

## EXPERIMENT NO. 07

Roll No.: 24141001

Batch: I1

**Title:** Optimal Binary Search trees using Dynamic Programming.

### Objectives:

- Understand the concept of Binary Search Trees (BST) and how search cost varies based on tree structure.
- Solve the Optimal Binary Search Tree (OBST) problem using Dynamic Programming.
- Minimize the expected search cost using given key frequencies and dummy keys.
- Implement the OBST algorithm in C.
- Analyze OBST's time and space complexity.

### Theory:

A Binary Search Tree (BST) allows efficient searching, insertion and deletion. However, the search cost depends on how balanced the tree is.

In many applications (e.g., compilers, dictionaries), each key has a probability of being searched. A poorly structured BST leads to high search time.

Optimal Binary Search Tree (OBST)

Given:

- Keys:  $K_1, K_2, \dots, K_n$  (sorted)
- Successful search probabilities:  $p_1, p_2, \dots, p_n$
- Unsuccessful search probabilities:  $q_0, q_1, \dots, q_n$

Goal:

Construct a BST that minimizes the expected cost of searching.

Dynamic Programming is used because:

- The structure has optimal substructure
- Overlapping subproblems exist
- A brute-force solution is exponential

OBST uses three DP tables:

- $\text{Cost}[i][j] \rightarrow$  Minimum cost between keys  $i$  to  $j$
- $\text{Weight}[i][j] \rightarrow$  Total frequency ( $p_i + q_i$ ) between  $i$  and  $j$
- $\text{Root}[i][j] \rightarrow$  Root key index producing minimum cost

### Algorithm:

OBST( $p[]$ ,  $q[]$ ,  $n$ ):

1. Create matrices:  $\text{cost}[n+1][n+1]$ ,  $\text{weight}[n+1][n+1]$ ,  $\text{root}[n+1][n+1]$

2. For  $i = 0$  to  $n$ :

$\text{cost}[i][i] = q[i]$

$\text{weight}[i][i] = q[i]$

3. For  $\text{length} = 1$  to  $n$ :

For  $i = 0$  to  $n - \text{length}$ :

$j = i + \text{length}$

$\text{weight}[i][j] = \text{weight}[i][j-1] + p[j] + q[j]$

$\text{cost}[i][j] = \infty$

For  $r = i+1$  to  $j$ :

$\text{temp} = \text{cost}[i][r-1] + \text{cost}[r][j] + \text{weight}[i][j]$

If  $\text{temp} < \text{cost}[i][j]$ :

$\text{cost}[i][j] = \text{temp}$

$\text{root}[i][j] = r$

4. Return  $\text{cost}[0][n]$  and  $\text{root}[][]$

### Program:

```
#include <stdio.h>
```

```
#include <limits.h>
```

```

int main() {
    int n, i, j, k;
    float p[20], q[20];
    float cost[20][20], weight[20][20];

    printf("Enter number of keys: ");
    scanf("%d", &n);

    printf("Enter successful search probabilities p[i]:\n");
    for (i = 1; i <= n; i++)
        scanf("%f", &p[i]);

    printf("Enter unsuccessful search probabilities q[i]:\n");
    for (i = 0; i <= n; i++)
        scanf("%f", &q[i]);

    // Initialize cost and weight
    for (i = 0; i <= n; i++) {
        cost[i][i] = q[i];
        weight[i][i] = q[i];
    }

    // OBST using Dynamic Programming
    for (int length = 1; length <= n; length++) {
        for (i = 0; i <= n - length; i++) {
            j = i + length;
            cost[i][j] = INT_MAX;
            weight[i][j] = weight[i][j - 1] + p[j] + q[j];

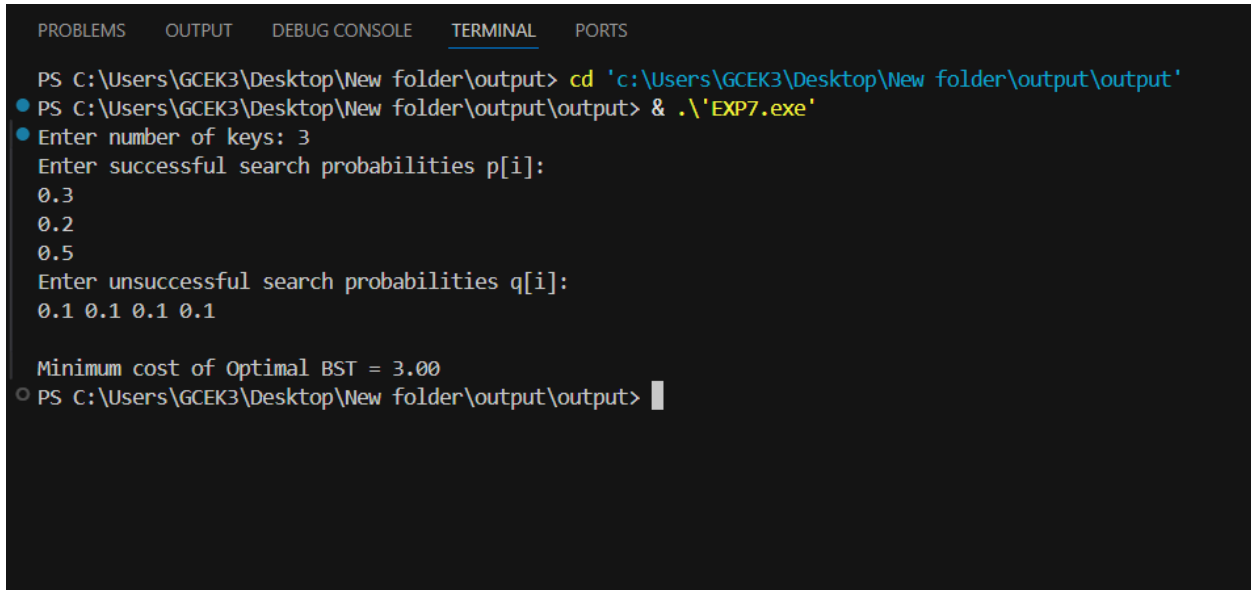
            for (k = i + 1; k <= j; k++) {
                float temp = cost[i][k - 1] + cost[k][j] + weight[i][j];
                if (temp < cost[i][j]) {
                    cost[i][j] = temp;
                }
            }
        }
    }

    printf("\nMinimum cost of Optimal BST = %.2f\n", cost[0][n]);
}

```

```
    return 0;
}
```

## OUTPUT:



```
PROBLEMS  OUTPUT  DEBUG CONSOLE  TERMINAL  PORTS

PS C:\Users\GCEK3\Desktop\New folder\output> cd 'c:\Users\GCEK3\Desktop\New folder\output\output'
PS C:\Users\GCEK3\Desktop\New folder\output\output> & .\EXP7.exe
Enter number of keys: 3
Enter successful search probabilities p[i]:
0.3
0.2
0.5
Enter unsuccessful search probabilities q[i]:
0.1 0.1 0.1 0.1

Minimum cost of Optimal BST = 3.00
PS C:\Users\GCEK3\Desktop\New folder\output\output> 
```

## Applications of OBST:

### 1. Compiler Design

Used in constructing optimal parsing tables and identifying common keywords efficiently.

### 2. Database Indexing

Minimizes search time for keys accessed with different frequencies.

### 3. File Systems

Used for directory lookup optimization.

### 4. Artificial Intelligence

Search optimization in decision trees.

### 5. Data Compression

Forms the basis for efficient symbol lookup.

**Time and space complexity:**

Time Complexity:  $O(n^3)$

Space Complexity:  $O(n^2)$

**Conclusion:**

The construction of Optimal Binary Search Trees using Dynamic Programming ensures the minimum expected search cost based on search frequencies. It is a powerful optimization method used in compilers, databases, and search applications where the frequency of accesses is known. Though the DP solution has high time complexity ( $O(n^3)$ ), it provides the most optimal BST structure.