Limitations of Finite Automata.

Canonical example

B= {a b | n = 0} = {E, ab, aabb, aaabbb, ...}

Consider lte string and where n>k.

Fastote Pinm s.t.

$$|V| = j70.$$
  $\hat{S}(8, u) = P \hat{S}(P, v) = P$   
String -  $uw$   $\hat{S}(P, w) = vEF$ 

$$\hat{S}(S,UW) = \hat{S}(\hat{S}(S,U),W) = \hat{S}(P,W) = \forall \in F$$
  
 $UW \in L(M)$   $UW = a^{n-|V|}b^n \notin B$ 

$$\mathcal{U} \mathcal{V}^{2} \mathcal{W} = a^{n+|\mathcal{V}|} b^{n} \in \mathcal{L}(m)$$

$$\hat{S}(S, \mathcal{U} \mathcal{V} \mathcal{W}) = \hat{S}(\hat{S}(\hat{S}(\hat{S}(S, \mathcal{U}), \mathcal{V}), \mathcal{V}), \mathcal{W})$$

$$= \hat{S}(\hat{S}(\hat{S}(P, \mathcal{V}), \mathcal{V}), \mathcal{W})$$

$$= \hat{S}(\hat{S}(P, \mathcal{V}), \mathcal{W})$$

$$= \hat{S}(P, \mathcal{W}) = \mathcal{V} \in F.$$

$$a^{n+|\mathcal{V}|} b^{n} \notin \mathcal{B}.$$

Contradicts our assumption that L(M)=B.

Pumping Lemma.

Let A be a regular set. Then Ite following Property holds of A.

There exists  $k \ge 0$  such that for any string x, y, z with  $xyz \in A$  and  $|y| \ge k$ , there exists strings u, v, w s.t  $y = uvw, v \ne \epsilon$  and for all  $i \ge 0$ , the string  $x u v^i w z \in A$ .

Contrapositive: Suppose A satisfies Ite following:

For all  $k \ge 0$  Itere exists strings x, y, z such that

(7P)  $y \ge 2 \le A$ ,  $y \ge k$  and for all  $y \ge k$  and with  $y = x \le k$  and  $y \ne k$ , Itere exists  $y \ge 2 \le k \le k$  and  $y \ne k$ .

Then A is not regular.

Suppose A Satisfies Ite following.

For all k≥0 lhere exists strings x,y,z buchthd (7P)  $y \in A$ ,  $|y| \geq k$  and for all  $u, v, \omega$ with y=uvw and v+E, Itere exists izo st xurwz &A.

Then A is not regular.

Claim: A is not regular Example. A = {a b | n≥m}. To show A sot is fies 7P Consider any kzo, let x= ak y=bk & z= E. Then xyz EA. Consider any split of y= urw

with  $V \neq E$ . Soy  $y = b^{j}b^{m}b^{n}$ . k = j+m+n

Let i=2. xu82wz = arb36m6m6n  $= a^k b^{j+2m+n}$  $=a^{R}b^{R+m}\notin A$