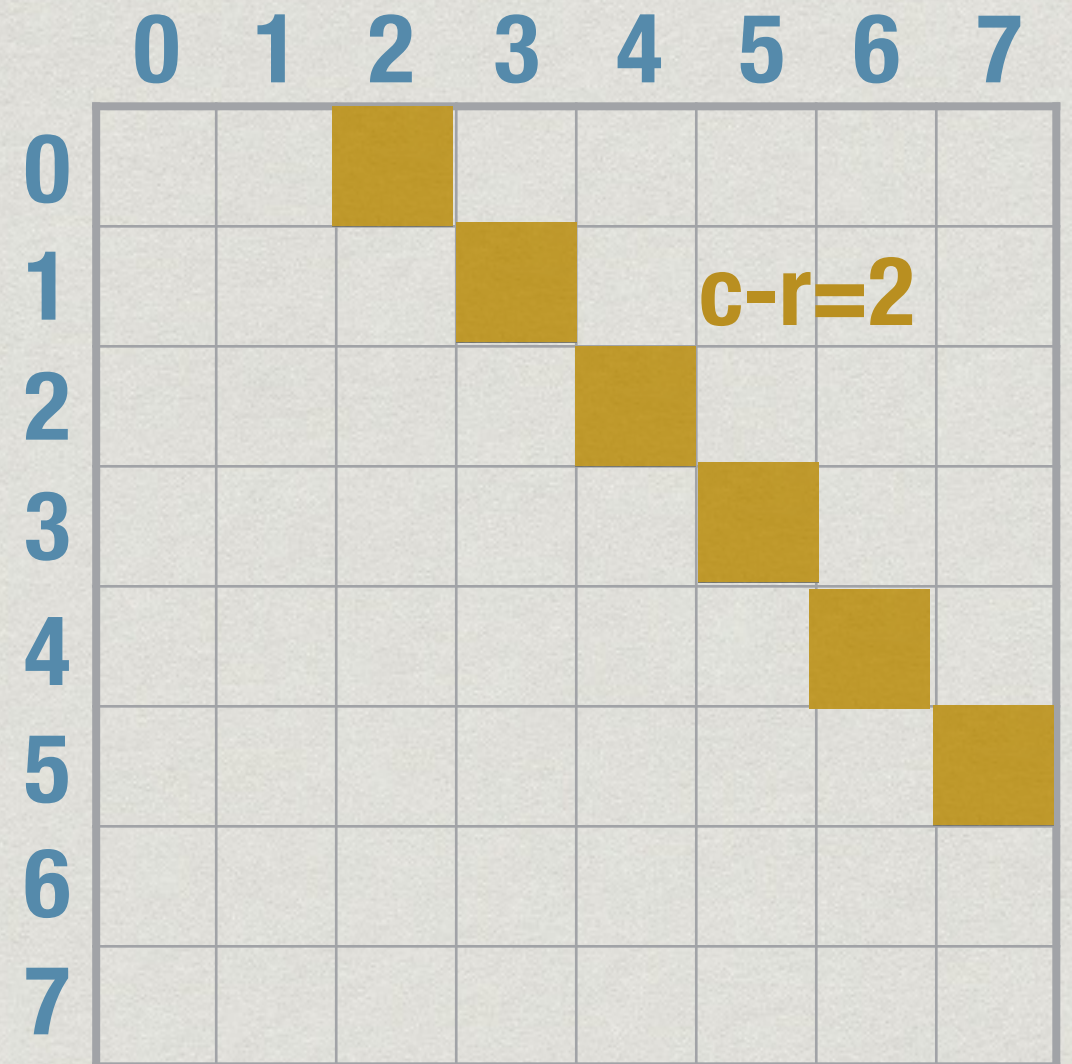
- * **attack** requires n^2 space
 - * Each update only requires $O(n)$ time
 - * Only need to scan row, column, two diagonals
- * Can we improve our representation to use only $O(n)$ space?

A better representation

- * How many queens attack row i ?
- * How many queens attack row j ?
- * An individual square (i,j) is attacked by upto 4 queens
 - * Queen on row i and on column j
 - * One queen on each diagonal through (i,j)

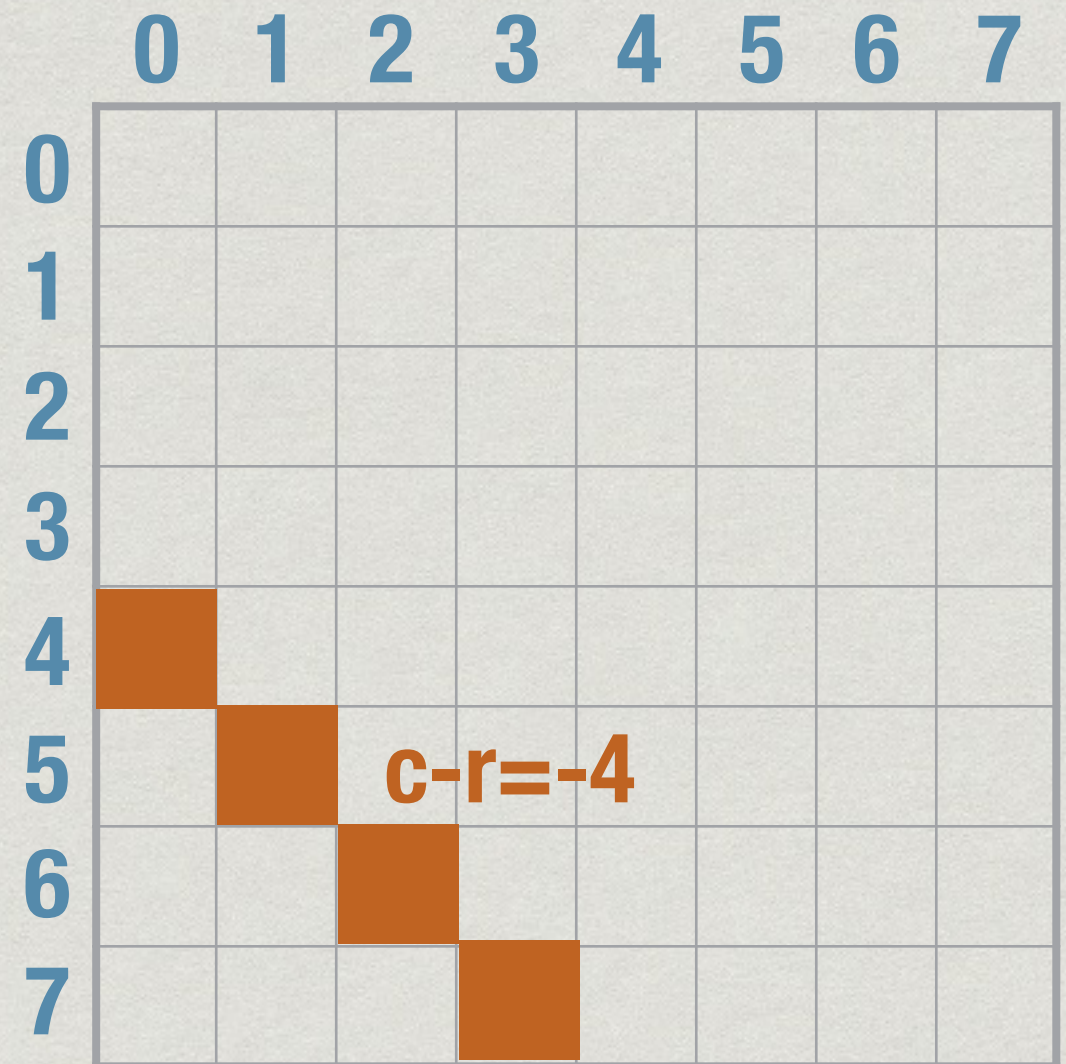
Numbering diagonals

- * Decreasing diagonal:
column - row is invariant



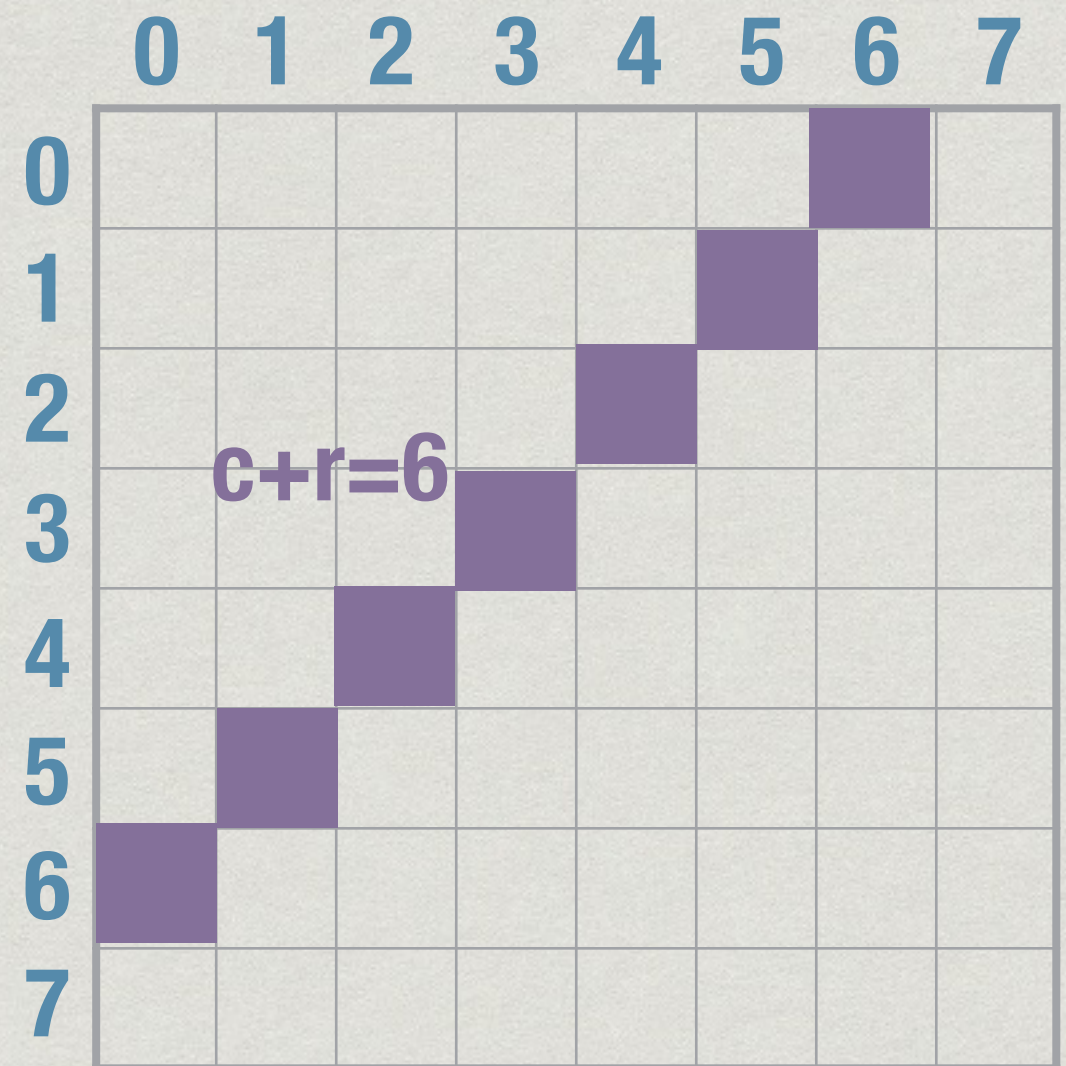
Numbering diagonals

- * Decreasing diagonal:
column - row is invariant



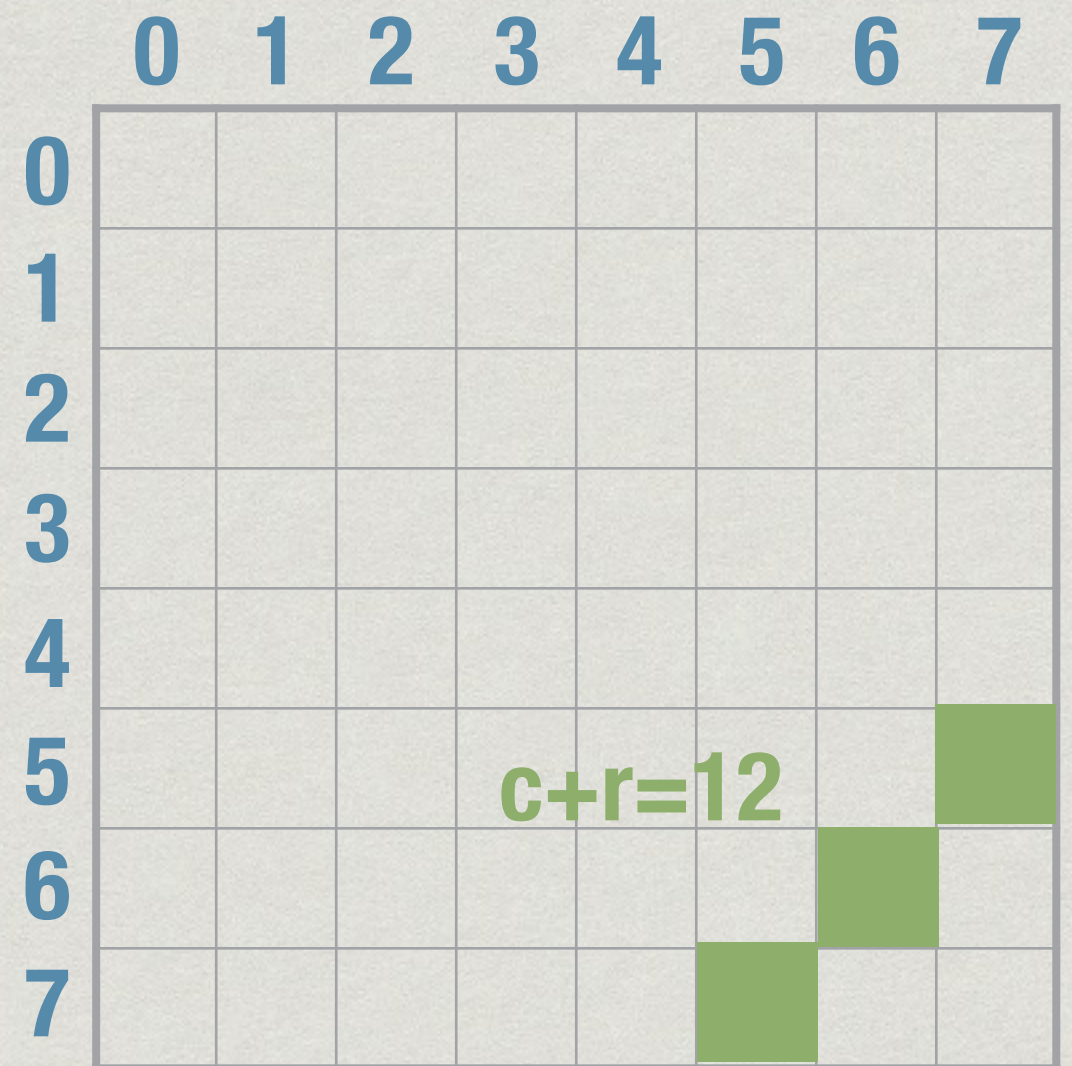
Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant



Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant



Numbering diagonals

- * Decreasing diagonal:
column - row is invariant
- * Increasing diagonal:
column + row is invariant
- * (i,j) is attacked if
 - * row i is attacked
 - * column j is attacked
 - * diagonal $j-i$ is attacked
 - * diagonal $j+i$ is attacked

	0	1	2	3	4	5	6	7
0								
1								
2								
3								
4								
5								
6								
7								

$c+r=12$

$O(n)$ representation

- * $\text{row}[i] == 1$ if row i is attacked, $0..N-1$
- * $\text{col}[i] == 1$ if column i is attacked, $0..N-1$
- * $\text{NWtoSE}[i] == 1$ if NW to SE diagonal i is attacked, $-(N-1)$ to $(N-1)$
- * $\text{SWtoNE}[i] == 1$ if SW to NE diagonal i is attacked, 0 to $2(N-1)$

Updating the board

- * (i, j) is free if
$$\text{row}[i] == \text{col}[j] == \text{NWtoSE}[j-i] == \text{SWtoNE}[j+i] == 0$$
- * Add queen at (i, j)
$$\begin{aligned} \text{board}[i] &= j \\ (\text{row}[i], \text{col}[j], \text{NWtoSE}[j-i], \text{SWtoNE}[j+i]) &= \\ &\quad (1, 1, 1, 1) \end{aligned}$$
- * Remove queen at (i, j)
$$\begin{aligned} \text{board}[i] &= -1 \\ (\text{row}[i], \text{col}[j], \text{NWtoSE}[j-i], \text{SWtoNE}[j+i]) &= \\ &\quad (0, 0, 0, 0) \end{aligned}$$

Implementation details

- * Maintain `board` as nested dictionary
 - * `board['queen'][i] = j` : Queen located at (i, j)
 - * `board['row'][i] = 1` : Row i attacked
 - * `board['col'][i] = 1` : Column i attacked
 - * `board['nwtose'][i] = 1` : NWtoSW diagonal i attacked
 - * `board['swtone'][i] = 1` : SWtoNE diagonal i attacked

Overall structure

```
def placequeen(i,board): # Trying row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            return(True) # Last queen has been placed
        else:
            extendsoln = placequeen(i+1,board)
            if extendsoln:
                return(True) # This solution extends fully
            else:
                undo this move and update board
    else:
        return(False) # Row i failed
```


All solutions?

```
def placequeen(i,board): # Try row i
    for each c such that (i,c) is available:
        place queen at (i,c) and update board
        if i == n-1:
            record solution # Last queen placed
        else:
            extendsoln = placequeen(i+1,board)
        undo this move and update board
```